GELFAND-TSETLIN POLYTOPES: A STORY OF FLOW & ORDER POLYTOPES

RICKY INI LIU, KAROLA MÉSZÁROS, AND AVERY ST. DIZIER

ABSTRACT. Gelfand-Tsetlin polytopes are prominent objects in algebraic combinatorics. The number of integer points of the Gelfand-Tsetlin polytope $GT(\lambda)$ is equal to the dimension of the corresponding irreducible representation of GL(n). It is well-known that the Gelfand-Tsetlin polytope is both a marked order polytope and a flow polytope. In this paper, we draw corollaries from this result and establish a general theory connecting marked order polytopes and flow polytopes.

1. Introduction

Given a partition $\lambda = (\lambda_1, \dots, \lambda_n) \in \mathbb{Z}_{>0}^n$, the **Gelfand-Tsetlin polytope** $GT(\lambda)$ is the set of all nonnegative triangular arrays

such that

$$x_{1i} = \lambda_i \text{ for all } 1 \le i \le n,$$

$$x_{i-1,j-1} \ge x_{ij} \ge x_{i-1,j} \text{ for all } 2 \le i \le j \le n.$$

The integer points of $GT(\lambda)$ are in bijection with semistandard Young tableaux of shape λ on the alphabet [n]. Moreover, the integer point transform of $GT(\lambda)$ projects to the Schur function s_{λ} . The latter beautiful result generalizes to Minkowski sums of Gelfand-Tsetlin polytopes and certain Schubert polynomials as well [5]. In this paper we will be interested in the Gelfand-Tsetlin polytope $GT(\lambda)$ from a purely discrete geometric point of view: we will explore it as a marked order polytope and as a flow polytope.

Ardila et al. introduced marked order polytopes and showed that Gelfand–Tsetlin polytopes are examples of them in [1]. That Gelfand-Tsetlin polytopes are also flow polytopes was shown by Danilov, Karzanov, and Koshevoy in [3]. Theorem 1.1 summarizes previous work by Postnikov [9, Theorem 15.1] on the Gelfand-Tsetlin polytope and also demonstrates how the Gelfand-Tsetlin polytopes being flow polytopes allows us to write the number of integer points and the volume of $GT(\lambda)$ —which are equal respectively to the dimension of the irreducible representation V_{λ} of GL(n) and the top homogeneous component of the dimension when viewed as a polynomial in $\lambda_1, \ldots, \lambda_n$ —in terms of Kostant partition functions:

Theorem 1.1. Let $\lambda \in \mathbb{Z}_{\geq 0}^n$ be a partition and $r = |V(G_{\lambda})| - 1 = \binom{n+2}{2} - 3$. The volume of $GT(\lambda)$ is given by

(1) Vol
$$GT(\lambda) = \prod_{1 \le i < j \le n} \frac{\lambda_i - \lambda_j}{j - i}$$

(2)
$$= \sum_{b_1, \dots, b_{n-1} > 0} \frac{(\lambda_1 - \lambda_2)^{b_1}}{b_1!} \cdots \frac{(\lambda_{n-1} - \lambda_n)^{b_{n-1}}}{b_{n-1}!} N(b_1, \dots, b_{n-1})$$

(2)
$$= \sum_{b_1,\dots,b_{n-1}\geq 0} \frac{(\lambda_1 - \lambda_2)^{b_1}}{b_1!} \cdots \frac{(\lambda_{n-1} - \lambda_n)^{b_{n-1}}}{b_{n-1}!} N(b_1,\dots,b_{n-1})$$

$$= \sum_{j} \frac{(\lambda_1 - \lambda_2)^{j_1}}{j_1!} \cdots \frac{(\lambda_{n-1} - \lambda_n)^{j_{n-1}}}{j_{n-1}!} K_{G_{\lambda}} (j_1 - 1,\dots,j_{n-1} - 1,-1,\dots,-1,0,\dots,0,0) .$$

Liu is partially supported by a National Science Foundation Grant (DMS 1700302). Mészáros is partially supported by a National Science Foundation Grant (DMS 1501059), a von Neumann Fellowship at the IAS funded by the Fund for Mathematics and the Friends of the Institute for Advanced Study, as well as a CAREER National Science Foundation Grant (DMS 1847284).

The integer point count of $GT(\lambda)$ is given by

(4)
$$|GT(\lambda) \cap \mathbb{Z}^{\binom{n+1}{2}}| = \prod_{1 \le i < j \le n} \frac{\lambda_i - \lambda_j + j - i}{j - i}$$
(5)
$$= \sum_{j} \binom{\lambda_1 - \lambda_2 + 1}{j_1} \cdots \binom{\lambda_{n-1} - \lambda_n + 1}{j_{n-1}} \binom{1}{j_n} \cdots \binom{1}{j_{\binom{n}{2}}} \binom{0}{j_{\binom{n}{2}+1}} \cdots \binom{0}{j_r}$$

$$\cdot K_{G_{\lambda}} \left(j_1 - 1, \dots, j_{n-1} - 1, j_n - 1, \dots, j_{\binom{n}{2}} - 1, j_{\binom{n}{2}+1}, \dots, j_r, 0 \right).$$

Equalities (3) and (5) follow from $GT(\lambda)$ being a flow polytope. The other equations are known and follow from the representation theory of GL(n) and from Postnikov's work [9, Theorem 15.1]. For the notation used in Theorem 1.1, we refer the reader to Section 2. We remark that from equations (2) and (3), we obtain that the evaluations $N(b_1, \ldots, b_{n-1})$ and $K_{G_{\lambda}}(j_1 - 1, \ldots, j_{n-1} - 1, -1, \ldots, -1, 0, \ldots, 0, 0)$ are equal. We additionally provide a bijective proof of this in Section 2.

Section 3 is devoted to marked order polytopes in general. In light of the work of the second author with Morales and Striker [7] where they show that order polytopes of strongly planar posets are flow polytopes, it is natural to wonder if the Gelfand–Tsetlin polytopes being both a marked order polytope and a flow polytope is part of a larger picture. Indeed, we show that marked order polytopes of strongly planar posets with certain conditions on the markings are flow polytopes:

Theorem 1.2. If (P, A, λ) is a marked poset admitting a bounded strongly planar embedding, then the marked order polytope $\mathcal{O}(P, A)_{\lambda}$ is integrally equivalent to the flow polytope $\mathcal{F}_{G_{(P,A,\lambda)}}$.

For the terminology used in Theorem 1.2, see Sections 3 and 4. There is a natural way of subdividing the marked order polytope $\mathcal{O}(P,A)_{\lambda}$ into products of simplices labeled by certain linear extensions of the poset (Theorem 3.4), and there is a natural way of subdividing the flow polytope $\mathcal{F}_{G(P,A,\lambda)}$ into products of simplices labeled by integer points of other flow polytopes (Section 5.1). In Section 5 we show that these subdivisions map to each other under the integral equivalence of Theorem 1.2, and we conclude by bijecting their combinatorial labelings in Corollary 5.6.

Roadmap of the paper. In Section 2 we define flow polytopes and show several consequences of Gelfand—Tsetlin polytopes being integrally equivalent to flow polytopes for their volume and Ehrhart polynomial formulas. It is well-known that the Gelfand—Tsetlin polytope is also a marked order polytope, and in Section 3 we define marked order polytopes as well as collect and extend some known results about them. Section 4 proves Theorem 1.2, which gives conditions under which marked order polytopes are integrally equivalent to flow polytopes. The Gelfand—Tsetlin polytopes appear as a special case in this more general theory. Finally, in Section 5 we review the subdivision methods for flow polytopes and (marked) order polytopes, and we show that they map to each other under the integral equivalence of Theorem 1.2. We conclude by bijecting the two sets of combinatorial labels coming from the subdivisions of flow and marked order polytopes in Corollary 5.6.

2. Gelfand-Tsetlin polytopes as flow polytopes

In this section, we recall the result from [3] that every Gelfand–Tsetlin polytope is integrally equivalent to a flow polytope. We then study the volume and Ehrhart polynomial of Gelfand–Tsetlin polytopes. We start by defining flow polytopes and providing the necessary background on them.

2.1. Background on flow polytopes. Let G be a directed acyclic connected (multi-)graph on the vertex set [n+1] with m edges. An integer vector $a = (a_1, \ldots, a_n, -\sum_{i=1}^n a_i) \in \mathbb{Z}^{n+1}$ is called a **netflow vector**. A pair (G, a) will be referred to as a **flow network**. To minimize notational complexity, we will typically omit the netflow a when referring to a flow network G, describing it only when defining G. When not explicitly stated, we will always assume vertices of G are labeled so that $(i, j) \in E(G)$ implies i < j.

To each edge (i,j) of G, associate the type A positive root $e_i - e_j \in \mathbb{R}^n$. Let M_G be the incidence matrix of G, the matrix whose columns are the multiset of vectors $e_i - e_j$ for $(i,j) \in E(G)$. A flow on a flow network G with netflow a is a vector $f = (f(e))_{e \in E(G)}$ in $\mathbb{R}^{E(G)}_{\geq 0}$ such that $M_G f = a$. Equivalently, for all

 $1 \leq i \leq n$, we have

$$\sum_{e=(k,i)\in E(G)} f(e) + a_i = \sum_{e=(i,k)\in E(G)} f(e).$$

The fact that the netflow of vertex n+1 is $-\sum_{i=1}^{n} a_i$ is implied by these equations. Define the **flow polytope** $\mathcal{F}_G(a)$ of a graph G with netflow a to be the set of all flows on G:

$$\mathcal{F}_G = \mathcal{F}_G(a) = \{ f \in \mathbb{R}_{>0}^{E(G)} \mid M_G f = a \}.$$

Given a graph G, the **Kostant partition function** K_G of G evaluated at a vector $b \in \mathbb{Z}^{n+1}$ is the number of ways to write b as a nonnegative integer combination of the multiset of vectors $\{e_i - e_j \mid (i, j) \in E(G)\}$, or equivalently

 $K_G(b) = \left| \mathcal{F}_G(b) \cap \mathbb{Z}^{E(G)} \right|.$

Remark 2.1. When G is a flow network (G, a), we will write \mathcal{F}_G for $\mathcal{F}_G(a)$. For any $b \in \mathbb{Z}^{n+1}$, we will write $\mathcal{F}_G(b)$ and $K_G(b)$ when we wish to use a vector possibly different from the netflow a associated to G.

The following remarkable theorem gives the volume and Ehrhart polynomial formulas for a family of flow polytopes.

Theorem 2.2 (Baldoni-Vergne-Lidskii formulas [2, Thm. 38]). Let G be a connected graph on the vertex set [n+1] with m edges, and let $a=(a_1,\ldots,a_n,-\sum_{i=1}^n a_i)$ with $a_i\in\mathbb{Z}_{\geq 0}$ for $i\in[n]$. Direct the edges of G by $i\to j$ if i< j, and assume there is at least one outgoing edge at vertex i for each $i=1,\ldots,n$. Then

(6)
$$\operatorname{Vol} \mathcal{F}_{G}(a) = \sum_{j} \frac{a_{1}^{j_{1}}}{j_{1}!} \cdots \frac{a_{n}^{j_{n}}}{j_{n}!} K_{G}(j_{1} - \operatorname{out}_{1}, \dots, j_{n} - \operatorname{out}_{n}, 0),$$

(7)
$$|\mathcal{F}_{G}(a) \cap \mathbb{Z}^{E(G)}| = \sum_{j} \binom{a_{1} + \operatorname{out}_{1}}{j_{1}} \cdots \binom{a_{n} + \operatorname{out}_{n}}{j_{n}} \cdot K_{G}(j_{1} - \operatorname{out}_{1}, \dots, j_{n} - \operatorname{out}_{n}, 0),$$

(8)
$$= \sum_{j} \begin{pmatrix} a_1 - i\mathbf{n}_1 \\ j_1 \end{pmatrix} \cdots \begin{pmatrix} a_n - i\mathbf{n}_n \\ j_n \end{pmatrix} \cdot K_G (j_1 - \mathbf{out}_1, \dots, j_n - \mathbf{out}_n, 0),$$

 $for \ out_i = outd_i - 1 \ and \ in_i = ind_i - 1, \ where \ outd_i \ and \ ind_i \ denote \ the \ outdegree \ and \ indegree \ of \ vertex$ i in G. Each sum is over weak compositions $j=(j_1,j_2,\ldots,j_n)$ of m-n that are greater than or equal to $(\operatorname{out}_1,\ldots,\operatorname{out}_n)$ in dominance order, and $\binom{n}{k}=\binom{n+k-1}{k}$.

2.2. The Gelfand-Tsetlin polytope as a flow polytope. Given any partition λ , we describe below a graph G_{λ} such that the following theorem holds.

Theorem 2.3 ([3, Proof of Theorem 2]). $GT(\lambda)$ is integrally equivalent to $\mathcal{F}_{G_{\lambda}}$.

Recall that two integral polytopes \mathcal{P} in \mathbb{R}^d and \mathcal{Q} in \mathbb{R}^m are **integrally equivalent** if there is an affine transformation $\varphi \colon \mathbb{R}^d \to \mathbb{R}^m$ whose restriction to \mathcal{P} is a bijection $\varphi \colon \mathcal{P} \to \mathcal{Q}$ that preserves the lattice, i.e., φ is a bijection between $\mathbb{Z}^d \cap \operatorname{aff}(\mathcal{P})$ and $\mathbb{Z}^m \cap \operatorname{aff}(\mathcal{Q})$, where $\operatorname{aff}(\cdot)$ denotes affine span. The map φ is called an integral equivalence. Note that integrally equivalent polytopes have the same Ehrhart polynomials and therefore the same volume.

We now define the flow network G_{λ} , describing the graph and its associated netflow (see Remark 2.1). For an illustration of G_{λ} , see Figure 1.

Definition 2.4. For a partition $\lambda \in \mathbb{Z}_{>0}^n$ with $n \geq 2$, let G_{λ} be defined as follows:

If n=1, let G_{λ} be a single vertex v_{22} defined to have flow polytope consisting of one point, 0. Otherwise, let G_{λ} have vertices

$$V(G_{\lambda}) = \{v_{i,j} \mid 2 \le i \le j \le n\} \cup \{v_{i,i-1} \mid 3 \le i \le n+2\} \cup \{v_{i,n+1} \mid 3 \le i \le n+1\}$$

and edges

$$\begin{split} E(G_{\lambda}) &= \{(v_{ij}, v_{i+1,j}) \mid 2 \leq i \leq j \leq n\} \cup \{(v_{i,n+1}, v_{i+1,n+1}) \mid 3 \leq i \leq n+1\} \\ &\quad \cup \{(v_{ij}, v_{i+1,j+1}) \mid 2 \leq i \leq j \leq n\} \cup \{(v_{i,i-1}, v_{i+1,i}) \mid 3 \leq i \leq n+1\}. \end{split}$$

The default netflow vector on G_{λ} is as follows:

- To vertex v_{2j} for $2 \le j \le n$, assign netflow $\lambda_{j-1} \lambda_j$.
- To vertex $v_{n+2,n+1}$, assign netflow $\lambda_n \lambda_1$.
- To all other vertices, assign netflow 0.

Given a flow on G_{λ} , denote the flow value on each edge $(v_{ij}, v_{i+1,j})$ by a_{ij} , and denote the flow value on each edge $(v_{ij}, v_{i+1,j+1})$ by b_{ij} .

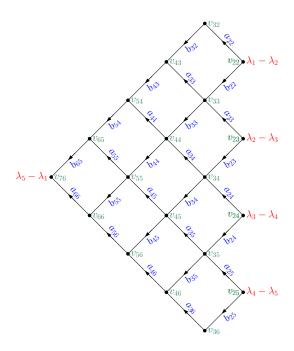


FIGURE 1. The flow network G_{λ} with $\ell(\lambda) = 5$.

We note that viewing $GT(\lambda)$ as a marked order polytope ([1]), Theorem 2.3 is also a special case of our more general Theorem 1.2.

Several expressions for the volume and integer point count of $GT(\lambda)$ are given in the following theorem. To apply the Lidskii formulas (Theorem 2.2) to G_{λ} , we list the vertices of G_{λ} in the following order: First are the vertices $\{v_{ij} \mid 2 \leq i \leq j \leq n\}$ ordered lexicographically; next, the vertices $\{v_{i,i-1} \mid 3 \leq i \leq n+1\}$ ordered lexicographically; then, the vertices $\{v_{i,n+1} \mid 3 \leq i \leq n+1\}$ ordered lexicographically; and lastly, the sink vertex $v_{n+2,n+1}$.

In (2) below, we need a few definitions: A shifted standard Young tableaux (shSYT) is a bijection $T: \{(i,j) \mid 1 \leq i \leq j \leq n\} \rightarrow \{1,2,\ldots,\binom{n+1}{2}\}$ such that T(i,j) < T(i+1,j) and T(i,j) < T(i,j+1). The diagonal vector of a shSYT T is $(T(1,1),T(2,2),\ldots,T(n,n))$. Denote by $N(b_1,\ldots,b_{n-1})$ the number of shSYT T with diagonal entries $T(i,i) = i + b_1 + \cdots + b_{i-1}$.

Theorem 1.1. Let $\lambda \in \mathbb{Z}_{\geq 0}^n$ be a partition and $r = |V(G_{\lambda})| - 1 = \binom{n+2}{2} - 3$. The volume of $GT(\lambda)$ is given by

(1) Vol
$$GT(\lambda) = \prod_{1 \le i < j \le n} \frac{\lambda_i - \lambda_j}{j - i}$$

(2)
$$= \sum_{b_1,\dots,b_{n-1}\geq 0} \frac{(\lambda_1 - \lambda_2)^{b_1}}{b_1!} \cdots \frac{(\lambda_{n-1} - \lambda_n)^{b_{n-1}}}{b_{n-1}!} N(b_1,\dots,b_{n-1})$$

(3)
$$= \sum_{j} \frac{(\lambda_1 - \lambda_2)^{j_1}}{j_1!} \cdots \frac{(\lambda_{n-1} - \lambda_n)^{j_{n-1}}}{j_{n-1}!} K_{G_{\lambda}} (j_1 - 1, \dots, j_{n-1} - 1, -1, \dots, -1, 0, \dots, 0, 0).$$

The integer point count of $GT(\lambda)$ is given by

$$(4) \qquad |\mathrm{GT}(\lambda)\cap\mathbb{Z}^{\binom{n+1}{2}}|=\prod_{1\leq i< j\leq n}\frac{\lambda_{i}-\lambda_{j}+j-i}{j-i}$$

(5)
$$= \sum_{j} {\lambda_{1} - \lambda_{2} + 1 \choose j_{1}} \cdots {\lambda_{n-1} - \lambda_{n} + 1 \choose j_{n}} \cdots {1 \choose j_{\binom{n}{2}}} {\binom{0}{j\binom{n}{2} + 1}} \cdots {\binom{0}{j_{r}}} \cdot K_{G_{\lambda}} \left(j_{1} - 1, \dots, j_{n-1} - 1, j_{n} - 1, \dots, j_{\binom{n}{2}} - 1, j_{\binom{n}{2} + 1}, \dots, j_{r}, 0\right).$$

Proof. In [9], (1), (2), and (4) are shown. Applying Theorem 2.2 to the flow network G_{λ} yields (3) and (5).

Corollary 2.5. Comparing the volume formulas (2) and (3), we obtain that

$$N(b_1,\ldots,b_{n-1})=K_{G_{\lambda}}(b_1-1,\ldots,b_{n-1}-1,-1,\ldots,-1,0,\ldots,0,0)$$

for all $b_1, \ldots, b_{n-1} > 0$.

One can also see Corollary 2.5 bijectively as follows. Given a shSYT T counted by $N(b_1, \ldots, b_{n-1})$, define for $1 \le i < j \le n$:

$$a_{j-i+1,j} = |\{(i',j') \mid T(i,j-1) < T(i',j') < T(i,j), \quad i' < i, \quad j' \ge j\}|,$$

$$b_{j-i+1,j} = |\{(i',j') \mid T(i,j) < T(i',j') < T(i+1,j), \quad i' \le i, \quad j' > j\}|.$$

We claim that these define a flow on G_{λ} with netflow $(b_1-1,\ldots,b_{n-1}-1,-1,\ldots,-1,0,\ldots,0,0)$.

- Note that $a_{ij} = 0$ for i = j, while $b_{ij} = 0$ for j = n. Thus there is no netflow at vertex v_{ij} unless $2 \le i \le j \le n$.
- At v_{2j} , the netflow is $a_{2j} + b_{2j}$, which counts pairs (i', j') such that T(j-1, j-1) < T(i', j') < T(j-1, j) or T(j-1, j) < T(i', j') < T(j, j), of which there are $T(j, j) T(j-1, j-1) 2 = b_{j-1} 1$.
- At any other vertex v_{ij} , the netflow is $a_{ij} + b_{ij} a_{i-1,j} b_{i-1,j-1}$. But $a_{ij} + b_{ij}$ and $a_{i-1,j} + b_{i-1,j-1}$ both count pairs (i',j') such that T(j-i+1,j-1) < T(i',j') < T(j-i+2,j) with the only difference being that (i',j') = (j-i+1,j) is not counted by the first quantity but is counted by the second. It follows that the netflow is -1, as desired.

For the inverse map, we can construct the shSYT T inductively: by removing the vertices v_{2j} and edges with flows a_{2j} and b_{2j} from G_{λ} , we arrive at the graph for a partition of length n-1 with netflow

$$(b_{22} + a_{23} - 1, b_{23} + a_{24} - 1, \dots, b_{2,n-1} + a_{2n} - 1, -1, \dots, -1, 0, \dots, 0).$$

By induction, we can construct from this a shSYT T' with side length n-1 whose ith diagonal entry T'(i,i) is

$$i + (b_{22} + a_{23}) + (b_{23} + a_{24}) + \dots + (b_{2i} + a_{2,i+1}) = i + (b_1 - 1) + (b_2 - 1) + \dots + (b_{i-1} - 1) + a_{2,i+1}$$

= $1 + b_1 + b_2 + \dots + b_{i-1} + a_{2,i+1}$.

Hence

$$1 + b_1 + b_2 + \dots + b_{i-1} \le T'(i, i) \le b_1 + b_2 + \dots + b_i$$
.

Then let us modify T' by adding 1 to the entries $1, \ldots, b_1$, adding 2 to the entries $1 + b_1, \ldots, b_1 + b_2$, and so forth, which in particular will add i to T'(i,i). We can then attach to T' a new diagonal with entries 1, $2 + b_1, 3 + b_1 + b_2, \ldots$, which will yield a shSYT T of side length n with the desired diagonal entries. It is straightforward to check that these two maps are inverses of one another, completing the bijection.

We will provide a bijective proof of a generalization of Corollary 2.5 in Section 5 in the more general setting of strongly planar marked order polytopes.

2.3. Kostka numbers and key polynomials. Since $GT(\lambda)$ is integrally equivalent to $\mathcal{F}_{G_{\lambda}}$, all concepts defined via the Gelfand–Tsetlin polytope may also be defined via the flow polytope. One such example is Kostka numbers $K_{\lambda\alpha}$, which count the number of integer points in $GT(\lambda)$ with weight α . Recall the weight wt of a point $(x_{ij})_{i,j}$ is the vector whose ith component is $\sum_{j=i}^{n} x_{ij} - \sum_{j=i+1}^{n} x_{i+1,j}$. Equivalently, $K_{\lambda\mu}$ counts the number of semistandard Young tableaux with shape λ and content α .

Kostka numbers can be given in terms of the flow polytope of G_{λ} by defining a weight map gwt on flows. Let $f \in \mathbb{R}^{E(G_{\lambda})}$ be a flow on G_{λ} corresponding to the point $P_f \in GT(\lambda)$. Define $gwt : \mathbb{R}^{E(G_{\lambda})} \to \mathbb{R}^n$ by setting

$$gwt(x_{(v_{ij},v_{i+1,j})}) = e_{i-1}$$
 and $gwt(x_{(v_{ij},v_{i+1,j+1})}) = 0$.

It is shown in [5, Proposition 3.9] that $wt(P_f)_i = gwt(f)_i + \lambda_n$ for each i.

Another example is the key polynomials, also known as Demazure characters. In [13], it is shown that every key polynomial appears as the integer point transform of the projection of a union of Kogan faces of $GT(\lambda)$. We define Kogan faces and describe them in the context of $\mathcal{F}_{G_{\lambda}}$.

Recall that every face of $GT(\lambda)$ is obtained by enforcing equality in some subset of the defining inequalities

$$x_{i-1,j-1} \ge x_{ij}$$
 and $x_{ij} \ge x_{i-1,j}$.

A face is called **Kogan** if it is given by equations of the form $x_{ij} = x_{i-1,j}$. A face F of $\mathcal{F}_{G_{\lambda}}$ corresponds to a Kogan face of $GT(\lambda)$ if and only if it is obtained by enforcing equations $b_{ij} = 0$ on $\mathcal{F}_{G_{\lambda}}$ for some pairs (i, j).

3. Marked Order Polytopes

In [1], Ardila, Bliem, and Salazar observed that the Gelfand–Tsetlin polytope is a section of an order polytope. Inspired, they introduced marked posets and marked poset polytopes, generalizing Stanley's notion of order and chain polytopes introduced in [11]. In this section we give background on unmarked and marked order polytopes, and we explain the generalizations of several results from order polytopes to marked order polytopes.

Definition 3.1. A marked poset (P, A, λ) consists of a finite poset P, a subposet $A \subseteq P$ containing all its extremal elements, and an order-preserving map $\lambda \colon A \to \mathbb{R}$. We identify (P, A, λ) with the marked Hasse diagram in which we label the elements $a \in A$ with $\lambda(a)$ in the Hasse diagram of P.

Definition 3.2. The marked order polytope of (P, A, λ) is

$$\mathcal{O}(P,A)_{\lambda} = \{x \in \mathbb{R}^P \mid x_p \le x_q \text{ for } p < q \text{ in } P \text{ and } x_a = \lambda(a) \text{ for } a \in A\}.$$

Let $\widetilde{\mathcal{O}}(P,A)_{\lambda}$ denote $\mathcal{O}(P,A)_{\lambda}$ projected onto the coordinates $P \setminus A$.

Stanley's construction of the **order polytope** $\mathcal{O}(P)$ [11] is a special case of a marked order polytope. Given a finite poset P, add a new smallest and largest element to obtain $\widehat{P} = P \sqcup \{\widehat{0}, \widehat{1}\}$. Let $A = \{\widehat{0}, \widehat{1}\}$ and $\lambda(\widehat{a}) = a$. Then

$$\mathcal{O}(P) = \widetilde{\mathcal{O}}(\widehat{P}, A)_{\lambda}.$$

In general, computing or finding a combinatorial interpretation for the volume of a polytope is a hard problem. Order polytopes are an especially nice class of polytopes whose volume has a combinatorial interpretation.

Theorem 3.3 (Stanley [11]). Given a poset P, we have that

- (i) the vertices of $\mathcal{O}(P)$ are in bijection with characteristic functions of complements of order ideals of P.
- (ii) the normalized volume of $\mathcal{O}(P)$ is e(P), where e(P) is the number of linear extensions of P, and
- (iii) the Ehrhart polynomial $L_{\mathcal{O}(P)}(m)$ of $\mathcal{O}(P)$ equals the order polynomial $\Omega(P, m+1)$ of P.

We now explain how Theorem 3.3 generalizes to the setting of marked order polytopes.

For part (i), the vertices and facial structure of marked order polytopes are described by Pegel in [8]. A point $x \in \mathcal{O}(P,A)_{\lambda}$ induces a partition π_x of P that is the transitive closure of the relation $p_1 \sim_x p_2$ if $x_p = x_q$ and p,q are comparable. A point $x \in \mathcal{O}(P,A)_{\lambda}$ is a vertex if and only if each block of π_x contains a marked point. In the case of an unmarked order polytope O(P), the blocks will be an order ideal and its complement, so the vertices are characteristic functions.

Part (ii) of Theorem 3.3 has a beautiful geometric justification: order polytopes have a canonical subdivision into e(P) unimodular simplices. Consider $\mathcal{O}(P)$ cut with all hyperplanes of the form $x_p = x_q$ where $p, q \in P$ with p and q incomparable. The regions of this arrangement correspond to the ways of totally ordering the coordinates x_p compatible with all inequalities of $\mathcal{O}(P)$, that is, linear extensions of P. Each region is defined by inequalities of the form $y_1 \leq y_2 \leq \cdots \leq y_{|P|}$ for $y_1, \ldots, y_{|P|}$ a permutation of $\{x_p\}_{p \in P}$, so each region is a simplex.

The following theorem generalizes part (ii) of Theorem 3.3 to marked order polytopes. For notational convenience, we will take a **linear extension** of a poset P with n elements to be an *order-reversing* bijection $\sigma: [n] \to P$, so for example $\sigma(1)$ will be a maximal element of P. We will generally label the elements of A as $\{p_1, p_2, \ldots, p_k\}$ such that $\lambda(p_1) \ge \cdots \ge \lambda(p_k)$ (and, additionally, if $p_i > p_j$ in P, then i < j).

Theorem 3.4. (cf. [10, Theorem 3.2]) If (P, A, λ) is a marked poset with marked elements $A = \{p_1, \ldots, p_k\}$ having markings $\lambda(p_1) \geq \cdots \geq \lambda(p_k)$ denoted $\lambda_1, \ldots, \lambda_k$, then

Vol
$$\widetilde{\mathcal{O}}(P,A)_{\lambda} = \sum_{a_1,\dots,a_{k-1}\geq 0} N_{P,A,\lambda}(a_1,\dots,a_{k-1}) \frac{(\lambda_1-\lambda_2)^{a_1}}{a_1!} \cdot \dots \cdot \frac{(\lambda_{k-1}-\lambda_k)^{a_{k-1}}}{a_{k-1}!},$$

where $N_{P,A,\lambda}(a_1,\ldots,a_{k-1})$ is the number of linear extensions of P such that elements of $A=\{p_1,\ldots,p_k\}$ occur at positions $1,2+a_1,\ldots,k+a_1+\cdots+a_{k-1}$, respectively.

We note that when the markings A are along a chain in the poset P, Stanley has shown the above theorem in his proof of a certain log-concavity conjecture which we explain below; see the proof of [10, Theorem 3.2]. His proof can be generalized to the above setting. We provide the proof here for completeness and take a slightly different perspective via hyperplane cuts, much like Postnikov does in [9] for $GT(\lambda)$.

Proof of Theorem 3.4. Consider $\widetilde{\mathcal{O}}(P,A)_{\lambda}$ cut with all hyperplanes of the form $x_p = x_q$ or $x_p = \lambda(a)$, where $p,q \in P \setminus A$ with p is incomparable with q, and $a \in A$ is incomparable with p. The regions of this arrangement correspond to the ways of totally ordering the coordinates $(x_p)_{p \in P}$ compatible with all inequalities and markings, that is, certain linear extensions of P. Let $\sigma \colon [n] \to P$ be a linear extension of P, say with $\sigma(i_j) = p_j$ for $j \in [k]$, $i_1 < i_2 < \dots < i_k$. Since A contains all minimal and maximal elements of P, note that $i_1 = 1$ and $i_k = n$. The associated region $\widetilde{\Pi}_{\sigma}$ in the subdivision is the projection of the region

$$\Pi_{\sigma} = \{ x \in \mathbb{R}^P \mid x_{\sigma(1)} \ge x_{\sigma(2)} \ge \dots \ge x_{\sigma(n)}, \ x_{\sigma(i_i)} = \lambda_i \}.$$

onto the coordinates $\mathbb{R}^{P\setminus A}$. If nonempty, $\widetilde{\Pi}_{\sigma}$ is the direct product

$$\widetilde{\Pi}_{\sigma} = \prod_{i=1}^{k-1} \{ \lambda_j \ge x_{\sigma(i_j+1)} \ge \dots \ge x_{\sigma(i_{j+1}-1)} \ge \lambda_{j+1} \}$$

where each term $\{\lambda_j \geq x_{\sigma(i_j+1)} \geq \cdots \geq x_{\sigma(i_{j+1}-1)} \geq \lambda_{j+1}\}$ is an $(i_{j+1}-i_j-1)$ -dimensional simplex with volume $\frac{(\lambda_j-\lambda_{j+1})^{i_{j+1}-i_j-1}}{(i_{j+1}-i_j-1)!}$. Thus

Vol
$$\widetilde{\Pi}_{\sigma} = \frac{(\lambda_1 - \lambda_2)^{a_1}}{a_1!} \cdots \frac{(\lambda_{k-1} - \lambda_k)^{a_{k-1}}}{a_{k-1}!},$$

where we set $a_j := i_{j+1} - i_j - 1$. Summing over all linear extensions σ of P, we obtain

$$\operatorname{Vol} \widetilde{\mathcal{O}}(P,A)_{\lambda} = \sum_{\sigma \in \mathcal{L}(P)} \operatorname{Vol} \widetilde{\Pi}_{\sigma}$$

$$= \sum_{a_{1},\dots,a_{k-1} \geq 0} N_{P,A,\lambda}(a_{1},\dots,a_{k-1}) \frac{(\lambda_{1} - \lambda_{2})^{a_{1}}}{a_{1}!} \cdots \frac{(\lambda_{k-1} - \lambda_{k})^{a_{k-1}}}{a_{k-1}!}.$$

Marked order polytopes also enjoy a Minkowski sum property and decomposition.

Theorem 3.5 ([4]). Let P be a poset and A a subposet. If $\lambda, \mu: A \to \mathbb{R}$ are markings, then

$$\mathcal{O}(P,A)_{\lambda+\mu} = \mathcal{O}(P,A)_{\lambda} + \mathcal{O}(P,A)_{\mu}.$$

Corollary 3.6. For (P, A, λ) a marked poset with marked elements $A = \{p_1, \ldots, p_k\}$ having markings $\lambda(p_1) \geq \cdots \geq \lambda(p_k)$, let $\omega_i \colon A \to \mathbb{R}$ be the map such that $\omega_i(p_j) = 1$ for $j \leq i$ and $\omega_i(p_j) = 0$ if j > i. Then, taking $\lambda(p_{k+1})$ to mean 0, $\mathcal{O}(P, A)_{\lambda}$ decomposes into the Minkowski sum

$$\mathcal{O}(P,A)_{\lambda} = \sum_{i=1}^{k} (\lambda(p_i) - \lambda(p_{i+1})) \mathcal{O}(P,A)_{\omega_i}.$$

3.1. A Log-Concavity Result. Recall that a sequence b_0, b_1, \ldots, b_m of non-negative real numbers is said to be log-concave if $b_i^2 \ge b_{i-1}b_{i+1}$ for $1 \le i \le m-1$. In particular, a log-concave sequence is **unimodal**, that is for some j, we have $b_1 \le b_2 \le \cdots \le b_j$ and $b_j \ge b_{j+1} \ge \cdots \ge b_m$.

Using the Alexandrov–Fenchel inequalities and the volume formula for order polytopes, Stanley proved the following log-concavity result in the special case where all marked elements of P lie on a chain in [10].

Theorem 3.7. Let (P, A, λ) be a marked poset with marked elements $A = \{p_1, \ldots, p_k\}$ having markings $\lambda(p_1) \geq \cdots \geq \lambda(p_k)$ denoted $\lambda_1, \ldots, \lambda_k$. If $a_1, \ldots, a_k \geq 0$ with $a_{j-1} > 1$ and $a_j > 1$ for some j, then

$$N_{P,A,\lambda}(a_1,\ldots,a_{k-1})^2 \ge N_{P,A,\lambda}(a_1,\ldots,a_{j-1},a_j-1,a_{j+1},\ldots,a_{k-1})N_{P,A,\lambda}(a_1,\ldots,a_{j-1},a_j+1,a_{j+1},\ldots,a_{k-1}).$$

Before proving Theorem 3.7, we give some background on the theory of mixed volumes and the Alexandrov–Fenchel inequalities following that of [10]. If K_1, \ldots, K_s are convex bodies (nonempty compact convex subsets) of \mathbb{R}^n , fix weights $r_1, \ldots, r_s \geq 0$ and let K denote the Minkowski sum

$$K = r_1 K_1 + \dots + r_s K_s = \{r_1 t_1 + \dots + r_s t_s \mid t_i \in K_i\}.$$

The volume V(K) of K is a homogeneous polynomial of degree n in r_1, \ldots, r_s :

$$V(K) = \sum_{i_1=1}^{s} \sum_{i_2=1}^{s} \cdots \sum_{i_n=1}^{s} V_{i_1,\dots,i_n} r_{i_1} \cdots r_{i_n}.$$

The coefficients $V_{i_1,...,i_n}$ are uniquely determined by requiring they be symmetric up to permutations of subscripts. The coefficient $V_{i_1,...,i_n}$ depends only on $K_{i_1},...,K_{i_n}$ and is called the **mixed volume** of $K_{i_1},...,K_{i_n}$. If we write $V(K_1^{a_1},...,K_s^{a_s})$ for

$$V\left(\underbrace{K_1,\ldots,K_1}_{a_1},\ldots,\underbrace{K_s,\ldots,K_s}_{a_s}\right),$$

then

$$V(K) = \sum_{a_1 + \dots + a_s = n} \binom{n}{a_1, \dots, a_s} V(K_1^{a_1}, \dots, K_s^{a_s}) r_1^{a_1} \cdots r_s^{a_s}.$$

The well-known result about mixed volumes needed for the proof of Theorem 3.7 is the following.

Theorem 3.8 (Alexandrov–Fenchel Inequalities, [12]). Given $0 \le m \le n$ and convex bodies $C_1, \ldots, C_{n-m}, K, L \subseteq \mathbb{R}^n$, the sequence (b_0, b_1, \ldots, b_m) defined by

$$b_i = V(C_1, \dots, C_{n-m}, K^{m-i}, L^i)$$

is log-concave.

We can now give the proof of Theorem 3.7.

Proof of Theorem 3.7. Corollary 3.6 yields the Minkowski sum

$$\widetilde{\mathcal{O}}(P,A)_{\lambda} = \sum_{i=1}^{k} (\lambda_i - \lambda_{i+1}) \widetilde{\mathcal{O}}(P,A)_{\omega_i},$$

so taking $\lambda_{k+1} = 0$,

$$\operatorname{Vol} \widetilde{\mathcal{O}}(P,A)_{\lambda} = \operatorname{Vol} \sum_{i=1}^{k} (\lambda_{i} - \lambda_{i+1}) \widetilde{\mathcal{O}}(P,A)_{\omega_{i}}$$

$$= \sum_{a_{1}+\dots+a_{k-1}=|P|-k} {|P|-k \choose a_{1},\dots,a_{k-1}} V(\widetilde{\mathcal{O}}(P,A)_{\omega_{1}}^{a_{1}},\dots,\widetilde{\mathcal{O}}(P,A)_{\omega_{k-1}}^{a_{k-1}}) (\lambda_{1} - \lambda_{2})^{a_{1}} \cdots (\lambda_{k-1} - \lambda_{k})^{a_{k-1}}.$$

Comparing this with the volume formula

Vol
$$\widetilde{\mathcal{O}}(P,A)_{\lambda} = \sum_{a_1,\dots,a_{k-1}\geq 0} N_{P,A,\lambda}(a_1,\dots,a_{k-1}) \frac{(\lambda_1 - \lambda_2)^{a_1}}{a_1!} \cdot \dots \cdot \frac{(\lambda_{k-1} - \lambda_k)^{a_{k-1}}}{a_{k-1}!}$$

of Theorem 3.4, we obtain

$$N_{P,A,\lambda}(a_1,\ldots,a_{k-1}) = (|P|-k)!V(\widetilde{\mathcal{O}}(P,A)^{a_1}_{\omega_1},\ldots,\widetilde{\mathcal{O}}(P,A)^{a_{k-1}}_{\omega_{k-1}}).$$

4. Marked order polytopes as flow polytopes

In this section we prove that for strongly planar posets with special markings, the marked order polytopes are integrally equivalent to flow polytopes. This generalizes a result of Mészáros-Morales-Striker [7, Theorem 3.14] for (unmarked) order polytopes, which we now review.

A poset P is **strongly planar** if the Hasse diagram of $\widehat{P} := P \sqcup \{\widehat{0}, \widehat{1}\}$ is planar and can be drawn in the plane so that the y-coordinates of vertices respect the order of P. When we refer to a **bounded embedding** of P, we will mean a strongly planar drawing of the Hasse diagram of \widehat{P} with an additional two edges between $\widehat{0}$ and $\widehat{1}$ added, one drawn to the left of \widehat{P} and the other drawn to the right (see Figure 2). We will view this embedding as a planar graph and discuss its (bounded) faces in the usual graph-theoretic sense.

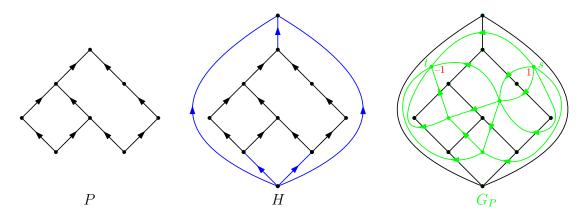


FIGURE 2. The Hasse diagram of a poset P (left), a bounded embedding H of P (middle), and the directed graph G_P drawn over H (right).

We begin by recalling the case of order polytopes. Given a strongly planar poset P, let H be a bounded embedding of P (viewed as a planar graph). Let H^* be the graph-theoretic dual of H. Define G_P to be the subgraph of H^* obtained by deleting the vertex corresponding to the unbounded face of H. Denote the two vertices of G that lie in faces of H containing the edges $(\hat{0}, \hat{1})$ by S and S with S on the right and S on the left

Assign each edge e in G_P an orientation by the following rule: in the construction of H^* the edge e crosses an edge p < q of H; orient e so that while traversing e, q is on the right and p is on the left. Make G_P into a flow network by assigning netflow 1 to s, -1 to t, and 0 to all other vertices.

Theorem 4.1 ([7, Theorem 3.14]). Let P be a strongly planar poset and G_P be the flow network constructed above. The polytopes $\mathcal{O}(P)$ and \mathcal{F}_{G_P} are integrally equivalent.

Proof sketch. The map from $\mathcal{O}(P) \to \mathcal{F}_{G_P}$ is given by $(x_p)_{p \in P} \mapsto f$ where $f(e) = x_q - x_p$ if e crosses the edge p < q in H and $x_{\hat{0}}$, $x_{\hat{1}}$ are taken to be 0 and 1 respectively. For the inverse, take a flow f on G_P . For each $p \in P$, choose any path in H from $\hat{0}$ to p. To define x_p , sum the flow values f(e) on each edge $e \in G_P$ crossing an edge in the chosen path from $\hat{0}$ to p in H.

We now generalize Theorem 4.1 to marked order polytopes. We begin with some terminology used to define the marked analogue of a strongly planar poset. If F is a bounded face of a bounded embedding H of P, let p denote the minimum element of F and let q denote the maximum. The graph $F \setminus \{p, q\}$ has two components whose unions with $\{p, q\}$ we will call the **left and right boundaries** of F.

Definition 4.2. A marked poset (P, A, λ) is called **strongly planar** if P is strongly planar as an unmarked poset and admits a bounded embedding H such that for each bounded face $F \subseteq \widehat{P}$ of H, if the left boundary of F (including $\min(F)$ and $\max(F)$) contains a marked element, then both $\min(F)$ and $\max(F)$ are marked. We will call such an embedding a **bounded strongly planar embedding.**

Remark 4.3. We note that in Definition 4.2 we made a choice to put conditions on the markings on the left boundaries of bounded faces. Of course we could have put those conditions on the right boundaries instead. Moreover, as Remark 4.4 explains, the definition can be relaxed by mixing and matching left and right boundaries of bounded faces under certain conditions in such a way that the main result, Theorem 1.2, still holds.

For a bounded strongly planar embedding H of the marked poset (P, A, λ) , we now construct a flow network $G_{(P,A,\lambda)}$ from G_P . Begin with a bounded strongly planar embedding H and the flow network G_P constructed from H, as in the case of order polytopes. View the markings λ as being on A inside of H, and add additional markings $\min\{\lambda(a) \mid a \in A\}$ on $\hat{0}$ and $\max\{\lambda(a) \mid a \in A\}$ on $\hat{1}$.

Recall that each vertex of G_P is naturally labeled by a bounded face of H. (In the rest of the paper, whenever we refer to a face of bounded strongly planar embedding H, we mean a bounded face.) Denote the vertex labeled by a face F by v_F . Starting from G_P , construct a flow network $G_{(P,A,\lambda)}$ by applying the following construction to v_F for each (bounded) face F of H. See Figure 3 for an illustration of this construction.

If F contains no marked elements on its left boundary, do nothing, and let v_F continue to have netflow 0. Otherwise, suppose the left boundary of F is composed of elements $p_1 > \cdots > p_k$ in H, with $\min(F) = p_k$ and $\max(F) = p_1$. Since some point on the left boundary of F is marked, so are $p_1 = q_1$ and $p_k = q_\ell$ by strong planarity. Suppose the marked elements among p_1, \ldots, p_k are $p_1 = p_{i_1} > p_{i_2} > \cdots > p_{i_g} = p_k$ marked by $a_1 \geq a_2 \geq \cdots \geq a_g$. Delete the edges outgoing from vertex v_F in G_P , and let v_F become a sink with netflow $-(p_1 - p_k)$, with the incoming edges as before. The edges previously outgoing from v_F in G_P that crossed the left boundary of F between marked elements p_{i_m} and $p_{i_{m+1}}$ will now be outgoing from the source vertex s_m^F , for $m \in [g-1]$. Assign s_m^F netflow $a_m - a_{m+1}$ for each m.

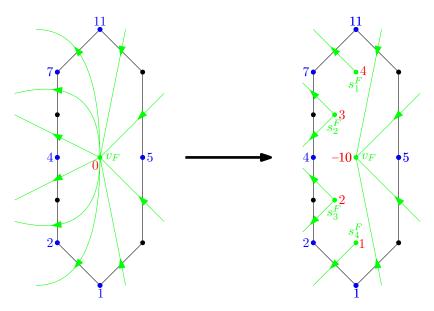


FIGURE 3. An illustration of the construction of $G_{(P,A,\lambda)}$ from G_P on a single face F.

Theorem 1.2. Given a bounded strongly planar embedding H of a marked poset (P, A, λ) , the marked order polytope $\mathcal{O}(P, A)_{\lambda}$ is integrally equivalent to the flow polytope $\mathcal{F}_{G_{(P,A,\lambda)}}$, where $G_{(P,A,\lambda)}$ is the flow network described above.

Proof. The integral equivalences between $\mathcal{O}(P,A)_{\lambda}$ and $\mathcal{F}_{G_{(P,A,\lambda)}}$ are exactly as in the order polytope case. The map $\Gamma \colon \mathcal{O}(P,A)_{\lambda} \to \mathcal{F}_{G_{(P,A,\lambda)}}$ is given by $(x_p)_{p \in P} \mapsto f$ where $f(e) = x_q - x_p$ if e crosses the edge p < q in H. The inverse map Γ^{-1} is given by $x_p = \sum_e f(e)$ over edges $e \in G_{(P,A,\lambda)}$ crossing any fixed path from $\hat{0}$ to p in H. (Note that from any marked point $p \in A$, there exists a path from p to $\hat{0}$ in H that only walks along the left boundaries of faces to the minimums of those faces.) The details of the proof are analogous to those in [7] and are left to the reader.

Theorem 1.2 provides a general framework for obtaining the graphs G_{λ} used in Theorem 2.3 and for proving Theorem 2.3. See Figure 4 for an example of this.

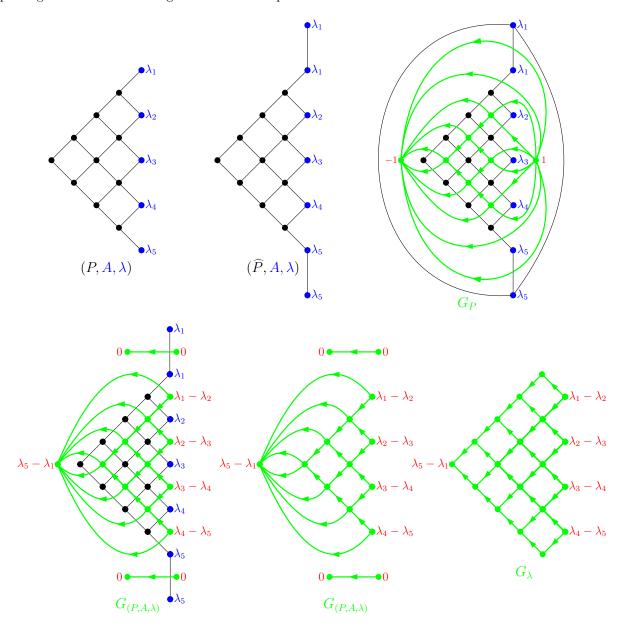


FIGURE 4. Top row: the marked poset (P, A, λ) with $\mathcal{O}(P, A)_{\lambda} = \operatorname{GT}(\lambda)$, $\widehat{P} = P \sqcup \{\widehat{0}, \widehat{1}\}$, and the flow network G_P ; Bottom row: the flow network $G_{(P,A,\lambda)}$ and the integrally equivalent flow network G_{λ} of Definition 2.4.

Remark 4.4. Note that Theorem 1.2 can be generalized in various ways. We can obtain slightly different conditions on the markings of strongly planar posets under which the statement of Theorem 1.2 as well as the map given in its proof are still correct. We picked the above particular definition for bounded strongly planar embeddings relying on conditions on the left boundaries of the bounded faces of the embedding as it seemed the least technical to state. We could have, of course, equally worked with right boundaries of the bounded faces of the embedding, or, we could mix and match as to when we consider the left or right boundary of a bounded face as long as we ensure that the flow conditions pick up the restriction coming from two marked points that are comparable but do not lie in a common face. Next, we give an example of

how relaxing the marking conditions in Theorem 1.2 yields a new proof that skew Gelfand–Tsetlin polytopes are flow polytopes.

Definition 4.5. Given partitions $\lambda, \mu \in \mathbb{Z}_{\geq 0}^n$ and $m \in \mathbb{N}$, the **skew Gelfand–Tsetlin polytope** $GT(\lambda/\mu, m)$ is the set of all arrays

$$x_{m1} \ x_{m2} \ \cdots \ x_{mn} \ \cdots \ x_{1n} \ x_{11} \ x_{12} \ \cdots \ x_{1n} \ x_{01} \ x_{02} \ \cdots \ x_{0n}$$

with top row λ and bottom row μ such that $x_{ij} \geq x_{i-1,j}$ and $x_{ij} \geq x_{i+1,j+1}$ for all i, j.

Proposition 4.6. Skew Gelfand–Tsetlin polytopes are marked order polytopes of strongly planar marked posets.

Proof. Given λ , μ , and m, begin with a skew Gelfand–Tsetlin array $(x_{ij})_{i,j}$. Replace each entry x_{ij} by a vertex, and each relation $x_{ij} \geq x_{i-1,j}$ or $x_{ij} \geq x_{i+1,j+1}$ by an edge between the corresponding vertices. Mark the top row of vertices with the corresponding entries of λ , and mark the bottom row of vertices with the corresponding entries of μ . Rotate the graph 90 degrees clockwise. The result is the Hasse diagram of a strongly planar marked poset (P, A, λ) with $\mathcal{O}(P, A, \lambda) = \operatorname{GT}(\lambda/\mu, m)$. See Figure 5 for an example of this construction.

Corollary 4.7 ([3]). Skew Gelfand-Tsetlin polytopes are integrally equivalent to flow polytopes.

Proof. Apply the generalization of Theorem 1.2 explained in Remark 4.4 to the poset constructed in Lemma 4.6. See Figure 5 for an example of the resulting flow network. \Box

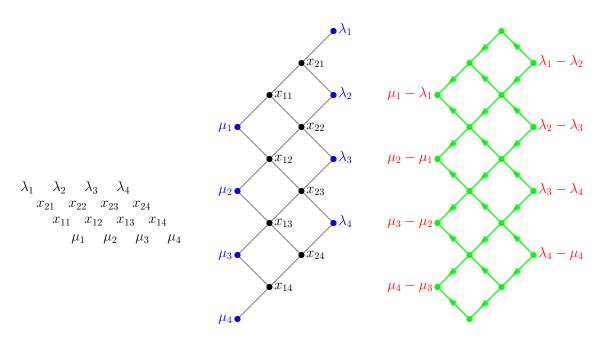


FIGURE 5. An example of the marked poset and corresponding flow network for recognizing a skew Gelfand–Tsetlin polytope $GT(\lambda/\mu, 3)$ with $\ell(\lambda) = 4$ as a marked order polytope and a flow polytope.

Remark 4.8. In [3], a family of polytopes containing the skew Gelfand–Tsetlin polytopes is considered, and all members of the family are shown to be flow polytopes. It is easy to see using the idea of Proposition 4.6 that all members of this family are also strongly planar marked order polytopes.

5. Subdivisions of marked order and flow polytopes

In this section, we will give subdivision procedures for $\mathcal{O}(P,A)_{\lambda}$ and $\mathcal{F}_{G_{(P,A,\lambda)}}$ and prove the two procedures are equivalent. In particular, this will yield a bijective proof of Corollary 2.5.

We start by reviewing the subdivision procedure for flow polytopes following the exposition of [6]. However, we will use a simplified version of the flow polytope subdivision method presented there, specialized to the types of graphs that appear in the present paper.

5.1. Subdividing flow polytopes into products of simplices. Flow polytopes admit a combinatorial iterative subdivision procedure. To describe the algorithm, we first introduce the necessary terminology and notation. A bipartite noncrossing tree is a tree with a distinguished bipartition of vertices into left vertices x_1, \ldots, x_ℓ and right vertices $x_{\ell+1}, \ldots, x_{\ell+r}$ with no pair of edges $(x_p, x_{\ell+q}), (x_t, x_{\ell+u})$ where p < t and q > u. Denote by $\mathcal{T}_{L,R}$ the set of bipartite noncrossing trees, where L and R are the ordered sets (x_1, \ldots, x_ℓ) and $(x_{\ell+1}, \ldots, x_{\ell+r})$ respectively. Note that $|\mathcal{T}_{L,R}| = {\ell+r-2 \choose \ell-1}$, since they are in bijection with weak compositions of r-1 into ℓ parts: a tree T in $\mathcal{T}_{L,R}$ corresponds to the composition $(b_1-1, \ldots, b_\ell-1)$ of r-1, where b_i denotes the number of edges incident to the left vertex $x_{\ell+i}$ in T.

The bipartite noncrossing tree encoded by the composition (0,2,1,1) is the following:



Consider a graph G on the vertex set [n+1] and an integer netflow vector $a=(a_1,\ldots,a_n,-\sum_i a_i)$. In this paper, we will assume that $a_i>0$ implies i has no incoming edges, $a_i<0$ implies i has no outgoing edges, and $a_i=0$ implies i has both incoming and outgoing edges. For these flow networks, the basic step of the subdivision method is the following:

Pick an arbitrary vertex i of G with netflow $a_i = 0$. Let $\mathcal{I}_i = \mathcal{I}_i(G)$ be the multiset of edges incoming to i, edges of the form (\cdot, i) . Let $\mathcal{O}_i = \mathcal{O}_i(G)$ be the multiset of outgoing edges from i, edges of the form (i, \cdot) .

Assign an ordering to the sets \mathcal{I}_i and \mathcal{O}_i , and consider a tree $T \in \mathcal{T}_{\mathcal{I}_i,\mathcal{O}_i}$. For each tree-edge (e_1,e_2) of T, where $e_1 = (r,i) \in \mathcal{I}_i$ and $e_2 = (i,s) \in \mathcal{O}_i$, let $\mathrm{edge}(e_1,e_2) = (r,s)$. Define a graph $G_T^{(i)}$ by starting with G, deleting vertex i and all incident edges $\mathcal{I}_i \cup \mathcal{O}_i$ of G, and adding the multiset of edges $\{\mathrm{edge}(e_1,e_2) \mid (e_1,e_2) \in E(T)\}$. See Figure 6 for an example.

Lemma 5.1 (Compounded Subdivision Lemma). Let G be a flow network on the vertex set [n+1] with netflow $a = (a_1, \ldots, a_n, -\sum_{i=1}^n a_i) \in \mathbb{Z}^{n+1}$ and a vertex $i \in \{2, \ldots, n\}$ with $a_i = 0$. Then, $\{\mathcal{F}_{G_T^{(i)}}(\hat{a})\}_{T \in \mathcal{T}_{\mathcal{I}_i, \mathcal{O}_i}}$ are top dimensional pieces in a subdivision of $\mathcal{F}_G(a)$, where \hat{a} equals a with ith coordinate deleted.

In order to view $\mathcal{F}_{G_T^{(i)}}$ as a subset of \mathcal{F}_G , label each edge e of G with a coordinate f_e . Label each new edge of $G_T^{(i)}$ by the formal sum of the coordinates of the edges of G that formed it. To get an inclusion $\mathcal{F}_{G_T^{(i)}} \subseteq \mathcal{F}_G$, simply add the flow value of each edge in $G_T^{(i)}$ to all edges of G appearing in the formal sum labeling it.

We refer to replacing G by $\{G_T^{(i)}\}_{T\in\mathcal{T}_{L,R}}$ as a **compounded reduction** on G. In order to fully subdivide \mathcal{F}_G into simplices, one performs a compounded reduction on G, then iteratively performs compound reductions on the graphs $G_T^{(i)}$. A series of these reductions can be efficiently encoded by a **compounded reduction tree**: the root of the tree is the original graph G; and the descendants of any node are the graphs obtained via a compounded reduction on that node. See Figure 6 for an example. The **canonical compounded reduction tree** of G is obtained by performing compounded reductions from highest to lowest index netflow zero vertices, as in Figure 6. There is a natural way of labeling the products of simplices into which we subdivide our flow polytope via the compounded reductions by integer points of other flow polytopes, as is explained in [6].

5.2. Subdividing order polytopes into products of simplices. Given a bounded strongly planar embedding H of a marked poset (P, A, λ) , consider the following method for subdividing $\mathcal{O}(P, A)_{\lambda}$: Consider any face F of H not containing an edge $(\hat{0}, \hat{1})$ (where, as previously, by face we mean bounded face). Suppose that F is bounded on the left by $p_1 > \cdots > p_k$ and on the right by $p_1 = q_1 > \cdots > q_\ell = p_k$. Replacing F

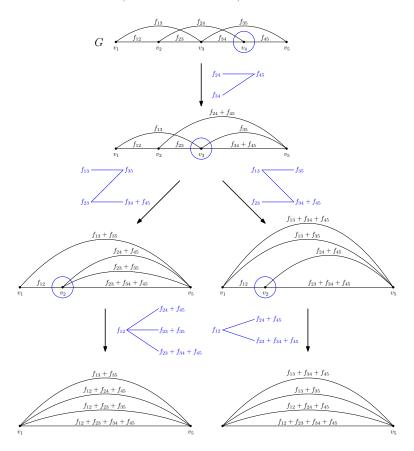


FIGURE 6. An example of a compounded reduction tree for a graph G with netflow (1,0,0,0,-1). The tree is built downward by performing a series of compounded reductions, starting with the root G. The labels on edges record the inclusion maps from the flow polytope of each graph to the flow polytope of G.

by any of the $\mathcal{N} = \binom{k+\ell-4}{\ell-2}$ linear extensions of $p_1, \ldots, p_k, q_2, \ldots, q_{\ell-1}$, we obtain \mathcal{N} strongly planar marked posets $(P_1, A, \lambda), \ldots, (P_{\mathcal{N}}, A, \lambda)$.

Lemma 5.2. The marked order polytopes $\mathcal{O}(P_1, A, \lambda), \ldots, \mathcal{O}(P_N, A, \lambda)$ described above form a subdivision of the order polytope $\mathcal{O}(P, A, \lambda)$.

Proof. This subdivision is obtained by cutting $\mathcal{O}(P,A)_{\lambda}$ by the hyperplanes $x_{p_i}=x_{q_j}$ for $i\in[1,k]$ and $j\in[1,\ell]$.

By the above lemma, applying the above construction iteratively to each face of the bounded strongly planar embedding H of the marked poset (P, A, λ) yields a subdivision of $\mathcal{O}(P, A)_{\lambda}$ into the marked poset polytopes of a set of marked chains, that is, into products of simplices.

5.3. Comparing the subdivisions of flow and order polytopes. Theorem 1.2 shows $\mathcal{O}(P,A)_{\lambda}$ is integrally equivalent to a flow polytope $\mathcal{F}_{G_{(P,A,\lambda)}}$. As we saw in Sections 5.1 and 5.2, both flow and order polytopes admit an iterative subdivision procedure. We show here that indeed those procedures can be considered identical.

Through a single application of Lemma 5.1 on $\mathcal{F}_{G_{(P,A,\lambda)}}$, we obtain the following.

Lemma 5.3. Given a bounded strongly planar embedding H of the marked poset (P, A, λ) , consider a face F of H which has no markings on its left boundary. Linearly order the k-1 outgoing and $\ell-1$ incoming edges of v_F from top to bottom. Performing a compounded reduction at v_F on $G_{(P,A,\lambda)}$ yields $\mathcal{N} = \binom{k+\ell-4}{\ell-2}$ flow networks $G_{(P,A,\lambda)}^1, \ldots, G_{(P,A,\lambda)}^{\mathcal{N}}$ such that the polytopes $\mathcal{F}_{G_{(P,A,\lambda)}^1}, \ldots, \mathcal{F}_{G_{(P,A,\lambda)}^{\mathcal{N}}}$ subdivide $\mathcal{F}_{G_{(P,A,\lambda)}}$.

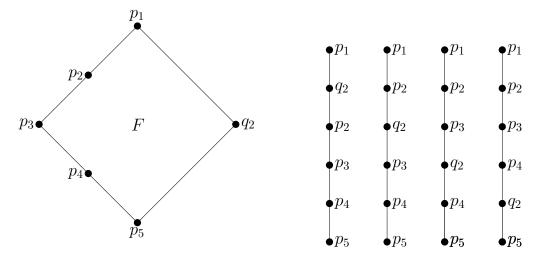


FIGURE 7. A face F with unmarked left boundary in the embedding of a strongly planar poset P (left) and the corresponding linear extensions of F that replace F to form the posets P_1, \ldots, P_N when subdividing $\mathcal{O}(P, A, \lambda)$ (right).

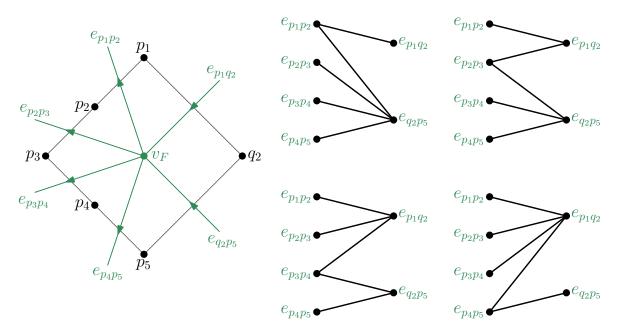


FIGURE 8. A netflow zero vertex v_F of $G_{(P,A,\lambda)}$ (left) and the corresponding bipartite noncrossing trees that can be used to perform a compounded reduction at v_F to produce flow networks $G^1_{(P,A,\lambda)}, \ldots, G^N_{(P,A,\lambda)}$ whose flow polytopes subdivide $\mathcal{F}_{G_{(P,A,\lambda)}}$ (right).

We now describe an equivalence between the subdivision procedures of $\mathcal{O}(P,A)_{\lambda}$ and $\mathcal{F}_{G(P,A,\lambda)}$ whose basic step is described in Lemma 5.2 and Lemma 5.3 respectively. We first focus on the case of a single step of both subdivisions.

Lemma 5.4. Given a bounded strongly planar embedding H of a marked poset (P, A, λ) , let F be a face of H with unmarked left boundary. Let Γ be the integral equivalence $\mathcal{O}(P, A)_{\lambda} \to \mathcal{F}_{G_{(P,A,\lambda)}}$ of Theorem 1.2. Then there is a bijection $\gamma \colon [\mathcal{N}] \to [\mathcal{N}]$ between the linear extensions of F and the bipartite noncrossing trees from a compounded reduction at v_F such that

$$\Gamma \colon \mathcal{O}(P_j, A)_{\lambda} \to \mathcal{F}_{G_{(P,A,\lambda)}^{\gamma(j)}}.$$

Proof. Let F be bounded by $p_1 > p_2 > \cdots > p_k$ on the left and $p_1 = q_1 > q_2 > \cdots > q_\ell = p_k$ on the right. To define the bijection, we start by drawing the bipartite noncrossing trees with the vertices arranged in vertical columns. Label each vertex of the tree by the edge of H dual to the edge of $G_{(P,A,\lambda)}$ it represents. Encase each tree in a bounding rectangle so that the vertex columns lie on the interiors of the sides of the rectangle. See Figure 9.

To construct a linear order from a tree, we will label the regions of the rectangle cut out by the tree. Label the top region p_1 and the bottom p_k . All intermediate regions are triangles with exactly one edge on the bounding rectangle. Label such regions by the common label of the endpoints of this edge. The result will be a linear order of the face F. Conversely, a linear ordering gives an ordering on the edge segments on each side of the rectangle. Build the tree top to bottom by adding in edges inside the bounding rectangle to cut out regions as specified by the linear order from top to bottom.

To see that γ has the property

$$\Gamma \colon \mathcal{O}(P_j, A)_{\lambda} \to \mathcal{F}_{G_{(P,A,\lambda)}^{\gamma(j)}}$$

it suffices to note that γ is constructed precisely so that $G_{(P_j,A,\lambda)} = G_{(P,A,\lambda)}^{\gamma(j)}$ for each $j \in [\mathcal{N}]$.

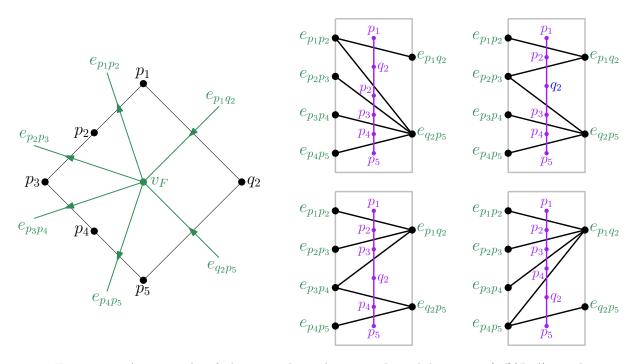


FIGURE 9. An example of the equivalence between the subdivisions of $\mathcal{O}(P,A)_{\lambda}$ and $\mathcal{F}_{G_{(P,A,\lambda)}}$ through a bijection of linear orders and bipartite noncrossing trees.

Theorem 5.5. Let H be a bounded strongly planar embedding of the marked poset (P, A, λ) . Choose an ordering F_1, \ldots, F_m of the faces of H that contain no marked elements on their respective left boundaries. Let Δ be the subdivision of $\mathcal{O}(P,A)_{\lambda}$ obtained by applying Lemma 5.2 to each of F_1, \ldots, F_m . Let Δ' be the subdivision of $\mathcal{F}_{G(P,A,\lambda)}$ obtained by applying Lemma 5.3 to each of v_{F_1}, \ldots, v_{F_m} in $G_{(P,A,\lambda)}$. Then the integral equivalence $\Gamma \colon \mathcal{O}(P,A)_{\lambda} \to \mathcal{F}_{G(P,A,\lambda)}$ induces a bijection γ from regions of Δ to regions of Δ' .

Proof. Apply Lemma 5.4 iteratively to each of P_1, \ldots, P_N .

5.4. Bijecting the combinatorial objects labeling the subdivisions of flow and order polytopes. Now we are ready to give a bijective proof of a generalization of Corollary 2.5. Figure 10 provides a detailed example of the bijection.

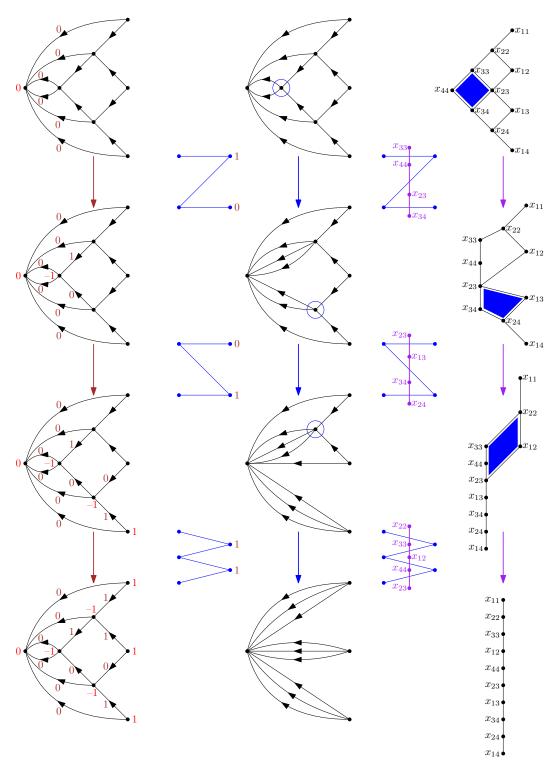


FIGURE 10. An example of the bijective proof of Corollary 5.6. In the left column, a flow on $G_{(P,A,\lambda)}$ is mapped to a leaf of the canonical compounded reduction tree of $G_{(P,A,\lambda)}$. Then, each step of the path (middle column) from the root to that leaf is mapped to a marked order polytope by γ (right column), the last of which is a linear order.

Corollary 5.6. Let H be a bounded strongly planar embedding of the marked poset (P, A, λ) with $A = \{p_1, \ldots, p_k\}$ such that $\lambda(p_1) \geq \cdots \geq \lambda(p_k)$. Additionally, assume P is marked in such a way that $G_{(P,A,\lambda)}$ has only one sink. Order the n vertices of $G_{P,A,\lambda}$ so that sources corresponding to p_1, \ldots, p_k are first, edges go from earlier to later vertices, and the sink is last. Then

$$N_{P,A,\lambda}(a_1,\ldots,a_{k-1}) = K_{G_{(P,A,\lambda)}}(a_1 - \text{out}_1,\ldots,a_{k-1} - \text{out}_{k-1},-\text{out}_k,\ldots,-\text{out}_{n-1},0),$$

where out_j is the outdegree of vertex j in $G_{(P,A,\lambda)}$ minus 1.

Proof. Choose an ordering of the vertices of $G_{(P,A,\lambda)}$ so that all edges go from earlier to later vertices in the order. Let $v_{F_{i_1}} < \cdots < v_{F_{i_m}}$ be the induced order of the vertices v_F corresponding to faces with unmarked left boundary. Let Δ be the subdivision of $\mathcal{O}(P,A)_{\lambda}$ and Δ' the subdivision of $\mathcal{F}_{G_{(P,A,\lambda)}}$ obtained by using Lemma 5.2 on F_{i_m}, \ldots, F_{i_1} and Lemma 5.3 on $v_{F_{i_m}}, \ldots, v_{F_{i_1}}$ respectively. The integral equivalence $\Gamma \colon \mathcal{O}(P,A)_{\lambda} \to \mathcal{F}_{G_{(P,A,\lambda)}}$ induces a bijection γ from regions of Δ to regions of Δ' .

As described in [6] Lemma 4.1, flows on $G_{P,A,\lambda}$ with netflow

$$(a_1 - \text{out}_1, \dots, a_{k-1} - \text{out}_{k-1}, -\text{out}_k, \dots, -\text{out}_{n-1}, 0)$$

are in bijection with leaves of the canonical compounded reduction tree of $G_{(P,A,\lambda)}$ with a_i edges outgoing from the *i*th source vertex. The flow values on edges incoming to each vertex are read off from the composition corresponding to the noncrossing bipartite tree chosen when reducing that vertex. The volume-preserving bijection γ provides a correspondence between these leaves and linear extensions of P with the marked elements in positions $1, 2 + a_1, \ldots, k + a_1 + \cdots + a_{k-1}$.

See Figure 10 for an illustration of Corollary 5.6.

References

- [1] F. Ardila, T. Bliem, and D. Salazar. Gelfand–Tsetlin polytopes and Feigin–Fourier–Littelmann–Vinberg polytopes as marked poset polytopes. *J. Combin. Theory Ser. A*, 118(8):2454–2462, November 2011.
- [2] W. Baldoni and M. Vergne. Kostant partitions functions and flow polytopes. Transform. Groups, 13(3-4):447–469, 2008.
- [3] V. Danilov, A. Karzanov, and G. Koshevoy. Discrete strip-concave functions, Gelfand-Tsetlin patterns, and related polyhedra. J. Combin. Theory Ser. A, 112(2):175–193, 2005.
- [4] X. Fang and G. Fourier. Marked chain-order polytopes. European J. Combin., 58:267 282, 2016.
- [5] R.I. Liu, K. Mészáros, and A. St. Dizier. Schubert polynomials as projections of Minkowski sums of Gelfand–Tsetlin polytopes, 2019. arXiv:1903.05548.
- [6] K. Mészáros and A. H. Morales. Volumes and Ehrhart polynomials of flow polytopes. Math. Z., to appear, 2019.
- [7] K. Mészáros, A. H. Morales, and J. Striker. On flow polytopes, order polytopes, and certain faces of the alternating sign matrix polytope. *Discrete Comput. Geom.*, to appear.
- [8] C. Pegel. The face structure and geometry of marked order polyhedra. Order, 35(3):467-488, Nov 2018.
- [9] A. Postnikov. Permutohedra, associahedra, and beyond. Int. Math. Res. Not. IMRN, 2009(6):1026-1106, 2005.
- [10] R. P. Stanley. Two combinatorial applications of the Aleksandrov–Fenchel inequalities. J. Combin. Theory Ser. A, 31(1):56–65, 1981.
- [11] R. P. Stanley. Two poset polytopes. Discrete Comput. Geom., 1(1):9–23, Mar 1986.
- [12] R. P. Stanley. Log-concave and unimodal sequences in algebra, combinatorics, and geometry. In Graph theory and its applications: East and West (Jinan, 1986), volume 576 of Ann. New York Acad. Sci., pages 500–535. New York Acad. Sci., New York, 1989.
- [13] V.Kiritchenko, E. Smirnov, and V. Timorin. Schubert calculus and Gelfand-Zetlin polytopes. Russian Math. Surveys, 67:685–719, 12 2012.

RICKY INI LIU, DEPARTMENT OF MATHEMATICS, NORTH CAROLINA STATE UNIVERSITY, RALEIGH, NC 27695. RILIU@NCSU.EDU

KAROLA MÉSZÁROS, DEPARTMENT OF MATHEMATICS, CORNELL UNIVERSITY, ITHACA, NY 14853 AND SCHOOL OF MATHEMATICS, INSTITUTE FOR ADVANCED STUDY, PRINCETON, NJ 08540. KAROLA@MATH.CORNELL.EDU

AVERY ST. DIZIER, DEPARTMENT OF MATHEMATICS, CORNELL UNIVERSITY, ITHACA NY 14853. AJS624@CORNELL.EDU