

# Alda: Integrating Logic Rules with Everything Else, Seamlessly (System Demonstration)\*

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Sets and rules have been used for easier programming since the late 1960s. While sets are central to database programming with SQL and are also supported as built-ins in high-level languages like Python, logic rules have been supported as libraries or in rule-based languages with limited extensions for other features. However, rules are central to deductive database and knowledge base programming, and better support is needed.

This system demonstration highlights the design of a powerful language, Alda [16, 14], that supports logic rules together with sets, functions, updates, and objects, all as seamlessly integrated built-ins, including concurrent and distributed processes. The key idea is to allow sets of rules to be defined in any scope, support predicates in rules as set-valued variables that can be used and updated directly, and support queries using rules as either explicit or implicit automatic calls to an inference function.

Alda has a formal semantics [15] and is implemented by building on an object-oriented language (DistAlgo [13, 3] extending Python [18]) and an efficient logic rule system (XSB [19, 20]). It has been used successfully on benchmarks and problems from a wide variety of application domains—including those in OpenRuleBench [6], role-based access control (RBAC) [1, 4], and program analysis—with generally good performance [17]. Our implementation and benchmarks are publicly available [23].

This system demonstration shows how Alda is used for OpenRuleBench benchmarks, ANSI standard for role-based access control, and program analysis for large Python programs, including with persistence support for large datasets, all programmed seamlessly without boiler-plate code. For comparisons with related work on rule languages and benchmarking, see [16, 14, 17].

## An example

Figure 1 shows an example program in Alda. It is for a small portion of the ANSI standard for role-based access control (RBAC) [1, 4]. It shows the uses (with line numbers in parentheses) of

- classes (1-8, 9-21) with inheritance (9, 11), and object creation (22) with setup (2-3, 10-12);
- sets, including relations (3, 12);
- methods, including procedures (5-6, 13-14) and functions (7-8, 18-19, 20-21), and calls (23, 24);
- updates, including initialization (3, 12) and membership changes (6, 14); and
- queries, including set queries (8, 19 after union “+”, 21) and queries using rules (19 before “+”);

where the rules are defined in a rule set (15-17), explained in the next part.

Note that queries using set comprehensions (e.g., on lines 8, 19, 21) can also be expressed by using rules and inference, even though comprehensions are more widely used. However, only some queries

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\*This work was supported in part by NSF under grants CCF-1954837, CCF-1414078, and IIS-1447549 and ONR under grants N00014-21-1-2719, N00014-20-1-2751, and N00014-15-1-2208.

```

1 class CoreRBAC:                                # class for Core RBAC component/object
2     def setup():                                # method to set up the object, with no arguments
3         self.USERS, self.ROLES, self.UR := {}, {}, {}
4         # set users, roles, user-role pairs to empty sets
5     def AddRole(role):                           # method to add a role
6         ROLES.add(role)                         # add the role to ROLES
7     def AssignedUsers(role):                     # method to return assigned users of a role
8         return {u: u in USERS | (u,role) in UR} # return set of users having the role
9         ...
10
11 class HierRBAC extends CoreRBAC: # Hierarchical RBAC extending Core RBAC
12     def setup():
13         super().setup()                        # call setup of CoreRBAC, to set sets as in there
14         self.RH := {}                         # set ascendant-descendant role pairs to empty set
15     def AddInheritance(a,d):                   # to add inherit. of an ascendant by a descendant
16         RH.add((a,d))                         # add pair (a,d) to RH
17     rules trans_rs:                           # rule set defining transitive closure
18         path(x,y) if edge(x,y)                # path holds for (x,y) if edge holds for (x,y)
19         path(x,y) if edge(x,z), path(z,y)     # ... if edge holds for (x,z) and for (z,y)
20     def transRH():                             # to return transitive RH and reflexive role pairs
21         return infer(path, edge=RH, rules=trans_rs) + {(r,r): r in ROLES}
22     def AuthorizedUsers(role):                 # to return users having a role transitively
23         return {u: u in USERS, r in ROLES | (u,r) in UR and (r,role) in transRH()}
24         ...
25
26 h = new(HierRBAC, [])                          # create HierRBAC object h, with no args to setup
27 h.AddRole('chair')                            # call AddRole of h with role 'chair'
28 ...
29 h.AuthorizedUsers('chair')                    # call AuthorizedUsers of h with role 'chair'
30 ...

```

Figure 1: An example program in Alda, for Role-Based Access Control (RBAC). In `rules trans_rs`, the first rule says there is a path from  $x$  to  $y$  if there is an edge from  $x$  to  $y$ , and the second rule says there is a path from  $x$  to  $y$  if there is an edge from  $x$  to  $z$  and there is an edge from  $z$  to  $y$ . The call to `infer` queries and returns the set of pairs for which path holds given that edge holds for exactly the pairs in set `RH`, by doing inference using rules in `trans_rs`.

using rules and inference can be expressed by using comprehensions; queries using recursive rules (e.g., on lines 16-17) cannot be expressed using comprehensions.

## Rules with sets, functions, updates, and objects

In Alda, rules are defined in rule sets, each with a name and optional declarations for the predicates in the rules.

```

ruleset ::= rules name (declarations): rule+
rule ::= p(arg1, ..., arga) if p1(arg11, ..., arg1a1), ..., pk(argk1, ..., argkak)

```

In the rule form,  $p$ ,  $p_1$ , ...,  $p_k$  denote predicates,  $p(arg_1, ..., arg_a)$  denotes that  $p$  holds for its tuple of arguments, and `if` denotes that its left-side conclusion holds if its right-side conditions all hold. In a rule set, predicates not in any conclusion are called base predicates; the other predicates are called derived predicates.

The key ideas of seamless integration of rules with sets, functions, updates, and objects are:

1. a predicate is a set-valued variable that holds the set of tuples for which the predicate is true;
2. queries using rules are calls to an inference function, [infer](#), that computes desired values of derived predicates using given values of base predicates;
3. values of base predicates can be updated directly as for other variables, whereas values of derived predicates can only be updated by [infer](#); and
4. predicates and rule sets can be object attributes as well as global and local names, just as variables and functions can.

Declarative semantics of rules are ensured by automatically maintaining values of derived predicates when values of base predicates are updated, by appropriate implicit calls to [infer](#).

For example, in Figure 1, one could use an object field `transRH` in place of calls to `transRH()` in `AuthorizedUsers(role)`, use the following rule set instead of `trans_rs`, and remove `transRH()`.

```
rules transRH_rs:                                # no need to call infer explicitly
    transRH(x,y) if RH(x,y)
    transRH(x,y) if RH(x,z), transRH(z,y)
    transRH(x,x) if ROLES(x)
```

Field `transRH` is automatically maintained at updates to `RH` and `ROLES` by implicit calls to [infer](#).

## Higher-order, patterns, distributed programming, and more

**Higher-order.** Note that predicates in [rules](#) as set-valued variables, e.g., `edge`, and calling [infer](#) to take or return values of set variables, e.g., `RH` in `edge=RH`, avoids the need of high-order predicates or other sophisticated features, e.g., [2], to reuse rules for different predicates in logic languages.

**Patterns.** Alda also supports tuple patterns for set elements in set queries (as in `DistAlgo` [13]) and in queries using rules, e.g., `(1,=x,y)` in `p` matches any triple in set `p` whose first element is 1 and whose second element equals the value of `x`, and binds `y` to the third element if such a triple exists.

**Distributed programming.** Of course, by building on `DistAlgo`, Alda also supports distributed programming with distributed processes, message passing, and high-level queries of message histories, e.g., for distributed RBAC [7, 9], also called trust management [5], in decentralized systems.

**Declarations for predicates in rules.** Declarations in [rules](#) could specify predicate types and scopes, but are designed more importantly for specifying assumptions about predicates being certain, complete, closed, or not [10, 11, 12]. This is to give respective desired semantics for rules with unrestricted negation, quantification, and aggregation.

**Python syntax.** Note that the examples discussed use an ideal syntax, while the Alda implementation supports the Python syntax. For example, `x := {}` is written as `x = set()` in Python syntax.

**Implementation.** The Alda implementation compiles rule sets in [rules](#) and queries using [infer](#) to XSB rules and queries, and compiles the rest to Python, which calls XSB to do the inference. The current implementation supports primarily Datalog rules, but also handles unrestricted negation by using XSB's computation of the well-founded semantics [24]. More general forms of rules and queries can be compiled to rules and queries in XSB or other rule systems using the same approach. In general, any efficient inference algorithm and implementation method can be used to compute the semantics of [rules](#) and [infer](#).

**Future work.** Future work includes (1) support for easy use of different desired semantics, especially with modular use of rules, similar to knowledge units in DA-logic [11]; and (2) efficient implementation with complexity guarantees [8, 21, 22] for computing different desired semantics.

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