THE *TESTIDEALS* PACKAGE FOR *MACAULAY2*

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ABSTRACT. This note describes a *Macaulay2* package for computations in prime characteristic commutative algebra. This includes Frobenius powers and roots, p^{-e}-linear and p^e-linear maps, singularities defined in terms of these maps, different types of test ideals and modules, and ideals compatible with a given p^{-e}-linear map.

1. INTRODUCTION

This paper describes methods for computing objects and numerical invariants in prime characteristic commutative algebra, implemented in the package *TestIdeals* for the computer algebra system *Macaulay2* [GS]. A ring R of prime characteristic p comes equipped with the Frobenius endomorphism

$$F : R \longrightarrow R \text{ given by } F(x) = x^p,$$

which is the basis for many constructions and methods. Notably, the Frobenius endomorphism can be used to detect whether a ring is regular [Kun69], and further, to quantify how far a ring is from being regular, measuring the severity of a singularity.

In this direction, two notable applications of the Frobenius endomorphism are the theory of tight closure (see [HH90, Hoc07] for an introduction) and the resulting theory of test ideals (see the survey [ST12]). These methods are used by a wide group of commutative algebraists and algebraic geometers.

The *TestIdeals* package was started during a *Macaulay2* development workshop in 2012, hosted by Wake Forest University. The package, at that time called *PosChar*, aimed at providing a unified and efficient set of tools to study singularities in characteristic $p > 0$, and in particular, collecting and implementing several algorithms that had been described in research papers. Since then, *PosChar* was split into two packages, *TestIdeals* and *FThresholds*, and much more functionality

has been added by many contributors, during several more *Macaulay2* development workshops.[1]

Starting at least with Kunz [Kun69] and Fedder [Fed83], it has been known that the Frobenius endormorphism offers *effective* methods for measuring singularities in positive characteristic. However, the algorithms that form the basis of this package are the methods for computing Frobenius roots, $\star$-closures, and parameter test ideals. These first appeared in [Kat08, BMS08, BMS09, Kat10]. Another algorithm in [KS12] for computing prime ideals compatible with a given p^{-e}-linear map was implemented and used to produce the examples in that paper. The methods for computing test ideals and test modules that were used implicitly in papers such as [BSTZ10, KLZ09, ST14a] became implementable via the Frobenius roots functionality. Algorithms for computing F-pure thresholds from [Her15], [Her14], and [HT17] form some of the key methods in the forthcoming *FThresholds* package.

2. Frobenius powers and Frobenius roots

Let R denote a commutative ring of prime characteristic p.

Definition 2.1. Given an ideal $I \subseteq R$ and an integer $e \geqslant 0$, we define the p^e-*th Frobenius power of* I, denoted $I^{[p^e]}$, to be the ideal generated by the p^e-th powers of all elements of I.

It is easy to see that if I is generated by $g_1, \ldots, g_\ell$, then $I^{[p^e]}$ is generated by $g_1^{p^e}, \ldots, g_\ell^{p^e}$.

Definition 2.2. Given an ideal $I \subseteq R$ and an integer $e \geqslant 0$, we define the p^e-*th Frobenius root of* I, denoted $I^{[1/p^e]}$, to be the smallest ideal J such that $I \subseteq J^{[p^e]}$, if such an ideal exists.

Frobenius roots always exist in polynomial and power series rings, and in F-finite regular rings (cf. [BMS08, §2] and [Kat08, §5]). Below is an example of a computation of a Frobenius root in a polynomial ring. In Section 2.1 we describe the main ideas behind such computations.

```
i1 : R = ZZ/5[x,y,z];

i2 : I = ideal(x^6*y*z + x^2*y^12*z^3 + x*y*z^18);

o2 : Ideal of R
```

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[1]Development workshops hosted by the University of California, Berkeley (2014, 2017), Boise State University (2015), and the University of Utah (2016).

```
i3 : frobeniusPower(1/5, I)
            2   3
o3 = ideal (x, y , z )
o3 : Ideal of R
```

2.1. The mathematics behind the computation of Frobenius roots.

We can also describe Frobenius roots as follows: In a (sufficiently local) regular ring, we have an identification of R with its canonical module ω_R. On the other hand, the Grothendieck dual to the e-iterated Frobenius map[2] $R \longrightarrow R^{1/p^e}$ provides a map

$$(1) \qquad\qquad T : \omega_{R^{1/p^e}} \longrightarrow \omega_R,$$

which using the identification $R \cong \omega_R$ gives us a map

$$T : R^{1/p^e} \to R.$$

It is not difficult to prove that $T(I^{1/p^e}) = I^{[1/p^e]}$, where $I^{1/p^e} \subseteq R^{1/p^e}$ is simply the ideal of p^e-th roots of the elements of I. In the case that

$$R = \mathbb{K}[x_1, \ldots, x_d],$$

where $\mathbb{K}$ is a perfect field, R^{1/p^e} is a free R-module with basis consisting of the monomials

$$\mathbf{x}^{\boldsymbol{\lambda}/p^e} = x_1^{\lambda_1/p^e} \cdots x_d^{\lambda_d/p^e},$$

where $0 \leqslant \lambda_i \leqslant p^e - 1$. Furthermore, the map $T : R^{1/p^e} \to R$ is simply the projection defined as follows:

$$T(x_1^{\lambda_1/p^e} \cdots x_d^{\lambda_d/p^e}) = \begin{cases} 1 & \text{if } \lambda_i = p^e - 1 \text{ for all } i \\ 0 & \text{otherwise} \end{cases}$$

Using this, it is not difficult to see that if

$$(2) \qquad\qquad f^{1/p^e} = \sum_{\boldsymbol{\lambda}} f_{\boldsymbol{\lambda}} \mathbf{x}^{\boldsymbol{\lambda}/p^e},$$

where $\boldsymbol{\lambda}$ runs over the tuples $(\lambda_1, \ldots, \lambda_d)$ with $0 \leqslant \lambda_j \leqslant p^e - 1$, and $f_{\boldsymbol{\lambda}} \in R$, then $(f)^{[1/p^e]} = T((f)^{1/p^e})$ is the ideal generated by the $f_{\boldsymbol{\lambda}}$. We can then compute $I^{[1/p^e]}$ for more general I by linearity: if $I = (f_1, \ldots, f_m)$, then $I^{[1/p^e]} = \sum (f_i)^{[1/p^e]}$.

2.1.1. *Complexity.* The computation of Frobenius roots is the workhorse behind many of the methods in *TestIdeals*. Hence it is important to understand how this is implemented, and its computational complexity.

The computation of Frobenius roots of ideals reduces to the case of principal ideals, and its complexity grows linearly with the number of generators of the ideal. Furthermore the calculation of $(f)^{[1/p^e]}$ reduces to finding the terms in the right hand side of (2), which essentially amounts to taking each term in f, computing the p^e-th root of its coefficient, and dividing the monomial exponent vector by p^e with remainder. Hence the complexity of computing $(f)^{[1/p^e]}$ is proportional to the number of terms in f, and is independent of its degree. We emphasize that the calculation of Frobenius roots *does not involve the calculation of Gröbner bases.*

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[2]In a reduced ring R, it is often convenient to express the Frobenius endomorphism $R \longrightarrow R$ as the inclusion $R \hookrightarrow R^{1/p}$, so that the source and target are distinguished.

2.2. **Dual to Frobenius on quotient rings and Frobenius on top local cohomology.** For *any* reduced ring R with finite e-iterated Frobenius map identified with $R \hookrightarrow R^{1/p^e}$, we always have a map

$$T_R : \omega_{R^{1/p^e}} \longrightarrow \omega_R.$$

At a maximal ideal $\mathfrak{m}$ of height d, this map is Matlis dual to the Frobenius map on $H^d_{\mathfrak{m}}(R)$. Hence we can study local cohomology by studying this map. If $R = S/I$, where S is a polynomial ring over a finite field, we can implement this map [Fed83, Kat08]. We briefly explain the case of a hypersurface here.

If $I = (f)$, then again $\omega_R \cong R = S/I$, and the action of T_R can be computed on S. In fact, if $\overline{J} \subseteq R \cong \omega_R$ has preimage $J \subseteq S$, then setting $u = f^{p^e-1}$, we have that $T_R(\overline{J}^{1/p^e})$ is the image of $(uJ)^{[1/p^e]}$ in R. More generally, for non-hypersurfaces, the analog to u is chosen to be an element that maps onto the generator of the cyclic R-module

$$\frac{(I^{[p^e]} : I) \cap (\Omega^{[p^e]} : \Omega)}{I^{[p^e]}},$$

where Ω is an ideal of S with $I \subseteq \Omega$ and whose image in R is isomorphic to the canonical module of R; see [Kat08] for details. Here we compute two examples.

```
i1 : S = ZZ/5[x,y,z];

i2 : f = x^3 + y^3 + z^3;

i3 : u = f^(5 - 1);

i4 : frobeniusRoot(1, ideal u)

o4 = ideal (z, y, x)

o4 : Ideal of S

i5 : S = ZZ/7[x,y,z];

i6 : f = x^3 + y^3 + z^3;

i7 : u = f^(7 - 1);

i8 : frobeniusRoot(1, ideal u)

o8 = ideal 1

o8 : Ideal of S
```

The above example has shown that the Frobenius map on top local cohomology of the cone over the Fermat elliptic curve is injective in characteristic 7 (the dual map $\omega_{R^{1/p}} \longrightarrow \omega_R$ is surjective) and not injective in characteristic 5 (the dual map $\omega_{R^{1/p}} \longrightarrow \omega_R$ is not surjective).

2.3. **A generalization of Frobenius powers and roots.** Let R be an F-finite regular ring. We can extend the definition of Frobenius powers as follows.

Definition 2.3 ([HTW18a]). Let $I \subseteq R$ be an ideal.

(a) If n is a positive integer with base p expansion $n = d_0 + d_1 p + \cdots + d_r p^r$, we define
$$I^{[n]} = I^{d_0}(I^{d_1})^{[p]} \cdots (I^{d_r})^{[p^r]}.$$

(b) If t is a nonnegative rational number of the form $t = a/p^e$, we define $I^{[t]} = \left(I^{[a]}\right)^{[1/p^e]}$.

(c) If t is any nonnegative rational number, and $\{a_n/p^{e_n}\}_{n \geqslant 1}$ is a sequence of rational numbers converging to t from above, we define $I^{[t]}$ to be the stable value of the non-decreasing chain of ideals $\{I^{[a_n/p^{e_n}]}\}_{n \geqslant 1}$.

```
i1 : p = 3;

i2 : R = ZZ/p[x,y];

i3 : m = monomialIdeal(x, y);

o3 : MonomialIdeal of R

i4 : I = m^5;

o4 : MonomialIdeal of R

i5 : t = 3/5 - 1/(5*p^3);

i6 : frobeniusPower(t, I)

            2       2
o6 = ideal (y , x*y, x )

o6 : Ideal of R

i7 : frobeniusPower(t - 1/p^5, I)

o7 = ideal (y, x)

o7 : Ideal of R
```

The generalized Frobenius powers of an ideal are test ideals (see Section 4) of sufficiently general linear combinations of generators of that ideal. Thus, the ideals computed above are test ideals of generic quintics in two variables in characteristic $p = 3$, and these computations suggest that $\frac{3}{5} - \frac{1}{5p^3}$ is a higher F-jumping exponent of such quintics, which is indeed the case for any $p \equiv 3 \mod 5$ (see [HTW18b]). To illustrate, compare the computations above with the following.

```
i8 : S = ZZ/3[a..f,x,y];

i9 : G = a*x^5 + b*x^4*y + c*x^3*y^2 + d*x^2*y^3 + e*x*y^4 + f*y^5;

i10 : testIdeal(t, G)

             2       2
o10 = ideal (y , x*y, x )

o10 : Ideal of S

i11 : testIdeal(t - 1/p^5, G)

o11 = ideal (y, x)

o11 : Ideal of S
```

Caveat. The computations of Frobenius powers with exponents whose denominators are not powers of the characteristic p are often very slow, and in several instances it is more efficient to introduce more variables and compute the test ideal of a "very general" linear combination of the generators of the ideal, as above. The code used for computations of Frobenius powers needs to undergo significant optimization.

2.4. **Frobenius roots of submodules of free modules.** Let S be a polynomial ring or power series ring. Given a submodule M of the free module S^k, there is a smallest submodule N of S^k that contains M, for which $M \subseteq N^{[p^e]}$. Here, $N^{[p^e]}$ is the submodule of S^k generated by the vectors of N with all coordinates raised to

the p^e-th power. (Cf. [KZ14]). Extending the previous definitions, we call this N the p^e-th Frobenius root of M, and denote it by $M^{[1/p^e]}$.

Example 2.4. Let $R = \mathbb{Z}/p\mathbb{Z}[a,b,c,d]$, and consider the ideals $\mathfrak{m} = (a,b,c,d)$ and

$$I = (a,b) \cap (a,c) \cap (c,d) \cap (c+d, a^3 + bd^2).$$

Then R/I is a 3-dimensional non-Cohen–Macaulay ring. Matlis Duality applied to $H^2_\mathfrak{m}(R/I)$ with its natural Frobenius map yields a p^{-1}-linear map U on $\operatorname{Ext}^2_R(R/I, R)$.

```
i1 : R = ZZ/2[a,b,c,d];

i2 : I = intersect(ideal(a,b),ideal(a,c),ideal(c,d),ideal(c+d,a^3+b*d^2));

o2 : Ideal of R

i3 : f = inducedMap(R^1/I, R^1/frobenius(I));

o3 : Matrix

i4 : E2 = Ext^2(f, R^1)

o4 = {-8}  | a4+abc2+abcd a2b               |
     {-10} | a2cd3        a3cd+a3d2+bcd3 |

o4 : Matrix

i5 : target E2

o5 = cokernel {-8}  | 0     a2 b2c2 |
              {-10} | c2+d2 d4 a4d2 |

                                    2
o5 : R-module, quotient of R

i6 : source E2

o6 = cokernel {-4} | 0   a  bc  |
              {-5} | c+d d2 a2d |

                                    2
o6 : R-module, quotient of R
```

The Frobenius map on $H^2_\mathfrak{m}(R/I)$ is injective if and only the image of $(\operatorname{Image} U)^{[1/p]}$ in $\operatorname{Ext}^2_R(R/I, R)$ as computed above is the whole of R^2.

```
i7 : U = matrix entries E2;

            2       2
o7 : Matrix R  <--- R

i8 : A = image matrix entries relations source E2;

i9 : frobeniusRoot(1, image U)

o9 = image {-2} | 1 0 0 |
           {0}  | 0 d a |

                                    2
o9 : R-module, submodule of R
```

The calculation above shows that R/I is not F-injective (cf. Section 3). Moreover, it shows that $H^2_\mathfrak{m}(R/I) \cong \operatorname{Ann}_{E^2} A^t$, where E is the injective hull of $R_\mathfrak{m}/\mathfrak{m}R_\mathfrak{m}$, and the Frobenius map on $H^2_\mathfrak{m}(R/I)$ induced from Frobenius on R is given by $U^t\Theta$, where Θ is the induced Frobenius on E. The submodule of nilpotent elements in $H^2_\mathfrak{m}(R/I)$ is given by $\operatorname{Ann}_{E^2} B^t$, where B is the smallest submodule of R^2 containing $\operatorname{Image} A + \operatorname{Image} U$ such that $UB \subseteq B^{[p]}$. The method `ascendModule` can be used to calculate B (see a detailed description of the similar method `ascendIdeal` in Section 4).

```
i10 : B = ascendModule(1, A, U)

o10 = image | 0   a  bc  |
            | c+d d2 a2d |

                                    2
o10 : R-module, submodule of R
```

3. F-SINGULARITIES

The *TestIdeals* package includes methods for determining if a ring is F-injective, F-pure, F-rational, or F-regular.

3.1. F-injectivity.

Definition 3.1. A local ring $(R, \mathfrak{m})$ is called *F-injective* if the map $H^i_{\mathfrak{m}}(R) \to H^i_{\mathfrak{m}}(R^{1/p})$ is injective for all $i > 0$. An arbitrary ring is called *F-injective* if its localization at each prime ideal of R is F-injective.

The function `isFInjective` determines whether the ring $R = S/I$ is F-injective, where S is a polynomial ring.

```
i1 : R = ZZ/7[x,y,z]/(x^3 + y^3 + z^3);

i2 : isFInjective R

o2 = true

i3 : R = ZZ/5[x,y,z]/(x^3 + y^3 + z^3);

i4 : isFInjective R

o4 = false
```

Equivalently, a ring is F-injective if the maps on the cohomology of the dualizing complex

$$h^{-i}\omega^{\bullet}_{R^{1/p}} \longrightarrow h^{-i}\omega^{\bullet}_{R}$$

surject for all i. Note that $h^{-i}\omega^{\bullet}_{R} \cong \operatorname{Ext}^{\dim S-i}(R, S)$, the latter of which *Macaulay2* readily computes. The algorithm used by `isFInjective` works by checking the surjectivity of the dual Frobenius map

$$\operatorname{Ext}^{\dim S-i}(R^{1/p}, S) \longrightarrow \operatorname{Ext}^{\dim S-i}(R, S).$$

We begin by computing the map $R \to R^{1/p}$ using the *PushForward* package [Rai]. Next the algorithm computes $\operatorname{Ext}^i(\_\_, S)$ applied to the map from the previous step. Then R is F-injective precisely when the cokernel of $\operatorname{Ext}^i(\_\_, S)$ is trivial for i.

The Frobenius action on top local cohomology (dual to ω_R) is usually computed in a different (faster) way than the other cohomologies, and this is modified by the `CanonicalStrategy` option. The default value for this option is `Katzman`, which instead of using the *PushForward* package, relies on the fact that we already know how to compute the Frobenius action on the canonical module, as described in Section 2.2.

The performance of the algorithm can be improved if the ring of interest is nice enough. If the ring is Cohen–Macaulay, then setting `AssumeCM => true` (the default value is `false`) lets the algorithm check the Frobenius action only on top cohomology (which is typically much faster, as explained above). When studying

a reduced ring, setting `AssumeReduced => true` (the default value) avoids computing the bottom local cohomology, and when studying a normal ring, setting `AssumeNormal => true` (the default is `false`) avoids computing the bottom two local cohomologies.

By default the algorithm checks for F-injectivity at all points of $\operatorname{Spec} R$. However, one can choose to check F-injectivity only at the origin by setting the option `AtOrigin` to `true`.

```
i1 : R = ZZ/7[x,y,z]/((x - 1)^5 + (y + 1)^5 + z^5);

i2 : isFInjective R -- R is not globally F-injective...

o2 = false

i3 : isFInjective(R, AtOrigin => true) -- but is F-injective at the origin

o3 = true
```

3.2. F-regularity.

Definition 3.2 ([HH90, HW02]). A ring R is called *strongly F-regular* if the (big) test ideal $\tau(R)$ is equal to R. Likewise a pair (R, f^t) is called *strongly F-regular* if $\tau(R, f^t) = R$.

The command `isFRegular` checks whether a ring or pair is strongly F-regular. Below are two examples, one of which is F-regular and the other one is not.

```
i1 : R = ZZ/5[x,y,z]/(x^2 + y*z);

i2 : isFRegular R

o2 = true

i3 : R = ZZ/7[x,y,z]/(x^3 + y^3 + z^3);

i4 : isFRegular R

o4 = false
```

We can also check whether a pair (R, f^t) is F-regular.

```
i1 : R = ZZ/5[x,y];

i2 : f = y^2 - x^3;

i3 : isFRegular(1/2, f)

o3 = true

i4 : isFRegular(5/6, f)

o4 = false

i5 : isFRegular(4/5, f)

o5 = false

i6 : isFRegular(4/5 - 1/100000, f)

o6 = true
```

All of these checks are done by actually computing the test ideal, as described in Section 4.

If the input ring is $\mathbb{Q}$-Gorenstein, then in each of the cases above, the output is a boolean indicating if the ring is strongly F-regular. If the input ring is not $\mathbb{Q}$-Gorenstein, then the algorithm can be used to determine if a ring is strongly

F-regular, but cannot prove that a ring is not strongly F-regular (this latter functionality can, however, be enabled by setting `QGorensteinIndex => infinity`).

In the case that R is $\mathbb{Q}$-Gorenstein, the algorithm works by computing the test ideal τ of the ring (or the pair) using `testIdeal` and checking whether $\tau = R$. In the non-$\mathbb{Q}$-Gorenstein case, the algorithm checks for strong F-regularity by computing better and better approximations of the test ideal, and checking whether any of these is the unit ideal. To compute approximations of the test ideal, the algorithm computes a test element c with `testElement` and then uses `frobeniusRoot` to compute the e-th root of $c(I^{[p^e]} : I)$; appropriate modifications are made for pairs. If at any step the approximation is the unit ideal, then then the algorithm returns `true`. Otherwise the algorithm continues checking for each e until a specified limit is reached. The default limit is 2, and can be changed using the option `DepthOfSearch`.

A number of options can be used to speed up the performance of some of the internal functions. The option `AssumeDomain` can be used if R is an integral domain, `FrobeniusRootStrategy` chooses a strategy for internal `frobeniusRoot` calls, `MaxCartierIndex` sets the maximum Gorenstein index to search for when working with a $\mathbb{Q}$-Gorenstein ambient ring, and `QGorensteinIndex` allows the user to specify the $\mathbb{Q}$-Gorenstein index of the ring.

The default behavior of `isFRegular` is that it checks for strong F-regularity globally. If the option `AtOrigin` is set to `true`, the algorithm will only check F-regularity at the origin, by checking whether the computed test ideal is in the irrelevant ideal. Below are examples for both a ring and a pair.

```
i1 : R = ZZ/7[x,y,z]/((x - 1)^3 + (y + 1)^3 + z^3);

i2 : isFRegular R -- R is not globally F-regular...

o2 = false

i3 : isFRegular(R, AtOrigin => true) -- but is F-regular at the origin

o3 = true

i4 : R = ZZ/13[x,y];

i5 : f = (y - 2)^2 - (x - 3)^3;

i6 : isFRegular(5/6, f) -- (R,f^(5/6)) is not F-regular...

o6 = false

i7 : isFRegular(5/6, f, AtOrigin => true) -- but is F-regular at the origin

o7 = true
```

3.3. F-purity.

Definition 3.3. A ring R is called F-*pure* if the inclusion $R \hookrightarrow R^{1/p^e}$ is a pure map, i.e., the tensor of this map with any R-module remains injective. If $R^{1/p}$ is a finite R-module, this is equivalent to requiring that the inclusion $R \hookrightarrow R^{1/p}$ split as a map of R-modules.

The function `isFPure` checks whether a ring is F-pure. Either a ring or a defining ideal can be input, as seen in the following example.

```
i1 : R = ZZ/5[x,y,z]/(x^2 + y*z);

i2 : isFPure R

o2 = true
```

```
i3 : R = ZZ/7[x,y,z]/(x^3 + y^3 + z^3);

i4 : isFPure R

o4 = true

i5 : S = ZZ/2[x,y,z];

i5 : isFPure ideal(y^2 - x^3)

o5 = false

i6 : isFPure ideal(z^2 - x*y*z + x*y^2 + x^2*y)

o6 = true
```

The algorithm works by applying Fedder's Criterion [Fed83], which states that a local ring $(R, \mathfrak{m})$ is F-pure if and only if $(I^{[p]} : I) \not\subseteq \mathfrak{m}^{[p]}$. When `AtOrigin` is set to `true`, the algorithm checks F-purity only at the origin, by explicitly checking the above containment. When `AtOrigin` is set to `false`, which is its default value, the algorithm computes the non F-pure locus, by applying `frobeniusRoot` to $I^{[p]} : I$. If the non F-pure locus is the whole ring, the algorithm returns `true`.

3.4. F-rationality.

Definition 3.4. Suppose that R is a Cohen–Macaulay ring and that $T_R : \omega_{R^{1/p}} \to \omega_R$ is the canonical dual to the Frobenius. We say that R has F-rational singularities, of simply that R is F-rational, if there are no non-zero proper submodules M of ω_R such that $T_R(M^{1/p}) \subseteq M$.

The command `isFRational` checks if a ring is F-rational.

```
i1 : S = ZZ/3[a,b,c,d,t];

i2 : M = matrix{{ a^2 + t^4, b, d }, { c, a^2, b^3 - d }};

            2        3
o2 : Matrix S  <--- S

i3 : I = minors(2, M);

o3 : Ideal of S

i4 : R = S/I;

i5 : isFRational R

o5 = true
```

The algorithm used by `isFRational` first checks if the ring is Cohen–Macaulay, unless the option `AssumeCM` is set to `true`. If the ring is not Cohen–Macaulay then `false` is returned. Next, the algorithm computes the test module $M \subseteq \omega_R$ and checks to see if $\omega_R \subseteq M$; see the next section for the description of a test module.

The options `AssumeDomain` and `FrobeniusRootStrategy` can be used to improve the speed of the `testModule` computation. By default, these options are set to `false` and `Substitution`, respectively. Finally, if `AtOrigin` is set to `true`, then F-rationality is checked only at the origin.

4. TEST IDEALS

In this section, we explain how to compute parameter test modules, parameter test ideals, test ideals, and HSLG modules[3].

––––––––––––

[3]HSLG modules can be used to give a scheme structure to the F-injective or F-pure locus.

4.1. **Parameter test modules.** Given a reduced ring R of finite type over a perfect field k, the Frobenius map $R \hookrightarrow R^{1/p^e}$ is dual to $T : \omega_{R^{1/p^e}} \longrightarrow \omega_R$. As in Section 2.2, we can represent the canonical module $\omega_R \subseteq R$ as an ideal, we can write $R = S/I$, and so we can find an element $u \in S^{1/p^e}$ representing the map $T : \omega_{R^{1/p^e}} \longrightarrow \omega_R$; see Section 2.2 or [Kat08].

Definition 4.1. The *parameter test submodule* is the smallest submodule $M \subseteq \omega_R$ (and hence ideal of R, since $M \subseteq \omega_R \subseteq R$) that agrees generically with ω_R and that satisfies

$$T(M^{1/p^e}) \subseteq M$$

for some $e > 0$ (or equivalently for all $e > 0$).

Using *Macaulay2*, we can compute this using the `testModule` command as follows.

```
i1 : R = ZZ/5[x,y,z]/(x^4 + y^4 + z^4);

i2 : N = testModule R;

i3 : N#0
            2          2       2
o3 = ideal (z , y*z, x*z, y , x*y, x )

o3 : Ideal of R

i4 : N#1

o4 = ideal 1

o4 : Ideal of R
```

The output of `testModule` is a sequence, consisting of three items:

(1) The test module itself, given as an ideal of R.
(2) The canonical module that contains the test module, given as an ideal of R. (Note the representation of the canonical module as an ideal is not unique, it is only unique up to isomorphism, hence it is important to return this module as well).
(3) The element u described above (not displayed above, since it takes a lot of space).

Note since this ring is Gorenstein, the canonical module is simply represented as the unit ideal. Here is another example, where the ring is not Gorenstein.

```
i1 : R = ZZ/5[x,y,z]/(y*z, x*z, x*y);

i2 : N = testModule R;

i3 : N#0
            2   2   2
o3 = ideal (z , y , x )

o3 : Ideal of R

i4 : N#1

o4 = ideal (y + 2z, x - z)

o4 : Ideal of R
```

We briefly explain how this is computed: First, we find a *test element*.

Remark 4.2 (Computation of test elements). We recall that, roughly, an element of the Jacobian ideal that is not contained in any minimal prime is a test element [Hoc07]. We compute test elements by computing random partial derivatives (and linear combinations thereof) until we find an element that does not vanish at all the minimal primes. This method is much faster than computing the entire Jacobian ideal. If it is known that the ring is a domain, setting `AssumeDomain => true` can speed this up further.

After the test element c has been identified, we pull back the ideal ω_R to an ideal $J \subseteq S$. Next, we compute the following ascending sequence of ideals where u represents $T : \omega_{R^{1/p^e}} \longrightarrow \omega_R$ as above:

$$
\begin{aligned}
J_0 &:= cJ \\
J_1 &:= J_0 + (uJ_0)^{[1/p]} \\
(3) \quad J_2 &:= J_1 + (uJ_1)^{[1/p]} = J_0 + (uJ_0)^{[1/p]} + (u^{1+p}J_0)^{[1/p^2]} \\
J_3 &:= J_2 + (uJ_2)^{[1/p]} = J_0 + (uJ_0)^{[1/p]} + (u^{1+p}J_0)^{[1/p^2]} + (u^{1+p+p^2}J_0)^{[1/p^3]} \\
&\vdots
\end{aligned}
$$

As soon as this ascending sequence of ideals stabilizes, we are done. In fact, because this strategy is used in several contexts, the user can call it directly for a chosen ideal J and u with the function `ascendIdeal` (this is done for test ideals below).

We can also compute parameter test modules of pairs (ω_R, f^t) with $t \in \mathbb{Q}_{\geqslant 0}$. This is done by modifying the element u when the denominator of t is not divisible by p. When t has p in its denominator, we rely on the fact (see [BMS08, ST14a]) that

$$
\begin{aligned}
T(\tau(\omega_R, f^a)) &= \tau(\omega_R, f^{a/p}) \\
\texttt{frobeniusRoot(1, u * I}_1) &= \texttt{I}_2
\end{aligned}
$$

where the second line roughly explains how this is accomplished internally. Here `I`$_1$ is $\tau(\omega_R, f^a)$ pulled back to S and, likewise, I_2 defines $\tau(\omega_R, f^{a/p})$ modulo the defining ideal of R.

Remark 4.3 (Optimizations in `ascendIdeal` and other `testModule` computations). Throughout the computations described above, we very frequently use the following fact:

$$(f^p \cdot J)^{[1/p]} = f \cdot (J^{[1/p]}).$$

To access this enhancement, one should try to pass functions like `ascendIdeal` and `frobeniusRoot` the elements and their exponents (see the documentation). In particular, when computing the p^e-th Frobenius root of an ideal of the form $f^n \cdot J$, one should not multiply out f^n and J, nor even raise f to the nth power, but rather call `frobeniusRoot(e, n, f, J)`.

4.2. **Parameter test ideals.** The parameter test ideal is simply the annihilator of $\omega_R/\tau(\omega_R)$. In other words, it is

$$(\tau(\omega_R) : \omega_R).$$

This can also be described as

$$\bigcap_I (I^* : I),$$

where I varies over ideals defined by a partial system of parameters, and I^* denotes its tight closure. This latter description is not computable, however.

Example 4.4. In this example we repeat the calculation in [Kat08, §9].

```
i1 : R = ZZ/2[a..e];

i2 : E = matrix {{a, b, b, e}, {d, d, c, a}};

              2      4
o2 : Matrix R  <--- R

i3 : I = minors(2, E);

o3 : Ideal of R

i4 : S = R/I;

i5 : J = parameterTestIdeal S

o5 = ideal (c + d, b, a)

o5 : Ideal of S

i6 : J = substitute(J, R);

o6 : Ideal of R

i7 : mingens(J + I)

o7 = | c+d b a de |

              1      4
o7 : Matrix R  <--- R
```

4.3. **Test ideals.** For an F-finite reduced ring $R = S/I$, where S is a regular ring, the (big) test ideal[4] of R is the smallest ideal J in R, not contained in any minimal prime, such that for all $e > 0$ and all $\phi \in \mathrm{Hom}_R(R^{1/p^e}, R)$ we have

$$\phi(J^{1/p^e}) \subseteq J.$$

In the case that R is Gorenstein, $\mathrm{Hom}_R(R^{1/p^e}, R)$ is a cyclic R^{1/p^e}-module generated by Φ_e, which corresponds with the map T above based on the identification $\omega_R \cong R$ [BS13]. More generally, if R is $\mathbb{Q}$-Gorenstein with index not divisible by p, then for at least sufficiently divisible $e > 0$, such a generating Φ_e still exists.

If such a Φ_e exists, it can be identified to a generator of the module $(I^{[p^e]} : I)/I^{[p^e]}$ by Fedder's Lemma [Fed83]. Hence we can find a corresponding[5] $u \in I^{[p^e]} : I$. In this case, if $c \in S$ is the pre-image of a test element of R, then setting $I_0 = cR$, it follows that

$$\tau(R) = \mathtt{ascendIdeal}(\mathtt{e}, \mathtt{u}, \mathtt{I_0}),$$

where $\mathtt{ascendIdeal}$ is the method explained above, in (3). The e here means all Frobenius roots are taken as multiples of e. In other words, we first compute $\mathtt{I_0} + (\mathtt{u} \cdot \mathtt{I_0})^{1/\mathtt{p}^{[e]}}$. Then we compute

$$\mathtt{I_0} + (\mathtt{u} \cdot \mathtt{I_0})^{1/\mathtt{p}^{[e]}} + (\mathtt{u}^{\mathtt{p}^e + 1} \cdot \mathtt{I_0})^{1/\mathtt{p}^{[2e]}}$$

etc.

Here is an example (a $\mathbb{Z}/3\mathbb{Z}$-quotient, where $3|(7-1)$), where exactly this logic occurs.

---

[4]The notion of test ideals was originally introduced in [HH90] in the context of tight closure in finitely generated modules, whereas our notion of test ideals arises from tight closure in possibly non finitely generated ("big") modules. Confusingly, big test ideals are included in (small) test ideals.

[5]This is done via the function $\mathtt{QGorensteinGenerator}$.

```
i1 : T = ZZ/7[x,y];

i2 : S = ZZ/7[a,b,c,d];

i3 : f = map(T, S, {x^3, x^2*y, x*y^2, y^3});

i4 : I = ker f;

i5 : R = S/I;

i6 : testIdeal R

o6 = ideal 1

o6 : Ideal of R
```

However, the term u constructed above can be quite complicated if $e > 1$ (which happens exactly when $(p-1)K_R$ is not Cartier). For instance, even in the above example we have an extremely complex u:

```
i7 : toString QGorensteinGenerator(1, R)

o7 = a^2*b^6*c^12+a^3*b^4*c^13+a^3*b^5*c^11*d+a^4*b^3*c^12*d+a^5*b*c^13*d+b^
     12*c^6*d^2+a^3*b^6*c^9*d^2+a^4*b^4*c^10*d^2+a^5*b^2*c^11*d^2+a^6*c^12*d
     ^2+b^13*c^4*d^3+a*b^11*c^5*d^3+a^2*b^9*c^6*d^3+a^4*b^5*c^8*d^3+a^5*b^3*
     c^9*d^3+a^6*b*c^10*d^3+a*b^12*c^3*d^4+a^2*b^10*c^4*d^4+a^3*b^8*c^5*d^4+
     a^4*b^6*c^6*d^4+a^5*b^4*c^7*d^4+a^6*b^2*c^8*d^4+a^7*c^9*d^4+a*b^13*c*d^
     5+a^2*b^11*c^2*d^5+a^3*b^9*c^3*d^5+a^4*b^7*c^4*d^5+a^5*b^5*c^5*d^5+a^6*
     b^3*c^6*d^5+a^7*b*c^7*d^5+a^2*b^12*d^6+a^3*b^10*c*d^6+a^4*b^8*c^2*d^6+a
     ^5*b^6*c^3*d^6+a^6*b^4*c^4*d^6+a^7*b^2*c^5*d^6+a^8*c^6*d^6+a^4*b^9*d^7+
     a^5*b^7*c*d^7+a^6*b^5*c^2*d^7+a^7*b^3*c^3*d^7+a^8*b*c^4*d^7+a^6*b^6*d^8
     +a^7*b^4*c*d^8+a^8*b^2*c^2*d^8+a^9*c^3*d^8+a^8*b^3*d^9+a^9*b*c*d^9+a^10
     *d^10
```

Therefore, we use a different strategy if either $(p-1)K_R$ is not Cartier or, more generally, if R is $\mathbb{Q}$-Gorenstein of index divisible by p. In these situations, this alternate strategy typically appears to be faster. We rely on the observation (see [BST15]) that
$$\tau(\omega_R, K_R) \cong \tau(R).$$
In fact, by embedding $\omega_R \subseteq R$, we can compute g so that $\tau(\omega_R, K_R) = g \cdot \tau(R)$. We can therefore find $\tau(R)$ if we can find $\tau(\omega_R, K_R)$. Next, if K_R is $\mathbb{Q}$-Cartier with $nK_R = \operatorname{div}_R(f)$ for some $f \in R$ and $n > 0$, then
$$\tau(\omega_R, K_R) = \tau(\omega_R, f^{1/n}).$$
Thus we directly compute $\tau(\omega_R, f^{1/n})$ via the command `testModule(1/n, f)`. Consider the following example, a μ_3-quotient, which uses the logic described above.

```
i1 : T = ZZ/3[x,y];

i2 : S = ZZ/3[a,b,c,d];

i3 : f = map(T, S, {x^3, x^2*y, x*y^2, y^3});

o3 : RingMap T <--- S

i4 : I = ker f;

o4 : Ideal of S

i5 : R = S/I;

i6 : testIdeal R

o7 = ideal 1

o7 : Ideal of R
```

Remark 4.5 (Non-graded caveats). It frequently happens that $(I^{[p^e]} : I)/I^{[p^e]}$ is principal but *Macaulay2* cannot identify the principal generator (since *Macaulay2* cannot always find minimal generators of non-graded ideals or modules). The same thing can happen when computing the element u corresponding to the map $T :$ $\omega_{R^{1/p}} \longrightarrow \omega_R$, as described in Section 2.2. In such situations, instead of a single u, *Macaulay2* will produce $u_1, \ldots, u_n$ (all multiples of u, and u is a linear combination of the u_i). Instead of computing the ideal

$$(u \cdot J)^{[1/p]},$$

we compute

$$(u_1 \cdot J)^{[1/p]} + \cdots + (u_n \cdot J)^{[1/p]},$$

which will produce the same answer.

We can similarly use the `testIdeal` command to compute test ideals of pairs, $\tau(R, f^t)$, and even mixed test ideals, $\tau(R, f_1^{t_1} \cdots f_n^{t_n})$.

4.4. **HSLG module; computing F-pure submodules of rank-1 Cartier modules.** Again, consider the maps $T^e : \omega_{R^{1/p^e}} \longrightarrow \omega_R$ we have considered throughout this section. It is a theorem of Hartshorne-Speiser, Lyubeznik, and Gabber [HS77, Lyu97, Gab04] that the descending images

$$(4) \qquad \omega_R \supseteq T(\omega_{R^{1/p}}) \supseteq \cdots \supseteq T^e(\omega_{R^{1/p^e}}) \supseteq T^{e+1}(\omega_{R^{1/p^{e+1}}}) \supseteq \cdots$$

stabilize for $e \gg 0$. The function `FPureModule` computes the stable submodule in this chain, called the *HSLG module*, and returns a sequence containing the following items:

(1) The HSLG module;
(2) The canonical module, embedded (non-uniquely) as an ideal;
(3) The element u representing the map on the canonical module (see Section 2.2);
(4) The value of $e > 0$ at which the descending images in (4) stabilize. This is sometimes also called the HSLG number of the canonical module as a Cartier module.

```
i1 : R = ZZ/3[x,y,z]/(x^3 + y^4 + z^5);

i2 : L = FPureModule R;

i3 : L#0

                    2       2   3
o3 = ideal (y*z, x*z, y , x*y, x , z )

o3 : Ideal of R

i4 : L#1

o4 = ideal 1

o4 : Ideal of R

i5 : L#3

o5 = 1

i6 : R = ZZ/3[x,y,z]/(x^3 + y^4 + z^5);

i7 : L = FPureModule R;

i8 : L#0

                    2       2   3
o8 = ideal (y*z, x*z, y , x*y, x , z )
```

```
o8 : Ideal of R

i9 : L#1

o9 = ideal 1

o9 : Ideal of R

i10 : L#3

o10 = 1
```

More generally, for any ideal J with a map $\phi : J^{1/p^e} \longrightarrow J$, we have that the images

$$J \supseteq \phi(J^{1/p^e}) \supseteq \phi^2(J^{1/p^{2e}}) \supseteq \cdots$$

stabilize as well (in fact, the analogous result even holds for modules [Gab04]).

4.5. **Computing the level of a polynomial.** Another interesting invariant that can be calculated using this package is the so–called *level* of a polynomial; more precisely:

Definition 4.6. Let K be an F–finite field of prime characteristic p, and let $f \in R = K[x_1, \ldots, x_d]$. We define the *level of* f as the smallest possible integer e where the descending chain

$$R = (f^{p^0-1})^{[1/p^0]} \supseteq (f^{p-1})^{[1/p]} \supseteq (f^{p^2-1})^{[1/p^2]} \supseteq \cdots \supseteq (f^{p^i-1})^{[1/p^i]} \supseteq \cdots$$

stabilizes.

This invariant was introduced in [ÀMBL05]. It is also essentially the same data as the HSLG number of the pair (R, f^1) as computed in (4) above (it is that number $+1$). One interesting particular case is when $K = \mathbb{F}_p$ and f is the defining homogeneous polynomial of a hyperelliptic curve of genus g; when $g = 1$, it was proved in [BDSV15] that the corresponding elliptic curve defined by f is ordinary if and only if its level is 1, (equivalently, if and only if $(R/(f))_{(x,y,z)}$ is F–pure,) and supersingular if and only if its level is 2. When the genus is at least 2, level 2 is a necessary (but not sufficient) condition for the curve for being ordinary [BCBFY18]; in this case, one also has that, if the curve is supersingular, then its level has to be at least 3, so the level can always distinguish these two properties in any genus. We illustrate these results by means of the following examples.

```
i1 : R = ZZ/2[x,y,z];

i2 : f = x^3 + y^2*z + y*z^2;

i3 : frobeniusPower(1/2, ideal f)

o3 = ideal (z, y, x)

o3 : Ideal of R

i4 : u = f^3;

i5 : frobeniusPower(1/4, ideal u)

o5 = ideal (z, y, x)

o5 : Ideal of R

i6 : ((FPureModule(1, f))#3) + 1

o6 = 2
```

This shows that this elliptic curve has level 2, hence it is supersingular. The next example shows a non-ordinary hyperelliptic curve of genus 2 with level 2.

```
i1 : R = ZZ/11[x,y,z];

i2 : f = y^2*z^3 - x^5 - 2*z^5;

i3 : frobeniusPower(1/11, ideal f^10)

             2        3
o3 = ideal (z , x*z, x )

o3 : Ideal of R

i4 : frobeniusPower(1/121, ideal f^120)

             2        3
o4 = ideal (z , x*z, x )

o4 : Ideal of R

i5 : ((FPureModule(1, f))#3) + 1

o5 = 2
```

5. Ideals compatible with a given p^{-e}-linear map

Throughout this section, let R denote a polynomial $\mathbb{K}[x_1, \ldots, x_n]$. In this section we address the following question: given a R-linear map $\phi : R^{1/p^e} \to R$, what are the ideals $I \subseteq R$ such that $\phi(I^{1/p^e}) \subseteq I$? We call these ideals ϕ-compatible.

Recall that $\mathrm{Hom}_R(R^{1/p^e}, R)$ is a principal R^{1/p^e}-module generated the *trace map* $T \in \mathrm{Hom}_R(R^{1/p^e}, R)$, constructed as in Section 2.1, (cf. [Fed83, Lemma 1.6] and [BK05, Example 1.3.1]).

We can now write our given ϕ as multiplication by some u^{1/p^e} followed by T and it is not hard to see that an ideal $I \subseteq R$ is ϕ-compatible if and only if $uI \subseteq I^{[p^e]}$.

Theorem 5.1. *If ϕ is surjective, there are finitely many ϕ-compatible ideals, consisting of all possible intersections of ϕ-compatible prime ideals (cf. [KM09], [Sch09], [Sha07], [EH08]). In general, there are finitely many ϕ-compatible prime ideals not containing $\sqrt{\mathrm{Image}\,\phi}$ (cf. [KS12]).*

The method `compatibleIdeals` produces the finite set of ϕ-compatible prime ideals in the second claim of Theorem 5.1.

```
i1 : R = ZZ/3[u,v];

i2 : u = u^2*v^2;

i3 : compatibleIdeals u

o3 = {ideal v, ideal (v, u), ideal u}

o3 : List
```

The defining condition $uI \subseteq I^{[p^e]}$ for ϕ-compatible ideals allows us to think of these in a dual form: write $\mathfrak{m} = (x_1, \ldots, x_n)$, $E = E_{R_\mathfrak{m}}(R_\mathfrak{m}/\mathfrak{m}) = H^n_\mathfrak{m}(R)$, and let $\Theta : E \to E$ be the p^e-linear map[6] induced from the Frobenius map on R. If $\psi = u\Theta$, then $\psi\,\mathrm{Ann}_E I \subseteq \mathrm{Ann}_E I$ if and only if $uI \subseteq I^{[p^e]}$. Thus finding all R-submodules of E compatible with $\psi = u\Theta$ also amounts to finding all ϕ-compatible ideals, where $\phi = T \circ u^{1/p^e}$.

---

[6]That is, Θ is additivive and $\Theta(ra) = r^{p^e}\Theta(a)$ for all $a \in E$ and $r \in R$.

Example 5.2. We return to Example 4.4. In (cf. [Kat08, §9]) it is shown that there is a surjection $\operatorname{Ann}_E I \to H^2_{\mathfrak{m}}(R/I)$ which is compatible with the induced p^1-linear map on $H^2_{\mathfrak{m}}(R/I)$ and the p^1-linear map $u\Theta$ on $\operatorname{Ann}_E I$, where u is computed as follows.

```
i1 : R = ZZ/2[a..f];

i2 : E = matrix {{a, b, b, e}, {d, d, c, a}};

              2        4
o2 : Matrix R   <--- R

i3 : I = minors(2, E);

o3 : Ideal of R

i4 : S = R/I;

i5 : isCohenMacaulay S

o5 = true
```

In [Kat08] it is shown that as R/I is Cohen–Macaulay, u can be taken as the generator of the cyclic module $(I^{[p]} : I) \cap (\Omega^{[p]} : \Omega)/I^{[p]}$ where $\Omega \subseteq R$ is an ideal whose image in R/I is a canonical module of that ring.

```
i6 : omega = canonicalIdeal S

o6 = ideal (e, d, a)

o6 : Ideal of S

i7 : omega = substitute(omega, R) + I;

o7 : Ideal of R

i8 : u = intersect((frobenius I):I, (frobenius omega) : omega);

o8 : Ideal of R

i9 : u = compress((gens u) % gens(frobenius I))

o9 = | a3bc+a3bd+a2cde+abcde+abd2e+b2d2e+cd2e2+d3e2 |

              1        1
o9 : Matrix R   <--- R

i10 :  u = first first entries u

         3        3       2                        2      2 2        2 2     3 2
o10 = a b*c + a b*d + a c*d*e + a*b*c*d*e + a*b*d e + b d e + c*d e   + d e

o10 : R
```

Now we can compute all annihilators of R-submodules of E stable under the p^1-linear map uT.

```
i11 : L = compatibleIdeals u;

i12 : print \ L;
            3        3       2                        2      2 2        2 2     3 2
ideal( a b*c + a b*d + a c*d*e + a*b*c*d*e + a*b*d e + b d e + c*d e   + d e )

                  2
ideal (a + b, a   + d*e)

ideal (e, a, d)

ideal (e, d, a, c)

ideal (e, d, c, b, a)

ideal (e, d, b, a)
```

```
ideal (e, b, a)

ideal (e, b, a, c + d)

                            2
ideal (c + d, a + b, b  + d*e)

ideal (d, a)

ideal (d, a, c)

ideal (d, c, b, a)

ideal (d, b, a)
```

We can also compute all annihilators of R-submodules of $H^2_{\mathfrak{m}}(R/I)$ stable under the p^1-linear map induced from the Frobenius map on R/I.

```
i13 : print \ unique apply(L, J -> J : omega);
          3       3     2                      2     2 2     2 2    3 2
ideal(a b*c + a b*d + a c*d*e + a*b*c*d*e + a*b*d e + b d e + c*d e  + d e )

                    2
ideal (a + b, b  + d*e)

ideal (e, d, a)

ideal 1

ideal (e, b, a)

ideal (e, c + d, b, a)

                            2
ideal (c + d, a + b, b  + d*e)

ideal (d, a)

ideal (d, c, a)

ideal (d, c, b, a)

ideal (d, b, a)
```

6. FUTURE PLANS

In [KZ14], the algorithms behind the methods in Section 5 were extended to compute prime annihilators of submodules of Artinian modules compatible with a given p^e-linear map. This would require, among other things, a faster implementation of our method for finding Frobenius roots of submodules of free modules.

On the other hand, it should be possible to compute test ideals of pairs $(R, \mathfrak{a}^t)$ even when $\mathfrak{a}$ is not principal. One strategy to do this is outlined in [ST14b] although certain improvement can be made.

We hope to achieve all these things during a future *Macaulay2* workshop. We also want to bring to the reader's attention the package *FThresholds*, which computes F-pure thresholds, F-thresholds, F-jumping numbers and more!

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