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# Gaussian Process Bandits for Top-k Recommendations

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## Abstract

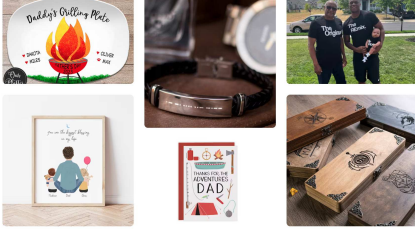
Algorithms that utilize bandit feedback to optimize top-k recommendations are vital for online marketplaces, search engines, and content platforms. However, the combinatorial nature of this problem poses a significant challenge, as the possible number of ordered top-k recommendations from  $n$  items grows exponentially with  $k$ . As a result, previous work often relies on restrictive assumptions about the reward or bandit feedback models, such as assuming that the feedback discloses rewards for all recommended items rather than offering a single scalar feedback for the entire set of top-k recommendations. We introduce a novel contextual bandit algorithm for top-k recommendations, leveraging a Gaussian process with a Kendall kernel to model the reward function. Our algorithm requires only scalar feedback from the top-k recommendations and does not impose restrictive assumptions on the reward structure. Theoretical analysis confirms that the proposed algorithm achieves sub-linear regret in relation to the number of rounds and arms. Also, empirical results using a bandit simulator show that the proposed algorithm surpasses other baselines across several scenarios.

## 1 Introduction

The top-k recommendation problem involves providing a ranked list of  $k$  items, such as news articles or products, from a pool of  $n$  items [35, 13]. Online algorithms must adapt to dynamic user preferences, making bandit algorithms suitable due to their use of limited feedback [1]. Developing bandit algorithms is challenging due to limited feedback and the need for computational efficiency in real-time recommendation environments. Recent research on user interfaces for recommendations shows that the overall layout of the recommendation page is crucial for user appeal as designs transition from simple dropdown lists to complex layouts [17, 13, 18]. As a result, bandit algorithms must comprehensively select and display all top-k items jointly rather than merely selecting the most relevant  $k$  items and displaying them in decreasing order of user relevance [32].

The joint consideration of top-k items makes the number of arms (possible actions for the bandit algorithm) combinatorially large, i.e.,  $\Theta(n^k)$ . Previous research on bandit algorithms often impose strict assumptions about feedback models [31, 21], e.g., *semi-bandit* feedback gives a scalar value for each of the top  $k$  items. Although semi-bandit feedback decomposes feedback from the combinatorial number of arms to feedback for every recommended item, it is often unavailable [33]. Another prevalent feedback assumption is *cascade* browsing [16], which posits that users examine items in a pre-determined order and stop searching upon finding a desirable item, which provides item-specific scalar feedback but does not fully capture possible non-linear interactions [27]. Figure 1 illustrates the limitations of the cascade model in accurately representing user interactions within contemporary top-k recommendation interfaces. These limitations motivate us to focus on a strictly more general setting of the *full-bandit* feedback, where a single value for the entire top-k set is assumed [24].

Beyond feedback assumptions, the reward of bandit algorithms must be decomposable into scalars over individual items to avoid a combinatorial explosion of arms, which is not always possible. For



**Figure 1:** A snapshot from Etsy showcases Father’s Day shopping recommendations. There is no obvious linear search order, which challenges the assumptions of the cascade model. Additionally, item proximity and arrangement are likely to influence clicks, suggesting a complex interaction pattern and advocating for full-bandit feedback without assumptions about user interaction with the recommended items.

**Table 1:** Compute and memory analysis for GP-TopK. Rows indicate different costs: overall *compute* and overall *memory* for  $T$  rounds, time for matrix-vector multiplication (**mvm**) with the kernel matrix  $K_{X_t}$  at time  $t$ , and time to update  $K_{X_t}$ . Columns indicate approaches: *kernel approach* uses full kernel matrices, while our novel *feature approach* performs the same operations via feature expansions and scales better with respect to  $T$ . The symbols  $c$ ,  $k$ , and  $T$  represent the embedding size for contexts, the number of items, and the number of rounds, respectively.

Tasks	<i>kernel approach</i>	<i>feature approach</i>
<i>compute</i>	$\mathcal{O}(T^3)$	$\mathcal{O}(c \cdot k^2 \cdot T^2)$
<i>memory</i>	$\mathcal{O}(T^2)$	$\mathcal{O}(c \cdot k^2 \cdot T)$
<b>mvm</b> ( $K_{X_t}$ )	$\mathcal{O}(t^2)$	$\mathcal{O}(c \cdot k^2 \cdot t)$
<i>compute</i> $K_{X_t}$	$\mathcal{O}((c + k^2) \cdot t)$	$\mathcal{O}(c \cdot k^2)$

instance, modern e-commerce platforms consider objectives such as diversity and fairness [1], which cannot be measured by focusing solely on individual items [15]. This necessitates algorithms for full-bandit feedback settings without assumptions about the objective or rewards [24].

This work develops a bandit algorithm that uses Gaussian processes (GPs) to model rewards under full-bandit (i.e., one scalar value) feedback. GPs are selected for their flexibility in modeling feedback for discrete, continuous, and mixed domains, such as continuous contexts and discrete rankings [34]. Additionally, unlike parametric models that require optimization for accumulated feedback, GP model updates are computationally inexpensive, involving only data updates [24]. While inference for GPs may generally face computational limits, we will develop efficient inference methods tailored to our proposed algorithm. Another challenge in developing GP-based bandit algorithms for top-k recommendations is creating expressive positive-definite kernels that capture the similarity between top-k recommendations [9].

GPs have been previously explored for bandit algorithms [28, 19]. Krause et al. [14] used GPs for contextual bandits in continuous domains; we focus on the discrete domain of top-k recommendations. Vanchinathan et al. [29] used GPs with a position-based feedback model, and Wang et al. [32] used GPs with semi-bandit feedback for recommending top-k items. In contrast, our work does not focus on a specific reward model or feedback assumption, and develops an efficient GP-based bandit algorithm for top-k recommendations.

## 1.1 Contributions

Our primary contribution is the GP-TopK algorithm, a contextual bandit algorithm for recommending top-k items. This algorithm operates in a full-bandit feedback setting without relying on assumptions on reward, offering broader applicability than prior works. We leverage GPs with variants of the *Kendall* kernel [12] to model the reward function and optimize the upper confidence bound (UCB) [28] acquisition function to select the next arm. Further, we give a novel weighted convolutional Kendall kernel for top-k recommendations that address pathologies in existing variants of the Kendall kernel applied to top-k recommendations.

Our second key contribution is to improve the scalability of the GP-TopK algorithm for longer time horizons. The initial computational demand for top-k ranking with GP-TopK is  $\mathcal{O}(T^4)$  for  $T$  rounds. We first reduce this to  $\mathcal{O}(T^3)$  using iterative algorithms from numerical linear algebra [26]. Then, we derive sparse feature representations for the novel weighted convolutional Kendall kernel, which, allows us to further improve the overall compute requirements from  $\mathcal{O}(T^4)$  to  $\mathcal{O}(T^2)$  and memory requirements from  $\mathcal{O}(T^2)$  to  $\mathcal{O}(T)$ . Table 1 summarizes these time and memory requirement improvements, including their dependence on other parameters.

We also provide a theoretical analysis showing that GP-TopK’s regret is sub-linear in  $T$  and benefits from the feature representations of the Kendall kernels we introduce. We show the regret’s upper

bound is almost quadratic in  $n$ , which improves significantly over the naive bound of  $\Theta(n^k)$  for top-k recommendations without using specialized kernels [28]. Finally, we empirically validate GP-TopK's regret through simulations on real-world datasets and show improved regret compared to baselines.

## 1.2 Organization

The remainder of this paper is as follows: Section 2 presents Kendall kernels for full and top-k rankings, including the novel weighted convolutional Kendall kernel. Section 3 presents faster matrix-vector multiplication (MVM) algorithms for Kendall kernels, making the proposed bandit algorithm faster, as detailed later in Section 4, along with the regret analysis. Lastly, Sections 5 and 6 present empirical results and discussion, respectively.

## 2 Kendall Kernels for Full and Top-k Rankings

This section overviews Kendall kernels and their extensions for top-k recommendations. First, we establish some notation. Let  $[n] = \{1, 2, \dots, n\}$ , and let  $\pi$  represent a top-k ranking, which is an ordered tuple of  $k$  distinct elements from  $[n]$ . We use  $\sigma$  to denote a full ranking ( $k = n$ ) and let  $\Pi^k$  represent the set of all possible top-k rankings. Note that  $|\Pi^k| = \Theta(n^k)$ . The vector  $\mathbf{p}^\sigma \in \mathbb{R}^n$  corresponds to a full ranking  $\sigma$  with entry  $\mathbf{p}_i^\sigma$  giving the rank of item  $i$ . For top-k rankings,  $\mathbf{p}^\pi \in \mathbb{R}^n$  is constructed by arbitrarily assigning distinct ranks to items not in the top  $k$ . Indicator functions  $\mathbf{p}_{i < j}^\sigma$  and  $\mathbf{p}_{i > j}^\sigma$  indicate whether item  $i$  is ranked before or after item  $j$ , respectively in  $\sigma$ . Also,  $\mathbf{p}_{i < j}^\pi$  and  $\mathbf{p}_{i > j}^\pi$  are indicator functions defined for top-k rankings.

### 2.1 Kendall Kernels for Full Rankings

Jiao et al. [9] showed that the Kendall tau rank correlation [12] is a positive definite (p.d.) kernel for full rankings, which we refer to as the standard Kendall (SK) kernel. The weighted Kendall (WK) kernel generalizes the SK kernel by differentially weighting item pairs [10]. Specifically, the SK and WK kernels for full rankings  $\sigma_1, \sigma_2$  are defined as:

$$k^{sk}(\sigma_1, \sigma_2) := \frac{1}{\binom{n}{2}} \sum_{i < j} \eta_{i,j}(\sigma_1, \sigma_2) \quad (1)$$

$$k^{wk}(\sigma_1, \sigma_2) := \frac{1}{\binom{n}{2}} \sum_{i < j} w((\mathbf{p}_i^{\sigma_1}, \mathbf{p}_j^{\sigma_1}), (\mathbf{p}_i^{\sigma_2}, \mathbf{p}_j^{\sigma_2})) \cdot \eta_{i,j}(\sigma_1, \sigma_2), \quad (2)$$

where  $\eta_{i,j}$  is 1 if the pair  $(i, j)$  is *concordant* (ordered the same in both rankings) and  $-1$  otherwise; concretely,  $\eta_{i,j}(\sigma_1, \sigma_2) := \mathbf{p}_{i < j}^{\sigma_1} \cdot \mathbf{p}_{i < j}^{\sigma_2} + \mathbf{p}_{i > j}^{\sigma_1} \cdot \mathbf{p}_{i > j}^{\sigma_2} - \mathbf{p}_{i < j}^{\sigma_1} \cdot \mathbf{p}_{i > j}^{\sigma_2} - \mathbf{p}_{i > j}^{\sigma_1} \cdot \mathbf{p}_{i < j}^{\sigma_2}$ ; and  $ww((\mathbf{p}_i^{\sigma_1}, \mathbf{p}_j^{\sigma_1}), (\mathbf{p}_i^{\sigma_2}, \mathbf{p}_j^{\sigma_2}))$  is the value of a positive definite weighting kernel  $w(\cdot, \cdot) : [n]^2 \times [n]^2 \mapsto \mathbb{R}$  that operates on pairs of ranks. The  $w_{i,j}$  allows flexibility and can assign varying importance to ranks, similar to the discounted cumulative gain (DCG) metric [7]. Note that both SK and WK kernels are p.d. and right-invariant with respect to  $\Pi^n$  [10]. In other words, they compute similarity based only on the relative ranks of pairs, not on the labels of items, as evident from Equations 1 and 2.

### 2.2 Kendall Kernels for Top-k Rankings

**Weighted Kendall (WK) and Convolutional Kendall (CK) kernels.** To adapt the WK kernel from full rankings to top-k rankings, Jiao et al. [10] set the weighting function  $w(i, j, \sigma_1, \sigma_2)$  to zero if either item is not in the top-k of either ranking. This scheme yields a p.d. kernel but does not consider items outside the intersection of top-k rankings. The convolutional operation offers an alternative for adapting the standard Kendall kernel to top-k rankings. Let  $B_\pi$  denote the set of full rankings consistent with the top-k ranking  $\pi$  (i.e., for every item  $i$  in  $\pi$ ,  $\forall \sigma \in B_\pi, \mathbf{p}_i^\sigma = \mathbf{p}_i^\pi$ ). The CK kernel is defined as:

$$k^{ck}(\pi_1, \pi_2) = \frac{1}{|B_{\pi_1}| \cdot |B_{\pi_2}|} \sum_{\sigma_1 \in B_{\pi_1}, \sigma_2 \in B_{\pi_2}} k^{sk}(\sigma_1, \sigma_2), \quad (3)$$

where  $k^{sk}$  is the standard Kendall kernel. The CK kernel is a p.d. kernel as it is a convolution of another p.d. kernel [5]. Unlike the WK kernel for top-k rankings, the CK kernel accounts for

114 items not in both top-k rankings. However, computing the CK kernel using Eq. (3) is expensive as it  
 115 requires exponentially many evaluations of the kernel  $k^{sk}$  in the double summation. Therefore, Jiao  
 116 et al. [9] developed an efficient algorithm to bypass this double summation and compute the kernel in  
 117  $\mathcal{O}(k \log k)$  time.

118 **Proposed Weighted Convolutional Kendall (WCK) Kernel.** To combine the strengths of the WK  
 119 and CK kernels for top-k rankings, we propose the weighted convolutional Kendall kernel for top-k  
 120 rankings  $\pi_1$  and  $\pi_2 \in \Pi^k$ :

$$k^{wck}(\pi_1, \pi_2) := \frac{1}{|B_{\pi_1}| \cdot |B_{\pi_2}|} \sum_{\sigma_1 \in B_{\pi_1}, \sigma_2 \in B_{\pi_2}} k^{wk}(\sigma_1, \sigma_2), \quad (4)$$

121 where  $k^{wk}$  represents the weighted Kendall kernel for full rankings  $\sigma_1, \sigma_2 \in \Pi^n$ . The proposed WCK  
 122 kernel combines the flexibility of weighting different ranks (among the top-k items) differently, like  
 123 the WK kernel, with the ability to account for items outside the intersection of both top-k rankings,  
 124 like the CK kernel. Also, since it's a convolution of a p.d. kernel it is also a p.d.. However, its  
 125 computation is again challenging as the RHS of Equation 4 evaluates  $k^{wk}$  exponentially many times.  
 126 To simplify, we focus on a specific form of rank weights for  $k^{wk}$ , which we call *product-symmetric*  
 127 rank weights:

$$w_{ps}((i_1, j_1), (i_2, j_2)) := w_s(i_1, j_1) \cdot w_s(i_2, j_2), \quad (5)$$

128 where,  $w_s(i, j) : [n] \times [n] \mapsto \mathbb{R}$  is a symmetric function, i.e.,  $w_s(i, j) = w_s(j, i)$ . Notably, the WCK  
 129 kernel can be computed efficiently for the case of  $w_{ps}$  weights (Claim 1 below).

130 The WCK kernel, even with the relatively  
 131 simple  $w_{ps}$  weights, has notable properties  
 132 as shown in Table 2. In this table we use  
 133  $w_s(i, j) = \frac{1}{\log(i+1)} \cdot \frac{1}{\log(j+1)}$  inspired by the  
 134 DCG metric used in recommendation sys-  
 135 tems [7]. Note that the WK kernel ranks two  
 136 rankings with no overlap ( $\pi_0$  and  $\pi_1$ ) as more  
 137 similar than two rankings with the same items  
 138 but reverses ordering ( $\pi_0$  and  $\pi_2$ ), a clear  
 139 pathology. On the other hand, the CK kernel  
 140 fails to differentiate between reversed pairs at  
 141 different ranks ( $k^{ck}(\pi_0, \pi_3) = k^{ck}(\pi_0, \pi_4)$ ),  
 142 another clear limitation. The WCK kernel  
 143 with product-symmetric ranking weights can  
 144 address these shortcomings and provide a  
 145 more nuanced similarity comparison for top-k rankings.

**Table 2:** Comparison of Kendall kernel similarities for top-k rankings. The table shows kernel values  $k(\pi_0, \cdot)$  for the top-k ranking  $\pi_0 = [1, 2, 3]$  with other rankings ( $\pi_1, \pi_2, \pi_3, \pi_4$ ) for  $n = 7$  and  $k = 3$ . Rankings are ordered left to right by increasing similarity to  $\pi_0$ . WCK kernel values, based on DCG rank weights, increase from left to right. All kernels are unit normalized.

Top-k Kernels	$\pi_1$ [4, 5, 6]	$\pi_2$ [3, 2, 1]	$\pi_3$ [2, 1, 3]	$\pi_4$ [1, 3, 2]
WK	0.00	-1.00	0.33	0.33
CK	-0.60	0.60	0.87	0.87
WCK	-0.38	0.09	0.46	0.87

**Claim 1.** *The weighted convolutional Kendall kernel (Equation 4) with product-symmetric rank weights (Equation 5) can be computed in  $\mathcal{O}(k^2)$  time.*

146 Appendix A provides the proof, which exploits the structure of product-symmetric rank weights  $w_{ps}$   
 147 to establish the existence of a feature representation for the WCK kernel, as given in Claim 3 below.  
 148 We then show that the inner product of these features, and thus the WCK kernel, can be computed  
 149 in  $\mathcal{O}(k^2)$  time (Algorithm 2 in the appendix). Similarly to the result of Jiao et al. [9], we can avoid  
 150 exponentially many evaluations of  $k^{wk}$  on the RHS of Equation (4) by computing the WCK kernel  
 151 directly by another means.

### 152 3 Fast Matrix-Vector Multiplication with Kendall Kernel Matrices

153 In GPs, inference can be accelerated by using iterative algorithms that take advantage of fast matrix-  
 154 vector-multiplications (MVMs) with the kernel matrix [3]. This section focuses on fast algorithms  
 155 for kernel MVMs that exploit the structure of Kendall kernel matrices. Specifically, let  $\text{mvm}(K_{X_t})$   
 156 denote the running time to multiply the  $t \times t$  kernel matrix  $K_{X_t} = (k(x_i, x_j))_{x_i, x_j \in X_t}$  by a vector.  
 157 Naively,  $\text{mvm}(K_{X_t}) = \mathcal{O}(t^2)$ . However, if  $k(x_i, x_j) = \phi^a(x_i)^T \phi^b(x_j)$  for vectors  $\phi^a(x_i)$  and  
 158  $\phi^b(x_j)$  with only  $z$  non-zero entries, then  $\text{mvm}(K_{X_t})$  reduces to  $\mathcal{O}(z \cdot t)$ , which is much faster than

159  $\mathcal{O}(t^2)$  when  $z \ll t$ . When  $\phi^a = \phi^b$ , we call  $\phi^a$  the *linear feature vector* for the kernel  $k$ . Before  
 160 focusing on top-k ranking kernels, we provide a linear feature vector for the WK kernel on full  
 161 rankings as defined in Equation 2.

**Claim 2.** Let  $\phi^{wk}(\sigma) : \Pi^n \mapsto \mathbb{R}^{\binom{n}{2}}$  be a vector indexed by unique item pairs  $(i, j)$ , defined as:

$$\phi_{i,j}^{wk}(\sigma) := \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_i^\sigma, \mathbf{p}_j^\sigma) \cdot (\mathbf{p}_{i < j}^\sigma - \mathbf{p}_{i > j}^\sigma),$$

where  $w_s$  is the symmetric weighting function in product-symmetric weights. Then,  $\phi^{wk}$  is a linear feature vector for the weighted Kendall kernel with product-symmetric weights  $w_{ps}$ .

162 Using Claim 2, the linear feature vector for the WK kernel can be extended to the WK top-k ranking  
 163 kernel utilizing the structure of the product-symmetric weights. Furthermore, the feature vector  
 164  $\phi^{wk}(\pi)$  contains only  $\mathcal{O}(k^2)$  non-zero entries due to the WK kernel's focus on item pairs within the  
 165 top-k, resulting in  $\text{mvm}(K_{X_t}) = \mathcal{O}(k^2 \cdot t)$  for the WK kernel matrix.

166 Next, we focus on fast MVMs with  
 167 the WCK kernel, which includes  
 168 the CK kernel as a particular case.  
 169 We observe that any convolutional  
 170 kernel inherits linear features from  
 171 its constituent kernel. Specifically,  
 172  $\sum_{\sigma \in B_\pi} \phi_{i,j}^{wk}(\sigma)$  forms a feature vec-  
 173 tor for the WCK kernel, which follows  
 174 from Equation 4 and Claim 2.

175 However, computing this feature vec-  
 176 tor explicitly requires the exponential  
 177 summation over all  $\sigma \in B_\pi$ . Claim 3  
 178 shows that the summation can be com-  
 179 puted analytically and provides ex-  
 180 plicit linear feature vectors for the  
 181 WCK and CK kernels. It also shows  
 182 that  $\phi^{wck}$  has only  $\mathcal{O}(k^2 + 2nk)$  non-  
 183 zeros among its  $\mathcal{O}(n^2)$  entries. Con-  
 184 sequently,  $\text{mvm}(K_{X_t})$  for the WCK  
 185 kernel requires  $\mathcal{O}((k^2 + 2nk) \cdot t)$  operations, which improves from  $\mathcal{O}(t^2)$  to linear in  $t$ . However,  
 186 this introduces a dependence on  $n$  and is beneficial only for  $n \leq t$ . We next leverage redundancy in  
 187  $\phi^{wck}$  to eliminate this dependence, leading to the following main theorem about the  $\text{mvm}(K_{X_t})$ .

**Claim 3.** Let  $\phi^{wck}(\pi) : \Pi^k \mapsto \mathbb{R}^{\binom{n}{2}}$  be a vector in-  
 dexed by unique item pairs  $(i, j)$  given as:  $\phi_{i,j}^{wck}(\pi) :=$   
 $\frac{1}{\sqrt{\binom{n}{2}}} \cdot \mathbf{w}_{i,j}^{wck}(\pi) \cdot (\mathbf{p}_{i < j}^\pi - \mathbf{p}_{i > j}^\pi)$ , where  $\mathbf{w}_{i,j}^{wck}(\pi)$  is  
 determined as follows:

$$\mathbf{w}_{i,j}^{wck}(\pi) = \begin{cases} w_s(\mathbf{p}_i^\pi, \mathbf{p}_j^\pi) & \text{if } \mathbf{p}_i^\pi \in [k] \text{ \& } \mathbf{p}_j^\pi \in [k] \\ w_s(\mathbf{p}_i^\pi, \cdot) & \text{else if } \mathbf{p}_i^\pi \in [k] \text{ \& } \mathbf{p}_j^\pi \notin [k], \\ w_s(\mathbf{p}_j^\pi, \cdot) & \text{else if } \mathbf{p}_i^\pi \notin [k] \text{ \& } \mathbf{p}_j^\pi \in [k], \\ 0 & \text{otherwise,} \end{cases}$$

where  $w_s$  denotes symmetric weights and  $w_s(\ell, \cdot) = \frac{1}{n-k} \sum_{j=k+1}^n w_s(\ell, j)$ . Then, the vector  $\phi^{wck}$  is a lin-  
 ear feature vector for the WCK kernel  $k^{wck}$ . By uni-  
 formly setting  $w_s(\cdot, \cdot) \equiv 1$  in the definitions above,  
 $\phi_{i,j}^{wck}(\pi)$  specializes to a linear feature vector for the  
 CK kernel.

**Theorem 1.** For the WCK kernel with product-symmetric weights  $w_{ps}$ , the computational  
 complexity of multiplying the kernel matrix  $K_{X_t}$  with any admissible vector is  $\mathcal{O}(k^2 t)$ , i.e.,  
 $\text{mvm}(K_{X_t}) = \mathcal{O}(k^2 t)$ , where  $X_t$  is any arbitrary set of  $t$  top-k rankings.

188 Appendix A provides the proof in two steps. First, we leverage the values of  $\phi^{wck}$  from Claim 3 and  
 189 categorize  $\phi^{wck}(\pi_1)^T \phi^{wck}(\pi_2)$  based on item pairs, as summarized in Table 4. Next, we show that  
 190 only five combinations yield non-zero values, i.e.,  $\phi^{wck}(\pi_1)^T \phi^{wck}(\pi_2) = \sum_{i=1}^5 s_i(\pi_1, \pi_2)$ . Each  
 191 term  $s_i(\pi_1, \pi_2)$  is a dot product of vectors  $\phi^{a_i}(\pi_1)^T \phi^{b_i}(\pi_2)$ , which contains at most  $\mathcal{O}(k^2)$  non-zero  
 192 entries. Thus, for WCK and CK kernels,  $\text{mvm}(K_{X_t}) = \mathcal{O}(k^2 t)$  as these vectors for all five terms  
 193 have only  $\mathcal{O}(k^2)$  non-zero entries. Consequently, Theorem 1 demonstrates that using these vector  
 194 representations for top-k rankings yields faster MVMs, i.e.,  $\text{mvm}(K_{X_t}) = \mathcal{O}(k^2 t) \ll \mathcal{O}(t^2)$ .

## 195 4 Proposed GP-TopK Bandit Algorithm

196 This section outlines the top-k recommendation problem and introduces a generic contextual bandit  
 197 algorithm for top-k recommendations. We then explain how the components of this algorithm

are instantiated using our GP approach. An analysis of the proposed algorithm’s computational complexity and regret follows this.

Let  $T$  be the number of rounds. Contexts  $\mathcal{C}$  are in a finite  $c$ -dimensional space,  $\mathcal{C} \subseteq \mathbb{R}^c$ . In the  $t^{\text{th}}$  round, we receive a context  $\mathbf{c}_t \in \mathcal{C}$  and select a top-k ranking  $\pi_t \in \Pi^k$ . We then obtain a noisy reward  $y_t = \hat{f}(\mathbf{c}_t, \pi_t) + \epsilon_t$ , where  $\hat{f}$  is the true reward function and  $\epsilon_t$  is round-independent noise. The regret is  $r_t := \max_{\pi' \in \Pi^k} \hat{f}(\mathbf{c}_t, \pi') - \hat{f}(\mathbf{c}_t, \pi_t)$ , with cumulative regret  $R_T := \sum_{t=1}^T r_t$ . The accumulated data at the  $t^{\text{th}}$  round is  $\mathcal{D}_t = (\mathbf{c}_i, \pi_i, y_i)_{i=1}^t$ . Algorithm 1 presents the bandit algorithm’s generic schematic, aiming to minimize cumulative regret while ensuring computational efficiency.

**Algorithm 1** Contextual Bandit Algorithm for Top-k Recommendations

**Input:** Total rounds  $T$ , initial reward model  $\mathcal{M}_0$ , and acquisition function  $\mathcal{A}\mathcal{F}$ .

- 1: **for**  $t = 1, \dots, T$  **do**
- 2:   Observe a context  $\mathbf{c}_t$  from the context space  $\mathcal{C}$ .
- 3:   Select a top-k ranking  $\pi_t$  that maximizes  $\mathcal{A}\mathcal{F}(\mathcal{M}_{t-1}(\mathbf{c}_t, \pi))$  for the context  $\mathbf{c}_t$ .
- 4:   Obtain the scalar reward  $y_t$ .
- 5:   Update the reward model  $\mathcal{M}_t$  using the accumulated feedback  $\mathcal{D}_t$ .
- 6: **end for**

The above algorithm requires two components: a) a reward model  $\mathcal{M}_t$  that can estimate the reward value given any context and top-k ranking utilizing the accumulated feedback  $\mathcal{D}_t$  and b) an acquisition function  $\mathcal{A}\mathcal{F}$  for selecting  $\pi_t$  given the current reward model  $\mathcal{M}_t$  and observed context  $\mathbf{c}_t$ .

**Reward model  $\mathcal{M}$  and acquisition function  $\mathcal{A}\mathcal{F}$ .** Our proposed GP-TopK bandit algorithm uses GP regression to model rewards for contexts and top-k rankings. GP regression is briefly covered in Section B.1 for completeness. In essence, the reward model  $\mathcal{M}$  maintains a distribution over functions  $f$ , i.e.,  $f \sim \mathcal{N}(0, k(\cdot, \cdot))$ , where  $k$  is a product kernel function over both contexts and top-k rankings ( $\mathcal{C} \otimes \Pi^k$ ). Specifically, the kernel function  $k$  is defined as follows:

$$k((\mathbf{c}_1, \pi_1), (\mathbf{c}_2, \pi_2)) := k^c(\mathbf{c}_1, \mathbf{c}_2) \cdot k^r(\pi_1, \pi_2), \quad (6)$$

where  $k^c(\mathbf{c}_1, \mathbf{c}_2) = \mathbf{c}_1^T \mathbf{c}_2$  is the dot-product kernel and  $k^r$  is a kernel for top-k rankings. We use variants of the Kendall kernel for  $k^r$  from Section 2. Updating the reward model  $\mathcal{M}_t$  at the  $t^{\text{th}}$  round involves adding new data points to our GP regression, which is computationally inexpensive compared to the fine-tuning steps required by parametric models to incorporate the latest feedback. We utilize the UCB function for the acquisition function as it effectively balances exploration and exploitation by selecting actions that maximize the upper confidence bound on the estimated reward [28]. The UCB acquisition function is  $\mathcal{A}\mathcal{F}(\mathcal{M}_t(\mathbf{c}_t, \pi)) := \mu_{f|\mathcal{D}}((\mathbf{c}_t, \pi)) + \beta^{\frac{1}{2}} \cdot \sigma_{f|\mathcal{D}}((\mathbf{c}_t, \pi))$ , where  $\sigma_{f|\mathcal{D}}((\mathbf{c}_t, \pi)) = \sqrt{k_{f|\mathcal{D}}((\mathbf{c}_t, \pi), (\mathbf{c}_t, \pi))}$ , and  $\beta$  controls the trade-off between exploration and exploitation.  $\mu_{f|\mathcal{D}}$  and  $k_{f|\mathcal{D}}$  are the GP posterior mean and covariance functions, as detailed in Section B.1. At the  $t^{\text{th}}$  round, the algorithm selects the top-k ranking  $\pi \in \Pi^k$  that maximizes  $\mathcal{A}\mathcal{F}(\mathcal{M}_t(\mathbf{c}_t, \pi))$ , which is performed using local search [19], as detailed in Appendix B.

**Computational complexity.** The GP-TopK bandit algorithm requires no compute for model updates. i.e., updating  $\mathcal{M}_t$  in the Line 5 of the Algorithm 1 requires only updating  $\mathcal{D}_t$ . The GP-TopK relies on local search for optimizing  $\mathcal{A}\mathcal{F}$ , so the compute requirements arise only from  $\mathcal{A}\mathcal{F}$  evaluations in the local search. As shown in Section B.1, computing the GP variance term for evaluating  $\mathcal{A}\mathcal{F}$ , i.e.,  $\sigma_{f|\mathcal{D}}((\mathbf{c}_t, \pi))$  involves solving  $[K_{X_t} + \sigma^2 I]^{-1} \mathbf{v}$  for a vector  $\mathbf{v}$ , where  $X_t = [(\mathbf{c}_1, \pi_1), \dots, (\mathbf{c}_t, \pi_t)]$ . Naively, this requires  $\mathcal{O}(t^3)$  time per round, leading to  $\mathcal{O}(T^4)$  over  $T$  rounds. Iterative algorithms can expedite the process using our results on fast MVMs with kernel matrices, as discussed previously in Section 3 [25]. Theorem 2 formalizes the computational demands of the GP-TopK algorithm.

**Theorem 2.** Assuming a fixed number of iterations required by the iterative algorithms, the total computational time for running the GP-TopK bandit algorithm for  $T$  rounds of top- $k$  recommendations, using the contextual product kernel (Equation 6), is  $\mathcal{O}(k^2 c \ell T^2)$ . This applies to WK, CK, and WCK top- $k$  ranking kernels, where  $\ell$  is the number of local search evaluations.

The proof of Theorem 2 is provided in Appendix B. It demonstrates efficiency gains from integrating feature representation with iterative algorithms, reducing the computational time from  $\mathcal{O}(T^3)$  to  $\mathcal{O}(T^2)$ . This is a significant improvement, as even one MVM with  $K_{X_t}$  using the full kernel matrix at each round requires  $\mathcal{O}(T^3)$  time. Furthermore, the theorem also shows that the running time of the GP-TopK algorithm does not explicitly depend on the number of items  $n$ .

**Regret analysis of the GP-TopK algorithm.** The cumulative regret of the proposed algorithm is  $R_T = \sum_{t=1}^T \max_{\pi'_t \in \Pi^k} \hat{f}(\mathbf{c}_t, \pi'_t) - \hat{f}(\mathbf{c}_t, \pi_t)$ , where  $\pi_t$  is the ranking chosen at round  $t$ . Optimizing cumulative regret for top- $k$  recommendations is challenging, requiring learning the context-arm relationship and outperforming the best mapping. To bound cumulative regret, regularity assumptions on the reward function  $\hat{f}$  are necessary [28, 14]. We consider the following two assumptions, either of which suffices.

**Assumption 1.**  $\mathcal{X}$  is finite, meaning that only finite contexts are considered ( $|\mathcal{C}| < \infty$ ), and the reward function  $\hat{f}$  is sampled from the GP prior with a noise variance of  $\xi^2$ .

**Assumption 2.**  $\mathcal{X}$  is arbitrary and the reward function  $\hat{f}$  has a bounded RKHS norm for the kernel  $k$ , i.e.,  $\|f\|_k \leq B$ . The reward noises  $\epsilon_t$  form an arbitrary martingale difference sequence (i.e., reward noise does not systematically depend on its past values) and are uniformly bounded by  $\xi$ .

Under either Assumption 1 or 2, we prove the following regret bound for GP-TopK:

**Theorem 3.** If either Assumptions 1 or 2 hold, setting  $\beta_t$  as  $2 \log \left( \frac{|\mathcal{C}| \cdot |\Pi^k| \cdot t^2 \cdot \pi^2}{6\delta} \right)$  and  $300\gamma_t \ln^3 \left( \frac{t}{\delta} \right)$  respectively, the cumulative regret  $\mathcal{R}_T$  of the GP-TopK bandit algorithm for top- $k$  recommendations can, with at least  $1 - \delta$  probability, be bounded by  $\tilde{\mathcal{O}}(n\sqrt{C_1 T c (\log|\mathcal{C}| + k + \log(T^2 \pi^2 / 6\delta))})$  under Assumption 1, and  $\tilde{\mathcal{O}}(n\sqrt{C_1 (2B^2 c + 300n^2 c^2 \ln^3(T/\delta))T})$  under Assumption 2. Here,  $C_1 = \frac{8}{\log(1+\xi^{-2})}$ , and  $\tilde{\mathcal{O}}$  excludes logarithmic factors related to  $n$ ,  $k$ , and  $T$ .

Appendix B.4 provides the proof, leveraging the insight that  $\log \det |I + \sigma^{-2} K_{X_T}|$  for any set  $X_T$  can be effectively bounded using the finite-dimensional feature vectors introduced in this work. Specifically, Proposition 2 utilizes the feature vectors from Section 2. Theorem 3 establishes that the cumulative regret of the GP-TopK bandit algorithm grows sublinearly in  $T$  with high probability under both assumptions. It also highlights the importance of using top- $k$  ranking kernels, which improve the asymptotic order  $\tilde{\mathcal{O}}$  concerning  $n$  to  $n^{k/2-1}$  and  $n^{k-1}$  under Assumptions 1 and 2, respectively, as compared with Srinivas et al. [28]. This improvement is significant even for small constant values of  $k$ , such as  $k = 6$ , as detailed below in Table 3.

**Table 3:** Comparison with Srinivas et al. (2010) for regret bounds of the bandit algorithm under both assumptions.

	Srinivas et al. (2010)	This work
Assumption 1	$\tilde{\mathcal{O}} \left( n^{\frac{k}{2}} \sqrt{C_1 T c (\log \mathcal{C}  + k + \log(T^2 \pi^2 / 6\delta))} \right)$	$\tilde{\mathcal{O}} \left( n \sqrt{C_1 T c (\log \mathcal{C}  + k + \log(T^2 \pi^2 / 6\delta))} \right)$
Assumption 2	$\tilde{\mathcal{O}} \left( n^{\frac{k}{2}} \sqrt{C_1 T c (2B^2 + 300n^2 c \ln^3(T/\delta))} \right)$	$\tilde{\mathcal{O}} \left( n \sqrt{C_1 T c (2B^2 + 300n^2 c \ln^3(T/\delta))} \right)$

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## 258 5 Experiments

259 This section empirically evaluates the proposed GP-TopK bandit algorithms for the top- $k$  recom-  
 260 mendations using a simulation based on the MovieLens dataset [4]. The reliance on simulation for

evaluating bandit algorithms is prevalent in the literature. It stems from the difficulty of conducting online evaluations in real-world bandit scenarios, mainly when there are combinatorial arms [29]. Next, we provide details of the simulation setup and considered reward settings. Following that, we present results for the empirical regret for small and large numbers of arms below, respectively.

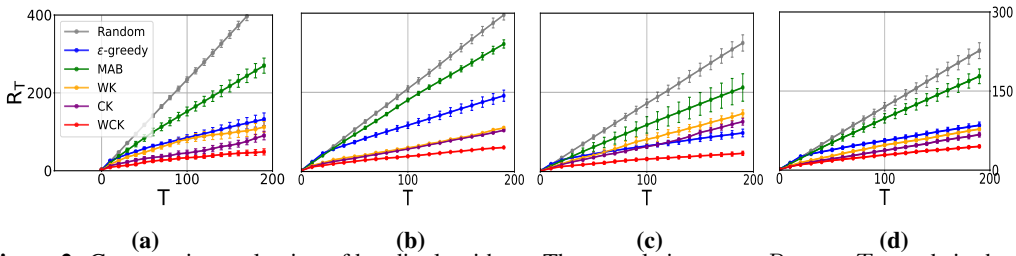
**Simulation setup and reward settings.** The bandit simulation setup follows the framework outlined by Jeunen et al. [8], utilizing real-world datasets on user-item interactions. Specifically, we train user and item embeddings using a collaborative filtering approach [6]. The user embeddings are accessed by the bandit algorithms as context embeddings, while the item embeddings remain hidden. In the non-contextual setup, the first user from the dataset is chosen as a fixed context throughout the bandit algorithm run, allowing us to use the same reward functions as the contextual bandit algorithm.

For setting up the reward functions, we utilize a similarity function  $s(\mathbf{c}, \theta) := \sigma(a \cdot (\mathbf{c}^T \theta) - b)$  to measure similarity between any user and item embeddings, where  $a$  and  $b$  are similarity score and shift scalars, respectively. The sigmoid function  $\sigma$  maps similarity scores to a range between 0 and 1, enhancing the interpretability of the reward signal [32]. We set  $a$  and  $b$  to 6 and 0.3, respectively, to fully utilize the range of the similarity function, as assessed by evaluating its value for many arms.

We set up two preliminary reward functions based on the similarity function  $s$ . The first is the DCG metric,  $\hat{f}_{\text{dcg}}(\mathbf{c}, \pi) = \sum_{i=1}^k \frac{1}{\log_2(i+1)} s(\mathbf{c}, \theta_{\pi_i})$ , where  $\mathbf{c}$  and  $\theta_{\pi_i}$  represent the context and item embeddings, respectively. The second is the diversity measure,  $\hat{f}_{\text{div}}(\pi) = \frac{1}{k^2} \sum_{i=1}^k \sum_{j=1}^k \theta_{\pi_j}^T \theta_{\pi_i}$ . These metrics quantify the relevance and diversity of top- $k$  recommendations, respectively.

We use these functions in two contextual reward settings. The first setting focuses on normalized-DCG (n-DCG),  $\hat{f}_{\text{ndcg}}(\mathbf{c}, \pi) = \frac{\hat{f}_{\text{dcg}}(\mathbf{c}, \pi)}{\max_{\pi'} \hat{f}_{\text{dcg}}(\mathbf{c}, \pi')}$  [7]. The second setting combines  $\hat{f}_{\text{ndcg}}$  and  $\hat{f}_{\text{div}}$  as  $\hat{f}_{\text{ndcgdiv}}(\mathbf{c}, \pi) = \lambda \cdot \hat{f}_{\text{ndcg}}(\mathbf{c}, \pi) + (1 - \lambda) \cdot \hat{f}_{\text{div}}(\pi)$ , evaluating the aggregate effect of relevance and diversity. We set  $\lambda = 0.25$  to emphasize relevance over diversity.

**Evaluation for small arm space.** This section presents empirical results for the cumulative regret of bandit algorithms with a limited number of arms. Specifically, with  $n = 20$  and  $k = 3$ , there are 6,840 top- $k$  rankings, allowing for an exhaustive search to optimize the acquisition function. All bandit algorithms run in batch mode, updating every five rounds. We consider both reward settings for contextual and non-contextual scenarios, using a subset of five users for the contextual setting. Several baselines are set to assess the benefits of ranking (Kendall) kernels. Section C details the remaining hyper-parameter configurations and details of other baseline bandit algorithms.



**Figure 2:** Comparative evaluation of bandit algorithms: The cumulative regret  $R_T$  over  $T$  rounds is shown. Lower values indicate better performance. Plots (a) and (b) represent non-contextual settings for nDCG ( $\hat{f}_{\text{ndcg}}$ ) and nDCG + diversity ( $\hat{f}_{\text{ndcgdiv}}$ ) rewards, respectively. Plots (c) and (d) show results for contextual settings for five users using the same rewards. The y-axis for (a) and (b) is on the left, and for (c) and (d) on the right. The GP-TopK algorithm with Kendall kernels, especially the weighted convolutional Kendall (WCK) kernel, outperforms others. Details on other algorithms are in the text. Results are averaged over six trials.

The *Random* algorithm randomly recommends any  $k$  items. The  $\epsilon$ -greedy algorithm alternates between recommending a random top- $k$  ranking with a probability of  $\epsilon$  and choosing the top- $k$  ranking with the highest observed mean reward. In contextual settings,  $\epsilon$ -greedy differentiates arms for each unique context. Similarly, *MAB-UCB* conceptualizes each ranking as an independent arm, an equivalent of using a direct delta kernel approach for GPs along with UCB  $\mathcal{AF}$ . In contextual scenarios, *MAB-UCB* also treats arms distinctly per context. Each variant of the top- $k$  ranking kernel yields one variation of the proposed GP-TopK algorithm, namely, WK, CK, and WCK. Figure 2

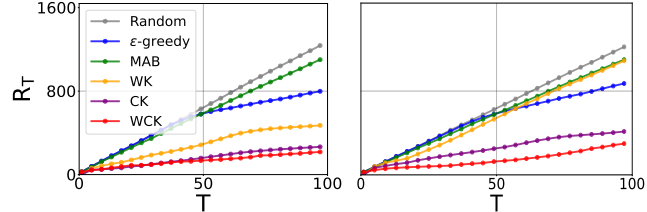


presents empirical values of the cumulative regrets for the above baseline and the proposed GP-TopK algorithms. In all cases, across both reward settings and in both contextual and non-contextual setups, the variants of the proposed GP-TopK algorithm outperform baselines that do not use Kendall kernels, highlighting the significance of top-k ranking kernels for full bandit feedback. Specifically, the CK and WCK kernels significantly outperform the WK kernel regarding the converged values of the regret, with the WCK kernel further improving on the CK kernel variant.

#### Evaluation for large arm space.

We evaluate bandit algorithms in a large arm space scenario with  $n = 50$  and  $k = 3$  and  $k = 5$ , resulting in  $1.1 \times 10^5$  and  $1.1 \times 10^{10}$  possible top-k rankings, respectively. Using local search, we focus on the nDCG reward. The remaining configuration is consistent with the small arm space setup. We use 10 restarts and 5 steps in each search direction for the local search, starting with 1000 initial candidates.

Figure 3 shows that the regret for the GP-TopK variants remains consistently lower even with a large arm space, despite the use of local search. The WCK approach significantly outperforms the CK, especially for  $k = 5$ , as illustrated in the right plot of Figure 3. Additional empirical results on the effectiveness of local search in a large arm space and other rewards are given in Appendix C.



**Figure 3:** Comparative evaluation of bandit algorithms for large arm spaces, with  $> 1.1 \times 10^5$  arms for the left plot and  $> 1.1 \times 10^{10}$  arms for the right plot. Cumulative regret with respect to the rounds of the bandit algorithm is depicted. Results are averaged over six trials. In both settings, the WCK approach outperforms other baselines. For more details, see the textual description.

## 6 Discussion

This work develops a contextual bandit algorithm for top-k recommendations using Gaussian processes with Kendall kernels in a full-bandit feedback setting. We make no restrictive assumptions about feedback or reward models. Gaussian processes allow computationally free model updates for accumulated feedback data, although inference remains challenging. We address this by providing features for Kendall kernels for top-k rankings, resulting in a faster inference algorithm that reduces the complexity from  $\mathcal{O}(T^4)$  to  $\mathcal{O}(T^2)$ . Additionally, we address issues with known variants and propose a more expressive Kendall kernel for top-k recommendations. Finally, we present theoretical and empirical results on cumulative regret to evaluate the proposed GP-TopK bandit algorithm.

**Future Directions and Limitations.** This work opens several research avenues. Efficient matrix-vector multiplication with Kendall kernel matrices can enable faster bandit algorithms with various acquisition functions, such as Thompson sampling and expected improvement. Exploring other kernels, like Mallow kernels, for top-k rankings and developing efficient algorithms for them is an intriguing direction, especially since the effectiveness of our algorithm depends on the function space induced by the RKHS of the underlying kernel. Assessing how well these kernels approximate various reward functions for top-k recommendations would provide valuable insights.

Exploring other bandit problem settings, such as stochastic item availability or delayed feedback, would enhance the applicability of this work to more complex scenarios. Extending the finite-dimensional GP framework to other acquisition functions using local search is another promising direction. One limitation of our regret analysis is that it does not account for approximations in the arm selection step due to local search [20]. This limitation is common in continuous domains, where optimizing acquisition functions often involves non-convex optimization [28].

**Impact.** This research advances bandit algorithms for top-k item recommendations. By improving recommendation efficiency and accuracy, our algorithms can enhance user experiences across platforms, promoting content relevancy and engagement. However, they may reinforce implicit biases in training data, limiting content diversity and entrenching prejudices. Therefore, monitoring over time is essential when deploying these algorithms in real-world environments.

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## A Kendall Kernels for Full and Top-k Rankings – Omitted Details

This section includes the proofs that were omitted from Section 2, presented in the following order:

- In Section A.1, we present proofs for Claims 2 and 3, which concern the feature representations of Kendall kernels.
- In Section A.2, we provide Algorithms 2 and a proof of its correctness for computing the WCK kernel in  $\mathcal{O}(k^2)$  time, thereby proving Claim 1. Additionally, we extend this proof to cover the proof of correctness for Algorithm 3, which can compute the CK kernel in  $\mathcal{O}(k \log k)$ , initially introduced by Jiao et al. [9]. The original paper presented the algorithm without a formal proof of correctness, a gap we address and fill in this section.
- Section A.3 details the proof for Theorem 1, discussing the matrix-vector multiplications with the Kendall kernel matrix for top-k rankings. This proof builds on the Algorithm 2 given for computing the WCK kernel for top-k rankings.

### A.1 Feature Representation for Kendall Kernels for Top-k Rankings

This section revisits the claims regarding the feature representations of the weighted Kendall kernel and the weighted convolutional Kendall kernel, subsequently providing the proofs for these claims mentioned earlier.

**Claim 2.** Let  $\phi^{wk}(\sigma) : \Pi^n \mapsto \mathbb{R}^{\binom{n}{2}}$  be a vector indexed by unique item pairs  $(i, j)$ , defined as:

$$\phi_{i,j}^{wk}(\sigma) := \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_i^\sigma, \mathbf{p}_j^\sigma) \cdot (\mathbf{p}_{i<j}^\sigma - \mathbf{p}_{i>j}^\sigma),$$

where  $w_s$  is the symmetric weighting function in product-symmetric weights. Then,  $\phi^{wk}$  is a linear feature vector for the weighted Kendall kernel with product-symmetric weights  $w_{ps}$ .

*Proof.* Following the definition of linear feature representation, we need to prove that  $k^{wk}(\sigma_1, \sigma_2) = \phi(\sigma_1)^T \phi(\sigma_2)$  for the product-symmetric weight kernel as given in Equation 5. Recalling from Equation 2, we have  $k^{wk}(\sigma_1, \sigma_2)$  as follows:

$$\begin{aligned} k^{wk}(\sigma_1, \sigma_2) &= \frac{1}{\binom{n}{2}} \cdot \sum_{i<j} w((\mathbf{p}_i^{\sigma_1}, \mathbf{p}_j^{\sigma_1}), (\mathbf{p}_i^{\sigma_2}, \mathbf{p}_j^{\sigma_2})) \cdot \eta_{i,j}(\sigma_1, \sigma_2), \\ &= \frac{1}{\binom{n}{2}} \cdot \sum_{i<j} w_s(\mathbf{p}_i^{\sigma_1}, \mathbf{p}_j^{\sigma_1}) \cdot w_s(\mathbf{p}_i^{\sigma_2}, \mathbf{p}_j^{\sigma_2}) \cdot \eta_{i,j}(\sigma_1, \sigma_2), \end{aligned} \quad (7)$$

where the second line incorporates the use of the product-symmetric weight kernel. Next, our focus shifts to the simplification of  $\eta_{i,j}(\sigma_1, \sigma_2)$ , which is elaborated as follows:

$$\begin{aligned} \eta_{i,j}(\sigma_1, \sigma_2) &= \mathbf{p}_{i<j}^{\sigma_1} \cdot \mathbf{p}_{i<j}^{\sigma_2} + \mathbf{p}_{i>j}^{\sigma_1} \cdot \mathbf{p}_{i>j}^{\sigma_2} - \mathbf{p}_{i<j}^{\sigma_1} \cdot \mathbf{p}_{i>j}^{\sigma_2} - \mathbf{p}_{i>j}^{\sigma_1} \cdot \mathbf{p}_{i<j}^{\sigma_2}, \\ &= \mathbf{p}_{i<j}^{\sigma_1} \cdot (\mathbf{p}_{i<j}^{\sigma_2} - \mathbf{p}_{i>j}^{\sigma_2}) + \mathbf{p}_{i>j}^{\sigma_1} \cdot (\mathbf{p}_{i>j}^{\sigma_2} - \mathbf{p}_{i<j}^{\sigma_2}), \\ &= (\mathbf{p}_{i<j}^{\sigma_1} - \mathbf{p}_{i>j}^{\sigma_1}) \cdot (\mathbf{p}_{i<j}^{\sigma_2} - \mathbf{p}_{i>j}^{\sigma_2}). \end{aligned}$$

Combining the above factorization of  $\eta_{i,j}$  with Equation 7, we get:

$$k^{wk}(\sigma_1, \sigma_2) = \frac{1}{\binom{n}{2}} \cdot \sum_{i<j} w_s(\mathbf{p}_i^{\sigma_1}, \mathbf{p}_j^{\sigma_1}) \cdot w_s(\mathbf{p}_i^{\sigma_2}, \mathbf{p}_j^{\sigma_2}) \cdot (\mathbf{p}_{i<j}^{\sigma_1} - \mathbf{p}_{i>j}^{\sigma_1}) \cdot (\mathbf{p}_{i<j}^{\sigma_2} - \mathbf{p}_{i>j}^{\sigma_2})$$

$$\begin{aligned}
&= \frac{1}{\binom{n}{2}} \cdot \sum_{i < j} \phi_{i,j}^{wk}(\sigma_1) \cdot \phi_{i,j}^{wk}(\sigma_2) \\
&= \phi(\sigma_1)^T \phi(\sigma_2).
\end{aligned}$$

459

□

**Claim 3.** Let  $\phi^{wck}(\pi) : \Pi^k \mapsto \mathbb{R}^{\binom{n}{2}}$  be a vector indexed by unique item pairs  $(i, j)$  given as:  $\phi_{i,j}^{wck}(\pi) := \frac{1}{\sqrt{\binom{n}{2}}} \cdot \mathbf{w}_{i,j}^{wck}(\pi) \cdot (\mathbf{p}_{i < j}^\pi - \mathbf{p}_{i > j}^\pi)$ , where  $\mathbf{w}_{i,j}^{wck}(\pi)$  is determined as follows:

$$\mathbf{w}_{i,j}^{wck}(\pi) = \begin{cases} w_s(\mathbf{p}_i^\pi, \mathbf{p}_j^\pi) & \text{if } \mathbf{p}_i^\pi \in [k] \text{ \& } \mathbf{p}_j^\pi \in [k] \\ w_s(\mathbf{p}_i^\pi, \cdot) & \text{else if } \mathbf{p}_i^\pi \in [k] \text{ \& } \mathbf{p}_j^\pi \notin [k], \\ w_s(\cdot, \mathbf{p}_j^\pi) & \text{else if } \mathbf{p}_i^\pi \notin [k] \text{ \& } \mathbf{p}_j^\pi \in [k], \\ 0 & \text{otherwise,} \end{cases}$$

where  $w_s$  denotes symmetric weights and  $w_s(\ell, \cdot) = \frac{1}{n-k} \sum_{j=k+1}^n w_s(\ell, j)$ . Then, the vector  $\phi^{wck}$  is a linear feature vector for the WCK kernel  $k^{wck}$ . By uniformly setting  $w_s(\cdot, \cdot) \equiv 1$  in the definitions above,  $\phi_{i,j}^{wck}(\pi)$  specializes to a linear feature vector for the CK kernel.

460 *Proof.* The main idea revolves around leveraging the feature representation of the Weighted Kendall  
461 kernel for a full ranking and the linearity of the convolution operation. It is already established  
462 that  $k^{wk}(\sigma_1, \sigma_2) = \phi^{wk}(\sigma_1)^T \phi^{wk}(\sigma_2)$ , as demonstrated in Claim 2. Recall that the WCK kernel  
463 requires a double summation over pairs of rankings from  $B_{\pi_1}$  and  $B_{\pi_2}$ , which represent the sets of  
464 full rankings consistent with their respective top-k rankings, as described in Equation 4. We simplify  
465 the WCK kernel as follows:

$$\begin{aligned}
k^{wck}(\pi_1, \pi_2) &= \frac{1}{|B_{\pi_1}|} \cdot \frac{1}{|B_{\pi_2}|} \cdot \sum_{\sigma_1 \in B_{\pi_1}} \sum_{\sigma_2 \in B_{\pi_2}} \phi^{wk}(\sigma_1)^T \phi^{wk}(\sigma_2) \\
&= \left( \frac{1}{|B_{\pi_1}|} \cdot \sum_{\sigma_1 \in B_{\pi_1}} \phi^{wk}(\sigma_1)^T \right) \cdot \underbrace{\left( \frac{1}{|B_{\pi_2}|} \cdot \sum_{\sigma_2 \in B_{\pi_2}} \phi^{wk}(\sigma_2) \right)}_{:= \phi^{wck}(\pi_2)} \\
&= \phi^{wck}(\pi_1)^T \phi^{wck}(\pi_2).
\end{aligned}$$

466 The simplification above reveals that the feature representation,  $\phi^{wck}$ , for the WCK kernel, is a  $\binom{n}{2}$   
467 dimensional vector and can be indexed by unique pairs of items  $(i, j)$ , much like the  $\phi^{wk}$ . However,  
468 the double summation is over an exponentially large number of pairs of rankings. Moving forward,  
469 we shift our focus to the individual entries of this representation involving this summation, elucidating  
470 the analytical values within the summation by exploring four unique cases, each dependent on whether  
471 these specific items fall within the top-k rankings.

472 In Case 1, we examine the scenario where items  $i$  and  $j$  are within the top-k ranking  $\pi$ . Here, the  
473 focus is on the feature representation of the pair, specifically when both elements are ranked among  
474 the top-k positions.

475

476 **Case 1:**  $\mathbf{p}_i^\pi \in [k]$  and  $\mathbf{p}_j^\pi \in [k]$ .

$$\begin{aligned}
\phi_{i,j}^{wck}(\pi) &= \frac{1}{|B_\pi|} \cdot \sum_{\sigma \in B_\pi} \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_i^\sigma, \mathbf{p}_j^\sigma) \cdot (\mathbf{p}_{i < j}^\sigma - \mathbf{p}_{i > j}^\sigma) \\
&= \frac{1}{|B_\pi|} \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_i^\pi, \mathbf{p}_j^\pi) \cdot \left( \sum_{\sigma \in B_\pi} \mathbf{p}_{i < j}^\sigma - \sum_{\sigma \in B_\pi} \mathbf{p}_{i > j}^\sigma \right)
\end{aligned}$$

$$\begin{aligned}
&= \frac{1}{|B_\pi|} \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_i^\pi, \mathbf{p}_j^\pi) \cdot (|B_\pi| \cdot \mathbf{p}_{i < j}^\pi - |B_\pi| \cdot \mathbf{p}_{i > j}^\pi) \\
&= \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_i^\pi, \mathbf{p}_j^\pi) \cdot (\mathbf{p}_{i < j}^\pi - \mathbf{p}_{i > j}^\pi). \tag{8}
\end{aligned}$$

477 The simplification in lines 3rd and 4th follows from the fact that any full ranking  $\sigma \in B_\pi$ , consistent  
478 with the top-k ranking  $\pi$ , the relative ranks and weights of items  $i$  and  $j$  remains unchanged, given  
479  $\mathbf{p}_i^\pi \in [k]$  and  $\mathbf{p}_j^\pi \in [k]$ . Concretely, this implies  $\mathbf{p}_{i < j}^\sigma = \mathbf{p}_{i < j}^\pi$  for all  $\sigma \in B_\pi$  and similar with the  
480 other term.

481 In Case 2, we analyze when item  $i$  is in the top-k ranking while item  $j$  is not.

482

483 Case 2:  $\mathbf{p}_i^\pi \in [k]$  and  $\mathbf{p}_j^\pi \notin [k]$ .

$$\begin{aligned}
\phi_{i,j}^{wck}(\pi) &= \frac{1}{|B_\pi|} \cdot \sum_{\sigma \in B_\pi} \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_i^\sigma, \mathbf{p}_j^\sigma) \cdot (\mathbf{p}_{i < j}^\sigma - \mathbf{p}_{i > j}^\sigma) \\
&= \frac{1}{|B_\pi|} \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot \sum_{\sigma \in B_\pi} w_s(\mathbf{p}_i^\sigma, \mathbf{p}_j^\sigma) \cdot (1 - 0) \quad (\text{since } \mathbf{p}_i^\pi \in [k] \text{ and } \mathbf{p}_j^\pi \notin [k]) \\
&= \frac{1}{|B_\pi|} \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot \sum_{\sigma \in B_\pi} w_s(\mathbf{p}_i^\pi, \mathbf{p}_j^\sigma).
\end{aligned}$$

484 Next, every possible consistent ranking is considered jointly while fixating on a specific rank outside  
485 top-k elements, leading to  $(n - k - 1)!$  different rankings. Given that  $|B_\pi| = (n - k)!$ , we can refine  
486 the above expression as follows:

$$\begin{aligned}
\phi_{i,j}^{wck}(\pi) &= \frac{1}{|B_\pi|} \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot \sum_{l=k+1}^n w_s(\mathbf{p}_i^\pi, l) \cdot (n - k - 1)! \\
&= \frac{(n - k - 1)!}{|B_\pi|} \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot \sum_{l=k+1}^n w_s(\mathbf{p}_i^\pi, l) \\
&= \frac{1}{\sqrt{\binom{n}{2}}} \cdot \frac{1}{n - k} \cdot \sum_{l=k+1}^n w_s(\mathbf{p}_i^\pi, l) \\
&= \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_i^\pi, \cdot). \tag{9}
\end{aligned}$$

487 In Case 3, we analyze when item  $i$  is not in the top-k ranking while item  $j$  is.

488

489 Case 3:  $\mathbf{p}_i^\pi \notin [k]$  and  $\mathbf{p}_j^\pi \in [k]$ . Similar to case 2, the simplification follows analogously, with the  
490 only change being  $\mathbf{1}_{\mathbf{p}_{i < j}^\sigma} - \mathbf{1}_{\mathbf{p}_{i > j}^\sigma} = -1$  instead of 1. Thus, by symmetry between  $i$  and  $j$ , we have  
491 the following:

$$\phi_{i,j}^{wck}(\pi) = \frac{-1}{\sqrt{\binom{n}{2}}} \cdot w_s(\cdot, \mathbf{p}_j^\pi) = \frac{-1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_j^\pi, \cdot) \quad (\text{using symmetry of } w_s). \tag{10}$$

492 Lastly, in Case 4, we analyze when items  $i$  and  $j$  are not in the top-k ranking.

493 Case 4:  $\mathbf{p}_i^\sigma \notin [k]$  and  $\mathbf{p}_j^\sigma \notin [k]$ .

$$\phi_{i,j}^{wck}(\pi) = \frac{1}{|B_\pi|} \cdot \sum_{\sigma \in B_\pi} \phi_{i,j}^{wk}(\sigma)$$

$$\begin{aligned}
&= \frac{1}{|B_\pi|} \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot \sum_{\sigma \in B_\pi} w_s(\mathbf{p}_i^\sigma, \mathbf{p}_j^\sigma) \cdot (\mathbf{p}_{i < j}^\sigma - \mathbf{p}_{i > j}^\sigma) \\
&= 0 \quad (\text{by symmetry}).
\end{aligned} \tag{11}$$

The result of zero arises from symmetry. Since  $\mathbf{p}_i^\sigma$  and  $\mathbf{p}_j^\sigma$  are not in the top- $k$  ranking, they are treated symmetrically in the summation overall rankings in  $B_\pi$ . For any ranking  $\sigma$ , suppose  $\mathbf{p}_i^\sigma = l$  and  $\mathbf{p}_j^\sigma = m$ , there exists a corresponding ranking  $\sigma'$  such that only the items  $i$  and  $j$  are swapped. Therefore, jointly, these two rankings yield  $w_s(l, m)$  and  $-w_s(l, m)$ . Since  $w_s$  is symmetric, the overall contribution from each pair of such rankings is zero. Hence, the entire summation nets to zero.

Thus, with the explanation provided for each case and combining results from Equations 8, 9, 10 and 11, it's trivial to validate the Claim 3, i.e.,  $\phi_{i,j}^{wck}(\pi) = \frac{1}{\sqrt{\binom{n}{2}}} \cdot \mathbf{w}_{i,j}^{wck}(\pi) \cdot (\mathbf{p}_{i < j}^\pi - \mathbf{p}_{i > j}^\pi)$  for all unique pair of items. From Case 4, we have  $\mathcal{O}((n-k)^2)$  entries leaving at max only  $\mathcal{O}(k^2 + 2nk)$  non-zero entries.  $\square$

## A.2 Algorithms for Computing Kendall Kernels for top-k Rankings

In this section, we provide and delve into the proofs of Algorithms 2 and 3 for the weighted convolutional Kendall kernel and the convolutional Kendall kernel, as previously discussed in Section 2. Section A.2.1 for valid both the correctness and computational complexity of Algorithm 2 as given earlier in Claim 1. Following this, Section A.2.2 revisits Algorithm 3, initially introduced by Jiao et al. [10]. The original publication presented the algorithm without formal proof of its correctness, which we rectify and offer in Section A.2.2.

### A.2.1 Efficiently Computing the Weighted Convolutional Kendall Kernel

This section provides a proof to Claim 1 to establish the efficiency and accuracy of Algorithm 2 in computing the weighted convolutional Kendall kernel, as specified in Equation 4, with a focus on its computational complexity.

**Claim 1.** *The weighted convolutional Kendall kernel (Equation 4) with product-symmetric rank weights (Equation 5) can be computed in  $\mathcal{O}(k^2)$  time.*

*Proof.* The claim is proven through Algorithm 2, where we establish its correctness and demonstrate its computation requirement is  $\mathcal{O}(k^2)$ . The essence of our proof centers on analyzing the feature representation of the WCK kernel,  $\phi^{wck}$ , as outlined in Claim 3. The feature vectors of  $\phi^{wck}$  reside in a  $\binom{n}{2}$  dimensional space, indexed by pairs of items. Our approach is to demonstrate that Algorithm 2 accurately computes the right-hand side (RHS) of the equation  $k^{wck}(\pi_1, \pi_2) = \phi^{wck}(\pi_1)^T \phi^{wck}(\pi_2)$ . This involves a summation over item pairs, expressed as  $k^{wck}(\pi_1, \pi_2) = \sum_{l < m} \phi_{l,m}^{wck}(\pi_1)^T \phi_{l,m}^{wck}(\pi_2)$ .

Our proof analyzes various scenarios: cases where pairs of items, namely  $l$  and  $m$ , fall within the top- $k$ , scenarios with one item within the top- $k$  and the other outside, and situations where neither item is within the top- $k$ . Each of these cases contributes distinctively to the computation of the overall kernel, resulting in different terms in the algorithmic computation. This is encapsulated in Algorithm 2, where  $k^{wck}(\pi_1, \pi_2) = \sum_{i=1}^5 s_i(\pi_1, \pi_2)$ , and each  $s_i$  corresponds to the terms given earlier in Algorithm 2 from Section 2.

Before proceeding with the cases of this summation as given in Table 4, we recall the notations utilized by Algorithm 3 in Definition 1. Also, remember that we will be proving for product-symmetric weights as given in Equation 5, where,  $w_s : [n] \times [n] \mapsto \mathbb{R}^n$  and its one-dimensional marginals are  $w_s(\ell, \cdot) = \frac{1}{n-k} \sum_{j=k+1}^n w_s(\ell, j)$  Table 4 shows how these cases are organized and relate to different  $s_i$  terms required for computing the WCK kernel. The key strategy involves breaking down the kernel's computation into cases based on the positioning of item pairs within the top- $k$  rankings. In case 1, we consider all the scenarios when both indices are within the set of items in both top- $k$  rankings, i.e., all items in the set  $I_1 \cup I_2$ .

Case	Description
1	Both items $(l, m)$ in $I_1 \cup I_2$ . Branches into the following three sub-cases based on the presence of items in $I_1 \cap I_2$ : 1-a: Both items in $I_1 \cap I_2$ . The concerned term is $s_1$ . 1-b: One item in $I_1 \cap I_2$ . Subdivided into 1-b-i (other in $I_1 \setminus I_2$ ) and 1-b-ii (other in $I_2 \setminus I_1$ ); concerned terms are $s_2$ and $s_3$ . 1-c: No item in $I_1 \cap I_2$ . Addresses cases where $l$ and $m$ are in different sets ( $I_1 \setminus I_2$ and $I_2 \setminus I_1$ ); concerned term is $s_4$ .
2	One item in $I_1 \cup I_2$ . I.e., either $l$ is in $I_1 \cup I_2$ or $m$ is in $I_1 \cup I_2$ , leading to sub-cases 2-a and 2-b; concerned term is $s_5$ .
3	No item in $I_1 \cup I_2$ . Addresses the scenario where neither $l$ nor $m$ is in $I_1 \cup I_2$ ; value trivially zero.

**Table 4:** Case categorization for the proof of Algorithms 2 and 3 based on item pair ranks, where  $I_1$  and  $I_2$  are the sets of items for top-k rankings  $\pi_1$  and  $\pi_2$ , respectively.

**Definition 1.** Algorithm 2 and 3 and utilize following notations.

- $I_1$  and  $I_2$  are the sets of items in rankings  $\pi_1$  and  $\pi_2$ , respectively.
- $\sigma_1 \in \Pi^{|I_1|}$  and  $\tau_1 \in \Pi^{|I_1 \cap I_2|}$  are the full rankings of  $I_1$  and  $I_1 \cap I_2$ , both consistent with the input top-k ranking  $\pi_1$ . I.e., relative ranks of items is same yielding  $\forall l, m \in I_1 \cap I_2, p_{i>j}^{\pi_1} = p_{i>j}^{\tau_1}$ .
- Analogously,  $\sigma_2$  and  $\tau_2$  are constructed utilizing the set  $I_2$  and ranking  $\pi_2$ .

**Algorithm 2** Computing Weighted Convolutional Kendall Kernel

**Input:** Two permutations  $\pi_1, \pi_2 \in \Pi^k$ . Ranking weighting function  $w_s : [n] \times [n] \mapsto \mathbb{R}^n$  and its one dimensional marginals are  $w_s(\ell, \cdot) = \frac{1}{n-k} \sum_{j=k+1}^n w_s(\ell, j)$ .

**Output:** Convolutional Weighted Kendall kernel  $k^{wck}(\pi_1, \pi_2)$ .

– Let  $I_1$  and  $I_2$  be the sets of items in rankings  $\pi_1$  and  $\pi_2$ , respectively.

```

1: if  $|I_1 \cap I_2| \geq 2$  then
2:    $s_1(\pi_1, \pi_2) = \left(\frac{1}{n}\right) \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} w_s(p_l^{\pi_1}, p_m^{\pi_1}) \cdot w_s(p_l^{\pi_2}, p_m^{\pi_2}) \cdot \eta_{l,m}(\pi_1, \pi_2)$ 
3: end if
4: if  $|I_1 \cap I_2| \geq 1$  and  $|I_1 \setminus I_2| \geq 1$  then
5:    $s_2(\pi_1, \pi_2) = \left(\frac{1}{n}\right) \cdot \sum_{l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} w_s(p_l^{\pi_1}, p_m^{\pi_1}) \cdot w_s(p_l^{\pi_2}, \cdot) \cdot (p_{l < m}^{\pi_1} - p_{l > m}^{\pi_1})$ 
6: end if
7: if  $|I_1 \cap I_2| \geq 1$  and  $|I_2 \setminus I_1| \geq 1$  then
8:    $s_3(\pi_1, \pi_2) = \left(\frac{1}{n}\right) \cdot \sum_{l \in I_1 \cap I_2 | m \in I_2 \setminus I_1} w_s(p_l^{\pi_1}, \cdot) \cdot w_s(p_l^{\pi_2}, p_m^{\pi_2}) \cdot (p_{l < m}^{\pi_2} - p_{l > m}^{\pi_2})$ 
9: end if
10: if  $|I_1 \setminus I_2| \geq 1$  and  $|I_2 \setminus I_1| \geq 1$  then
11:    $s_4(\pi_1, \pi_2) = -\left(\frac{1}{n}\right) \cdot \sum_{l \in I_1 \setminus I_2 | m \in I_2 \setminus I_1} w_s(p_l^{\pi_1}, \cdot) \cdot w_s(p_m^{\pi_2}, \cdot)$ 
12: end if
13: if  $|I_1 \cap I_2| \geq 1$  and  $|[n] \setminus (I_1 \cup I_2)| \geq 1$  then
14:    $s_5(\pi_1, \pi_2) = \left(\frac{1}{n}\right) \cdot (n - |I_1 \cup I_2|) \cdot \sum_{l \in I_1 \cap I_2} w_s(p_l^{\pi_1}, \cdot) \cdot w_s(p_l^{\pi_2}, \cdot)$ 
15: end if
16:  $k^{wck}(\pi_1, \pi_2) = s_1(\pi_1, \pi_2) + s_2(\pi_1, \pi_2) + s_3(\pi_1, \pi_2) + s_4(\pi_1, \pi_2) + s_5(\pi_1, \pi_2)$ 

```

536 **Case 1:** The pair  $(l, m) \in I_1 \cup I_2$  falls within the top-k, leading to three distinct cases. Below, we  
537 provide  $s_i$  terms for each case as given in Table 4.



538 **Case 1-a:** Two items in  $I_1 \cap I_2$ , meaning both  $l$  and  $m$  belong to  $I_1 \cap I_2$ . Using Claim 3 regarding  
 539 the feature vector  $\phi^{wck}$ , we simplify  $s_1$  as follows:

$$\begin{aligned}
 s_1(\pi_1, \pi_2) &= \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} \phi_{l,m}^{wck}(\pi_1) \cdot \phi_{l,m}^{wck}(\pi_2) \\
 &= \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \\
 &\quad \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2}) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}) \\
 &= \frac{1}{\binom{n}{2}} \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2}) \cdot \eta_{l,m}(\pi_1, \pi_2). \quad (12)
 \end{aligned}$$

540 **Case 1-b:** When one item is in  $I_1 \cap I_2$ , the other must reside either in  $I_1 \setminus I_2$  or  $I_2 \setminus I_1$ , thus leading  
 541 to two distinct sub-cases. This is specified in Table 4. Concretely, if the other item is in  $I_1 \setminus I_2$ , it  
 542 contributes to the  $s_2$  terms, whereas if it's in  $I_2 \setminus I_1$ , it contributes to the  $s_3$  terms.

543 Corresponding to Case 1-b-i, when the other item is in  $I_1 \cap I_2$ , i.e.,  $s_2$  is the term corresponding to  
 544 indices where  $l$  is in  $I_1 \cap I_2$  and  $m$  in  $I_1 \setminus I_2$ , or the reverse, represented by partial sums  $u$  and  $v$ . For  
 545 the partial sum  $u$ , with  $l$  in  $I_1 \cap I_2$  and  $m$  in  $I_1 \setminus I_2$ , we find that  $\mathbf{p}_l^{\pi_2}$  is in  $[k]$ , while  $\mathbf{p}_m^{\pi_2}$  is not. The  
 546 simplification of  $u$  proceeds using Claim 3 as follows:

$$\begin{aligned}
 u &= \sum_{1 \leq l < m \leq n | l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \\
 &\quad \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}) \\
 &= \frac{1}{\binom{n}{2}} \sum_{1 \leq l < m \leq n | l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}).
 \end{aligned}$$

547 Similarly, the partial sum  $v$  can be simplified as follows:

$$\begin{aligned}
 v &= \sum_{1 \leq l < m \leq n | m \in I_1 \cap I_2 | l \in I_1 \setminus I_2} \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \\
 &\quad \cdot \frac{-1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}) \\
 &= \frac{-1}{\binom{n}{2}} \sum_{1 \leq l < m \leq n | m \in I_1 \cap I_2 | l \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \\
 &= \frac{-1}{\binom{n}{2}} \sum_{1 \leq m < l \leq n | l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} w_s(\mathbf{p}_m^{\pi_1}, \mathbf{p}_l^{\pi_1}) \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot) \cdot (\mathbf{p}_{m < l}^{\pi_1} - \mathbf{p}_{m > l}^{\pi_1}) \\
 &= \frac{1}{\binom{n}{2}} \sum_{1 \leq m < l \leq n | l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}).
 \end{aligned}$$

548 In the above, the first two lines use results from Claim 3 and use similarity of  $w_s$ . In the following  
 549 line,  $l$  and  $m$  are exchanged. Lastly, the negative sign is pushed into the indicator functions to make  
 550 the summand function of this partial sum  $v$  similar to the partial sum  $u$ , and the similarity of the  $w_s$   
 551 is utilized. The above partial sums simplify  $s_2$  as follows:

$$s_2(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \cdot \sum_{l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}). \quad (13)$$

552 Analogously, in Case 1-b-ii, we deduce the corresponding term  $s_3$  for the pair of indices as described  
 553 in Table 4 through symmetry. Specifically, the term  $s_3$  can be outlined as follows:

$$s_3(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \cdot \sum_{l \in I_1 \cap I_2 | m \in I_2 \setminus I_1} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2}) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}). \quad (14)$$

554 **Case 1-c:** Both items are outside  $I_1 \cap I_2$ , specifically,  $l \in I_1 \setminus I_2$  and  $m \in I_2 \setminus I_1$  or the reverse.  
 555 Like Case 1-b-i, we divide  $s_4$  into partial summations  $u$  and  $v$ . Now, we calculate  $u$  under the  
 556 condition that  $l \in I_1 \setminus I_2$  and  $m \in I_2 \setminus I_1$ .

$$\begin{aligned} u &= \sum_{1 \leq l < m \leq n | l \in I_1 \setminus I_2 | m \in I_2 \setminus I_1} \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \\ &\quad \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}), \\ &= \frac{1}{\binom{n}{2}} \cdot \sum_{1 \leq l < m \leq n | l \in I_1 \setminus I_2 | m \in I_2 \setminus I_1} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot (1 - 0) \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot) \cdot (0 - 1), \\ &= \frac{-1}{\binom{n}{2}} \cdot \sum_{1 \leq l < m \leq n | l \in I_1 \setminus I_2 | m \in I_2 \setminus I_1} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot). \end{aligned}$$

557 Similarly, we can estimate partial sum  $v$  for the set  $l \in I_2 \setminus I_1$  &  $m \in I_1 \setminus I_2$ . Using calculations  
 558 similar to Case-1-b-i for summing  $u$  and  $v$ , we have:

$$s_4(\pi_1, \pi_2) = \frac{-1}{\binom{n}{2}} \cdot \sum_{l \in I_1 \setminus I_2 | m \in I_2 \setminus I_1} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot). \quad (15)$$

559 **Case 2:** One item exists in  $I_1 \cap I_2$ , the other in  $[n] \setminus (I_1 \cap I_2)$ . It branches into two sub-cases: Case  
 560 2-a with one item in  $I_1 \cup I_2$ , and Case 2-b, where one item outside  $I_1 \cap I_2$  but is in  $I_1 \cup I_2$ . Focusing  
 561 on Case 2-a, represented by  $s_5$ , we simplify as follows. This involves two index scenarios, either  
 562  $l \in I_1 \cap I_2$  and  $m \notin I_1 \cup I_2$  or vice versa, represented by partial sums  $u$  and  $v$ . We now simplify  $u$   
 563 below:

$$\begin{aligned} u &= \frac{1}{\binom{n}{2}} \sum_{1 \leq l < m \leq n | l \in I_1 \cap I_2 | m \notin I_1 \cup I_2} \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \\ &\quad \cdot \frac{1}{\sqrt{\binom{n}{2}}} \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}), \\ &= \frac{1}{\binom{n}{2}} \sum_{1 \leq l < m \leq n | l \in I_1 \cap I_2 | m \notin I_1 \cup I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot), \\ &= \frac{1}{\binom{n}{2}} \sum_{1 \leq l < m \leq n | l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot (n - |I_1 \cup I_2|). \end{aligned}$$

564 Using steps similar to the previous case, we get the following value for  $s_5$ :

$$s_5(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \cdot (n - |I_1 \cup I_2|) \cdot \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot). \quad (16)$$

565 For Case 2-b,  $l$  or  $m$  are absent from  $I_1$  or  $I_2$ , leading to two sub-scenarios. Consequently, either  
 566  $\phi_{l,m}^{wck}(\pi_1)$  is zero or  $\phi_{l,m}^{wck}(\pi_2)$  is zero. Therefore, these terms don't contribute to the overall WCK  
 567 kernel value.

568 **Case 3:** No item is in the top-k, i.e., both  $l, m \notin I_1 \cup I_2$ . As both items are absent from the top-k  
 569 in either ranking, the value trivially reduces to zero.

570 After covering all configurations of  $l$  and  $m$ , we incorporate results from Equations 12, 13, 14, 15,  
 571 and 16. This integration yields the expression  $k^{wck}(\pi_1, \pi_2) = \sum_{i=1}^5 s_i(\pi_1, \pi_2)$ , where, each term  $s_i$   
 572 matches precisely with its corresponding expression in Algorithm 2. The proof for the correctness of  
 573 Algorithm 2 is complete, as each term  $s_i$  corresponds to its respective expression in the algorithm.  
 574 Regarding the time complexity of Algorithm 2, each term  $s_i$  sums at most  $k^2$  quantities, and each  
 575 quantity summed can be computed in  $\mathcal{O}(1)$  time. Therefore, the computation time required for  
 576 Algorithm 2 is  $\mathcal{O}(k^2)$ .  $\square$

### 577 A.2.2 Efficiently Computing the Convolutional Kendall Kernel

578 This section provides Algorithm 3 for computing the convolutional Kendall kernel, as specified in  
 579 Equation 3. Later, its efficiency and accuracy are proved in Claim 4.

#### Algorithm 3 Computing Convolutional Kendall Kernel [10]

**Input:** Two top-k rankings  $\pi_1, \pi_2 \in \Pi^k$ .

**Output:** Convolutional Kendall kernel  $k^{ck}(\pi_1, \pi_2)$ .

- Let  $I_1$  and  $I_2$  be the sets of items in rankings  $\pi_1$  and  $\pi_2$ , respectively.
- Let  $\sigma_1 \in \Pi^{|I_1|}$  and  $\tau_1 \in \Pi^{|I_1 \cap I_2|}$  be the full rankings of  $I_1$  and  $I_1 \cap I_2$ , both consistent with the input top-k ranking  $\pi_1$ .
- Analogously, construct  $\sigma_2$  and  $\tau_2$  utilizing the set  $I_2$  and ranking  $\pi_2$ .

- 1: **if**  $|I_1 \cap I_2| \geq 2$  **then**
- 2:    $s_1(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \cdot \binom{|I_1 \cap I_2|}{2} \cdot k^{sk}(\tau_1, \tau_2)$
- 3: **end if**
- 4: **if**  $|I_1 \cap I_2| \geq 1$  and  $|I_1 \setminus I_2| \geq 1$  **then**
- 5:    $s_2(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \cdot \sum_{l \in I_1 \cap I_2} 2 \cdot (\sigma_1(l) - \tau_1(l)) - k + |I_1 \cap I_2|$
- 6: **end if**
- 7: **if**  $|I_1 \cap I_2| \geq 1$  and  $|I_2 \setminus I_1| \geq 1$  **then**
- 8:    $s_3(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \cdot \sum_{l \in I_1 \cap I_2} 2 \cdot (\sigma_2(l) - \tau_2(l)) - k + |I_1 \cap I_2|$
- 9: **end if**
- 10:  $s_4(\pi_1, \pi_2) = -\frac{1}{\binom{n}{2}} \cdot |I_1 \setminus I_2| \cdot |I_1 \setminus I_2|$
- 11:  $s_5(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \cdot |I_1 \cap I_2| \cdot |[n] \setminus (I_1 \cup I_2)|$
- 12:  $k^{ck}(\pi_1, \pi_2) = s_1(\pi_1, \pi_2) + s_2(\pi_1, \pi_2) + s_3(\pi_1, \pi_2) + s_4(\pi_1, \pi_2) + s_5(\pi_1, \pi_2)$

**Claim 4.** Algorithm 3 computes the convolutional Kendall kernel (as given in the Equation 3) with a computational complexity of  $\mathcal{O}(k^2)$ .

580 *Proof.* To establish the correctness of Algorithm 3, we will adopt the same proof approach as the one  
 581 used for Claim 1 concerning Algorithm 2. Specifically, we will adhere to the earlier categorization in  
 582 Table 4 and notations given in Definition 1. Since the CK kernel can be derived by uniformly setting  
 583 the weight function  $w_s(i, j) = 1$ , we will insert them in  $s_i$  terms as given in Algorithm 2. These  
 584 cases will be revisited and simplified by applying the condition  $w_s(i, j) = 1$ . Note that this also  
 585 implies its one-direction marginal weights to be 1, i.e.,  $w_s(i, \cdot) = 1$

586 **Simplifying the  $s_1$  Term:** For the WCK kernel, Case 1-a leads to the expression of  $s_1$  as stated in  
 587 Equation 12. In this case, when two items, specifically  $l$  and  $m$ , are both in the intersection  $I_1 \cap I_2$ ,  
 588 it implies that  $\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}, \mathbf{p}_l^{\pi_2}$ , and  $\mathbf{p}_m^{\pi_2}$  all rank within the top-k, denoted as  $[k]$ . We simplify the  $s_1$   
 589 term for CK kernel as follows:

$$\begin{aligned}
s_1(\pi_1, \pi_2) &= \frac{1}{\binom{n}{2}} \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2}) \cdot \eta_{l,m}(\pi_1, \pi_2) \\
&= \frac{1}{\binom{n}{2}} \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} \eta_{l,m}(\pi_1, \pi_2) \\
&= \frac{1}{\binom{n}{2}} \sum_{1 \leq l' < m' \leq |I_1 \cap I_2|} \eta_{l',m'}(\tau_1, \tau_2) = \frac{\binom{|I_1 \cap I_2|}{2}}{\binom{n}{2}} k^{sk}(\tau_1, \tau_2). \tag{17}
\end{aligned}$$

590 The simplification process begins by assigning unit rank weights in the first line, i.e.,  $w_i = 1$ .  
591 Following this, by relabeling the items in  $I_1 \cap I_2$  and using  $\tau_1$  and  $\tau_2$ , which are the rankings of  
592  $\pi_1$  and  $\pi_2$  limited to the set  $I_1 \cap I_2$  as defined in Definition 1, it is established that  $\eta_{l',m'}(\tau_1, \tau_2) =$   
593  $\eta_{l,m}(\pi_1, \pi_2)$ . This is because the relative order of any pair of items is maintained in  $\tau_1$  and  $\tau_2$ .  
594 Consequently, this leads to the final simplification to a scaled value of the standard Kendall kernel  
595  $k^{sk}$ , as given in Equation 1.

596 **Simplifying the  $s_2$  and  $s_3$  Terms:** The  $s_2$  and  $s_3$  terms are obtained for Case 1-b, which is for  
597 case when one item is in  $I_1 \cap I_2$  and the other item is either in  $I_1 \setminus I_2$  or  $I_2 \setminus I_1$ . We divide this into  
598 two sub-cases. Case 1-b-i: The other item is in  $I_1 \setminus I_2$ , with  $s_2$  representing the summation terms  
599 derived from the CK's inner product. Case 1-b-ii: The other item is  $I_2 \setminus I_1$ , where  $s_3$  denotes the  
600 summation terms. We simplify the  $s_2$  term for the CK kernel as follows:

$$\begin{aligned}
s_2(\pi_1, \pi_2) &= \frac{1}{\binom{n}{2}} \sum_{l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \\
&= \frac{1}{\binom{n}{2}} \sum_{l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \\
&= \frac{1}{\binom{n}{2}} \underbrace{\sum_{l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} \mathbf{p}_{l < m}^{\pi_1}}_{:=u} - \frac{1}{\binom{n}{2}} \underbrace{\sum_{l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} \mathbf{p}_{l > m}^{\pi_1}}_{:=v}.
\end{aligned}$$

601 Next, we examine the terms  $u$  and  $v$  in detail, starting with  $u$ . The term  $u$ , which corresponds to  
602  $\mathbf{p}_{l < m}^{\pi_1}$ , signifies instances where item  $l$  is ranked before item  $m$  in the top-k ranking  $\pi_1$ . This can be  
603 derived from the observation that  $\sigma_1(l) - 1$  items are positioned before item  $l$  in the set  $I_1$ . Out of  
604 these items,  $\tau_1(l) - 1$  also belong to the intersection  $I_1 \cap I_2$ . This follows from the definition of  
605 the full rankings  $\sigma_1$  and  $\tau_1$  on the set  $I_1$  and the intersection  $I_1 \cap I_2$ , respectively. Consequently, it  
606 can be concluded that  $\sigma_1(l) - \tau_1(l)$  items from the set difference  $I_1 \setminus I_2$  are ranked before item  $l$ .  
607 The second term,  $v$ , corresponds to  $\mathbf{p}_{l > m}^{\pi_1}$  and involves a calculation that takes into account the items  
608 ranked after the  $l$ -th item in the set  $I$ . Specifically, there are  $k - \sigma_1(l)$  items following the  $l$ -th item.  
609 Within the intersection  $I_1 \cap I_2$ , the number of items before  $l$  is given by  $|I_1 \cap I_2| - \tau_1(l)$ . Therefore,  
610 the expression  $(k - \sigma_1(l)) - (|I_1 \cap I_2| - \tau_1(l))$  represents the count of elements that are positioned  
611 after  $l$  in the set difference  $I_1 \setminus I_2$ .

612 Combining the above calculations for both terms  $u$  and  $v$ , the  $s_2$  term for the CK kernel can be  
613 simplified as follows:

$$s_2(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \sum_{l \in I_1 \cap I_2} 2 \cdot (\sigma_1(l) - \tau_1(l)) - k + |I_1 \cap I_2|. \tag{18}$$

614 Using the symmetry between Case 1-b-i and Case 1-b-ii, we can simplify  $s_3$  for the CK kernel as  
615 follows:

$$s_3(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \sum_{l \in I_1 \cap I_2} 2 \cdot (\sigma_2(l) - \tau_2(l)) - k + |I_1 \cap I_2|. \tag{19}$$

616 **Simplifying the  $s_4$  and  $s_5$  Terms:** We simplify the  $s_4$  and  $s_5$  terms for the CK kernel starting from  
 617 Equation 15 and Equation 16, respectively, as follows:

$$s_4(\pi_1, \pi_2) = \frac{-1}{\binom{n}{2}} \cdot \sum_{l \in I_1 \setminus I_2, m \in I_2 \setminus I_1} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot) = \frac{-|I_1 \setminus I_2| \cdot |I_2 \setminus I_1|}{\binom{n}{2}} \quad (20)$$

$$s_5(\pi_1, \pi_2) = \frac{1}{\binom{n}{2}} \cdot (n - |I_1 \cup I_2|) \cdot \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) = \frac{|I_1 \cap I_2| \cdot |[n] \setminus (I_1 \cup I_2)|}{\binom{n}{2}}. \quad (21)$$

618 We have obtained the values of all the simplified  $s_i$  terms for the CK kernel in Equations 17, 18,  
 619 19, 20, and 21. By combining these terms, we get  $k^{ck}(\pi_1, \pi_2) = \sum_{i=1}^5 s_i(\pi_1, \pi_2)$ , where each term  
 620  $s_i$  precisely matches its corresponding expression in Algorithm 3. This completes the proof of the  
 621 correctness of Algorithm 3. Regarding its time complexity, each term  $s_i$  sums at most  $k^2$  quantities,  
 622 and each quantity can be computed in  $\mathcal{O}(1)$  time. Therefore, the time required for Algorithm 3 to  
 623 compute the CK kernel is  $\mathcal{O}(k^2)$ .  $\square$

### 624 A.3 Fast Matrix-Vector Multiplication with Kendall Kernel Matrix on Top-k Rankings

625 This section revisits Theorem 1 about the fact matrix-vector multiplication time for the Kendall kernel  
 626 matrix for top-k rankings. Specifically, we aim to eliminate the  $\mathbf{mvm}(K_X)$ 's dependence on the  
 627 number of items, i.e.,  $n$  on and linear dependence in the number of rounds, i.e.,  $T$ , as claimed in  
 628 Theorem 1.

**Theorem 1.** *For the WCK kernel with product-symmetric weights  $w_{ps}$ , the computational complexity of multiplying the kernel matrix  $K_{X_t}$  with any admissible vector is  $\mathcal{O}(k^2 t)$ , i.e.,  $\mathbf{mvm}(K_{X_t}) = \mathcal{O}(k^2 t)$ , where  $X_t$  is any arbitrary set of  $t$  top-k rankings.*

629 *Proof.* The cornerstone of this proof lies in the computation of the WCK kernel, as delineated in  
 630 Algorithm 2. This algorithm requires only  $\mathcal{O}(k^2)$  computation. For brevity, we write  $X$  to represent  
 631  $X_T$ , and the proof follows for any  $X_t$ , i.e., any value of  $t$ , not just  $T$ .

632 As also suggested previously, we will demonstrate through the equation  $K_X = (\Phi_X^a)^T \Phi_X^b$ , where  
 633 both matrices  $\Phi_X^a$  and  $\Phi_X^b$  have columns with only  $\mathcal{O}(k^2)$  non-zero entries. Consequently, this  
 634 leads to the computational complexity of matrix-vector multiplication, denoted as  $\mathbf{mvm}(K_X)$ , being  
 635  $\mathcal{O}(k^2 \cdot T)$ .

636 From Algorithm 2, we know that each entry of the kernel matrix  $k(\pi_1, \pi_2)$ , can be expressed as a sum  
 637  $\sum_{i=1}^5 s_i(\pi_1, \pi_2)$ . Assuming each  $s_i(\pi_1, \pi_2)$  equals  $\phi^{a_i}(\pi_1)^T \phi^{b_i}(\pi_2)$ , and considering that all vectors  
 638  $\phi^{a_i}$  and  $\phi^{b_i}$  exhibit this property, we can express  $K_X$  as  $(\Phi_X^a)^T \Phi_X^b$ . Here, the  $i^{th}$  row of  $(\Phi_X^a)^T$   
 639 and the  $j^{th}$  column of  $\Phi_X^b$  are represented by  $[\phi^{a_1}(\pi_i)^T, \dots, \phi^{a_5}(\pi_i)^T]$  and  $[\phi^{b_1}(\pi_j), \dots, \phi^{b_5}(\pi_j)]$ ,  
 640 respectively. Therefore, the overall mvm complexity can be characterized by the sparsity of the  
 641 vectors  $\phi^{a_i}$  and  $\phi^{b_i}$ , as is formalized in the claim presented below.

**Claim 5.** *Consider a kernel matrix  $K_X$  corresponding to any set  $X$  of cardinality  $T$ . Each entry of  $K_X$ , denoted as  $k(x_1, x_2)$ , is defined by the sum  $\sum_{i=1}^5 s_i(x_1, x_2)$ , where each  $s_i(x_1, x_2)$  is the result of the dot product  $\phi^{a_i}(x_1)^T \phi^{b_i}(x_2)$ , where,  $\phi^{a_i}$  and  $\phi^{b_i}$  are vectors characterized by having  $\mathcal{O}(z)$  non-zero entries. Given this structure, the matrix-vector multiplication complexity for  $K_X$  is  $\mathcal{O}(nnz \cdot T)$ , i.e.,  $\mathbf{mvm}(K_X) = \mathcal{O}(z \cdot T)$ .*

642 *Proof.* We will demonstrate this in the following discussion by concentrating on the  $k^{th}$  entry of the  
 643 output vector, specifically  $K_X \mathbf{v}$ , for any arbitrary vector  $\mathbf{v}$ :

$$(K_X \mathbf{v})_k = \sum_j K_X(k, j) v_j = \sum_j \left( \sum_{i=1}^5 s_i(\pi_k, \pi_j) \right) v_j = \sum_j \left( \sum_{i=1}^5 \phi^{a_i}(\pi_k)^T \phi^{b_i}(\pi_j) \right) v_j,$$

$$= \sum_{i=1}^5 \left( \sum_j \phi^{a_i}(\pi_k)^T \phi^{b_i}(\pi_j) v_j \right) = \sum_{i=1}^5 \phi^{a_i}(\pi_k)^T \left( \sum_j \phi^{b_i}(\pi_j) v_j \right).$$

644 Given that for all  $i$ ,  $\phi^{b_i}$  possesses only  $\mathcal{O}(z)$  non-zero entries for any  $\pi_j$ , the computation of  
 645  $\sum_j \phi^{b_i}(\pi_j) v_j$  requires  $\mathcal{O}(z)$  operations. This implies that the expression  $\sum_j \phi^{b_i}(\pi_j) v_j$  also neces-  
 646 sitates  $\mathcal{O}(z)$  computation. Applying a similar rationale to  $\phi^{a_i}$ , it follows that computing  $(K_X v)_k$   
 647 demands only  $\mathcal{O}(z)$  operations. Extending this argument to all entries of the output vector, it is  
 648 evident that computing  $K_X \mathbf{v}$  requires only  $\mathcal{O}(z \cdot T)$  computation  $\square$

649 Utilizing Claim 5, it suffices to complete the proof by showcasing that these exist vectors  $\phi^{a_i}$  and  
 650  $\phi^{b_i}$ , each with only  $\mathcal{O}(k^2)$  non-zero elements, corresponding to each  $s_i$  as specified in Algorithm 2.  
 651 Additionally, these vectors ensure that  $s_i(\pi_1, \pi_2) = \phi^{a_i}(\pi_1)^T \phi^{b_i}(\pi_2)$ . We will next establish such  
 652 vectors for all  $s_i$  terms. Starting with the  $s_1$  term below.

653 **Showcasing**  $s_1(\pi_1, \pi_2) = \phi^{a_1}(\pi_1)^T \phi^{b_1}(\pi_2)$  for sparse  $\phi^{a_1}(\pi_1)$  and  $\phi^{b_1}(\pi_2)$  vectors. We begin  
 654 by manipulating  $s_1$ , as defined in Equation 12. For the sake of brevity, their scalar factors will be  
 655 omitted in the following explanation.

$$\begin{aligned} s_1(\pi_1, \pi_2) &= \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2}) \cdot \eta_{l,m}(\pi_1, \pi_2), \\ &= \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2}) \cdot (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}), \\ &= \sum_{1 \leq l < m \leq n | l, m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot (\mathbf{p}_{i < j}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2}) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}), \\ &= \sum_{1 \leq l < m \leq n} \underbrace{w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot (\mathbf{p}_{i < j}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \cdot \mathbf{1}_{\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1} \in [k]}}_{:= \phi_{l,m}^{a_1}(\pi_1)} \\ &\quad \cdot \underbrace{w_s(\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2}) \cdot (\mathbf{p}_{l < m}^{\pi_2} - \mathbf{p}_{l > m}^{\pi_2}) \cdot \mathbf{1}_{\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2} \in [k]}}_{:= \phi_{l,m}^{b_1}(\pi_2)}, \\ &= (\phi^{a_1}(\pi_1)^T \phi^{b_1}(\pi_2)). \end{aligned} \tag{22}$$

656 Both  $\phi^{a_1}$  and  $\phi^{b_1}$  are sparse by design, taking non-zero values only when  $l$  and  $m$  appear in the top- $k$   
 657 rankings. This demonstrates the existence of sparse vectors for the  $s_1$  term. Next, we will establish  
 658 the same for the  $s_2$  and  $s_3$  terms.

659 **Showcasing sparse vectors for  $s_2$  and  $s_3$ .** We begin by manipulating  $s_2$ , as defined in Equation 13,  
 660 while ignoring its scalar factor. We will exploit symmetry between  $s_2$  and  $s_3$  terms.

$$\begin{aligned} s_2(\pi_1, \pi_2) &= \sum_{l \in I_1 \cap I_2 | m \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}), \\ &= \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_2}, \cdot) \sum_{m \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}), \\ &= \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_2}, \cdot) \left( \sum_{m \in I_1} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) - \sum_{m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \right), \\ &= \sum_{l \in [n]} \underbrace{\mathbf{1}_{\mathbf{p}_l^{\pi_2} \in [k]} w_s(\mathbf{p}_l^{\pi_2}, \cdot)}_{:= \phi_l^{b_2}(\pi_2)} \underbrace{\mathbf{1}_{\mathbf{p}_l^{\pi_1} \in [k]} \sum_{m \in I_1} w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1})}_{:= \phi_l^{a_2}(\pi_1)} \end{aligned}$$

$$- \sum_{l,m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_2}, \cdot) w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}), \quad (23)$$

$$\begin{aligned} &= \phi^{a_{21}}(\pi_1)^T \phi^{b_{21}}(\pi_2) - \sum_{l,m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_2}, \cdot) w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}), \\ &= \phi^{a_{21}}(\pi_1)^T \phi^{b_{21}}(\pi_2) + \sum_{l,m \in [n]} \underbrace{-w_s(\mathbf{p}_l^{\pi_2}, \cdot) \mathbf{1}_{\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2} \in [k]}}_{:= \phi_{l,m}^{b_{22}}} \\ &\quad \cdot \underbrace{w_s(\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1}) (\mathbf{p}_{l < m}^{\pi_1} - \mathbf{p}_{l > m}^{\pi_1}) \mathbf{1}_{\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1} \in [k]}}_{:= \phi_{l,m}^{a_{22}}}, \end{aligned} \quad (24)$$

$$\begin{aligned} &= \phi^{a_{21}}(\pi_1)^T \phi^{b_{21}}(\pi_2) + \phi^{a_{22}}(\pi_1)^T \phi^{b_{22}}(\pi_2), \\ &= \underbrace{[\phi^{a_{21}}(\pi_1); \phi^{a_{22}}(\pi_2)]^T}_{:= \phi^{a_2}(\pi_1)^T} \underbrace{[\phi^{a_{21}}(\pi_2); \phi^{b_{22}}(\pi_2)]}_{:= \phi^{b_2}(\pi_2)} = \phi^{a_2}(\pi_1)^T \phi^{b_2}(\pi_2). \end{aligned} \quad (25)$$

Equation 25 demonstrates the existence of vectors  $\phi^{a_2}$  and  $\phi^{b_2}$  for the  $s_2$  term. The vectors  $\phi^{a_{21}}$  and  $\phi^{a_{22}}$ , possessing  $\mathcal{O}(k)$  and  $\mathcal{O}(k^2)$  non-zero entries respectively, are defined in Equations 23 and 24. Consequently, the  $\phi^{a_2}$  vector has  $\mathcal{O}(k^2)$  non-zero entries. Similarly, it can be shown that  $\phi^{b_2}$  contains  $\mathcal{O}(k^2)$  non-zero entries, thus fulfilling the proof requirements for proving the  $s_2$  term. For the  $s_3$  term, we observe a symmetry between  $s_2$  and  $s_3$ , namely  $s_3(\pi_1, \pi_2) = s_2(\pi_2, \pi_1)$ . This symmetry makes it trivial to satisfy the requirements, as further highlighted by the following equation:

$$s_3(\pi_1, \pi_2) = s_2(\pi_2, \pi_1) = \phi^{a_2}(\pi_2)^T \phi^{b_2}(\pi_1) = \underbrace{\phi^{b_2}(\pi_1)^T}_{:= \phi^{a_3}(\pi_1)} \underbrace{\phi^{a_2}(\pi_2)}_{:= \phi^{b_3}(\pi_2)} = \phi^{a_3}(\pi_1)^T \phi^{b_3}(\pi_2). \quad (26)$$

**Showcasing sparse vectors**  $s_4(\pi_1, \pi_2) = \phi^{4a}(\pi_1)^T \phi^{4b}(\pi_2)$ . We begin by manipulating the  $s_4$  term without scalar, as defined in Equation 15.

$$\begin{aligned} s_4(\pi_1, \pi_2) &= - \sum_{l \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot), \\ &= - \sum_{l \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot \left( \sum_{m \in I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot) - \sum_{m \in I_1 \cap I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot) \right). \end{aligned}$$

Observing that  $\bar{w} := \sum_{m \in I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot)$  represents a constant value that does not depend on  $I_2$ , we can further simplify the above expression for  $s_4$  as follows:

$$\begin{aligned} s_4(\pi_1, \pi_2) &= - \sum_{l \in I_1 \setminus I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot \left( \bar{w} - \sum_{m \in I_1 \cap I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot) \right), \\ &= - \left( \bar{w} - \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \right) \cdot \left( \bar{w} - \sum_{m \in I_1 \cap I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot) \right), \\ &= -\bar{w}^2 + \bar{w} \left( \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) + \sum_{m \in I_1 \cap I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot) \right) \\ &\quad - \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \sum_{m \in I_1 \cap I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot). \end{aligned} \quad (27)$$

Next, to simplify the above equation, we first focus on the second term and have the following:

$$\bar{w} \left( \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) + \sum_{m \in I_1 \cap I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot) \right)$$

$$\begin{aligned}
&= \sum_{l \in [n]} \underbrace{\mathbf{1}_{\mathbf{p}_l^{\pi_1} \in [k]} w_s(\mathbf{p}_l^{\pi_1}, \cdot)}_{:= \phi_l^{4a_1}(\pi_1)} \underbrace{\mathbf{1}_{\mathbf{p}_l^{\pi_2} \in [k]} \bar{w}}_{:= \phi_l^{4b_1}(\pi_2)} + \sum_{m \in I_1 \cap I_2} \bar{w} \cdot w_s(\mathbf{p}_m^{\pi_2}, \cdot), \\
&= \phi^{4a_1}(\pi_1)^T \phi^{4b_1}(\pi_2) + \sum_{m \in [k]} \underbrace{\mathbf{1}_{\mathbf{p}_m^{\pi_1} \in [k]} \bar{w}}_{:= \phi_m^{4a_2}(\pi_1)} \underbrace{w_s(\mathbf{p}_m^{\pi_1}, \cdot)}_{:= \phi_m^{4b_2}(\pi_2)} \\
&= \phi^{4a_1}(\pi_1)^T \phi^{4b_1}(\pi_2) + \phi^{4a_2}(\pi_1)^T \phi^{4b_2}(\pi_2). \tag{28}
\end{aligned}$$

Next, we simplify the third and last term in the Equation 27 as follows:

$$\begin{aligned}
\sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \sum_{m \in I_1 \cap I_2} w_s(\mathbf{p}_m^{\pi_2}, \cdot) &= \sum_{l \in [n], m \in [n]} \underbrace{w_s(\mathbf{p}_l^{\pi_1}, \cdot) \mathbf{1}_{\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1} \in [k]}}_{:= \phi_{l,m}^{4a_3}(\pi_1)} \underbrace{w_s(\mathbf{p}_m^{\pi_2}, \cdot) \mathbf{1}_{\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2} \in [k]}}_{:= \phi_{l,m}^{4b_3}(\pi_2)}, \\
&= \phi^{4a_3}(\pi_1)^T \phi^{4b_3}(\pi_2). \tag{29}
\end{aligned}$$

Next, combining the results from Equations 27, 28, and 29, we obtain the following:

$$\begin{aligned}
s_4(\pi_1, \pi_2) &= \underbrace{[\bar{w}, \phi^{4a_1}(\pi_1); \phi^{4a_1}(\pi_1); \phi^{4a_3}(\pi_1)]^T}_{:= \phi^{4a}(\pi_1)^T} [-\bar{w}; \phi^{4b_1}(\pi_2); \phi^{4b_2}(\pi_2); -\phi^{4b_3}(\pi_2)]_{:= \phi^{4b}(\pi_2)} \\
&= \phi^{4a}(\pi_1)^T \phi^{4b}(\pi_2). \tag{30}
\end{aligned}$$

Equation 30 showcases both  $\phi^{4a}$  and  $\phi^{4b}$  has three components with having only  $\mathcal{O}(k^2)$  non-zero entries, thus fulfilling the requirements for the  $s_4$  term. Next, we focus on the  $s_5$  term.

**Showcasing sparse vectors**  $s_5(\pi_1, \pi_2) = \phi^{5a}(\pi_1)^T \phi^{5b}(\pi_2)$ . We begin by examining the  $s_5$  term, excluding its scalar component, as outlined in Equation 16.

$$\begin{aligned}
s_5(\pi_1, \pi_2) &= (n - |I_1 \cup I_2|) \cdot \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot), \\
&= (n - (2k - |I_1 \cap I_2|)) \cdot \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot), \\
&= (n - 2k) \cdot \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) + |I_1 \cap I_2| \cdot \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot), \\
&= \sum_{l \in I_1 \cap I_2} \sqrt{n - 2k} \cdot w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot \sqrt{n - 2k} \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \\
&\quad + \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot |I_1 \cap I_2|, \tag{31} \\
&= \sum_{l \in [n]} \underbrace{\sqrt{n - 2k} \cdot w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot \mathbf{1}_{\mathbf{p}_l^{\pi_1} \in [k]}}_{:= \phi_l^{5a_1}(\pi_1)} \underbrace{\sqrt{n - 2k} \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot \mathbf{1}_{\mathbf{p}_l^{\pi_2} \in [k]}}_{:= \phi_l^{5b_1}(\pi_2)} \\
&\quad + \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot |I_1 \cap I_2|, \\
&= \phi^{5a_1}(\pi_1)^T \phi^{5b_1}(\pi_2) + \sum_{l \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot \sum_{m \in I_1 \cap I_2} 1, \\
&= (\phi^{5a_1}(\pi_1)^T \phi^{5b_1}(\pi_2) + \sum_{l \in I_1 \cap I_2, m \in I_1 \cap I_2} w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot w_s(\mathbf{p}_l^{\pi_2}, \cdot)), \\
&= \phi^{5a_1}(\pi_1)^T \phi^{5b_1}(\pi_2) + \sum_{l \in [n], m \in [n]} \underbrace{w_s(\mathbf{p}_l^{\pi_1}, \cdot) \cdot \mathbf{1}_{\mathbf{p}_l^{\pi_1}, \mathbf{p}_m^{\pi_1} \in [k]}}_{:= \phi_{l,m}^{5a_2}(\pi_1)} \underbrace{w_s(\mathbf{p}_l^{\pi_2}, \cdot) \cdot \mathbf{1}_{\mathbf{p}_l^{\pi_2}, \mathbf{p}_m^{\pi_2} \in [k]}}_{:= \phi_{l,m}^{5b_2}(\pi_2)}, \tag{32}
\end{aligned}$$



$$\begin{aligned}
&= \phi^{5a_1}(\pi_1)^T \phi^{5b_1}(\pi_2) + \phi^{5a_2}(\pi_1)^T \phi^{5b_2}(\pi_2), \\
&= \underbrace{[\phi^{5a_1}(\pi_1); \phi^{5a_2}(\pi_1)]^T}_{:=\phi^{5a}(\pi_1)^T} \underbrace{[\phi^{5b_1}(\pi_2) + \phi^{5b_2}(\pi_2)]}_{:=\phi^{5b}(\pi_2)} = \phi^{5a}(\pi_1)^T \phi^{5b}(\pi_2).
\end{aligned} \tag{33}$$

The equation shows that  $s_5(\pi_1, \pi_2) = \phi^{5a}(\pi_1)^T \phi^{5b}(\pi_2)$ , where both  $\phi^{5a}$  and  $\phi^{5b}$  possess components with a maximum number of non-zero entries, as indicated in Equations 31 and 32. This completes the proof requirements for the  $s_5$  term.

By combining the results from Equations 22, 25, 26, 30, and 33, we have demonstrated the existence of vectors  $\phi^{a_i}$  and  $\phi^{b_i}$ , each containing only  $\mathcal{O}(k^2)$  non-zero elements, and have established that  $s_i(\pi_1, \pi_2) = \phi^{a_i}(\pi_1)^T \phi^{b_i}(\pi_2)$  for each  $i \in 1, 2, 3, 4, 5$ . In conjunction with Claim 5, this completes the proof.  $\square$

## B Proposed GP-TopK Bandit Algorithm– Omitted Details

This section includes the proofs that were omitted from Section 4, presented in the following order:

- Section B.1 outlines a brief of Gaussian process regression for any domain.
- Section B.2 summarizes the committed details about the local search utilized for optimizing the UCB function.
- Section B.3 provides the removed proof for the Theorem 2 concerning the overall time for the bandit algorithm.
- Section B.4 provides the proof for Theorem 3 concerning regret analysis of the proposed bandit algorithm.

### B.1 Gaussian Process Regression

In GP regression [22], the training data are modeled as noisy measurements of a random function  $f$  drawn from a GP prior, denoted  $f \sim \mathcal{N}(0, k(\cdot, \cdot))$ , where  $k : \mathcal{X} \times \mathcal{X} \rightarrow \mathbb{R}$  is a kernel function over any domain  $\mathcal{X}$ . The observed training pairs  $(\mathbf{x}_i, y_i)$  are collected as  $X = [\mathbf{x}_1, \dots, \mathbf{x}_T]$  and  $\mathbf{y} = [y_1, \dots, y_T] \in \mathbb{R}^T$ , where, for an input  $\mathbf{x}_i$ , the observed value is modeled as  $y_i = f(\mathbf{x}_i) + \epsilon$ , with  $\epsilon_i \sim \mathcal{N}(0, \sigma^2)$ . The kernel matrix on data is  $K_X = [k(\mathbf{x}_i, \mathbf{x}_j)]_{i,j=1}^T \in \mathbb{R}^{T \times T}$ . The posterior mean  $\mu_{f|\mathcal{D}}$  and variance  $\sigma_{f|\mathcal{D}}$  functions for GPs are:

$$\mu_{f|\mathcal{D}}(\mathbf{x}) := \mathbf{k}_{\mathbf{x}}^T \mathbf{z} \tag{34}$$

$$\sigma_{f|\mathcal{D}}(\mathbf{x}) := k(\mathbf{x}, \mathbf{x}) - \mathbf{k}_{\mathbf{x}}^T (K_X + \sigma^2 I)^{-1} \mathbf{k}_{\mathbf{x}} \tag{35}$$

where  $\mathbf{k}_{\mathbf{x}} \in \mathbb{R}^T$  has as its  $i^{th}$  entry  $k(\mathbf{x}, \mathbf{x}_i)$ ,  $\mathbf{z} = (K_X + \sigma^2 I)^{-1} \mathbf{y}$ , and  $I$  is an identity matrix. For GP regression on an arbitrary domain  $\mathcal{X}$ , the kernel function must be a p.d. kernel [23].

Naive approaches rely on the Cholesky decomposition of the matrix  $K_X + \sigma^2 I$ , which takes  $\Theta(T^3)$  time [23]. To circumvent the  $\Theta(T^3)$  runtime, recent works use iterative algorithms such as the conjugate gradient algorithm, which facilitate GP inference by exploiting fast kernel matrix-vector multiplication (MVM) algorithms, i.e.,  $\mathbf{v} \mapsto \mathbf{K}_{\mathbf{X}} \mathbf{v}$  [3]. In practice, these methods yield highly accurate approximations for GP posterior functions with a complexity of  $\Theta(p \cdot T^2)$  for  $p$  CG iterations, as  $\text{mvm}(K_X) = T^2$ , and  $\text{mvm}(M)$  is the operation count for multiplying matrix  $M$  by a vector.  $p \ll T$  proves to be efficient in practical application [3].

### B.2 Contextual GP Reward Model

Optimizing the  $\mathcal{AF}$ , i.e., UCB function, poses a significant challenge due to its enormous size of  $\Pi^k$ . Drawing inspiration from prior research on Bayesian optimization within combinatorial spaces, we employ a breadth-first local search (BFLS) to optimize the UCB acquisition function [2, 19]. The BFLS begins with the selection of several random top-k rankings. Subsequently, each specific top-k ranking is compared with the UCB values of its neighboring rankings, proceeding to the one with the highest UCB value.

717 The neighbors of a top-k ranking include all its permutations and the permutations of modified top-k  
 718 rankings obtained by swapping one item with any of the remaining items. For any top-k ranking,  
 719 there are  $(n - k) \cdot k! + k!$  neighbors, which is often not huge as  $k$  is often  $\leq 6$ . This search continues  
 720 until no neighboring top-k ranking with a higher value is discovered. Although BFLS is a local  
 721 search, the initial random selection and multiple restart points help it evade local minima, a strategy  
 722 that previous studies have corroborated [19].

### 723 B.3 Assessing GP-TopK Compute Requirements

**Theorem 2.** *Assuming a fixed number of iterations required by the iterative algorithms, the total computational time for running the GP-TopK bandit algorithm for  $T$  rounds of top-k recommendations, using the contextual product kernel (Equation 6), is  $\mathcal{O}(k^2 c \ell T^2)$ . This applies to WK, CK, and WCK top-k ranking kernels, where  $\ell$  is the number of local search evaluations for selecting the next arm in every round.*

724 *Proof.* The proof can be straightforwardly derived by combining the results presented in Table 1,  
 725 which succinctly summarizes the time complexities for each step of computing the UCB using both  
 726 feature and kernel approaches. It is important to emphasize that iterative algorithms enhance results  
 727 from  $\mathcal{O}(T^4)$  to  $\mathcal{O}(T^3)$  in computational complexity. Furthermore, these algorithms can further  
 728 reduce complexity to  $\mathcal{O}(T^2)$  when used with the feature approach.

729 The results presented in Table 1 can be validated through straightforward observations and by  
 730 leveraging findings from previous Sections 2. Specifically, Section 2 offers proof for the `mvm`( $K_X$ )  
 731 row explicitly. For the `compute`  $K_{X_t}$  row, the complexity of kernel approaches is deduced from  
 732 Algorithms 2 and 3. For feature approaches, the `compute`  $K_{X_t}$  row is inferred from the sparsity of the  
 733 feature representations as stated in Claim 3. Lastly, the `memory` row is straightforwardly deduced for  
 734 the kernel approach by counting its entries. For the feature approach, it is derived from the sparsity of  
 735 the feature representations.  $\square$

### 736 B.4 Regret Analysis

737 In this section, we revisit Theorem 3 and provide its proof. The proofs build on the work by Krause  
 738 et al. [14], delivering results for bounding the contextual regret in the context of the top-k ranking  
 739 problem. To set the stage for our regret analysis, let's first define the critical term *maximum mutual*  
 740 *information*, denoted by  $\gamma_t$ , is given below:

$$\gamma_t := \max_{X \subseteq \mathcal{X}: |X|=t} I(y_X; f), \quad I(y_X; f) = H(y_X) - H(y_X|f),$$

741 where  $I(y_X; f)$  quantifies the reduction in uncertainty (measured in terms of differential Shannon  
 742 entropy) about  $f$  achieved by revealing  $y_A$  [28]. In Gaussian observation case, the entropy can be  
 743 computed in closed form:  $H(N(\mu, \Sigma)) = \frac{1}{2} \log |2\pi e \Sigma|$ , so that  $I(y_X; f) = \frac{1}{2} \log |I + \sigma^{-2} K_X|$ ,  
 744 where  $K_X = [k(x, x')]_{x, x' \in X}$  is the Gram matrix of  $k$  evaluated on set  $X \subseteq \mathcal{X}$ . For the contextual  
 745 bandit algorithm,  $X$  represents contexts and arms considered until round  $t$ .

746 Before proving Theorem 3, we align the Krause et al. [14] results with our notation for consistency.  
 747 Furthermore, we modify  $\beta_t$  to accommodate embeddings encompassing negative values, aligning  
 748 with the fact that contextual embeddings may exhibit negative dimensions.

**Proposition 1** (Theorem 1, [14]). *Let  $\delta \in (0, 1)$ , and the unknown reward function  $\hat{f}$  be sampled*

749

from the known GP prior with known noise variance  $\sigma^2$ . Suppose one of the following holds:

1. Assumption 1 holds and set  $\beta_t = 2 \log(|\mathcal{X}|t^2\pi^2/6\delta)$ .
2. Assumption 2 holds and set  $\beta_t = 2B^2 + 300\gamma_t \ln^3(t/\delta)$ .

Then the cumulative regret  $\mathcal{R}_T$  of the contextual GP bandit algorithm with the UCB acquisition function is bounded by  $\tilde{\mathcal{O}}(\sqrt{C_1 T \gamma_T \beta_T})$  w.h.p. Precisely,  $\Pr\{R_T \leq \sqrt{C_1 T \gamma_T \beta_t} + 2 \quad \forall T \geq 1\} \geq 1 - \delta$ , where,  $C_1 = 8/\log(1 + \sigma^{-2})$  and the notation  $\tilde{\mathcal{O}}$  hides logarithmic factors in  $n$ ,  $\frac{1}{\delta}$  and  $T$ .

750

Proposition 1 shoes that the regret  $\mathcal{R}_T$  for the contextual GP bandit algorithm, utilizing the UCB acquisition function is bounded with high probability within  $\tilde{\mathcal{O}}(\sqrt{C_1 T \gamma_T \beta_T})$ , where the notation  $\tilde{\mathcal{O}}$  hides logarithmic factors in  $n$ ,  $\frac{1}{\delta}$  and  $T$ . To ascertain the  $\tilde{\mathcal{O}}$  order for  $\mathcal{R}_T$ , it is imperative to first bound the  $\tilde{\mathcal{O}}$  order of  $\gamma_T \beta_t$ . We begin by examining the  $\gamma_T$  term in the subsequent proposition.

**Proposition 2.** Under the assumptions of Theorem 3,  $\gamma_T$  can be succinctly characterized as  $\gamma_T = \mathcal{O}(n^2 c \log(n^2 T) + c \log T)$ , which also simplifies to  $\tilde{\mathcal{O}}(n^2 c)$ , where the  $\tilde{\mathcal{O}}$  notation omits logarithmic factors in  $n$  and  $T$ .

*Proof.* For the GP bandit algorithm with the UCB acquisition function,  $\gamma_T = C \cdot \log(|I + \sigma^{-2} K_{X_T}|)$ , where  $C$  equals  $(1/2) \cdot (1 - 1/e)^{-1}$  and  $K_{X_T}$  represents the kernel matrix computed over contexts and arms across  $T$  rounds [28, 14]. Precisely,  $K_{X_T}$  is calculated using the contextual kernel defined in Equation 6. It is applied to contexts and top-k ratings from the feedback data  $\mathcal{D}_t$ , corresponding to Line 6 of the generic contextual bandit Algorithm 1.

Next, we leverage the characteristic of the contextual kernel being a product kernel. Consequently, the maximum mutual information term for the joint kernel,  $\gamma_T$ , can be upper bounded by  $c \cdot (\gamma_T^\pi + \log T)$ , where  $c$  denotes the dimensionality of contexts and  $\gamma_T^\pi$  represents the maximum information gain in a non-contextual setting [14]. Specifically,  $\gamma_T^\pi$  is computed similarly but is confined to top-k rankings. That is,  $\gamma_T^\pi = C \cdot \log(|I + \sigma^{-2} K_{X_T^\pi}|)$ , with  $K_{X_T^\pi}$  being calculated exclusively using the top-k kernels on the top-k rankings as selected by the bandit algorithm.  $X_T^\pi$  represents the top-k rankings selected by the bandit algorithm, i.e., excluding the contexts from the collected feedback.

Recalling the formulation for top-k rankings kernels, we have  $K_{X_T} = \Phi_{X_T^\pi}^T \Phi_{X_T^\pi}$ , where  $\Phi_{X^\pi} \in \mathbb{R}^{\binom{n}{2} \times T}$  comprises feature columns pertinent to the top-k ranking kernels, as elucidated in Section A. Utilizing the Weinstein–Aronszajn identity,  $\gamma_T^\pi$  is expressed as  $C \cdot \log(|I + \sigma^{-2} \Phi_{X_T^\pi} \Phi_{X_T^\pi}^T|)$ . Further,

we deduce that  $\gamma_T^\pi \leq C \cdot \sum_{i=1}^{\binom{n}{2}} \log(|1 + \sigma^{-2} \lambda_i|)$ , where  $\lambda_i$  is an eigenvalue of  $\Phi_{X_T^\pi} \Phi_{X_T^\pi}^T$ . Given the Gershgorin circle theorem, which bounds all eigenvalues of a matrix by the maximum absolute sum of its rows, therefore we can conclude that  $\gamma_T^\pi = \mathcal{O}(n^2 \log(n^2 T))$ , as for all the columns of the  $\Phi_{X^\pi}$  have bounded normed as given in Claims 2 and 3, i.e.,  $\|\phi(\pi)\|_2^2 \leq 1$  [30].

By combining  $\gamma_T^\pi = \mathcal{O}(n^2 \log(n^2 T))$  with the contextual product kernel, we obtain  $\gamma_T = \mathcal{O}(n^2 c \log(n^2 T) + c \log T)$ , thereby providing the claimed bound in the proposition.  $\square$

Next, we build on Propositions 1 and 2 to prove the main theorem regarding the regret of the proposed GP-TopK bandit algorithm for top-k recommendations.

**Theorem 3.** If either Assumptions 1 or 2 hold, setting  $\beta_t$  as  $2 \log\left(\frac{|\mathcal{C}| \cdot |\Pi^k| \cdot t^2 \cdot \pi^2}{6\delta}\right)$  and  $300\gamma_t \ln^3\left(\frac{t}{\delta}\right)$  respectively, the cumulative regret  $\mathcal{R}_T$  of the GP-TopK bandit algorithm for top-k recommendations can, with at least  $1 - \delta$  probability, be bounded by  $\tilde{\mathcal{O}}(n\sqrt{C_1 T c(\log|\mathcal{C}| + k + \log(T^2 \pi^2/6\delta))})$  under Assumption 1, and  $\tilde{\mathcal{O}}(n\sqrt{C_1(2B^2 c + 300n^2 c^2 \ln^3(T/\delta))T})$  under Assumption 2. Here,  $C_1 = \frac{8}{\log(1+\xi^{-2})}$ , and  $\tilde{\mathcal{O}}$  excludes logarithmic factors related to  $n$ ,  $k$ , and  $T$ .

778 *Proof.* We will prove the above theorem for both cases separately.

779 **For Assumption-1.** Given  $|\mathcal{C}|$  is finite and  $\beta_T = 2 \log(|\mathcal{D}|T^2\pi^2/6\delta)$ . First, we focus on bounding  
780  $\beta_T$  as follows:

$$\begin{aligned}\beta_T &= 2 \log(|\mathcal{D}|T^2\pi^2/6\delta) \\ &= \mathcal{O}(\log|\mathcal{C}| + \log|\Pi^k| + \log(T^2\pi^2/6\delta))\end{aligned}$$

781 As  $\binom{n}{k} \leq n^k$  and  $k! \leq k^k$ , we also have  $\log|\Pi^k| = \log\left(\binom{n}{k}k!\right) \leq \log(n^k k^k) = \mathcal{O}(k \log(nk))$ ,  
782 which implies that  $\beta_T = \mathcal{O}(\log|\mathcal{C}| + k \log(nk) + \log(T^2\pi^2/6\delta))$ . Combining this with Proposition  
783 2, we have following:

$$\begin{aligned}\mathcal{O}(\gamma_T \beta_T) &= \mathcal{O}\left((n^2 c \log(n^2 T) + c \log T)(\log|\mathcal{C}| + k \log(nk) + \log(T^2\pi^2/6\delta))\right) \\ &= \mathcal{O}\left(n^2 c \log(n^2 T)(\log|\mathcal{C}| + k \log(nk) + \log(T^2\pi^2/6\delta))\right) \quad (\text{Ignoring } c \log T \text{ term}) \\ &= \tilde{\mathcal{O}}\left(n^2 c (\log|\mathcal{C}| + k + \log(T^2\pi^2/6\delta))\right).\end{aligned}$$

784 Thus, we showcase the asserted bound for the regret  $\mathcal{R}_T$  as  $\tilde{\mathcal{O}}(\sqrt{C_1 T \gamma_T \beta_T}) =$   
785  $\tilde{\mathcal{O}}\left(n \sqrt{C_1 T c (\log|\mathcal{C}| + k + \log(T^2\pi^2/6\delta))}\right)$ .

786 **For Assumption-2.** Given  $\|f\|_k \leq B$  and  $\beta_t = 2B^2 + 300\gamma_t \ln^3(t/\delta)$ . First, we bound the  $\beta_T$  term  
787 using Proposition 2 as follows:

$$\begin{aligned}\beta_T &= 2B^2 + 300 \cdot \gamma_T \cdot \ln^3(T/\delta), \\ &= 2B^2 + 300 \cdot (n^2 c \log(n^2 T) + c \log T) \cdot \ln^3(T/\delta).\end{aligned}$$

788 Using the above result, we have the following:

$$\begin{aligned}\mathcal{O}(\sqrt{C_1 T \gamma_T \beta_T}) &= \mathcal{O}\left(\sqrt{C_1 T \gamma_T \cdot (2B^2 + 300 \cdot \gamma_T \cdot \ln^3(T/\delta))}\right), \\ &= \mathcal{O}\left(\sqrt{C_1 T n^2 c \log(n^2 T) \cdot (2B^2 + 300 \cdot n^2 c \log(n^2 T) \cdot \ln^3(T/\delta))}\right), \\ &= \tilde{\mathcal{O}}\left(n \sqrt{C_1 T c (2B^2 + 300 n^2 c \ln^3(T/\delta))}\right).\end{aligned}$$

789

□

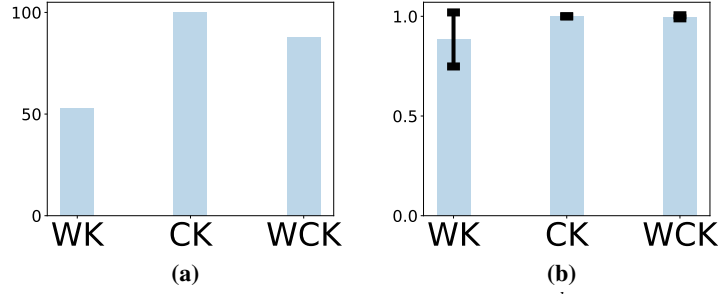
790 **Comparison with Srinivas et al. (2010).** Using the identity kernel for top-k rankings, we can  
791 develop a finite-dimensional feature for the contextual kernel and apply Theorem 5 by Srinivas et al.  
792 (2010). Given that  $\gamma_T = \mathcal{O}(n^k c \log T)$ , the regret bounds are as follows under both assumptions. For  
793 instance, the calculations for the  $\mathcal{O}(\sqrt{C_1 T \gamma_T \beta_T})$  under the Assumption 2 are as follows:

$$\begin{aligned}\mathcal{O}(\sqrt{C_1 T \gamma_T \beta_T}) &= \mathcal{O}\left(\sqrt{C_1 T \gamma_T \cdot (2B^2 + 300 \cdot \gamma_T \cdot \ln^3(T/\delta))}\right), \\ &= \mathcal{O}\left(\sqrt{C_1 T (n^k c \log T) \cdot (2B^2 + 300 \cdot (n^k c \log T) \cdot \ln^3(T/\delta))}\right), \\ &= \tilde{\mathcal{O}}\left(n^{\frac{k}{2}} \sqrt{C_1 T c (2B^2 + 300 n^k c \ln^3(T/\delta))}\right).\end{aligned}$$

794 Similarly, we can analogously perform the analysis for Assumption 1 and combine it with Proposi-  
795 tion 1 to obtain the regret bounds mentioned in the Table 3.

## 796 C Experiments – Omitted Details

797 This section presents omitted details from the main body of the text.



**Figure 4:** Local search results for optimizing combinatorial objectives in  $\Pi^k$  for  $n = 50$  and  $k = 6$ . For details, see the textual description. Left (a) shows how many times out of 100 trials the local search recovers the exact maximizer, i.e.,  $\pi^*$ , and right plot (b) shows the average value of the objective for the returned maximizer. These results indicate that the local search utilized in this work is effective.

### 798 C.1 Compute resources

799 We utilized multiple NVIDIA Tesla M40 GPUs with 40 GB RAM on our in-house cluster for our  
800 experiments. The experiments in Section 5 required approximately 5 GPU-hours for small arm  
801 space and 24 GPU-hours per iteration for large arm space. We conducted about 50 to 100 iterations  
802 throughout the project. The results reported in Section C.3 required the same computational resources  
803 as the large arm space experiments.

### 804 C.2 Bandit Simulation and Hyper-parameter Configurations – Omitted Details

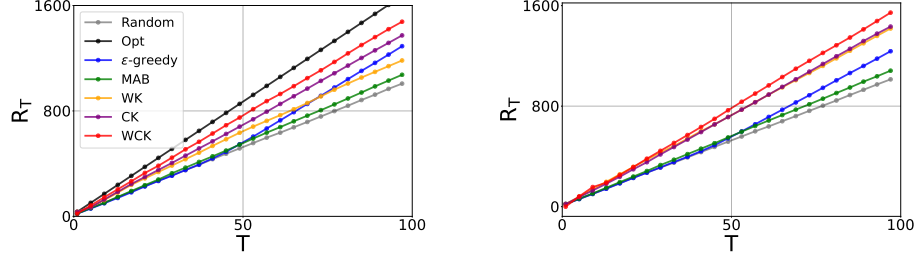
805 To set up the simulation, we utilized embeddings trained on the MovieLens dataset using a collabora-  
806 tive filtering approach [6]. We consider a  $1M$  variant of the MovieLens dataset, which contains 1  
807 million ratings from 6040 users for 3677 items. Specifically, we train user embeddings  $\mathbf{c}_u$  and item  
808 embeddings  $\theta_i$  such that the user’s attraction to the items are captured by the inner product of the user  
809 embedding with the item embeddings, respectively. Both context and item embeddings, i.e.,  $\mathbf{c}_u$  and  
810  $\theta_i$ , are 5-dimensional, optimized by considering the 5-fold performance on this dataset. The reward  
811 provided in our experiments is contaminated with zero mean and standard deviation equals 0.05.

812 For the  $\epsilon$ -greedy baselines, we considered various values of  $\epsilon$  are considered, specifically  $\epsilon =$   
813  $\{0.01, 0.05, 0.1\}$ . The outcomes are presented for the configuration that demonstrates optimal  
814 performance. For *MAB-UCB* baseline, the algorithm has an upper confidence score  $ucb(i) =$   
815  $\bar{\mu}_i + \beta_{mab} \sqrt{\frac{2 \ln(t+1)}{n_i}}$  [11]. Here,  $\bar{\mu}_i$  represents the average reward,  $n$  denotes the total number of  
816 rounds, and  $n_i$  signifies the frequency of arm  $i$  being played.  $\beta_{mab}$  is a hyper-parameter. We evaluate  
817  $\beta_{mab}$  values within the set  $\{0.1, 0.25, 0.5\}$  and disclose results for the best-performing configuration.  
818 For the parameters of proposed GP-TopK bandit algorithms, we set  $\beta_t = \beta_{gp} \cdot \log(|\mathcal{X}| \cdot t^2 \cdot \pi^2)$  with  
819  $\beta_{gp} \in \{0.05, 0.1, 0.5\}$ , reporting results the value that yields the best performance. The choice of  $\beta_t$   
820 is informed by prior work in GP bandits [28]. The selection of  $\sigma$  for all variants is determined by  
821 optimizing the log-likelihood of the observed after very 10 rounds by considering values in the set  
822  $\{0.01, 0.05, 0.1\}$ .

### 823 C.3 Additional results

824 **Local search** results for optimizing combinatorial objectives in  $\Pi^k$  for  $n = 50$  and  $k = 6$ . Specif-  
825 ically,  $\pi^* = \max_{\pi} \phi^r(\pi)^T \phi^r(\pi')$ , where  $\phi^r(\pi')$  represents the feature vector for Kendall kernels  
826 on top-k rankings. Notably, for this optimization problem, it is known that the optimal value is 1  
827 obtained by only  $\pi'$ . Figure 4 shows results for this optimization problem when applied to WK, CK,  
828 and WCK kernels.

829 **Reward results** for large arm space for the nDCG + diversity reward. Similar to Figure 3, a large  
830 setup with  $n = 50$  for  $k = 3$  and  $k = 6$ , is considered. For  $k = 6$ , the possible arms are over  
831  $1.1 \times 10^{10}$  possible top-k rankings. Given the vastness of this arm space, computing the optimal  
832 arm for the diversity reward is not straightforward. Therefore, we focus on reporting the cumulative



**Figure 5:** Comparative evaluation of bandit algorithms for large arm spaces for the nDCG + diversity reward, with  $> 1.1 \times 10^5$  for the left plot and  $> 1.1 \times 10^{10}$  for the right plot, respectively. Cumulative reward with respect to the rounds of the bandit algorithm is depicted. Results are averaged over 6 trials. In both settings, the WCK approach outperforms other baselines. For more details, see the textual description.

833 reward in Figure 5. We implement this setup using a Local search in batch mode, updating every 5  
834 round and considering a substantial horizon of  $T = 100$  rounds. Specifically, we use 5 restarts, 5  
835 steps in every search direction, and start with 1000 initial candidates. Figure 5 shows that the WCK  
836 approach demonstrates superior performance, continuing to learn effectively even after extensive  
837 rounds.

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 919 institution) were obtained?  
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