

Fast signal recovery from quadratic measurements

Miguel Moscoso, Alexei Novikov, George Papanicolaou and Chrysoula Tsogka

Abstract—We present a novel approach for recovering a sparse signal from quadratic measurements corresponding to a rank-one tensorization of the data vector. Such quadratic measurements, referred to as interferometric or cross-correlated data, naturally arise in many fields such as remote sensing, spectroscopy, holography and seismology. Compared to the sparse signal recovery problem that uses linear measurements, the unknown in this case is a matrix formed by the cross correlations of the sought signal. This creates a bottleneck for the inversion since the number of unknowns grows quadratically with the dimension of the signal. The main idea of the proposed approach is to reduce the dimensionality of the problem by recovering only the diagonal of the unknown matrix, whose dimension grows linearly with the size of the signal, and use an efficient *Noise Collector* to absorb the cross-correlated data that come from the off-diagonal elements of this matrix. These elements do not carry extra information about the support of the signal, but significantly contribute to these data. With this strategy, we recover the unknown matrix by solving a convex linear problem whose cost is similar to the one that uses linear measurements. Our theory shows that the proposed approach provides exact support recovery when the data is not too noisy, and that there are no false positives for any level of noise. It also demonstrates that the level of sparsity that can be recovered scales almost linearly with the number of data. The numerical experiments presented in the paper corroborate these findings.

Index Terms—quadratic data, ℓ_1 -minimization, noise, dimension reduction

I. INTRODUCTION

Reconstruction of signals or images from cross correlations has interesting applications in many fields of science and engineering such as optics, quantum mechanics, electron microscopy, antenna testing, seismic interferometry, or imaging in general [14], [18], [35], [31]. Using cross correlations of measurements collected at different locations presents several advantages since the inversion does not require knowledge of the emitter positions, or the probing pulses shapes as only time differences matter. Cross correlations have been used, for example, when imaging is carried out with opportunistic sources whose properties are mainly unknown [13], [11], [19].

In many applications, we seek information about an object or a signal $\rho \in \mathcal{C}^K$ given data $\mathbf{b} \in \mathcal{C}^N$ most often related through a linear transformation

$$\mathcal{A}\rho = \mathbf{b}, \quad (1)$$

M. Moscoso is with the Department of Mathematics, Universidad Carlos III de Madrid, Leganes, Madrid 28911, Spain e-mail: moscoso@math.uc3m.es.

A. Novikov is with the Mathematics Department, Penn State University, University Park, PA 16802, USA e-mail: novikov@psu.edu

G. Papanicolaou is with the Department of Mathematics, Stanford University, Stanford, CA 94305 USA e-mail: papanicolaou@stanford.edu

C. Tsogka is with the Department of Applied Mathematics, University of California, Merced, 5200 North Lake Road, Merced, CA 95343 USA e-mail: ctsogka@ucmerced.edu

where $\mathcal{A} \in \mathcal{C}^{N \times K}$ is the measurement or model matrix. When the signal ρ is compressed or when the data is scarce, $N < K$, in which case (1) is underdetermined and infinitely many signals or objects match the data. However, if the signal ρ is sparse so only $M \ll K$ components are different than zero, ℓ_1 -minimization algorithms that solve

$$\rho_{\ell_1} = \operatorname{argmin} \|\rho\|_{\ell_1}, \text{ subject to } \mathcal{A}\rho = \mathbf{b} \quad (2)$$

can recover the true signal efficiently even when $N \ll K$.

On the other hand, there are situations in which it is difficult or impossible to record high quality data \mathbf{b} , and it is more convenient to use the cross-correlated data contained in the matrix

$$\mathbf{B} = \mathbf{b}\mathbf{b}^* \in \mathcal{C}^{N \times N} \quad (3)$$

to find the desired information about the object or signal ρ (see [12] and references therein). We will refer to (3) as the quadratic data.

One way to address this problem is to lift it to the matrix level and reformulate it as a low-rank matrix linear system, which can be solved by using nuclear norm minimization as it was suggested in [7], [4] for imaging with intensities-only. This makes the problem convex over the appropriate matrix vector space and, thus, the unique true solution can be found using well established algorithms with increased storage requirements as they involve the SVD of the iterate matrix unknown [1]. Thus, the big caveat is that the computational cost and memory requirements rapidly become prohibitive because the dimension of the problem increases quadratically with K , making its solution infeasible for large scale problems.

To avoid the bottleneck arising from squaring the dimensionality of the inverse problem, non-convex approaches that use gradient descent iterations and operate on the original signal domain have also been recently proposed [6], [10], [34], [5]. The premise of these methods is that if the initial guess obtained by means of a spectral method is sufficiently accurate, the iterates provably converge to the correct solution (up to a global phase) with a geometric rate and algorithmic step-size $O(1/K)$. Other interesting non-convex methods with some theoretical guarantees for the matrix completion problem are the *OptSpace* algorithm [21] that also computes the initial guess by means of a spectral method, and the *AltMinPhase* algorithm [29] that alternates between updates of the phase and the signals.

In this paper we suggest a different approach. We propose to consider the linear matrix equation

$$\mathcal{A}X\mathcal{A}^* = \mathbf{B} \quad (4)$$

for the correlated signal $X = \rho\rho^* \in \mathcal{C}^{K \times K}$, vectorize both sides so that

$$\operatorname{vec}(\mathcal{A}X\mathcal{A}^*) = \operatorname{vec}(\mathbf{B}), \quad (5)$$

and use the Kronecker product \otimes , and its property $\text{vec}(PQR) = (R^T \otimes P)\text{vec}(Q)$, to express the matrix multiplications as the linear transformation

$$(\bar{\mathcal{A}} \otimes \mathcal{A}) \text{vec}(X) = \text{vec}(B). \quad (6)$$

Thus, we can promote the sparsity of the sought signal using ℓ_1 -minimization algorithms that are more efficient than nuclear norm minimization ones, as they do not require heavy operations that involve matrix factorizations. However, the dimension of the unknown $\text{vec}(X)$ in (6) also increases quadratically with K , so this approach by itself would still be impractical when K is not small.

Furthermore, one might think that it is not a good idea to go quadratic and solve (6) instead of (1) for two reasons. First, because the condition number of $(\bar{\mathcal{A}} \otimes \mathcal{A})$ increases quadratically, so the conditioning of the problem can worsen substantially and, thus, its numerical solution can be less reliable. Second, and perhaps more importantly, because the mutual coherence of the matrix $(\bar{\mathcal{A}} \otimes \mathcal{A})$ is exactly the same as the one of matrix \mathcal{A} [20], i.e. $\mu(\bar{\mathcal{A}} \otimes \mathcal{A}) = \mu(\mathcal{A})$, so recovery guarantees do not get better. However, it follows from [20] that if one chooses only a few specific columns of the matrix $(\bar{\mathcal{A}} \otimes \mathcal{A})$, then $\mu(\bar{\mathcal{A}} \otimes \mathcal{A}) = (\mu(\mathcal{A}))^2$. The right columns to choose are those that correspond to the diagonal entries of the matrix X .

Hence, we propose to use a *Noise Collector* to reduce the dimensionality of problem (6). The *Noise Collector* was introduced in [28] to eliminate the clutter in the recovered signals when the data are contaminated by additive noise. In this paper, we use the *Noise Collector* to absorb part of the data instead. Specifically, we treat the data vector that corresponds to the $K^2 - K$ off-diagonal entries in the matrix X as noise. Using the *Noise Collector* allows us to ignore these entries and construct a linear system with the same number of unknowns as the original problem (1) that uses linear data. Thus, a dimension reduction from K^2 to K unknowns is achieved with almost no extra computational cost. This is because the number of operations used by the *Noise Collector* is of the same order as the ones needed to solve the original problem with linear measurements (see Section IV-A). In other words, the cost of solving the sparse signal recovery problem using cross-correlated data also grows linearly with K using the proposed approach. We point out that the *Noise Collector* also works as a regularization of the inverse problem, so its condition number becomes $O(1)$.

The main result of this paper is Theorem 3 which says that under certain decoherence conditions on the matrix \mathcal{A} , we can efficiently find the support of an M -sparse signal exactly if the cross-correlated data is noise-free or the noise is low enough. Furthermore, Theorem 3 shows that the level of sparsity M that can be recovered is $O(N/\sqrt{\ln N})$.

The numerical experiments included in this paper support the results of Theorem 3. They show that the support of a signal can be found exactly if the noise in the data is not too large with almost no extra computational cost with respect to the original problem (1) that considers linear data with no correlations. Once the support has been found, a trivial second step allows us to find the signal, including its phases.

The reconstruction is exact when there is no noise in the data and the results are very satisfactory even for noisy data with low signal to noise ratios. That is, our numerical experiments suggest that the approach presented here is robust with respect to additive noise. Additional properties are that for any level of noise the solution has no false positives, and that the algorithm is parameter-free, so it does not require an estimation of the energy of the *off-diagonal signal* that we need to absorb, or of the level of noise in the data.

The paper is organized as follows. In Section II, we summarize the model used to generate the signals to be recovered, which in our case are images. In Section III, we present the theory that supports the proposed strategy for dimension reduction when correlated data are used to recover the signals. Section IV explains the algorithm for carrying out the inversion efficiently. Section V shows the numerical experiments. Section VI summarizes our conclusions. The proofs of the theorems are given in A.

II. PASSIVE ARRAY IMAGING

We consider processing of passive array signals where the object to be imaged is a set of point sources at positions \vec{z}_j and (complex) amplitudes α_j , $j = 1, \dots, M$. The data used to image the object are collected at several sensors on an array; see Figure 1. The imaging system is characterized by the array aperture a , the distance L to the sources, the bandwidth B and the central wavelength λ_0 of the signals.

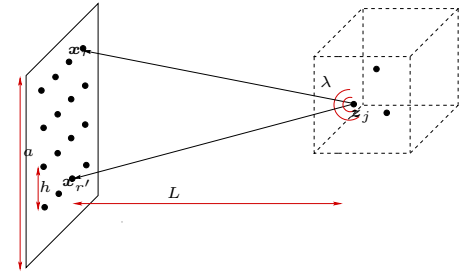


Fig. 1. General setup for passive array imaging. The source at \vec{z}_j emits a signal that is recorded at all array elements \vec{x}_r , $r = 1, \dots, N_r$.

The sources are located inside an image window IW discretized with a uniform grid of points \vec{y}_k , $k = 1, \dots, K$. Thus, the signal to be recovered is the source vector

$$\tilde{\rho} = [\tilde{\rho}_1, \dots, \tilde{\rho}_K]^T \in \mathbb{C}^K, \quad (7)$$

whose components $\tilde{\rho}_k$ correspond to the amplitudes of the M sources at the grid points \vec{y}_k , $k = 1, \dots, K$, with $K \gg M$. This vector has components $\tilde{\rho}_k = \alpha_j$ if $\vec{y}_k = \vec{z}_j$ for some $j = 1, \dots, M$, while the others are zero.

Denoting by $G(\vec{x}, \vec{y}; \omega)$ the Green's function for the propagation of a wave of angular frequency ω from point \vec{y} to point \vec{x} , we define the single-frequency Green's function vector that connects a point \vec{y} in the IW with all the sensors on the array located at points \vec{x}_r , $r = 1, \dots, N_r$, so

$$\mathbf{g}(\vec{y}; \omega) = [G(\vec{x}_1, \vec{y}; \omega), G(\vec{x}_2, \vec{y}; \omega), \dots, G(\vec{x}_{N_r}, \vec{y}; \omega)]^T \in \mathbb{C}^{N_r}.$$

In three dimensions, $G(\vec{x}, \vec{y}; \omega) = \frac{\exp\{i\omega|\vec{x} - \vec{y}|/c_0\}}{4\pi|\vec{x} - \vec{y}|}$ if the medium is homogeneous. Hence, the signals of frequencies ω_l recorded at the sensors locations \vec{x}_r are

$$b(\vec{x}_r, \omega_l) = \sum_{j=1}^M \alpha_j G(\vec{x}_r, \vec{z}_j; \omega_l), \quad r = 1, \dots, N_r.$$

They form the single-frequency data vector $\mathbf{b}(\omega_l) = [b(\vec{x}_1, \omega_l), b(\vec{x}_2, \omega_l), \dots, b(\vec{x}_{N_r}, \omega_l)]^\top \in \mathbb{C}^{N_r}$. As several frequencies ω_l , $l = 1, \dots, N_f$, are used to recover (7), all the recorded data are stacked in the multi-frequency column data vector

$$\mathbf{b} = [\mathbf{b}(\omega_1)^\top, \mathbf{b}(\omega_2)^\top, \dots, \mathbf{b}(\omega_{N_f})^\top]^\top \in \mathbb{C}^N, \text{ with } N = N_r N_f. \quad (8)$$

A. The inverse problem with linear data

When the data (8) are available and reliable, one can form the linear system

$$\mathcal{A} \boldsymbol{\rho} = \mathbf{b} \quad (9)$$

to recover (7). Here, \mathcal{A} is the $N \times K$ measurement matrix whose columns \mathbf{a}_k are the multi-frequency Green's function vectors

$$\mathbf{a}_k = \frac{1}{c_k} [g(\vec{y}_k; \omega_1)^\top, g(\vec{y}_k; \omega_2)^\top, \dots, g(\vec{y}_k; \omega_{N_f})^\top]^\top \in \mathbb{C}^N, \quad (10)$$

where c_k are scalars that normalize these vectors to have ℓ_2 -norm one, and

$$\boldsymbol{\rho} = \text{diag}(c_1, c_2, \dots, c_K) \tilde{\boldsymbol{\rho}}, \quad (11)$$

where $\tilde{\boldsymbol{\rho}}$ is given by (7). Then, one can solve (9) for the unknown vector $\boldsymbol{\rho}$ using a number of ℓ_2 and ℓ_1 inversion methods to find the sought image. In general, ℓ_2 methods are robust but the resulting resolution is low. On the other hand, ℓ_1 methods provide higher resolution but they are much more sensitive to noise in the data [27]. Hence, they cannot be used with poor quality data unless one carefully takes care of the noise.

B. The inverse problem with quadratic cross correlation data

In many instances, imaging with cross correlations helps to form better and more robust images. This is the case, for example, when one uses high frequency signals and has a low-budget measurement system with inexpensive sensors that are not able to resolve the signals well. Another situation is when the raw data (8) can be measured but it is more convenient to image with cross correlations because they help to mitigate the effects of the inhomogeneities of the medium between the sources and the sensors [3], [15]

Assume that all the cross-correlated data contained in the matrix

$$\mathbf{B} = \mathbf{b} \mathbf{b}^* \in \mathbb{C}^{N \times N} \quad (12)$$

are available for imaging. Then, one can consider the linear system

$$\mathcal{A} \mathcal{X} \mathcal{A}^* = \mathbf{B}, \quad (13)$$

and seek the matrix $\mathcal{X} = \boldsymbol{\rho} \boldsymbol{\rho}^* \in \mathbb{C}^{K \times K}$ that solves it. The unknown matrix \mathcal{X} is rank-one and, hence, one possibility is to look for a low-rank matrix by using nuclear norm minimization as it was suggested for imaging with intensities-only in [7], [4]. This is possible in theory, but it is unfeasible when the problem is large because the number of unknowns grows quadratically and, therefore, the computational cost rapidly becomes prohibitive. For example, to form an image with 1000×1000 pixels one would have to solve a system with 10^{12} unknowns.

Instead, we suggest the following strategy. We propose to vectorize both sides of (13) so

$$\text{vec}(\mathcal{A} \mathcal{X} \mathcal{A}^*) = \text{vec}(\mathbf{B}), \quad (14)$$

where $\text{vec}(\cdot)$ denotes the vectorization of a matrix formed by stacking its columns into a single column vector. Then, we use the Kronecker product \otimes , and its property $\text{vec}(PQR) = (R^T \otimes P) \text{vec}(Q)$, to express the matrix multiplications as the linear transformation

$$(\bar{\mathcal{A}} \otimes \mathcal{A}) \text{vec}(\mathcal{X}) = \text{vec}(\mathbf{B}). \quad (15)$$

With this formulation of the problem we can use an ℓ_1 minimization algorithm to form the images, which is much faster than a nuclear norm minimization algorithm that needs to compute the SVD of the iterate matrices. However, with just this approach the main obstacle is not overcome, as the dimensionality still grows quadratically with the number of unknowns K . To this effect, we propose here a dimension reduction strategy that uses the *Noise Collector* [28] to absorb a component of the data vector that does not provide extra information about the signal support. We point out that this component is not a gaussian random vector as in [28], but a deterministic vector resulting from the off-diagonal terms of \mathcal{X} that are neglected. Hence, the use of the *Noise Collector* as an effective dimension reduction tool is not a straightforward application of [28]. We need to show why, and under what conditions, it can indeed absorb the interference terms, i.e., the off-diagonal elements $\rho_i \rho_j^*$, $i \neq j$. This is the purpose of Theorems 2 and 3 given in the next section.

We point out that other ℓ_1 regularizations used in signal processing, as lasso [32], [9] or square-root lasso [2], could be used as well to absorb the contribution to the data of the off-diagonal interference terms and, thus, reduce the dimensionality of the problem. The advantage of square-root lasso over the commonly used lasso regularization is that the former, as the *Noise Collector*, does not need to know (or estimate) the level of noise, so its solution does not depend on the choice of the penalty level.

III. THE NOISE COLLECTOR AND DIMENSION REDUCTION

We explain in this section how the *Noise Collector* can be used to effectively reduce the dimensionality of problem (15). We first present it in Section III-A as an effective denoising tool for a linear problem. Theorem 1 considers the case of a generic gaussian noise vector \mathbf{e} . This is the result of [28], recalled here for completeness. Theorem 2 is then introduced to address the problem of absorbing a deterministic noise

vector. This is a new original result with slightly weaker assumptions than Theorem 1. In particular, Theorem 2 requires the noise vector to be decoherent only with respect to the columns of the sensing matrix that do not correspond to the signal support. Although this is not essential when the noise is random, it turns out to be crucial for problem (15) because, in this case, the deterministic noise vector is not decoherent with respect to the columns that correspond to the signal support. The specific assumptions needed for the *Noise Collector* to be employed as an effective dimension reduction tool for problem (15) are given in Theorem 3, whose proof relies in verifying that the assumptions of Theorem 2 indeed hold in this case.

A. The Noise Collector

The *Noise Collector* [28] is a method to find the vector $\chi \in \mathbb{C}^K$ in

$$T\chi = d_0 + e, \quad (16)$$

from highly incomplete measurement data $d = d_0 + e \in \mathbb{C}^N$ possibly corrupted by noise $e \in \mathbb{C}^N$, where $1 \ll N < K$. Here, T is a general measurement matrix of size $N \times K$, whose columns have unit length. The main results in [28] ensure that we can still recover the support of χ when the data is noisy by looking at the support of χ_τ found as

$$(\chi_\tau, \eta_\tau) = \arg \min_{\chi, \eta} (\tau \|\chi\|_{\ell_1} + \|\eta\|_{\ell_1}), \quad (17)$$

subject to $T\chi + C\eta = d$,

with an $O(1)$ no-phantom weight τ , and a *Noise Collector* matrix $C \in \mathbb{C}^{N \times \Sigma}$ with $\Sigma = N^\beta$, for $\beta > 1$. If the noise e is Gaussian, then the columns of C can be chosen independently and at random on the unit sphere \mathbb{S}^{N-1} . The weight $\tau > 1$ is chosen so it is expensive to approximate e with the columns of T , but it cannot be taken too large because then we lose the signal χ that gets absorbed by the *Noise Collector* as well. Intuitively, τ is a measure of the rate at which the signal is lost as the noise increases. For practical purposes, τ is chosen as the minimal value for which $\chi = 0$ when the data is pure noise, i.e., when $d_0 = 0$. The key property is that the optimal value of τ does not depend on the level of noise and, therefore, it is chosen in advance, before the *Noise Collector* is used for a specific task. We have the following result.

Theorem 1: [28] Fix $\beta > 1$, and draw $\Sigma = N^\beta$ columns to form the Noise Collector C , independently, from the uniform distribution on \mathbb{S}^{N-1} . Let χ be an M -sparse solution of the noiseless system $T\chi = d_0$, and χ_τ the solution of (17) with gaussian noise e so $d = d_0 + e$. Denote the ratio of minimum to maximum significant values of χ as

$$\gamma = \min_{i \in \text{supp}(\chi)} \frac{|\chi_i|}{\|\chi\|_{\ell_\infty}}. \quad (18)$$

Assume that the columns of T are incoherent, so that

$$|\langle t_i, t_j \rangle| \leq \frac{1}{3M} \text{ for all } i \text{ and } j. \quad (19)$$

Then, for any $\kappa > 0$, there are constants $\tau = \tau(\kappa, \beta)$, $c_1 = c_1(\kappa, \beta, \gamma)$, and $N_0 = N_0(\kappa, \beta)$ such that, if the noise level satisfies

$$\max(1, \|e\|_{\ell_2}) \leq c_1 \frac{\|d_0\|_{\ell_2}^2}{\|\chi\|_{\ell_1}} \sqrt{\frac{N}{\ln N}}, \quad (20)$$

then $\text{supp}(\chi_\tau) = \text{supp}(\chi)$ for all $N > N_0$ with probability $1 - 1/N^\kappa$.

To gain a better understanding of this theorem, let us consider the case where T is the identity matrix (the classical denoising problem) and all coefficients of $d_0 = \chi$ are either 1 or 0. Then $\|d_0\|_{\ell_2}^2 = \|\chi\|_{\ell_1} = M$. In this case, an acceptable level of noise is

$$\|e\|_{\ell_2} \lesssim \|d_0\|_{\ell_2} \sqrt{\frac{N}{M \ln N}} \sim \sqrt{\frac{N}{\ln N}}. \quad (21)$$

The estimate (21) implies that we can handle more noise as we increase the number of measurements. This holds for two reasons. Firstly, a typical noise vector e is almost orthogonal to the columns of T so, using bounds on the maximum of a family of sub-gaussian random variables, we obtain

$$|\langle t_i, e \rangle| \leq c_0 \sqrt{\frac{\ln N}{N}} \|e\|_{\ell_2} \quad (22)$$

for some $c_0 = c_0(\kappa)$ with probability $1 - 1/N^\kappa$. In particular, a typical noise vector e is almost orthogonal to the signal subspace V . More formally, suppose V is the M -dimensional subspace spanned by the column vectors t_j with j in the support of χ , and let $W = V^\perp$ be the orthogonal complement to V . Consider the orthogonal decomposition $e = e^v + e^w$, such that e^v is in V and e^w is in W . Then,

$$\|e^v\|_{\ell_2} \lesssim \sqrt{\frac{M}{N}} \|e\|_{\ell_2}$$

with high probability that tends to 1, as $N \rightarrow \infty$. In Theorem 1, a quantitative estimate of this convergence is $1 - 1/N^\kappa$. It means that if a signal is sparse so $M \ll N$, then we can recover it for very low signal-to-noise ratios. Secondly, and more importantly, if the columns of the noise collector C are also almost orthogonal to the signal subspace, then it is too expensive to approximate the signal d_0 with the columns of C and, hence, we have to use the columns of the measurement matrix T . If we draw the columns of C , independently, from the uniform distribution on \mathbb{S}^{N-1} , then they will be almost orthogonal to the signal subspace with high probability. It is again estimated as $1 - 1/N^\kappa$ in Theorem 1. Finally, the incoherence condition (19) implies that it is too expensive to approximate the signal d_0 with columns T that are not in the support of χ and, hence, there are no false positives.

In Theorem 1 we used randomness twice: the noise vector e was random and the columns of the noise collector were drawn at random. Note that in both cases randomness could be replaced by deterministic conditions requiring that e and the columns of C are almost orthogonal to the signal subspace. It is natural to assume that the noise vector e is a random variable and, as we explain in [28], the columns of C are random because it is hard to construct a deterministic C that satisfies the almost orthogonality conditions. In the present work we still construct the matrix C randomly, but we treat the vector e as deterministic (see Theorem 2). Inspection of the proofs in [28] shows that the only condition on e we need to verify from Theorem 1 is (22). Thus, the next Theorem is a reformulation of Theorem 1 for a deterministic noise vector e . The proof is given in Appendix B.

Theorem 2: Assume conditions on χ , T , and \mathcal{C} are as in Theorem 1 and define γ as in (18). Then, for any $\kappa > 0$, there are constants $\tau_0 = \tau_0(\kappa, \beta)$, $c_0 = c_0(\kappa, \beta)$, and $\mathcal{N}_0 = \mathcal{N}_0(\kappa, \beta, \gamma)$, $\alpha = \alpha(c_0, \kappa, \beta)$ such that the following two claims hold.

(i) If e satisfies (22) for all t_i , $i \notin \text{supp}(\chi)$; all columns of T satisfy

$$|\langle t_i, t_j \rangle| \leq c_0 \frac{\sqrt{\ln \mathcal{N}}}{\sqrt{\mathcal{N}}} \quad (23)$$

for all i and j ; the sparsity M is such that

$$M \leq \alpha \frac{\sqrt{\mathcal{N}}}{\sqrt{\ln \mathcal{N}}}; \quad (24)$$

and $\tau \geq \tau_0$, then $\text{supp}(\chi_\tau) \subset \text{supp}(\chi)$ with probability $1 - 1/\mathcal{N}^\kappa$.

(ii) If, in addition, the noise is not large, so

$$|\langle t_m, e \rangle| \leq \min_{i \in \text{supp}(\chi)} |\chi_i|/2 \quad (25)$$

for all t_m , $m \in \text{supp}(\chi)$, and

$$\|e\|_{\ell_2} \leq c_1 \|\chi\|_{\ell_1} \quad (26)$$

for some c_1 , then $\text{supp}(\chi) = \text{supp}(\chi_\tau)$ for all $\mathcal{N} > \mathcal{N}_0$ with probability $1 - 1/\mathcal{N}^\kappa$.

In contrast to Theorem 1, we require in Theorem 2 condition (22) to hold only for t_i , $i \notin \text{supp}(\chi)$, that is for the columns of T outside the support of χ . For the columns inside the support, $i \in \text{supp}(\chi)$, we relax condition (22) to condition (25) which is easier to be satisfied. Thus Theorem 2 has slightly weaker assumptions than Theorem 1. For a random e this weakening is not essential, because one needs to know the support of χ in advance. It turns out that for our e this weakening will become important (see Remark 3 in the end of Appendix C).

B. Dimension reduction for cross correlation data

The $N^2 \times K^2$ linear problem (15) that uses quadratic cross correlation data is notoriously hard to solve due to its high dimensionality. Therefore, we propose the following strategy for robust dimensionality reduction. The idea is to treat the contribution of the off-diagonal elements of $X = \rho \rho^* \in \mathbb{C}^{K \times K}$ as *noise* and, thus, use the *Noise Collector* to absorb it. Namely, we define

$$\chi = \text{diag}(X) = [|\rho_1|^2, |\rho_2|^2, \dots, |\rho_K|^2]^T, \quad (27)$$

and re-write (15) as

$$T\chi + \mathcal{C}\eta = d, \quad (28)$$

where we replace the off-diagonal elements by the Noise Collector term $\mathcal{C}\eta$ and

$$T = (\bar{A} \otimes A)\chi \quad (29)$$

contains only the K columns of $\bar{A} \otimes A$ corresponding to χ . Thus, the size of χ is \mathcal{K} and the size of T is $\mathcal{N} \times \mathcal{K}$, with $\mathcal{K} = K$ and $\mathcal{N} = N^2$. In practice, the measurements may be subsampled as well, so the size of the system can be further reduced to $\mathcal{N} \times \mathcal{K}$, with $\mathcal{N} = O(N)$ and $\mathcal{K} = K$. Notice

that, as a result of the proposed dimension reduction strategy, the resulting χ has nonnegative entries and, hence, specific algorithms for sparse recovery of nonnegative signals could be applied [16], [22], [23].

Problem (28) can be understood as an exact linearization of the classical phase retrieval problem, where all the interference terms $\rho_i \rho_j^*$ for $i \neq j$ are absorbed in $\mathcal{C}\eta$, with η being an unwanted vector considered to be noise in this formulation. In other words, the phase retrieval problem with K unknowns has been transformed to the linear problem (28) that also has K unknowns. Note, though, that in phase retrieval only autocorrelation measurements are considered, while in (28) we also use cross-correlated measurements.

In the next theorem we use all the measurements $d \in \mathbb{C}^{\mathcal{N}}$, so $\mathcal{N} = N^2$ in (28). This is done for simplicity of presentation, but in practice $\mathcal{N} = O(N)$ measurements are enough. We will choose a solution of (28) using (17). As in Theorems 1 and 2, the vector η in (28) has \mathcal{N}^β entries that do not have physical meaning. Its only purpose is to absorb the off-diagonal contributions in $e = d - T\chi$. We point out that the magnitude of e is not small if $M \geq 2$. Indeed, the contribution of $\chi = \text{diag}(X)$ to the data d is of order M , while the contribution of the off-diagonal terms of X is of order M^2 . Furthermore, the vector e is not independent of χ anymore.

Theorem 3: Fix $|\rho_i|$. Suppose the phases $\rho_i/|\rho_i|$ are independent and uniformly distributed on the (complex) unit circle. Suppose X is a solution of (15), $\chi = \text{diag}(X)$ is M -sparse, and $T = (\bar{A} \otimes A)\chi : \mathbb{C}^{\mathcal{K}} \rightarrow \mathbb{C}^{\mathcal{N}}$, $\mathcal{K} = K$ and $\mathcal{N} = N^2$. Fix $\beta > 1$, and draw $\Sigma = \mathcal{N}^\beta$ columns for \mathcal{C} , independently, from the uniform distribution on \mathbb{S}^{N^2-1} . Denote

$$\Delta = \sqrt{\mathcal{N}} \max_{i \neq j} |\langle a_i, a_j \rangle|, \quad (30)$$

and define γ as in (18). Then, for any $\kappa > 0$, there are constants $\alpha = \alpha(\kappa, \gamma, \Delta)$, $\tau = \tau(\kappa, \beta)$, and $\mathcal{N}_0 = \mathcal{N}_0(\kappa, \beta, \gamma, \Delta)$ such that the following holds. If

$$M \leq \alpha N / \sqrt{\ln \mathcal{N}} \quad (31)$$

and χ_τ is the solution of (17), then $\text{supp}(\chi) = \text{supp}(\chi_\tau)$ for all $\mathcal{N} > \mathcal{N}_0$ with probability $1 - 1/\mathcal{N}^\kappa$.

Note that the main role of \mathcal{N}_0 is to absorb all constants so that simple expressions of the form $1 - 1/\mathcal{N}^\kappa$ arise (see also Remark 1 in Appendix A). All the key parameters κ, β, γ and Δ are of order $O(1)$. Indeed, Δ is $O(1)$ because $|\langle a_i, a_j \rangle|$ is $O(1/\sqrt{\mathcal{N}})$, γ is also typically $O(1)$ as it is the ratio between the smallest and the largest value of the unknown, β is an $O(1)$ parameter that determines the dimension $\Sigma = \mathcal{N}^\beta$ of the *Noise Collector*, and κ is the algebraic decay rate of the probability estimate $1 - 1/\mathcal{N}^\kappa$, which is also $O(1)$.

The proof of Theorem 3 is given in Appendix C. In Theorem 3 the scaling for sparse recovery is (31). This linear in M scaling is in good agreement with our numerical experiments, see Figure 7. In order to obtain this scaling we introduced our probabilistic framework in Theorem 3 assuming that the phases of the signals are random. The idea is that a vector with random phases better describes a typical signal in many applications. Our assumption about the type of randomness is not essential and it is only used to simplify

the proof of this theorem. It is, however, essential to assume randomness so as to obtain the linear in M scaling (31). The dimension reduction still could be done without introducing the probabilistic framework. If no randomness is assumed, then the scaling for sparse recovery is not linear as in (31). It is more conservative: $M \leq \alpha\sqrt{N}/\sqrt{\ln N}$, and it does not agree with our numerical experiments. We state and prove a deterministic version of Theorem 3 in Appendix D for completeness.

IV. ALGORITHMIC IMPLEMENTATION

A key point of the proposed strategy is that the M -sparse solution of (28) can be effectively found by solving the minimization problem

$$(\chi_\tau, \eta_\tau) = \arg \min_{\chi, \eta} (\tau \|\chi\|_{\ell_1} + \|\eta\|_{\ell_1}), \quad (32)$$

subject to $T\chi + C\eta = d$,

with an $O(1)$ no-phantom weight τ . Here, T is an $\mathcal{N} \times \mathcal{K}$ matrix, C an $\mathcal{N} \times \mathcal{N}^\beta$ matrix, with β close to one, χ is a $\mathcal{K} \times 1$ vector and η is a $\mathcal{N}^\beta \times 1$ vector. The main property of this approach is that if the matrix T is incoherent enough, so its columns satisfy assumption (19) of Theorem 1, the ℓ_1 -norm minimal solution of (32) has a zero false discovery rate for any level of noise, with probability that tends to one as the dimension of the data \mathcal{N} increases to infinity.

To find the minimizer in (32), we define the function

$$F(\chi, \eta, z) = \lambda (\tau \|\chi\|_{\ell_1} + \|\eta\|_{\ell_1}) \quad (33)$$

$$+ \frac{1}{2} \|T\chi + C\eta - d\|_{\ell_2}^2 + \langle z, d - T\chi - C\eta \rangle$$

for a no-phantom weight τ , and determine the solution as

$$\max_z \min_{\chi, \eta} F(\chi, \eta, z). \quad (34)$$

This strategy finds the minimum in (32) exactly for all values of the regularization parameter λ . Thus, the method is fully automated, meaning that it has no tuning parameters. To determine the exact extremum in (34), we use the iterative soft thresholding algorithm GeLMA [26] that works as follows.

Pick a value for the no-phantom weight τ ; for optimal results calibrate τ to be the smallest value for which $\chi = 0$ when the algorithm is fed with pure noise. In our numerical experiments we use $\tau = 2$. Next, pick a value for the regularization parameter, for example $\lambda = 1$, and choose step sizes $\Delta t_1 < 2/\|T\|^2$ and $\Delta t_2 < \lambda/\|T\|$. Set $\chi_0 = 0$, $\eta_0 = 0$, $z_0 = 0$, and iterate for $k \geq 0$:

$$\begin{aligned} r &= d - T\chi_k - C\eta_k, \\ \chi_{k+1} &= \mathcal{S}_{\tau \lambda \Delta t_1}(\chi_k + \Delta t_1 T^*(z_k + r)), \\ \eta_{k+1} &= \mathcal{S}_{\lambda \Delta t_1}(\eta_k + \Delta t_1 C^*(z_k + r)), \\ z_{k+1} &= z_k + \Delta t_2 r, \end{aligned} \quad (35)$$

where $\mathcal{S}_r(y_i) = \text{sign}(y_i) \max\{0, |y_i| - r\}$. Terminate the iterations when the distance $\|\chi_{k+1} - \chi_k\|$ between two consecutive iterates is below a given tolerance.

¹Choosing two step sizes instead of the smaller one Δt_1 improves the convergence speed.

A. The Noise Collector: construction and properties

To construct the *Noise Collector* matrix $C \in \mathbb{C}^{\mathcal{N} \times \mathcal{N}^\beta}$ that satisfies the assumptions of Theorem 1 one could draw \mathcal{N}^β normally distributed \mathcal{N} -dimensional vectors, normalized to unit length. Thus, the additional computational cost incurred for implementing the *Noise Collector* in (35), due to the terms $C\eta_k$ and $C^*(z_k + r)$, would be $O(\mathcal{N}^{\beta+1})$, which is not very large as we use $\beta \approx 1.5$ in practice. The computational cost of (35) without the *Noise Collector* mainly comes from the matrix vector multiplications $T\chi_k$ which can be done in $O(\mathcal{N}\mathcal{K})$ operations and, typically, $\mathcal{K} \gg \mathcal{N}$.

To further reduce the additional computational time and memory requirements we use a different construction procedure that exploits the properties of circulant matrices. The idea is to draw instead a few normally distributed \mathcal{N} -dimensional vectors of length one, and construct from each one of them a circulant matrix of dimension $\mathcal{N} \times \mathcal{N}$. The columns of these matrices are still independent and uniformly distributed on $\mathbb{S}^{\mathcal{N}-1}$, so they satisfy the assumptions of Theorem 1. The full *Noise Collector* matrix is then formed by concatenating these circulant matrices together.

More precisely, the *Noise Collector* construction is done in the following way. We draw $\mathcal{N}^{\beta-1}$ normally distributed \mathcal{N} -dimensional vectors, normalized to unit length. These are the generating vectors of the *Noise Collector*. To these vectors are associated $\mathcal{N}^{\beta-1}$ circulant matrices $C_i \in \mathbb{C}^{\mathcal{N} \times \mathcal{N}}$, $i = 1, \dots, \mathcal{N}^{\beta-1}$, and the *Noise Collector* matrix is constructed by concatenation of these $\mathcal{N}^{\beta-1}$ matrices, so

$$C = [C_1 | C_2 | C_3 | \dots | C_{\mathcal{N}^{\beta-1}}] \in \mathbb{C}^{\mathcal{N} \times \mathcal{N}^\beta}.$$

We point out that the *Noise Collector* matrix C is not stored, only the $\mathcal{N}^{\beta-1}$ generating vectors are saved in memory. On the other hand, the matrix vector multiplications $C\eta_k$ and $C^*(z_k + r)$ in (35) can be computed using these generating vectors and FFTs [17]. This makes the complexity associated to the *Noise Collector* $O(\mathcal{N}^\beta \log(\mathcal{N}))$.

To explain this further, we recall briefly below how a matrix vector multiplication can be performed using the FFT for a circulant matrix. For a generating vector $c = [c_0, c_1, \dots, c_{\mathcal{N}-1}]$, the C_i circulant matrix takes the form

$$C_i = \begin{bmatrix} c_0 & c_{\mathcal{N}-1} & \dots & c_1 \\ c_1 & c_0 & \dots & c_2 \\ \vdots & & \ddots & \vdots \\ c_{\mathcal{N}-1} & c_{\mathcal{N}-2} & \dots & c_0 \end{bmatrix}.$$

This matrix can be diagonalized by the Discrete Fourier Transform (DFT) matrix, i.e.,

$$C_i = \mathcal{F} \Lambda \mathcal{F}^{-1}$$

where \mathcal{F} is the DFT matrix, \mathcal{F}^{-1} is its inverse, and Λ is a diagonal matrix such that $\Lambda = \text{diag}(\mathcal{F}c)$, where c is the generating vector. Thus, a matrix vector multiplication $C_i\eta$ is performed as follows: (i) compute $\hat{\eta} = \mathcal{F}^{-1}\eta$, the inverse DFT of η in $\mathcal{N} \log(\mathcal{N})$ operations, (ii) compute the eigenvalues of C_i as the DFT of c , and component wise multiply the result with $\hat{\eta}$ (this step can also be done in $\mathcal{N} \log(\mathcal{N})$ operations),

and (iii) compute the FFT of the vector resulting from step (ii) in, again, $\mathcal{N} \log(\mathcal{N})$ operations.

Consequently, the cost of performing the multiplication $\mathcal{C}\eta_k$ is $\mathcal{N}^{\beta-1} \mathcal{N} \log(\mathcal{N}) = \mathcal{N}^\beta \log(\mathcal{N})$. As the cost of finding the solution without the *Noise Collector* is $O(\mathcal{N}\mathcal{K})$ due to the terms $T\chi_k$, the additional cost due to the *Noise Collector* is negligible since $\mathcal{K} \gg \mathcal{N}^{\beta-1} \log(\mathcal{N})$ because, typically, $\mathcal{K} \gg \mathcal{N}$ and $\beta \approx 1.5$.

We emphasise that in the dimension reduced signal recovery problem (28) that uses cross-correlated data, $\mathcal{K} = K$ and $\mathcal{N} = O(N)$, so the computational cost of finding its sparsest solution is $O(NK)$, as the original problem that uses linear data.

V. NUMERICAL RESULTS

We consider processing of passive array signals. We seek to determine the positions \mathbf{z}_j and the complex amplitudes α_j of M point sources, $j = 1, \dots, M$, from measurements of polychromatic signals on an array of receivers; see Figure 1. The source imaging problem is considered here for simplicity. The active array imaging problem can be cast under the same linear algebra framework even when multiple scattering is important [8].

The array consists of $N_r = 21$ receivers located at $x_r = -\frac{a}{2} + \frac{r-1}{N_r-1}a$, $r = 1, \dots, N_r$, where $a = 100\lambda$ is the array aperture. The imaging window (IW) is at range $L = 100\lambda$ from the array and the bandwidth $B = f_0/3$ of the emitted pulse is $1/3$ of the central frequency f_0 , so the resolution in range is $c/B = 3\lambda$ while in cross-range it is $\lambda L/a = \lambda$. We consider a high frequency microwave imaging regime with central frequency $f_0 = 60\text{GHz}$ corresponding to $\lambda_0 = 5\text{mm}$. We make measurements for $N_f = 21$ equally spaced frequencies spanning a bandwidth $B = 20\text{GHz}$. The array aperture is $a = 50\text{cm}$, and the distance from the array to the center of the IW is $L = 50\text{cm}$. Then, the resolution is $\lambda_0 L/a = 5\text{mm}$ in the cross-range (direction parallel to the array) and $c_0/B = 15\text{mm}$ in range (direction of propagation). These parameters are typical in microwave scanning technology [24].

We consider an IW with $K = 1681$ pixels which makes the dimension of $X = \rho\rho^*$ equal to $K^2 = 2825761$. The pixel dimensions, i.e., the resolution of the imaging system, is $5\text{mm} \times 15\text{mm}$. The total number of measurements is $N = N_r N_f = 441$. Thus, we can form $N^2 = 194481$ cross-correlations over frequencies and locations.

Let us first note that with these values for N and K , which in fact are not big, we cannot form the full $K^2 \times K^2$ matrix $(\bar{\mathcal{A}} \otimes \mathcal{A})$ so as to solve (15) for the $K^2 \times 1$ vector $\text{vec}(X)$ because of its huge dimensions. Instead, we propose to reduce the dimensionality of the problem to K unknowns corresponding to $\text{diag}(X)$, neglecting all the off-diagonal terms that correspond to the interference terms $\rho_k \rho_{k'}^*$ for $k \neq k'$. Their contributions to the cross-correlated data are treated as *noise*, which is absorbed in a fictitious vector η using a *Noise Collector*. We stress that this *noise* is never small if $M \geq 2$, as its contribution to the the cross-correlated data is of order $O(M^2)$, while the contribution of $\text{diag}(X)$ is only of order $O(M)$.

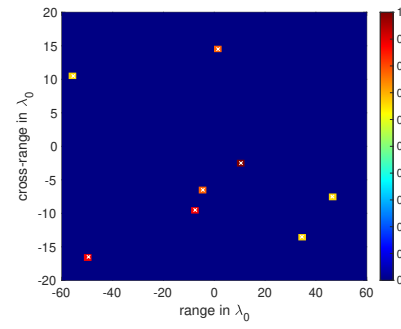


Fig. 2. The true $\chi = \text{diag}(X) = \text{diag}(\rho\rho^*)$, i.e., the absolute values squared of the point sources amplitudes.. The dimension of the image is $K = 1681$.

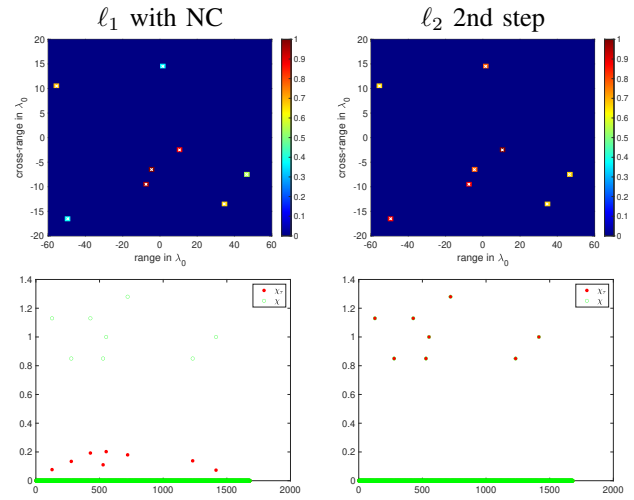


Fig. 3. Imaging $M = 8$ sources using correlations and the NC. The dimension of the image is $K = 1681$. The dimension of the linear data is $N = 441$. The ℓ_1 images are obtained using $21N$ of the N^2 correlation data. Noise free data.

In the following examples, we consider imaging of $M = 8$ point sources that are on the grid; see Fig. 2. If the sources are off-grid the resulting modeling error is also absorbed by the *Noise Collector*. In the following numerical experiments, we only use $\mathcal{N} = 21N$ cross-correlated data picked at random instead of the N^2 cross-correlated data which are in principle available. This reduces even more the dimensionality of the problem that we are solving.

In Fig. 3, we present the results when the used data is noise-free. The left column shows the results when we use the ℓ_1 algorithm (35); the top plot is the recovered image and the bottom plot the recovered $\chi = \text{diag}(X) = \text{diag}(\rho\rho^*)$ vector. The support of the sources is exact but the amplitudes are not. If it is important for an application to recover the amplitudes with precision, one can consider in a second step the full problem (15) for $\text{vec}(X)$ with all the interference terms $\rho_k \rho_{k'}^*$ for $k \neq k'$, but restricted to the exact support found in the first step. If there is no noise in the data, this second step finds the exact values of the amplitudes efficiently using an ℓ_2 minimization method; see the right column of Fig. 3.

In Figs. 4 and 5 we consider the same configuration of sources but we add white Gaussian noise to the data. The

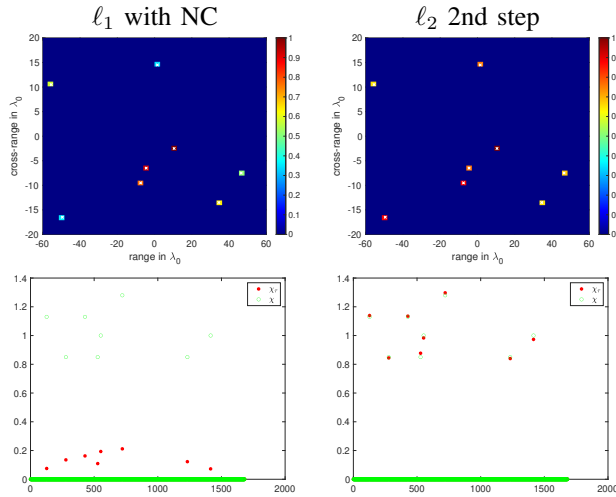


Fig. 4. Imaging $M = 8$ sources using correlations and the NC. The dimension of the image is $K = 1681$. The dimension of the linear data is $N = 441$. The ℓ_1 images are obtained using $21N$ of the N^2 correlation data. Data with 10dB SNR.

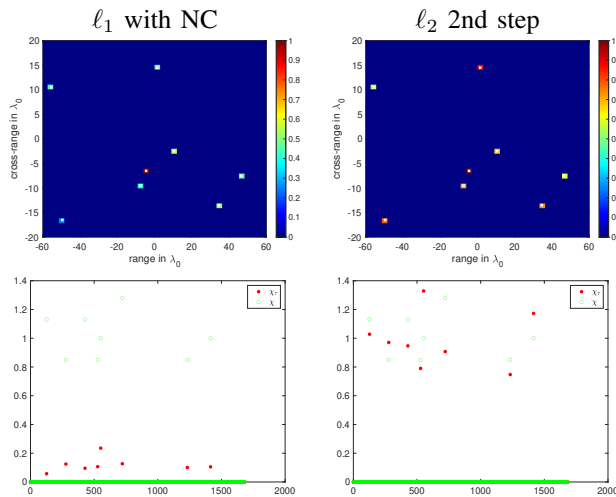


Fig. 5. Imaging $M = 8$ sources using correlations and the NC. The dimension of the image is $K = 1681$. The dimension of the linear data is $N = 441$. The ℓ_1 images are obtained using $21N$ of the N^2 correlation data. Data with 0dB SNR.

resulting SNR values are 10dB and 0dB, respectively. In both cases, the solutions obtained in the first step look very similar to the one obtained in Fig. 3 for noise free data. This is so, because the noise in the data is dominated by the neglected interference terms. The actual effect of the additive noise is only seen in the 2nd step when we solve for $\text{vec}(X)$, restricted to the support, using an ℓ_2 minimization method. Indeed, when the data are noisy we cannot recover the exact values of the amplitudes. Still, since an ℓ_2 method is used on the correct support, the reconstructions are extremely robust and give very good results, even when the SNR is 0dB.

To illustrate the robustness of the reconstructions of the entire matrix $X = \rho\rho^*$ we also plot in Fig. 6 the angle of X_τ compared to the angle of X restricted on the support recovered during the first step. We get an exact reconstruction for noise-free data. The error in the reconstruction increases as the SNR

decreases but the results are very satisfactory even for the 0dB SNR case.

Again, the big advantage of the proposed ℓ_1 minimization approach that seeks only for the components of $\text{diag}(X)$, and uses a *Noise Collector* to absorb the interference terms that are treated as noise, is that it is linear in the number of pixels K instead of quadratic. This allows us to consider large scale problems. Moreover, as we observed in the results of Figs. 3 to 6, the number of data \mathcal{N} used to recover the images do not need to be N^2 , but only a multiple of N . Note that in the examples shown here $K = 1681$ while $N = 441$. The full matrix $N^2 \times K^2$ would require 9000 gigabytes of RAM which makes the problem of solving for X practically unfeasible.

In Fig. 7 we illustrate the performance of the proposed ℓ_1 approach for different sparsity levels M and data sizes \mathcal{N} . There is no additive noise added to the data in this figure. Success in recovering the true support of the unknown χ corresponds to the value 1 (yellow) and failure to 0 (blue). The small phase transition zone (green) contains intermediate values. The red line is the estimate $\sqrt{\mathcal{N}}/(2\sqrt{\ln \mathcal{N}})$. These results are obtained by averaging over 10 realizations.

VI. CONCLUSIONS

In this paper, we consider the problem of sparse signal recovery from cross correlation measurements. The unknown in this case is the correlated matrix signal $X = \rho\rho^*$ whose dimension grows quadratically with the size K of ρ and, hence, inversion becomes computationally unfeasible as K becomes large. To overcome this issue, we propose a novel dimension reduction strategy. Specifically, we vectorize the problem and consider only the diagonal terms $|\rho_i|^2$ of X as unknown. The contribution of the off-diagonal interference terms $\rho_i\rho_j^*$ for $i \neq j$ is treated as noise, which is absorbed using the *Noise Collector* approach introduced in [28]. In this way, we are able to relate the (noisy) data to the diagonal terms $|\rho_i|^2$ of X through a linear transformation. This allows us to recover the signal exactly using efficient ℓ_1 -minimization algorithms. The cost of solving this dimension reduced problem is similar to the one using linear data. Furthermore, our numerical experiments show that the suggested approach is robust with respect to additive noise in the data.

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APPENDIX

A. A special case of the Dvoretzky-Milman's theorem

The proof of Theorem 2 uses the following special case of the Dvoretzky-Milman's theorem [25] proved, e.g. in [28] as Lemma 1. For clarity of presentation we work in \mathbb{R} instead of \mathbb{C} .

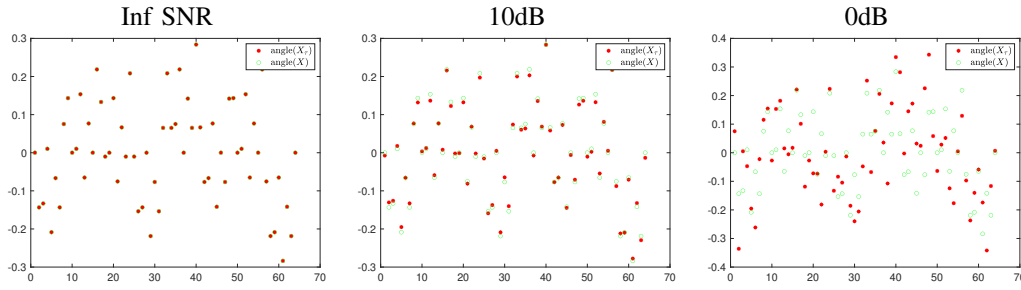


Fig. 6. Imaging $M = 8$ sources using correlations and the NC. The dimension of the image is $K = 1681$. The dimension of the linear data is $N = 441$. The angle of the components of X_τ compared to angle of the components of the true X restricted on the support.

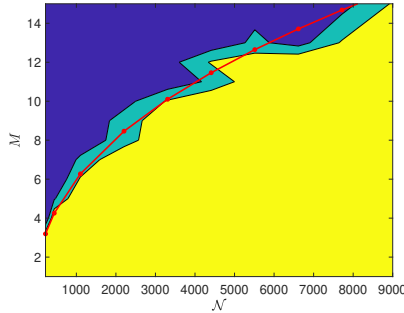


Fig. 7. Algorithm performance for exact support recovery during the first step using ℓ_1 and the Noise Collector. Success corresponds to the value 1 (yellow) and failure to 0 (blue). The small phase transition zone (green) contains intermediate values. The red line is the estimate $\sqrt{N}/(2\sqrt{\ln N})$. Ordinate and abscissa are the data used $N = N^2$ and the sparsity M .

Theorem 4: Fix $\beta > 1$, a unit vector \mathbf{b} , and draw $\Sigma = \mathcal{N}^\beta$ columns of \mathcal{C} , independently, from the uniform distribution on \mathbb{S}^{N-1} . Define H_1 as a convex hull of the columns of \mathcal{C} . More formally,

$$H_1 = \left\{ x \in \mathbb{R}^N \left| x = \sum_{i=1}^{\Sigma} \xi_i \mathbf{c}_i, \sum_{i=1}^{\Sigma} |\xi_i| \leq 1 \right. \right\},$$

Then

$$c\sqrt{\frac{\ln N}{N}} \mathbf{b} \notin H_1, \text{ with probability } 1 - 2\mathcal{N}^{\beta - \frac{c^2}{2}}, \quad (36)$$

and

$$\theta \mathbf{b} \in H_1, \text{ with probability } 1 - 4\mathcal{N} \exp\left(-\mathcal{N}^{\phi^2}/12\right), \quad (37)$$

where

$$\theta = \phi\sqrt{\frac{\ln N}{N}}, \phi = \sqrt{(\beta - 1)/2}.$$

Proof of Theorem 4. Using the rotational invariance of all our probability distributions, inequality (36) is true if

$$\mathbb{P}(\max_i |\langle \mathbf{c}_i, \mathbf{b} \rangle| \geq c\sqrt{\ln N}/\sqrt{N}) \leq 2\mathcal{N}^{\beta - \frac{c^2}{2}},$$

where $\mathbf{c}_i, i = 1, 2, \dots, \mathcal{N}^\beta$ are uniformly distributed on \mathbb{S}^{N-1} , and we can assume $\mathbf{b} = (1, 0, \dots, 0)$. Denote the event

$$\Omega_t = \left\{ \max_i |\langle \mathbf{c}_i, \mathbf{b} \rangle| \geq t/\sqrt{N} \right\}.$$

Since $\mathbb{P}(|\langle \mathbf{c}_i, \mathbf{b} \rangle| \geq t/\sqrt{N}) \leq 2\exp(-t^2/2)$ for each \mathbf{c}_i , we can use the union bound to obtain $\mathbb{P}(\Omega_t) \leq 2\mathcal{N}^\beta \exp(-t^2/2)$. Choosing $t = c\sqrt{\ln N}$ we get (36).

We will now prove (37). The idea is to find sufficient conditions so that $3\mathcal{N}$ columns, say, $\mathbf{c}_j, j \in S$ of \mathcal{C} satisfy

$$\mathbf{b} = \sum_{j \in S} \alpha_j \mathbf{c}_j, \sum_{j \in S} |\alpha_j| \leq \frac{1}{\theta}. \quad (38)$$

The first of these sufficient conditions is:

$$\min_{j \in S} \langle \mathbf{c}_j, \mathbf{b} \rangle \geq \theta. \quad (39)$$

Let us discuss this condition (39). It means that the vectors $\mathbf{c}_j, j \in S$ and the vector \mathbf{b} lie on the same side of the plane $x_1 = \theta$. We now observe that vectors with (39) will satisfy (38) if

$$0 = \sum_{j \in S} \alpha_j \tilde{\mathbf{c}}_j, \alpha_j \geq 0, \quad (40)$$

where $\tilde{\mathbf{c}}_j$ are projections of \mathbf{c}_j onto the plane $x_1 = 0$. In turn, if there does not exist a vector \mathbf{e} such that $\mathbf{e} \perp \mathbf{b}$ and $\langle \mathbf{c}_j, \mathbf{e} \rangle \geq 0$ for all $j \in S$, then condition (40) holds. Our vectors $\mathbf{c}_j, j \in S$ are independent (as it follows from their selection below) with rotationally invariant probability measure. By Wendel's Theorem [36] the second sufficient condition holds with probability

$$\frac{1}{2^{3\mathcal{N}-1}} \sum_{k=0}^{2\mathcal{N}-1} \binom{3\mathcal{N}-1}{k} \geq 1 - e^{-2\mathcal{N}/5}. \quad (41)$$

In the last inequality we used the probability that out of $2m$ coin tosses, the number of heads is less than $m - t$, is at most $e^{-t^2/(m+t)}$. It remains to compute the probability (39) holds.

For each \mathbf{c}_i

$$\mathbb{P}\left(\langle \mathbf{c}_i, \mathbf{b} \rangle \geq \frac{t}{\sqrt{N}}\right) = \frac{1}{\sqrt{2\pi}} \int_t^\infty e^{-\frac{x^2}{2}} dx \geq \frac{1}{4} e^{-t^2}.$$

Split the index set $1, 2, \dots, \Sigma$ into $3\mathcal{N}$ non-overlapping subsets $S_k, k = 1, 2, \dots, 3\mathcal{N}$ of size $\mathcal{N}^{\beta-1}/3$. For each S_k

$$\mathbb{P}\left(\max_{i \in S_k} \langle \mathbf{c}_i, \mathbf{b} \rangle \leq \frac{\phi\sqrt{\ln N}}{\sqrt{N}}\right) \leq \left(1 - \frac{1}{4\mathcal{N}^{\phi^2}}\right)^{\frac{\mathcal{N}^{\beta-1}}{3}} \leq e^{-\frac{1}{12} \mathcal{N}^{\phi^2}}$$

for $\phi = \sqrt{(\beta - 1)/2}$. For the set S in (39) we will take \mathbf{c}_i with the largest $\langle \mathbf{c}_i, \mathbf{b} \rangle$ from each $S_k, k = 1, 2, \dots, 3\mathcal{N}$. Note that they are independent, as required above. By independence

$$\mathbb{P}((39) \text{ holds}) = \Pi_{k=1}^{3\mathcal{N}} \mathbb{P}(\max_{i \in S_k} \langle \mathbf{c}_i, \mathbf{b} \rangle \geq \phi\sqrt{\ln N}/\sqrt{N}).$$

Then $\mathbb{P}((39) \text{ holds}) \geq (1 - e^{-\frac{1}{12}\mathcal{N}^{\phi^2}})^{3\mathcal{N}} \geq 1 - 3\mathcal{N}e^{-\frac{1}{12}\mathcal{N}^{\phi^2}}$. Recalling (41) we replace 3 by 4 in the last estimate and obtain (37). \square

Remark 1: Theorem 4 states that if we take an arbitrary 1-dimensional subspace V then the intersection of a convex hull, say H_1 , with V is approximately a ball. More precisely, it contains a small ball and it is contained in a bigger ball:

$$B_{r^-} \subset V \cap H_1 \subset B_{r^+}, \quad (42)$$

where $r^\pm \sim \sqrt{\frac{\ln \mathcal{N}}{\mathcal{N}}}$. This statement is probabilistic, meaning that the arbitrary subspace V could be chosen with some probability that tends to 1 as $\mathcal{N} \rightarrow \infty$, as it follows from (37) and (36). While we are mostly interested in the result for large \mathcal{N} , our results are not asymptotic - one can deduce an explicit probabilistic estimate of success for any \mathcal{N} . In the main Theorems we choose to demonstrate the behavior of this probability of success for large \mathcal{N} . Namely, we show that convergence to 1 of this probability happens faster than any inverse polynomial rate: for any $\kappa > 0$ there exists \mathcal{N}_0 so that the probability of success is more than $1 - 1/\mathcal{N}^\kappa$ for all $\mathcal{N} \geq \mathcal{N}_0$. The main role of \mathcal{N}_0 in all our Theorems is to absorb all constants so that simple expressions like $1 - 1/\mathcal{N}^\kappa$ arise.

Remark 2: We use inclusion (42) in dimension $k = 2$ only. In other words, we apply it to families of two-dimensional subspaces V^i . Since the number of these subspaces is small, we still can claim $1 - 1/\mathcal{N}^\kappa$ probability of success using the union bound. The upper bound of inclusion (42) is a variant the Dvoretzky-Milman Theorem. It could be proved in a more general setup as Theorem 11.3.3 in the form of exercise 11.3.5 in [33]. The lower bound in (42) is sharper, because the general Dvoretzky-Milman bound may degenerate to 0. The lower bound relies on the construction of our noise collector \mathcal{C} .

B. Proof of Theorem 2

Proof: To prove the first claim, we repeat the proof of Theorem 2 from [28]. Suppose the $(M+1)$ -dimensional space V is spanned by e and the column vectors t_j , with j in the support of χ . For notational convenience and without loss of generality assume the support of χ is $1, 2, \dots, M$. Define several auxiliary convex hulls:

$$H_2 = \left\{ x \in \mathbb{R}^{\mathcal{N}} \left| x = \sum_{i=1}^M \xi_i t_i, \sum_{i=1}^M |\xi_i| \leq 1 \right. \right\}$$

and

$$\tilde{H}_\varepsilon = \left\{ x \in \mathbb{R}^{\mathcal{N}} \left| x = (1 + \varepsilon) \sum_{i=M+1}^K \xi_i t_i, \sum_{i=M+1}^K |\xi_i| \leq 1 \right. \right\}.$$

Further define convex hulls

$$H^\tau = \left\{ \frac{\xi}{\tau} h_1 + (1 - \xi) h_2, 0 \leq \xi \leq 1, h_i \in H_i, i = 1, 2 \right\}$$

and

$$H_\varepsilon^\tau = \left\{ \xi h + (1 - \xi) \tilde{h}, 0 \leq \xi \leq 1, h \in H^\tau, \tilde{h} \in \tilde{H}_\varepsilon \right\}.$$

The key observation is that $\text{supp}(\chi_\tau) \subset \text{supp}(\chi)$ if there exists a (sufficiently small) positive ε such that H^τ and H_ε^τ coincide in the subspace V . In other words, if there is $\varepsilon > 0$ such that

$$V \cap H_\varepsilon^\tau = V \cap H^\tau, \quad (43)$$

then $\text{supp}(\chi_\tau) \subset \text{supp}(\chi)$.

In turn, a sufficient condition for (43) could be formulated in terms of orthogonal projections of t_i on V . Denote by t_i^v the orthogonal projections of t_i on V for $i \geq M+1$. If $t_i^v \subset H^\tau$ strictly (i.e. $t_i^v \cap \partial H^\tau = \emptyset$) for all $i \geq M+1$, then (43) must hold for some $\varepsilon > 0$.

Therefore it remains to show $t_j^v \subset H^\tau$ strictly for all $i \geq M+1$.

Fix $j \geq M+1$, and suppose

$$t_j^v = \xi_0 t_0 + \sum_{n=1}^M \xi_n t_n, \text{ where } t_0 = \frac{e}{\|e\|}. \quad (44)$$

Suppose k is the index of the entry of ξ with the largest absolute value, i.e. $|\xi_k| = \max_{n \leq M} |\xi_n|$. Multiply (44) by $|\xi_k| t_k / \xi_k$, take absolute values, and obtain

$$|\langle t_j^v, t_k \rangle| \geq |\xi_k| - \sum_{n \neq k, n \leq M} |\xi_n \langle t_n, t_k \rangle|.$$

Using (22) for $i \leq M$ and (23) we obtain

$$c_0 \frac{\sqrt{\ln \mathcal{N}}}{\sqrt{\mathcal{N}}} \geq |\xi_k| \left(1 - M c_0 \frac{\sqrt{\ln \mathcal{N}}}{\sqrt{\mathcal{N}}} \right).$$

Choose α in (24) so that

$$M c_0 \frac{\sqrt{\ln \mathcal{N}}}{\sqrt{\mathcal{N}}} \leq \frac{1}{4}. \quad (45)$$

Then,

$$\left(1 - M c_0 \frac{\sqrt{\ln \mathcal{N}}}{\sqrt{\mathcal{N}}} \right) \geq \frac{3}{4},$$

and therefore,

$$|\xi_k| \leq \frac{4c_0}{3} \frac{\sqrt{\ln \mathcal{N}}}{\sqrt{\mathcal{N}}}$$

for all $k = 0, 1, 2, \dots, M$. Hence, $\sum_{k=1}^M |\xi_k| \leq 1/3$.

By (37) we can find $\tau_0 = 4c_0/\phi = O(1)$ so that

$$\frac{\tilde{t}_0}{\tau_0} \in H_1 \text{ for } \tilde{t}_0 = 4c_0 \frac{\sqrt{\ln \mathcal{N}}}{\sqrt{\mathcal{N}}} t_0$$

with probability $1 - 4\mathcal{N} \exp(-\mathcal{N}^{\phi^2}/12)$. Therefore,

$$t_j^v = \tilde{\xi}_0 \tilde{t}_0 + \sum_{k=1}^M \xi_k t_k, \text{ and } |\tilde{\xi}_0| + \sum_{k=1}^M |\xi_k| \leq 1/3 + 1/3 \leq 2/3$$

and $\tilde{t}_0/\tau \in H_1$ for all $\tau \geq \tau_0$. Therefore, $t_j^v \subset H^\tau$ strictly with probability $1 - 4\mathcal{N} \exp(-\mathcal{N}^{\phi^2}/12)$. For a given $\kappa > 0$ we now can choose \mathcal{N}_0 large enough so that $4\mathcal{N} \exp(-\mathcal{N}^{\phi^2}/12) \leq 1/\mathcal{N}^\kappa$ for all $\mathcal{N} \geq \mathcal{N}_0$, and this completes the proof of the first part.

To prove the second claim, we repeat the proof of Theorem 3 from [28]. Let us reduce the general case to the case when \mathcal{A} is the first M columns of the identity matrix:

$\mathcal{A} = [e_1, e_2, \dots, e_M]$. By the first claim the last $K - M$ columns of \mathcal{A} are not used in approximating \mathbf{d} . Therefore we can assume for simplicity and without loss of generality that \mathcal{A} has only the first M columns. Construct a square matrix \mathcal{M} so that its first M columns are \mathbf{t}_i $i \leq M$, and the remaining $N - M$ columns form an orthonormal basis of V^\perp - the orthogonal complement of V . Gershgorin circle theorem applied to $\mathcal{M}^* \mathcal{M}$ with estimate (45) implies that $\sqrt{3}/2$ is a lower bound for singular values of \mathcal{M} . Suppose we use $\tilde{\mathcal{A}} = \mathcal{M}^{-1} \mathcal{A}$ and $\tilde{\mathcal{C}} = \mathcal{M}^{-1} \mathcal{C}$ instead of \mathcal{A} and \mathcal{C} , respectively. Then all assumptions of the Theorem will remain essentially the same, up to this universal constant $\sqrt{3}/2$. For example, the lengths of \mathbf{c}_j may become at most $\sqrt{3}/2$, or the constant c_1 in (26) may be increased to $2c_1/\sqrt{3}$.

Thus, we can assume for simplicity and without loss of generality that \mathcal{A} is the first M columns of the identity matrix.

Suppose V^i are the 2-dimensional spaces spanned by \mathbf{e} and \mathbf{t}_i for $i \in \text{supp}(\chi)$. We will denote by λH a hull H rescaled by λ . By (42) all $\lambda H^\tau \cap V^i$ look like rounded rhombi depicted on Fig. 8, and $\lambda H_1^\tau \cap V^i \subset B_{\lambda\tau}^i$ with probability $1 - N^{-\kappa}$, where $B_{\lambda\tau}^i$ is a 2-dimensional ℓ_2 -ball of radius $\lambda\tau c_0 \sqrt{\ln N} / \sqrt{N}$. Thus $\lambda H^\tau \cap V^i \subset H_{\lambda\tau}^i$ with probability $1 - N^{-\kappa}$, where $H_{\lambda\tau}^i$ is the convex hull of $B_{\lambda\tau}^i$ and a vector $\lambda \mathbf{f}_i$, $\mathbf{f}_i = \chi_i \|\chi\|_{\ell_1}^{-1} \mathbf{t}_i$. Then $\text{supp}(\chi_\tau) = \text{supp}(\chi)$, if there exists λ_0 so that $\chi_i \mathbf{t}_i + \mathbf{e}$ lies on the flat boundary of $H_{\lambda_0}^i$ for all $i \in \text{supp}(\chi)$.

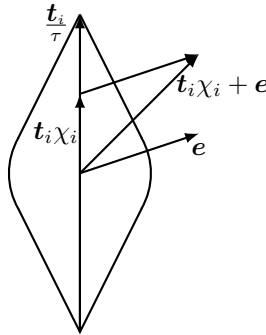


Fig. 8. An intersection of H^τ with the $\text{span}(\mathbf{t}_i, \mathbf{e})$ is a rounded rhombus.

If $\min_{i \in \text{supp}(\chi)} |\chi_i| \geq \gamma \|\chi\|_\infty$, then there exists a constant $c_2 = c_2(\gamma)$ such that if $\chi_i \mathbf{t}_i + \mathbf{e}$ lies on the flat boundary of H_{λ}^i for some i and some λ , then there exists λ_0 so that $\chi_i \mathbf{t}_i + c_2 \mathbf{e}$ lies on the flat boundary of $H_{\lambda_0}^i$ for all $i \in \text{supp}(\chi)$. If

$$\frac{|\langle \mathbf{t}_i, \chi_i \mathbf{t}_i + \mathbf{e} \rangle|}{\|\chi_i \mathbf{t}_i + \mathbf{e}\|_{\ell_2}} \geq \frac{\tau c_0 \sqrt{\ln N}}{\sqrt{N} \|\mathbf{f}_i\|_{\ell_2}} = \frac{\tau c_0 \|\chi\|_{\ell_1} \sqrt{\ln N}}{\sqrt{N} |\chi_i|}, \quad (46)$$

then $\chi_i \mathbf{t}_i + c_2 \mathbf{e}$ lies on the flat boundary of H_{λ}^i .

Since $|\langle \mathbf{t}_i, \chi_i \mathbf{t}_i + \mathbf{e} \rangle| \geq |\chi_i|/2$ by (25), inequality (46) holds if

$$\frac{|\chi_i|}{\|\chi_i \mathbf{t}_i + \mathbf{e}\|_{\ell_2}} \geq \frac{2\tau c_0 \|\chi\|_{\ell_1} \sqrt{\ln N}}{\sqrt{N} |\chi_i|}.$$

By (26) and using $\|\chi\|_{\ell_1} \leq M$ the last inequality is true if

$$M \leq \frac{\sqrt{N}}{\tau c_1 c_0 \sqrt{\ln N}}.$$

The last inequality is true if α in (24) is small enough. Thus, $\text{supp}(\chi_\tau) = \text{supp}(\chi)$. ■

In the proof of Theorem 3 we will use the following variant of Hanson-Wright inequality.

Theorem 5: (Hanson-Wright inequality for bounded symmetric random variables) Suppose X_i are independent symmetric random variables, with $\|X_i\|_{\ell_\infty} \leq K$. Let $\Xi = \sum_{i \neq j} X_i X_j m_{ij}$. Then

$$\mathbb{P}(|\Xi| > t) \leq 2 \exp\left(-\frac{t^2/32}{K^4 \|\mathbf{M}\|_F^2}\right). \quad (47)$$

where \mathbf{M} is a matrix with components m_{ij} , $\|\mathbf{M}\|_F$ is its Frobenius (Hilbert-Schmidt) norm.

Theorem 3 can be proved using standard arguments from high-dimensional probability. For example one can modify proof of Theorem 6.2.1 in [30] to bounded random variables to obtain 3.

C. Proof of Theorem 3

Proof: Theorem 3 immediately follows from Theorem 2 as soon as we check that conditions (22), (23), (25) and (26) are satisfied. Choose c_0 , τ_0 , \mathcal{N}_0 and α so that Theorem 2 is satisfied with probability $1 - \frac{1}{3N^\kappa}$. Note that we can increase c_0 , τ_0 , \mathcal{N}_0 and decrease α in this proof if necessary.

By rescaling we can and we do assume that $\|\rho\|_{\ell_\infty} = 1$ and therefore $\|\chi\|_{\ell_\infty} = 1$.

The idea is to prove that each (22), (23), (25) and (26) hold with probability close to 1. Then we use the union bound to conclude that all these conditions hold simultaneously with probability close to 1. We prove (23), (26), (22), and (25) in paragraphs a), b), c) and d) below, respectively.

a) We start with verifying (23). We denote by $(\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l}$ the column of $\bar{\mathcal{A}} \otimes \mathcal{A}$ that arises from a tensor product $\bar{\mathbf{a}}_k \otimes \mathbf{a}_l$. If we use all $\mathcal{N} = N^2$ of the data then

$$\begin{aligned} \langle (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l}, (\bar{\mathcal{A}} \otimes \mathcal{A})_{m,n} \rangle &= \sum_{i=1}^N \sum_{j=1}^N \bar{a}_{k,i} a_{l,j} a_{m,i} \bar{a}_{n,j} \\ &= \langle \mathbf{a}_m, \mathbf{a}_k \rangle \langle \mathbf{a}_l, \mathbf{a}_n \rangle. \end{aligned}$$

In particular, all columns of $\bar{\mathcal{A}} \otimes \mathcal{A}$ have length 1. Therefore,

$$|\langle \mathbf{t}_i, \mathbf{t}_j \rangle| = |(\bar{\mathcal{A}} \otimes \mathcal{A})_{i,i}, (\bar{\mathcal{A}} \otimes \mathcal{A})_{j,j}| = |\langle \mathbf{a}_i, \mathbf{a}_j \rangle|^2 \leq \frac{\Delta^2}{N},$$

and condition (23) is verified because $\Delta^2 < c_0 \sqrt{\ln N}$ for $\mathcal{N} > \mathcal{N}_0$ if we choose \mathcal{N}_0 large enough.

b) The next condition is (26). Note that

$$\lambda_1 M \leq \|e\|_{\ell_2} \leq \lambda_2 M \quad (48)$$

implies (26) because $\gamma M \leq \|\chi\|_{\ell_1} \leq M$. To prove (48) with high probability we write

$$\|e\|_{\ell_2}^2 = \|\chi\|_{\ell_1}^2 + 2\|\chi\|_{\ell_1} \Xi_1 + \Xi_2,$$

where we define random variables

$$\Xi_1 = \sum_{k,l,k \neq l} \bar{\rho}_k \rho_l \langle \mathbf{a}_k, \mathbf{a}_l \rangle, \quad (49)$$

and

$$\Xi_2 = \sum_{\text{all indices different}} \rho_k \bar{\rho}_l \bar{\rho}_m \rho_n \langle \mathbf{a}_m, \mathbf{a}_k \rangle \langle \mathbf{a}_l, \mathbf{a}_n \rangle. \quad (50)$$

We first estimate the random variable Ξ_1 from (49). By Hanson-Wright inequality (47)

$$\mathbb{P}(|\Xi_1| > t) \leq 2 \exp\left(-\frac{t^2/32}{\|\mathbf{M}\|_F^2}\right)$$

where \mathbf{M} is a matrix with components $|\rho_k \rho_l| \langle \mathbf{a}_k, \mathbf{a}_l \rangle$. Since $|\langle \mathbf{a}_k, \mathbf{a}_l \rangle| \leq \Delta/\sqrt{N}$, we obtain $\|\mathbf{M}\|_F \leq \Delta M/\sqrt{N}$. Take $t = \gamma M/8 \leq \|\chi\|_{\ell_1}/8$ and obtain

$$\mathbb{P}(|\Xi_1| > \gamma M/8) \leq 2 \exp(-cN), c = c(\gamma),$$

which is negligible for large N . Thus

$$|\Xi_1| \leq \frac{\|\chi\|_{\ell_1}}{8} \quad (51)$$

with probability $1 - 2 \exp(-cN)$.

We now estimate the random variable Ξ_2 from (50). Observe that $\Xi_2 = (\Xi_1)^2 - \Xi_3$, where the random variable

$$\Xi_3 = \sum_{m=l \text{ or } k=n \text{ or both}} \rho_k \bar{\rho}_l \bar{\rho}_m \rho_n \langle \mathbf{a}_m, \mathbf{a}_k \rangle \langle \mathbf{a}_l, \mathbf{a}_n \rangle.$$

For Ξ_3 we can use a deterministic estimate:

$$|\Xi_3| \leq 2 \frac{c_0^2 M^3}{N} \leq c\alpha \frac{\|\chi\|_{\ell_1}^2}{\sqrt{\ln N}} \leq \frac{\|\chi\|_{\ell_1}^2}{16}.$$

For the random variable $(\Xi_1)^2$ we use (51).

We now combine estimates for Ξ_1 and Ξ_2 to obtain (48). Using the union bound, we obtain

$$\frac{1}{2} \|\chi\|_{\ell_1}^2 \leq \|e\|_{\ell_2}^2 \leq \frac{3}{2} \|\chi\|_{\ell_1}^2 \quad (52)$$

with probability $1 - 2 \exp(-cN)$. Thus, (48) holds with probability $1 - 2 \exp(-cN)$.

c) We will now prove (22). For $m \notin \text{supp}(\chi)$, consider a random variable

$$\begin{aligned} \Theta_m &= \langle \mathbf{t}_m, \mathbf{e} \rangle \\ &= \sum_{k,l,k \neq l} \bar{\rho}_k \rho_l \langle \mathbf{t}_m, (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l} \rangle \\ &= \sum_{k,l,k \neq l} \bar{\rho}_k \rho_l \langle \mathbf{a}_m, \mathbf{a}_k \rangle \langle \mathbf{a}_m, \mathbf{a}_l \rangle. \end{aligned} \quad (53)$$

We have

$$|\langle \mathbf{t}_m, (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l} \rangle| = |\langle \mathbf{a}_m, \mathbf{a}_k \rangle \langle \mathbf{a}_m, \mathbf{a}_l \rangle| \leq \frac{\Delta^2}{N} \quad (54)$$

if $m \neq k$, and $m \neq l$. If \mathbf{M} is a matrix with components $|\rho_k \rho_l| \langle \mathbf{a}_m, \mathbf{a}_k \rangle \langle \mathbf{a}_m, \mathbf{a}_l \rangle$, then $\|\mathbf{M}\|_F \leq \Delta^2 M/N$. Using (48) choose $t = c_0 \frac{\sqrt{\ln N^2}}{N} \|e\|_{\ell_2} > c_0 \frac{\gamma}{2} \frac{M \sqrt{\ln N}}{N}$ in Hanson-Wright inequality (47) to obtain:

$$\begin{aligned} \mathbb{P}\left(|\Theta_m| > c_0 \frac{\sqrt{\ln N^2}}{N} \|e\|_{\ell_2}\right) &\leq \mathbb{P}\left(|\Theta_m| > c_0 \frac{\gamma}{2} \frac{M \sqrt{\ln N}}{N}\right) \\ &\leq 2 \exp\left(-\frac{\gamma^2 c_0^2 \ln N}{128 \Delta^4}\right). \end{aligned}$$

Then (22) holds with probability $1 - \frac{1}{3N^\kappa}$ if c_0 is large enough.

d) We will now prove (25). For $m \in \text{supp}(\chi)$ decompose

$$\Theta_m = \langle \mathbf{t}_m, \mathbf{e} \rangle = \Theta_m^1 + \Theta_m^2$$

where the random variables

$$\Theta_m^1 = \sum_{k,l,k \neq l, k \neq m, l \neq m} \bar{\rho}_k \rho_l \langle \mathbf{a}_m, \mathbf{a}_k \rangle \langle \mathbf{a}_m, \mathbf{a}_l \rangle,$$

and

$$\Theta_m^2 = \sum_{k,k \neq m} (\bar{\rho}_m \rho_k + \bar{\rho}_k \rho_m) \langle \mathbf{a}_m, \mathbf{a}_k \rangle. \quad (55)$$

The distribution of the random variable Θ_m^1 has exactly the same behavior as Θ_m for $m \notin \text{supp}(\chi)$. We therefore have

$$\begin{aligned} \mathbb{P}(|\Theta_m^1| > \frac{\gamma}{4} \|\chi\|_{\ell_\infty}) &\leq 2 \exp\left(-c \frac{N^2}{\Delta^4 M^2}\right) \\ &\leq 2 \exp(-\bar{c} \ln N / \alpha^2) \\ &\leq \frac{1}{6 N^\kappa}, \end{aligned}$$

by Hanson-Wright inequality (47) if α is small enough. If $m = l$ (or $m = k$) then

$$|\langle \mathbf{t}_m, (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,m} \rangle| = |\langle \mathbf{a}_m, \mathbf{a}_k \rangle| \leq \frac{\Delta}{\sqrt{N}}. \quad (56)$$

We will now estimate Θ_m^2 given by (55). If we condition on ρ_m , then Θ_m^2 is a sum of independent random variables. Therefore by Hoeffding's inequality

$$\mathbb{P}(|\Theta_m^2| > t) \leq 2 \exp\left(-c \frac{t^2}{b^2}\right), \text{ where } b^2 \leq \frac{c_0^2 M}{N} \leq \frac{\Delta^2 \alpha}{\ln N}.$$

Choosing t appropriately we obtain

$$\mathbb{P}\left(|\Theta_m^2| > \frac{\gamma}{4} \|\chi\|_{\ell_\infty}\right) \leq \frac{1}{6 N^\kappa}.$$

by choosing α small enough.

Using the union bound we combine estimates for Θ_m^1 and Θ_m^2 to conclude that

$$\mathbb{P}\left(|\Theta_m| > \frac{\gamma}{2} \|\chi\|_{\ell_\infty}\right) \leq \frac{1}{3 N^\kappa}$$

for $m \in \text{supp}(\chi)$.

All the conditions (23) (26), (22), (25) are now verified. Applying the union bound we conclude that estimates in Theorem 2 hold with probability $1 - \frac{1}{N^\kappa}$. This completes the proof of Theorem 3. ■

Remark 3: The proof of Theorem 3 reveals why we had to assume (25) for $m \in \text{supp}(\chi)$. When $m \notin \text{supp}(\chi)$ then $\langle \mathbf{t}_m, \mathbf{e} \rangle$ is estimated in (53) using (54). When $m \in \text{supp}(\chi)$ then $\langle \mathbf{t}_m, \mathbf{e} \rangle$ contains Θ_m^2 given by (55). For Θ_m^2 we cannot use (54), and we have to use a weaker estimate (56).

D. A deterministic version of Theorem 3

Theorem 6: Suppose X is a solution of (15), $\chi = \text{diag}(X)$ is M -sparse, $\mathbf{d} \in \mathbb{C}^N$, $N = N^2$, and $T = (\bar{\mathcal{A}} \otimes \mathcal{A})\chi : \mathbb{C}^K \rightarrow \mathbb{C}^N$. Fix $\beta > 1$, and draw $\Sigma = \mathcal{N}^\beta$ columns for \mathcal{C} , independently, from the uniform distribution on \mathbb{S}^{N-1} and define γ as in (18) and Δ as in (30). Then, for any $\kappa > 0$, there are constants $\alpha = \alpha(\kappa, \gamma, \Delta)$, $\tau = \tau(\kappa, \beta)$, and $N_0 = N_0(\kappa, \beta, \gamma, \Delta)$ such that the following holds. If $M \leq \alpha \sqrt{N}$ and χ_τ is the solution (17), then $\text{supp}(\chi) = \text{supp}(\chi_\tau)$ for all $N > N_0$ with probability $1 - 1/N^\kappa$.

Proof: We need to verify that all conditions of Theorem 2 are satisfied non-probabilistically. Conditions (23) is already

verified in the proof of Theorem 3 under even weaker assumptions than in Theorem 6. Therefore we only need to verify estimates (26), (25) and (22).

Since

$$e = \sum_{k \neq l} \bar{\rho}_k \rho_l (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l},$$

we have

$$\begin{aligned} \|e\|_{\ell_2}^2 &\leq \frac{2\Delta}{\sqrt{N}} \sum_{\text{all indices}} \chi_k |\rho_{m_1}| |\rho_{m_2}| \\ &+ \frac{\Delta^2}{N} \sum_{\text{all indices}} |\rho_{m_1}| |\rho_{m_2}| |\rho_{k_1}| |\rho_{k_2}| \\ &+ \sum_{k,m} \chi_k \chi_m \end{aligned}$$

and thus

$$\begin{aligned} \|e\|_{\ell_2}^2 &\leq 2\Delta\alpha \|\chi\|_{\ell_1} \|\rho\|_{\ell_2}^2 + \Delta^2 \alpha^2 \|\rho\|_{\ell_2}^4 + \|\chi\|_{\ell_1}^2 \\ &= (1 + \Delta\alpha)^2 \|\chi\|_{\ell_1}^2. \end{aligned}$$

Therefore estimate (26) holds.

A non-probabilistic version of estimate (25) is as follows.

For $m \in \text{supp}(\chi)$ we have

$$\begin{aligned} |\langle t_m, e \rangle| &= \left| \sum_{k,l,k \neq l} \bar{\rho}_k \rho_l \langle t_m, (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l} \rangle \right| \\ &\leq 2 \sum_{k,k \neq m} |\rho_k| |\rho_m| |\langle t_m, (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,m} \rangle| \\ &+ \sum_{k,l,k \neq l \neq m} |\rho_k| |\rho_l| |\langle t_m, (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l} \rangle| \\ &\leq \frac{2\Delta}{\sqrt{N}} \sum_k |\rho_k| |\rho_m| + \frac{\Delta^2}{N} \sum_{k,l} |\rho_k| |\rho_l| \\ &\leq \left(\frac{2\Delta M}{\sqrt{N}} + \frac{\Delta^2 M^2}{N} \right) \|\rho\|_{\ell_\infty}^2 \\ &= \left(\frac{2\Delta M}{\sqrt{N}} + \frac{\Delta^2 M^2}{N} \right) \|\chi\|_{\ell_\infty} \\ &\leq \frac{\gamma}{2} \|\chi\|_{\ell_\infty}. \end{aligned}$$

if α is small enough.

We now obtain a lower bound on $\|e\|_{\ell_2}$. For Ξ_1 and Ξ_2 in (49) and (50), respectively, we have

$$|\Xi_1| \leq \frac{\Delta M^2}{\sqrt{N}} \leq \Delta\alpha M, \quad |\Xi_2| \leq \frac{\Delta^2 M^4}{N} \leq \Delta^2 \alpha^2 M^2.$$

Since

$$\|e\|_{\ell_2}^2 = \|\chi\|_{\ell_1}^2 + 2\|\chi\|_{\ell_1} \Xi_1 + \Xi_2, \text{ and } \|\chi\|_{\ell_1} = M$$

we can choose α so that

$$M/2 = \|\chi\|_{\ell_1}/2 \leq \|e\|_{\ell_2}^2.$$

To show (22) observe that

$$|\langle t_m, (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l} \rangle| = |\langle a_m, a_k \rangle \langle a_m, a_l \rangle| \leq \Delta^2/N,$$

because $m \neq k$, and $m \neq l$. Therefore

$$\begin{aligned} |\langle t_m, e \rangle| &= \left| \sum_{k,l,k \neq l} \bar{\rho}_k \rho_l \langle t_m, (\bar{\mathcal{A}} \otimes \mathcal{A})_{k,l} \rangle \right| \\ &\leq \frac{\Delta^2 M^2}{N} \leq \frac{\Delta^2 \alpha \|e\|_{\ell_2}}{\sqrt{N}}, \end{aligned}$$

and (22) follows either for choosing α small or $\ln N$ large. ■

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