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Highlights:

- Importance of inspections in binary systems based on the Value of Information;
- Two metrics for component-level and for system-level maintenance actions;
- Inspection priorities for series and parallel systems;
- A heuristic proposed for reducing the computation complexity in general systems;

Optimal Inspection of Binary Systems via Value of Information Analysis

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Abstract

We develop computable metrics to assign priorities for information collection on binary systems composed of binary components. Components are worth inspecting because their condition states are uncertain, and system functioning depends on them. The Value of Information (VoI) enables assessment of the impact of information in decision making under uncertainty, including the component's reliability and role in the system, the precision of the observation, the available maintenance actions and the expected economic loss. We introduce the VoI-based metrics for system-level ("global") and component-level ("local") maintenance actions, analyze the properties of these metrics, and apply them to series and parallel systems. We discuss their computational complexity in applications to general network systems and, to tame the complexity for the local metric assessment, we present a heuristic and assess its performance on some case studies.

Keywords:

Binary Networks, Importance Measure, Inspections, Value of Information

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1 1. Introduction

2 Many civil infrastructures (e.g. transportation and gas pipeline networks)
3 consist of multiple binary components, arranged in a system to fulfill various
4 functions [21] [23] [28]. The binary states of the components, either intact or
5 damaged, determine the system condition. The belief of the agent controlling
6 the maintenance process can be described by a probabilistic distribution on
7 the possible states of the components. Maintenance actions are selected to
8 trade the risk of system malfunctioning for the cost of maintenance (includ-
9 ing repair and retrofitting actions). Observations of the components' states
10 can improve decision making and reduce the uncertainty and the mainte-
11 nance cost. However, because of budget constraints, it is often impossible
12 to inspect all components in a system. Therefore it is important to assign
13 inspection priorities for the components. Intuitively, many factors can af-
14 fect the inspection preferences, such as the probabilities of failure events,
15 the maintenance costs and the role of each component in the system. These
16 factors can be integrated in an Importance Measure (IM) for inspections, i.e.
17 a value assigned to each component to summarize the benefit of observing
18 the state of a component.

19 To introduce the problem, consider a binary system composed of N com-
20 ponents: $\{c_1, c_2, \dots, c_N\}$. Let $s = [s_1, s_2, \dots, s_N] \in S$ denote the states of the
21 components, with $s_j = 1$ indicating that component c_j is working, and $s_i = 0$
22 that it fails, where $S = \mathbb{B}^N$ and $\mathbb{B} = \{0, 1\}$. The system state $u = \phi(s)$ is
23 also a binary variable, where $\phi : S \rightarrow \mathbb{B}$ is the component-to-system function.
24 State s is unknown to the agent who manages the system. Instead, the agent

25 optimizes the measurement and maintenance plans on the basis of her belief
 26 of s . The prior probability distribution of s is denoted as $p_s : S \rightarrow [0, 1]$,
 27 and $p_i = \mathbb{P}[s_i = 0]$ indicates the prior marginal failure probability of c_i . The
 28 failure probability of the system is $p_u = \mathbb{P}[u = 0]$, and we use p_π and $p_{\omega|E}$ for
 29 the prior value of p_u and its posterior value given event E , respectively.

30 In this paper, we develop metrics to assess the importance of inspecting
 31 any component. We assume that the outcome of the inspection is also binary.
 32 If component c_i is inspected, $y_i = 0$ indicates an “alarm”, i.e. a symptom
 33 that c_i is not working, whereas $y_i = 1$ indicates that c_i seems to work, and we
 34 define this outcome as a “silence”. If the inspection is perfect, then $y_i = s_i$.

35 On the basis of the measurement outcome, we can update the prior dis-
 36 tribution of random variables s to posterior distribution $p_{s|y_i}$ and, the system
 37 level failure probability to $p_{\omega|y_i}$. When the components are interdependent,
 38 the measurement of one component may also affect the failure probability of
 39 other components.

40 Birnbaum was first to introduce Importance Measures (IMs) [4] to eval-
 41 uate the contribution of each component to a system’s performance, such
 42 as the system connectivity. Birnbaum’s Measure (BM) [4] evaluates the im-
 43 portance of a component by the difference in the posterior system failure
 44 probability when it is damaged or intact (i.e., in our framework, when the
 45 inspection outcome is alarm or silence):

$$\text{BM}(i) = p_{\omega|y_i=0} - p_{\omega|y_i=1} \quad (1)$$

46 Other IMs are discussed in Appendix B. Most of them focus on the marginal
 47 or conditional probability of the failure events, and they do not explicitly
 48 include any evaluation of the maintenance cost and risk. In maintenance

49 problems, a component need high attention because of its topological func-
50 tion in the system and because of its high probabilities of failure. To assess
51 this need of attention, Wu and Coolen [29] extended the BM to a cost-based
52 IM. Zio and Podofillini [30] presented an approach for optimizing multiple
53 objectives (such as system risk and maintenance costs), and they developed
54 generic algorithms to reduce the computation time. Der Kiureghian et al.
55 [6] modeled the component failures as independent Poisson events and devel-
56 oped IMs for long-term maintenance of series, parallel and general systems
57 based on the system unavailability, mean rate of failure and mean duration
58 of downtime.

59 To compare and rank the impact of inspections, one can assess their
60 Value of Information (VoI). VoI assessment is based on Bayesian pre-posterior
61 analysis, as introduced by [10], who integrated the probabilistic knowledge
62 about the system with the economic factors related to the available actions.
63 In the maintenance process of infrastructure systems, the economic costs are
64 related to the system malfunctioning, the execution of inspections, and repair
65 or replacement actions.

66 VoI has been studied intensively in the area of Structural Health Moni-
67 toring (SHM). Straub and Faber [26] integrated VoI for risk-based inspection
68 scheduling and maintenance planning of structural systems. Pozzi and Der
69 Kiureghian [18] provided a framework for assessing VoI for the long-term
70 SHM, and proposed a Monte Carlo approach to reduce the computation com-
71 plexity. They also investigated how the imperfect measurements affected the
72 posterior decisions. Straub et al. [25] illustrated how to model the stochastic
73 dependencies of component deterioration, the failure consequences and the

74 inspection cost. The VoI has also been applied to long-term decision making
75 problems. Miller [17] extended VoI analysis to optimize not only static one-
76 shot inspection, but also to optimize sequentially dependent observations.
77 [Srinivasan and Parlikad \[24\]](#), [Memarzadeh and Pozzi \[16\]](#) and [Andriotis et al.](#)
78 [\[1\]](#) applied the component-wise VoI metric to sequential decision making in
79 the management of infrastructure systems, modeled by Partially Observable
80 Markov Decision Process (POMDP). [Thöns \[27\]](#) used decision trees to assess
81 long-term VoI. [Bensi et al. \[3\]](#) developed Bayesian Networks and Influence
82 Diagrams to evaluate post-event inspections, and they proposed VoI-based
83 heuristic for optimal inspection sequences. [Sensitivity analysis of the process](#)
84 [parameters with respect to the optimal maintenance actions was presented by](#)
85 [\[31\] and \[5\]](#). The complexity of computing VoI can grow exponentially with
86 the number of components in a system [14]. Even worse, the VoI generally
87 lacks the property of submodularity [15], so that the application of greedy ap-
88 proaches does not provide certain guarantees of near-optimal solutions [20].
89 Effective strategies have been proposed for efficient VoI computation in some
90 special cases [12].

91 In this paper, we investigate VoI-based metrics related to system-level
92 (“global”) and component-level (“local”) decision making after component
93 inspections, for systems with various topologies, and compare these results
94 with traditional IMs. A recent paper [8] also focuses on inspections for net-
95 worked systems, developing an approach to identify the components most in
96 need of inspection, similar to what we define as the local metric. We also
97 derive simple optimal rules for series and parallel systems. For general sys-
98 tems, we discuss the computational complexity of the problem and provide

99 a heuristic approach. In Section 2, we introduce the global and local metrics
 100 for evaluating the components' VoI. Section 3 describes rules for optimiz-
 101 ing these metrics to typical systems such as series and parallel systems. In
 102 Section 4, we propose approximated approaches to simplify the optimization
 103 complexity, and in Section 5 we examine different applications of global, local
 104 and heuristic approaches to some system examples.

105 2. Global and local VoI metrics

106 2.1. Principles of VoI

107 Fig 1a illustrates the decision graph for the process of inspecting and
 108 maintaining the system. Continuous arrows from one set of nodes to one node
 109 indicate that the probability of latter variable is defined conditional to the
 110 former ones. Double arrows indicate deterministic relations. Dashed arrows
 111 from random variables to decision variables indicate that the former ones are
 112 observed before the latter is selected. Let \mathcal{A} denote the set of all possible
 113 maintenance plans, that we simply call "actions". Action $A \in \mathcal{A}$ transforms
 114 current components' state $s \in S$ into state $s' \in S$, via transition distribution
 115 $p_{s'|s,A} : S \times \mathcal{A} \times S \rightarrow [0, 1]$. Loss function $\mathcal{L}(s', A) = \mathcal{L}_I(\phi(s')) + \mathcal{L}_{II}(A) :$
 116 $S \times \mathcal{A} \rightarrow \mathbb{R}$ summarizes the overall cost: $\mathcal{L}_I(\phi(s')) = C_F(1 - u')$ adds failure
 117 costs C_F if the system is not functioning, depending on system state u' after
 118 taking action A , which is associated with implementing cost $\mathcal{L}_{II}(A)$.

119 The prior loss L_π is the minimum expected cost among all possible ac-
 120 tions, before any inspection:

$$L_\pi = \min_A \mathbb{E}_s \mathbb{E}_{s'|s,A} \mathcal{L}(s', A) = \min_A \mathbb{E}_{s'|A} \mathcal{L}(s', A) \quad (2)$$

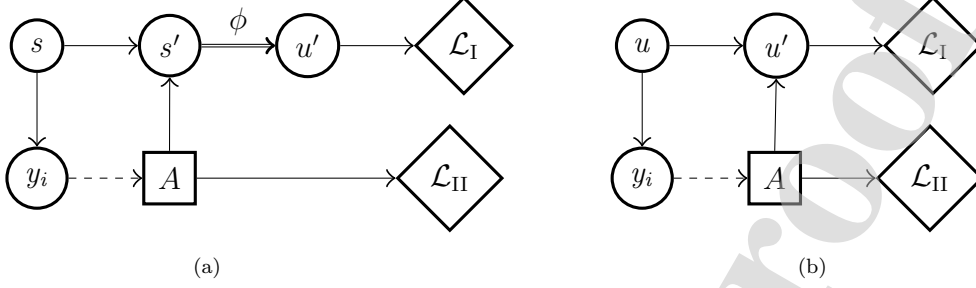


Figure 1: Decision graph for the general problem (a), and for the global metric (b).

121 where $\mathbb{E}_{s'|A}[\cdot] = \mathbb{E}_s \mathbb{E}_{s'|s,A}[\cdot]$ denotes the statistical expectation depending on
 122 distributions $p_{s'|s,A}$ and p_s .

123 Inspecting component c_i , the agent collects observation y_i distributed
 124 according to function $p_{y_i} : \mathbb{B} \rightarrow [0, 1]$, and the belief of the components' state
 125 s is updated to posterior distribution $p_{s|y_i} : S \times \mathbb{B} \rightarrow [0, 1]$. These functions
 126 are obtained by Bayes' rule:

$$p_{y_i} = \sum_s p_{y_i|s} p_s \quad p_{s|y_i} = \frac{p_{y_i|s} p_s}{p_{y_i}} \quad (3)$$

127 where $p_{y_i|s} : \mathbb{B} \times S \rightarrow [0, 1]$ is the likelihood function related to observation
 128 y_i .

129 The corresponding expected posterior loss is:

$$L_\omega(i) = \mathbb{E}_{y_i} \min_A \mathbb{E}_{s'|y_i,A} \mathcal{L}(s', A) \quad (4)$$

130 where $\mathbb{E}_{s'|y_i,A}[\cdot] = \mathbb{E}_s \mathbb{E}_{s'|s,A}[\cdot]$ is the posterior expectation, related to distri-
 131 bution $p_{s|y_i}$, and $\mathbb{E}_{y_i}[\cdot]$ is related to distribution p_{y_i} .

132 The VoI for inspecting c_i is the expected loss reduction due to the in-
 133 spection, i.e. the difference between the prior and posterior loss functions
 134 [10]:

$$\text{VoI}(i) = L_\pi - L_\omega(i) \quad (5)$$

135 Loss function \mathcal{L} does not include the cost of monitoring, and the VoI is
 136 always not negative. However, if such cost is uniform among components, the
 137 VoI is a rational IM that assesses the relevance of inspections. The optimal
 138 component to inspect, c_{i^*} , is the argument that maximizes Eq.(5):

$$i^* = \arg \max_i \text{VoI}(i) \quad (6)$$

139 The VoI depends on the specific number N of components, the action
 140 domain \mathcal{A} , the loss function \mathcal{L} (in turn defined by the component-to-system
 141 function ϕ , the failure cost C_F , and the implementing cost \mathcal{L}_{II}), the prior
 142 probability p_s , the transition probability $p_{s'|s,A}$ and the likelihood function
 143 $p_{y_i|s}$ adopted, as apparent in Fig 1a. In the following Sections, we describe
 144 a form of the likelihood function for binary components, and then we focus
 145 on two classes of losses and transitions, related to global and local decision
 146 making.

147 2.2. Modeling imperfect inspections

148 The VoI analysis also depends on the specific assumed likelihood function.
 149 If the binary outcome y_i , of inspecting component c_i , depends only on the
 150 state s_i of that component, likelihood function $p_{y_i|s}$ in Eq.(3) is reduced to a
 151 4-entry emission table $p_{y_i|s_i} : \mathbb{B} \times \mathbb{B} \rightarrow [0, 1]$, shown in Table 1.

152 Observations of components' states are prone to error, and the inaccuracy
 153 can be formulated by two parameters $\epsilon_{\text{FA}} = \mathbb{P}[y_j = 0|s_j = 1]$ and $\epsilon_{\text{FS}} =$
 154 $\mathbb{P}[y_i = 1|s_i = 0]$, which are the probability of type I error: having an "alarm"
 155 when the component is undamaged, and of type II error, a silence when
 156 the component is damaged. Although these probabilities can depend on

Actual state	Observation	
	Silence $y_i = 1$	Alarm $y_i = 0$
Undamaged $s_i = 1$	$1 - \epsilon_{\text{FA}}$	ϵ_{FA}
Damaged $s_i = 0$	ϵ_{FS}	$1 - \epsilon_{\text{FS}}$

Table 1: Emission probability table for observation y_i given state s_i .

157 the specific component, in the following discussion, we assume that all the
 158 components have identical ϵ_{FS} and ϵ_{FA} .

159 Inspection outcomes probability function $p_{y_i} : \mathbb{B} \rightarrow [0, 1]$, is related to a
 160 single value: the probability $h_i = \mathbb{P}[y_i = 0]$ of receiving an alarm on c_i , which
 161 is:

$$h_i = (1 - \epsilon_{\text{FS}})p_i + \epsilon_{\text{FA}}(1 - p_i) = \epsilon_{\text{FA}} + Kp_i \quad (7)$$

162 where constant $K = 1 - \epsilon_{\text{FA}} - \epsilon_{\text{FS}}$ is strictly positive, because we assume that
 163 both ϵ_{FA} and ϵ_{FS} are less than $1/2$.

164 2.3. Global metric

165 We define the global metric assuming that action A affects the system
 166 state u . In this setting, for any of the two values of the binary variable u , an
 167 expected loss value can be assigned to any action A , regardless of the details
 168 of components' conditions (e.g., the damage location) described by variable
 169 s . Fig 1b shows the corresponding decision graph, in which the loss is a
 170 function of system state u' after the taken action: $l(u', A) = \mathcal{L}(s', A)$, with
 171 $u' = \phi(s')$. Transition function $p_{s'|s,A}$ is now converted into function $p_{u'|u,A} :$
 172 $\mathbb{B} \times \mathcal{A} \times \mathbb{B} \rightarrow [0, 1]$, in turn defined by a pair of values: $p'_{\omega|A,u=0}$ and $p'_{\omega|A,u=1}$,

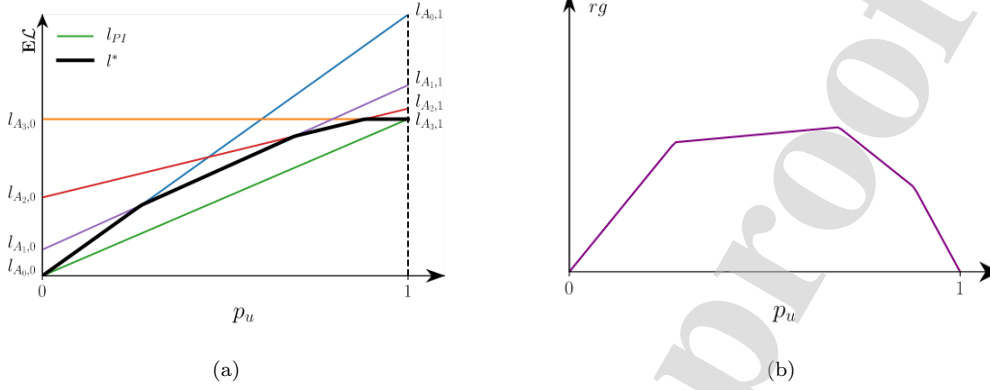


Figure 2: Expected loss function (a) and corresponding regret (b) for a global problem with 4 possible actions.

173 which are the probabilities that $u' = 0$ given action A and given $u = 0$ or
 174 $u = 1$, respectively. Then, $l_{A,0} = p'_{\omega|A,u=0}C_F + C_A$ and $l_{A,1} = p'_{\omega|A,u=1}C_F + C_A$,
 175 with $C_A = \mathcal{L}_{\Pi}(A)$, represent the expected losses when u is zero and one (i.e.
 176 when the system is not working and is working), respectively, for action A .
 177 For each pair of losses $l_{A,0}$ and $l_{A,1}$, with $0 \leq l_{A,0} - l_{A,1} \leq C_F$, one can
 178 find a pair of values $C_A = l_{A,1}$ and $p'_{\omega|A,u=0} = (l_{A,0} - l_{A,1})/C_F$, to represent
 179 the target losses, assuming that no maintenance action makes the system
 180 degrade, so $p'_{\omega|A,u=1} = 0$. In this interpretation C_A is the cost for repairing,
 181 and $p'_{\omega|A,u=0}$ is the probability that the repair is ineffective. The agent has
 182 to find an optimal trade-off between implementing more expensive actions
 183 related to a low risk, and less expensive actions related to a higher risk.

184 The corresponding expected loss under action A is a linear function of
 185 the system failure probability p_u :

$$l_A(p_u) = \mathbb{E}_u \mathbb{E}_{u'|u,A} l(u', A) = p_u l_{0,A} + (1 - p_u) l_{1,A} \quad (8)$$

186 By taking the minimum among available actions in domain \mathcal{A} , we define

187 the optimal loss by concave function $l^*(p_u) = \min_A l_A(p_u)$. Thus, the prior
 188 expected loss of Eq.(2) for the global metric is $L_\pi^G = l^*(p_\pi)$ and, following
 189 Eq.(2), the posterior loss inspecting c_i is:

$$L_\omega^G(i) = h_i l^*(p_{\omega|y_i=0}) + (1 - h_i) l^*(p_{\omega|y_i=1}) \quad (9)$$

190 and the VoI, following Eq.(5), is $\text{VoI}_G(i) = L_\pi^G - L_\omega^G(i)$.

191 As a function of p_u , the expected loss with perfect information of u is the
 192 linear function $l_{\text{PI}}(p_u) = p_u l_0^* + (1 - p_u) l_1^*$, with $l_0^* = \min_A l_{0,A} = l^*(1)$ and
 193 $l_1^* = \min_A l_{1,A} = l^*(0)$, and the ‘‘regret’’ is the concave function $rg(p_u) =$
 194 $l^*(p_u) - l_{\text{PI}}(p_u)$, with $rg(0) = rg(1) = 0$. The corresponding prior regret is
 195 $\text{RG}_\pi = rg(p_\pi)$. Because function l_{PI} is linear, the expected posterior loss with
 196 perfect information is $L_{\text{PI}} = l_{\text{PI}}(p_\pi)$, and expected posterior regret inspecting
 197 c_i is $\text{RG}_\omega(i) = L_\omega^G(i) - L_{\text{PI}} = -\text{VoI}_G(i) + L_\pi^G - L_{\text{PI}}$. Hence, component c_{i^*} ,
 198 that maximizes the VoI identified in Eq.(6), also minimizes the expected
 199 posterior regret:

$$i^* = \arg \min_i \text{RG}_\omega(i) \quad (10)$$

200 The global metric depends on the set of pairs of expected losses for all actions
 201 $\{l_{0,A_0}, l_{0,A_0}, l_{0,A_1}, l_{1,A_1}, \dots, l_{A_0,|\mathcal{A}|}, l_{1,A,|\mathcal{A}|}\}$, where $|\mathcal{A}|$ is the cardinality of set \mathcal{A} ,
 202 or, equivalently, on the concave function l^* . Fig 2 shows an example with
 203 $|\mathcal{A}| = 4$ actions available. The binary case is when only $|\mathcal{A}| = 2$ actions are
 204 available: doing-nothing, accepting the risk of paying cost C_F if the system is
 205 not working, with $A = 0$, or repairing the system at cost C_R , with $A = 1$. As
 206 shown in Fig 3a, this setting is defined by $l_{1,0} = 0, l_{0,1} = C_F, l_{0,1} = l_{1,1} = C_R$,
 207 and the corresponding normalized regret function rg/C_R is bi-linear, with
 208 peak $(1 - \tilde{p})$ at $p_u = \tilde{p} = C_R/C_F$.

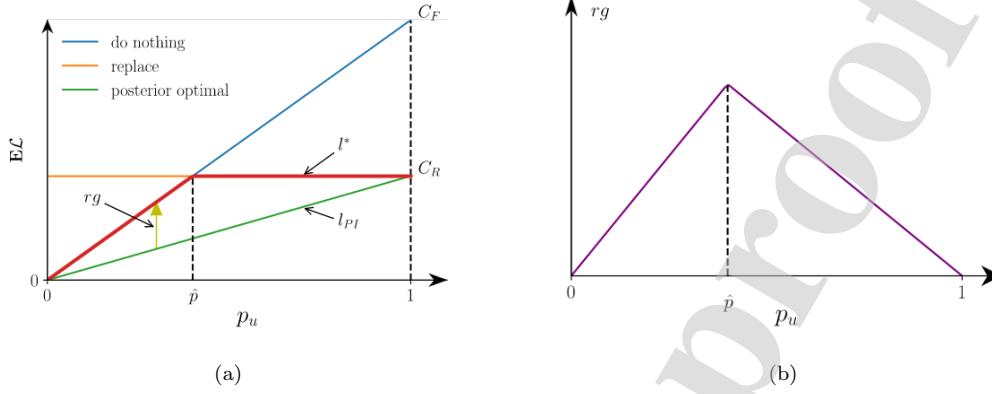


Figure 3: Expected loss (a) and corresponding regret (b) for the binary actions case.

209 *2.4. Local metric*

210 The local metric refers to actions at component level, whose effects de-
 211 pend on components' state s . For this approach, we define each action A
 212 as a vector $\{a_1, a_2, \dots, a_N\}$ of N binary entries, where $a_i = 1$ if the agent
 213 repairs c_i , and $a_i = 0$ otherwise. Hence the cardinality of the **action set** is
 214 $|\mathcal{A}| = 2^N$. We assume that the components' repairs are perfect so that tran-
 215 sition function $p_{s'|s,A}$ is defined as follows: in the vector $s' = [s'_1, s'_2, \dots, s'_N]$
 216 of states after maintenance, $s'_i = 1$ if $a_i = 1$, and $s'_i = s_i$ if $a_i = 0$. Func-
 217 tion $\mathcal{L}_I(\phi(s'))$ is defined as in Section 2.1, while $\mathcal{L}_{II}(A) = C_R^\top \cdot A$, where
 218 $C_R = [C_{R,1}, C_{R,2}, \dots, C_{R,N}]^\top$ is the repair cost vector and $C_{R,i}$ is the cost
 219 of repairing c_i . This model assumes that the accumulated cost is the sum
 220 of repair costs for the individual components. Other cost models, assum-
 221 ing a more complex cost interaction among component' costs, can also be
 222 implemented.

223 After the inspection, the agent selects the optimal subset of components
 224 to repair. **When the inspection outcome is $y_i = c$, the corresponding posterior**

225 expected loss is:

$$L_{\omega|y_i=c}^L = \mathbb{E}_{s|y_i=c} \min_A \mathbb{E}_{s'|s,A} \mathcal{L}(s', A) \quad (11)$$

226 Following Eq.(4), the corresponding expected posterior loss is:

$$L_{\omega}^L(i) = (1 - h_i)L_{\omega|y_i=1}^L + h_i L_{\omega|y_i=0}^L \quad (12)$$

227 and the VoI according to the local metric is $\text{VoI}_L(i) = L_{\pi}^L - L_{\omega}^L(i)$, where
 228 prior loss L_{π}^L is computed as in Eq.(2).

229 3. Metric properties and inspection priorities on typical systems

230 3.1. Nested posterior intervals for global metric

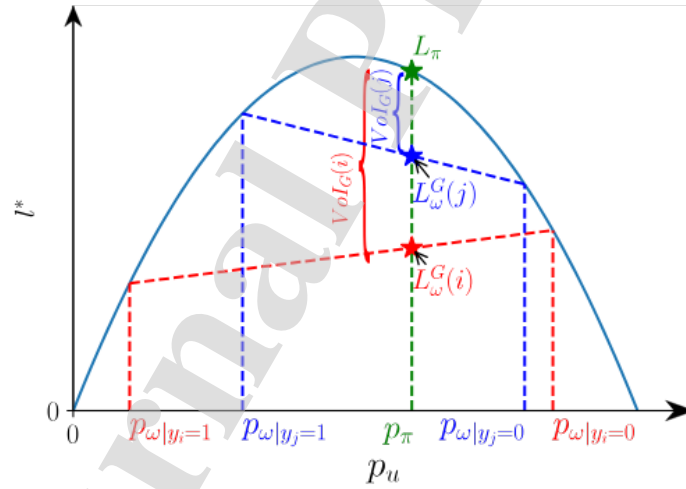


Figure 4: Example of expected loss for the global metric, with nested posterior intervals.

231 As we discussed in Section 2.3, the global metric adopts a univariate
 232 concave function l^* , or rg , of p_u . An example of such a function is shown in
 233 Fig 4, which can also be interpreted as regret, because it is zero at the limits

234 of the probability domain. Inspecting every component c_i , the posterior
 235 system failure probability after an alarm is higher than the priori, which is
 236 in turn higher than the posterior system failure probability after a silence:
 237 $p_{\omega|y_i=1} \leq p_{\pi} \leq p_{\omega|y_i=0}$.

238 Now consider two components c_i and c_j . Suppose that a silence on c_i is
 239 more reassuring than a silence on c_j and an alarm from c_i is more worrying
 240 than an alarm from c_j , i.e. $p_{\omega|y_i=1} \leq p_{\omega|y_j=1}$ and $p_{\omega|y_i=0} \geq p_{\omega|y_j=0}$. Then, for
 241 any concave function l^* (or rg), the posterior loss of inspecting c_i is lower
 242 than the loss of inspecting c_j and the VoI of inspecting c_i is higher than that
 243 of inspecting c_j i.e. $L_{\omega}^G(i) \leq L_{\omega}^G(j)$ and $\text{VoI}_G(i) \geq \text{VoI}_G(j)$. The proof of
 244 this implication is intuitive by examining Fig 4, and it is given formally in
 245 Appendix C.

246 We can also reformulate the implication in terms of “posterior intervals”.
 247 Let us define the posterior interval for c_i as $I_i = [p_{\omega|y_i=1}, p_{\omega|y_i=0}]$. If that
 248 posterior interval contains the corresponding interval for c_j , i.e. if $I_i \supseteq$
 249 I_j , then $\text{VoI}_G(i) \geq \text{VoI}_G(j)$. Hence, the importance ranking is invariant
 250 with respect to the choice of l^* , and all possible global metrics prioritize the
 251 component with larger interval to inspect, consistently with BM defined in
 252 Eq.(1), i.e. $I_i \supseteq I_j \Rightarrow \text{BM}(i) \geq \text{BM}(j)$. However, the reverse implication is
 253 not guaranteed, and Birnbaum’s measure is not necessarily consistent with
 254 the global metric.

255 Moreover, if the posterior intervals are not nested, one can always find
 256 a pair of loss functions $\{l_{\alpha}^*, l_{\beta}^*\}$, so that c_i has a higher VoI than c_j under
 257 l_{α}^* , but a lower VoI under l_{β}^* . For proof, refer to the bi-linear loss function
 258 plotted in Fig 3. If probability \tilde{p} is not in posterior interval I_i (i.e., I_i is on

259 one side of \tilde{p}), then the corresponding VoI, $\text{VoI}_G(i)$, is zero, because the loss
 260 function is linear in that range. If intervals I_i and I_j are not nested, we can
 261 find two disjoint intervals: interval $I_{i \setminus j}$ belongs to I_i but not to I_j , interval
 262 $I_{j \setminus i}$ belongs to I_j but not to I_i . If \tilde{p} is in $I_{i \setminus j}$, then $\text{VoI}_G(i) \geq \text{VoI}_G(j) = 0$,
 263 while if \tilde{p} is in $I_{j \setminus i}$, then $\text{VoI}_G(j) \geq \text{VoI}_G(i) = 0$. This argument shows that,
 264 for not nested posterior intervals, the priority **order** depends on the adopted
 265 loss function.

266 3.2. Global metric for parallel systems

267 A parallel system will function if at least one of its components is intact.
 268 For such systems, the global metric will always give the highest priority to
 269 the most reliable component (i.e., to the one with the lowest marginal failure
 270 probability), independent of the specific loss function l^* adopted, when the
 271 inspection quality is the same for all components. The proof is simple for the
 272 special case of perfect sensors, i.e. when ϵ_{FA} and ϵ_{FS} are zero. In that case,
 273 if a silence is detected for any component, then the posterior system failure
 274 probability is zero. Because the failure of the system implies the failure of
 275 all components, after an alarm on component c_i , p_u becomes $p_{\omega|s_j=0} = p_\pi/p_i$.
 276 Hence, if $p_i \leq p_j$, then $I_i \supseteq I_j$ and, according to the rule illustrated in Section
 277 3.1, we conclude that $\text{VoI}_G(i) \geq \text{VoI}_G(j)$.

278 When sensors are imperfect, the proof is still based on Bayes' formula
 279 (i.e., on the ratio between joint and marginal probabilities). After a silence
 280 on c_i , p_u becomes:

$$p_{\omega|y_i=1} = \frac{p_\pi \epsilon_{\text{FS}}}{1 - h_i} = \frac{p_\pi \epsilon_{\text{FS}}}{1 - \epsilon_{\text{FA}} - K p_i} \quad (13)$$

281 where the second identity follows from Eq.(7), and we note again that K is

282 strictly positive. The corresponding probability after an alarm is:

$$p_{\omega|y_i=0} = \frac{p_{\pi}(1 - \epsilon_{FS})}{h_i} = \frac{p_{\pi}(1 - \epsilon_{FS})}{\epsilon_{FA} + Kp_i} \quad (14)$$

283 The denominator of Eq.(13) decreases monotonically with p_i , and the de-
 284 nominator of Eq.(14) increases monotonically with p_i . Hence, as in the case
 285 of perfect sensors, if $p_i \leq p_j$, then $I_i \supseteq I_j$ and, $\text{VoI}_G(i) \geq \text{VoI}_G(j)$.

286 In summary, the ranking of importance measures follows the opposite of
 287 the marginal failure probability of the components (i.e., the ranking follows
 288 component reliability). Hence, in a parallel system, the component, c_{i^*} ,
 289 with highest VoI is the most reliable component. This result holds for any
 290 interdependence between components' states, that is for any distribution p_s ,
 291 when the inspection quality, defined by parameters ϵ_{FA} and ϵ_{FS} , is the same
 292 for all components.

293 3.3. Global metric for series systems

294 A series system works only if all components function properly. In that
 295 case, the global metric always prioritizes the most vulnerable component, i.e.
 296 the component with the highest prior failure probability, regardless of the
 297 adopted function l^* or the interdependence among components. The proof
 298 is similar to that related to parallel systems. Let us start with the case of
 299 perfect sensors. The posterior system failure probability will become 1 after
 300 an alarm on any component, and will become $p_{\omega|s_j=1} = 1 - (1 - p_{\pi})/(1 - p_i)$
 301 after a silence on component c_i , which monotonically increases with marginal
 302 component failure probability p_i . Hence the most vulnerable component
 303 should be inspected.

304 For imperfect sensors, after a silence on c_i , p_u is (again using Eq.(7)):

$$p_{\omega|y_i=1} = 1 - \frac{(1-p_\pi)(1-\epsilon_{FA})}{1-h_i} = 1 - \frac{(1-p_\pi)(1-\epsilon_{FA})}{1-\epsilon_{FA}-Kp_i} \quad (15)$$

305 After an alarm, that probability is:

$$p_{\omega|y_i=0} = 1 - \frac{(1-p_\pi)\epsilon_{FA}}{h_i} = 1 - \frac{(1-p_\pi)\epsilon_{FA}}{\epsilon_{FA}+Kp_i} \quad (16)$$

306 The denominator of the fraction in Eq.(15) monotonically decreases with
 307 p_i , and the denominator of Eq.(16) monotonically increases with p_i . Hence,
 308 if $p_i \geq p_j$, then $I_i \supseteq I_j$ and $\text{VoI}_G(i) \geq \text{VoI}_G(j)$, as in the case of perfect
 309 sensors. So, in a series system, regardless of the interdependence between
 310 components, the inspection ranking follows the marginal component failure
 311 probability, and c_{i^*} is the most vulnerable component.

312 In other words, the most vulnerable component, c_{i^*} , is the one to inspect
 313 because detecting a silence on that component (i.e. $y_{i^*} = 1$) induces the
 314 highest reduction of p_u , and an alarm (i.e. $y_{i^*} = 0$) induces the highest
 315 increment in that probability. Although the former property is almost trivial,
 316 the latter may be less intuitive. After all, c_{i^*} was (relatively) likely to be
 317 damaged; thus, why does an alarm on that component produce the more
 318 “surprising” result on the system reliability (compared with alarms on less
 319 vulnerable components)? For imperfect inspections, two factors affect the
 320 posterior probability. On one hand, after detecting an alarm on c_{i^*} , the
 321 system can still count on the other components, which are more reliable than
 322 c_{i^*} (instead, after an alarm on a safer component, the system can only count
 323 on more vulnerable components). Hence, this factor suggests that an alarm
 324 of c_{i^*} is less worrying than an alarm on others. Conversely, following Bayes’
 325 rule, an alarm on c_{i^*} produces a relatively high posterior failure probability

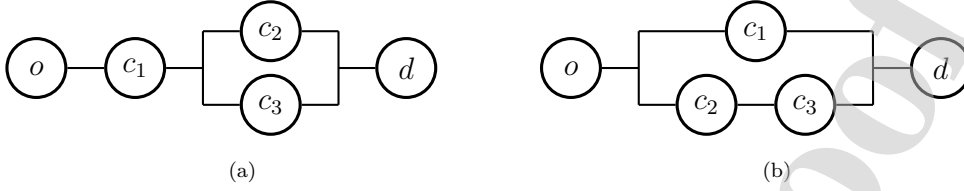


Figure 5: Block diagrams for series-parallel (a) and parallel-series (b) 3-component system.

326 (at component and at system level), because of the high prior probability that
 327 c_{i^*} is damaged. For safer components, the impact of the alarm is diluted by
 328 the more optimistic prior information, and the posterior failure probability
 329 after an alarm is lower at component level (obviously) and at system level, as
 330 formally proved by Eq.(16). Hence, this latter factor dominates the former
 331 factor, and c_{i^*} has the highest VoI. This result depends on the assumption
 332 that the sensor accuracy is uniform among components. If the accuracy was
 333 higher for a specific component, that component could have the highest VoI,
 334 even if it was not the most vulnerable component.

335 3.4. Global metric for general systems

336 If the posterior probability interval related to one component nests all
 337 the others, then the rule of Section 3.1 identifies the optimal component
 338 to inspect. For general systems, the global metric does not always select
 339 the most vulnerable or the most reliable component, because the posterior
 340 intervals may not be nested, and the rule does not apply.

341 We illustrate this by discussing two simple examples of 3-component sys-
 342 tems, with perfect sensors as shown in Fig 5. The system functions if there
 343 is an intact path from origin node o to destination node d . Fig 5a shows a
 344 system in which component c_1 is in series with a parallel subsystem composed

345 of components c_2 and c_3 . Intuitively, component c_1 should be inspected, be-
 346 cause it is a “bottleneck” of the system; thus, it seems topologically more
 347 important. Detecting that c_1 is not working takes p_u from one, i.e., to a
 348 higher value related to an alarm on c_2 or on c_3 . A silence detected on c_1
 349 takes p_u to the joint failure probability, $p_{\omega|s_1=1}$, which is determined by com-
 350 ponents c_2 and c_3 , and is (p_2p_3) if they are independent. Instead, a silence
 351 detected on c_2 (or on c_3) takes p_u to $p_{\omega|s_2=1} = p_1$. Hence, posterior interval
 352 I_1 contains the other two if $p_{\omega|s_1=1}$ is less than p_1 . Conversely, if p_1 is less
 353 than $p_{\omega|s_1=1}$, the posterior intervals are not nested, and the priority depends
 354 on the selected loss function l^* . This result confirms the intuition that if c_1
 355 is much safer than the other components, it may not be the most important
 356 component to inspect (trivially, if p_1 is zero while p_2 and p_3 are positive, then
 357 c_1 has the lowest priority).

358 In the example of Fig 5b, component c_1 is parallel with a series subsystem
 359 composed of components c_2 and c_3 . Again, c_1 seems topologically more
 360 important. After a silence on c_1 , p_u is zero, a value lower than the value
 361 related to silence on c_2 or on c_3 . An alarm on c_1 takes p_u to $1 - r_{2,3}$, where
 362 $r_{2,3}$ is the joint survival probability of the other two components, that is
 363 $(1 - p_2)(1 - p_3)$ for independent components, whereas an alarm on c_2 (or on
 364 c_3) takes p_u to p_1 . Hence, posterior interval I_1 contains the others if p_1 is less
 365 than $1 - r_{2,3}$ i.e., for independent components, if p_1 is less than $p_2 + p_3 - p_2p_3$.
 366 Approximating this latter value with $p_2 + p_3$, we conclude that the global
 367 metric gives higher priority to c_1 when p_1 is less than $(p_2 + p_3)$. If p_1 is
 368 higher than $(p_2 + p_3)$, priority depends on the selected loss function l^* . This
 369 **conclusion** confirms the intuition that, if c_1 is much more vulnerable than the

370 other components, it is better to inspect others (in the limit case where p_1 is
 371 one, $\text{VoI}_G(1)$ is zero). These two examples illustrate how the topological role
 372 of a component matters, but also its failure probability: in some schemes a
 373 high failure probability guarantees a high priority.

374 We discuss now a more general example, focusing on two components, c_1
 375 and c_2 . The components' roles are described completely by the system fail-
 376 ure probability for each of the 2^2 joint conditions of the pair of components,
 377 that we assume as $p_{\omega|s_1=1,s_2=1} = 0.5\%$, $p_{\omega|s_1=1,s_2=0} = p_{\omega|s_1=0,s_2=1} = 2.5\%$,
 378 $p_{\omega|s_1=0,s_2=0} = 90\%$ (so the roles played by the two components are the
 379 same). We also assume that $p_1 = 1\%$ and $p_2 = 20\%$ (so that the c_1 is
 380 significantly more reliable than c_2), the states of all components are inde-
 381 pendent, and inspections are perfect (i.e., $y_i = s_i$). Fig 6 shows the diagram
 382 of a system consistent with these values. The interval of posterior proba-
 383 bilities I_i is $[0.90\%, 20.0\%]$ for $i = 1$ and $[0.52\%, 3.38\%]$ for $i = 2$, whereas
 384 p_π is 1.09% (these results are directly related to the assumed values, e.g.
 385 $p_{\omega|y_2=1} = p_{\omega|s_1=0,s_2=1}p_1 + p_{\omega|s_1=1,s_2=1}(1 - p_1)$). The intervals are not nested;
 386 hence, the rule in Section 3.1 does not apply, and the VoI depends on the
 387 specific function l^* . Fig 7 refers to the bi-linear regret function for binary
 388 actions plotted in Fig 3, and mentioned in Section 2.2, with peak at \tilde{p} . The
 389 figure shows how the VoI related to each component, normalized by prior
 390 regret RG_π , varies as a function of \tilde{p} . When \tilde{p} is below $p_{\omega|y_2=1} = 0.52\%$ (i.e.,
 391 when C_R is below 0.52% of C_F), the VoI of each component is nil, because the
 392 posterior decision is always to repair. Then, $\text{VoI}_G(2)$ increases up to about
 393 42% of RG_π when $\tilde{p} = p_\pi$ (i.e., for that condition observing y_2 is worth 42%
 394 of the value of observing u), then it decreases to zero at $\tilde{p} = p_{\omega|y_2=0} = 3.38\%$.

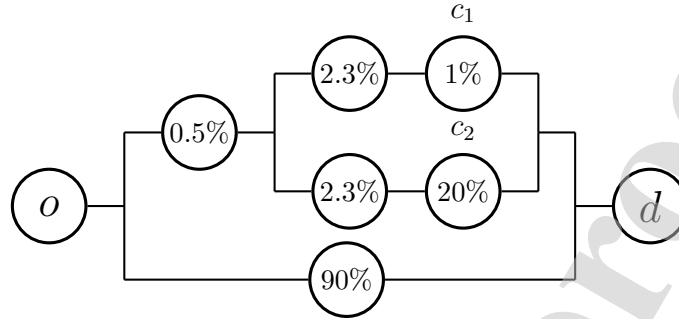


Figure 6: Block diagram for a system where posterior probability intervals for components c_1 and c_2 are not nested

395 For \tilde{p} higher than $p_{\omega|y_2=0}$, the posterior decision is always to accept the risk.
 396 The behavior of $\text{VoI}_G(1)$ is similar; it is zero outside I_1 , and it peaks at p_π ,
 397 where it is about 17% of RG_π . Clearly, the optimal inspection decision de-
 398 pends on \tilde{p} , i.e. on the decision-making problem shaping function l^* . This
 399 conclusion is apparent in Fig 7, if the repair cost is cheaper than 2.5% of the
 400 cost of failure, it is more convenient to inspect the more reliable c_2 , whereas
 401 it is better to inspect the less reliable c_1 for a higher repair cost.

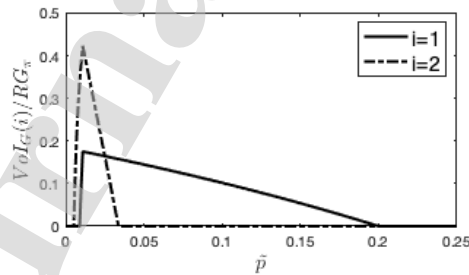


Figure 7: Normalized VoI depending on peak probability \tilde{p} .

402 *3.5. Local metric on parallel systems*

403 The local metric, as defined in Section 2.4, will select the most reliable
 404 component in a parallel system, consistently with the global metric. This
 405 selection behavior is because, [in a parallel system](#), repairing one component
 406 guarantees the functioning of the system. Hence, the agent faces a binary
 407 decision: do nothing or repair the less expensive component, at cost $\min_i C_{R,i}$.
 408 This problem setting is equivalent to that of the global metric, with the bi-
 409 linear function l^* of Fig 3a. Therefore, the local and global metrics have
 410 identical conclusions about the optimal inspection.

411 *3.6. Local metric on series systems*

412 With the local metric, the optimal component to inspect in a series system
 413 is not always the most vulnerable component, i.e. the component identified
 414 by the global metric. We start discussing the case of a system with two
 415 components, c_1 and c_2 , with identical repair costs, $C_{R_1} = C_{R_2} = C_R$, and
 416 equipped with perfect sensors. Let us also assume that $C_R \leq C_F/2$, so
 417 that the cost for repairing both components is less than the failure cost.
 418 Hence, if any component c_i is detected as damaged, it is necessary to repair it
 419 ($A_i = 1$), to avoid paying the failure cost. After the repair, the system failure
 420 probability is the posterior failure probability of the uninspected component.
 421 That uninspected component should be also repaired if the corresponding risk
 422 is above the repair cost, so that the posterior expected maintenance cost for
 423 that component is $R(i, x) = \min\{C_R, p_{\omega|s_i=x, A_i=1-x} C_F\}$, with $x = 0$. Instead,
 424 if the inspected component works, it has not to be repaired, and the state of
 425 the uninspected component is decided by comparing repair cost and system
 426 failure risk, so that the expected posterior cost is $R(i, 1)$. Hence, the expected

427 posterior loss is $L_\omega^L(i) = p_i C_R + p_i R(i, 0) + (1 - p_i) R(i, 1)$. In the special
 428 case of independent components, for any outcome x , probability $p_{\omega|s_i=x}$ is
 429 identical to the prior failure probability p_j of the uninspected component
 430 c_j , so that $R(i, 0) = R(i, 1)$, and $L_\omega^L(i) = p_i C_R + \min\{C_R, p_j C_F\}$. If we
 431 refer to bi-linear regret function rg of Fig 5b, we conclude that, for each
 432 component c_i , $\text{VoI}_L(i) = \text{RG}_\omega(i)$; thus, we should inspect the component
 433 with the higher value of $\text{RG}_\omega(i)$. If both prior failure probabilities are below
 434 $\tilde{p} = C_R/C_F$, the local metric will prioritize the more vulnerable component.
 435 However, if the failure probability of a component is above \tilde{p} , then the higher
 436 that probability, the lower the corresponding VoI. Fig 8 shows the optimal
 437 inspection policies for $\tilde{p} = 0.2$, $p_1 \geq p_2$, and different correlation coefficient ρ
 438 between variables s_1 and s_2 . The joint probability can be defined given the
 439 correlation coefficient ρ and the marginal probability p_1 and p_2 . For example,
 440 the joint probability of both components work is:

$$\mathbb{P}[s_1 = 1, s_2 = 1] = \rho \sqrt{p_1(1-p_1)p_2(1-p_2)} + (1-p_1)(1-p_2) \quad (17)$$

441 . We have discussed the case when ρ is zero. When it is positive, the domain
 442 of feasible pairs (p_1, p_2) shrinks but, inside the feasible domain, the region ex-
 443 pands where it is more convenient to inspect the more vulnerable component.
 444 When the correlation is negative, for any feasible pair $\{p_1, p_2\}$, the VoI is the
 445 same for both components. We can provide a simple approximation for se-
 446 ries system if their states are independent and the failure probabilities are
 447 relatively low. In that case, the risk $\mathbb{E}[\mathcal{L}_I]$ can be approximated as that of
 448 a ‘‘cumulative system’’ [13]. In a cumulative system, individual costs are as-
 449 sociated with the failure of each component, and the costs are accumulated
 450 to obtain the system-level cost (hence, no component-to-system function ϕ

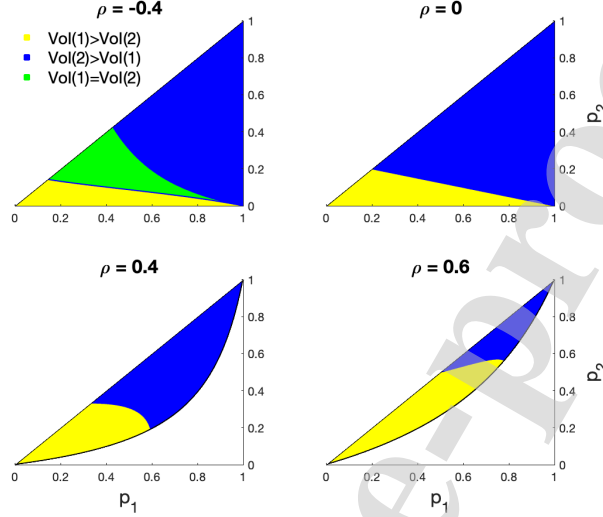


Figure 8: Optimal inspection i^* when two components are dependent, depending on correlation ρ .

451 is defined for these systems). To illustrate this approach, we recall that, for
 452 a series system with independent components, the risk is:

$$\mathbb{E}_{\text{ser.}}[\mathcal{L}_I] = C_F \left[1 - \prod_i (1 - p_i)^{1-A_i} \right] \quad (18)$$

453 For a cumulative system with component failure cost C_F , it is:

$$\mathbb{E}_{\text{cum.}}[\mathcal{L}_I] = C_F \sum_i (1 - A_i) p_i \quad (19)$$

454 By linearizing the former expression (neglecting higher order terms), the
 455 two risks become identical. For a cumulative system, it is straightforward
 456 to evaluate the benefit of inspecting component c_i which is related to the
 457 selection of action A_i , and when sensors are perfect, it is $\text{VoI}_L(i) = \text{RG}_\omega(i)$
 458 (for imperfect sensors, one has to subtract the posterior regret). These results
 459 are also consistent with the case in which $N = 2$.

460 *3.7. Connections and differences between local and global metrics*

461 Whereas the local metric follows the traditional Bayesian pre-posterior
 462 analysis for a selected class of actions and losses, the global metric takes
 463 inspiration from the BM, and it focuses on the impact of information on
 464 the system failure probability, neglecting the impact on damage localization.
 465 This latter approach allows for a simpler optimization and VoI analysis, and
 466 the intuition supporting it is that a component should receive a high priority
 467 if its inspection outcome is highly informative on the system state.

468 The two metrics refer to different problem classes, which are not nested
 469 one into another (given the restrictive rules we impose to the local met-
 470 ric). On the one hand, information in the local metric keeps referring to
 471 local quantities, i.e. the damage condition of individual components, and
 472 therefore is key to supporting local repairing, i.e. the repairing of those com-
 473 ponents. In the global metric approach, local information is neglected, and
 474 the optimal posterior action is only a function of the posterior system failure
 475 probability. This indicates that the local metric cannot be generally reduced
 476 to an equivalent global one. On the other hand, the global metric cannot be
 477 generally reduced to the local one either. This is because, for example, in
 478 the local metric we assume that each action identifies whether to repair or
 479 not for each component; hence, for a system with N components, there are
 480 2^N available actions in local metric, whereas the global metric can define an
 481 arbitrary number of actions (each associated with different repair cost C_A
 482 and posterior probability $p'_{\omega|A,u=0}$). Also, the assumptions of perfect repairs
 483 and of additive repair costs impose constraints to the available actions in the
 484 local metric, while the global metric is not subject to these limitations.

485 We note that for both metrics, following Section 2.1, one can define a
486 concave function l^* on the belief p_s of the joint condition state s of all com-
487 ponents, i.e. on a N dimensional domain (with the linear constraint that
488 $\sum_j p_s(j) = 1$). However, this function can be transformed into a univariate
489 function of the system failure probability p_u , as illustrated in Section 2.3,
490 only in the the global metric.

491 These major difference between the local and global metrics can be high-
492 lighted on a paradigmatic case. Consider a two-component series sub-system
493 in a larger system, where one of the two components is working and the other
494 is not, and the working component is one (or the other), with probability $1/2$.
495 Hence, the states of the two components are perfectly negatively correlated.
496 The perfect inspection of one component enables the agent to identify the
497 malfunctioning component (it is the inspected one, if the outcome is “alarm”,
498 or the other, if the outcome is “silence”). From the local metric perspective,
499 this information is relevant because it enables repairing the malfunctioning
500 component with perfect information. However, from the perspective of the
501 global metric, the VoI of inspecting any of the two components is nil, because
502 the posterior system failure probabilities are identical to the prior one. From
503 this latter perspective information has value only if the inspection outcome
504 has an impact on the system failure probability: a component is important
505 if silence is good news and alarm bad news for system reliability. The exam-
506 ple shows how information can have value for supporting local maintenance,
507 regardless of its impact on system reliability.

508 4. Computational complexity and Heuristic

509 4.1. Complexity of VoI computation

510 The computational complexity of solving Eq.(6) varies with different met-
 511 rics, but is generally intimidating for large systems. The core step of the com-
 512 putational process is solving the reliability problem, identifying the system
 513 failure probability p_u , depending on actions and observations. This analysis
 514 is nested into the optimization of the maintenance actions.

515 The general network reliability problem is NP-hard [2] [19], but numerous
 516 approximations and bounds have been proposed. To compute the risk $\mathbb{E}[\mathcal{L}_I]$,
 517 one has to assess the system connectivity for each of the 2^N system states. A
 518 matrix-based method was proposed to compute system reliability based on
 519 a components' condition matrix with each row representing one state, and
 520 a binary condition vector with each entry representing whether the system
 521 is functioning at that specific components' state [21]. The general computa-
 522 tion complexity of the method is $\mathcal{O}(N \times 2^N)$. When the joint distribution
 523 of the components' states is known, with the components' condition matrix
 524 and the binary system condition vector previously computed, the computa-
 525 tional complexity of system reliability is linear with respect to the system
 526 states. An approximate estimation can also be achieved based on Monte
 527 Carlo simulations [11].

528 For the global metric, $\mathbb{E}[\mathcal{L}_{II}]$ can be determined in $\mathcal{O}(1)$ time once we
 529 compute the posterior system failure probabilities. Therefore, to select the
 530 component with highest VoI among N components will cost $\mathcal{O}(N \times 2^N)$.

531 For the local metric, there is an additional computation step before as-
 532 sessing the VoI. $\mathbb{E}[\mathcal{L}]$ is optimized among 2^N combinations of maintenance

533 actions. Suppose that, on the basis of different inspection outcomes, the
 534 agent can select an arbitrary subset of the components to repair; then, the
 535 computation **complexity** is generally $\mathcal{O}(N \times 2^N \times 2^N) = \mathcal{O}(N \times 2^{2N})$.

536 4.2. Approximation for local metric

537 In this section, we propose a simple heuristic approach for approximat-
 538 ing the local metric, to reduce the computational complexity related to the
 539 optimization of maintenance actions depending on the inspection outcome.

540 Let us define $A_\pi = \{a_{\pi,1}, a_{\pi,2}, \dots, a_{\pi,N}\}$ as the prior maintenance plan,
 541 $A_\omega = \{a_{\omega,1}, a_{\omega,2}, \dots, a_{\omega,N}\}$ as the posterior one, and L_π is the prior optimal
 542 loss related to A_π , as defined in Section 2.1. We assume that A_π and L_π have
 543 been identified. Consider inspecting component c_i . The proposed heuristic
 544 assumes that the agent confirms all actions for uninspected components (i.e.,
 545 $\forall j \neq i, a_{\omega,j} = a_{\pi,j}$). Only the posterior action on the inspected component,
 546 $a_{\omega,i}$, depends on the inspection's outcome, y_i . If the prior action for c_i is to
 547 do-nothing (i.e., if $a_{\pi,i} = 0$) and the inspection's outcome is silence (i.e., if
 548 $y_i = 1$), or if the prior action is to repair (i.e., if $a_{\pi,i} = 1$) and the inspection
 549 produces an alarm (i.e., if $y_i = 0$), then the agent will confirm the prior
 550 action also for the inspected component (i.e., if $y_i \neq a_{\pi,i}$, then $A_\omega = A_\pi$).
 551 Instead, if an alarm is detected on a **previously** unrepaired component, or if
 552 a silence is detected on a **previously** repaired component (i.e., if $y_i = a_{\pi,i}$),
 553 then the agent **considers** the two alternatives: repair or not repair c_i . One
 554 of the two alternatives is, again, to completely confirm the prior plan (i.e.
 555 $A_\omega = A_\pi$); thus, the prior loss L_π associated with this option is already
 556 known. The agent computes the expected cost of the alternative plan (in
 557 which only action $a_{\omega,i}$ is reversed), and executes the best option, i.e. the

558 option related to the minimum expected cost. The computational saving is
 559 related to the avoidance of [searching for optimal posterior action in the full](#)
 560 [set \$\mathcal{A}\$](#) .

561 One argument supporting the choice of this heuristic is [the consistency](#)
 562 [with the optimal behavior in some special cases, for example, when the high](#)
 563 [penalty of a system collapse forces the agent to be conservative](#). To model this
 564 scenario, suppose that (i) the prior decision is to do-nothing (i.e., $\forall i, a_{\pi,i} = 0$),
 565 that (ii) a detected silence cannot increase the system failure probability (i.e.
 566 $\forall i, p_{\omega|y_i=1} \leq p_{\pi}$), that (iii) the do-nothing option is still optimal when the
 567 probability of failure decreases and that (iv) a component sending an alarm
 568 must be repaired, because its posterior failure probability is too high to
 569 be tolerated. Condition (iii) is not obviously satisfied even if the first two
 570 conditions are satisfied, because the prior decision might also be doing noth-
 571 ing for another reason, i.e., the agent is pessimistic about the components'
 572 conditions. For such a pessimistic agent, it is not worth repairing any set
 573 of components; repairing few components may be ineffective, and repairing
 574 many components may be too expensive. However, detecting a functioning
 575 component may [improve the expectation of the system and persuade the](#)
 576 [pessimistic agent to invest in repairing other components](#). Condition (iii)
 577 forbids the occurrence of this process, by assuming [the agent's optimism](#)
 578 [about the system condition](#). To prove that the heuristic is optimal under
 579 conditions (i-iv), we must show that the optimal response to an alarm on
 580 component c_i cannot be to repair any other component. Because of (iv), c_i
 581 must be repaired. Now suppose that component c_j is also to be repaired.

582 This condition implies the following inequality:

$$C_{R,i} + C_{R,j} + C_F p_{\omega|y_i=0,a_i=1,a_j=1} \leq C_{R,i} + C_F p_{\omega|y_i=0,a_i=1} \quad (20)$$

583 If, as assumed before, repairs are perfect and components' states are indepen-
 584 dent, then $p_{\omega|y_i=0,a_i=1} = p_{\omega|s_i=1} = p_{\omega|y_i=1}$, and $p_{\omega|y_i=0,a_i=1,a_j=1} = p_{\omega|s_i=1,s_j=1} =$
 585 $p_{\omega|y_i=1,a_j=1}$, so the inequality Eq.(20) can be re-written, subtracting $C_{R,i}$ from
 586 both terms as:

$$C_{R,j} + C_F p_{\omega|y_i=1,a_j=1} \leq C_F p_{\omega|y_i=1} \quad (21)$$

587 This equation indicates that repairing c_j should be the optimal response
 588 to a silence on c_i , but this response violates conditions (i-iii), which show
 589 that only c_i should be repaired after receiving an alarm on it. Of course, if
 590 conditions (i-iv) are not satisfied, there is no guarantee that the heuristic is
 591 truly optimal.

592 The VoI defined by the heuristic is certainly non-negative, as the prior
 593 maintenance plan can be confirmed, if the collected observations do not sug-
 594 gest any improvement. Moreover, given that the heuristic limits the domain
 595 of the posterior actions, the corresponding VoI cannot be higher than the
 596 VoI assessed by the local metric.

597 5. Examples of System Analysis

598 We analyze three examples of systems. The first example is the 6-
 599 component system in Fig 9, in which the failure probability of each compo-
 600 nent is listed inside the corresponding node. We start by considering perfect
 601 inspections and independent components. The corresponding values of the
 602 BMs are shown in Fig 10a, and c_2 has the highest importance in BM.

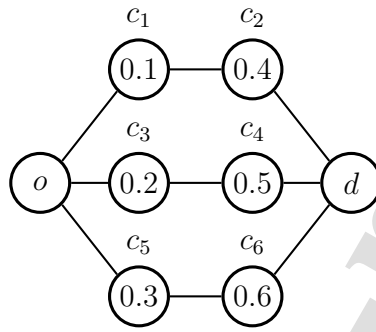


Figure 9: Block diagram for the counter-intuitive example 1

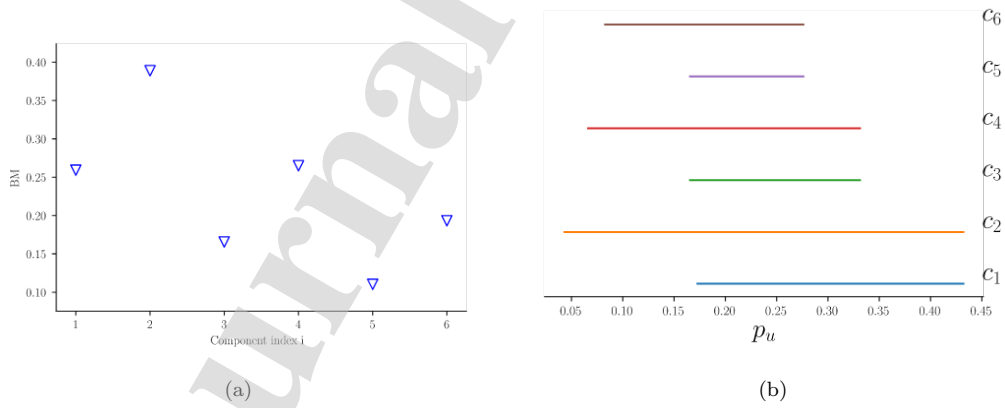


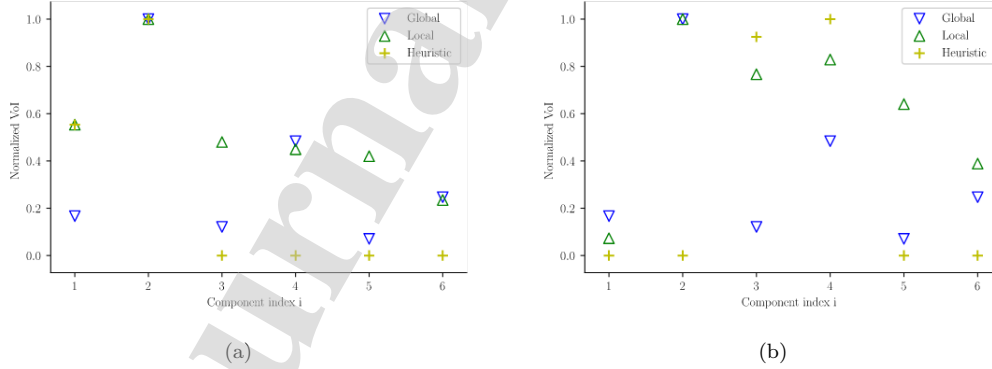
Figure 10: BM for the system in Fig 9 (a), and corresponding posterior intervals (b).

603 Fig 10b shows the posterior probabilities intervals I_i for all components.
 604 All intervals are nested in I_2 . Thus, component c_2 has the highest VoI,
 605 according to the global metric, regardless of the loss function we adopt (and
 606 so it has also the highest BM, as noted above). We divide the VoI of each
 607 component by the maximum VoI of all the components under the same metric
 608 obtain a normalized VoI. The normalized VoI under the global metric, for
 609 loss function $l^*(p_u) = p_u(1 - p_u)$, is shown in Fig 11. For the local metric, we
 610 assume that $C_F/C_{R,i} = 10$, for every component c_i , i.e. the cost of system
 611 failure is ten times the cost of repairing one component. The optimal prior
 612 maintenance action is to [repair](#) component c_2 . As shown in Fig 11a, the local
 613 metric and the heuristic both identify c_2 as the component with the highest
 614 VoI.

615 However, if the maintenance cost for c_2 increases to $C_F/C_{R,2} = 5$ while
 616 the cost for the others remains the same, the optimal prior action becomes
 617 repairing c_4 . Table 2 reports the optimal posterior actions depending on
 618 the inspection outcome, for this new assumption on the costs. As shown in
 619 Fig 11b, the local metric still gives the highest inspection priority to c_2 (as
 620 the global metric does), but the heuristic selects c_4 instead. The difference
 621 between the local metric and heuristic approach is because the posterior
 622 optimal action may not include repairing c_4 (e.g. after a silence on c_2), or it
 623 may include repairing uninspected components (e.g. after an alarm on c_1, c_3
 624 is to be repaired). [Though the VoI calculated from the heuristic approach is
 625 no higher than that from the true optimal solution, the heuristic approach
 626 overestimates the priority of inspecting \$c_4\$ over \$c_2\$, which is inconsistent with
 627 the local metric.](#)

Inspected component	Inspection outcome	
	Silence ($y_i = 1$)	Alarm ($y_i = 0$)
c_1	$\{c_4\}$	$\{c_3, c_4\}$
c_2	\emptyset	$\{c_3, c_4\}$
c_3	$\{c_4\}$	$\{c_3, c_4\}$
c_4	\emptyset	$\{c_4\}$
c_5	$\{c_6\}$	$\{c_4\}$
c_6	$\{\emptyset\}$	$\{c_6\}$

Table 2: Posterior subset of components to be repaired for the system in Fig 9.

Figure 11: Normalized VoI for the system in Fig 9, with $C_F/C_{R_i} = 10$ (a), and with $C_F/C_{R_2} = 5, C_F/C_{R_i} = 10, \forall i \neq 2$ (b).

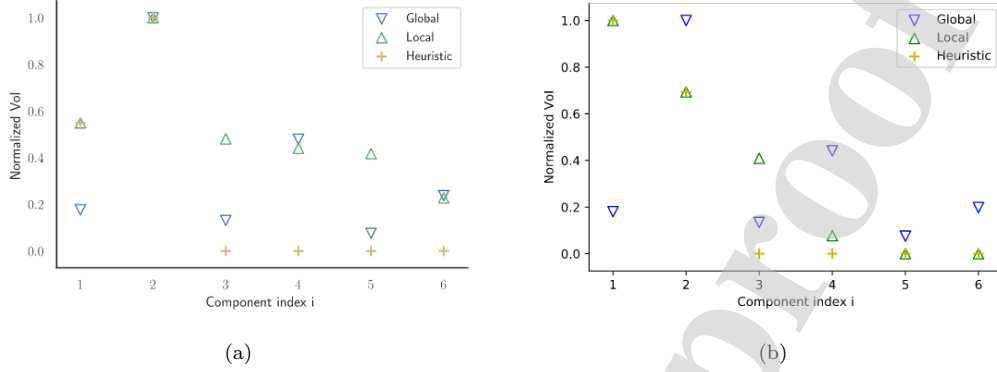


Figure 12: Normalized VoI for the system in Fig 9, with $\epsilon_{FA} = \epsilon_{FS} = 0.01$ (a), and $\epsilon_{FA} = 0.01, \epsilon_{FS} = 0.40$ (b).

628 Error rates in imperfect inspections also affect the optimal decision. We
 629 now assume, again, that $C_F/C_{R,i} = 10$ for every component c_i , but inspec-
 630 tions are imperfect; when $\epsilon_{FA} = \epsilon_{FS} = 0.01$, the corresponding VoI, shown
 631 in Fig 12a, is similar to the perfect inspection case shown in Fig 11a, and
 632 c_2 has the highest VoI. But when the type II error rate ϵ_{FS} is increased to
 633 0.40, the VoI changes to Fig 12b, and component c_1 gains the highest priority
 634 according to the local metric and heuristic approach.

635 The second example is a 16-component system represented in Fig 13. The
 636 components have different topological importance: components c_1, c_4 and c_8 ,
 637 and the ones symmetric to them, can be considered as “bottlenecks”, with
 638 respect to other components.

639 We assume that the marginal probability of failure is $p_i = 0.01$ for every
 640 component c_i . For the global metric, we use $l^*(p_u) = p_u(1 - p_u)$ as a loss
 641 function, and the corresponding VoI is shown in Fig 14a. For the local metric,
 642 we assume the cost ratio between system failure penalty and component

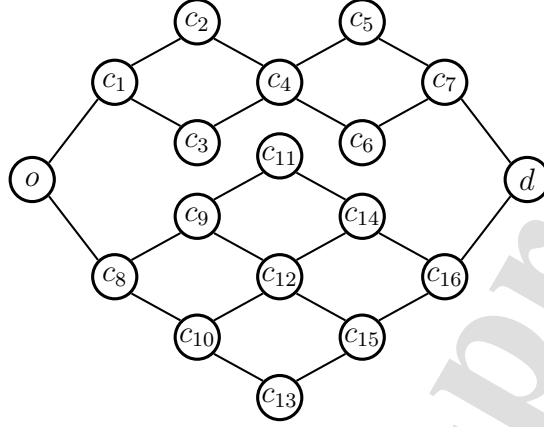


Figure 13: Block diagram of 16 component system

643 maintenance is $C_F/C_{R,i} = 10^3$ for every component c_i , so that the resulting
 644 optimal prior maintenance action is to repair no component. Under the local
 645 metric, c_8 and c_{16} have the highest VoI, followed by c_1 , c_4 and c_7 . The
 646 heuristic approach gives the same result as the [local metric and the global](#)
 647 [metric](#).

648 However, if $C_F/C_{R,i}$ [increases](#) to 10^4 for every component c_i , the new
 649 optimal prior maintenance action becomes repair the symmetric bottlenecks
 650 c_8 and c_{16} . The VoI for the global and local metrics and the heuristic with
 651 this new assumption on costs is illustrated in Fig 14b. The local VoI of
 652 inspecting c_2 , c_9 and the components symmetric to them is now nil, because
 653 the cost for system failure is so (relatively) high, that the agent will not alter
 654 the prior action even if a silence is received on these components.

655 Depending on the setting, the bottleneck components may not always
 656 have the highest VoI. If $p_{11} = 0.5, p_{12} = 0.4, p_{13} = 0.3, p_i = 0.01, i \neq 11, 12, 13$
 657 and $C_F/C_{R,i} = 1000$ for every component c_i , the VoI is that shown in Fig 15.

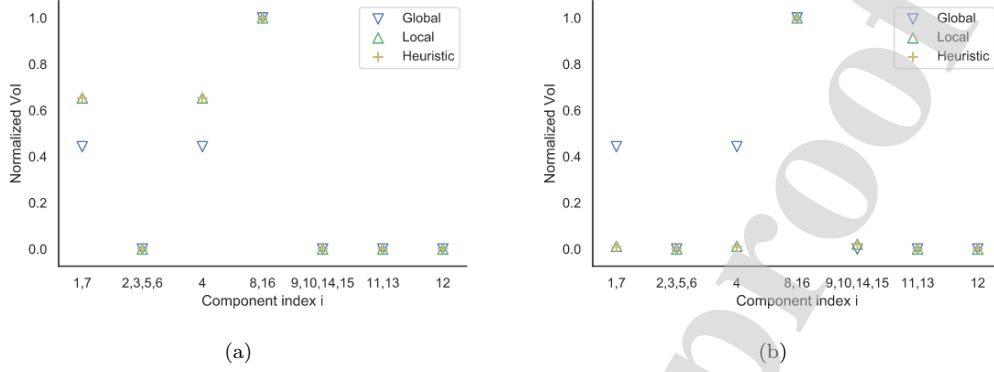


Figure 14: Normalized VoI for the system in Fig 13 with different maintenance cost, with and $C_F/C_{R,i} = 10^3$ (a) and $C_F/C_{R,i} = 10^4$ (b)

658 Now the optimal prior action is to repair c_{12} . The global metric prioritizes c_1 ,
 659 c_4 and c_7 for inspection, but the local metric prioritizes c_{13} , even though it is
 660 not the most vulnerable component (which is c_{11}). After c_{13} , the components
 661 with high VoI under local metric will be c_{12} and c_{11} . Instead, the heuristic
 662 approach assigns the highest VoI to c_{12} . This assignment occurs because,
 663 when the inspection of c_{11} or c_{13} receives silence, the optimal action is to do
 664 nothing, but the heuristic approach forces the agent to at least execute the
 665 prior plan.

666 The third example is taken from [21] and it represents a two-line electrical
 667 substation with 12 components with 6 different functions as illustrated in Fig
 668 16: DS - Disconnect Switch, CB - Circuit Breaker, PT - Power Transformer,
 669 DB - Drawout Breaker, TB - Tie Breaker, FB - Feeder Breaker. We assume
 670 that the marginal failure probability of the components with function DS,
 671 CB or DB is 9.53×10^{-3} , and that of components with function FB, PT
 672 and TB is 2.32×10^{-3} . For every component c_i , costs are defined by ratio

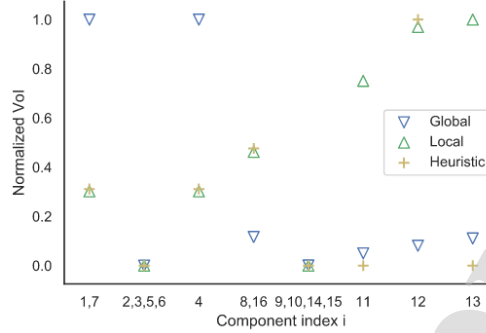


Figure 15: Normalized VoI for the system in Fig 13 with $p_i = 0.01, i \neq 11, 12, 13$, $p_{11} = 0.5, p_{12} = 0.4, p_{13} = 0.3$ and $C_F/C_{R,i} = 10^3$

673 $C_F/C_{R_i} = 1000$. In this example, we investigate how the correlation between
674 a component's state affects the VoI. If all the components are statistically de-
675 pendent, complexity of computing the system failure probability may become
676 intractable. Conditional independence between component events given the
677 outcomes of a few random variables representing the source of common effects
678 was assumed in [22], and a matrix-based method based on this assumption
679 was developed to compute the system reliability. Following that work, we
680 assume interdependence among the components' states, but only for com-
681 ponents with the same function. For group k of components with the same
682 function, let x_k denote a binary variable which indicates the occurrence of
683 an external event relevant for the group if $x_k = 1$, and it is $x_k = 0$ oth-
684 erwise. The Bernoulli probability of such variable is defined by probability
685 $\alpha_k = \mathbb{P}[x_k = 0]$. For any component c_i within that group, if $x_k = 1$, then the
686 component is surely functioning, i.e. $\mathbb{P}[s_i = 1|x_k = 1] = 1$; while if $x_k = 0$
687 component c_i fails with conditional probability $\beta_k = \mathbb{P}[s_i = 0|x_k = 0]$. Let ρ_k
688 be the correlation coefficient between the states of any pair of components

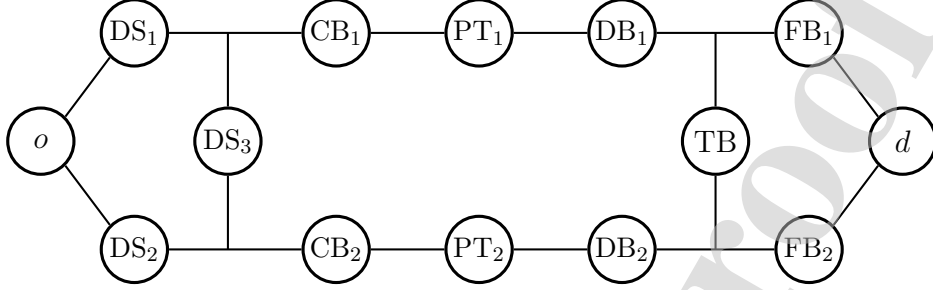


Figure 16: Block diagram of a two-transmission-line substation system

689 within group k , and \bar{p}_k the marginal failure probability for any component
 690 in the group. The corresponding factors are:

$$\begin{cases} \beta_k = \rho_k(1 - \bar{p}_k) + \bar{p}_k \\ \alpha_k = \bar{p}_k / \beta_k \end{cases} \quad (22)$$

691 The example of Fig 16 is defined by 6 pairs of features $\{\bar{p}_1, \rho_1, \dots, \bar{p}_6, \rho_6\}$
 692 corresponding to coefficients $\{\alpha_1, \beta_1, \dots, \alpha_6, \beta_6\}$.

693 When all the components are independent, the prior action is to do noth-
 694 ing, and the optimal posterior action is to repair the inspected component
 695 after an alarm, except for DS₃ and TB. Thus, the local metric and heuristic
 696 give identical results. Although CB and DB have relatively higher failure
 697 probability compared with other components, the cost reduction by repair-
 698 ing the damaged components CB or DB is significantly higher than repairing
 699 others. This is why CB and DB components have the highest VoI according
 700 to the local metric and the heuristic, as shown in Fig 17a. For the global
 701 metric, with loss function $l^*(p_u) = p_u(1 - p_u)$, the posterior system failure
 702 probability given an alarm from components CB or DB is the highest, and the
 703 probability given a silence from those components is the lowest, i.e. poste-

rior intervals $I_{CB} = I_{DB}$ contain the corresponding intervals of all the others; thus, those components have the highest VoI, according to the global metric.

When the correlation among states in DS components grows, while other groups remain independent (and the marginal probability remains the same), the VoI favors the group of correlated components. The prior action becomes repairing DS_1 or DS_2 when the correlation coefficient ρ is above 0.4. The optimal action is shown in Table 3. Components DS_1 or DS_2 should be kept functioning, depending on which link set the inspected component is in. One exception is DS_3 , which has different VoI for the local metric and the heuristic. As shown in Fig 17b, when the correlation coefficient ρ for the states of components DS increases, inspecting one of them reveals additional information about the other two, [making the VoI of inspecting DS components higher than the VoI of other components with different functions](#). When ρ is close to one, DS_1 and DS_2 act like one bottleneck component, which dominates the VoI as shown in Fig 17c.

6. Discussion and Conclusions

We have derived metrics based on the VoI to assign priorities among component inspections in networked systems. The VoI analysis can be applied to any setting, but its computational complexity depends on the complexity of the ingredients that define the problem. We have restricted the attention to binary components, binary inspection outcomes in a binary system. In this setting, we have introduced two metrics, local and global, that assume different sets of available actions and [different loss functions](#). The problems modeled by the global metric do not form a [strict subclass of that mod-](#)

Component	Insp. outcome	
	Silence ($y_i = 1$)	Alarm ($y_i = 0$)
DS ₁	\emptyset	DS ₁
DS ₂	\emptyset	DS ₂
DS ₃	\emptyset	DS ₃
CB ₁	DS ₁	DS ₁ , CB ₁
CB ₂	DS ₂	DS ₂ , CB ₂
PT ₁	DS ₁	DS ₁ , PT ₁
PT ₂	DS ₂	DS ₂ , PT ₂
DB ₁	DS ₁	DS ₁ , DB ₁
DB ₂	DS ₂	DS ₂ , DB ₂
TB	DS ₁	DS ₁
FB ₁	DS ₁	DS ₁ , FB ₁
FB ₂	DS ₂	DS ₂ , FB ₂

Table 3: Optimal posterior action for the system in Fig 16 when $\rho = 0.4$

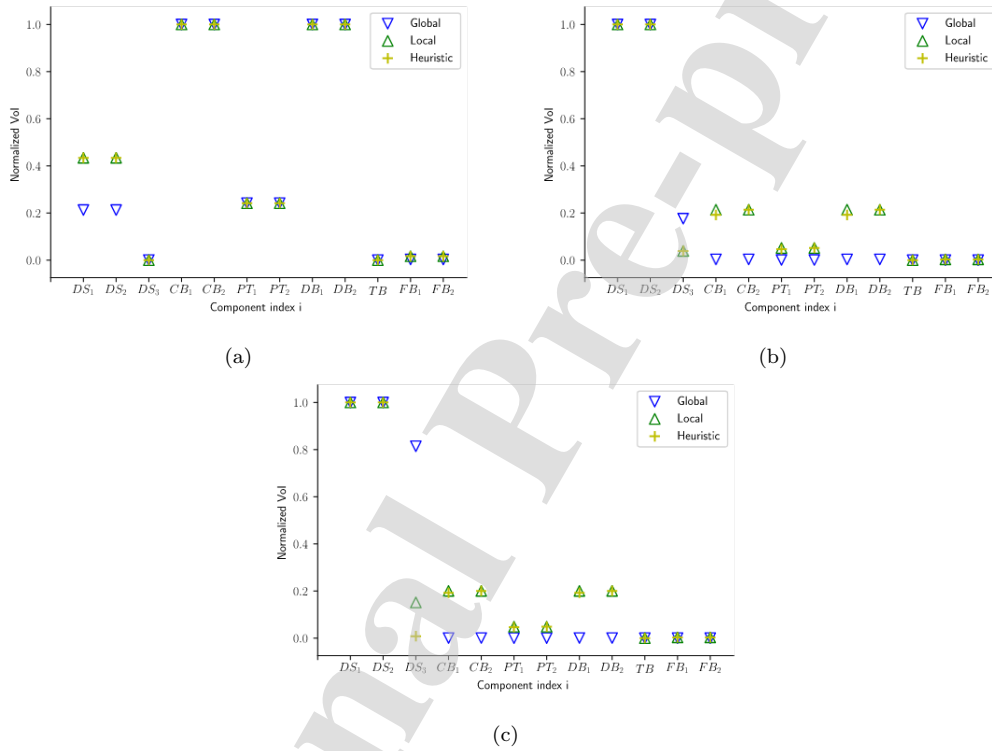


Figure 17: Normalized Vol for the system in Fig 16, with correlation among the DS component of $\rho = 0$ (a), $\rho = 0.4$ (b), $\rho = 0.9$ (c)

728 eled by the local metric. We have proven general rules for identifying what
729 components have higher importance for those metrics in series and parallel
730 systems. The evaluation of the global metric is generally less complex than
731 evaluation of the local metric, because of the underlying optimization of the
732 maintenance actions of different scales. The selection of the appropriate met-
733 ric should be based on the actual set of actions available. However, when
734 only limited computational resources are available, a simpler metric such as
735 the global metric or the heuristic approach may be appropriate.

736 We have proposed a heuristic approach to approximate the local metric,
737 by simplifying the corresponding optimization of maintenance actions. We
738 have illustrated the heuristic's performance in some examples, but there is
739 no guarantee that the heuristic captures the exact local metric. The VoI
740 assessed by the heuristic is surely non-negative, and no higher than the VoI
741 of the original local metric; however, the ranking can be arbitrarily different.

742 The distinction between local and global metrics can be extended to the
743 case of multiple values (more than binary) for the state of the components
744 and of the system, and for inspection outcomes. However, some concepts
745 are defined only for the binary case, e.g. the posterior intervals in the global
746 metric are defined only for binary inspection outcomes in a binary system (if
747 the system state dimension is higher than the component state dimension,
748 then the posterior interval can be generalized into the concept of "posterior
749 polyhedron").

750 We have limited the analysis to the "static" optimization of the inspection
751 of one component. Several more complex problems can be built on this
752 optimization. One complex problem is the off-line or on-line optimization of

753 multiple inspections for a system with static condition states [13], that can
 754 be based on a greedy sequential approximation. The same global and local
 755 approaches can be adopted in the greedy approach. Among these two options,
 756 the on-line setting is generally simpler; the off-line option is generally more
 757 expensive because M binary inspections produce 2^M joint outcomes that
 758 must be analyzed in an exhaustive pre-posterior analysis. Another extension
 759 is related to temporal problems, in which the components' condition degrades
 760 in time, and they can be periodically and sequentially inspected and repaired
 761 [15, 16]. Given the complex interplay between present and future decisions
 762 and costs, we cannot predict the effectiveness of the metrics proposed when
 763 applied to those dynamic settings.

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 768 ond author acknowledges the support of Visiting Faculty Fellows Program
 769 of the Wilton E. Scott Institute for Energy Innovation at Carnegie Mellon
 770 University.

771 Appendix A. Notation table

Notation	Meaning
N	Number of components
$c_i, i = 1, \dots, N$	Component index
$s_i \in \mathbb{B}, i = 1, \dots, N$	Binary component state

$s = [s_1, s_2, \dots, s_N]$	System state vector
$u \in \mathbb{B} = \phi(s)$	Binary system state
p_s	Probability distribution of s
$p_i, i = 1, \dots, N$	Marginal failure probability of component c_i
p_u	Failure probability of the system
$y_i \in \mathbb{B}, i = 1, \dots, N$	Inspection outcome on component c_i
p_{y_i}	Prior probability of receiving y_i
$p_{s y_i}$	Posterior probability distribution of s given y_i
$p_{\omega y_i}$	Posterior system failure probability given y_i
A	Maintenance action that changes the state s
$p_{s' s,A}$	Posterior probability distribution of s' given s and A
$\mathcal{L}(s', A)$	Expected loss given action A and posterior state s'
$\mathcal{L}_I(s')$	Expected system failure penalty
$\mathcal{L}_{II}(A)$	Expected action cost
L_π	Minimum expected cost before inspection
$L_\omega(i)$	Minimum expected cost after inspecting c_i
$\text{VoI}(i)$	Value of Information for inspecting c_i
$\epsilon_{FS}, \epsilon_{FA}$	False silence and alarm inspection error rate
h_i	Probability of receiving alarm on c_i
$l_{A,u}$	Cost given prior state u and action A
$q_{u,A}$	System failure probability given prior state u and action A
$p'_{\omega A,u}$	Posterior probability of system failure given prior state u and action A
C_A	Action cost
C_F	System failure cost

$l_A(p_u)$	Expected loss when taking action A given prior belief p_u
$L_\omega^G(i)$	VoI of inspecting c_i under global metric
$L_\omega^L(i)$	VoI of inspecting c_i under local metric
$l_{PI}(p_u)$	Value of Perfect Information (VoPI) given prior belief p_u
$rg(p_u)$	Difference between the VoI and VoPI

Table A.4: Major notations

772 Appendix B. Importance Measures

773 Similar to the Birnbaum's measure, the Criticality IM [7], evaluates the
 774 importance of c_i with the approximated conditional component failure prob-
 775 ability given that the system has failed:

$$\text{CRT}(i) = (p_{\omega|y_i=0} - p_{\omega|y_i=1}) \frac{p_i}{p_\pi} \propto \text{BM}(i) \cdot p_i \quad (\text{B.1})$$

776 Some IMs emphasize on the topology structure of the system. Based on
 777 the cut sets, [9] evaluates the importance of c_i by the number of cut sets it
 778 belongs to and the accumulated appearance probability of such cut sets.

779 To use IMs as utility-based applications, the risk achievement worth
 780 (RAW) and the risk reduction worth (RRW) are developed. RAW evalu-
 781 ates the component with the contributions of maintaining a certain level of
 782 reliability of the component to the system reliability, i.e. for component c_j ,
 783 its importance can be measured as:

$$\text{RAW}(i) = \frac{1 - p_{\omega|y_i=1}}{p_\pi} \quad (\text{B.2})$$

784 So, between two components c_i and c_j , $\text{RAW}(i) \geq \text{RAW}(j) \Leftrightarrow p_{\omega|y_i=1} \leq$
 785 $p_{\omega|y_j=1}$. RRW evaluates a component by the decrease of system failure risk

786 given that the component is intact:

$$\text{RRW}(j) = \frac{p_\pi}{1 - p_{\omega|y_i=0}} \quad (\text{B.3})$$

787 So $\text{RRW}(i) \geq \text{RRW}(j) \Leftrightarrow p_{\omega|y_i=0} \geq p_{\omega|y_j=0}$.

788 Appendix C. Nested posterior intervals in the global metric

789 To prove the lemma in Section 3.1, we now write $p_{\omega|y_a=b}$ as $x_{a,b}$ for simplic-
 790 ity. We assume that $I_i \supseteq I_j$, we have that $0 \leq x_{i,1} \leq x_{j,1} \leq x_{j,0} \leq x_{i,0} \leq 1$.
 791 Because of the law of expectation, we have:

$$p_\pi = p_1 x_{i,1} + (1 - p_1) x_{i,0} = p_2 x_{j,1} + (1 - p_2) x_{j,0} \quad (\text{C.1})$$

792 We prove that:

$$L_\omega^G(1) = p_1 l(x_{i,1}) + (1 - p_1) l(x_{i,0}) \leq p_2 l(x_{j,1}) + (1 - p_2) l(x_{j,0}) = L_\omega^G(2) \quad (\text{C.2})$$

793 Because $x_{j,1} = \frac{x_{i,0} - x_{j,1}}{x_{i,0} - x_{i,1}} x_{i,1} + \frac{x_{j,1} - x_{i,1}}{x_{i,0} - x_{i,1}} x_{i,0}$ and $x_{j,0} = \frac{x_{i,0} - x_{j,0}}{x_{i,0} - x_{i,1}} x_{i,1} + \frac{x_{j,0} - x_{j,1}}{x_{i,0} - x_{i,1}} x_{i,0}$,
 794 and l is a concave function, we have:

$$\begin{aligned} p_2 l(x_{j,1}) + (1 - p_2) l(x_{j,0}) &\geq p_2 \left[\frac{x_{i,0} - x_{j,1}}{x_{i,0} - x_{i,1}} l(x_{i,1}) + \frac{x_{j,1} - x_{i,1}}{x_{i,0} - x_{i,1}} l(x_{i,0}) \right] \\ &\quad + (1 - p_2) \left[\frac{x_{i,0} - x_{j,0}}{x_{i,0} - x_{i,1}} l(x_{i,1}) + \frac{x_{j,0} - x_{j,1}}{x_{i,0} - x_{i,1}} l(x_{i,0}) \right] \\ &= p_1 l(x_{i,1}) + (1 - p_1) l(x_{i,0}) \end{aligned} \quad (\text{C.3})$$

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Authorship Conformation Form

For the manuscript “Optimal inspection of network systems via Value of Information Analysis”, by Chaochao Lin, Junho Song and Matteo Pozzi, submitted to RESS, the authors agree that:

- All authors have participated in (a) conception and design, or analysis and interpretation of the data; (b) drafting the article or revising it critically for important intellectual content; and (c) approval of the final version.

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Conflict of Interest

For the manuscript “Optimal inspection of network systems via Value of Information Analysis”, by Chaochao Lin, Junho Song and Matteo Pozzi, submitted to RESS, the authors agree that:

- The authors have no affiliation with any organization with a direct or indirect financial interest in the subject matter discussed in the manuscript.

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