

On amalgamation in NTP_2 theories and generically simple generics

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Abstract We prove a couple of results on NTP_2 theories. First, we prove an amalgamation statement and deduce from it that the Lascar distance over extension bases is bounded by 2. This improves previous work of Ben Yaacov and Chernikov. We propose a line of investigation of NTP_2 theories based on S1 ideals with amalgamation and ask some questions. We then define and study a class of groups with generically simple generics, generalizing NIP groups with generically stable generics.

Introduction

The class of NTP_2 theories contains both simple and NIP theories. It is probably the largest class where forking is sufficiently well behaved to be taken seriously. A couple of important facts are known: over extension bases, forking equals dividing ([CK12]) and the non-forking ideal is S1 ([BYC14]). In addition, some theorems on groups generalizing similar results for simple and NIP theories have been proved: Hempel and Onshuus [HO17] construct definable envelopes for abelian and solvable subgroups; [CKS15] studies chain conditions and [MOS18] sets the foundations for the theory of definably amenable NTP_2 groups. More recently [KS17] explores analogues of some NIP-like phenomena.

After a first section of preliminaries, the second section of this paper improves some results from [BYC14]: we give a stronger, more natural, amalgamation theorem and, using an argument from Itai Ben Yaacov, we give the optimal bound of 2 for Lascar distance over extension bases. We also speculate on a strategy for developing the theory of NTP_2 . We observe that in simple theories, one usually works with a type p over a small set and consider its non-forking extensions all at once. In NIP however, we prefer to fix some global non-forking (or invariant) extension \tilde{p} of p and study it, possibly using compactness of the space of non-forking extensions at the end of our construction to obtain a result on p itself. Our idea is that in NTP_2 , one would have to do a mixture of those two things and we suggest that the class

of non-forking S1 ideals with amalgamation will replace the ideal of non-forking extensions in simple theories and the use of global invariant types in NIP.

We then turn our attention to definable groups. In the NIP setting, groups with an invariant measure play an important role. Some results concerning them have been generalized to NTP_2 in [MOS18]. In this paper, we pursue this enterprise by generalizing a subclass: that of groups with a generically stable generic. Such groups play an important role in Hrushovski and Rideau's study of metastable groups [HR] (where they appear with the stronger property of having a stably dominated generic type). Moving from NIP to NTP_2 , the condition becomes that of having a *generically simple generic*. We show that under this assumption all forking-generic types are generically simple and non-forking over any extension base. We leave open the questions of generalizing the classes of fsg groups and compactly dominated groups. There are some natural candidates for these, but we were not able to prove convincing statements about them.

1 Preliminaries

Our notations are standard. We work in a complete theory T which has a monster model \mathcal{U} . Usually, A, B, C, \dots will denote small subsets of \mathcal{U} and M, N, \dots small submodels. The group of automorphisms of \mathcal{U} fixing A pointwise is denoted by $\text{Aut}(\mathcal{U}/A)$. We will often assume that T is NTP_2 . For definitions and basic facts about this condition see [CK12]. We will actually never use the definition of NTP_2 , but only certain properties that we recall here.

A subset A of the monster model is an *extension base* if no type $p \in S(A)$ forks over A . It is proved in [CK12] that if A is an extension base in an NTP_2 theory, then forking and dividing over A coincide.

We use the notation $a \downarrow_C b$ to mean that $\text{tp}(a/Cb)$ does not fork over C . We know from [CK12] that if C is an extension base, then this relation satisfies extension on both sides: If $a \downarrow_C b$ (resp. $a \downarrow_C b$) and d is any tuple, then there is $d' \equiv_{Ca} d$ such that $d' \downarrow_C b$ (resp. $b \downarrow_C d'$). Also, in any theory, non-forking satisfies left transitivity: if $a \downarrow_{Cd} b$ and $d \downarrow_C b$, then $ad \downarrow_C b$, as well as base monotonicity: if $a \downarrow_C bd$, then $a \downarrow_{Cd} b$.

A *Morley sequence* over A is a sequence $(a_i : i < \omega)$ which is A -indiscernible and such that $a_i \downarrow_A a_{<i}$ for all $i < \omega$. Recall that Lascar equivalence over a set A is defined as the finest bounded A -invariant equivalence relation. A class of this equivalence relation is called a Lascar strong type. We let $\text{Lstp}(a/A)$ denote the Lascar strong type of a over A and \equiv_A^L denote equality of Lascar strong types over A . If a, b have the same Lascar strong type over A , then there are $a = a_0, a_1, \dots, a_n = b$ such that (a_i, a_{i+1}) starts an A -indiscernible sequence for all i . The minimal such n is called the Lascar distance of a and b over A and denoted $d_A(a, b)$.

Fact 1.1 ([BYC14], Theorem 3.3) *Let T be NTP_2 and let A be an extension base. Assume that $c \downarrow_A ab$, $a \downarrow_A bb'$ and $b \equiv_A^L b'$. Then there is c' such that $c' \downarrow_A ab'$, $c'a \equiv_A ca c'b' \equiv_A cb$.*

The following lemma will be useful in the next section.

Lemma 1.2 *Let T be NTP_2 and let A be an extension base. Let $a \downarrow_A b$ and $c \in \mathcal{U}$. Then there is $ac' \equiv_A^L ac$ such that $ac' \downarrow_A b$.*

Proof Let C be a set of tuples of same size as c such that for any $c' \in C$, $ac' \equiv_A ac$ and the set $\{\text{Lstp}(ac'/A) : c' \in C\}$ is maximal. By maximality, it contains all possible Lascar strong types over A of some ac' , with $ac' \equiv_A ac$. By left-extension, there is $C' \equiv_{Aa} C$ such that $aC' \perp_A b$. Then $\{\text{Lstp}(ac'/A) : c' \in C'\}$ is also maximal, hence we can find the c' we are looking for in C' . \square

1.1 Measures and ideals A (Keisler) measure $\mu(x)$ on a definable set D over A is a finitely additive probability measure on A -definable subsets of D (in the variable x). Measures play an important role in NIP theories. They are useful for the general theory (for instance distality can be defined via properties of measures) and as well as to study definable groups, where translation-invariant measure are an essential tool. In NTP_2 theories, we believe a similar role will be played by ideals.

By an ideal, we always mean an ideal on the boolean algebra of definable sets over some A . If I is such an ideal, we say that a type p is I -wide if it does not imply a formula in I .

A measure $\mu(x)$ over \mathcal{U} is M -invariant if $\mu(\phi(x;a))$ depends only on $\text{tp}(a/M)$, equivalently μ is invariant under $\text{Aut}(\mathcal{U}/M)$. Similarly, an ideal is M -invariant if it is invariant under $\text{Aut}(\mathcal{U}/M)$. The following definition comes from [Hru12].

Definition 1.3 Let I be an M -invariant ideal on definable sets. We say that I is S1 if given an M -indiscernible sequence $(a_i : i < \omega)$ and a formula $\phi(x;y)$, if $\phi(x;a_0)$ is I -wide, then so is $\phi(x;a_0) \wedge \phi(x;a_1)$.

Note that one then has the stronger property that the full partial type $\{\phi(x;a_i) : i < \omega\}$ is I -wide. It follows that if I is A -invariant and S1, then I contains all formulas which divide over A .

If μ is an M -invariant measure, then we can associate to it an ideal I_μ on definable sets by defining $X \in I_\mu$ if $\mu(X) = 0$. Then I_μ is S1 (see for example [Hru12] or [Sim15, Lemma 7.5]). Another source of S1 ideals comes from the following important fact.

Fact 1.4 ([BYC14], Theorem 2.9) *Let T be NTP_2 and let A be an extension base. Then the ideal of formulas that fork over A is (A -invariant and) S1.*

2 Amalgamation

Throughout this section, we assume that T is NTP_2 . We will improve Fact 1.1. We first show that we can always amalgamate a type with itself.

Proposition 2.1 *Let A be an extension base. Let $\phi(x;a)$ be non-forking over A and assume that $b \perp_A a$ with $b \equiv_A^L a$. Then $\phi(x;a) \wedge \phi(x;b)$ is non-forking over A .*

Proof Let $b \perp_A a$, with $b \equiv_A^L a$. Build an indiscernible sequence $(b_i : i < \omega)$ in $\text{tp}(b/Aa)$ which is Morley over Aa and with $b_0 = b$. We then have $b_1 \perp_A ba$, $b \equiv_A^L a$ and by the S1 property, $\phi(x;b) \wedge \phi(x;b_1)$ is non-forking over A . By Fact 1.1, $\phi(x;a) \wedge \phi(x;b_1)$ is non-forking over A . As $\text{tp}(b_1/Aa) = \text{tp}(b/Aa)$, also $\phi(x;a) \wedge \phi(x;b)$ is non-forking over A . \square

We deduce the following strengthening of Fact 1.1.

Theorem 2.2 *Let A be an extension base and $\phi(x;y), \psi(x;y)$ over A . Let $a, b, b' \in \mathcal{U}$, $b \equiv_A^L b'$ and either $a \perp_A b'$ or $b' \perp_A a$. Assume that $\phi(x;a) \wedge \psi(x;b)$ is non-forking over A , then so is $\phi(x;a) \wedge \psi(x;b')$.*

Proof Assume first that $b' \downarrow_A a$. Let $M \supseteq A$ be a model such that $b' \downarrow_A Ma$. Find $b'' \equiv_{Ma} b'$ such that $b'' \downarrow_A Mab$. Then we have $b'' \equiv_{Aa}^L b'$ and replacing b' by b'' , we may assume that $b' \downarrow_A ab$. By Lemma 1.2, we can then find a' such that $a'b' \equiv_A^L ab$ and $a'b' \downarrow_A ab$. By the previous proposition, $(\phi(x;a) \wedge \psi(x;b)) \wedge (\phi(x;a') \wedge \psi(x;b'))$ is non-forking over A and a fortiori so is $\phi(x;a) \wedge \psi(x;b')$.

Now assume that $a \downarrow_A b'$. Let σ be a Lascar-strong automorphism over A sending b to b' and set $a' = \sigma(a)$. Then $a \equiv_A^L a'$, $\phi(x;a') \wedge \psi(x;b')$ does not fork over A and $a \downarrow_A b'$. We can then apply the previous paragraph to conclude that $\phi(x;a) \wedge \psi(x;b')$ is non-forking over A . \square

Here is another way to state the result of the theorem.

Corollary 2.3 *Let A be an extension base. Let $a, b, b', c \in \mathcal{U}$, $b \equiv_A^L b'$ and either $a \downarrow_A b'$ or $b' \downarrow_A a$. If $c \downarrow_A ab$, then there is c' such that $c' \downarrow_A ab'$, $c'a \equiv_A ca$ and $c'b' \equiv_A cb$.*

Proof Write $p(x;a) = \text{tp}(c/a)$ and $q(x;b) = \text{tp}(c/b)$. By Theorem 2.2 the partial type $p(x;a) \wedge q(x;b')$ does not fork over A . Take c' to realize a completion of that type over Aab' which is still non-forking over A . \square

The following consequence of Theorem 2.2 is due to Itai Ben Yaacov. It answers Question 3.8 from [BYC14].

Theorem 2.4 *Let A be an extension base and $b \equiv_A^L b'$. Then $d_L(b, b') \leq 2$. Furthermore, if $b \downarrow_A b'$, then b, b' start a Morley sequence over A .*

Proof It is sufficient to prove the furthermore part, since we can then take $b'' \downarrow_A bb'$ with $b'' \equiv_A^L b \equiv_A^L b'$ (by Lemma 1.2 say) and the sequence (b, b'', b') witnesses that $d_L(b, b') \leq 2$.

Fix some large enough cardinal κ and we build by induction a sequence $(b_i : i < \kappa)$ such that for each $i < \kappa$, $b_i \downarrow_A b_{<i}$ and for $i < j < \kappa$, $b_i b_j \equiv_A bb'$. We will also ensure that the sequence $\text{tp}(b_i/b_{<i})$ is increasing. Start by setting $b_0 = b$ and $b_1 = b'$. At some limit λ , let b_λ realize $\bigcup_{i < \lambda} \text{tp}(b_i/b_{<i})$. This satisfies the conditions. Assume we have constructed b_i for $i \leq \alpha$ and we look for $b_{\alpha+1}$. Let $p(x; b_{<\alpha}) = \text{tp}(b_\alpha/b_{<\alpha})$ and $q(x; b_0) = \text{tp}(b_\alpha/b_0) = \text{tp}(b'/b_0)$. Then $p(x; b_{<\alpha}) \cup q(x; b_0) \subseteq \text{tp}(b_\alpha/b_{<\alpha})$ is non-forking over A . Since $b_0 b_\alpha \equiv_A bb'$, we have $b_\alpha \equiv_A^L b_0$. Also $b_\alpha \downarrow_A b_{<\alpha}$ so by Theorem 2.2, the type $p(x; b_{<\alpha}) \cup q(x; b_\alpha)$ does not fork over A . Take $b_{\alpha+1}$ to realize it so that $b_{\alpha+1} \downarrow_A b_{\leq \alpha}$. This finishes the construction. Finally, using Erdős-Rado, we extract from the sequence $(b_i : i < \kappa)$ an indiscernible subsequence. This gives what we were looking for. \square

2.1 Some speculations and questions Let $\text{Autf}(\mathcal{U}/A)$ be the group of automorphisms of \mathcal{U} which fix every Lascar-strong type over A .

Definition 2.5 Let A be an extension base and let $B, C \subseteq \mathcal{U}$ contain A . Let $p \in S(B)$ and $q \in S(C)$ both non-forking over A . We say that p and q are compatible over A if for some/every $\sigma \in \text{Autf}(\mathcal{U}/A)$ such that either $\sigma(B) \downarrow_A C$ or $C \downarrow_A \sigma(B)$, the type $\sigma(p)(x) \cup q(x)$ is non-forking over A .

Note that the independence on the automorphism follows from Theorem 2.2.

We find the following question very appealing.

Question 2.6 Assume that p and q are M -invariant types such that $p^{(\omega)}|_M = q^{(\omega)}|_M$. Does it follow that p and q are compatible?

This is true in simple theories because any two M -invariant types having the same restriction to M are compatible. It also holds in NIP theories because the condition $p^{(\omega)}|_M = q^{(\omega)}|_M$ implies $p = q$ ([Sim15, Proposition 2.36]).

Definition 2.7 Let A be an extension base. An $\text{Autf}(\mathcal{U}/A)$ -invariant S1 ideal $\mu(x)$ has amalgamation if any two μ -wide types are compatible over A .

Example 2.8 The dual ideal of a global type non-forking over A has amalgamation over A . In an NIP theory those are the only ones since two different non-forking types are never compatible.

In a simple theory, if $p(x) \in \text{Lstp}(A)$, the ideal of formulas $\phi(x)$ such that $p(x) \wedge \phi(x)$ forks over A has amalgamation.

Let A be an extension base and for simplicity assume Lascar strong types and types over A coincide. We speculate that A -invariant S1 ideals with amalgamation could play in NTP_2 theories the same role that A -invariant types play in NIP. Of particular importance should be the minimal A -invariant S1 ideals with amalgamation. In a simple theory, there is only one such ideal: the ideal of all forking formulas. In NIP, those are precisely duals of A -invariant types since two different invariant types can never be amalgamated. In both cases, we see that those minimal ideals partition A -invariant types (in two opposite trivial ways). We ask whether this holds in all NTP_2 theories.

Question 2.9 Let A be as above. Is the compatibility relation on A -invariant types an equivalence relation? Is it the case that if μ and ν are two distinct minimal A -invariant S1 ideals with amalgamation, then no A -invariant type can be wide for both μ and ν ?

3 Groups

Recall that a group is definably amenable if it admits a translation-invariant measure, that is a measure $\mu(x)$ on G over some model M with $\mu(g \cdot X) = \mu(X)$ for any $g \in G(M)$ and M -definable set X . There is a rich theory of NIP definably amenable groups. As shown in [CS18], a group which is not definably amenable cannot admit any notion of *generic type* (this includes for instance strongly f -generic types, defined below, or types having a small orbit under translation). In NTP_2 theories, definable amenability is slightly too strong and the right condition that generalizes jointly simple groups and definably amenable NIP groups is the existence of strongly f -generic types. This was studied in [MOS18] and we recall the main results here.

Let G be a group definable in an NTP_2 structure M .

Definition 3.1 A global type $p \in S_G(\mathcal{U})$ is *strongly (left) f -generic* over A if for all $g \in G(\mathcal{U})$, $g \cdot p$ does not fork over A .

It is *strongly bi- f -generic* if for all $g, h \in G(\mathcal{U})$, $g \cdot p \cdot h$ does not fork over A .

If G admits a strongly f -generic type over some extension base A , then it admits a bi- f -generic type over any extension base. When this is the case, we say that G has strong f -generics. Any group definable in a simple theory is such, as is any definably amenable NTP_2 group. In NIP, this condition is equivalent to definable amenability.

Definition 3.2 Let $\phi(x) \in L(A)$ be a formula. We say that $\phi(x)$ is *f-generic* over A if no (left) translate of $\phi(x)$ forks over A . We say that $\phi(x)$ *G-divides* over A if for some A -indiscernible sequence $(g_i : i < \omega)$ of elements of G , the partial type $\{g_i \cdot \phi(x) : i < \omega\}$ is inconsistent.

Fact 3.3 ([MOS18], Lemma 3.7) Let A be an extension base and $\phi(x) \in L(A)$. Then $\phi(x)$ is *f-generic* over A if and only if it does not *G-divide* over A .

Fix some model M and let μ_M be the ideal of formulas which do not extend to a global type strongly *f-generic* over M . A definable set is *wide* if it does not lie in μ_M (i.e., if it extends to a global type strongly *f-generic* over M). A type is *wide* if all formulas in it are *wide*. A type over M is *wide* precisely if it is *f-generic*, that is all formulas in it are *f-generic*.

For any wide type p , let $St_{\mu_M}(p) = \{g \in G : gp \cup p \text{ is wide}\}$. We have the following stabilizer theorem.

Fact 3.4 ([MOS18], Theorem 3.18) Assume that G has strong *f-generics*. Let $p \in S_G(M)$ be *f-generic* and define μ_M as above.

Then $G_M^{00} = G_M^\infty = St_{\mu_M}(p)^2 = (pp^{-1})^2$ and $G_M^{00} \setminus St_{\mu_M}(p)$ is contained in a union of non- μ_M -wide M -definable sets.

3.1 Generically simple types In [Che14], Chernikov defines *simple types* in NTP_2 theories (see Definition 6.1 there). We define here a weaker notion of generically simple types. We prove their basic properties following essentially the arguments in [Che14].

Definition 3.5 Let A be any set and $p \in S(A)$. We say that p is *generically simple* if for every $b \in \mathcal{U}$ and $a \models p$, $b \downarrow_A a \implies a \downarrow_A b$.

If $A \subseteq B$, $p \in S(B)$ does not fork over A and $p|_A$ is generically simple, we say that p is *generically simple over A* .

Lemma 3.6 Assume that $\text{tp}(a/A)$ is generically simple and $b \in \text{dcl}(Aa)$, then $\text{tp}(b/A)$ is generically simple.

Proof Let $c \downarrow_A b$. By taking a non-forking extension of $\text{tp}(c/Ab)$ to Aa , we may assume that $c \downarrow_A a$. Then $a \downarrow_A c$ and in particular $b \downarrow_A c$. \square

Lemma 3.7 If $p, q \in S(A)$ are generically simple, $a \models p$, $a' \models q$ with $a \downarrow_A a'$, then $\text{tp}(a, a'/A)$ is generically simple.

Proof Let $b \downarrow_A aa'$. We have $a \downarrow_A a'$, hence by left transitivity, $ba \downarrow_A a'$. As $\text{tp}(a'/A)$ is generically simple, $a' \downarrow_A ba$. On the other hand, we have $b \downarrow_A a$, hence $a \downarrow_A b$ by generic simplicity of $\text{tp}(a/A)$. Therefore by left transitivity again, $aa' \downarrow_A b$ as required. \square

Lemma 3.8 Assume that $\text{tp}(a/A)$ is generically simple and $\text{tp}(b/Aa)$ is generically simple. Then $\text{tp}(ab/A)$ is generically simple.

Proof Let $c \downarrow_A ab$. Then $c \downarrow_{Aa} b$ and hence $b \downarrow_{Aa} c$. On the other hand, as $\text{tp}(a/A)$ is generically simple, $a \downarrow_A c$. By left transitivity, $ab \downarrow_A c$ as required. \square

Lemma 3.9 Let $(a_i : i < n)$ be tuples, possibly of different sizes. Assume that for each i , $\text{tp}(a_i/A)$ is generically simple and $a_i \downarrow_A a_{<i}$. Then for any two disjoint subsets $I, J \subseteq n$, we have $a_I \downarrow_A a_J$ (where $a_I = (a_i : i \in I)$).

Proof First, by Lemma 3.7 and induction on $|I|$, $\text{tp}(a_I/A)$ is generically simple for all $I \subseteq n$. Fix $k \leq n$. Then we have $a_{\geq k} \downarrow_A a_{<k}$ and $a_{<k} \downarrow_A a_{\geq k}$ by generic simplicity. It follows that $a_{<k} \downarrow_{Aa_{>k}} a_k$. As also $a_{>k} \downarrow_A a_k$, we have by left transitivity $a_{\neq k} \downarrow_A a_k$ and $a_k \downarrow_A a_{\neq k}$ by generic simplicity. This shows that the hypothesis of the lemma is stable under permutation of the indices of the a_i 's. The result then follows from the fact that $a_{\leq k} \downarrow_A a_{>k}$ for all $k < n$. \square

Note that the lemma goes through with infinitely many tuples.

Proposition 3.10 *Let A be an extension base, and assume that $p \in S(A)$ is generically simple. Let $a \downarrow_A b$ where $a \models p$. Then $b \downarrow_A a$.*

Proof The proof of [Che14, Section 6.2] of the analogue result for simple types goes through using the lemmas above. More precisely, Lemma 6.1 there follows from Fact 1.4 (with no assumption of simplicity). In Lemma 6.13, simplicity is only used for checking that property (3) holds. Note that each sequence \bar{b}_i realizes a generically simple type over the base A by Lemma 3.7. We then have by construction $\bar{b}_i \downarrow_{a_{>i}} \bar{b}_{<i}$. So the sequence $(\bar{b}_0, \dots, \bar{b}_i, a_{i+1}, a_{i+2}, \dots)$ satisfies the hypothesis of Lemma 3.9 and we conclude $a_{>i+1} \bar{b}_{\leq i} \downarrow_{a_{i+1}}$ as required. Lemma 6.14 only uses generic simplicity, then Proposition 6.15 goes through unchanged. \square

Corollary 3.11 *If $q \in S(B)$ is generically simple over A , A an extension base, then it is generically simple itself.*

In fact, if $a \models q$ and $b \downarrow_B a$, then $a \downarrow_A Bb$.

Proof Let b with $b \downarrow_B a$. We have $B \downarrow_A a$ as $\text{tp}(a/A)$ is generically simple and $a \downarrow_A B$ by assumption. By left transitivity, $Bb \downarrow_A a$, hence $a \downarrow_A Bb$. In particular, $a \downarrow_B b$, which shows that q is generically simple. \square

Lemma 3.12 *Assume that $\text{tp}(a/A)$ is generically simple. Let $A \subseteq B \subseteq C$ with $a \downarrow_A B$ and $a \downarrow_B C$. Then $a \downarrow_A C$.*

Proof As $a \downarrow_A B$, Corollary 3.11 implies that $\text{tp}(a/B)$ is generically simple. Therefore we have $C \downarrow_B a$. Then again by Corollary 3.11, $a \downarrow_A C$. \square

The following is the analogue of Problem 6.6 in [Che14].

Question 3.13 *Assume that $q \in S(B)$ is generically simple and does not fork over A , then is $q|_A$ generically simple?*

Note that by Lemma 3.8, this is true if $\text{tp}(B/A)$ is generically simple.

3.2 Generically simple generics We now define a notion of generically simple generic type, which is stronger than that of strong f-generics. We will then prove that if a definable group admits such a type, then all its f-generic types are such and do not fork over any extension base, similarly to what happens with groups in simple theories. In particular, in those groups, any global f-generic type is a strong f-generic.

In what follows, G is again a group definable in an NTP_2 structure; $S_G(A)$ denotes the space of types over A that concentrate on G .

We adopt the convention that if $a, b \in G$, then ab always denotes the product $a \cdot b$ (as opposed to concatenation of tuples, which will be denoted by $\hat{a}\hat{b}$).

Definition 3.14 A type $p \in S_G(A)$ is generically simple generic (gsg) if p is generically simple and for any $B \supseteq A$ and $b \in G$, we have $a \downarrow_A Bb \implies b \cdot a \downarrow_A Bb$.

The type p will be said to be a two-sided gsg if it is generically simple and for any $B \supseteq A$ and $b, c \in G$, we have $a \downarrow_A Bb^{\wedge}c \implies b \cdot a \cdot c \downarrow_A Bb^{\wedge}c$.

Note that if p is generically simple generic, then any non-forking extension of p is again generically simple generic. (If say $a \models p$, $a \downarrow_A B$ and $a \downarrow_B b$, then $a \downarrow_A Bb$ by Lemma 3.12.)

Now if p is gsg and \tilde{p} is a non-forking extension of p to \mathcal{U} , then no translate $g \cdot \tilde{p}$, $g \in G(\mathcal{U})$ forks over A . Hence \tilde{p} is strongly f-generic over A .

Lemma 3.15 Let A be an extension base and $\text{tp}(a/A)$ be gsg, $a \downarrow_A b$, then $\text{tp}(ba/A)$ is generically simple.

Proof Let $d \downarrow_A ba$. We can assume furthermore that $d \downarrow_A b\hat{a}b^{\wedge}a$. Hence $d \downarrow_{Ab} a$. As $a \downarrow_A b$, by Corollary 3.11, $a \downarrow_A b^{\wedge}d$. As $\text{tp}(a/A)$ is gsg, we conclude $ba \downarrow_A d$ as required. \square

It will follow from the statements proved below that in fact $\text{tp}(ba/A)$ is gsg (take a model M such that $a \downarrow_A Mb$, then $\text{tp}(ba/M)$ is f-generic and by Proposition 3.19, $\text{tp}(ba/A)$ is gsg).

Lemma 3.16 Let $p \in S_G(A)$ be two-sided gsg and $b, c \in A$. Then $\text{tp}(bac/A)$ is two-sided gsg.

Proof Assume that $bac \downarrow_A B\hat{d}d'$. Then $a \downarrow_A Bb^{\wedge}c^{\wedge}d^{\wedge}d'$ since $b, c \in A$. Therefore $dbacd' \downarrow_A Bb^{\wedge}c^{\wedge}d^{\wedge}d'$ as $\text{tp}(a/A)$ is two-sided gsg. \square

Lemma 3.17 Let A be an extension base, $p \in S_G(A)$ be gsg and take $a, b \models p$, $a \downarrow_A b$. Then $\text{tp}(ab^{-1}/A)$ is two-sided gsg.

Proof By Lemma 3.7, $\text{tp}(a, b/A)$ is generically simple, therefore so is $\text{tp}(ab^{-1}/A)$. Now let $B \supseteq A$ and $c, d \in G$ such that $ab^{-1} \downarrow_A Bc^{\wedge}d$. By left extension, we may assume that $a^{\wedge}b \downarrow_A Bc^{\wedge}d$. Then $Bc^{\wedge}d \downarrow_A a^{\wedge}b$ by generic simplicity. Since also $b \downarrow_A a$, by left transitivity, $Bb^{\wedge}c^{\wedge}d \downarrow_A a$ and therefore $ca \downarrow_A Bb^{\wedge}c^{\wedge}d$ as $\text{tp}(a/A)$ is gsg. Similarly, $db \downarrow_A Ba^{\wedge}c^{\wedge}d$. By left transitivity, $ca, db \downarrow_A Bc^{\wedge}d$ and in particular $cab^{-1}d \downarrow_A Bc^{\wedge}d$ as required. \square

Proposition 3.18 Let $A \subseteq M$ an extension base. Assume that $p \in S_G(A)$ and $q \in S_G(M)$ are two-sided gsg. Then q does not fork over A and $q|_A$ is (two-sided) gsg.

Proof Let $M \prec N$, N sufficiently saturated. Recall that μ_M is the ideal of formulas which do not extend to a strongly f-generic type. It is invariant by translation on the left. Let $d_0 \in N$ be a realization of p and take $e_0 \in N$, such that $d_0 \downarrow_A e_0$ and $d := e_0 d_0$ lies in the same G_M^{00} -coset as q^{-1} . By Lemma 3.15, $\text{tp}(d/A)$ is generically simple. Let $b \models q$, with $\text{tp}(b/N)$ μ_M -wide and set $b' = db$. Then $\text{tp}(b'/N)$ is μ_M -wide and lies in G_M^{00} .

Let p' be a non-forking extension of p to M . Then p' is generically simple over A and gsg. By Fact 3.4, the type $b'p' \cup p'$ is μ_M -wide, hence so is $p' \cup b'^{-1}p'$ by left-invariance of μ_M . Let a realize the latter type such that $\text{tp}(a/Nb)$ is μ_M -wide. In

particular $a \downarrow_M Nb$ and $Nb \downarrow_M a$ as $\text{tp}(a/M) = p'$ is generically simple. By Corollary 3.11,

$$(0) \quad a \downarrow_A Nb.$$

As $\text{tp}(a/M)$ is gsg over A , we deduce

$$(1) \quad b'a \downarrow_A Nb.$$

On the other hand, we deduce from $a \downarrow_M Nb$ that $a \downarrow_M Nb'$ and then $a \downarrow_N b'$ by base monotonicity. As $\text{tp}(b'/N)$ is two-sided gsg by Lemma 3.16, it follows that

$$(2) \quad b'a \downarrow_N a.$$

By construction, $\text{tp}(b'a/M) = p'$ and thus $\text{tp}(b'a/A)$ is generically simple. By (1), (2) and Lemma 3.12, $b'a \downarrow_A Na$. Now $\text{tp}(b'a/A) = p$ is two-sided gsg, so we have $b' \downarrow_A Na$ and $b \downarrow_A Na$. In particular, $b \downarrow_A M$. Therefore q does not fork over A .

It remains to see that $\text{tp}(b/A)$ is two-sided gsg. We first show that it is generically simple. We know that $\text{tp}(b'a/A)$ is generically simple and $b'a \downarrow_A Na$, hence $b'a \downarrow_A d'a$. We deduce that $\text{tp}(b'a/Ad'a)$ is generically simple, hence so is $\text{tp}(b/Aa)$ by Lemma 3.6. Since $a \downarrow_A N$, $\text{tp}(d'a/A)$ is generically simple and Lemma 3.8 gives that $\text{tp}(b/A)$ is generically simple.

Now assume that $b \downarrow_A Bc\hat{c}'$ and we want to show that $cbc' \downarrow_A Bc\hat{c}'$. By generic simplicity, $M \downarrow_A b$. Moving M over Ab , we may assume that $M \downarrow_A Bb\hat{c}'c'$. As $Bc\hat{c}' \downarrow_A b$, by left transitivity, $MBc\hat{c}' \downarrow_A b$. It follows that $b \downarrow_M Bc$. As $\text{tp}(b/M)$ is gsg, $cbc' \downarrow_M Bc\hat{c}'$. Now $\text{tp}(cbc'/Mc\hat{c}')$ is also two-sided gsg by Lemma 3.16, hence by what we have already proved, it is generically simple over A . We conclude that $cbc' \downarrow_A MBc\hat{c}'$. \square

Proposition 3.19 *Let $A \subseteq M$ an extension base. Assume that $p \in S_G(A)$ is two-sided gsg and $q \in S_G(M)$ is f -generic. Then q does not fork over A and $q|_A$ is two-sided gsg.*

Proof Let $a \models p$ such that $a \downarrow_A M$ and let $b \models q$ with $\text{tp}(b/Ma)$ μ_M -wide. Then $ab \downarrow_M a$ and also $ab \downarrow_A Mb$. Since $b \downarrow_M a$ and $\text{tp}(a/M)$ is generically simple over A , we have $a \downarrow_A Mb$. If N is a model containing Mb such that $a \downarrow_A N$, then by Lemma 3.16, $\text{tp}(ab/N)$ is gsg. By Proposition 3.18, it does not fork over A and is gsg over A . Hence $ab \downarrow_M a$ implies $ab \downarrow_A Ma$ and then $b \downarrow_A Ma$ by gsg. In particular, $b \downarrow_A M$ and q does not fork over A . It remains to see that $q|_A$ is two-sided gsg. We know that $\text{tp}(ab/Ma)$ is two-sided gsg, and then so is $\text{tp}(b/Ma)$. By the previous proposition, $\text{tp}(b/A)$ is two-sided gsg. \square

Corollary 3.20 *Assume that G has a gsg type and let A be an extension base. Then any f -generic type of G is non-forking over A .*

Proof Assume that G has a gsg type over some $B \supseteq A$. Let $p \in S_G(N)$ be an f -generic type. By Proposition 3.19, it is gsg over B . Let $a \models p$ and take $B' \downarrow_A aN$, $B' \equiv_A B$. As $\text{tp}(a/N)$ is generically simple over B , we have $a \downarrow_B NB'$. As B' is a conjugate of B , there is a gsg type over B' and therefore also a two-sided gsg type. Proposition 3.19 implies that $a \downarrow_{B'} N$. As $B' \downarrow_A N$, we have $a \downarrow_A N$. \square

Question 3.21 Given G an arbitrary definable groups in an NTP_2 theory, assume that $p \in S_G(A)$ is generically simple and f -generic, then is it gsg?

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