AutoGR: Automated Geo-Replication with Fast System Performance and Preserved Application Semantics

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ABSTRACT

Geo-replication is essential for providing low latency response and quality Internet services. However, designing fast and correct geo-replicated services is challenging due to the complex trade-off between performance and consistency semantics in optimizing the expensive cross-site coordination. State-of-the-art solutions rely on programmers to derive sufficient application-specific invariants and code specifications, which is both time-consuming and error-prone. In this paper, we propose an end-to-end geo-replication deployment framework AutoGR (AUTOmated Geo-Replication) to free programmers from such label-intensive tasks. AutoGR enables the geo-replication features for non-replicated, serializable applications in an automated way with optimized performance and correct application semantics. Driven by a novel static analyzer Rigi, AutoGR can extract application invariants by verifying whether their geo-replicated versions obey the serializable semantics of the non-replicated application. Rigi takes application codes as inputs and infers a set of side effects and path conditions possibly leading to consistency violations. Rigi employs the Z3 theorem prover to identify pairs of conflicting side effects and feed them to a geo-replication framework for automated across-site deployment. We evaluate AutoGR by transforming four serializable and originally non-replicated DB-compliant applications to geo-replicated ones across 3 sites. Compared with state-of-the-art human-intervention-free automated approaches (e.g., strong consistency), AutoGR reduces up to 61.8% latency and achieves up to 2.12× higher peak throughput. Compared with state-of-the-art approaches relying on a manual analysis (e.g., PoR), AutoGR can quickly enable the geo-replication feature with zero human intervention while offering similarly low latency and high throughput.

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1 INTRODUCTION

To offer fast responses for users worldwide [2, 5, 47], many Internet services replicate data across multiple geographically dispersed data centers [23, 24, 26, 40, 50]. However, geo-replicated services face an expensive wide-area communication problem, since concurrent user requests across sites must be coordinated for precluding incorrect behaviors, e.g., states may permanently diverge between sites, and/or application-specific invariants may be violated.

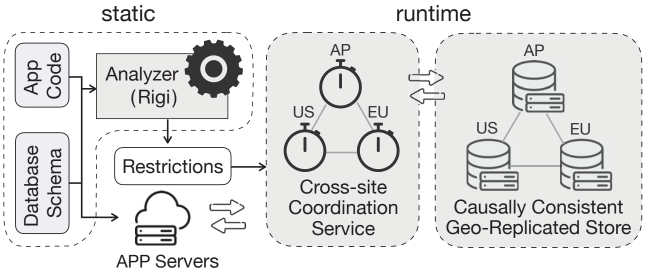

Figure 1: An overview of the proposed end-to-end AutoGR solution. AP, US, and EU stand for data centers in Singapore, Oregon, and Frankfurt, respectively. The runtime library is co-located with Server, omitted from the graph.

Various fine-grained consistency models have been proposed for safely avoiding unnecessary coordination [30, 36, 42, 56]. PoR consistency [36] and Generic Broadcast [42] express the consistency semantics as the *ordering restrictions* over pairs of operations so that weakening or strengthening the consistency semantics can be achieved by imposing fewer or more restrictions. Tools like Indigo [20] and SIEVE [34] can assist developers in identifying necessary coordination, though it requires programmers to write application-specific invariants and code specifications [18]. Such a manual process is often very time-consuming and error-prone.

In this paper, we present an end-to-end framework named AutoGR (Automated Geo-Replication), which automatically deploys geo-replicated applications based on their non-replicated serializable versions without manually dealing with restriction-based fine-grained consistency. AutoGR requires *zero programmer efforts* to write any form of specifications while *minimizing* cross-site coordination costs and *preserving* correct application semantics.

Our work is based on the key observation that the side effects produced by user requests and the application-specific invariants that programmers need to maintain are already implicitly reflected in the logic of the programs. Therefore, it is possible to automatically infer through a static program analysis of the code behaviors and the invariants to detect potential conflicting operations that need to be coordinated. For instance, to implement the *withdraw* operation of a banking system, we must subtract from the current balance a *delta*, which is specified by users. This subtraction summarizes the intended code behaviors. To make this subtraction happen, we must check the code to make sure the amount to be withdrawn should be

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less than or equal to the current balance. This checking reflects the intended invariant that the balance must always be non-negative.

Following this insight, we propose, as one of our key contributions – a novel static analyzer called Rigi. As shown in Figure 1, Rigi forms the static part of the AutoGR framework. In this paper, we focus on the DB-compliant applications that implement their logics in high-level language such as Java and store their states in transactional databases. Rigi takes the non-replicated, serializable application code and the corresponding database schema as inputs. Without requiring any form of user specifications, it traverses the corresponding control flow graph and extracts the side effects of each operation over the shared data and the corresponding path conditions that need to be met before the effects can be made to ensure correct application semantics. With the help of our formal database abstraction, we make those side effects and conditions verifiable by further translating them into the code of Z3, a popular and powerful theorem prover [27].

Rigi then employs Z3 to automatically check the commutativity of the side effects of every pair of operations, and the compatibility of their corresponding path conditions. If pairs of conflicting side effects are detected, Rigi generates ordering restrictions as its outputs. The ordering restrictions are enforced by the runtime of AutoGR to ensure the geo-replicated execution obeys the serializable semantics of the non-replicated version of the application [41].

AutoGR integrates Rigi with an existing geo-replication and coordination system Olisipo [7, 36], see Figure 1. The runtime of AutoGR automatically deploys the geo-replicated application through carefully coordinating the operations confined by the restrictions generated by Rigi between sites. This guarantees that the generation of their conflicting side effects are serialized, and these effects are propagated and replicated across all sites according to the determined proper orders.

To demonstrate the power of AutoGR, we transform four non-replicated applications including *SmallBank* [53], *RUBiS* [29], *Seats reservation* [54] , and *HealthPlus* [6] into their geo-replicated versions. Our static analysis results suggest that the automatic reasoning cost is low and it is able to find the same minimal set of restrictions as state-of-the-art solutions that need extensive efforts from the programmer (e.g., PoR consistency [36]). The experimental results with a 3-site geo-replication setting highlight that compared with state-of-the-art human-intervention-free approaches (e.g., strong consistency), AutoGR reduces up to 61.8% latency and achieves up to 2.12× higher peak throughput; compared with manual analysis using state-of-the-art approaches (e.g., PoR), AutoGR can quickly enable geo-replication with zero human intervention while offering similarly low latency and high peak throughput.

We summarize our main contributions below:

- To the best of our knowledge, AutoGR is the first end-to-end geo-replication framework that can automatically transform non-replicated, serializable applications into their fast geo-replicated version with little user interventions and without compromising important system properties such as state convergence and invariant preservation.
- We offer an analyzer tool Rigi for performing static analysis over source codes written in Java with manipulating data via SQL interface. Rigi suggests a minimal set of ordering restrictions over pairs of conflicting operations, requiring zero efforts from developers. The core of Rigi consists of the database abstraction in Z3, covering a substantial range of SQL features, and the general commutativity and semantics checks.
- We evaluate AutoGR on three widely used benchmark applications and a real-world large web application to demonstrate AutoGR's generality and practicability for a wide range of applications, and its strengths over existing works in automatically enabling the geo-replication feature with superior performance on application latency and system throughput.

2 PRELIMINARIES

2.1 System model

We assume a distributed and replicated system consisting of multiple replicas, spanning over geographically dispersed data centers (a.k.a. *sites*). User requests (a.k.a. *operations*) are first submitted to a nearby replica (denoted as *primary*) and are executed on that replica against its local state S. As the first step, this execution only performs condition checks (taking banking as an example, a *withdraw* request checks whether the amount of money requested is greater than 0 and the *balance* is sufficient) and identifies the corresponding side effect, but does not commit it. We model a side effect as a deterministic function, which takes some arguments as the input and transition the system from one state to another. Upon completion of the primary execution, the collected side effect is then propagated and replicated to all replicas, including the primary one. Note that the state at a site that a side effect is applied to might differ from S observed when the corresponding side effect was created. Following this, all sites together establish a global partial order on the generation of side effects. Each site applies those side effects in a total order, which may differ from site to site but must be compatible with the global partial order.

2.2 Restriction-based consistency model

Data copies at different replicas must be synchronized for maintaining "consistency." Traditionally, consistency means *serializability* [21], i.e., the results obtained by concurrently executing a set of operations/transactions should be equivalent to one of some serial orders. However, the synchronization cost across wide-area nodes is extremely high; moreover, a few industrial studies point out even a slight increase in user experienced latency can lead to significant revenue loss [2, 5, 47]. To offer *low latency responses*, several fine-grained consistency models have been proposed towards minimizing the expensive cross-site coordination in geo-replicating services [17, 20, 39, 45, 46, 52]. More recently, PoR consistency [36], expresses consistency semantics as a set of ordering restrictions *(for short, restrictions)* over pairs of operations. Formally, for any two operations u and v, if there exists a restriction over them, then u and v must be serialized so that their side effects are produced and executed on all sites following the same order. For example, in a banking system, to avoid over-withdrawal, one has to impose a restriction over any pair of *withdraws*.

The key ideas behind the above proposals are to make the behaviors delivered by the concurrent geo-replicated executions be explained by some serial executions. In particular, they guarantee that the following two crucial system properties are maintained under relaxed consistency models. Informally, the first property is called *state convergence*, i.e., when beginning at the same initial state, after applying the same set of side effects (possibly in a different order), all sites reach the same final state. Besides, the second property is called *invariant preservation*, which says that at each site, every state transition triggered by applying a side effect should not violate application-specific invariants defined by developers, assuming the initial state preserving those invariants (We call invariant-preserving state *valid*).

For geo-replicating an application, the key to making the best use of the restriction-based fine-grained consistency models is to identify all pairs of conflicting operations, each of which will be confined by a restriction. AutoGR is built atop of this line of work.

2.3 Tool support and limitations

Some works have formulated a set of principles to guide developers to identify restrictions [17, 20, 32, 36]. However, most of them require extensive domain expertise and manual efforts to reason about all possible concurrent executions of user requests. Such a manual reasoning process is time-consuming and error-prone.

To relieve the burden imposed on programmers, many attempts have been made towards offering automated tools for completing this task [20, 34]. SIEVE computes the weakest preconditions for each side effect, which summarizes when they cannot be executed without coordination [34]. This solution is conservative, as it requires coordinating the generation and replication of all problematic side effects, whose weakest preconditions are evaluated to `False`. To improve the limitation of SIEVE, Indigo performs static analysis of operation post-conditions against invariants to infer the pairs of concurrent operations that may lead to invariant violations [20]. One of the major drawbacks of this line of work is that they all require programmers to write correct and sufficient specifications about application-specific invariants and pre- or post-conditions of their code in logic formulas. Another major drawback is that they assume that the specifications written by programmers are always correct and sufficient. Incomplete or too weak invariants can cause incorrect system behaviors, while too strong or unnecessary invariants lead to performance penalty. As a consequence, it's desired to completely release this burden from programmers by not requiring them to specify application-specific invariants.

2.4 The Z3 theorem prover

Z3 is a state-of-the-art theorem prover developed by Microsoft Research [27], and has been widely used to check the satisfiability of logical formulas [37, 49]. Those formulas are input constraints given by users and can be written as a Boolean combination of atomic formulas, which are defined over a rich set of theories supported by Z3, such as integer and real arithmetic, bit-vectors, arrays, quantifiers, and functions. Given a formula in the first-order logic calculus with free variables, Z3 searches a set of assignments to its variables, which satisfy that formula, and reports `unsat`, if the search fails. We can take advantage of the rich set of theories in Z3 to model the changes to the shared state of target applications, while using the solver to reason about the consistency constraints encoded in the original code.

```
void withdraw(Connection conn, String custName,
                      double amount) throws Exception {
    PreparedStatement stmt = conn.prepareStatement(
                  "SELECT * FROM ACCOUNTS WHERE name = ?");
    stmt.setString(1, custName);
    ResultSet rs = stmt.executeQuery();
    if (rs.next() == false) throw new Exception("Invalid account");
    long custId = rs.getLong(1);
    stmt = conn.prepareStatement("
           SELECT bal FROM SAVINGS WHERE custid = ?");
    stmt.setLong(1, custId);
    rs = stmt.executeQuery();
    if (rs.next() == false) throw new Exception("No saving account");
    double balance = rs.getDouble(1) − amount;
    if (balance < 0) throw new Exception("Negative balance");
    stmt = conn.prepareStatement(
                  "UPDATE SAVINGS SET bal =? WHERE custid =?");
    stmt.setDouble(1, balance); stmt.setLong(2, custId);
    stmt.executeUpdate();
    conn.commit();
}
```

Figure 2: Pseudocode of a *Withdraw* transaction

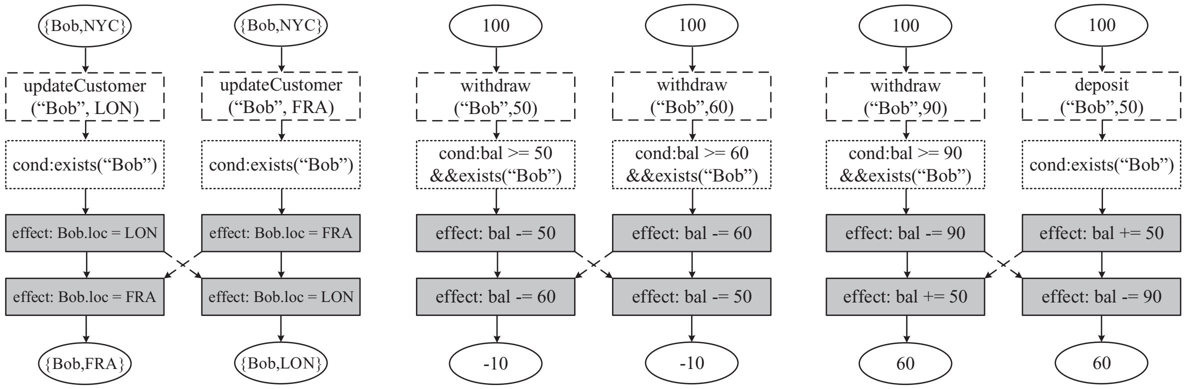
3 OVERVIEW

The goal of AutoGR is to develop an end-to-end framework to automatically deploy non-replicated serializable code in a geo-replicated setting to explore performance benefits while maintaining the correct application semantics. The performance gains enabled by AutoGR come from two parts. First, AutoGR leverages on an existing geo-replication framework Olisipo that enables fine-grained coordination over pairs of operations that produce conflicting side effects. Second, AutoGR integrates Olisipo with a static analyzer Rigi to identify a minimal set of ordering restrictions that must be ensured so that the intended semantics of the corresponding geo-replication applications are not violated.

3.1 Target applications

We target and transition classic three-tier non-replicated applications [1, 10]. One of such applications spans its logic across client, application server, and database tiers, and assumes a single-copy database (often a relational database). The application server contains the functional business logic, defined as a set of transactions. Each transaction takes requests from clients as inputs, and intellectually manipulates the state stored in the database via SQL queries. These queries are statically defined as parameterized functions, and instantiated by user inputs at runtime.

Here, we briefly elaborate on how the applications work. When a transaction starts, it takes inputs from users (e.g., the key to the target objects), and executes a `SELECT` query to retrieve shared objects from the backend relational database. Then, it performs certain condition checks to determine the side effects. For instance, in Figure 2, the *Withdraw* transaction checks the existence of the target account and the balance in that account fetched from the database is greater than or equal to the amount supplied by users. Depending on the checking results, the execution of that transaction can go to different branches, which leads to different side effects. For instance, if the check fails, then the *Withdraw* transaction generates no side effects. Otherwise, it may perform changes to the shared objects and write those side effects back to the database by executing



(a) Violating execution 1

(b) Violating execution 2

(c) Correct execution.

Figure 3: Three possible executions of geo-replicated services without coordination. Ellipses represent states. Dashed boxes represent user requests, while the dotted boxes and gray boxes indicate their path conditions and side effects, respectively. The serial order established by solid arrows corresponds to the execution of generating and applying side effects at a site. The dashed arrows indicate the propagation and replication of side effects from a site to another.

update queries such as `INSERT`, `UPDATE`, and `DELETE`. In the end, it commits all side effects and makes the changes persistent.

3.2 A motivating example

Next, we use a simplified banking system as a motivating example to explain the design principles behind Rigi for identifying conflicting operations and generating ordering restrictions. The system has three operations, namely, *Withdraw*, *Deposit*, and *UpdateCustomer*. The pseudocode of *Withdraw* is shown in Figure 2, and we omit the other two in the interest of space. Figure 3 gives three concurrent executions of these operations with no coordination placed.

Figure 3a illustrates an example of replicating non-commutative side effects. Two concurrent requests of the *UpdateCustomer* operation are submitted to two data centers to change the location of “Bob” from “NYC” to “LON” and “FRA,” respectively. When the side effects of the two requests were replicated across the two sites in different orders, the resulting states would differ. This violates a critical system property called *state convergence* [43]. To preclude this kind of anomalies, we need to check the commutativity of side effects produced by operations in a pair-wise fashion and add a restriction to coordinate the generation of any pairs of side effects that do not commute.

Figure 3b shows an example using the side effects of two concurrent *Withdraw* requests. Although the two side effects commute, this execution is still invalid. This is because the path conditions established at the primary replica could be invalidated by the concurrent execution of other operations. In this example, the path condition for the side effect of a successful withdrawal would be $balance - money \geq 0$. Clearly, the condition holds in the primary replica of the two simultaneous *Withdraw* requests. Nevertheless, it becomes false when the side effect reaches the other (remote) site. Application of the side effects of both requests results in a negative balance value, violating the expected invariant. To preclude this type of violations, we need to derive a semantics check to determine if we should add a restriction between the pair of side effects that are not invariant preserving.

Unlike the above two harmful concurrent executions, Figure 3c represents a correct one, where a *Withdraw* can run concurrently with a *Deposit*, for which the reason is two-fold. First, their side effects are commutative. Second, the conditions corresponding to the two identified side effects are compatible, i.e., the side effect of *Deposits* will not invalidate the condition of successful *Withdraws*, and vice versa. This indicates that no coordination should be placed to restrict the generation and replication of the two side effects.

3.3 Challenges and opportunities

Inspired by the examples shown in Figure 3, the correctness requirement implies that the geo-replicated version should not generate more states than the original, non-replicated version. Therefore, we may need to output a restriction over a pair of side effects if they meet either of the following two conditions: a) they are non-commutative side effects and may cause replicas to terminate in a divergent state that cannot be generated in any serializable order; and b) the condition leading to one of the two side effects might no longer hold when seeing the other side effect.

To detect the two cases without the programmer’s inputs, we need to “understand” the operations solely from the source code to infer side effects and invariants. First, due to the code structure of the target applications, for inferring side effects, it is sufficient to understand the specifications of SQL queries and the changes induced on arguments that instantiate the queries. Second, application-specific invariants are implicitly encoded into the conditional checks taken at each branch along a control flow path of each transaction. As a result, to automatically infer the invariants, we need to extract and understand the semantics of the path condition leading to a particular side effect. This step requires traversing all distinct code paths. It may also involve the understanding of the semantics of `SELECT` queries, since the variables in the branch conditional check formula may be fetched from the database.

Identifying conflicting pairs of side effects requires us to make the gathered side effects and path conditions verifiable. Here, we translate them into the code of Z3 [27], and then ask Z3 to check the commutativity of the side effects and the compatibility of their

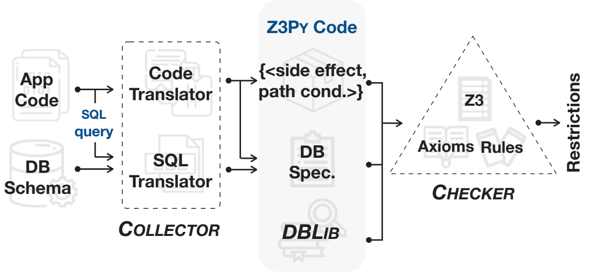


Figure 4: The design of the Rigi analyzer.

corresponding conditions by proving certain formulas. The checking provably identifies a minimal set of non-commutative or non-invariant-compliant pairs of side effects as necessary ordering restrictions that need to be enforced by the runtime part of AutoGR.

4 Rigi: A STATIC ANALYZER

As illustrated in Figure 4, Rigi is comprised of a **database specification library** (*or short, DBLib*), a **side effect collector** (*or short, Collector*), and a **rule checker** (*or short, Checker*).

(1) DBLib represents the abstraction of relational databases in Z3Py, a Z3 API in Python. It covers a substantial range of SQL features and serves as a building block for Rigi.
(2) Collector inspects the source code and identifies the side effects generated in different code paths. It gathers the path conditions leading to particular side effects. It also translates the side effects and their path conditions into code in Z3Py.
(3) Checker calls the Z3 theorem prover [27] and generates restrictions through commutativity and semantics checks over pairs of operations' side effects and their path conditions.

4.1 Database abstraction in Z3

Table 1 shows DBLib, our database abstraction in Z3, covers a substantial range of SQL features. The basic features cover the table schema definition and simple queries such as `SELECT`, `UPDATE`, etc. The first row of Table 1 shows the translation of a table definition in SQL into an extensional array in Z3Py, which consists of pairs of < `KEY`, `VALUE` >, where `KEY` represents the unique index of the corresponding table that may include one or more attributes, and `VALUE` covers the remaining attributes. Both `KEY` and `VALUE` are defined as an algebraic data type in Z3Py where each of their attributes is declared as a primitive type such as `IntSort`, `StringSort`, `BoolSort`, and `RealSort` (We also enable to use `FPSort` instead of `RealSort` for stricter floating-point verification). For instance, the primary key of the table `RSVN` is translated into a customized data structure of type `K_RSVN` to include two attributes, namely `R_C_ID` and `R_F_ID`. Considering that the data for each row of the table can be fetched by a primary key or a unique index key, we additionally allow both of them to point to the same row of the table.

The second four lines of Table 1 represent the translation of basic SQL queries. First, the `SELECT` and `UPDATE` queries are translated into `Select` and `Store` operations on Z3Py arrays, parameterized by the primary key. We additionally allow specifying which set of attributes of the target record a query intends to fetch or update. Regarding `INSERT`, we create a key value pair, and make the corresponding key point to the value through the `Store` function. For `DELETE`, we just set the target record to be `nil`.

The last three lines of Table 1 show how we model the advanced SQL features, ranging from comparison operator to aggregation. First, for the comparison operator, we fetch the target attribute by the corresponding key and compare it with the supplied argument, e.g., `end_date` in the example. Second, we translate cross-table join queries into multiple sub-selections, each of which fetches records of interest from a particular table with the same unique key. Third, for `ORDER BY` on an attribute, we specify a relation over the list of records that for any adjacent pairs of records, the value of the target attribute of the former record is always either greater or less than the counterpart of the latter record.

4.2 Collecting side effects and conditions

The mission of Collector is to collect side effects and their corresponding path conditions, and to translate them into Z3Py code, on a per path basis. It first builds an Abstract Syntax Tree (AST) from the source code of the target application to obtain the control flow graph. Then, it traverses the entire control flow graph to identify paths via a depth-first-search algorithm. The path analysis faces an infinite path problem, in the presence of loops. To address this challenge, we take inspiration from an existing work [34] that many real-world applications (including those we analyze) satisfy a loop iteration independence property, which says that the parts of the state modified in each iteration are disjoint. In this case, we handle loops by first verifying this independence property and then unfolding the independent loops so that the number of paths becomes finite. We explain the handling of complex loops in Section 4.3.

We skip paths that are either read-only or leading to *abort*, since they do not generate side effects. Then, we iterate all remaining code paths, and collect the side effect for a path by performing one-to-one direct translation from Java code to Z3Py code with the following two exceptions. First, we ignore control flow condition checks along the target path, which will be handled by the condition collecting phase. Second, for every encountered SQL statement, we translate it into its corresponding form using our extended DBLib APIs, presented in Table 1. It is worth mentioning that this one-to-one translation is feasible since Z3 itself already supports common data types and the associated computations, while our DBLib covers the database statements. Take `Withdraw` for an example. To update the account balance, the transaction assigns a local variable *balance* by the value fetched from the backend database. Then it performs a local computation to subtract *amount* from *balance*, where *amount* is supplied by the user. At the end, the transaction sets the final value of *balance* to instantiate a `UPDATE` query "update set bal = ? where custId = ?". After the translation, the resulting Z3Py code preserves the semantics of the original code. This translation enables our Z3-based Checker to infer the numeric changes on the *balance* variable, i.e., *decrement* here, to extract commutative side effects.

In addition to side effects, Rigi gathers the path condition by making conjunction of the Boolean formulas of every branch taken by the corresponding path. The individual formula may involve both shared objects fetched from the database or arguments from user inputs, which requires Rigi to substitute the variables involved in each formula with their source. For example, in the banking system, a formula "*balance* − *amount* >= 0" determining whether a *withdraw* can be admitted, is translated into "table_accounts(userId).bal

Table 1: Specifications in Z3Py automatically derived for database APIs which cover a substantial range of SQL features.

Category	Description	Example SQL statement	Z3Py specification

Basic features	Data model	CREATE TABLE RSVN (
		R_C_ID BIGINT NOT NULL,
		R_F_ID BIGINT NOT NULL,
		R_SEAT BIGINT NOT NULL, ...
		PRIMARY KEY (R_C_ID,R_F_ID))
			K_RSVN = Datatype(...)
			K_RSVN.declare(...,(R_C_ID, R_F_ID))
			V_RSVN = Datatype(...)
			V_RSVN.declare(...)
			TABLE_RSVN = Array(..., K_RSVN, V_RSVN)

	Basic queries	SELECT NAME FROM ACCOUNTS
		WHERE CUSTID = cid
			V_ACCOUNTS.NAME(Select(
			TABLE_ACCOUNTS, K_ACCOUNTS.new(cid)))

		INSERT INTO USERS
		(ID,FIRSTNAME,LASTNAME,...)
		VALUES (uid,fname,lname,...)
			Store(TABLE_USERS, K_USERS.new(uid),
			V_USERS.new(uid, fname, lname, ...))

		UPDATE CUSTOMER
		SET C_BASE_AP_ID = aid,
		LOCATION = loc
		WHERE C_ID = cid
			Store(TABLE_CUSTMER, T_CUSTMER.new(cid),
			V_CUSTMER.new(V_CUSTMER.BALANCE(
			Select(TABLE_CUSTMER,T_CUSTMER.new(cid)))
			,aid,...,loc,...))

		DELETE FROM ACCOUNTS
		WHERE CUSTID = cid
			Store(TABLE_ACCOUNTS,
			K_ACCOUNTS.new(cid), V_ACCOUNTS.nil)

Advanced features	Comparison operator	SELECT ID FROM ITEMS
		WHERE ID=id AND END_DATE<=end_date
			V_ITEMS.END_DATE(Select(
			TABLE_ITEMS,K_ITEMS.new(id))) <= end_date

	Inner join	SELECT BALANCE,F_SEATS_LEFT,
		R_ID, R_SEAT, R_PRICE
		FROM T_CUSTOMER,T_FLIGHT,T_RSVN
		WHERE C_ID=cid AND C_ID=R_C_ID
		AND F_ID=fid AND F_ID=R_F_ID
			V_CUSTOMER.BALANCE(Select(T_CUSTOMER,K_CUSTOMER.new(cid)))
			V_FLIGHT.F_SEATS_LEFT(Select(T_FLIGHT,K_FLIGHT.new(fid)))
			V_RSVN.R_ID(Select(T_RSVN,K_RSVN.new(cid,fid)))
			V_RSVN.R_SEAT(Select(T_RSVN,K_RSVN.new(cid,fid)))
			V_RSVN.R_PRICE(Select(T_RSVN,K_RSVN.new(cid,fid)))

	Aggregation	SELECT USER_ID, BID
		FROM BIDS
		WHERE ITEM_ID = iid
		ORDER BY bid
		DESC LIMIT 1
			ForAll([k_bids_1,k_bids_2], Implies(And(V_BIDS.ITEM_ID(
			Select(TABLE_BIDS,k_bids_1)) == iid, V_BIDS.ITEM_ID(
			Select(TABLE_BIDS,k_bids_2)) == iid, V_BIDS.USER_ID(
			Select(TABLE_BIDS,k_bids_2)) == winner),V_BIDS.BID(
			Select(TABLE_BIDS,k_bids_1)) <= V_BIDS.BID(Select(
			TABLE_BIDS,k_bids_2))))

- amount >= 0", where *userId* is a user's customer id and *bal* represents the balance attribute she wants to retrieve. $rs.next() == false$ is a commonly used conditional check to decide whether a particular record exists or not. In this case, Rigi first determines which SELECT query *rs* refers to, and then generates a Z3Py formula "table_[name](pkey) == nil", where *name* and *pkey* stand for the corresponding table name and primary key in that query.

In the end, the output of Collector is a set of tuples, each of which consists of *o*, *pcond*, and *arg*, standing for the side effect, path condition, and arguments, respectively. Examples are in Table 2.

4.3 Checking rules

As depicted in Figure 4, Checker takes the database specifications and the $\langle o, pcond, arg\rangle$ tuples generated by Collector as input, and performs two checks, namely the commutativity and semantics check, to figure out the conflicting side effect pairs.

Commutativity check. For two operations u_1 and u_2, assume their side effects and arguments are $\langle o_1, arg_1\rangle$ and $\langle o_2, arg_2\rangle$, respectively. The commutativity check is relatively straightforward and well studied in the literature [25, 55], and concentrates only on side effects. Here, we check, for any reachable state *s* of a target application, whether the following assertion is satisfied:

$$o_1(o_2(s, arg_2), arg_1) = o_2(o_1(s, arg_1), arg_2)$$

This assertion says that two side effects are commutative if applying them against the same initial state in different orders leads to the same final state. Take the problematic concurrent execution shown in Figure 3a, for example. The side effects of *UpdateCustomer* are not self-commutative; thus, Rigi adds a restriction over any pair of *UpdateCustomer* operations.

Additionally, for a transaction containing complex loops with non-independent iterations, it is hard for Rigi to infer their side effects. In this case, we resort to a conservative but safe strategy, where the corresponding transaction will be marked as conflicting with all other side effects that have write-write conflicts with it.

Semantics check. To ensure the transitioning from a non-replicated serializable code into its geo-replicated version is correct, we perform semantics check by asserting for any concurrent execution of two operations u_1 and u_2, the application of their side effects against a valid state in a serial order will not lead to an unreachable state in the non-replicated version. Formally, let o_1 be the side effect of u_1, o_2 and o_2' be the side effect of u_2 generated when u_2 misses or sees o_1, we should assert that the two operations u_1 and u_2 can be executed without restriction if (1) o_2 is a no-op side effect (i.e., missing the side effect of u_1, u_2 produces no side effect); or (2) $o_2 = o_2'$ (i.e., u_2 produces the same side effect regardless of the existence of u_1). Finally, we translate the above principle into checking the following formula:

$$\textit{For any reachable state } s, (s, arg_2) \textit{ satisfying } pcond_2 \Rightarrow (o_1(s, arg_1), arg_2) \textit{ satisfies } pcond_2.$$

We use two following counter-examples to illustrate this idea. First, as shown in Figure 3b, imagine we run two withdraws simultaneously against an initial state (balance is 100) where they want to deduct 50 (denoted as u_1) and 60 (denoted as u_2) from the shared account, respectively. If u_1 sees the side effect of u_2 (the remaining balance is 40), then the side effect of u_1 will never be generated. In this case, without ordering the execution of these two operations may lead to violations in the intended application semantics that balance should never be negative. Thus, Rigi automatically adds a restriction over u_1 and u_2.

Take *RUBiS* as another counter-example, *placebid* is to bid on an item when the corresponding auction is still open, while *closeAuction* terminates an auction and declares the winner. When *placebid* did not see the side effects of *closeAuction*, it would continue to bid for an auction that was already closed. As *closeAuction* may miss the side effects of *placebid* which issues a higher bid than those observed by *closeAuction*, the winner declaration would violate an intended application semantic where the winner must be the bidder associated with the highest accepted bit. Following a similar logic above, Rigi produces a restriction over *placebid* and *closeAuction*.

Furthermore, a transaction containing complex loops with non-independent iterations is marked conflicting with all other side effects whose generation phases need to read objects written by it.

In summary, Rigi faithfully follows the program logic of the non-replicated serializable code and may create false positives if the side effect analysis becomes conservative and less accurate, due to complex loops. False positives will negatively impact performance, but they will not cause state divergence or invariant violations. In contrast, it produces no false negatives since we assume that the non-replicated code to be analyzed must be serializable.

4.4 Correctness

To ensure our proposal is correct, we formulate the correctness into the following two theorems, and provide only the proof sketches, in the interest of space.

Definition 1 (Idleness). A k-site geo-replicated system is idle, if k sites begin at the same valid initial state, and after operating for a while, all side effects ever generated are applied at all k sites while no new side effects are being generated at the same time.

Theorem 2 (State convergence). *For a k-site geo-replicated system beginning at the same valid initial state, if the generation and application of side effects is coordinated by the restrictions identified by Rigi and executed by AutoGR, then when the system is idle, the reached final states of all k sites converge.*

Proof. Assume n side effects are generated in the target geo-replicated system, denoted by $o_1, o_2, \ldots, o_n$. Let $\langle o_1, o_2, \ldots, o_n \rangle$ denotes the state obtained by applying these n side effects sequentially against the initial state. We show by induction on n that for all n, all traces of n side effects are either not admitted by the target due to the Rigi's static checks and AutoGR's runtime coordination, or will lead to the same final state when the system is idle.

(1) $n = 0, 1$. Trivially convergent.
(2) $n = 2$. There are two possible traces, (o_1, o_2) and (o_2, o_1). If $\langle o_1, o_2 \rangle \neq \langle o_2, o_1 \rangle$, then one of the two traces will be eliminated by Rigi as the two side effects are not commutative.
(3) For case $n + 1$, we assume one admissible trace is $(o_1, o_2, \ldots, o_n, o_{n+1})$. We then consider an arbitrary trace $(p_1, \ldots, p_l, o_{n+1}, p_r, \ldots, p_n)$, where $r = l + 1$ and the sequence (p_i) is an arbitrary permutation of the sequence (o_j) for $j \leq n$. If $l = n$, by the induction hypothesis, the proposition holds. If $l < n$, for $\langle p_1, \ldots, o_{n+1}, \ldots, p_n \rangle = \langle o_1, \ldots, o_{n+1} \rangle$ to hold, the only possibility is to move o_{n+1} to the end. And this condition is only satisfied when o_{n+1} commutes with all p_i for $i \geq r$, which can be shown by induction on $n - r$. □

Theorem 3 (Invariant preservation). *For a k-site geo-replicated system with a valid initial state, if the generation and application of side effects are coordinated by the restrictions identified by Rigi and executed by AutoGR, then at each site, every state before generating or after applying a side effect must be valid.*

Proof. Let I denote the application-specific invariants encoded in the non-replicated, serializable application. For a side effect o, we say a state s is o-compatible, if o can be generated against state s, i.e., s satisfies both I and the path condition leading to o. For a k-site geo-replicated version of the target non-replicated application, we need to reason if the states before (denoted by s_i) and after (denoted by s_i') applying an admitted side effect o_i are both *valid*. Without loss of generality, we perform the following analysis against a site A, which receives two side effects o_1 and o_2 from any site in the geo-replicated system (A itself included). There are three major cases to be considered.

(1) o_2 arrived at A after o_1, and $s_2 = s_1'$ (o_2 is right after o_1).
   (a) The generation of o_2 sees the side effect of o_1, and thus I holds at s_1, s_2, s_2' by the correctness of the non-replicated version.
   (b) The generation of o_2 misses the side effect of o_1. In this case, s_1' must be o_2-compatible, by the logic of Rigi's semantics checks.
(2) o_2 arrived at A after o_1, but there are several side effects $o_3, o_4, \ldots, o_i, \ldots, o_{n+2}$ between o_1 and before o_2. We prove this by induction on n.
   (a) The base case $n = 0$ is exactly the first case.
   (b) For case $n + 1$, by the induction hypothesis, s_{n+2}' must be o_2-compatible, and for each $2 < j < n+2$, s_j is o_j-compatible.
(3) o_1 arrived at A after o_2. We prove this by a similar logic in the above case (1) and (2).

By the case analysis above, we therefore show that I is preserved at every moment before and after applying a side effect. □

4.5 Optimizations

Here, we extend Rigi to incorporate the following advanced features for practical considerations.

CRDTs support. Conflict-free Replicated Data Types (CRDTs) [44, 48] design operations that commute by construction. Thus, applications can leverage CRDTs to increase the space of commutative operations/side effects, in order to relax strong consistency constraints imposed by potential state divergence in the presence of uncoordinated concurrent request executions. To cope with the increasing attention on CRDTs [14, 16], we adapt Rigi to work with CRDTs. This adaption only requires us to extend Collector to generate commutative side effects according to the target CRDT's intended rules for merging concurrent updates, while leaving the commutativity checks to remain the same. To demonstrate this extensibility, we use a simple "Last-Writer-Win (LWW)" strategy as an example, which resolves conflicting updates by determining an arbitration order over those updates and pick up the latest write to be the final value [33, 43, 48]. To support this strategy, Collector assigns a unique timestamp *ts* to the corresponding side effect by including *ts* in its argument list. Then an axiom is added to Checker saying that the side effect will be applied when *ts* is greater than the local observed timestamp.

Uniqueness. The non-replicated, serializable applications often make use of the database's `AUTOINCREMENTAL` feature to ensure that a unique number is generated automatically whenever a new record is inserted into the table. When migrating to geo-replication, restrictions must be placed between pairs of side effects, each of which needs to put a record into the same database table, to preclude possible commutativity and semantic violations. To avoid the costly cross-site coordination for the unique id assignment, popular web services choose to partition the key space and assign different ranges to different replicas [4]. To support this feature, we add another optional axiom to Checker saying that the *id* obtained in the corresponding side effects is unique, i.e., $\forall id', id \neq id'$.

Developers can enable the commutative transformation by turning on the $-DOpt$ option. Note that although these transformations introduce new behaviors, i.e., some serializable execution orders previously rejected by the corresponding non-replicated application are now accepted, theorems 2 and 3 are still valid.

4.6 AutoGR integration

The implementation of Rigi consists of 2.1 *k* lines of Java code, 3.4 *k* lines of Kotlin code, and 521 lines of Z3Py code, where the Z3Py code performs the commutativity and semantics checks shared across applications. Rigi uses a Java parser [9] and a SQL statement parser [8] to generate database specifications, analyze the code paths, and collect side effects plus path conditions, for applications that are written in Java and using the SQL interface for data access. Note that we specialize our implementation on Java web applications since Java is one of the most popular languages used for implementing back-end server applications in major Internet service providers [13]. We further take advantage of the fact that the application code is often straightforward with few advanced programming features, thus considering only a subset of Java, e.g., control flow constructs (`if`, `while`, `try-catch`, etc.), built-in data types (`Int`, `String`, etc.), method calls, and basic OOP features.

We integrate Rigi with Olisipo, an existing geo-replication and coordination system built by Li et al. [7, 36], to form our end-to-end geo-replication AutoGR framework, see Figure 1. Olisipo consists of three key system components. First, there is a runtime library attached with application servers deployed at each site for communicating with the rest of Olisipo for extracting runtime side effects through local transaction execution, and for initiating cross-site replication and necessary coordination. Second, the global coordination service of Olisipo consists of a set of servers, which are spanning over different sites, running a BftSmart [51] consensus protocol. This service will tell whether a side effect that needs coordination can be admitted or not. If so, the dependencies are determined, which implies that this operation is admitted to the global partial order. Otherwise, the coordination service suggests aborting this operation due to conflicts. Finally, at the low level of the runtime there is a geo-replicated data store with replicas spanning over different sites. The store offers causally consistent replication, in which it receives the side effects with ordering dependencies and replicates the side effects at every replica when the specified dependencies are locally satisfied. This replication strategy ensures that the serial order of applying side effects at every site is compatible with the global partial order of the corresponding operations, as pointed by the system model description presented in Section 2. However, there is no guarantee from the geo-replicated store that the final state will converge and the correct application semantics are always preserved. This is all done by the joint work of the coordination layer and the static analyzer Rigi.

This integration requires Rigi to generate a static and unique signature, for each side effect that requires coordination, and the list of static signatures will be consumed by the runtime library for initiating the coordination procedure. In addition, the restrictions identified by Rigi are kept and looked up by the coordination service to serialize pairs of conflicting side effects. To minimize the code changes to make target non-replicated applications geo-replicated, the logic of the runtime library is implemented in the database connector. Following this, the code adaption is just to replace the original connector with the one offered by AutoGR. Olisipo offers "Last-Writer-Win (LWW)" and "UniqueIdGeneration" strategies for transforming the non-commutative side effects into commutative ones to enable geo-replicated services to maximize the performance, which match Rigi's optimizations. When the $-DOpt$ is enabled, AutoGR automatically assigns the logical timestamps to corresponding attributes to take advantage of LWW, and enables UniqueIdGeneration on all primary keys associated with the `AUTOINCREMENTAL` keyword.

Finally, the workflow in the AutoGR runtime is as follows: user requests are sent to a web server at a site in close proximity to them. At the primary site, the server executes the user request against the local database replica to collect side effects. Then, the web server consults the coordination service for admitting or rejecting the request via the runtime library. When accepting this request, the web server sends the side effect associated with the dependencies to the geo-replicated store for applying side effects across sites. Finally, results and decisions are delivered back to the end users.

4.7 Discussions

Application languages. Though the current implementation of Rigi only supports applications written in basic Java, the design principles of Rigi are general and can be extended to support advanced Java features and other languages. Such extensions can be made by enhancing Collector to gradually take more Java constructs and library functions into consideration or use language-specific parsers. To link other language applications plus the analysis results achieved by Rigi to the runtime of AutoGR, we can modify the language-specific database connectors to intercept database calls for creating runtime path signatures, and extracting read/write sets for cross-site replication and coordination, as described above.
Backend storage systems. AutoGR relies on serializability and transactional support offered by the backend storage systems, since the abstraction of transactions eases the consistency reasoning. In addition, the Z3 abstraction over SQL queries already expresses the behaviors of key-value stores since the former covers the latter. However, compared to SQL-compliant databases, the use of key-value stores would lead us to make Collector stronger since the application logic needs to perform computations that can be handled by SQL queries.
Incremental analysis. Currently, Rigi does not assume incremental analysis when facing changes in the target code base. It suggests that upon code changes, the developers would need to run Rigi against the whole code base again to infer possible new restrictions for safety. This would result in additional efforts to link the updated application with AutoGR, as newly introduced side effects need to generate unique signatures for runtime lookup. Despite that we believe this overhead is manageable, adapting Rigi to support incremental analysis would be a wise option.
Failure handling. At runtime, we rely on Olisipo to handle failures. The state of the site coordinators is replicated through a

Table 2: Conditions and side effects corresponding to a few simplified examples in the four benchmark applications. Note that the values of parameters such as *itemId* are part of the corresponding side effect, which are obtained when the side effect is generated at the primary site and encoded into that side effect, thus precluding non-determinism.

	Operation	Condition	Side effect
RUBiS	PlaceBid	table_item(*itemId*) != nil ∧ table_item(*itemId*).closed==False ∧ table_user(*uId*) != nil ∧ table_item(*itemId*).max_bid<*price*	table_bids(*bid*) = (*itemId*, *uId*, *price*) table_item(*itemId*).nb_of_bids++ table_item(*itemId*).max_bid = *price*
Small-Bank	Send-Payment	*amount* ≥ 0 ∧ table_checking(*sendAcct*).balance ≥ *amount*	table_checking(*sendAcct*).balance -= *amount* table_checking(*destAcct*).balance += *amount*

Seats	NewRe-servation	table_flight(*fid*) != nil ∧ table_flight(*fid*).seat_left > 0 ∧ table_res(*fid*, *seat*) == nil ∧ table_res(*fid*, *cid*) == nil	InsertReservation(*table_res*, *cid*, *fid*, *seat*, *price*)
			table_flight(*fid*).seat_left -= 1
			UpdateCustomer(*table_customer*, *cid*, *info*, *price*)
			AddFrequentFlyer(*table_fflyer*, *cid*, *alid*)

Health-Plus	makeLabAp-pointment	table_lab_appointment(*laId*) == nil ∧ table_lab_test(*testId*) != nil ∧ table_tmp_bill(*patientId*) == nil	InsertLabAppintment(*laId*, *testId*, *patientId*, ...)
			InsertTmpBill(*billId*, *patientId*, *fee*)

Table 3: The basic statistics of applications and the lines of code generated by Rigi for applications.

	Code-base	#op/ #update_op/ #loop_op	Code generated by Rigi		
			DB	Cond.	Side effects
SmallBank	*2.5k*	6/5/0	29	38	119
RUBiS	*9.8k*	28/6/0	113	62	191
Seats	*5.0k*	7/4/2	267	65	207
HealthPlus	*15.7k*	157/50/40	524	1,113	1,387

Paxos-like protocol, thus tolerating up to f failures. However, the vulnerability in AutoGR lies in the operation propagation. Currently, AutoGR relies on each site to send its own local side effects to all remote sites. A pair-wise network outage or failure of a site would possibly result in partial replication, e.g., side effects were missing by some sites. However, this can be addressed using standard techniques for exchanging causal logs or reliable multicast.

5 EVALUATION

We exercise four web applications: *SmallBank* [53], *RUBiS* [29], *Seats Reservation* [54], and *HealthPlus* [6]. The first three have been extensively used in related work and the last one is a real-world deployable application. They are designed by having web servers running application logic to interact with a shared backend database. Table 3 shows their basic statistics. The simplest application, *SmallBank*, simulates an online banking system where four out of the total five transactions contain updates. *RUBiS* emulates an eBay-like [3] online auction website, where six out of the total 16 transactions are update transactions. *Seats Reservation* (*or short, Seats*) models an electronic Airline ticketing service. It consists of six transactions, four out of which contain updates. HealthPlus is a real-world deployable management system for health care facility. It consists of 157 transactions in total and 50 of which are updates. Given the large code base and the prohibitive number of checks (5112) for manual analysis in HealthPlus, we use it to evaluate the generality and practicability of our approach.

5.1 Static analysis

Here, we report the Rigi's main results and refer readers to our online documents [12] for other detailed results.

Side effects and path conditions. Table 3 summarizes the number of code automatically generated by Rigi. First, it reads the database organization files (e.g., SQL schema) to produce the Z3Py specifications about the table data structures. For this purpose, Rigi generates 29, 113, 267, and 524 lines of database specifications for *SmallBank*, *RUBiS*, *Seats*, and *HealthPlus*, respectively. *SmallBank* has the least amount of database specifications, as it has only three tables, while *HealthPlus* is at the other end with 37 tables.

Second, Rigi traverses all code paths, and creates 119-1387 and 38-1113 lines of side effect and path condition specifications for the four applications, respectively. We show the exemplified extracted side effects and path conditions in Table 2. The number of side effect and path condition specifications is proportional to the codebase size for applications except for *Seats*. Though *Seats* has fewer tables and transactions than *RUBiS*, they are more complex (e.g., containing more attributes, or more state changes), leading to 136.28% and 8.38 % more DB and side effect specifications, respectively. Finally, we also examine the number of paths and loops in those update transactions. In *SmallBank* and *RUBiS*, each update transaction does not contain loops and has only a single path leading to side effects. By contrast, there are 2 and 40 transactions that contain loops in *Seats* and *HealthPlus*, respectively. Rigi further discovers that the two *Seats*' loops obey the loop iteration independence property, thus can be unfolded one time. For *HealthPlus*, 24 out of 40 loop transactions obey the independence property. Overall, each of these transactions of the four applications leads to between 1 and 9 pairs of conditions and side effects.

Table 4: Statistics of checks, where #failed_commu and #failed_sem correspond to the number of failed commutativity and semantics checks, respectively.

		#checks	#failed commu.	#failed sem.	#restr.
SmallBank	normal	20	0	4	4
	opt.		0	4	4
RUBiS	normal	90	19	15	21
	opt.		4	8	8
Seats	normal	90	32	26	35
	opt.		25	26	28
HealthPlus	normal	5112	148	76	148
	opt.		61	47	74

Checks and restrictions. Table 4 summarizes the number of restrictions identified by Rigi for four applications. By taking pairs of side effects and path conditions extracted above as the input, Rigi performs 20, 90, 90, and 5112 checks for *SmallBank*, *RUBiS*, *Seats*, and *HealthPlus*, respectively, where half are commutativity checks, and the other half are semantics checks. Based on the pair-wise checking presented in Section 4.3, the number of checks is roughly the square of the number of distinct code paths. This explains the substantial gap in the number of checks between *HealthPlus* and the other three applications.

We configure Rigi with two modes: `Normal` or `Opt`, corresponding to disable or enable the optimizations presented in Section 4.5. Note that switching on optimizations may affect the number of identified restrictions, rather than the number of checks. We first focus on the the `Normal` mode results. For *SmallBank*, all side effects of its four update transactions are commutative w.r.t. each other, thus it does not report any commutativity violation restrictions. Besides, it identifies four restrictions related to the failures of semantics checks. Take the restriction between a pair of *SendPayment* operations as an example. The semantics check performed by Rigi takes as input the path condition and side effect of *SendPayment* as shown in Table 2, and identifies a possible violation: when two *SendPayment* operations are running concurrently, the side effect of one operation possibly invalidates a path condition associated with the side effect of the other, which asserts the remaining balance is enough. The remaining three restrictions follow a similar logic.

Rigi suggests 21 restrictions for *RUBiS*, out of which 19 and 15 correspond to commutativity and semantic violations, respectively. There are 35 restrictions identified for *Seats*, where 32 and 26 are related to commutativity and semantic violations, respectively. *HealthPlus* has 148 restrictions, out of which 76 correspond to both commutativity and semantic violations, while the remaining 72 are only related to failed commutativity checks. Interestingly, there are 13, 23, and 76 pairs of side effects and path conditions fail in both commutativity and semantics checks, for *RUBiS*, *Seats*, and *HealthPlus*, respectively. Take the pair `(RegisterItem, StoreBid)` in *RUBiS* as an example. Its commutativity check fails because `RegisterItem` overwrites `nb_of_bids`, which is increased by `StoreBid`. Its semantics check fails as `StoreBid` asserts the new bid is larger than `max_bid`, which can be invalidated by `RegisterItem` when inserting a new item. Among the four applications, *Seats* exhibits the highest ratio of restricted pairs of side effects. The reason is two-fold. First, side effects of *Seats*' transactions are complex and with limited commutativity, due to either conflicting writes on shared attributes or concurrent addition and removal from a shared database table. Second, all reservation-related side effects have to ensure a strict invariant, encoded in the corresponding path conditions, that one seat cannot be sold to two different users.

Next, we evaluate Rigi under the `Opt` mode, which enables the two commutative transforming features introduced in Section 4.5. When switching from `Normal` to `Opt`, the results of *SmallBank* remain unchanged, since all its side effects commute w.r.t each other. The numbers of restrictions of *RUBiS* and *HealthPlus* are greatly reduced from 21 to 8, and 148 to 74, respectively. For *RUBiS*, the UniqueIdGeneration feature eliminates 15 and 7 failed commutativity and semantics checks. Similarly, the last-writer-win and UniqueIdGeneration strategy together introduce a 58.78% drop in the number of failed commutativity checks for *HealthPlus*, and UniqueIdGeneration additionally eliminates 29 failed semantics checks.

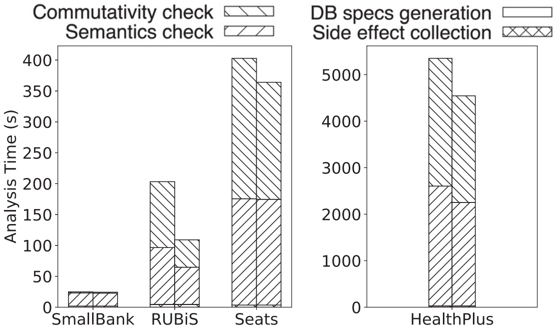


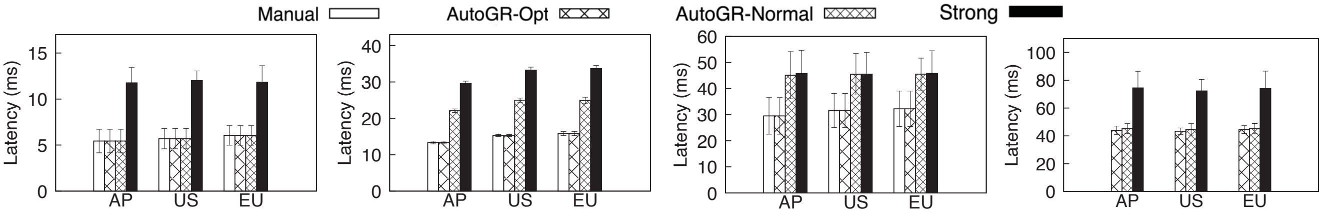
Figure 5: Analysis Cost breakdown for static analysis of four representative applications. In each bar cluster, the bars on the left and right correspond to the time cost of running Rigi under `Normal` or `Opt` mode, respectively.

Nevertheless, we observe the limited improvement for *Seats* even with optimizations, i.e., only 7 failed commutativity checks have been removed while the failed semantics checks remain the same. This is because only commutativity violations corresponding to the *UpdateCustomer* side effects can be eliminated by the last-writer-win strategy. We also see the opportunity to apply some other advanced CRDT solutions, e.g., add-win or remove-win set [48], to eliminate the commutativity violations between *NewReservation* and *DeleteReservation*. Considering the correctness reasoning of such transformation, we leave this exploration as future work.

Comparison to other analysis methods. Here, we compare our work with the most relevant tool Indigo. However, as Indigo offers only two simple APIs (*Decrement* and *Increment*), it is impossible to write specifications for the three case study applications. Therefore, to evaluate the effectiveness of Rigi, we choose to compare its results with manual efforts to achieve PoR consistency, which identifies a set of rules to find a minimal set of restrictions [36]. We find that the restrictions automatically identified by Rigi match those found by the manual process for applications except *HealthPlus*. This shows that Rigi can find a minimal set of sufficient restrictions without requiring any explicit invariants and conditions describing the target operations' behaviors. However, with more than 5000 checks required by *HealthPlus* (see Table 4), it is impractical to perform the reasoning manually.

Performance implications of Rigi. We measure Rigi's time cost on a server, which has 16 Intel(R) Xeon(R) CPU E5-2620 cores and 64GB RAM, and runs CentOS 7, Python 3.6.6, Z3 4.8.10, and Java 8. The single-thread performance numbers are summarized in Figure 5 under the two aforementioned modes. Overall, within the `Normal` mode, Rigi takes from 23.8 seconds to 1.5 hours to analyze all four applications. As expected, it spends the least amount of the analysis time for the simplest *SmallBank*, and the most for *HealthPlus* with the largest codebase.

Figure 5 presents the analysis cost breakdown for the four steps of the static analysis: generating DB specifications, collecting side effects plus conditions, checking semantic violations, and checking commutativity. Among the four steps, the time cost of the first two steps is almost negligible. For instance, it only takes 29 seconds to complete the DB specification generation and side effect/path condition even for the largest *HealthPlus*. The commutativity and



(a) SmallBank (b) RUBiS (c) Seats (d) HealthPlus

Figure 6: Average latency perceived by users at different sites across four system configurations (lower is better)

semantics checks dominate the analysis cost and the cost of performing these checks exhibits quadratic growth. Among the four applications, it takes the least time cost to analyze *SmallBank*, as the side effects and conditions of this application are relatively straightforward, plus it requires the least number of checks. Although *Seats* contains fewer transactions than *RUBiS*, it takes a longer time to complete the static analysis. This is due to the following two reasons. First, some transactions in *Seats* have more paths than *RUBiS*. Second, some paths of *Seats* contain more update queries than *RUBiS*. Thus, the analysis of Seats is more complex than *RUBiS*. *HealthPlus* observes the longest analysis time due to its larger transaction space.

When switching to `Opt`, the time for computing the DB specifications and collecting pairs of conditions and side effects remains almost the same across the four applications. *SmallBank* observes no change in the checking time since its checks are not affected (see Table 4). However, we observe an up to 58.38% reduction in the commutativity checking time for the remaining three applications. This result is consistent with the elimination of failed commutativity checks in Table 4. Furthermore, the semantics checking time has also been reduced for *RUBiS* and *HealthPlus*, except *Seats*. This is because the semantics checks of *Seats* remain the same.

Additionally, with 16 threads, Rigi under the `Normal` mode introduces a 3.2-11.5*X* speedup of the static analysis time for the four examined applications, e.g., in the largest improvement, it reduces the time cost of analyzing *HealthPlus* from 5351.209s to 465.402s. Thus, we conclude that given the analysis is a one-time and offline job, the cost is moderate, and Rigi scales well to applications with larger sizes. Note that even for experienced experts with state-of-the-art tools like PoR, it may take a few days for small-scale applications and usually it is infeasible for large ones. Thus the analysis cost saving is tremendous, not to mention that the whole process is fully automated with little manual effort.

5.2 Geo-replication

System configurations. We run geo-replicated experiments on EC2 m4.2xlarge virtual machines (VM) across three sites: Asia-Pacific-Southeast (AP), US-West (US), and EU-Central (EU). Each VM has 8 vCPUs, 32 GB of RAM, and runs Ubuntu 16.04, MySQL 5.5, Tomcat 6.0, and Java 8. The average round-trip latency between any pair of the three sites ranges from 157 to 169ms.

We deploy AutoGR across the three sites, where each site has a database server host a copy of replicated data, a coordinator to participate in a Paxos-like consensus protocol [22], and a proxy that is attached to the original application code to execute user requests via AutoGR. We run applications with two different configurations: (a) we configure AutoGR with the set of restrictions identified by Rigi without the commutative transformation (denoted as "AutoGR-Normal"); and (b) we configure AutoGR with the set of restrictions identified when the commutative transformation is enabled (denoted as "AutoGR-Opt").

Baselines. We deploy two baselines: (1) a geo-replicated service spanning three sites but serializing all updates, which is a standard deployment offering strong consistency and requiring no manual work, denoted as "Strong"; and (2) a three-site geo-replicated service coordinating use requests guided by the aforementioned manually identified restrictions (denoted as "Manual").

Datasets and workloads. The dataset of *RUBiS* is generated using the following parameters: 500,000 old items, 33,000 active items, and 1,000,000 users. We populate a 1.3-2.5 GB database with randomly generated records for the remaining three applications. *RUBiS* runs the representative bidding mix workload, which consists of 85% read-only transactions and 15% update transactions. For the other three applications, we generate workloads for them by following the *RUBiS*'s setup. Users are evenly distributed across the three sites, issuing requests in a close loop to the replica at the closest proximity. We measure average user-observed latency per site and the aggregated throughput across sites by adding up their individual values. For each evaluation point, we repeat experiments three times and report the average.

User-observed latency. The primary goal of our proposal is to automatically transition non-replicated applications to be geo-replicated with enhanced performance in terms of low user perceived latency. Figure 6 summarizes the average latency observed by users at different sites across all system configurations. The error bars represent standard deviation. For all applications except *HealthPlus*, compared to "Manual", "AutoGR-Opt" achieves similar latency performance. This finding is consistent with the results obtained in the case studies section that Rigi with the commutative transformation option enabled can identify the same restriction set as PoR consistency does. We do not show such a comparison for *HealthPlus* since the manual analysis is infeasible for such a large codebase. In comparison with "Strong", "AutoGR-Opt" reduces the user observed latency from 39.8% to 61.8% for all four applications. This is because the set of identified restrictions is minimal and only a small fraction of requests that need cross-site coordination, while in "Strong" all update requests are coordinated. Among the *four* applications, the latency improvement of *RUBiS* is the best, while that of *Seats* is the worst. This is because with the AutoGR-Opt setting, *RUBiS* imposes the fewest number of restrictions, but *Seats* needs the most.

"AutoGR-Normal" achieves 24.84% to 54.1% lower user observed latency than "Strong" for *RUBiS*, *SmallBank*, and *HealthPlus*, but behaves similarly as "Strong" for *Seats*. This result is consistent

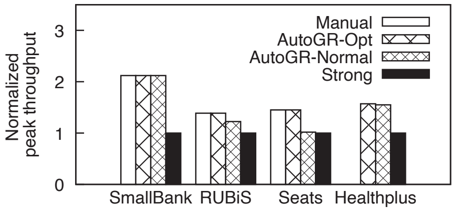


Figure 7: Normalized peak throughput numbers

with the restrictions identified by Rigi without the commutative transformation. In that case, "AutoGR-Normal" almost serializes all update transactions for *Seats*. Finally, "AutoGR-Opt" achieves the same performance as "AutoGR-Normal" for *SmallBank* and *HealthPlus* for different reasons. For *SmallBank*, "Opt" does not eliminate its restrictions. Unlike this, even the number of restrictions drops in *HealthPlus*, the limited performance improvement is because the workload is random and the number of user requests, whose side effects are restricted by AutoGR-Normal but freed by AutoGR-Opt, accounts for a quite small fraction. Therefore, we expect visible improvements in the presence of more skewed workloads. In contrast, "AutoGR-Opt" reduces latency from 24.8% to 26.0% than "AutoGR-Normal" for both RUBiS.

Peak throughput. Next, we shift our attention to the implication of peak throughput. Figure 7 illustrates the peak throughput of different system configurations, which is normalized to the peak throughput of "Strong". These results are consistent with latency observed in Figure 6. "AutoGR-Opt" performs the best for all applications. It achieves as good performance as "Manual", which represents the best performance gain with the most human intervention for manually identifying restrictions, for all applications except *HealthPlus* (Again, we haven't conducted the manual analysis for *HealthPlus* since it can be extremely time-consuming and also error-prone). "AutoGR-Opt" achieves a speedup ranging from 1.39 to 2.12 times, compared to "Strong". Finally, if we disable the commutative transformation, "AutoGR-Normal" still introduces a 1.22-2.12× speedup of peak throughput for all applications except *Seats*. The reasons for throughput comparison are similar to those of the latency comparison we presented above.

6 RELATED WORK

Geo-Replicated Systems. Many cloud storage systems offer geo-replication features [11, 15, 24, 26, 28], assuming either strong or eventual consistency [24, 26, 28]. Recently, Azure CosmosDB [15] and Google Cloud DataStore [11] provide a set of APIs that read data with different consistency guarantees, but they still serialize updates for strong consistency. Unlike them, AutoGR follows restriction-based consistency models, and automatically enables the geo-replication feature for un-replicated applications with enhanced performance without requiring developers to identify manually problematic concurrent executions that should be coordinated.

Consistency models and fine-grained coordination. To close the gap between strong and eventual consistency, hybrid consistency models, such as RedBlue consistency [35], have been proposed to allow strongly and causally consistent operations to co-exist in a single system. Recently, to minimize the coordination in geo-replication, fine-grained consistency proposals such as PoR consistency [36] completely drop the definition of consistency levels and instead map the consistency semantics into a set of ordering restrictions over pairs of operations (analogous to "conflict relation" in Generic Broadcast [42]). Our approach is built on top of this line of work and extends them for two main aspects – a) we build an automated analysis tool, which identifies the required restrictions and free programmers from the error-prone and time-consuming task; and b) we offer an end-to-end solution for automatically deploying non-replicated serializable applications as geo-replicated with enhanced performance and preserved application semantics.

Tool supports. Some recent proposals aim to relieve developers' burdens to adapt their applications to use fine-grained consistency models in geo-replicated scenarios [17, 30, 34]. SIEVE statically computes the weakest preconditions which lead to side effects that always preserve invariants [34]. Indigo [20] checks whether the concurrent execution of a pair of operations would violate invariants and provides a set of mechanisms to resolve conflicts. Hamsaz takes a sequential object as input and automatically synthesizes a replicated object that avoids unnecessary coordination [31]. These works rely on specifications and application-specific invariants that programmers need to write manually. Unlike them, Rigi does not assume the existence of specifications; instead, it infers the side effects and their path conditions solely from the source code. MixT is a language for designing geo-distributed transactions [38], where different levels of consistency are associated with the attributes of database tables. However, this approach requires programmers to re-implement their applications using MixT. The goal of IPA [19] is orthogonal to ours, as it ensures invariants on an eventually consistent replicated store by repairing and compensating violations introduced by running conflicting operations in parallel.

CRDTs. We significantly differ from CRDTs [44, 48] by (1) the primary interest of AutoGR and Rigi is to identify pairs of non-commutative side effects based on the original code rather than improving operation commutativity; and (2) convergence alone cannot prevent invariant violations, which can be avoided by Rigi's semantics checks and AutoGR's runtime coordination. Furthermore, CRDTs are complementary to our proposal, and AutoGR can be extended to incorporate more CRDTs, as some case-study applications show limited commutativity and the use of LWW indeed leads to fewer restrictions and better geo-replication performance.

7 CONCLUSION

We present AutoGR – an automated end-to-end framework for deploying non-replicated applications as geo-replicated. The core of AutoGR is a static analysis tool Rigi that automatically identifies necessary restrictions without user interventions. Experimental results show that AutoGR has the following benefits: (1) low reasoning cost, (2) optimal performance, and (3) little manual effort.

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