

A Generalized False Data Injection Attacks Against Power System Nonlinear State Estimator and Countermeasures

Junbo Zhao, *Student Member, IEEE*, Lamine Mili, *Life Fellow, IEEE*, and Meng Wang, *Member, IEEE*

Abstract—This paper develops a generalized framework that allows us to investigate the vulnerability of the power system nonlinear state estimator to false data injection attacks (FDIAs) from the operator's perspective and to initiate some countermeasures. Unlike most existing FDIA methods, which assume a perfect knowledge of the system measurements and topology by a hacker, we derive and analyze the uncertainties for launching successful FDIA along with their upper bounds. To effectively defend against an FDIA, we propose a robust detector that checks the measurement statistical consistency using a subset of secure PMU measurements. We first show that if these secure PMU measurements are free of bad data while making the system observable, the FDIA is detectable. We then show that detectability is also ensured if these conditions are relaxed while using alternative redundant measurements from short-term nodal synchrophasor predictions together with the robust Huber M-estimator. Numerical simulation results on the IEEE 30-bus and 118-bus systems demonstrate the effectiveness and robustness of the proposed method even the secure measurements contain noise and bad data.

Index Terms—Cyber security, false data injection attacks (FDIAs), power system nonlinear state estimation, robust estimation, phasor measurement units, Neyman-Pearson detector.

I. INTRODUCTION

Due to a strong reliance of smart grid functions on communication networks, cyber attacks have become a major concern among power researchers. The analysis of cyber attacks on power system state estimation (SE) was pioneered by Liu *et al.* [1], where the so-called false data injection attack (FDIA) was introduced. Following this work, three types of FDIA were pinpointed and investigated, including state attacks [1], [2], topology attacks [3], [4] and load redistribution attacks [5]. Their impacts on the electricity markets were also analyzed [6], [7].

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J. Zhao and L. Mili are with the Bradley Department of Electrical Computer Engineering, Virginia Polytechnic Institute and State University, Northern Virginia Center, Falls Church, VA 22043 USA (e-mail: zjunbo@vt.edu; lmili@vt.edu).

M. Wang is with the Department of Electrical, Computer, and Systems Engineering, Rensselaer Polytechnic Institute, Troy, NY 12180 USA (e-mail: wangm7@rpi.edu).

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To safeguard the system operation and control against cyber attacks, various detectors and mitigation methods have been proposed. They include measurement protection-based approaches [8]–[11], sparse optimization or game theory-based approaches [12], [13], innovation-based approaches [14], [15], robust estimation-based approaches [16], to name a few. However, with the exception of [17]–[20], the bulk of the literature focused on the linear DC rather than on the AC state estimator. In [17], the vulnerability of the nonlinear SE to FDIA was analyzed and discussed. It was shown that a hacker should know perfectly all the state variables in the attack subgraph to conduct a successful attack. In other words, she should have exactly the same information about the power system as the operators of the control center, including an exact knowledge about the measurements and the system topology. In this paper, this type of attack will be called the perfect FDIA attack. The work in [17] was later extended by Rahman *et al.* [18], Zhao *et al.* [19], and Xuan *et al.* [20] to take into account the uncertainties in the system information that may be gathered by a hacker, resulting in an imperfect FDIA. However, no analytical investigations were carried out to explain why an imperfect attack may succeed and under which conditions it will be detected by the operators of the control center. Furthermore, defense approaches have been studied extensively for a linear DC model-based FDIA. But little work has been done for a nonlinear AC model-based FDIA. Note that the practical control center uses nonlinear power system state estimator for monitoring and control. It is thus of vital importance to ensure the security and reliability of that estimator.

In this paper, an analytical framework is proposed to investigate the vulnerability of power system nonlinear state estimator to an FDIA from the operator's perspective. In particular, we propose a generalized FDIA framework against the nonlinear state estimator. In this framework, the perfect knowledge of the system information is relaxed to account for measurement, parameter and topology uncertainties. The latter may be induced by the hacker's limited real-time knowledge of the status of various grid elements or restricted access to communication channels [14]. The upper bounds of these uncertainties for launching a successful FDIA are quantified and analyzed as well. To effectively detect an FDIA, we propose a robust detector by checking the measurement statistical consistency using a subset of secure PMU measurements. It is shown that these secure measurements allow us to detect an FDIA if they are free of gross errors while making the system observable. These con-

ditions are further relaxed by using a robust Huber M-estimator together with alternative redundant measurements from short-term nodal synchrophasor predictions. Interestingly, robust state estimates provided by Huber M-estimator is shown to follow a Gaussian distribution, which enables us to derive the analytical form of the Neyman-Pearson detector.

The remainder of this paper is organized as follows: Section II introduces the existing FDIA against nonlinear state estimator and presents the problem statement. Section III presents the proposed generalized FDIA framework, while Section IV presents the proposed robust FDIA detector. The simulation results are analyzed in Section V, and finally Section VI concludes the paper.

II. PROBLEM FORMULATION

A. Power System Nonlinear State Estimation

As shown in [21], for an N -bus power system using an AC power flow model, the relationship between the vector of measurements $\mathbf{z} \in \mathbb{R}^m$ obtained from the supervisory control and data acquisition (SCADA) system and the state vector $\mathbf{x} \in \mathbb{R}^n$, which contains the nodal voltage magnitudes and phase angles, yielding $n = 2N - 1 < m$, is given by

$$\mathbf{z} = \mathbf{h}(\mathbf{x}) + \mathbf{e}, \quad (1)$$

where $\mathbf{h}(\cdot) : \mathbb{R}^n \rightarrow \mathbb{R}^m$ is a vector-valued nonlinear function; $\mathbf{e} \in \mathbb{R}^m$ is the measurement error vector that is assumed to follow a Gaussian distribution with zero mean and covariance matrix $\mathbf{R} \in \mathbb{R}^{m \times m}$, i.e., $\mathbf{e} \sim \mathcal{N}(\mathbf{0}, \mathbf{R})$. The state estimator is solved by minimizing the weighted least squares criterion, yielding

$$\hat{\mathbf{x}} = \arg \min_{\mathbf{x}} [\mathbf{z} - \mathbf{h}(\mathbf{x})]^T \mathbf{R}^{-1} [\mathbf{z} - \mathbf{h}(\mathbf{x})]. \quad (2)$$

Let us apply the Gauss-Newton iterative algorithm [21] to solve for the state vector. Formally, we have

$$\mathbf{x}^{k+1} = \mathbf{x}^k + \Delta \mathbf{x}^k, \quad k = 1, 2, \dots, \quad (3)$$

$$\Delta \mathbf{x}^k = (\mathbf{H}(\mathbf{x}^k)^T \mathbf{R}^{-1} \mathbf{H}(\mathbf{x}^k))^{-1} \mathbf{H}(\mathbf{x}^k)^T \mathbf{R}^{-1} (\mathbf{z} - \mathbf{h}(\mathbf{x}^k)), \quad (4)$$

where $\mathbf{H}(\mathbf{x}^k) = \partial \mathbf{h}(\mathbf{x}) / \partial \mathbf{x}|_{\mathbf{x}=\mathbf{x}^k} \in \mathbb{R}^{m \times n}$ is the Jacobian matrix. The algorithm converges once the norm of $\Delta \mathbf{x}^k$ is smaller than a pre-specified threshold. After estimation, the ℓ_2 -norm detector is applied to detect the existence of bad data by checking if the following inequality holds [6], [21]:

$$\|\mathbf{r}\| = \|\mathbf{z} - \mathbf{h}(\hat{\mathbf{x}})\| \geq \tau, \quad (5)$$

where τ is a detection threshold of the ℓ_2 -norm detector. Note that $\|\cdot\|$ is used to represent the ℓ_2 -norm throughout the paper, where the subscript 2 in (5) has been dropped for simplicity.

B. Attack Model of the Nonlinear State Estimator

To perform an FDIA, we make the same assumptions as that in [6], [17], that is: i) an attacker could access the real-time measurements in a small area \mathcal{S} bounded by buses, where the measurement and state indices in \mathcal{S} are denoted as \mathcal{M}_s and \mathcal{I}_s , respectively; ii) the hacker could change all the measurements

in \mathcal{S} ; iii) the hacker might have an a priori knowledge of the system topology, including the line parameters of the area \mathcal{S} . Thus, for the i th measurement, z_i , the attack model is

$$z_i^{(a)} = \begin{cases} z_i & \text{if } i \notin \mathcal{M}_s \\ z_i + a_i & \text{if } i \in \mathcal{M}_s \end{cases}, \quad (6)$$

where a_i is the i th element of the attack vector \mathbf{a} .

Lemma 1: Let us now assume that the hacker has obtained the same $z_i, i \in \mathcal{M}_s$ and $\hat{\mathbf{x}}_i$ as the operators of the control center. If the original measurement $z_i, i \in \mathcal{M}_s$, could bypass the ℓ_2 -norm detector, the malicious measurement $z_i^{(a)}$ could also pass this detector under the condition $a_i = \mathbf{h}(\hat{\mathbf{x}}_i + c_i) - \mathbf{h}(\hat{\mathbf{x}}_i)$, where c_i represents the changes in the i th-attacked state variable.

Proof: Since we are interested in the area \mathcal{S} , the index i is omitted for simplicity. Because \mathbf{z} can bypass the ℓ_2 -norm detector, $\|\mathbf{r}\| = \|\mathbf{z} - \mathbf{h}(\hat{\mathbf{x}})\| \leq \tau$ holds. The ℓ_2 -norm of the attacked measurement residual \mathbf{r}_a is given by

$$\begin{aligned} \|\mathbf{r}_a\| &= \|\mathbf{z}^a - \mathbf{h}(\hat{\mathbf{x}}_a)\| = \|\mathbf{z} + \mathbf{a} - \mathbf{h}(\hat{\mathbf{x}} + \mathbf{c})\| \\ &= \|\mathbf{z} + \mathbf{a} - \mathbf{h}(\hat{\mathbf{x}} + \mathbf{c}) + \mathbf{h}(\hat{\mathbf{x}}) - \mathbf{h}(\hat{\mathbf{x}})\| \\ &= \|\mathbf{r} + \mathbf{a} - \mathbf{h}(\hat{\mathbf{x}} + \mathbf{c}) + \mathbf{h}(\hat{\mathbf{x}})\| \\ &= \|\mathbf{r}\| \leq \tau, \end{aligned} \quad (7)$$

which means that the attacked measurements could also avoid the detection. Note that $\hat{\mathbf{x}}_a = \hat{\mathbf{x}} + \mathbf{c}$ and $\mathbf{r} = \mathbf{z} - \mathbf{h}(\hat{\mathbf{x}})$ is the measurement residual vector. ■

When implementing an FDIA for practical power systems, Lemma 1 intrinsically assumes that the hacker has enough computational capability to estimate the local state vector $\hat{\mathbf{x}}_i, i \in \mathcal{I}_s$ so that the attack vector $\mathbf{a} = \mathbf{h}(\hat{\mathbf{x}}_i + c_i) - \mathbf{h}(\hat{\mathbf{x}}_i)$ can be constructed. This assumption is acceptable given that the SCADA measurements are non-synchronized while the collection rates of measurements differ from one region to another one. In addition, a hacker may intentionally attack the communication system to delay the SCADA measurements for some parts of the system so that the local state x_i can be estimated by the hacker [6], [17].

C. Determining the Attack Graph of the Target Buses

Let $\mathcal{S} = \{\mathcal{B}, \Omega\}$ denote the attack graph, where \mathcal{B} and Ω are the sets of buses and transmission lines, respectively; let \mathcal{K} denote a set of bus indices for power injection buses, including the load and the generator buses. In [17], a topographical analysis was proposed to determine the attack graph of a single target bus. That approach is summarized below:

- *Step 1:* Let $i \in \mathcal{K}$ be the i th targeted power injection bus, the first step is to include bus i into the subgraph \mathcal{S}_i ;
- *Step 2:* Extend \mathcal{S}_i to include all the buses and branches Ω_i that are connected to bus i , where Ω_i is the set of adjacent branches connected to bus i ;
- *Step 3:* If there exists any zero injection Bus j not connected to either the load or the generator on the boundary of \mathcal{S}_i , extend \mathcal{S}_i to include Ω_j and continue Step 4; otherwise go to Step 5;

- *Step 4:* Repeat Step 3 until all buses on the boundary belong to the set \mathcal{K} ;
- *Step 5:* Obtain the final attack subgraph as $\mathcal{S} = \bigcup_{i \in \mathcal{K}} \mathcal{S}_i$.

174 The above procedures can be simply summarized by the following lemma:

175 **Lemma 2:** For a bus provided with a power injection, its
176 adjacent buses provided with power injections must be changed
177 accordingly so that specified state changes can be made by the
178 hacker. For a zero-injection bus, its adjacent power flows must
179 sum to zero. This means that the measurements of the power
180 injection buses adjacent to a zero power injection bus must be
181 changed accordingly so that the equality constraints are satisfied.
182 Consequently, buses that belong to \mathcal{S} are bounded by buses in \mathcal{K} .

183 Using Lemma 2, we are able to determine the attack graph
184 of multi-buses. The difference is that the size of the set \mathcal{S} is
185 increased with additional buses bounded by power injections.

187 D. Problem Statement

188 With the obtained attack graph \mathcal{S} and all the assumptions
189 stated in Section II-B satisfied, a perfect FDIA against a nonlin-
190 ear state estimator can be achieved. However, because a hacker
191 has typically a lack of real-time knowledge of the status of grid
192 elements such as the position of circuit breaker switches and
193 transformer tap changers, and also because she is restricted to
194 access partial measurement channels, it is thus impossible for
195 her to obtain the same state estimates as the operators of the
196 control center in the attack graph \mathcal{S} . In other words, a perfect
197 FDIA approach proposed in the literature [6], [17] seems to be
198 impractical for realistic power systems. This is because with
199 uncertain information of the system, $\hat{x}_i, i \in \mathcal{I}_s$ obtained by the
200 hacker is different from the state estimate $\hat{x}_i^w, i \in \mathcal{I}_s$, and bias ζ_i
201 exists, i.e., $\hat{x}_i = \hat{x}_i^w + \zeta_i$. Note that $\hat{x}_i^w, i \in \mathcal{I}_s$ is the i -th state
202 estimate calculated by the control center without an FDIA. As a
203 result, when an FDIA occurs, the inequality constraint (7) may
204 not hold true anymore. Interestingly, simulations carried out
205 in [18], [20] reveal that even with some uncertain information
206 about the system, FDIA can be successful without being de-
207 tected by the control center. Furthermore, the expected changes
208 on the target state variables are not equal to c . However, no
209 analytical investigations were carried out to explain why this
210 imperfect attack can succeed and under which conditions it will
211 be detected by the operators of the control center.

212 In this paper, an analytical investigation will be performed to
213 show how the inequality constraint (7) can be satisfied in pres-
214 ence of system uncertainties to avoid the detection of an FDIA
215 by the control center. In addition, we will quantify the maxi-
216 mum uncertainties a hacker can have so as to perform imperfect
217 FDIA. The trade-off between attack magnitudes on the target
218 state variables and the system uncertainties will be analyzed
219 as well. Finally, to detect this type of FDIA, we will propose a
220 measurement statistical consistency-based robust detector using
221 a subset of secure PMU measurements.

222 *Remark:* To avoid the confusion between the bias terms ζ
223 and c , we make the following clarifications: \hat{x} is the estimated
224 state vector before an FDIA and it is equal to \hat{x}^w obtained by
225 the control center if the hacker has the same information of the
226 system as the control center. Otherwise, there is a difference

227 between \hat{x} and \hat{x}^w caused by information uncertainties, which
228 is the bias ζ . By contrast, c is the expected bias by the hacker
229 when performing an FDIA.

230 III. PROPOSED GENERALIZED FDIA FRAMEWORK AGAINST 231 THE NONLINEAR STATE ESTIMATOR

232 An FDIA is in fact a type of perfect interacting and conform-
233 ing bad data [22]. Therefore, the statistical tests applied to the
234 weighted or the normalized residuals or the sum of the squared
235 residuals (ℓ_2 detector) are unable to detect them. Without loss of
236 generality, we consider in the sequel only the ℓ_2 detector when
237 deriving the generalized FDIA framework.

238 In the developed generalized FDIA framework, we first pro-
239 vide a sufficient condition in Section III-A to justify theoretically
240 how the imperfect FDIA can bypass the detector at the control
241 center. This allows us to derive the upper bound of the uncertain-
242 ties the hacker can have so as to launch a successful imperfect
243 FDIA. Therefore, we are able to analyze the trade-off between
244 attack magnitudes on the target state variables and the system
245 uncertainties. To our best knowledge, this is the first attempt
246 to provide theoretical justification to an imperfect FDIA and to
247 quantify the tradeoff between the attack magnitude and the state
248 bias caused by system uncertainties.

249 A. Sufficient Condition for an Imperfect FDIA

250 As clarified before, an adversary cannot obtain the same es-
251 timated state \hat{x}^w (the subscript i is dropped for simplicity) as
252 the operators of the control center. Here, we provide a sufficient
253 condition for an imperfect FDIA to succeed subject to the state
254 bias. This is shown in the following Lemma:

255 **Lemma 3:** If the true measurement residual $\|r\| =$
256 $\|z - h(\hat{x}^w)\| \leq \tau$ holds, a sufficient condition for the mea-
257 surement z subject to attack a to pass ℓ_2 detector is

$$258 \|a - h(\hat{x}^w + c) + h(\hat{x}^w)\| \leq \gamma = \tau - \|r\|. \quad (8)$$

259 *Proof:* When there are no bad data in the original measure-
260 ments, $\|r\| \leq \tau$ is always satisfied. The measurement residual
261 under FDIA can be derived as

$$262 \begin{aligned} \|r_a\| &= \|z_a - h(\hat{x}_a)\| = \|z + a - h(\hat{x}^w + c)\| \\ &= \|z + a - h(\hat{x}^w + c) + h(\hat{x}^w) - h(\hat{x}^w)\| \\ &= \|r + a - h(\hat{x}^w + c) + h(\hat{x}^w)\| \\ &\leq \|r\| + \|a - h(\hat{x}^w + c) + h(\hat{x}^w)\| \leq \tau, \end{aligned} \quad (9)$$

263 which means that if the constraint (8) holds, an FDIA can not be
264 detected by the residual statistical bad data detection test. ■

265 Although Lemma 3 looks straightforward, it provides a suf-
266 ficient condition that the attack vector a should satisfy to avoid
267 her detection by the operators. In addition, it serves as the foun-
268 dation for the derivation of the upper bound of uncertainties
269 the hacker can have when implementing a successful imper-
270 fect FDIA. On the other hand, it is easy to verify that a perfect
271 attack with $\hat{x}^w = \hat{x}$ mentioned in Lemma 1 is just a special
272 case here. Finally, since \hat{x}^w is unknown to the hacker due to
273 the limited knowledge of the system, conditions with the con-
274 sideration of system uncertainties should be derived. That is,
275 what is the tradeoff between system uncertainties and attack

274 magnitude? This question will be investigated and analyzed in
275 Section III-B.

276 *B. Tradeoff Between System Uncertainties and
277 Attack Magnitude*

278 Due to the existence of uncertain system information obtained
279 by an attacker, the initial state vector used to construct an attack
280 vector has uncertainties as well. This in turn yields biases on
281 the target state variables. The larger uncertainties the hacker
282 has, the less possibilities she can change the attack magnitudes,
283 and vice versa. In other words, there exists a tradeoff between
284 system uncertainties and attack magnitude. To quantify that, we
285 define $\mathbf{a} = \mathbf{h}(\hat{\mathbf{x}} + \mathbf{c}) - \mathbf{h}(\hat{\mathbf{x}})$ and $\hat{\mathbf{x}} = \hat{\mathbf{x}}^w + \zeta$. The ℓ_2 -norm of
286 the measurement residual becomes

$$\begin{aligned} \|r_a\| &= \|z + \mathbf{a} - \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c})\| \\ &= \|z + \mathbf{a} - \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c}) + \mathbf{h}(\hat{\mathbf{x}}^w) - \mathbf{h}(\hat{\mathbf{x}}^w)\| \\ &= \|z - \mathbf{h}(\hat{\mathbf{x}}^w) + (\mathbf{a} - \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c}) + \mathbf{h}(\hat{\mathbf{x}}^w))\| \\ &\leq \|z - \mathbf{h}(\hat{\mathbf{x}}^w)\| + \|\mathbf{a} - \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c}) + \mathbf{h}(\hat{\mathbf{x}}^w)\| \\ &= \|r\| + \|\mathbf{h}(\hat{\mathbf{x}} + \mathbf{c}) - \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c}) - (\mathbf{h}(\hat{\mathbf{x}}) - \mathbf{h}(\hat{\mathbf{x}}^w))\|. \end{aligned} \quad (10)$$

287 Performing Taylor series expansions of $\mathbf{h}(\hat{\mathbf{x}} + \mathbf{c})$ and $\mathbf{h}(\hat{\mathbf{x}})$ at
288 $\hat{\mathbf{x}}^w + \mathbf{c}$ and $\hat{\mathbf{x}}^w$, respectively, we obtain

$$\begin{aligned} \mathbf{h}(\hat{\mathbf{x}} + \mathbf{c}) - \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c}) &= \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c}) + \mathbf{H}_1(\hat{\mathbf{x}} - \hat{\mathbf{x}}^w - \mathbf{c}) + \mathbf{o}_1 - \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c}) \\ &= \mathbf{H}_1(\zeta - \mathbf{c}) + \mathbf{o}_1 \\ \mathbf{h}(\hat{\mathbf{x}}) - \mathbf{h}(\hat{\mathbf{x}}^w) &= \mathbf{h}(\hat{\mathbf{x}}^w) + \mathbf{H}_2(\hat{\mathbf{x}} - \hat{\mathbf{x}}^w) + \mathbf{o}_2 - \mathbf{h}(\hat{\mathbf{x}}^w) \\ &= \mathbf{H}_2\zeta + \mathbf{o}_2, \end{aligned} \quad (11)$$

289 where $\mathbf{H}_1 = \partial\mathbf{h}/\partial\mathbf{x}|_{\mathbf{x}=\hat{\mathbf{x}}^w+\mathbf{c}}$ and $\mathbf{H}_2 = \partial\mathbf{h}/\partial\mathbf{x}|_{\mathbf{x}=\hat{\mathbf{x}}^w}$ are
290 Jacobian matrices; \mathbf{o}_1 and \mathbf{o}_2 are the higher order Taylor expansion
291 terms. Since only the first order approximation is used in
292 the WLS based state estimation algorithm, all the higher order
293 terms are neglected during the iteration. In other words, \mathbf{o}_1 and
294 \mathbf{o}_2 tend to 0 faster than the convergence of state estimation.
295 Therefore,

$$\begin{aligned} \|\mathbf{h}(\hat{\mathbf{x}} + \mathbf{c}) - \mathbf{h}(\hat{\mathbf{x}}^w + \mathbf{c}) - (\mathbf{h}(\hat{\mathbf{x}}) - \mathbf{h}(\hat{\mathbf{x}}^w))\| &= \|(\mathbf{H}_1 - \mathbf{H}_2)\zeta - \mathbf{H}_1\mathbf{c} + (\mathbf{o}_1 - \mathbf{o}_2)\| \\ &\cong \|(\mathbf{H}_1 - \mathbf{H}_2)\zeta - \mathbf{H}_1\mathbf{c}\| \\ &\leq \|\mathbf{H}_1 - \mathbf{H}_2\| \|\zeta\| + \|\mathbf{H}_1\| \|\mathbf{c}\|. \end{aligned} \quad (12)$$

296 By combining (10) and (12), we get

$$\|r_a\| \leq \|r\| + \|\mathbf{H}_1 - \mathbf{H}_2\| \|\zeta\| + \|\mathbf{H}_1\| \|\mathbf{c}\|. \quad (13)$$

297 In order not to be detected by the operators of the control cen-
298 ter, the right-hand side of (13) must be less than the detection
299 threshold τ , that is,

$$0 \leq \|\mathbf{H}_1 - \mathbf{H}_2\| \|\zeta\| + \|\mathbf{H}_1\| \|\mathbf{c}\| \leq \gamma. \quad (14)$$

300 The above equation shows the tradeoff between the attack mag-
301 nitude and the estimation error $\|\zeta\|$ of the state variables in the

attack graph \mathcal{S} . Note that $\|\zeta\|$ is caused by an uncertain infor-
302 mation about the system state. If the estimation error is fixed,
303 i.e., $\|\zeta\| = \beta \neq 0$, the attack magnitude is bounded by
304

$$0 \leq \|\mathbf{c}\| \leq -\beta + \frac{\gamma + \|\mathbf{H}_2\| \beta}{\|\mathbf{H}_1\|}. \quad (15)$$

305 If $\|\zeta\| = \beta = 0$, which means that the hacker can get exactly
306 the same state estimates as the operators of the control center in
307 the attack graph \mathcal{S} , the attack reduces to the perfect FDIA. The
308 attack magnitude is bounded according to either (14) or (15) by
309 setting $\beta = 0$. Formally, we have

$$0 \leq \|\mathbf{c}\| \leq \frac{\gamma}{\|\mathbf{H}_1\|}. \quad (16)$$

310 *Remark:* Note that the main scope of this paper is to investi-
311 gate the vulnerability of a nonlinear WLS state estimator to
312 imperfect FDIA from the operator's perspective. To be more
313 specific, given the current states estimated from the measure-
314 ments and the assumed attack magnitudes, the operator knows
315 matrices \mathbf{H}_1 and \mathbf{H}_2 . Then he can analyze how large uncertain-
316 ties the hacker can have so that a success FDIA is launched under
317 this condition. He may vary the assumed attack magnitudes to
318 assess how the maximum uncertain information of the hacker
319 changes if a successful FDIA is initiated. As a result, the vul-
320 nerability of the estimator can be assessed. On the other hand, due
321 to the existence of uncertain system information and the limited
322 access to measurements, the hacker is unable to know the exact
323 matrices \mathbf{H}_1 and \mathbf{H}_2 . However, as long as the inequality (14)
324 holds true, he can initiate successful FDIA with inexact \mathbf{H}_1 and
325 \mathbf{H}_2 . This analysis can warn the operator to pay attention to the
326 potential FDIA as the hacker is able to launch successful FDIA
327 even with uncertain system information and limited measure-
328 ments. To this end, corresponding effective countermeasures
329 can be proposed.

IV. PROPOSED ROBUST FDIA DETECTOR

330 In this section, we first present the motivations of designing
331 a robust FDIA detector with a limited number of secure PMUs.
332 The challenges and solution methodologies associated with the
333 detector are discussed thoroughly. Then, the robust FDIA de-
334 tector using measurement statistical consistency is proposed.
335 To derive this detector, we enhance the data redundancy of the
336 PMU measurements by short-term measurement forecasting,
337 which allows us to handle noise and outliers in secure PMU
338 measurements. We show through Theorem 1 that our robust
339 state estimates follow a Gaussian distribution even when the
340 PMU measurement errors are not normally distributed. This en-
341 ables us to derive the Neyman-Pearson detector for an FDIA
342 detection.

A. Motivations and Challenges

344 Recall that the hacker's objective is to change the estimated
345 state vector by injecting malicious measurements. Once some
346 measurements are compromised, the distribution of the esti-
347 mated state vector will be perturbed by the attack [2], [14]. If
348 one can find a set of measurements that will produce close ap-
349 proximations to the true estimated state vector and its probability
350 distribution, then this statistical information can be further used
351

352 to double check the estimation results obtained from the remaining
 353 measurements. This strategy shares similar characteristics
 354 of the machine learning techniques, where partial data is used
 355 for learning and training while the remaining data is leveraged
 356 for validation. On the other hand, with the increasing deployment
 357 of PMUs in power systems, most of these systems are
 358 expected to be observable by PMUs. Indeed, many real systems
 359 have been observed by PMUs such as the Virginia Dominion
 360 Power [23], the 765/345/230 kV power grid in New York (NY),
 361 and the 345 kV power grid in New England (NE) [24], [25], to
 362 cite a few. In addition, the PMU observability of a given power
 363 system has been widely assumed in the literature, see [24], [29]
 364 for example. Thus, in this paper, we assume as suggested in [9]
 365 that the system is observable by a minimal set of PMUs that are
 366 made secure against cyber attacks. These PMUs that are usually
 367 installed at transmission system substations can be protected
 368 by encryptions, advanced fire walls and data package anomaly
 369 detectors, to name a few [26]–[28]. In addition, limited number
 370 of PMUs are assumed to be secure in the sense that they cannot
 371 be controlled by the hacker. On the other hand, unlike [9], we
 372 propose a robust detector based on an AC power system model
 373 that is able to handle outliers. Indeed, the authors in [9] make
 374 use of a DC not an AC state estimation model. Furthermore, they
 375 suppose that the PMU measurements are free of bad data, which
 376 is unrealistic in practice since impulsive communication noise
 377 and faulty GPS synchronization may corrupt the metered values.
 Note that strongly biased state estimates may result, which
 378 will mislead the operators of a control center; consequently, they
 379 may take wrong decisions based on them [30].

381 *Remark:* We assume the number of secure PMUs only guarantees the observability of the system, yielding no measurement
 382 redundancy. As a result, the PMU-based linear state estimator is
 383 not able to filter out noise as well as bad data. On the other hand,
 384 the state estimation model itself is an approximate model with
 385 uncertainties in the parameters, the topology and the measurements.
 Thus, a high measurement redundancy is required to reliably estimate
 386 the system state vector. This motivates us to validate and correct the remaining SCADA measurements, yielding
 387 improved state estimation results and system visualization.

391 B. Robust Detector Using Measurement Statistical 392 Consistency

393 Similarly to the strategy proposed in [9], we assume in the
 394 proposed detector that the system is observed by a minimal set
 395 of secure PMU measurements. On the other hand, we advocate
 396 to enhance the data redundancy of PMUs by short-term mea-
 397 surement forecasting as proposed by Yacine *et al.* [31]. This
 398 allows our robust estimator to handle outliers in secure PMU
 399 measurements. It should be noted that it is in general chal-
 400 lenging to forecast the future operating conditions due to many
 401 changing factors. However, we focus only on a very short-term
 402 forecast of the PMU metered variables, where the system oper-
 403 ating conditions vary slowly, which is a reasonable assumption
 404 for practical power systems. Furthermore, system loads and re-
 405 newable energy-based distributed generations change continu-
 406 ously from time to time, exhibiting temporal correlations. These
 407 changes in turn affect other generators and loads within the same

408 geographic area, yielding spatial correlations. As a result, the
 409 nodal voltage and current phasors of the system exhibit similar
 410 statistical properties, which can be easily proved through the
 411 power flow equations. Thanks to these temporal and spatial cor-
 412 relations, we are able to use time series analysis technique to
 413 perform a short-term forecasting of the PMU metered variables.
 414 Interestingly, the temporal and spatial correlations of the nodal
 415 voltage and current phasors have been proved by [31] through
 416 field measurements. In that reference, an effective forecast of
 417 the PMU metered variables using a vector autoregressive model
 418 has been demonstrated as well. Following [31], we consider a
 419 vector autoregressive model of first order and dimension D at
 420 time instant k , i.e.,

$$421 \mathbf{y}_k = \Phi_k \mathbf{y}_{k-1} + \varepsilon_k, \quad (17)$$

422 where $\mathbf{y}_k \in \mathbb{R}^D$ is the vector of secure PMU measurements;
 423 $\Phi_k \in \mathbb{R}^{D \times D}$ represents the transition matrix; $\varepsilon_k \in \mathbb{R}^D$ is the
 424 Gaussian noise and $\varepsilon_k \sim \mathcal{N}(\mathbf{0}, \mathbf{S}_k)$, where $\mathbf{S}_k \in \mathbb{R}^{D \times D}$ is a
 425 non-diagonal error covariance matrix due to temporal and spatial
 426 correlations among PMU measurements. Using the Yule-Walker
 427 method, $\hat{\Phi}_k$ are estimated using historical measurements [31].
 Then, the forecasted PMU metered values are obtained through
 428 $\mathbf{y}_k^f = \hat{\Phi}_k \mathbf{y}_{k-1}$, while its covariance matrix is given by $\mathbf{P}_k =$
 429 $\hat{\Phi}_k \mathbf{P}_{k-1} \hat{\Phi}_k^T + \mathbf{S}_k$, where \mathbf{P}_{k-1} is the error covariance matrix
 430 of the filtered PMUs at time instant $k-1$. By processing the
 431 metered and the forecasted PMU values simultaneously and
 432 then performing the data prewhitening, we get the following
 433 regression form

$$434 \mathbf{z}_k = \mathcal{H}_k \mathbf{x}_k + \eta_k, \quad (18)$$

435 where $\mathbf{z}_k = \mathbf{L}_k^{-1} \left[(\mathbf{y}_k^s)^T \left(\mathbf{y}_k^f \right)^T \right]^T \in \mathbb{R}^l$ is the extended mea-
 436 surement vector that contains the forecasted measurements \mathbf{y}_k^f
 437 and received PMU measurements \mathbf{y}_k^s ; $l = 2D$; \mathbf{L}_k is a matrix
 438 for prewhitening that is determined by applying a Cholesky
 439 decomposition to the augmented error covariance matrix $\Gamma =$
 440 $\text{diag}[\Lambda_k \mathbf{P}_k] = \mathbf{L}_k \mathbf{L}_k^T$. Here, Λ_k denotes the measurement er-
 441 ror covariance matrix; $\mathcal{H}_k = \mathbf{L}_k^{-1} [\mathbf{A}_k^T \mathbf{A}_k^T]^T$; \mathbf{A}_k is a con-
 442 stant admittance matrix; η_k is the normalized error vector; and
 443 $\eta_k \sim \mathcal{N}(\mathbf{0}, \mathbf{I}_k)$, where \mathbf{I}_k is an identity matrix. Note that,
 444 only the current phasor measurement on a given line having
 445 an impedance that is very different from the rest of the lines can
 446 induce leverage point in the regression equation (18). By using
 447 the scaling technique proposed in [29], those leverage points
 448 can be eliminated and only vertical outliers are of concern.

449 Note that there always exists the case that system operation
 450 conditions vary abruptly, yielding unreliable predicted PMU
 451 measurements. In addition, the received PMU measurements
 452 can be corrupted with gross errors as well due to the impulsive
 453 communication noise, loss of communications, etc. To handle
 454 them while achieving good statistical efficiency, we advocate the
 455 use of the robust Huber-estimator [32]. This estimator minimizes
 456 the following objective function:

$$457 \mathcal{J}(\mathbf{x}) = \sum_{i=1}^l \rho(r_{S_i}), \quad (19)$$

456 where $\rho(\cdot)$ is the Huber convex cost function defined as:

$$\rho(r_{S_i}) = \begin{cases} r_{S_i}^2/2 & \text{for } |r_{S_i}| \leq \lambda \\ \lambda|r_{S_i}| - \lambda^2/2 & \text{for } |r_{S_i}| > \lambda \end{cases}, \quad (20)$$

457 where the parameter λ is typically set to be between 1.5 and 3
458 for achieving high statistical efficiency when the measurement
459 errors are Gaussian [32]; $r_{S_i} = r_i/s$ is the standardized residual;
460 $r_i = \mathbf{z}_i - \boldsymbol{\alpha}_i^T \hat{\mathbf{x}}$ and $\boldsymbol{\alpha}_i^T$ is the i th column vector of the matrix
461 \mathbf{H}_k^T ; $s = 1.4826 c_m$ median $|r_i|$ is the robust scale estimate and
462 c_m is a correction factor [33].

463 To solve (19), the following necessary condition must be
464 satisfied:

$$\frac{\partial \mathcal{J}(\mathbf{x})}{\partial \mathbf{x}} = \sum_{i=1}^l -\frac{\boldsymbol{\alpha}_i}{s} \psi(r_{S_i}) = \mathbf{0}, \quad (21)$$

465 where $\psi(r_{S_i}) = \partial \rho(r_{S_i}) / \partial r_{S_i}$. Multiplying and dividing r_{S_i}
466 on both sides of (21), we obtain

$$\sum_{i=1}^l \boldsymbol{\alpha}_i \frac{\psi(r_{S_i})}{r_{S_i}} \cdot \frac{r_{S_i}}{s} = \mathbf{0}, \quad (22)$$

467 which can be arranged in a matrix form as

$$\mathbf{H}_k^T \mathbf{Q} (\mathbf{z}_k - \mathbf{H}_k \hat{\mathbf{x}}) = \mathbf{0}, \quad (23)$$

468 where $q(r_S) = \psi(r_S)/r_S$ and $\mathbf{Q} = \text{diag}(q(r_S))$. Finally, using
469 the iteratively re-weighted least square (IRLS) algorithm
470 [32], the solution can be obtained through

$$\hat{\mathbf{x}}_k^{(\ell+1)} = (\mathbf{H}_k^T \mathbf{Q}^{(\ell)} \mathbf{H}_k)^{-1} \mathbf{H}_k^T \mathbf{Q}^{(\ell)} \mathbf{z}_k, \quad (24)$$

471 where ℓ is the iteration counter. The algorithm converges if

$$\left\| \hat{\mathbf{x}}_k^{(\ell+1)} - \hat{\mathbf{x}}_k^{(\ell)} \right\|_{\infty} \leq \varsigma, \text{ e.g., } 10^{-2}. \quad (25)$$

472 *Theorem 1:* The state estimation error by the Huber M-
473 estimator above has an asymptotic normal probability distribution
474 with zero mean and covariance matrix \mathbf{V}_k given by

$$\mathbf{V}_k = \frac{\mathbb{E}[\psi^2(r_S)]}{(\mathbb{E}[\psi'(r_S)])^2} (\mathbf{H}_k^T \mathbf{H}_k)^{-1}. \quad (26)$$

475 *Proof:* Let us define $T_l = T(F_l)$ as the statistic estimates
476 of $T(F)$ and consider the ϵ -contaminated distribution $F_{\epsilon} = (1 - \epsilon)F + \epsilon\delta_{\hat{x}}$, where F is the true distribution while $\delta_{\hat{x}}$ is
477 the point mass at \hat{x} with an unknown distribution for the outliers
478 or thick-tailed distributions. By virtue of the Glivenko-Cantelli
479 theorem, $T(F_{\epsilon}) \rightarrow T(F)$ as $\epsilon \rightarrow 0$. Taking Taylor series ex-
480 pansion on $\psi(r_{S_i}; \hat{x})$ of (21) about \mathbf{x} , we obtain
481

$$\begin{aligned} & \sqrt{l} \left\{ \frac{1}{l} \sum_{i=1}^l \boldsymbol{\alpha}_i \cdot \psi(r_{S_i}; \mathbf{x}) \right\} \\ & + \mathbf{H}_k^T \sqrt{l} (\hat{\mathbf{x}} - \mathbf{x}) \left\{ \frac{1}{l} \sum_{i=1}^l \psi'(r_{S_i}; \mathbf{x}) \right\} \\ & + \boldsymbol{\alpha}_i \cdot O_p(1/\sqrt{l}) = \mathbf{0}, \end{aligned} \quad (27)$$

482 where $O_p(\cdot)$ represents the higher order error terms.

Using the central limit theorem, we have

$$\sqrt{l} \left\{ \frac{1}{l} \sum_{i=1}^l \psi(r_{S_i}; \mathbf{x}) \right\} \xrightarrow{d} \mathbf{Z} \sim \mathcal{N} \left(0, E_F \left[\psi(r_{S_i}; \mathbf{x})^2 \right] \right). \quad (28)$$

By the weak law of large numbers, we obtain

$$\left\{ \frac{1}{l} \sum_{i=1}^l \psi'(r_{S_i}; \mathbf{x}) \right\} \xrightarrow{p} E_F \left[\psi'(r_{S_i}; \mathbf{x}) \right]. \quad (29)$$

Thus, by Slutsky's theorem, we get

$$\mathbf{H}_k^T \sqrt{l} (\hat{\mathbf{x}} - \mathbf{x}) \xrightarrow{d} \frac{-\mathbf{Z}}{\mathbb{E}_F [\psi'(r_{S_i}; \mathbf{x})]} \sim \mathcal{N} (0, \eta^2) \quad (30)$$

486 where $\eta^2 = \frac{\mathbb{E}_F [\psi(r_{S_i}; \mathbf{x})^2]}{\mathbb{E}_F [\psi'(r_{S_i}; \mathbf{x})]^2}$. Thus, we can obtain

$$\boldsymbol{\mu}_k = \lim_{l \rightarrow \infty} \mathbb{E} \left[\sqrt{l} (\hat{\mathbf{x}} - \mathbf{x}) \right] = \mathbf{0}, \quad (31)$$

$$\mathbf{V}_k = \lim_{l \rightarrow \infty} \text{Var} \left[\sqrt{l} (\hat{\mathbf{x}} - \mathbf{x}) \right] = \frac{\mathbb{E} [\psi^2(r_S)]}{(\mathbb{E} [\psi'(r_S)])^2} (\mathbf{H}_k^T \mathbf{H}_k)^{-1}, \quad (32)$$

which complete the proof. ■ 487

488 After a state estimation is carried out, the estimated/
489 interpolated SCADA measurements can be calculated through
490 $\hat{\mathbf{z}} = \mathbf{h}(\hat{\mathbf{x}})$, where the subscript k is omitted for simplicity. Define
491 the difference between the interpolated and the received
492 SCADA measurements as a new residual, i.e., $\mathbf{v} = \mathbf{z} - \hat{\mathbf{z}}$, we
493 have the following theorem:

494 *Theorem 2:* With the assumption that the measurement er-
495 rors of the received SCADA measurements follow a Gaussian
496 distribution, i.e., $\mathbf{e} \sim \mathcal{N}(\mathbf{0}, \mathbf{R})$ defined in (1), the new residual
497 is normally distributed with zero mean and the covariance ma-
498 trix $\mathbf{C} = \mathbf{R} + \mathbf{H} \mathbf{V} \mathbf{H}^T$, where \mathbf{H} is the Jacobian matrix of the
499 vector-valued function $\mathbf{h}(\cdot)$ evaluated at $\hat{\mathbf{x}}$. 499

500 *Proof:* By taking the first order Taylor series expansion of
501 $\mathbf{h}(\mathbf{x})$ at $\hat{\mathbf{x}}$, the innovation vector \mathbf{v} can be expressed as

$$\begin{aligned} \mathbf{v} &= \mathbf{z} - \hat{\mathbf{z}} = \mathbf{h}(\mathbf{x}) + \mathbf{e} - \mathbf{h}(\hat{\mathbf{x}}) \\ &= \mathbf{h}(\hat{\mathbf{x}}) + \mathbf{H}(\mathbf{x} - \hat{\mathbf{x}}) + \mathbf{e} - \mathbf{h}(\hat{\mathbf{x}}) \\ &= \mathbf{H}(\mathbf{x} - \hat{\mathbf{x}}) + \mathbf{e}. \end{aligned} \quad (33)$$

502 Thus, $\mathbb{E}[\mathbf{v}] = \mathbf{H}\mathbb{E}[\mathbf{x} - \hat{\mathbf{x}}] + \mathbb{E}[\mathbf{e}] = \mathbf{0}$ and the covariance ma-
503 trix is $\mathbb{E}[\mathbf{v}\mathbf{v}^T] = \mathbf{H}\text{cov}(\mathbf{x} - \hat{\mathbf{x}})\mathbf{H}^T + \mathbf{R} = \mathbf{H}\mathbf{V}\mathbf{H}^T + \mathbf{R}$. 504

505 Note that the asymptotic zero mean of the innovation vector
506 \mathbf{v} is valid under the condition that none of the received SCADA
507 measurements is attacked. Otherwise, $\mathbb{E}[\mathbf{v}] = \mathbf{a}$, where \mathbf{a} is the
508 measurement bias injected by the hacker. Therefore, by checking
509 the zero mean hypothesis of \mathbf{v} , we can determine whether an
510 FDIA has been conducted or not. To this end, a binary hypothesis
511 test on the measurement consistency can be developed; it is as
512 follows:

$$\begin{cases} \mathcal{H}_0 : \mathbf{v} \sim \mathcal{N}(\mathbf{0}, \mathbf{C}) \\ \mathcal{H}_1 : \mathbf{v} \sim \mathcal{N}(\mathbf{a}, \mathbf{C}) \end{cases}, \quad (34)$$

513 where hypothesis \mathcal{H}_0 and \mathcal{H}_1 represent no FDIA and the occurrence
 514 of FDIA, respectively. By using the log-likelihood ratio
 515 test, we have

$$\xi = \mathbf{v}^T \mathbf{C}^{-1} \mathbf{a} - \frac{1}{2} \mathbf{a}^T \mathbf{C}^{-1} \mathbf{a} \begin{cases} \stackrel{\mathcal{H}_1}{>} \\ \stackrel{\mathcal{H}_0}{<} \end{cases} \kappa, \quad (35)$$

516 where κ is a decision threshold. Since ξ is a linear combination
 517 of \mathbf{v} , according to (34), it is expressed as

$$\begin{cases} \mathcal{H}_0 : \xi \sim \mathcal{N}(\mu_0, \sigma_\xi^2) \\ \mathcal{H}_1 : \xi \sim \mathcal{N}(\mu_1, \sigma_\xi^2); \end{cases} \quad (36)$$

518 and where

$$\begin{aligned} \mu_0 &= -\frac{1}{2} \mathbf{a}^T \mathbf{C}^{-1} \mathbf{a} \\ \mu_1 &= \frac{1}{2} \mathbf{a}^T \mathbf{C}^{-1} \mathbf{a} \\ \sigma_\xi^2 &= \mathbf{a}^T \mathbf{C}^{-1} \mathbf{a}. \end{aligned} \quad (37)$$

519 Therefore, given a false alarm rate P_{fa} , the relationship between
 520 the detection threshold κ and P_{fa} of the Neyman-Pearson
 521 detector is given by

$$P_{fa} = P(\xi \geq \kappa | \mathcal{H}_0) = \frac{1}{2} \operatorname{erfc}\left(\frac{\kappa - \mu_0}{\sqrt{2}\sigma_\xi}\right), \quad (38)$$

522 where $\operatorname{erfc}(\cdot)$ represents the complementary error function of
 523 the normal distribution. Thus, the threshold of the detector can
 524 be calculated as

$$\kappa = \sqrt{2}\sigma_\xi \operatorname{erfc}^{-1}(2P_{fa}) + \mu_0. \quad (39)$$

525 Finally, the detection probability of an FDIA using the hypothesis
 526 testing is given by

$$P_d = P(\xi \geq \kappa | \mathcal{H}_1) = \frac{1}{2} \operatorname{erfc}\left(\frac{\kappa - \mu_1}{\sqrt{2}\sigma_\xi}\right). \quad (40)$$

527 *Remark:* the aim of above analysis is to investigate the theoretical
 528 tradeoff between detection probability and false alarm of the proposed detector given the specified attack vector \mathbf{a} .
 529 In practice, the operator does not need to know the attack vector \mathbf{a} . Instead, the following equivalent normalized innovation
 530 vector-based statistical test is used:

$$v_{Ni} = \frac{|v_i|}{\sqrt{C(i,i)}} \leq \Theta, \quad (41)$$

533 where v_i and v_{Ni} are the i th element of the innovation vector \mathbf{v}
 534 and its normalized value, respectively; $C(i,i)$ is the i th diagonal
 535 element of the derived covariance matrix; Θ is the detection
 536 threshold that is determined by the given confidence level of
 537 the Gaussian distribution, e.g., 3 for 99.7% confidence level.
 538 Note that we have proved that the innovation vector follows a
 539 Gaussian distribution with zero mean and the covariance matrix
 540 \mathbf{C} . The occurrence of the FDIA violates this fact and will be
 541 detected by our proposed robust detector with 99.7% confidence
 542 level.

V. NUMERICAL RESULTS

543

In this section, we use Sections V-A and V-B to demonstrate the validity of the proposed analytical FDIA framework against nonlinear state estimator, where the trade-off between attack magnitudes and information uncertainties is analyzed as well; by contrast, Sections V-C and V-D demonstrate the effectiveness of our robust FDIA detector in presence of measurement noise and bad data.

Specifically, extensive numerical simulations are carried out on the IEEE 30-bus and the 118-bus test systems. The measurement configurations of two test systems are as follows: 1) the IEEE 30-bus system is measured by 93 SCADA measurements, including 18 pairs of active and reactive power injections, 28 pairs of power flows and voltage magnitude of Bus 1; 2) the 118-bus system has 150 pairs of SCADA measurements, including 39 pairs of injection measurements and 111 pairs of flow measurements. The detailed measurement placements and topology of both test systems can be found in [34]. The meter errors of SCADA and PMU measurements follow the normal distribution with zero mean and standard deviations of 10^{-2} and 10^{-3} , respectively. The detection threshold for the normalized residual test is set to 3 with 99.7% confidence level. Two types of AC FDIA are considered: i) *Perfect Attack*: the adversary does not have estimation errors corrupting the state variables in the attack graph $\mathcal{S} = \{\mathcal{N}, \Omega\}$; ii) *Imperfect Attack*: the adversary has estimation errors corrupting the state variables in the attack graph \mathcal{S} . One hundred Monte Carlo simulations are carried out to estimate the average value of the state estimation errors. The effectiveness of the proposed method will be first validated on the IEEE 30-bus system in Sections V-A–V-C, and then its scalability and robustness for larger-scale system will be tested using the 118-bus system.

A. Validation of the Imperfect State Variable Attack

575

A hacker is assumed to change the electricity consumption at Bus 26 through cyber attacks. To this end, she only needs to compromise the measurements P_{25-27} , Q_{25-27} , P_{25-24} , Q_{25-27} , P_{25} and Q_{25} so as to estimate V_{25} and θ_{25} . After that, the following two types of attacks are conducted without being detected by the operators of the control center.

Case 1: Perfect attack where the phase angle at Bus 26 is changed from $\theta_{26} = -0.2990$ to -0.0990 radians.

Case 2: Imperfect attack where the phase angle at Bus 26 is changed from $\theta_{26} = -0.2990$ to -0.0990 radians; here, the estimated state variables in the attack graph \mathcal{S} have 5% errors, which are simulated using 100 Monte Carlo simulations. To be specific, a random sample taken from the uniform distribution $[-\Delta\tau_j + \hat{x}_j, \Delta\tau_j + \hat{x}_j]$, $j \in \Omega_i$ is used to represent the j th state variable to be estimated by the hacker; the final results are obtained by taking the average value of the 100 estimation errors.

Fig. 1 displays the results for Case 1 and Case 2. We can observe from this figure that the perfect attack has successfully changed the nonlinear SE results to the target value within 4 iterations. By contrast, the imperfect attack is not able to change the compromised state variable to the exact target value due to the uncertain knowledge of the system by the hacker. However, the difference between the hacker's target value and the

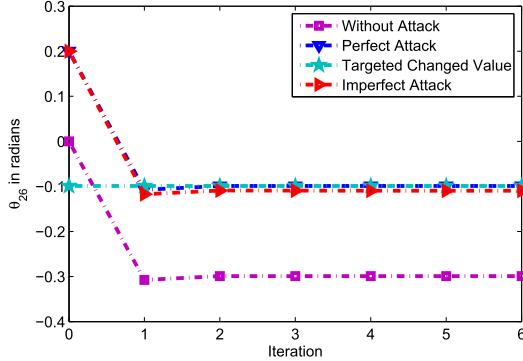


Fig. 1. Performance of perfect and imperfect FDIA over iterations of the nonlinear state estimator using weighted least squares algorithm.

599 achieved value is quite small. Therefore, it can be concluded
600 that although the system information obtained by a hacker is
601 inaccurate or contains uncertainties, she is capable of performing
602 an imperfect attack with consequences that are close to her
603 expectations without being detected by the traditional residual
604 bad data detection tests.

605 B. Tradeoff Between Attack Magnitude and Uncertainties

606 It has been verified in Section V-A that a hacker is able
607 to launch an imperfect attack to change some state variables
608 close to her desire subject to system uncertainties. Now, what
609 is the relationship between system uncertainties and the attack
610 magnitude? To answer that question, let us first note that the
611 system uncertainties will lead to estimation errors of the state
612 variables in graph \mathcal{S} , which further affects the construction of
613 the attack vector. To show this, we first resort to our theoretical
614 results given by (14). It is clear that there is a tradeoff between the
615 attack magnitude and the estimation errors of the state variables
616 in graph \mathcal{S} . To further demonstrate this tradeoff, some simulation
617 results are conducted and analyzed. To this end, we implement
618 our imperfect FDIA to attack a single state variable and multiple
619 state variables with varying errors in graph \mathcal{S} . The former case is
620 similar to Case 2 while the latter one is similar to Case 3, which
621 are imperfect attacks aimed at changing the phase angles at
622 Buses 26 and 24 to -0.0990 and -0.1886 radians, respectively.
623 The traditional normalized residual statistical test is used with
624 a detection threshold of 3.

625 The test results are displayed in Fig. 2. It can be seen from
626 Fig. 2 that with increased estimation errors of the state variables
627 in the attack graph, the largest normalized residual continues to
628 increase and finally exceeds the detection threshold. The largest
629 estimation errors of the state variables in the attack graph when
630 imperfect attacks are implemented successfully are 18.4% and
631 16.5% for single and multiple state attacks, respectively. This
632 means that if a hacker is unable to estimate the state variables
633 in the attack graph within a certain error tolerance, the attack
634 will be detected. To show how this error tolerance threshold
635 is affected by the attack magnitude, we consider changing θ_{26}
636 and θ_{24} to -0.05 and -0.0886 radians, respectively. It is found
637 that the error tolerance threshold decreases to 9.25% due to the
638 increase of attack magnitude. Thus, the tradeoff between attack
639 magnitude and system uncertainties is validated.

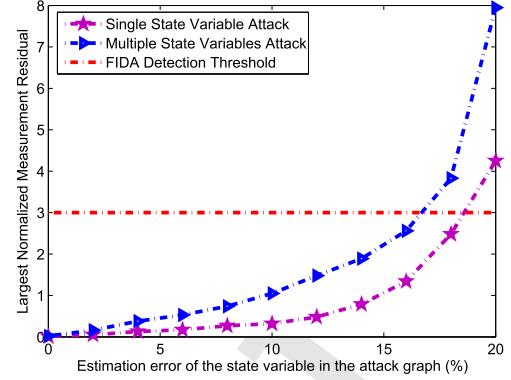


Fig. 2. Largest normalized residue vs. estimation error of the state variables in the attack graph when implementing imperfect attacks.

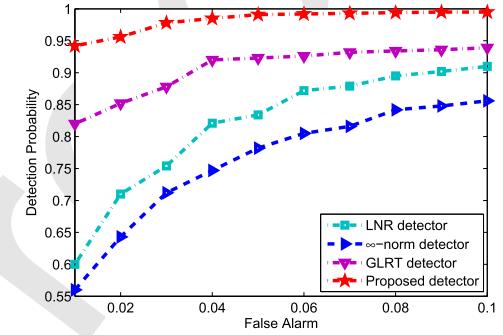


Fig. 3. Performance comparisons of different detectors for imperfect FDIA in Case 2.

640 C. Detection of an FDIA on the Nonlinear State Estimator

641 As revealed in Fig. 2, if the estimation errors of the state vari-
642 ables in the attack graph \mathcal{S} are less than 15% for Cases 2 and
643 3, the largest normalized residual is unable to reveal an FDIA.
644 By contrast, our proposed detector can detect them with a high
645 probability. To show this, we implement the proposed detector
646 for both cases with a limited number of secure PMU measure-
647 ments, say 10 PMUs, which are assumed to be installed at Buses
648 1, 7, 8, 10, 11, 12, 18, 23, 26 and 30. This is the minimal number
649 of PMUs to observe the 30-bus system. All the measurements
650 from these PMUs are assumed to be protected by the control
651 center. The estimation errors of the state variables in the at-
652 tack graph \mathcal{S} vary randomly between 10% and 20%. 100 Monte
653 Carlo simulations are carried out to assess the performance of
654 each detector.

655 Figs. 3 and 4 show the receiver operating characteristic
656 (ROC) curves for different detectors, including the generalized
657 likelihood ratio test (GLRT)-based detector, infinite norm-based
658 detector, and the largest normalized residual (LNR)-based
659 detector [2].

660 It is observed from these two figures that the proposed detector
661 outperforms the other three detectors. In particular, the infinite-
662 norm, the LNR, and the GLRT detectors have more difficulties in
663 detecting the single state variable attack than the multiple state
664 variable attack. By contrast, our proposed detector is slightly
665 affected. These results actually validate the fact that the more
666 uncertainty the information a hacker obtains, the easier she will
667 be detected by the operators of the control center. Thus, from

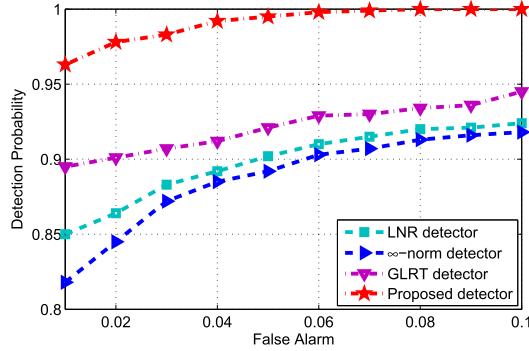


Fig. 4. Performance comparisons of different detectors for imperfect FDIA in Case 3.

668 the hacker's point of view, attacking a bus with less adjacent
669 buses needs less information than other buses and therefore,
670 will increase the probability of successful attack.

671 Note that a perfect FDIA is less likely to occur in practice,
672 which represents certainly the worst case for the detectors at the
673 control center. To evaluate the performance of our detector in
674 this situation, the perfect attack of Case 1 is used. It is found
675 that the infinite-norm, the LNR, and the GLRT detectors are
676 unable to detect this attack. By contrast, the proposed detector
677 is capable of detecting it with a similar performance as that
678 of the imperfect attack of Case 2. Upon a closer look, this is
679 not a surprising result. The reason is that by using the secure
680 PMU measurements as well as the predicted metered values, a
681 relative reliable set of system state information can be obtained;
682 then the robust statistical test of measurement consistency is
683 able to detect the compromised SCADA measurements with a
684 high probability. Note that the secure PMU measurements and
685 the predicted measurements are not from the same source or
686 devices as the SCADA measurements, and thus are not affected
687 by the attacks. In addition, as long as the attack magnitude on
688 SCADA measurements exceed 3 times the standard deviation of
689 the measurement error, FDIA will be detected by our detector
690 with a probability of 95%.

691 D. Scalability and Robustness of the Proposed Method

692 To evaluate the scalability and the robustness of the proposed
693 method, numerical tests are performed on the IEEE 118-bus
694 system. It is assumed that 29 secure PMUs are deployed to
695 make the system observable [9]; the adversary aims to attack the
696 state variable θ_5 with 10 times the standard deviation error; the
697 estimation errors of the state variables in the attack graph \mathcal{S} are
698 randomly varied between 5% and 15%. Note that only the single
699 state attack scenario is tested. This is because from the hacker's
700 point of view, she must try to launch a successful imperfect FDIA
701 with as a small uncertainty as possible. In other words, she will
702 be able to attack a small power system area with high confidence.
703 The larger area she wants to attack, the more uncertainties about
704 the system information are, yielding higher probability of being
705 detected by the operators at the control center. Thus, from the
706 operator's point of view, if the attacks on a small area with
707 small uncertainties can be effectively detected, there is no need
708 to worry about the risk of attacks on large areas with very high
709 uncertainties. Fig. 5 displays the detection probability versus the

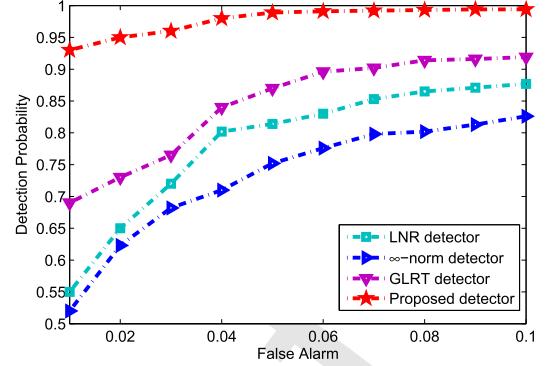


Fig. 5. Performance comparisons of different detectors for imperfect FDIA in IEEE 118-bus system.

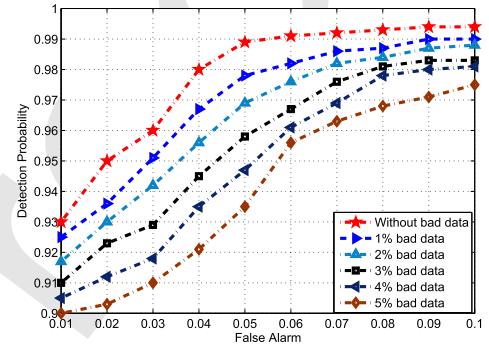


Fig. 6. Performance of the proposed detector for imperfect FDIA with various percentage of bad PMU measurements in IEEE 118-bus system.

710 false alarm probability for all the detectors. It can be observed
711 that the performance of the proposed detector is superior to the
712 other three methods. In addition, compared with the results on
713 the 30-bus test system, our detector is slightly affected by the
714 increased size of the test system.

715 To further investigate the impact of bad received and fore-
716 casted PMU measurements on the proposed detector, 1% to 5%
717 of them are contaminated by adding errors with 10 to 15 times
718 their standard deviations. The test results are shown in Fig. 6.
719 Thanks to the robustness of the Huber M-estimator and the en-
720 hanced measurement redundancy with forecasted metered val-
721 ues, the influence of gross errors are bounded, yielding a slightly
722 decreased performance of detecting attacks. However, the least
723 detection probability of the proposed method is still greater than
724 90%. Note that for the bus with several adjacent buses, it has
725 high local redundancy and the proposed detector can suppress
726 several bad PMU measurements, while for those who have only
727 one adjacent bus, it can handle fewer bad PMU measurements
728 associated to that bus.

VI. CONCLUSION

730 The first contribution of this paper is to extend the existing
731 perfect FDIA model by developing a generalized FDIA frame-
732 work against nonlinear SE that accounts for the uncertainties
733 in the measurements or in system topology. The upper bounds
734 of these uncertainties for performing successful FDIA are in-
735 vestigated analytically. They provide the operators with a better
736 understanding of the vulnerability of a nonlinear SE to FDIA.

737 and thus may facilitate the adoption and implementation of ef-
 738 fective defense methods. The second contribution of this paper
 739 is the development of a robust FDIA detection method that
 740 checks the measurement statistical consistency using a limited
 741 number of secure PMU measurements. Numerical results are
 742 provided to demonstrate the effectiveness and robustness of the
 743 proposed method. Future work will concentrate on the detection
 744 of topology attacks caused by the change of parameter values
 745 in the system. In addition, we will evaluate the sensitivity of
 746 the proposed approach to the accuracy of the forecasted PMU
 747 measurements and the change of system operation conditions
 748 in a short timeframe. Corresponding mitigation methods will be
 749 proposed if needed.

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Junbo Zhao (S'13) received the Bachelor's degree in electrical engineering
 from Southwest Jiaotong University, Chengdu, China, in 2012. He is currently
 working toward the Ph.D. degree in Bradley Department of Electrical and Com-
 puter Engineering, Virginia Polytechnic Institute and State University, Falls
 Church, VA, USA. He did a summer internship at Pacific Northwest National
 Laboratory from May to August 2017. He is now the Chair of the IEEE Task
 Force on Power System Dynamic State and Parameter Estimation, the secretary
 of the IEEE Working Group on State Estimation Algorithms, and the IEEE Task
 Force on Synchrophasor Applications in Power System Operation and Control.
 His research interests include power system real-time monitoring, operations,
 and security.

Lamine Mili (LF'17) received the Electrical Engineering Diploma from the
 Swiss Federal Institute of Technology, Lausanne, Switzerland, in 1976, and the
 Ph.D. degree from the University of Liège, Liège, Belgium, in 1987. He is cur-
 rently a Professor with the Department Electrical and Computer Engineering,
 Virginia Polytechnic Institute and State University, Falls Church, VA, USA. He
 has five years of industrial experience with the Tunisian electric utility, STEG.
 He was a Visiting Professor with the Swiss Federal Institute of Technology,
 Lausanne, the Grenoble Institute of Technology, the École Supérieure
 D'électricité, France, and the École Polytechnique de Tunisie, Tunisia, and
 did consulting work for the French Power Transmission company, RTE. His
 research has focused on power system planning for enhanced resiliency and
 sustainability, risk management of complex systems to catastrophic failures,
 robust estimation and control, nonlinear dynamics, and bifurcation theory.

Meng Wang (M'12) received the B.S. and M.S. degrees from Tsinghua Uni-
 versity, Beijing, China, and the Ph.D. degree from Cornell University, Ithaca,
 NY, USA, in 2012. She is currently an Assistant Professor with the Depart-
 ment of Electrical, Computer, and Systems Engineering, Rensselaer Polytechnic
 Institute, Troy, NY, USA. Her research interests include high-dimensional
 data analysis and their applications in power systems monitoring and network
 inference.