POSTER: Asynchrony versus Bulk-Synchrony for a Generalized N-body Problem from Genomics

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Abstract

This work examines a data-intensive irregular application from genomics, a long-read to long-read alignment problem, which represents a kind of Generalized N-Body problem, one of the "seven giants" of the NRC Big Data motifs [5]. In this problem, computations (genomic alignments) are performed on sparse and data-dependent pairs of inputs, with variable cost computation and variable datum sizes. In particular, there is no inherent locality in the pairwise interactions, unlike simulation-based N-Body problems, and the interaction sparsity depends on particular parameters of the input, which can also affect the quality of the output. We examine two extremes to distributed memory parallelization for this problem, bulk-synchrony and asynchrony, with real workloads. Our bulk-synchronous implementation, uses collective communication in MPI, while our asynchronous implementation uses cross-node RPCs in UPC++. We show that the asynchronous version effectively hides communication costs, with a memory footprint that is typically much lower than the bulk-synchronous version. Our application, while simple enough to be a kind of proxy for genomics or data analytics applications more broadly, is also part of a real application pipeline. It shows good scaling on real input problems, and at the same time, reveals some of the programming and architectural challenges for scaling this type of data-intensive irregular application.

CCS Concepts: • Computing methodologies \rightarrow Parallel computing methodologies; Distributed computing methodologies; • Applied computing \rightarrow *Computational genomics; Bioinformatics; Molecular sequence analysis;* Computational proteomics.

Keywords: HPC, genomics, bioinformatics, big data

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1 Extended Abstract

Per the 2013 National Academies report on Big Data Frontiers, "Generalized N-body problems (Gray and Moore, 2001) include virtually any problem involving distances, kernels, or other similarities between (all or many) pairs (or higher-order n-tuples) of points". This describes classic, well-understood N-Body problems like particle simulation. It also describes currently challenging problems with non-Euclidean or higherdimensional similarity metrics, proliferating from bioinformatics, computer vision, computational chemistry, and other Big Data domains [3, 5, 15]. Long-read to long-read alignment, a genomics problem, is representative of this challenging subclass.

To achieve acceptable fidelity for genomic analysis in the presence of sequencing errors, pairwise alignment is used to compute similarity or "distance" between reads (strings representing DNA fragments) from a set of input reads. For a read pair with maximum length n, pairwise alignment is $O(n^2)$ in general (Needleman-Wunsch 1970, Smith-Waterman 1981) [10, 12]. However, average-case O(n) methods, such as X-drop [16] can achieve good accuracy for longread to long-read alignment in particular [7]. Yet, individual long reads vary widely in length, roughly $n \in (10^2, 10^5)$, and while algorithms like X-drop offer significantly lower expected complexity, they also entail much more irregularity due to their dynamic adaptation and early-termination heuristics. The overall problem can be solved by computing the pairwise alignment of each input read to each other input read in $O(N^2 \times n^2)$. However, the size of the input read sets vary with the size of the underlying genomes and the sequencing depth; in practice, the all-to-all approach is intractable for all but the smallest workloads. Domain-specific analysis [2, 4, 7-9, 14] can effectively reduce the number of interactions (alignments). However, these analyses also typically yield unstructured sparse (data-dependent) interaction

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patterns across the input reads, over which the many-tomany pairwise alignment must be performed. This poses challenges for parallel load balancing and communication cost minimization, and remains the dominant scalability bottleneck of practical pipelines [6]. Similar challenges appear in other computational biology problems as well, such as protein alignment [11] and protein searches in massive data sets [13].

This work examines practical approaches, near the bulksynchronous and asynchronous extremes, for many-to-many long read alignment. Our bulk-synchronous approach in MPI 2.0 uses collective communication for the many-to-many read exchange, supporting independently parallel computation of many alignments in subsequent compute steps. It utilizes all available memory to minimize the number of supersteps and maximize bandwidth-utilization. The dataintensive nature of the application and the fact that there is no perfect partitioning strategy (in general) as we illustrate, motivates alternative one-sided or asynchronous approaches. Our asynchronous approach, retrieves reads required for sets of local alignment tasks asynchronously, and overlaps retrievals with unrelated computations. Given its performance and programmability for the target use-case, the asynchronous version is implemented in UPC++ [1] with RPCs. From the outset, it is unclear whether the latency of many finegrained messages, with asynchronous framework overheads, can compete with a small number of bisection-bandwidthbound exchanges. Our analysis considers both communication performance and memory footprint for fixed workloads.

Results were collected on a Cray XC40, Cori KNL at NERSC, using three real workloads. Strong scaling results up to 32K cores are shown, with runtime breakdowns and memory footprints. Within a node, both codes effectively reduce the runtime from just under 1 hour on a single core to just under 1 minute on 64 cores. For another larger workload, both codes reduce the runtime from an estimated \approx 7 hours on a single core to under 10 seconds on 4K cores (64 nodes).

While both versions achieve good scalability, the asynchronous version effectively hides communication costs, despite the irregularity of remote-lookups and alignment computations, resulting in up to 20% better efficiency. The asynchronous version also maintains a typically lower memory footprint over the bulk-synchronous version. Across workloads, implementations, and architectures, we expect that bandwidth for many-to-many exchanges supporting bulksynchronous approaches, versus the (balance of) alignment computation to one-sided or asynchronous message latency supporting asynchronous approaches, will determine the relative performance and scalability of each approach.

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