

Capacities and Optimal Input Distributions for Particle-Intensity Channels

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Abstract—This work introduces the particle-intensity channel (PIC) as a new model for molecular communication systems that includes imperfections at both transmitter and receiver and provides a new characterization of the capacity limits as well as properties of the optimal (capacity-achieving) input distributions for such channels. In the PIC, the transmitter encodes information, in symbols of a given duration, based on the probability of particle release, and the receiver detects and decodes the message based on the number of particles detected during the symbol interval. In this channel, the transmitter may be unable to control precisely the probability of particle release, and the receiver may not detect all the particles that arrive. We model this channel using a generalization of the binomial channel and show that the capacity-achieving input distribution for this channel always has mass points at probabilities of particle release of zero and one. To find the capacity-achieving input distributions, we develop a novel and efficient algorithm we call dynamic assignment Blahut-Arimoto (DAB). For diffusive particle transport, we also derive the conditions under which the input with two mass points is capacity-achieving.

Index Terms—Molecular communication, particle intensity channel, channel models, channel capacity, optimal input, optimization.

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I. INTRODUCTION

IN MOLECULAR communication (MC) transmitters convey information by releasing small particles, typically from nanometers to tens of micrometers in diameter. Information may be contained in the number or type of released particles or in the time of release [3], [4]. These particles travel to the receiver where they are detected and the message decoded. The stochastic nature of the transport process introduces uncertainty about the time of particle release and even the number of particles released during a given symbol interval.

One approach to understanding the capacity limits of molecular channels investigated in prior work assumes information is encoded in the time instants at which particle(s) are released. Such channels are called molecular timing channels (MTCs). In particular, the additive inverse Gaussian noise channel is presented in [5], [6], and upper and lower bounds on capacity are derived. These works assume a system where information is encoded in the release time of a *single* particle. Molecular timing channels where information is encoded via the release times of *multiple* particles are considered in [7], which presents upper and lower bounds on capacity, and [8] introduces a MTC where particles decay after a finite interval and derives upper and lower bounds on the associated capacity.

Another approach to MC encodes information through the number of particles released at the transmitter and decodes based on the number of particles that arrive at the receiver during the symbol interval. We focus on this type of modulation scheme and call it *particle-intensity modulation* (PIM).¹

In [9], [10], this concentration-based channel is considered with a receiver equipped with ligand receptors. The process of molecule reception of a ligand receptor is modeled as a Markov chain and the capacity in bits per channel use is analyzed. The results are extended to multiple access channels in [11]. In [12], a binomial distribution is used to model a system where the transmitter can perfectly control the release of particles and the receiver can perfectly detect the number of particles that arrive. It is assumed that the channel has finite memory and particle transport is assisted by flow. Using this

¹This is related to the concentration-shift-keying or the amplitude-modulation in previous work. However, the channel input X for concentration-shift-keying or amplitude modulation is an integer while for the PIM, X is a real number between zero and 1, i.e., a probability.

model, bounds on the capacity are derived, and the capacity for different memory lengths is analyzed. Reference [13] assumes that the channel input is the rate of particle release. The channel is represented as a Poisson channel with finite memory, and upper and lower bounds on capacity per channel use are presented. Finally, different channel coding schemes are compared for MC systems that employ PIM in [14].

This article extends our conference papers in [1], [2] where we considered molecular channels with imperfect PIM and imperfect detection. We call this channel the particle-intensity channel (PIC). Specifically, in the PIC the sender releases particles independently and probabilistically (i.e., information is encoded in the release probability), and the destination may not detect all the particles that arrive. Note that this is a different formulation than [1] in that the channel input is continuous as opposed to discrete. We assume that the duration of the symbol is long enough that particles from one symbol have a negligible effect on future symbols. This model is reasonable if particles diffuse beyond the receptor or disappear in some other fashion, for instance through degradation [15]. Under this assumption, the PIC is memoryless. Finally, we assume that particles can be generated at a constant fixed rate at the transmitter.

For this model, we show that the PIC can be represented with a channel model similar to the binomial channel [16] where the input is the probability of success and the output is the number of successes in a fixed number of Bernoulli trials. Like the binomial channel, the PIC channel input (probability of particle release) is continuous over the interval $[0, 1]$. However, unlike the original binomial channel, the probability of success (i.e., of the particle being detected by the receiver) is smaller than channel input so that the maximum probability of success is less than 1. This introduces asymmetry in the behavior induced by the extreme channel inputs; a zero induces a deterministic result at the receiver but a one does not.

Another difference from the original binomial channel is the introduction of a symbol duration. The number of trials in the PIC, which is the maximum number of particles that can be released by the transmitter, changes as a function of symbol duration because particles are generated at a constant rate at the transmitter.

Our contributions in this work are as follows.

- This article defines the capacity of the PIC channel in bits per second and as a function of symbol duration. We show that this channel is related to the binomial channel. To the best of our knowledge this is the first time that a channel model for molecular communication is presented that includes imperfections at *both* transmitter and receiver.
- This article demonstrates that the optimal input distribution for the PIC channel, for any symbol duration, always has mass points at the two extremes 0 and 1. We also derive an expression for the capacity when the input is binary, and present the conditions under which binary input achieves capacity.
- This article presents the *dynamic assignment Blahut-Arimoto* (DAB) algorithm as a new algorithm for finding the capacity and the optimal finite-support input distribution for many channels with continuous input alphabets, including the PIC channel. This algorithm

converges much faster than the ellipsoid method and finds the minimum-cardinality capacity-achieving input distribution. Using DAB, this article calculates the capacities and minimum-cardinality capacity-achieving input distributions for a wide range of channel parameters.

- Although the binary input distribution (i.e., on-off keying) is capacity-achieving for a large class of PICs based on diffusive particle transport, this article shows that the capacity-achieving input distribution has more than two mass points when the probability of particle arrival is sufficiently high.

The rest of this article is organized as follows. In Section II we present the PIC. Then in Section III, we formulate the capacity, investigate characteristics of the optimal input distribution, and derive the capacity of the binary input PIC. Section IV presents the ellipsoid method for finding the optimal input distribution for the binomial channel, while in Section V, the DAB algorithm is developed. We present numerical results in Sections VI, and in Section VII we discuss the concluding remarks.

II. THE PARTICLE INTENSITY CHANNEL (PIC)

The PIC is an MC channel in which information is communicated through PIM, i.e., the channel input X is the probability of particle release by the transmitter. The transmitter controls the intensity of the released particles by controlling this probability. The particles themselves are assumed to be *identical and indistinguishable* at the receiver, and no other properties (such as the time-of-release) are used for encoding information. The receiver then counts the number of particles that arrive during the symbol duration to produce the channel output Y . The particles released by the transmitter travel to the receiver through a propagation mechanism (e.g., diffusion). We assume that the particles *travel independently of each other*, and are *detected independently of each other*. This model is used in many previous works [3].

According to the PIC model, particles are released *instantly and simultaneously* at the beginning of the symbol interval. Particles are released independently of each other, and the transmitter's intended probability of release is X , which is the channel input. The transmitter controls the number of particles that are released by changing X , e.g., by controlling the degree of opening in a nozzle or porous membrane. Therefore, the channel input is continuous over the interval $X \in [0, 1]$.

Figure 1 shows the stochastic release, arrival, and detection of a single particle. We define the probability of release failure to be $1 - \alpha$ where $0 < \alpha \leq 1$. Thus the actual probability that each particle is released is αx , when the transmitter wants to transmit with probability x .

We now consider the stochastic particle transport. Each released particle will arrive at the receiver at some independent identically distributed random time $T \sim f_T(\cdot)$. Let $f_T(t)$ denote the PDF of the time the particle arrives, and $F_T(t)$ denote its corresponding CDF. Then the probability that a released particle arrives during a symbol duration τ is given by

$$\rho = F_T(\tau), \quad (1)$$

and the probability that it never arrives (and is by assumption never detected) is $1 - \rho$.

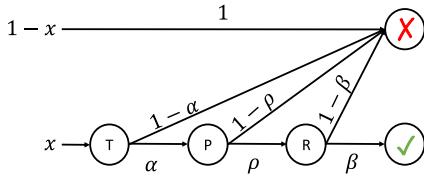


Fig. 1. The stochastic transmission (T), propagation (P) and reception (R) of a single particle. At the transmitter, particles are selected to be released with probability x . A selected particle is actually released with probability α . Through propagation, a released particle arrives at the receiver within the symbol duration with probability ρ . A particle that arrives at the receiver is detected with probability β . Thus, a particle selected for release is detected at the receiver with probability $\theta_\rho = \alpha\rho\beta$.

In the PIC model, each particle that arrives at the receiver is detected with probability β . This detection process is i.i.d. In this way, the PIC model incorporates the receiver's inability to perfectly detect the particles that arrive, owing to sensitivity or uncertainty in the detection process.

Thus, as shown in Figure 1 a particle selected for release is detected at the receiver with probability $\theta_\rho = \alpha\rho\beta$. Let Y be the number of particles that are actually detected during the corresponding symbol duration. The end-to-end channel between the input X and the output Y is the binomial $\mathcal{B}(m_\rho, x\theta_\rho)$:

$$P(Y = y|X = x) = \binom{m_\rho}{y} (x\theta_\rho)^y (1 - x\theta_\rho)^{m_\rho - y}, \quad (2)$$

where $\mathcal{B}(n, p)$ indicates the binomial distribution with parameters n and p , and m_ρ is the number of particles available for release, which we now discuss.

The channel is used in a time-slotted fashion, where τ is the symbol duration. We define a parameter λ as a constant fixed rate at which the transmitter can generate particles. We assume that $m_\tau = \lfloor \lambda\tau \rfloor$ particles are available to be released by the transmitter at the beginning of each time slot, and the transmitter releases each with channel input probability X . Note that in this model, the number of particles that can be released at the beginning of each time slot can change with the symbol duration τ .

If we assume $F_T(t)$ is strictly monotone, the symbol duration τ can be obtained from ρ by using the inverse CDF (iCDF) function, i.e., $\tau = F_T^{-1}(\rho)$. Using the iCDF, we can also rewrite m_τ as a function of ρ :

$$m_\rho = \lfloor \lambda F_T^{-1}(\rho) \rfloor. \quad (3)$$

Particles that neither arrive at the receiver nor dissipate over the symbol duration could interfere with detection during future channel uses. Such intersymbol interference (ISI) should be incorporated into deriving the channel capacity. While such ISI represents an interesting area for future investigation, in this article we assume that particles with transit times exceeding τ dissipate or are otherwise inactivated. That is, particles are assumed to have a finite lifetime of duration τ . This assumption seems reasonable since particles could be rendered undetectable either naturally or by design (via denaturing or gettering/enzyme reactions [15]). Under this assumption, channel uses are independent and the maximum mutual information

between input and output during a single channel use defines the channel capacity.

The PIC channel is an accurate model for certain molecular communication channels. Specifically, the transmitter in the PIC can represent a large class of transmitters that control the particle-intensity through gated opening(s) on the surface of the transmitter. One example of such transmitters is the ion-pump transmitter discussed in [17]. In addition, the memoryless characteristic of the PIC channel is an accurate model for relatively large symbol durations using a single particle type, when different particle types are released during consecutive channel uses [18], or when enzymes in the environment degrade the particles that do not arrive at the receiver by the end of the symbol duration [19]. The PIC receiver also accurately models, at a high level, the ligand receptors on the surface of a synthetic cell [20].

An important observation here is that ρ and m_ρ change as a function of the symbol duration τ . In this work, we incorporate the optimization of the symbol duration into the formulation of capacity to determine the channel capacity of the memoryless PIC in bits per second. This is one important distinction between this and previous work such as [5], [6], [12], [13], where the channel capacity is typically defined in bits per channel use.

For the case when m_ρ is large and θ_ρ is small, the system can be well approximated by the Poisson distribution [21]

$$P(Y = y|X = x) = \frac{(x\theta_\rho m_\rho)^y e^{x\theta_\rho m_\rho}}{y!}. \quad (4)$$

We write this as $P(y|x; \rho) \sim \mathcal{P}(x\theta_\rho m_\rho)$, where $\mathcal{P}(a)$ indicates the Poisson distribution with parameter a .

Remark 1: Using the Poisson approximation, the PIC in MC systems can be viewed as a more general formulation of the discrete-time Poisson channel used to model optical intensity channels [22]–[24]. Because a finite number of particles are released, particle arrival rate does not increase linearly with time and thus neither does the capacity. This is in contrast to the discrete-time Poisson channel in optical communications, where photon arrival rate increases linearly with the symbol duration [24]. PIC symbol durations that are too long can reduce the information rate, as demonstrated in Section VI. Note that although we do not consider interfering particles, they can be introduced to the Poisson model in (4) by adding an extra term similar to the dark current in optical communications [22]–[24].

III. CHANNEL CAPACITY AND OPTIMAL INPUT

We now characterize the channel capacity of the PIC. Let $f_X(x)$ be the channel input PDF and let \mathcal{F} be the set of all valid input PDFs. Then the capacity of the channel in (2) as a function of the particle arrival probability ρ and having units of [bits per second] is defined as

$$C(\rho) = \max_{f_X(x) \in \mathcal{F}} \frac{I(X; Y|\rho)}{F_T^{-1}(\rho)}, \quad (5)$$

where $F_T^{-1}(\cdot)$ is the iCDF of the particle detection time. Since the channel changes as a function of the symbol duration, the

fundamental limit of this channel is then

$$C^* = \max_{\rho} C(\rho). \quad (6)$$

We now investigate the characteristics of the optimal input distribution, determine capacity under a binary input constraint, and investigate settings for which the capacity-achieving distribution is binary. First, observe that although the input distribution is over a finite interval, the capacity-achieving input distribution has finite support, requiring at most $m_\rho + 1$ mass points for the PIC in (2). This was proven in [25] using Dubin's theorem [26]. See also [27, Corollary 3 in Ch. 4.5].

Throughout the paper we will use a tilde over a letter to indicate that this is an ordered set or vector whose cardinality is equal to the number of mass points being used by the associated input distribution. Let $\tilde{\mathcal{X}} = \{x_0, x_1, \dots, x_{m_\rho}\}$ be the location of the $m_\rho + 1$ mass points with $0 \leq x_0 < x_1 < \dots < x_{m_\rho} \leq 1$. Let $\tilde{P} = \{p_{x_0}, p_{x_1}, \dots, p_{x_{m_\rho}}\}$ be the probabilities corresponding to each mass point. Then the corresponding input distribution is given by

$$f_X(x) = \sum_{i=0}^{m_\rho} p_{x_i} \delta(x - x_i). \quad (7)$$

The optimal input distribution always has the form of (7). We now show that the optimal input always has non-zero mass points at the two extremes of $X = 0$ and $X = 1$.

Theorem 1: For a given symbol duration τ , and hence a given ρ , the mutual information given in (5) is maximized by a PDF $f_X^*(x)$, where $p_0^* > 0$ and $p_1^* > 0$.

Proof: The first inequality $p_0^* > 0$ can be proved by using [24, Lemma 1] and the second inequality $p_1^* > 0$ can be proved using [24, Lemma 3]. ■

We now derive the capacity in (5) as a function of ρ for binary input PIC (i.e., a system that is limited to on-off-keying). Note that conveniently for the PIC in (2), on-off-keying is equivalent to $X = 1$ or $X = 0$.

Theorem 2: Let $\tilde{\mathcal{X}}_b = \{0, 1\}$ be the selected input alphabet for the PIC in (2) with $p_1 = P(X = 1)$, and $\varphi_\rho = (1 - \theta_\rho)^{m_\rho}$. The optimal input distribution p_1 under this binary input constraint is given by

$$p_1^* = \frac{1}{\frac{\varphi_\rho}{\varphi_\rho - 1} - \varphi_\rho + 1}, \quad (8)$$

and the capacity of (5), in bits per second, is given by

$$C^b(\rho) = \frac{1}{F_T^{-1}(\rho)} \log \left(1 + (1 - \varphi_\rho) \varphi_\rho^{\frac{\varphi_\rho}{1 - \varphi_\rho}} \right). \quad (9)$$

Proof: Define Y^+ as the indicator function

$$Y^+ = \mathbf{1}(Y > 0). \quad (10)$$

Y^+ is a sufficient statistic of Y for X [28] so that the binary-input PIC is equivalent to a Z channel [29]. Thus, the mutual information in (5) can be written as a function of p_1 using

$$I_\rho(X^b; Y) = I_\rho(p_1) \quad (11)$$

$$= H(p_1(1 - \varphi_\rho)) - p_1 H(\varphi_\rho). \quad (12)$$

Setting the derivative of $I_\rho(p_1)$ with respect to p_1 equal to zero yields (8). Substituting (8) into (12) and using (5) we obtain the capacity expression in (9). ■

An interesting question arises here as to when the binary input alphabet $\tilde{\mathcal{X}}_b$ is optimal for the PIC. In the following proposition, we provide a guideline for the optimality of the binary input for a subclass of PICs.

Proposition 1: For the PIC in (2) where m_ρ is large and θ_ρ is small such that the Poisson approximation in (4) is accurate, the binary input distribution given in (8) is optimal if $m_\rho \theta_\rho < 3.3679$.

Proof: Using the same technique presented in [22] for the optical channels, the proposition can be proved. See [22, Sec. 3] for a detailed derivation of this result. ■

Note that this condition may be satisfied in many practical systems where the radius of the receiver is much smaller than the distance between the transmitter and the receiver, hence the probability of particles arriving is small. Upper bounds on the total variation between binomial and Poisson distributions can be used to show that this variation is small for small θ_ρ [21].

IV. COMPUTING CAPACITY OF THE BINOMIAL CHANNEL

The previous section provides an expression for the capacity when the optimal input distribution is binary. This section and the next address how to compute the capacity of the PIC channel when the optimal input is not binary. Several papers have addressed similar capacity computations including [16], [30], [31], [24], [32], [33], [34], [35], [36], and [37].

This section lays the foundation for Section V to introduce the DAB algorithm as a general solution technique for computing the capacity and optimal finite-support input distribution for channels with continuous input alphabets. Section IV-A introduces the binomial channel, which is a corner case of the PIC channel. Section IV-B formulates the binomial channel capacity problem as a convex optimization problem, and presents its dual. Section IV-C explores the Ellipsoid method as one technique to solve the dual problem.

A. The Binomial Channel

To simplify our initial development, we focus on the corner case of the PIC channel (2) where $\theta_\rho = 1$. In this case the PIC channel is the binomial channel (of parameter $n = m_\rho$) and has a channel law defined by the binomial probability distribution of order n [16]. For each channel use, the input X is the probability of success of a Bernoulli trial. The channel output Y is the number of successes observed during n Bernoulli trials. Thus the channel transition probability law is described as

$$P_{Y|X}^{(n)}(y|x) = \binom{n}{y} x^y (1-x)^{n-y} \quad (13)$$

where the possible y values are the integers zero through n . The channel output could also be the ordered list of Bernoulli trial outcomes, but since Y is a sufficient statistic [38] of those outcomes for estimating X , the capacity is the same.

Figure 2 shows these finite-support capacity-achieving distributions for $1 \leq n \leq 50$, with the area of the circle indicating

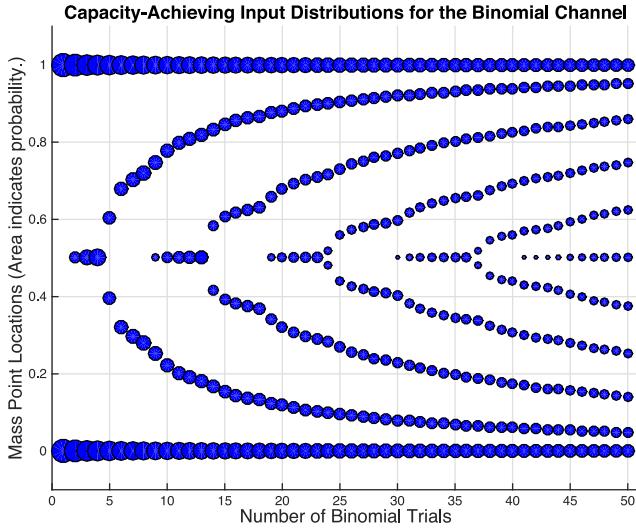


Fig. 2. Capacity-achieving input distributions for the binomial channel with $1 \leq n \leq 50$ obtained by the Dynamic Assignment Blahut-Arimoto algorithm described in Section V.

the probability of the mass point. For $n = 1$, the mass points are at zero and one, resulting in a noiseless binary channel. At $n = 2$, a mass point is introduced at 0.5, growing in probability as n increases, until at $n = 5$ that mass point splits into two. As n increases, these two mass points move away from 0.5, and when, at $n = 9$, they are far enough away, a new mass point is born at 0.5.

B. A Convex Optimization Problem and Its Dual

Despite the fact that finite-support distributions achieve capacity of the PIC, direct application of the Blahut-Arimoto algorithm [39] is complicated because the locations of the support points within the unit interval are not known. Reasonable approximations can be obtained by applying Blahut-Arimoto with mass points closely spaced along the entire unit interval, with most of these having zero probability. However, we are interested in algorithms that identify the capacity more precisely and that explicitly identify the mass points of the capacity-achieving distribution.

In [16], the capacity of the binomial channel is computed by first formulating the problem as a convex optimization problem and then solving it by using the Ellipsoid method. Assuming that X has discrete support, capacity C_n is

$$C_n = \max_{f_X(x)} I(X; Y) \quad (14)$$

$$= H(Y) - H(Y|X) \quad (15)$$

$$= \max_{f_X(x)} \left\{ H(Y) - \int_{x=0}^1 f_X(x) H(Y|X=x) f_X(x) dx \right\}. \quad (16)$$

Despite the fact that the capacity-achieving distribution on X has at most $n + 1$ mass points, the distribution on X is expressed as a density function $f_X(x)$ (and an integral is used in (16)) because the positions of the support points are located anywhere in the uncountable set of the unit interval.

Thus, $f_X(x)$ consists of a countable number of delta functions located anywhere in the unit interval.

The optimization problem of (16) can be formulated as a convex optimization problem in a vector space with uncountably infinite dimension. For more mathematical precision, replace $f_X(x)dx$ with $dF(x)$ where $F(x)$ is the cumulative distribution. We allow $dF(x) \in \mathcal{F}$, the set of signed measures on the unit interval and include additional constraints to force $dF(x)$ to be a valid probability distribution. Introducing the additional variables q_y for $y \in \{0, 1, \dots, n\}$ and appropriate equality constraints that force the q_y values to be the output probability distribution $P_Y(y)$ induced by the input distribution yields the following convex optimization **primal problem**:

$$\text{minimize} \sum_{y=0}^n q_y \log q_y + \int_{x=0}^1 f_X(x) H(Y|X=x) dF(x)$$

$$\text{subject to } -dF(x) \leq 0, \forall x \in [0, 1]$$

$$\int_{x=0}^1 dF(x) - 1 = 0$$

$$q_y - \int_{x=0}^1 P_{Y|X}^{(n)}(y|x) dF(x) = 0, \\ y \in \{0, \dots, n\}.$$

The infinite dimensional $dF(x)$ makes the problem intractable. We create a Lagrangian dual problem that can be solved with traditional methods. We introduce Lagrange multipliers $\mu, z_0, z_1, \dots, z_n$ for the equality constraints and the measurable mapping $v(x)$ of $[0, 1]$ to the one-dimensional real space \mathbf{R} for the inequality constraint producing Lagrangian $L(\mathbf{q}, dF(x), v(x), \mathbf{z}, \mu)$ [40]:

$$L = \sum_{y=0}^n q_y \log q_y + \int_{x=0}^1 H(Y|X=x) dF(x) \\ - \int_{x=0}^1 v(x) dF(x) + \mu \left(\int_{x=0}^1 dF(x) - 1 \right) \\ + \sum_{y=0}^n z_y \left(q_y - \int_{x=0}^1 P_{Y|X}^{(n)}(y|x) dF(x) \right)$$

which is the cost function augmented with the weighted sum of the constraints. Minimizing $L(\mathbf{q}, dF(x), v(x), \mathbf{z}, \mu)$ with respect to primal variables \mathbf{q} and $dF(x)$ gives the dual function $g(v(x), \mathbf{z}, \mu)$ as follows:

$$g = \inf_{\mathbf{q}, dF(x)} \left\{ \sum_{y=0}^n q_y (z_y + \log q_y) - \mu + \int_{x=0}^1 \gamma(x) dF(x) \right\},$$

where

$$\gamma(x) = H(Y|X=x) - v(x) + \mu - \sum_{y=0}^n z_y P_{Y|X}^{(n)}(y|x). \quad (17)$$

Because $dF(x)$ is an unconstrained unsigned measure, $g(v(x), \mathbf{z}, \mu) = -\infty$ unless $\gamma(x) \geq 0$ for all $x \in [0, 1]$, in

which case we have

$$g(v(x), \mathbf{z}, \mu) = \inf_{\mathbf{q}} \left\{ \sum_{y=0}^n q_y (z_y + \log q_y) - \mu \right\}. \quad (18)$$

Setting d/dq_y of the summation in (18) to zero yields the minimizing value of $q_y = \frac{2^{-z_y}}{e}$ so that

$$g(v(x), \mathbf{z}, \mu) = \frac{-\log e}{e} \sum_{y=0}^n 2^{-z_y} - \mu. \quad (19)$$

The **dual problem** for our primal problem maximizes this $g(v(x), \mathbf{z}, \mu)$ subject to constraints on the slack variables:

$$\text{maximize } \frac{-\log e}{e} \sum_{y=0}^n 2^{-z_y} - \mu$$

subject to $v(x) \geq 0 \quad \forall x \in [0, 1]$,

$$\begin{aligned} H(Y|X=x) - v(x) + \mu \\ - \sum_{y=0}^n z_y P_{Y|X}^{(n)}(y|x) = 0. \end{aligned}$$

Combining these two constraints eliminates the cumbersome infinite-dimensional $v(x)$ producing

$$\text{minimize } \sum_{y=0}^n 2^{-z_y} + \frac{\mu e}{\log e}$$

$$\text{subject to } H(Y|X=x) + \mu - \sum_{y=0}^n z_y P_{Y|X}^{(n)}(y|x) \geq 0 \\ \forall x \in [0, 1].$$

Minimizing the objective function requires the minimum possible value of μ that satisfies

$$\mu \geq \sum_{y=0}^n z_y P_{Y|X}^{(n)}(y|x) - H(Y|X=x) \quad \forall x \in [0, 1] \quad (20)$$

which leads to the final formulation of the dual problem, in which only the variables z_y remain:

$$\begin{aligned} \min_{\mathbf{z}} \quad & \sum_{y=0}^n \frac{2^{-z_y}}{e} + \frac{1}{\log e} \max_{x \in [0, 1]} \\ & \times \left\{ \sum_{y=0}^n z_y P_{Y|X}^{(n)}(y|x) - H(Y|x) \right\}. \end{aligned}$$

The dual problem is a finite variable convex optimization problem over the vector \mathbf{z} , which can be solved using a variety of techniques.

Once the minimizing \mathbf{z} vector is obtained, complementary slackness indicates that the capacity-achieving mass points are the x values that maximize $\sum_{y=0}^n z_y P_{Y|X}^{(n)}(y|x) - H(Y|x)$. The output distribution is recovered using $P_Y(y) = q_y = \frac{2^{-z_y}}{e}$, and the probability $P_X^*(x)$ associated with each mass point can be found by solving the equations:

$$P_Y(y) = \sum_{x \in A} P_{Y|X}^{(n)}(y|x) \tilde{P}(x) \quad \forall y \in \{0, \dots, n\}, \quad (21)$$

where A is the set of maximizing x values. The mutual information induced by this P_X is also the capacity.

C. The Ellipsoid Method

In [16], the ellipsoid method was used to solve the dual problem identified in Section IV-B. The ellipsoid method was developed by Shor, Nemirovski, and Yudin in the 1970's and used by Khachiyan [41] in 1979 to show the polynomial solvability of linear programs. See [42] for an excellent survey. One conclusion of [42] is that the ellipsoid method, while of academic interest, is often not the fastest way to solve a convex problem and can have stability issues as well. However, it is straightforward to program.

The method begins with an initial ellipsoid $\mathcal{E}^{(0)} \in \mathbb{R}^N$ centered at a point z_0 , which is defined as

$$\mathcal{E}^{(0)} = \left\{ z \in \mathbb{R}^N : (z - z_0)^T P_0^T (z - z_0) \leq 1 \right\} \quad (22)$$

and is known to contain the optimizing point z^* . At the k^{th} iteration, the point z_k is at the center of the ellipsoid

$$\mathcal{E}^{(k)} = \left\{ z \in \mathbb{R}^N : (z - z_k)^T P_k^T (z - z_k) \leq 1 \right\}. \quad (23)$$

To compute the $\mathcal{E}^{(k+1)}$ we need the subgradient $g_{k+1} \in \mathbb{R}^N$ which is a vector that satisfies $g_{k+1}^T (z^* - z_k) \leq 0$, so that

$$z^* \in \mathcal{E}^{(k)} \cap \left\{ z : g_{k+1}^T (z - z_k) \leq 0 \right\}. \quad (24)$$

The following computations create a new ellipsoid that contains the half-ellipsoid described above:

$$\tilde{g}_{k+1} = \left(\sqrt{g_{k+1}^T P_k g_{k+1}} \right)^{-1} g_{k+1} \quad (25)$$

$$z_{k+1} = z_k - \frac{1}{N+1} P_k \tilde{g}_{k+1} \quad (26)$$

$$P_{(k+1)} = \frac{N^2}{N^2-1} \left(P_k - \frac{2}{N+1} P_k \tilde{g}_{k+1} \tilde{g}_{k+1}^T P_k \right). \quad (27)$$

The ellipsoid method stopping criterion computes $\sqrt{g_k^T P_k g_k}$ which is an upper bound on the error in the objective function. To apply the ellipsoid algorithm to the dual problem of Section IV-B, $N = n + 1$ and the subgradient is the vector with elements

$$-\frac{2^{-z_y} - e P_{Y|X}^{(n)}(y|x^*)}{\log e} \quad (28)$$

where x^* is any value of x that maximizes

$$\sum_{y=0}^n z_y P_{Y|X}^{(n)}(y|x) - H(Y|x) \quad (29)$$

for the current set of z_y values. Also needed are z_0 and P_0 that create an initial ellipse that contains z^* . The simplest approach is to select the origin for z_0 and choose P_0 to be the identity scaled by a value that is larger than the square of the optimizing z^* . Figure 3 shows for the first 25 values of n these squared distances, which grow to over 350 by $n = 25$.

However, there is a difficulty in knowing what the squared distance is before the problem has been solved. This problem is avoided by selecting the initial z vector for the $(n+1)^{\text{st}}$ case by using the solution obtained for the n^{th} case as follows:

$$q_y^{(n+1, \text{initial})} = \sum_{x \in A} P_{Y|X}^{(n)}(y|x) P_X^{(n, *)}(x) \quad \forall y \in 0, \dots, n,$$

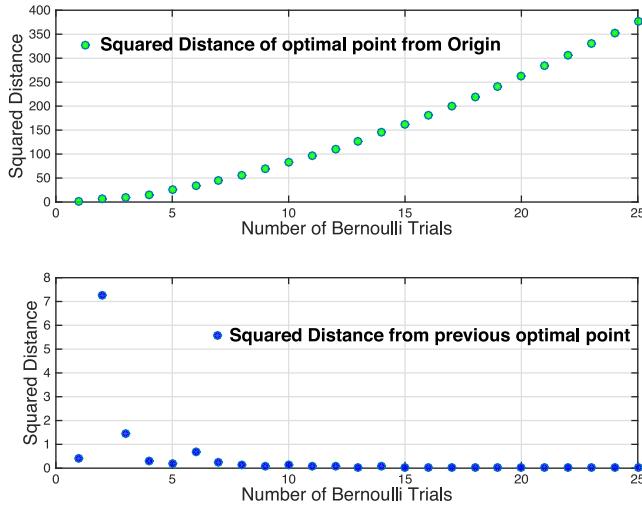


Fig. 3. Squared distance of optimal points from the origin (top) and from the optimal point from the previous iteration (bottom).

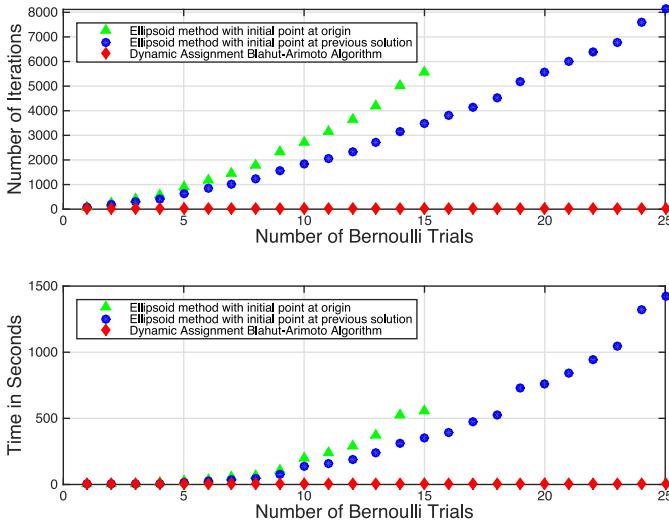


Fig. 4. Number of iterations and computation time in seconds (on a 2.5 GHz MacBook Pro purchased in 2014) running the Ellipsoid and Dynamic Assignment Blahut-Arimoto (DAB) algorithms implemented in MATLAB to compute the binomial channel capacity: Ellipsoid method with initial point at the origin (green triangles), Ellipsoid method with initial point at position indicated by the solution for the previous n (blue dots), and the DAB algorithm with initialization as described in Section V-A (red diamonds).

$$z_y^{(n+1, \text{initial})} = -\log e q_y^{(n+1, \text{initial})}.$$

In this case, as shown in Figure 3, P_0 can often be the unscaled identity (or the identity scaled by a value less than one). With such a close starting value, one would expect that the ellipsoid method would converge much more quickly. However, as shown in Figure 4, initializing the z vector to the previously optimal point does not significantly improve performance, highlighting the slow convergence of the ellipsoid algorithm even when initialized to a favorable point.

V. DYNAMIC ASSIGNMENT BLAHUT-ARIMOTO

This section introduces the Dynamic Assignment Blahut-Arimoto (DAB) algorithm, which computes the capacity and associated capacity-achieving distribution when the input

alphabet is continuous but a capacity-achieving (or capacity-approaching) distribution is known to have finite support.

The DAB approach alternates between a Blahut-Arimoto step that optimizes the allocation of probability to a fixed set of mass points and a second step that adjusts the placement (and possibly the number) of mass points given the PMF identified by the previous Blahut-Arimoto step.

Algorithm 1 summarizes DAB. This iterative algorithm starts by initializing the number of mass points $N^{(1)}$ and their locations $\tilde{\mathcal{X}}^{(1)}$. DAB can increase the number of mass points if necessary. During the k^{th} iteration of the algorithm, first $\tilde{\mathcal{X}}^{(k)}$ is used with Blahut-Arimoto algorithm to maximize MI and find the corresponding maximizing distribution $\tilde{p}^{(k)}$. This provides a lower bound on capacity. Then, the distribution $p(y)$ induced by $\tilde{\mathcal{X}}^{(k)}$ and $\tilde{p}^{(k)}$ are used to compute an upper bound on capacity. If these bounds are within a specified threshold, then DAB has found the capacity (to a specified level of precision) and the associated capacity-achieving (or capacity-approaching) input distribution. Otherwise the location of the mass points and/or the number of mass points needs to be updated. These updates occur in steps 4, 5, and 6 of Algorithm 1. The subsections that follow explore specific steps of Algorithm 1 in more detail.

A. Initialization

Unlike the Ellipsoid method, proper initialization can dramatically reduce computation time for DAB. Consider again Figure 2. The mass points often move only slightly as n progresses; significant changes occur only when the mass point at $x = 1/2$ splits or when a new mass point is born at $x = \frac{1}{2}$. Notice that the gentle evolution of the capacity-achieving distribution as a function of n is such that the number of mass points never increases by more than one.

DAB allows this behavior to be exploited. A clever approach to initialization uses the previously computed capacity-achieving distribution of a channel in the same “family” with a slightly lower capacity as a starting point for its optimization. For the binomial channel, the starting distribution $\tilde{\mathcal{X}}^{(1)}$ may be selected as the capacity-achieving distribution for the binomial channel with one fewer Bernoulli trial. Figure 4 shows how this initialization approach (and the general efficiency of DAB) lead to dramatically smaller values for run times and required iterations as compared to the Ellipsoid method.

B. Upper Bound via Csiszar’s Min-Max Capacity Theorem

Step 2 of DAB relies on Csiszar’s Min-Max Capacity Theorem [43, Th. 3.4], which states:

$$C = \min_{p(y) \in \{P_Y\}} \max_x D(p(y|x) \parallel p(y)) \quad (32)$$

where $\{P_Y\}$ is the set of distributions on Y that can be induced by a valid input distribution. In fact, we can restate the dual problem found in Section IV-B in terms of Csiszar’s Min-Max Capacity Theorem as follows:

$$\min_{p(y)} \left\{ \sum_y P_Y(y) - 1 + \max_x D(p(y|x) \parallel p(y)) \right\}. \quad (33)$$

Algorithm 1 Dynamic Assignment Blahut Arimoto

Initialization: Select $\tilde{\mathcal{X}}^{(1)} = [x_1 \ x_2 \ \cdots \ x_{N^{(1)}}]$, the vector specifying the $N^{(1)}$ mass point locations in increasing order. Select the tolerance ϵ controlling the accuracy of the final capacity. Set $k = 1$.

Iterations: Determination of the optimal $\tilde{\mathcal{X}}^*$, $\tilde{\mathcal{P}}^*$, and the capacity C (within ϵ bits) proceeds as follows:

- 1) Given $\tilde{\mathcal{X}}^{(k)}$, use Blahut-Arimoto to compute the MI maximizing distribution $\tilde{\mathcal{P}}^{(k)}$ and the corresponding MI $I^{(k)}$, which is a lower bound on C .
- 2) Use the distribution $p(y)$ induced by $\tilde{\mathcal{X}}^{(k)}$ and $\tilde{\mathcal{P}}^{(k)}$ to compute the capacity upper bound

$$D_{\max}^{(k)} = \max_{x \in \mathcal{X}} D(p(y|x) \parallel p(y)).$$

- 3) If $D_{\max}^{(k)} - I^{(k)} < \epsilon$ conclude by reporting $\tilde{\mathcal{X}}^* = \tilde{\mathcal{X}}^{(k)}$, $\tilde{\mathcal{P}}^* = \tilde{\mathcal{P}}^{(k)}$, and $C = I^{(k)}$. Otherwise continue.
- 4) Determine whether $N^{(k+1)} = N^{(k)}$ or $N^{(k)} + 1$ and if the latter, update $\tilde{\mathcal{X}}^{(k+1)}$ to include the additional location.
- 5) Determine direction vector $\tilde{\mathcal{D}}^k$ to adjust $\tilde{\mathcal{X}}$.
- 6) Compute

$$\tilde{\mathcal{X}}^{(k+1)} = \tilde{\mathcal{X}}^{(k)} + \lambda^* \tilde{\mathcal{D}}^k, \quad (30)$$

where

$$\lambda^* = \arg \max_{\lambda} I(\tilde{\mathcal{X}}^{(k)} + \lambda \tilde{\mathcal{D}}^k, \tilde{\mathcal{P}}^{(k)}), \quad (31)$$

and $I(\tilde{\mathcal{X}}, \tilde{\mathcal{P}})$ is the mutual information that results from an input distribution with mass points whose locations are described by the vector $\tilde{\mathcal{X}}$ with corresponding probabilities are described by the vector $\tilde{\mathcal{P}}$.

- 7) $k = k + 1$.
- 8) Go to 1.

An upper bound on capacity follows directly from (32): For any valid output distribution on \mathcal{Y} ,

$$C \leq \max_x D(P_{Y|X=x} \parallel P_Y). \quad (32)$$

C. Determining Direction Vector $\tilde{\mathcal{D}}^k$

In this step of DAB (step 5 of Alg. 1), a direction is selected along which $\tilde{\mathcal{X}}^k$ will be varied in step 6 to increase the mutual information $I(\tilde{\mathcal{X}}, \tilde{\mathcal{P}})$. This article considers two approaches to selecting the direction: moving a single mass point (or symmetric mass-point pair) and moving all points by setting the direction to be the relevant gradient.

1) *Moving a Single Mass Point (or a Symmetric Pair):* The technique considered in the original DAB [2] was to select and move a single mass point so that $\tilde{\mathcal{D}}^k = e_j$, which the vector with all elements set to zero except the j^{th} element, which is set to one. When the capacity-achieving distribution is known to have symmetry about its center, $\tilde{\mathcal{D}}^k$ should be selected as a symmetric pair of mass points so that

$$\tilde{\mathcal{D}}^k = e_j + e_{N^{(k)}+1-j}. \quad (35)$$

The binomial channel is an example where $\tilde{\mathcal{X}}^*$ is symmetric (about 1/2) so that a symmetric pair of mass point locations are set to 1 to form $\tilde{\mathcal{D}}^k$. The PIC does not have such symmetry.

a) *Proximity to x_{\max} :* In [2], motivated by reducing the upper bound of (34), for the specific case of the binomial channel, the original DAB sets $\tilde{\mathcal{D}}^k = e_j$ where the mass point x_j is the point in the interval bounded by x_{\max} and 1/2 that is closest to x_{\max} , where

$$x_{\max}^{(k)} = \arg \max_x D(p(y|x) \parallel p(y)). \quad (36)$$

b) *Maximum derivative:* In this article we also consider the alternative where DAB selects the mass point x_j to maximizes the partial derivative

$$\frac{\partial I(\tilde{\mathcal{X}}^{(k)}, \tilde{\mathcal{P}}^{(k)})}{\partial x_j} = p_j \frac{\partial D(p(y|x_j) \parallel p(y))}{\partial x_j}. \quad (37)$$

c) *Round robin:* As a third alternative DAB could select the mass points (or symmetric pairs) in a round robin fashion. This is a conservative approach that makes sure that every pair gets a chance to adjust its position.

2) *Moving Along the Gradient:* Considering (37), rather than selecting any single mass point $\tilde{\mathcal{D}}^k = e_j$, this approach sets $\tilde{\mathcal{D}}^k$ to the gradient:

$$\tilde{\mathcal{D}}^k = \nabla_{\tilde{\mathcal{X}}} I^{(k)} = \begin{bmatrix} \frac{\partial I(\tilde{\mathcal{X}}^{(k)}, \tilde{\mathcal{P}}^{(k)})}{\partial x_1} \\ \frac{\partial I(\tilde{\mathcal{X}}^{(k)}, \tilde{\mathcal{P}}^{(k)})}{\partial x_2} \\ \vdots \\ \frac{\partial I(\tilde{\mathcal{X}}^{(k)}, \tilde{\mathcal{P}}^{(k)})}{\partial x_N} \end{bmatrix}. \quad (38)$$

When optimizing along the direction of the gradient, the value of scalar λ in (30) is limited so that the mass point locations cannot cross each other. Figure 5 compares values for the running time and the number of iterations required to calculate the capacity and the optimal input distribution for the Binomial channel in Figure 2 based on methods 1b) and 2. As can be seen moving a single mass point tends to result in faster convergence. Therefore, in the rest of the paper we will use this method for numerical evaluations.

D. Determining λ to Maximize $I(\tilde{\mathcal{X}}, \tilde{\mathcal{P}})$

A line search routine (such as fminbnd in MATLAB) determines the value of λ that maximizes $I(\tilde{\mathcal{X}}, \tilde{\mathcal{P}})$ in step 6 of Algorithm 1. Since mutual information is a concave function of the input distribution [28], this line search is guaranteed to find the globally optimal point. As the number of mass points grows, fully optimizing λ to maximize $I(\tilde{\mathcal{X}}, \tilde{\mathcal{P}})$ provides more consistent performance than the approach in [2] of moving the mass point in the direction of x_{\max} using a step size δ .

E. Determining When to Increment $N^{(k)}$

There are three possible approaches to deciding when to increment $N^{(k)}$ in step 4 of Algorithm 1.

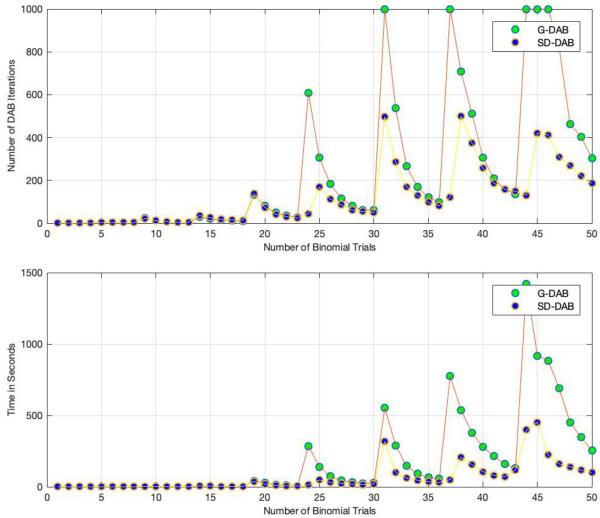


Fig. 5. Number of iterations and computation time in seconds (on a 2.5 GHz MacBook Pro purchased in 2018) running the Dynamic Assignment Blahut-Arimoto (DAB) algorithms implemented in MATLAB to compute the binomial channel capacity with the single-derivative (SD) DAB described in Section V-C1 and the Gradient (G) DAB described in Section V-C2 Overall, SD-DAB proves to be superior to G-DAB in computational speed and number of iterations.

1) *“Missing Mass Point” Approach:* When applying approach 1a of Sec V-C for the binomial channel, an additional mass point is added whenever none of the mass points in $\tilde{\mathcal{X}}^{(k)}$ lies in the interval bounded by x_{\max} and $1/2$. If $N^{(k)}$ is even, a new point is added at $1/2$. If $N^{(k)}$ is odd, a new point is added by splitting the point at $1/2$ into two points that will then be pulled away from $1/2$ by the line search of step 7.

2) *Minimum Derivative Test:* When applying approaches 1b or 2 of Sec V-C, additional mass point is added whenever the largest derivative is small enough that further improvement requires an additional mass point. The test in step 4 reveals when further mutual information increase is possible. When the derivatives all become too small to allow the potential improvement identified in step 4, this is a clear indication that an additional mass point is needed. If $N^{(k)}$ is odd, a new point is added by splitting the central point. If $N^{(k)}$ is even, a new point is added in between the two central points.

3) *Negligible Rate of Change of $I^{(k)}$:* When applying any of the approaches of Section V-C, simply tracking the increase in mutual information can be an effective way to determine when an additional mass point is needed. If the $I^{(k)} - I^{(k-1)}$ becomes negligible, then a new mass point is added. As above, if $N^{(k)}$ is odd, a new point is added by splitting the central point. If $N^{(k)}$ is even, a new point is added in between the two central points.

VI. NUMERICAL EVALUATIONS

This section uses DAB to evaluate the capacity and the optimal input distribution for several instances of the PIC. The assumed mechanism for particle transport is diffusion with coefficient d from a point-source transmitter to the surface of a spherical receiver with radius r . Let ℓ be the shortest distance between the point source and the receiver surface. Under this

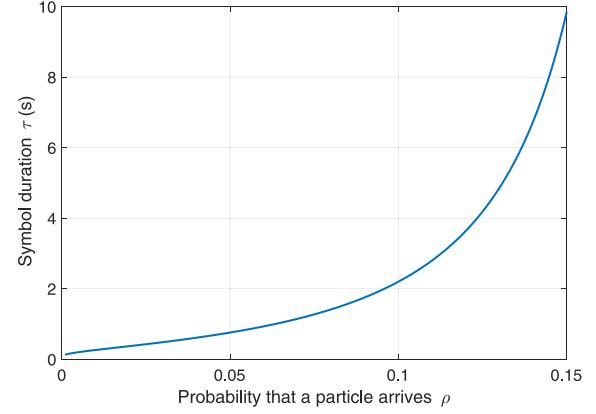


Fig. 6. Plot of the iCDF of the $\tilde{\mathcal{X}}$ Lévy-distribution for $c = 1$ and $\eta = 0.2$.

model, the motion of each particle can be represented using a random Brownian path in 3D space.

Since we assume that the particles are either detected when they arrive at the receiver or they are never detectable, the time of arrival is given by the first time the particle reaches the receiver. For Brownian motion in 3D space, the first arrival time T to the spherical receiver is a scaled Lévy-distributed random variable where the scale constant is $\eta = \frac{r}{\ell+r}$ [44]. This means that there is a non-zero probability that a particle never arrives at the receiver. Note that for Brownian motion in 1D space, $\eta = 1$. Using the iCDF of the scaled Lévy distribution, we obtain

$$\tau = F_T^{-1}(\rho) = \frac{c}{2 \operatorname{erfcinv}^2(\rho/\eta)}. \quad (39)$$

where $c = \frac{\ell^2}{2d}$ and $\operatorname{erfcinv}(.)$ is the inverse of the complementary error function $\operatorname{erfc}(.)$. We call a channel that relies on this diffusive transport the *diffusion-based PIC (DBPIC)*.

Remark 2: Substituting (39) into (5), we observe that the diffusion coefficient d has no effect on the optimal input distribution and the optimal ρ . This is despite the fact that the capacity increases linearly with d . This means that if the type of particle is changed, so long as the distance between the transmitter and the receiver is the same, and the receiver has the same radius, the optimal distribution and the optimal ρ values will remain the same. Note that the change in capacity is due to the fact that a shorter or a longer symbol duration is required to achieve the same ρ value.

Remark 3: If we consider a 1D environment² (i.e., $\eta = 1$), we observe that the capacity decreases as $\frac{1}{l^2}$, and the distance l does not affect the optimal input distribution and the optimal ρ . For a 3D environment however, changing the distance l and the radius r could affect the optimal ρ and $P(x)$ values through η .

Figure 6 shows the iCDF of the scaled Lévy distribution in (39) for $c = 1$ and $\eta = 0.2$. Note that a small increase in ρ can require a large increase in symbol duration. Therefore, larger ρ values may not necessarily result in higher information rate in bits per second. This effect is verified in Figure 7 where

²Note that a 1D environment is a good approximation if the system is confined inside a very narrow and long physical channel.

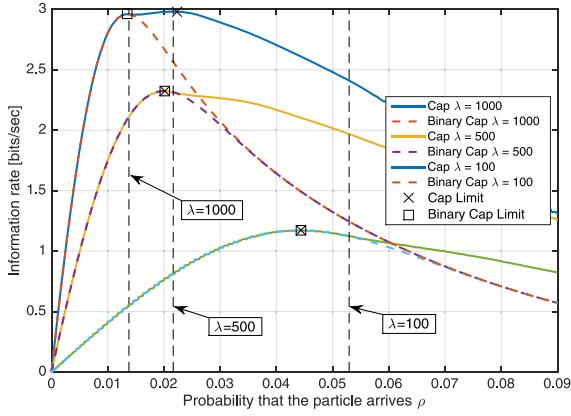


Fig. 7. Information rates for the PIC channel with $c = 1$, $\eta = 0.2$, $\alpha = 0.9$, and $\beta = 0.9$ for binary-input (dashed lines) and information rates with the optimal input (solid lines) for different λ . The three vertical dashed lines indicate the ρ value after which $m\rho\theta\rho$ in Proposition 1 is greater than 3.3676.

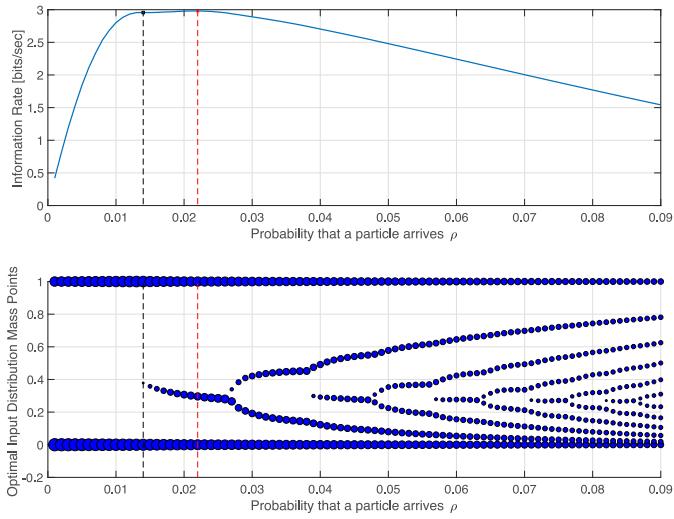


Fig. 8. Information rates and capacity-achieving input distributions as a function of ρ for the PIC with $c = 1$, $\eta = 0.2$, $\alpha = 0.9$, $\beta = 0.9$ and $\lambda = 1000$. Finite-support capacity-achieving distributions were obtained by the algorithm described in Section V. To generate the results of this graph, we use methods described in Sections V-C1b and V-E3.

the information rate is plotted for three different particle generation rates, λ . The scaling factor for the Lévy distribution is $\eta = 0.2$, which means the distance between the transmitter and the receiver is four times the radius of the receiver.

The square markers indicate the maximum value of $C^b(\rho)$ of (9) for the binary input distribution, and the \times -markers indicated C^* in (6) for the optimal input distribution. For the case of $\lambda = 1000$, Figure 7 shows that the binary input distribution does not maximize C^* . The three vertical dashed lines indicate, for each choice of λ , the ρ value after which $m\rho\theta\rho$ in Proposition 1 is greater than 3.3676. We observe that for the ρ values smaller than this critical value, the binary input is the optimal input distribution.

DAB provides the capacity and optimal mass point locations for a sequence of ρ values for the PIC channel. Unlike the binomial channel, there is no longer any assumption of symmetry in the input distribution. Therefore, the new mass points are introduced at a location between the two middle

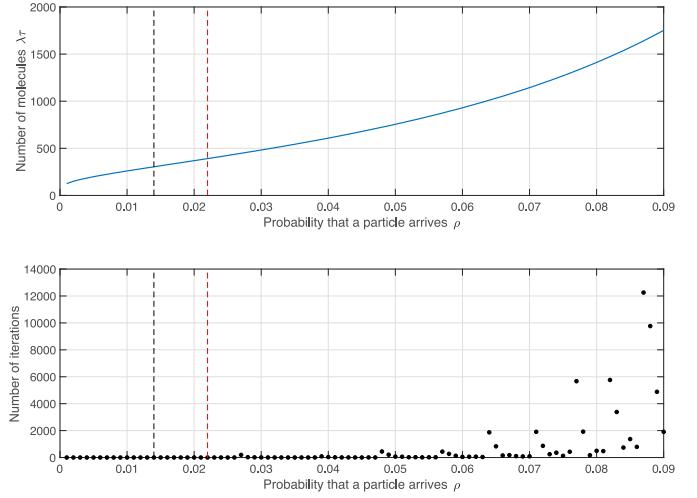


Fig. 9. Number of available particles vs. ρ and number of DAB iterations until the information difference $D(P_{Y|X=x} || P_Y) - I(X; Y)$ is below 10^{-5} as a function of ρ for the PIC with the same parameters and algorithm as in Figure 8.

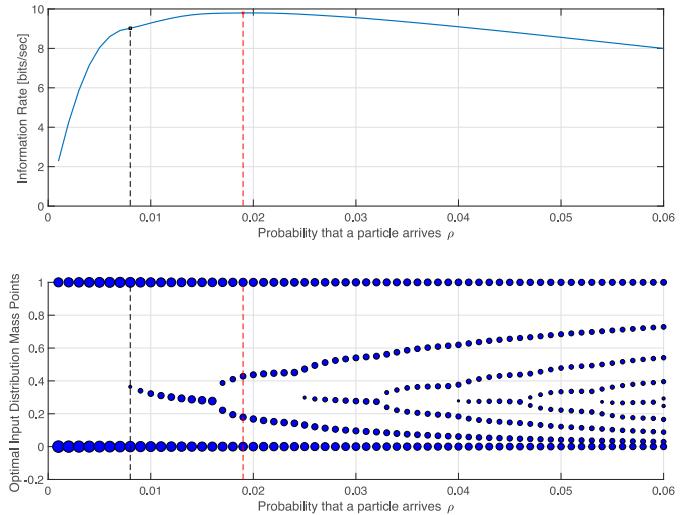


Fig. 10. Information rates and capacity-achieving input distributions as a function of ρ for the PIC with $c = 0.5$, $\eta = 0.3$, $\alpha = 0.95$, $\beta = 0.95$ and $\lambda = 5000$. Finite-support capacity-achieving distributions were obtained by the algorithm described in Section V. To generate the results of this graph, we use methods described in Section V-C1b and V-E3.

mass points, or the middle mass point is split. Figures 8 and 9 summarize the results of application of DAB to the PIC with parameters $c = 1$, $\eta = 0.2$, $\alpha = 0.9$, $\beta = 0.9$ and $\lambda = 1000$. Particularly, the top plot in Figure 8 shows the channel capacity, and the bottom plot shows the capacity achieving distribution corresponding to the each ρ value. The black dashed line indicates the first ρ value for which binary distribution is no longer capacity achieving, and the red dashed line indicates the location of the ρ value that achieves the largest information rate. Again we can verify that Proposition 1 holds, and that only 3 mass points are required to achieve the highest information rate.

The top of Figure 9 plots the number of particles that can be released by the transmitter as a function of ρ . The bottom shows the number of iterations required until

$D(P_{Y|X=x} \| P_Y) - I(X; Y)$ is below 10^{-5} . We observe that even with multiple mass points, the number of iterations can be quite small when no new mass point has been recently introduced. As can be seen in Fig. 8, in these cases the new mass point positions are small modifications to the previous mass point positions, which were used to initialize DAB. However, as the number of mass points increases the maximum number of iterations required by DAB (typically right after a mass point is introduced) does increase.

To demonstrate that the input distribution that achieves capacity can have more than 3 mass points, in Figure 10, we consider a system with $c = 0.5, \eta = 0.3, \alpha = 0.95, \beta = 0.95$ and $\lambda = 5000$. Recall that the black dashed line indicates the first ρ value for which binary distribution is no longer capacity achieving, and the red dashed line indicates the location of the ρ value that achieves the largest information rate. As can be seen, for this system, the largest information rate is achieved when the number of mass points is equal to 4. Again the results verify that Proposition 1 holds.

VII. CONCLUSION

This article introduces the PIC and analyzes its capacity and the associated capacity-achieving distribution. We show that the optimal input distribution for this channel always has mass points at probabilities of release having values of zero and one. For diffusion-based propagation, the diffusion coefficient, and hence the type of the particles used, does not affect the optimal input distribution. We then derived capacity for the binary input diffusion-based PIC and present conditions under which a binary input is optimal for this channel.

This article also introduces DAB as an efficient algorithm for finding the capacity and associated capacity-achieving distribution when the input alphabet is continuous but a capacity-achieving (or capacity-approaching) distribution is known to have finite support. This article applies DAB to the binomial channel and the PIC. DAB provides numerical results illustrating that binary input is optimal for systems where the transmitter cannot generate particles at rates that satisfy Proposition 1. This can be thought of as the low SNR regime. Future work on the PIC could explore the effect of ISI. Future applications of DAB include a wide range of channels including peak and power limited additive white Gaussian noise channels and optical channels.

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