

# The Schur-Erdős problem for semi-algebraic colorings

Jacob Fox\*

János Pach<sup>†</sup>

Andrew Suk<sup>‡</sup>

## Abstract

We consider  $m$ -colorings of the edges of a complete graph, where each color class is defined semi-algebraically with bounded complexity. The case  $m = 2$  was first studied by Alon et al., who applied this framework to obtain surprisingly strong Ramsey-type results for intersection graphs of geometric objects and for other graphs arising in computational geometry. Considering larger values of  $m$  is relevant, e.g., to problems concerning the number of distinct distances determined by a point set.

For  $p \geq 3$  and  $m \geq 2$ , the classical Ramsey number  $R(p; m)$  is the smallest positive integer  $n$  such that any  $m$ -coloring of the edges of  $K_n$ , the complete graph on  $n$  vertices, contains a monochromatic  $K_p$ . It is a longstanding open problem that goes back to Schur (1916) to decide whether  $R(p; m) \leq 2^{cm}$ , where  $c = c(p)$ . We prove that this is true if each color class is defined semi-algebraically with bounded complexity, and that the order of magnitude of this bound is tight. Our proof is based on the Cutting Lemma of Chazelle *et al.*, and on a Szemerédi-type regularity lemma for multicolored semi-algebraic graphs, which is of independent interest. The same technique is used to address the semi-algebraic variant of a more general Ramsey-type problem of Erdős and Shelah.

## 1 Introduction

The Ramsey number  $R(p; m)$  is the smallest integer  $n$  such that any  $m$ -coloring on the edges of the complete  $n$ -vertex graph contains a monochromatic copy of  $K_p$ . The existence of  $R(p; m)$  follows from the celebrated theorem of Ramsey [20] from 1930, and for the special case when  $p = 3$ , Issai Schur proved the existence of  $R(3; m)$  in 1916 in his work related to Fermat’s Last Theorem [21]. He showed that

$$\Omega(2^m) \leq R(3; m) \leq O(m!).$$

While the upper bound has remained unchanged over the last 100 years, the lower bound was successively improved and the current record is  $R(3; m) = \Omega(3.199^m)$  due to Xiaodong *et al.* [25]. It is a major open problem in Ramsey theory, for which Erdős offered some prize money, to close the gap between the lower and upper bounds for  $R(3; m)$ .

In this paper, we study edge-colorings of complete graphs where each color class is defined algebraically with bounded complexity. Over the last decade, several researchers have shown that

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\*Stanford University, Stanford, CA. Supported by a Packard Fellowship and by NSF grant DMS-1855635. Email: [jacobfox@stanford.edu](mailto:jacobfox@stanford.edu).

<sup>†</sup>EPFL, Lausanne and Rényi Institute, Budapest. Supported by Swiss National Science Foundation Grants 200020-162884 and 200021-175977. Email: [pach@cims.nyu.edu](mailto:pach@cims.nyu.edu).

<sup>‡</sup>Department of Mathematics, University of California at San Diego, La Jolla, CA, 92093 USA. Supported by an NSF CAREER award and by an Alfred Sloan Fellowship. Email: [asuk@ucsd.edu](mailto:asuk@ucsd.edu).

some of the classical theorems in extremal combinatorics can be significantly improved if the underlying graphs are intersection graphs of geometric objects of bounded “description complexity” or bounded VC-dimension, graphs of incidences between points and hyperplanes, distance graphs, or, more generally, semi-algebraic graphs [1, 11, 5, 10, 22]. To make this statement more precise, we need to introduce some terminology. Let  $V$  be an ordered point set in  $\mathbb{R}^d$ , and let  $E \subset \binom{V}{2}$ . We say that  $E$  is a *semi-algebraic* relation on  $V$  with *complexity* at most  $t$  if there are at most  $t$  polynomials  $g_1, \dots, g_s \in \mathbb{R}[x_1, \dots, x_{2d}]$ ,  $s \leq t$ , of degree at most  $t$  and a Boolean formula  $\Phi$  such that for vertices  $u, v \in V$  such that  $u$  comes before  $v$  in the ordering,

$$(u, v) \in E \quad \Leftrightarrow \quad \Phi(g_1(u, v) \geq 0; \dots; g_s(u, v) \geq 0) = 1.$$

At the evaluation of  $g_\ell(u, v)$ , we substitute the variables  $x_1, \dots, x_d$  with the coordinates of  $u$ , the variables  $x_{d+1}, \dots, x_{2d}$  with the coordinates of  $v$ . We may assume that the semi-algebraic relation  $E$  is *symmetric*, i.e., for all points  $u, v \in \mathbb{R}^d$ ,  $(u, v) \in E$  if and only if  $(v, u) \in E$ . Indeed, given such an ordered point set  $V \subset \mathbb{R}^d$  and a not necessarily symmetric semi-algebraic relation  $E$  of complexity at most  $t$ , we can define  $V^* \subset \mathbb{R}^{d+1}$  with points  $(v, i)$  where  $v \in V$  and  $v$  is the  $i$ th smallest element in the given ordering of  $V$ . Then we can define a symmetric semi-algebraic relation  $E^*$  on the pairs of  $V^*$  with complexity at most  $2t + 2$ , by comparing the value of the last coordinates of the two points, and checking the relation  $E$  using the first  $d$  coordinates of the two points. We will therefore assume throughout this paper that all semi-algebraic relations we consider are symmetric, and the vertices are not ordered. Hence, all edges are unordered and we denote  $uv = \{u, v\}$ . We also assume that the dimension  $d$  and complexity  $t$  are fixed parameters, and  $n = |V|$  tends to infinity.

Let  $R_{d,t}(p; m)$  be the minimum  $n$  such that every  $n$ -element point set  $V$  in  $\mathbb{R}^d$  equipped with  $m$  semi-algebraic binary relations (edge-colorings)  $E_1, \dots, E_m \subset \binom{V}{2}$ , each of complexity at most  $t$ , where  $E_1 \cup \dots \cup E_m = \binom{V}{2}$ , contains a subset  $S \subset V$  of size  $p$  such that  $\binom{S}{2} \subset E_k$  for some  $k$ . Note that the relations  $E_i$  are not necessarily disjoint, that is, an edge  $uv$  may have several colors. Clearly  $R_{d,t}(p; m) \leq R(p; m)$ . It was known that for fixed  $d, t \geq 1$ ,  $R_{d,t}(3; m) = 2^{O(m \log \log m)}$ , which is much smaller than Schur’s bound  $R(3; m) = O(m!)$  mentioned in the first paragraph of the Introduction; see [22]. In this paper, we completely settle Schur’s problem for semi-algebraic graphs, by showing that in this setting Schur’s lower bound (which is semi-algebraic with bounded complexity) is tight in the sense that these Ramsey numbers grow exponentially in the number of colors. In fact, we prove this in a more general form, for any  $p \geq 3$ .

**Theorem 1.1.** *For integers  $d, t \geq 1$  and  $p \geq 3$ , there is a constant  $c = c(d, t, p)$  such that*

$$R_{d,t}(p; m) \leq 2^{cm}.$$

Our proof uses geometric techniques and is based on the Cutting Lemma of Chazelle, Edelsbrunner, Guibas, and Sharir [3] described in Section 2.

Edge-colorings of semi-algebraic graphs with  $m$  colors can be used, e.g., for studying problems concerning the number of distinct distances determined by a point set; see [12]. One can explore the fact that multicolored semi-algebraic graphs have a very nice structural characterization, reminiscent of Szemerédi’s classic regularity lemma for general graphs [24], but possessing much stronger homogeneity properties. Our next theorem provides such a characterization, which is of independent interest. To state our result, we need some notation and terminology.

A partition is called *equitable* if any two parts differ in size by at most one. According to Szemerédi's lemma, for every  $\varepsilon > 0$  there is a  $K = K(\varepsilon)$  such that the vertex set of every graph has an equitable partition into at most  $K$  parts such that all but at most an  $\varepsilon$ -fraction of the pairs of parts are  $\varepsilon$ -regular.<sup>1</sup> It follows from Szemerédi's proof that  $K(\varepsilon)$  may be taken to be an exponential tower of 2s of height  $\varepsilon^{-O(1)}$ . Gowers [14] used a probabilistic construction to show that such an enormous bound is indeed necessary.

Alon *et al.* [1] (see also Fox, Gromov *et al.* [9]) established a strengthening of the regularity lemma for point sets in  $\mathbb{R}^d$  equipped with a semi-algebraic relation  $E$ . It was shown in [1] that for any semi-algebraic graph of bounded complexity defined on the vertex set  $V \subset \mathbb{R}^d$  (that is, for any semi-algebraic binary relation  $E \subset \binom{V}{2}$ ),  $V$  has an equitable partition into a bounded number of parts such that all but at most an  $\varepsilon$ -fraction of the pairs of parts  $(V_1, V_2)$  behave not only regularly, but *homogeneously* in the sense that either  $V_1 \times V_2 \subseteq E$  or  $V_1 \times V_2 \cap E = \emptyset$ . The first proof of this theorem was essentially qualitative: it gave a poor estimate for the number of parts in such a partition. Fox, Pach, and Suk [11] gave a stronger quantitative form of this result, showing that the number of parts can be taken to be polynomial in  $1/\varepsilon$ .

Let  $V$  be an  $n$ -element point set in  $\mathbb{R}^d$  equipped with  $m$  semi-algebraic relations  $E_1, \dots, E_m$  such that  $E_1 \cup \dots \cup E_m = \binom{V}{2}$  of bounded complexity. In other words, suppose that the edges of the complete graph on  $V$  are colored with  $m$  colors, where each color class is semi-algebraic. Then, for any  $\varepsilon > 0$ , an  $m$ -fold repeated application of the result of Fox, Pach, and Suk [11] gives an equitable partition of  $V$  into at most  $K \leq (1/\varepsilon)^{cm}$  parts such that all but an  $\varepsilon$ -fraction of the pairs of parts are complete with respect to some relation  $E_k$ , i.e., all edges between the two parts are of color  $k$ , for some  $k$ . In Section 4, we strengthen this result by showing that the number of parts can be taken to be *polynomial* in  $m/\varepsilon$ .

**Theorem 1.2.** *For any positive integers  $d, t \geq 1$  there exists a constant  $c = c(d, t) > 0$  with the following property. Let  $0 < \varepsilon < 1/2$  and let  $V$  be an  $n$ -element point set in  $\mathbb{R}^d$  equipped with semi-algebraic relations  $E_1, \dots, E_m$  such that each  $E_k$  has complexity at most  $t$  and  $\binom{V}{2} = E_1 \cup \dots \cup E_m$ . Then  $V$  has an equitable partition  $V = V_1 \cup \dots \cup V_K$  into at most  $4/\varepsilon \leq K \leq (m/\varepsilon)^c$  parts such that all but an  $\varepsilon$ -fraction of the pairs of parts are complete with respect to some relation  $E_k$ .*

In Section 5, we apply this result to solve a problem of Erdős and Shelah [6] in the semi-algebraic setting. Let  $d, t, p, q, n$  be positive integers,  $p \geq 3$ , and  $2 \leq q \leq \binom{p}{2}$ . Let  $f_{d,t}(n, p, q)$  be the minimum  $m$  such that there exists a semi-algebraic  $m$ -coloring of the edges of the complete graph of  $n$  vertices (with parameters  $d$  and  $t$ , as above) with the property that each edge has exactly one color and any set of  $p$  vertices induces at least  $q$  distinct colors. Notice that here we must assume that the color classes  $E_i$  are disjoint, since otherwise  $f_{d,t}(n, p, q) = q$  by assigning all  $q$  colors to every edge. Our next theorem precisely determines the smallest  $q$  for a given  $p$ , where  $f_{d,t}(n, p, q)$  changes from a log  $n$  to a power of  $n$ .

**Theorem 1.3.** *For fixed integers  $d, t \geq 1$ , there is a  $c = c(d, t) > 0$  such that for  $p \geq 3$ , we have*

$$f_{d,t}(n, p, \lceil \log p \rceil + 1) = \Omega \left( n^{\frac{1}{c \log^2 p}} \right).$$

Moreover, for  $t \geq 4$ ,

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<sup>1</sup>For a pair  $(V_i, V_j)$  of vertex subsets,  $e(V_i, V_j)$  denotes the number of edges in the graph running between  $V_i$  and  $V_j$ . The density  $d(V_i, V_j)$  is defined as  $\frac{e(V_i, V_j)}{|V_i||V_j|}$ . The pair  $(V_i, V_j)$  is called  $\varepsilon$ -regular if for all  $V'_i \subset V_i$  and  $V'_j \subset V_j$  with  $|V'_i| \geq \varepsilon|V_i|$  and  $|V'_j| \geq \varepsilon|V_j|$ , we have  $|d(V'_i, V'_j) - d(V_i, V_j)| \leq \varepsilon$ .

$$f_{d,t}(n, p, \lceil \log p \rceil) = O(\log n).$$

In [12], we studied a geometric instance of this problem, where every set of  $p$  points induces at least  $q$  *distinct distances*.

Our paper is organized as follows. In the next section, we describe the Cutting Lemma of Chazelle *et al.*, which is the main geometric tool used in all proofs. In Section 3, we establish Theorem 1.1. Section 4 contains the proof of our multicolored semi-algebraic regularity lemma, Theorem 1.2, which is applied in the following section to deduce Theorem 1.3. We end this paper with some concluding remarks.

## 2 The cutting lemma

The main tool we use to prove Theorems 1.1 and 1.2 is commonly referred to as the *cutting lemma*, which we now recall. A set  $\Delta \subset \mathbb{R}^d$  is *semi-algebraic* if there are polynomials  $g_1, \dots, g_t$  and a boolean formula  $\Phi$  such that

$$\Delta = \{x \in \mathbb{R}^d : \Phi(g_1(x) \geq 0; \dots; g_t(x) \geq 0) = 1\}.$$

We say that a semi-algebraic set in  $d$ -space has *description complexity* at most  $t$  if the number of inequalities is at most  $t$ , and each polynomial  $g_i$  has degree at most  $t$ . Let  $\sigma \subset \mathbb{R}^d$  be a *surface* in  $\mathbb{R}^d$ , that is,  $\sigma$  is the zero set of some polynomial  $h \in \mathbb{R}[x_1, \dots, x_d]$ . The *degree* of a surface  $\sigma = \{x \in \mathbb{R}^d : h(x) = 0\}$  is the degree of the polynomial  $h$ . We say that the surface  $\sigma \subset \mathbb{R}^d$  *crosses* a semi-algebraic set  $\Delta$  if  $\sigma \cap \Delta \neq \emptyset$  and  $\Delta \not\subset \sigma$ .

Let  $\Sigma$  be a collection of surfaces in  $\mathbb{R}^d$ , each having bounded degree. A  $(1/r)$ -*cutting* for  $\Sigma$  is a family  $\Psi$  of disjoint (possibly unbounded) semi-algebraic sets of bounded complexity such that

1. each  $\Delta \in \Psi$  is crossed by at most  $|\Sigma|/r$  surfaces from  $\Sigma$ , and
2. the union of all  $\Delta \in \Psi$  is  $\mathbb{R}^d$ .

In [3], Chazelle *et al.* (see also [15]) proved the following.

**Lemma 2.1** (Cutting lemma). *Let  $\Sigma$  be a multiset of  $N$  surfaces in  $\mathbb{R}^d$ , each surface having degree at most  $t$ , and let  $r$  be an integer parameter such that  $1 \leq r \leq N$ . Then there is a constant  $c_1 = c_1(d, t)$  such that  $\Sigma$  admits a  $(1/r)$ -cutting  $\Psi$ , where  $|\Psi| \leq c_1 r^{2d}$ , and each semi-algebraic set  $\Delta \in \Psi$  has complexity at most  $c_1$ .*

We note that the original statement of Chazelle *et al.* [3] and Koltun [15] is stronger. Namely, they also guarantee that the number of cells in the cutting  $\Psi$  is at most  $r^{2d-4+\epsilon}$  for  $d \geq 4$ . Here, for simplicity, we use the weaker bound of  $c_1 r^{2d}$ , as stated above.

## 3 Multicolor Ramsey numbers for small cliques

### –Proof of Theorem 1.1

Theorem 1.1 will easily follow from Theorem 3.1 below. For integers  $p_1, \dots, p_m \geq 2$ ,  $d, t \geq 1$ , let  $R_{d,t}(p_1, \dots, p_m)$  be the minimum integer  $n$  with the following property. Every complete graph  $K_n$ ,

whose  $n$  vertices lie in  $\mathbb{R}^d$  and whose edges are colored with  $m$  colors such that each color class is defined by a semi-algebraic relation of description complexity  $t$ , contains a monochromatic copy of  $K_{p_k}$  in color  $k$  for some  $1 \leq k \leq m$ .

**Theorem 3.1.** *For any  $d, t \geq 1$  and  $p \geq 3$ , there exists a constant  $c = c(d, t, p)$  satisfying the following condition. For any  $m$  integers  $p_1, \dots, p_m \leq p$ , we have*

$$R_{d,t}(p_1, \dots, p_m) \leq 2^{c \sum_{k=1}^m p_k}.$$

*Proof.* Fix  $d, t \geq 1$ ,  $p \geq 3$  and set  $c = c(d, t, p)$  to be a large constant that will be determined later. We will show that  $R_{d,t}(p_1, \dots, p_m) \leq 2^{c \sum_{k=1}^m p_k}$  by induction on  $s = \sum_{k=1}^m p_k$ . The base case  $s \leq 10 \cdot 2^{10dt}$  follows for  $c$  sufficiently large.

Now assume that the statement holds for  $s' < s$ . Set  $n = 2^{cs}$  and let  $V$  be an  $n$ -element point set in  $\mathbb{R}^d$  equipped with semi-algebraic relations  $E_1, \dots, E_m \subset \binom{V}{2}$  such that  $\binom{V}{2} = E_1 \cup \dots \cup E_m$  and each  $E_k$  has complexity at most  $t$ . Recall that an edge  $uv$  may have several colors. We will show that there is a subset  $S \subset V$  of size  $p_k$  such that  $\binom{S}{2} \subset E_k$  for some  $k \in [m]$ , in other words, we will find a monochromatic copy of  $K_{p_k}$  in color  $k$  for some  $k \in [m]$ . Throughout the proof, we will let  $c_1$  be as defined in Lemma 2.1.

For each relation  $E_k$ , there are  $t$  polynomials  $g_{k,1}, \dots, g_{k,t}$  of degree at most  $t$ , and a Boolean function  $\Phi_k$  such that

$$uv \in E_k \iff \Phi_k(g_{k,1}(u, v) \geq 0, \dots, g_{k,t}(u, v) \geq 0) = 1.$$

For  $1 \leq k \leq m, 1 \leq \ell \leq t, v \in V$ , we define the surface  $\sigma_{k,\ell}(v) = \{x \in \mathbb{R}^d : g_{k,\ell}(v, x) = 0\}$ .

Before we continue, let us briefly sketch the idea of the proof. We start by applying Lemma 2.1 (the cutting lemma) to  $\Sigma = \{\sigma_{k,\ell}(v) : k \in [m], \ell \in [t], v \in V\}$ , the set of surfaces which determines the neighborhoods of each vertex, and obtain a space partition which induces a partition of the vertex set  $V = V_1 \cup \dots \cup V_K$ . If there is a “large” part  $V_j$  with many distinct colors appearing in  $V_j \times (V \setminus V_j)$ , then we show that  $V_j$  induces few distinct colors, and by induction we can find a monochromatic copy of  $K_{p_k}$  for some  $k \in [m]$ . If none of the “large” parts has the above property, the colors of nearly all edges can be defined by much fewer polynomial inequalities, i.e., by a much smaller set of surfaces  $\Sigma' \subset \Sigma$ . Now we can repeat.

In what follows, we spell out these ideas in full detail. Set  $m_0 = m$  and define  $m_i = 4d \log(c_1 m_{i-1} t)$  for  $i > 0$ . We will establish the following claim.

**Claim 3.2.** *Let  $V$  and  $E_1, \dots, E_m \subset \binom{V}{2}$  be defined as above. Then for  $i \geq 0$  we will recursively find either*

1. *a monochromatic copy of  $K_{p_k}$  in color  $k$  for some  $k \in [m]$ , or*
2. *a function  $\chi_i : V \rightarrow 2^{[m]}$  such that  $|\chi_i(v)| \leq m_i$ , and the number of edges  $uv \in \binom{V}{2}$  with the property that for one of its endpoints, say  $u$ , no color assigned to  $uv$  belongs to  $\chi_i(u)$ , is at most  $\frac{4n^2}{tm_{i-1}}$ . We will refer to these edges as bad at stage  $i$ . All edges that are not bad are called good at this stage, meaning that, there is a color  $k$  appearing on  $uv$  such that  $k \in \chi_i(u)$ , and there is a color  $k'$  appearing on  $uv$  such that  $k' \in \chi_i(v)$ .*

*Proof.* We start by setting  $\chi_0(v) = [m]$  for all  $v \in V$ , and  $m_0 = m$ . Having found  $\chi_i$  with the properties above, we will produce  $\chi_{i+1}$  as follows. We have  $m_{i+1} = 4d \log(c_1 m_i t)$ , and let us assume

that  $m_i > (8c_1 dtp)^2$ . Hence, there are at most  $\frac{4n^2}{tm_{i-1}}$  bad edges. Let  $\Sigma$  be the set of surfaces  $\sigma_{k,\ell}(v)$ , where  $v \in V$ ,  $k \in \chi_i(v)$ , and  $1 \leq \ell \leq t$ . This implies that  $|\Sigma| \leq nm_i t$ .

We apply Lemma 2.1 to  $\Sigma$  with parameter  $r = (tm_i)^2$  to obtain a  $(1/(tm_i)^2)$ -cutting  $\Psi = \{\Delta_1, \Delta_2, \dots, \Delta_{K_0}\}$ , such that  $K_0 \leq c_1(tm_i)^{4d}$ . Hence, we have a partition  $\mathcal{P}_0 : V = V_1 \cup \dots \cup V_{K_0}$ , where  $V_j = V \cap \Delta_j$  for  $\Delta_j \in \Psi$ . For each part  $V_j$  of size greater than  $2n/(tm_i)$ , we (arbitrarily) partition  $V_j$  into parts of size  $\lfloor 2n/(tm_i) \rfloor$  and possibly one additional part of size less than  $2n/(tm_i)$ . Let  $\mathcal{P} : V = V_1 \cup \dots \cup V_K$  be the resulting partition, where  $K \leq 2c_1(tm_i)^{4d}$  and  $|V_j| \leq 2n/(tm_i)$  for all  $j$ .

Now we define  $\chi_{i+1}(v)$  for all  $v \in V$ .

*Case 1.* If  $v \in V_j$  for some  $V_j$  with  $|V_j| < \frac{n}{2c_1(tm_i)^{4d+1}}$ , we set  $\chi_{i+1}(v) = \emptyset$ .

*Case 2.* Suppose  $v \in V_j$  such that  $|V_j| \geq \frac{n}{2c_1(tm_i)^{4d+1}}$ . In order to define  $\chi_{i+1}(v)$ , we need some preparation. Let  $\Delta_j \in \Psi$  such that  $V_j \subset \Delta_j$ . We define  $X_j \subset V \setminus V_j$  to be the set of vertices from  $V \setminus V_j$  that gives rise to a surface in  $\Sigma$  that crosses  $\Delta_j$ . By Lemma 2.1, the cutting lemma, we have  $|X_j| \leq n/(tm_i)$ .

Fix a vertex  $v \in V \setminus \{V_j, X_j\}$ . Since none of the surfaces of the form  $\sigma_{k,\ell}(v)$ , where  $k \in \chi_i(v)$  and  $\ell \in \{1, \dots, t\}$ , cross  $\Delta_j$ , either  $v \times V_j$  is monochromatic with color  $k$  for some  $k \in \chi_i(v)$ , or none of the colors in  $\chi_i(v)$  appear in  $v \times V_j$ . Let  $S_j$  be the set of vertices  $v \in V \setminus \{V_j, X_j\}$  satisfying the former condition and let  $T_j$  denote a set of vertices  $v \in V \setminus \{V_j, X_j\}$  satisfying the latter one. Since there are at most  $4n^2/(tm_{i-1})$  bad edges, we have

$$|T_j| \frac{n}{2c_1(tm_i)^{4d+1}} \leq \frac{4n^2}{tm_{i-1}},$$

which implies

$$|T_j| \leq \frac{8nc_1(tm_i)^{4d+1}}{tm_{i-1}} \leq \frac{n}{tm_i},$$

where the last inequality follows from the fact that  $m_{i-1} = 2^{m_i/4d}/(c_1 t)$ , and the assumption  $m_i > (8c_1 dtp)^2$ . Now, suppose there are at least  $m_{i+1} = 4d \log(c_1 tm_i)$  distinct colors between  $V_j$  and  $S_j$ . Let  $I = \{k_1, \dots, k_{m_{i+1}}\} \subset [m]$  be the set of these  $m_{i+1}$  distinct colors. Then there are  $m_{i+1}$  vertices  $v_1, \dots, v_{m_{i+1}} \in S_j$ , possibly with repetition, such that  $v_w \times V_j$  is monochromatic with color  $k_w \in I$ , for each  $w \in \{1, \dots, m_{i+1}\}$ . Hence, if  $V_j$  contains a monochromatic copy of  $K_{p_k-1}$  in color  $k \in I$ , we would have a monochromatic copy of  $K_{p_k}$  in color  $k$ . On the other hand, if  $V_j$  does not contain a monochromatic copy of  $K_{p_k-1}$  in color  $k$  for no  $k \in I$ , then, using that

$$|V_j| \geq \frac{n}{2c_1(tm_i)^{4d+1}} > 2^{cs-8d \log(c_1 m_i t)} > 2^{c(s-m_{i+1})} = 2^{c(\sum_{k \in I} (p_k-1) + \sum_{k \notin I} p_k)}$$

for a sufficiently large  $c$ , we obtain by induction that there is a monochromatic copy of  $K_{p_k}$  in color  $k$  where  $k \notin I$ .

Therefore, we can assume that the number of distinct colors between  $V_j$  and  $S_j$  is less than  $m_{i+1} = 4d \log(c_1 m_i t)$ . For every vertex  $v \in V_j$ , define  $\chi_{i+1}(v)$  as the set of all colors that appear on the edges belonging to  $v \times S_j$ .

Now that we have defined  $m_{i+1}$  and  $\chi_{i+1}$  such that  $|\chi_{i+1}(v)| \leq m_{i+1}$  for all  $v \in V$ , it remains to show that the number of edges  $uv \in \binom{V}{2}$  with the property that for one of its endpoints, say

$u$ , no color assigned to  $uv$  belongs to  $\chi(u)$ , is at most  $\frac{4n^2}{tm_i}$ . Let  $B \subset \binom{V}{2}$  be the collection of such edges. Notice that if  $uv \in B$ , then either

1. both  $u$  and  $v$  lie inside the same part in the partition  $\mathcal{P}$ , or
2.  $u$  or  $v$  lies inside a part  $V_j$  such that  $|V_j| < \frac{n}{2c_1(tm_i)^{4d+1}}$ , or
3.  $u \in V_j$  with  $|V_j| \geq \frac{n}{2c_1(tm_i)^{4d+1}}$  and  $v \in X_j \cup T_j$ , or
4.  $v \in V_j$  with  $|V_j| \geq \frac{n}{2c_1(tm_i)^{4d+1}}$  and  $u \in X_j \cup T_j$ .

Since each part in  $\mathcal{P}$  has at most  $2n/(tm_i)$  vertices, the number of edges of type 1 is at most  $\frac{tm_i}{2} \left( \frac{2n}{tm_i} \right)^2 / 2 = n^2/(tm_i)$ . The number of edges of type 2 is at most

$$\left( \frac{n}{2c_1(tm_i)^{4d+1}} \right) K \cdot n \leq \frac{n^2}{tm_i}.$$

Since  $|X_j|, |T_j| \leq n/(tm_i)$ , the number of edges of types 3 and 4 is at most

$$\sum_j |V_j| \frac{2n}{tm_i} \leq \frac{2n^2}{tm_i}.$$

Hence,  $|B| \leq \frac{4n^2}{tm_i}$ . Therefore, either we have found a monochromatic copy of  $K_{p_k}$  in color  $k$  for some  $k \in [m]$ , or we have found  $m_{i+1}$  and  $\chi_{i+1}$  with the desired properties.  $\square$

Let  $w$  be the minimum integer such that  $m_w \leq (8c_1dtp)^2$ . Then either we have found a monochromatic copy of  $K_{p_k}$  in color  $k$  for some  $k \in [m]$ , or we have obtained  $m_w$  and  $\chi_w$  with the desired properties. Since there are at most  $4n^2/(tm_{w-1}) < n^2/8$  bad edges, there is a vertex  $v \in V$  incident to at least  $n/2$  good edges. Moreover, since  $|\chi_w(v)| \leq m_w \leq (8c_1dtp)^2$ , at least  $\frac{n}{2(8c_1dtp)^2}$  of these edges incident to  $v$  have color  $k'$  for some color  $k' \in \chi_w(v)$ . Let  $S \subset V$  be the set of endpoints of these edges. If  $S$  contains a monochromatic copy of  $K_{p_{k'}-1}$  in color  $k'$ , then we are done. On the other hand, if  $S$  does not contain a monochromatic copy of  $K_{p_{k'}-1}$  in color  $k'$ , and using the lower bound

$$|S| \geq \frac{n}{2(8c_1dtp)^2} = \frac{2^{cs}}{2(8c_1dtp)^2} \geq 2^{c(\sum_{k \neq k'} p_k + (p_{k'} - 1))},$$

for  $c = c(d, t, p)$  sufficiently large, we conclude by induction that  $S$  contains a monochromatic copy of  $K_{p_k}$  for some  $k \neq k'$ . This completes the proof of Theorem 3.1.  $\square$

## 4 Multicolor semi-algebraic regularity lemma

### –Proof of Theorem 1.2

First, we prove the following variant of Theorem 1.2, which easily implies Theorem 1.2.

**Theorem 4.1.** *For any  $\varepsilon > 0$ , every  $n$ -element point set  $V \subset \mathbb{R}^d$  equipped with semi-algebraic binary relations  $E_1, \dots, E_m \subset \binom{V}{2}$  such that  $\binom{V}{2} = E_1 \cup \dots \cup E_m$  and each  $E_k$  has complexity at most  $t$ , can be partitioned into  $K \leq c_2(\frac{m}{\varepsilon})^{5d^2}$  parts  $V = V_1 \cup \dots \cup V_K$ , where  $c_2 = c_2(d, t)$ , such that*

$$\sum \frac{|V_i||V_j|}{n^2} \leq \varepsilon,$$

where the sum is taken over all pairs  $(i, j)$  such that  $(V_i, V_j)$  is not complete with respect to  $E_k$  for all  $k = 1, \dots, m$ .

*Proof.* For each relation  $E_k$ , let  $g_{k,1}, \dots, g_{k,t} \in \mathbb{R}[x_1, \dots, x_{2d}]$  be polynomials of degree at most  $t$ , and let  $\Phi_k$  be a boolean formula such that

$$uv \in E_k \iff \Phi_k(g_{k,1}(u, v) \geq 0; \dots; g_{k,t}(u, v) \geq 0) = 1.$$

For each point  $x \in \mathbb{R}^d$ ,  $k \in \{1, \dots, m\}$ , and  $\ell \in \{1, \dots, t\}$ , we define the surface

$$\sigma_{k,\ell}(x) = \{y \in \mathbb{R}^d : g_{k,\ell}(x, y) = 0\}.$$

Let  $\Sigma$  be the family of  $tmn$  surfaces in  $\mathbb{R}^d$  defined by

$$\Sigma = \{\sigma_{k,\ell}(u) : u \in V, 1 \leq k \leq m, 1 \leq \ell \leq t\}.$$

We apply Lemma 2.1 to  $\Sigma$  with parameter  $r = tm/\varepsilon$  to obtain a  $(1/r)$ -cutting  $\Psi$ , where  $|\Psi| = s \leq c_1 \left(\frac{tm}{\varepsilon}\right)^{2d}$ , such that each semi-algebraic set  $\Delta_i \in \Psi$  has complexity at most  $c_1$ , where  $c_1$  is defined in Lemma 2.1. Hence, at most  $tmn/r = \varepsilon n$  surfaces from  $\Sigma$  cross  $\Delta_i$  for every  $i$ . This implies that at most  $\varepsilon n$  points in  $V$  give rise to at least one surface in  $\Sigma$  that crosses  $\Delta_i$ .

Let  $U_i = V \cap \Delta_i$  for each  $i \leq s$ . We now partition  $\Delta_i$  as follows. For  $k \in \{1, \dots, m\}$  and  $j \in \{1, \dots, s\}$ , define  $\Delta_{i,j,k} \subset \mathbb{R}^d$  by

$$\Delta_{i,j,k} = \{x \in \Delta_i : \sigma_{k,1}(x) \cup \dots \cup \sigma_{k,t}(x) \text{ crosses } \Delta_j\}.$$

That is,  $\Delta_{i,j,k}$  will correspond to the vertices in  $\Delta_i$  that may not be homogeneous to the set of vertices in  $\Delta_j$  with respect to color  $k$ .

**Observation 4.2.** For any  $i, j$ , and  $k$ , the semi-algebraic set  $\Delta_{i,j,k}$  has complexity at most  $c_3 = c_3(d, t)$ .

*Proof.* Set  $\sigma_k(x) = \sigma_{k,1}(x) \cup \dots \cup \sigma_{k,t}(x)$ , which is a semi-algebraic set with complexity at most  $c_4 = c_4(d, t)$ . Then

$$\Delta_{i,j,k} = \left\{ x \in \Delta_i : \begin{array}{l} \exists y_1 \in \mathbb{R}^d \text{ s.t. } y_1 \in \sigma_k(x) \cap \Delta_j, \text{ and} \\ \exists y_2 \in \mathbb{R}^d \text{ s.t. } y_2 \in \Delta_j \setminus \sigma_k(x). \end{array} \right\}.$$

We can apply quantifier elimination (see Theorem 2.74 in [2]) to make  $\Delta_{i,j,k}$  quantifier-free, with description complexity at most  $c_3 = c_3(d, t)$ .  $\square$

Set  $\mathcal{F}_i = \{\Delta_{i,j,k} : 1 \leq k \leq m, 1 \leq j \leq s\}$ . We partition the points in  $U_i$  into equivalence classes, where two points  $u, v \in U_i$  are equivalent if and only if  $u$  belongs to the same members of  $\mathcal{F}_i$  as  $v$  does. Since  $\mathcal{F}_i$  gives rise to at most  $c_3|\mathcal{F}_i|$  polynomials of degree at most  $c_3$ , by the Milnor-Thom theorem (see [18] Chapter 6), the number of distinct sign patterns of these  $c_3|\mathcal{F}_i|$  polynomials is at most  $(50c_3(c_3|\mathcal{F}_i|))^d$ . Hence, there is a constant  $c_5 = c_5(d, t)$  such that  $U_i$  is partitioned into at most  $c_5(m s)^d$  equivalence classes. After repeating this procedure to each  $U_i$ , we obtain a partition of our point set  $V = V_1 \cup \dots \cup V_K$  with

$$K \leq sc_5(ms)^d = c_5m^d s^{d+1} \leq c_5t^{2d(d+1)}c_1^{d+1} \left(\frac{m}{\varepsilon}\right)^{5d^2} = c_2 \left(\frac{m}{\varepsilon}\right)^{5d^2},$$

where we define  $c_2 = c_5t^{2d(d+1)}c_1^{d+1}$ .

For fixed  $i$ , consider the part  $V_i$ . Then there is a semi-algebraic set  $\Delta_{w_i}$  obtained from Lemma 2.1 such that  $U_{w_i} = V \cap \Delta_{w_i}$  and  $V_i \subset U_{w_i} \subset \Delta_{w_i}$ . Now consider all other parts  $V_j$  such that not all of their elements are related to every element of  $V_i$  with respect to any relation  $E_k$  where  $1 \leq k \leq m$ . Then each point  $u \in V_j$  gives rise to a surface in  $\Sigma$  that crosses  $\Delta_{w_i}$ . By Lemma 2.1, the total number of such points in  $V$  is at most  $\varepsilon n$ . Therefore, we have

$$\sum_j |V_i||V_j| = |V_i| \sum_j |V_j| \leq |V_i|\varepsilon n,$$

where the sum is over all  $j$  such that  $V_i \times V_j$  is not contained in the relation  $E_k$  for any  $k$ . Summing over all  $i$ , we have

$$\sum_{i,j} |V_i||V_j| \leq \varepsilon n^2,$$

where the sum is taken over all pairs  $i, j$  such that  $(V_i, V_j)$  is not complete with respect to  $E_k$  for all  $k$ .  $\square$

*Proof of Theorem 1.2.* Apply Theorem 4.1 with approximation parameter  $\varepsilon/2$ . Hence, there is a partition  $\mathcal{Q} : V = U_1 \cup \dots \cup U_{K'}$  into  $K' \leq (m/\varepsilon)^c$  parts with  $c = c(d, t)$  and  $\sum |U_i||U_j| \leq (\varepsilon/2)|V|^2$ , where the sum is taken over all pairs  $(i, j)$  such that  $(U_i, U_j)$  is not complete with respect to  $E_k$  for all  $k$ .

Let  $K = 8\varepsilon^{-1}K'$ . Partition each part  $U_i$  into parts of size  $|V|/K$  and possibly one additional part of size less than  $|V|/K$ . Collect these additional parts and divide them into parts of size  $|V|/K$  to obtain an equitable partition  $\mathcal{P} : V = V_1 \cup \dots \cup V_K$  into  $K$  parts. The number of vertices of  $V$  which are in parts  $V_i$  that are not contained in a part of  $\mathcal{Q}$  is at most  $K'|V|/K$ . Hence, the fraction of pairs  $V_i \times V_j$  with not all  $V_i, V_j$  are subsets of parts of  $\mathcal{Q}$  is at most  $2K'/K = \varepsilon/4$ . As  $\varepsilon/2 + \varepsilon/4 < \varepsilon$ , we obtain that less than an  $\varepsilon$ -fraction of the pairs of parts of  $\mathcal{P}$  are not complete with respect to any relation  $E_1, \dots, E_m$ .  $\square$

## 5 Generalized Ramsey numbers for semi-algebraic colorings

### –Proof of Theorem 1.3

Due to the lack of understanding of the classical Ramsey number  $R(p; m)$ , Erdős and Shelah (see [6]) introduced the following generalization, which was studied by Erdős and Gyárfás in [7].

**Definition 5.1.** For integers  $p$  and  $q$  with  $2 \leq q \leq \binom{p}{2}$ , a  $(p, q)$ -coloring is an edge-coloring of a complete graph in which every  $p$  vertices induce at least  $q$  distinct colors.

Let  $f(n, p, q)$  be the minimum integer  $m$  such that there is a  $(p, q)$ -coloring of  $K_n$  with at most  $m$  colors. Here, both  $p$  and  $q$  are considered fixed integers, where  $p \geq 3$ ,  $2 \leq q \leq \binom{p}{2}$ , and  $n$  tends to infinity. Trivially, we have  $f(n, p, \binom{p}{2}) = \binom{n}{2}$ , and at the other end, estimating  $f(n, p, 2)$  is equivalent to estimating  $R(p; m)$  since  $f(n, p, 2)$  is the inverse of  $R(p; m)$ . In particular,

$$\Omega\left(\frac{\log n}{\log \log n}\right) \leq f(n, 3, 2) \leq O(\log n). \quad (1)$$

Erdős and Gyárfás [7] determined certain ranges for  $q \in \{2, 3, \dots, \binom{p}{2}\}$  for which  $f(n, p, q)$  is quadratic, linear, and subpolynomial in  $n$ . In particular, they showed that

$$\Omega\left(n^{\frac{1}{p-2}}\right) \leq f(n, p, p) \leq O\left(n^{\frac{2}{p-1}}\right),$$

which implies that  $f(n, p, q)$  is polynomial in  $n$  for  $q \geq p$ . Surprisingly, estimating  $f(n, p, p-1)$  is much more difficult. They [7] asked for  $p$  fixed if  $f(n, p, p-1) = n^{o(1)}$ . The trivial lower bound is  $f(n, p, p-1) \geq f(n, p, 2) \geq \Omega\left(\frac{\log n}{\log \log n}\right)$ , which was improved by several authors [17, 13], and it is now known [4] that  $f(n, p, p-1) \geq \Omega(\log n)$ . In the other direction, Mubayi [19] found an elegant construction which implies  $f(n, 4, 3) \leq e^{O(\sqrt{\log n})}$ , and later, Conlon *et al.* [4] gave another example which implies  $f(n, p, p-1) \leq e^{(\log n)^{1-1/(p-2)+o(1)}}$ . Hence, it is now known that  $f(n, p, p-1)$  does not grow as a power in  $n$ .

Here, we study the variant of the function  $f(n, p, q)$  for point sets  $V \subset \mathbb{R}^d$  equipped with semi-algebraic relations. Let  $f_{d,t}(n, p, q)$  be the minimum  $m$  such that there is a  $(p, q)$ -coloring of  $K_n$  with  $m$  colors, whose vertices can be chosen as points in  $\mathbb{R}^d$ , and each color class can be defined by a semi-algebraic relation on the point set with complexity at most  $t$ . We note that here we require that each edge receives *exactly* one color. Clearly, we have  $f(n, p, q) \leq f_{d,t}(n, p, q)$ . Theorem 1.3 stated in the Introduction shows the exact value of  $q$  for which  $f_{d,t}(n, p, q)$  changes from  $\log n$  to a power of  $n$ .

In the rest of this section, we prove Theorem 1.3. Let  $V$  be a set of points in  $\mathbb{R}^d$  equipped with semi-algebraic relations  $E_1, \dots, E_m$  such that each  $E_k$  has complexity at most  $t$ ,  $\binom{V}{2} = E_1 \cup \dots \cup E_m$ , and  $E_k \cap E_\ell = \emptyset$  for all  $k \neq \ell$ . Let  $S_1, S_2 \subset V$  be  $q$ -element subsets of  $V$ . We say that  $S_1$  and  $S_2$  are *isomorphic*, denoted by  $S_1 \simeq S_2$ , if there is a bijective function  $h : S_1 \rightarrow S_2$  such that for  $u, v \in S_1$  we have  $uv \in E_k$  if and only if  $h(u)h(v) \in E_k$ .

Let  $S \subset V$  be such that  $|S| = 2^s$  for some positive integer  $s$ . We say that  $S$  is *s-layered* if  $s = 1$  or if there is a partition  $S = S_1 \cup S_2$  such that  $|S_1| = |S_2| = 2^{s-1}$ ,  $S_1$  and  $S_2$  are  $(s-1)$ -layered,  $S_1 \simeq S_2$ , and for all  $u \in S_1$  and  $v \in S_2$  we have  $uv \in E_k$  for some fixed  $k$ . Notice that given an  $s$ -layered set  $S$ , there are at most  $s$  relations  $E_{k_1}, \dots, E_{k_s}$  such that  $\binom{S}{2} \subset E_{k_1} \cup \dots \cup E_{k_s}$ . Hence, the lower bound in Theorem 1.3 is a direct consequence of the following result.

**Theorem 5.2.** *Let  $s \geq 1$  and let  $V$  be an  $n$ -element point set in  $\mathbb{R}^d$  equipped with semi-algebraic relations  $E_1, \dots, E_m$  such that each  $E_k$  has complexity at most  $t$ ,  $E_1 \cup \dots \cup E_m = \binom{V}{2}$ , and  $E_k \cap E_\ell = \emptyset$  for all  $k \neq \ell$ . If  $m \leq n^{\frac{1}{cs^2}}$ , then there is a subset  $S \subset V$  such that  $|S| = 2^s$  and  $S$  is  $s$ -layered, where  $c = c(d, t)$ .*

*Proof.* We proceed by induction on  $s$ . The base case  $s = 1$  is trivial. For the inductive step, assume that the statement holds for  $s' < s$ . We will specify  $c = c(d, t)$  later. We start by applying Theorem 1.2 with parameter  $\varepsilon = \frac{1}{m^s}$  to the point set  $V$ , which is equipped with semi-algebraic relations  $E_1, \dots, E_m$ , and obtain an equitable partition  $\mathcal{P} : V = V_1 \cup \dots \cup V_K$ , where

$$K \leq c_2 \left(\frac{m}{\varepsilon}\right)^{5d^2} \leq c_2 m^{10sd^2},$$

and  $c_2 = c_2(d, t)$ . Since all but an  $\varepsilon$  fraction of the pairs of parts in  $\mathcal{P}$  are complete with respect to  $E_k$  for some  $k$ , by Turán's theorem, there are  $m^{s-1} + 1$  parts  $V'_i \in \mathcal{P}$  such that each pair  $(V'_i, V'_j) \in \mathcal{P} \times \mathcal{P}$  is complete with respect to some relation  $E_k$ . Since  $\mathcal{P}$  is an equitable partition, we have  $|V'_i| \geq \frac{n}{c_2 m^{10d^2 s}}$ . By picking  $c = c(d, t)$  sufficiently large, we have

$$|V'_i|^{\frac{1}{c(s-1)^2}} \geq \left( \frac{n}{c_2 m^{10d^2 s}} \right)^{\frac{1}{c(s-1)^2}} \geq m^{\frac{cs^2 - 10c_2 d^2 s}{c(s-1)^2}} \geq m.$$

By the induction hypothesis, each  $V'_i$  contains an  $(s-1)$ -layered set  $S_i$  for  $i \in \{1, \dots, m^{s-1} + 1\}$ . By the pigeonhole principle, there are two  $(s-1)$ -layered sets  $S_i, S_j$  such that  $S_i \simeq S_j$ . Since  $S_i \times S_j \subset E_k$  for some  $k$ , the set  $S = S_i \cup S_j$  is an  $s$ -layered set. This completes the proof.  $\square$

To prove the upper bound for  $f_{d,t}(n, p, \lceil \log p \rceil)$ , when  $d \geq 1$  and  $t \geq 100$ , it is sufficient to construct a  $2^m$ -element point set  $V \subset \mathbb{R}$  equipped with  $m$  distinct semi-algebraic relations  $E_1, \dots, E_m$  that is  $m$ -layered. More precisely, for each integer  $m \geq 1$ , we construct a set  $V_m$  of  $2^m$  points in  $\mathbb{R}$  equipped with semi-algebraic relations  $E_1, \dots, E_m$  such that

1.  $V_m$  with respect to relations  $E_1, \dots, E_m$  is  $m$ -layered,
2.  $E_1 \cup \dots \cup E_m = \binom{V_m}{2}$  is a partition,
3. each  $E_i$  has complexity at most four, and
4. each  $E_i$  is *shift invariant*, that is  $uv \in E_i$  if and only if  $(u + c, v + c) \in E_i$  for  $c \in \mathbb{R}$ .

We start by setting  $V_1 = \{1, 2\}$  and defining  $E_1 = \{u, v \in V_1 : |u - v| = 1\}$ . Having defined the point set  $V_i$  and relations  $E_1, \dots, E_i$ , we define  $V_{i+1}$  and  $E_{i+1}$  as follows. Let  $C = C(i)$  be a sufficiently large integer such that  $C > 10 \max_{u \in V_i} u$ . Then we have  $V_{i+1} = V_i \cup (V_i + C)$ , where  $V_i + C$  is a translated copy of  $V_i$ . We now define the relation  $E_{i+1}$  by

$$uv \in E_{i+1} \iff C/2 < |u - v| < 2C.$$

Hence,  $V_{i+1}$  with respect to relations  $E_1, \dots, E_{i+1}$  satisfies the properties stated above and is clearly  $(i+1)$ -layered. One can easily check that any set of  $p$  points in  $V_m$  induces at least  $\lceil \log p \rceil$  distinct relations (colors).

Let us remark that the arguments above hold for semi-algebraic relations  $E_1, \dots, E_m$  that are not necessarily disjoint if one defines a  $(p, q)$ -coloring as follows. Given a coloring  $\chi : \binom{V(K_n)}{2} \rightarrow 2^{[m]}$  on the edges of  $K_n$ , where each edge receives *at least* one color among  $[m]$ ,  $\chi$  is a  $(p, q)$ -coloring if for every set  $S \subset V$  of size  $p$ , no matter how you choose one color in  $\chi(uv)$  for each edge  $uv \in \binom{S}{2}$ ,  $S$  will induce at least  $q$  distinct colors.

## 6 Concluding remarks

In [22], it was shown that  $R_{1,t}(3; m) > (1681)^{m/7}$  for  $t > 5$ , thus implying that the upper bound in Theorem 1.1 is tight up to a constant factor in the exponent. This can be improved as follows. Let  $C(p) = \lim_{m \rightarrow \infty} R(p; m)^{1/m}$ . Note that this limit exists by considering product colorings, but may be finite or infinite. Then for each  $C < C(p)$ , there is a  $t = t(C, p)$ , such that for all  $m$  sufficiently large we have

$$R_{1,t}(p; m) > C^m.$$

Indeed, take a fixed coloring of the edges of  $K_N$  which realizes  $R(p; m_0) > C^{m_0}$ , and recursively blow up this graph by introducing  $m_0$  new colors at each stage. Then this coloring can be realized semi-algebraically in  $\mathbb{R}$  with  $t = O(m_0^2)$  linear constraints for each color class based on distances.

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