

# Linear Complexity Entropy Regions

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**Abstract**—A great many problems in network information theory, both fundamental and applied, involve determining a minimal set of inequalities linking the Shannon entropies of a certain collection of subsets of random variables. In principle this minimal set of inequalities could be determined from the entropy region, whose dimensions are all the subsets of random variables, by projecting out dimensions corresponding to the irrelevant subsets. As a general solution technique, however, this method is plagued both by the incompletely known nature of the entropy region as well as the exponential complexity of its bounds. Even worse, for four or more random variables, it is known that the set of linear information inequalities necessary to completely describe the entropy region must be uncountably infinite. A natural question arises then, if there are certain nontrivial collections of subsets where the inequalities linking only these subsets is both completely known, and have inequality descriptions that are linear in the number of random variables. This paper answers this question in the affirmative. A decomposition expressing the collection of inequalities linking a larger collection of subsets from that of smaller collections of subsets is first proven. This decomposition is then used to provide systems of subsets for which it both exhaustively determines the complete list of inequalities, which is linear in the number of variables.

**Index Terms**—information inequalities, complexity of entropy region

## I. INTRODUCTION

For two sets  $A$  and  $B$ , we will denote by  $A^B = \{\mathbf{x} | \mathbf{x} : B \rightarrow A\}$ . For two subsets  $\mathcal{P}_1 \subseteq \mathbb{R}^{A_1}$  and  $\mathcal{P}_2 \subseteq \mathbb{R}^{A_2}$  of Euclidean space, define  $\mathcal{P}_1 \bowtie \mathcal{P}_2$  as,

$$\mathcal{P}_1 \bowtie \mathcal{P}_2 = \left\{ \mathbf{f} \in \mathbb{R}^{A_1 \cup A_2} \left| \begin{array}{l} \exists \mathbf{g}_1 \in \mathcal{P}_1, \mathbf{g}_2 \in \mathcal{P}_2 \text{ s.t. :} \\ \mathbf{f}(a) = \mathbf{g}_1(a), \forall a \in A_1 \\ \mathbf{f}(b) = \mathbf{g}_2(b), \forall b \in A_2 \end{array} \right. \right\} \quad (1)$$

Let the ground set  $E \subseteq \mathbb{N}$  index a set of  $|E|$  finite discrete random variables (RVs), and let,

$$\Delta_E = \left\{ p_E \left| \begin{array}{l} \mathcal{X} \subseteq \mathbb{N}, p_E : \mathcal{X}^E \rightarrow [0, 1], \\ \sum_{\mathbf{x}_E \in \mathcal{X}^E} p_E(\mathbf{x}_E) = 1 \end{array} \right. \right\} \quad (2)$$

be a set of valid joint probability mass functions (PMFs) for these RVs paired with their ranges. For any such joint PMF  $p_E \in \Delta_E$  and any subset  $A \subseteq E$  define the generalized marginalization operator  $M : \Delta_E \times 2^E \rightarrow \bigcup_{A \subseteq E} \Delta_A$  with,

$$M(p_E, A) = p_A, \quad A \neq \emptyset \quad (3)$$

where,

$$p_A(\mathbf{x}_A) = \sum_{\mathbf{x}_E \in \mathcal{X}^E \text{ s.t.}} p_E(\mathbf{x}_E) \quad (4)$$

$$\mathbf{x}_E(e) = \mathbf{x}_A(e) \quad \forall e \in A$$

Given two PMFs  $p_{E_1}$  and  $p_{E_2}$  such that  $E_1 \cap E_2 \neq \emptyset$ , we write

$$p_{E_1} \stackrel{M}{=} p_{E_2} \quad (5)$$

if and only if

$$M(p_{E_1}, E_1 \cap E_2) = M(p_{E_2}, E_1 \cap E_2) \quad (6)$$

Define the function  $\hat{h}_E : \Delta_E \rightarrow \mathbb{R}^{2^E}$  such that for any  $A \subseteq E$ ,

$$\hat{h}_E(p_E) : A \mapsto - \sum_{\mathbf{x}_A \in \mathcal{X}^A} p_A(\mathbf{x}_A) \log_2(p_A(\mathbf{x}_A)) \quad (7)$$

with again  $M(p_E, A) = p_A$ . Here, based on preference, we can either think of  $\hat{h}_E(p_E)$  as a vector whose elements are indexed by all subsets  $2^E$ , or equivalently as a function  $\hat{h}_E(p_E) : 2^E \rightarrow \mathbb{R}$  assigning a number to each such subset of  $E$ . Either way, the value associated with the subset  $A$  is the joint Shannon entropy of that subset of the random variables.

The entropy region for the ground set  $E$  will be the image of the set of all joint PMFs  $\Delta_E$  under this map  $\hat{h}_E$

$$\Gamma_E^* = \hat{h}_E(\Delta_E) \quad (8)$$

denoted as  $\mathbf{h} = [\mathbf{h}(A) | A \subseteq E]$  an entropic vector in  $\Gamma_E^*$  and let,

$$\Xi = \hat{h}_E^{-1}(\mathbf{h}) \quad (9)$$

be the collection of PMFs that is associated with  $\mathbf{h}$ .

$\Gamma_E^*$  was first introduced in [1], motivated by the desire to develop a new framework to proof information inequalities. By definition,  $\Gamma_E^*$  gives a complete characterization of all information inequalities and thus is closely related to some important problems in probability theory and information theory. The attempt to fully characterize  $\Gamma_E^*$  started from the so-called basic information inequalities [2], which are implied from the non-negativity of the conditional mutual informations and are of the form:

$$\mathbf{h}(A) + \mathbf{h}(B) \geq \mathbf{h}(A \cup B) + \mathbf{h}(A \cap B), \quad \forall A, B \subseteq E \quad (10)$$

Each vector  $\mathbf{h} \in \Gamma_E^*$  must satisfies all the inequalities of the form (10), which means, when  $\Gamma_E$  is defined as,

$$\Gamma_E = \{\mathbf{h} \in \mathbb{R}^{2^E} | \mathbf{h} \text{ satisfies (10), } \forall A, B \subseteq E\} \quad (11)$$

we have an outer bound  $\Gamma_E^* \subseteq \Gamma_E$ .  $\overline{\Gamma_E^*}$  denotes the closure of  $\Gamma_E^*$ , the natural question of whether or not  $\overline{\Gamma_E^*} = \Gamma_E$  motivates the finding of non-Shannon type inequalities [3]–[5], which are defined as the information inequalities that can not be implied from basic inequalities in (10). The closure of the entropy region  $\overline{\Gamma_E^*}$  was proved to be a closed convex cone [3], and for  $|E| \leq 3$  the cone is polyhedral and  $\overline{\Gamma_E^*} = \Gamma_E$ . The equality doesn't hold, meaning  $\overline{\Gamma_E^*} \neq \Gamma_E$ , for  $|E| \geq 4$ , moreover  $\overline{\Gamma_E^*}$  is not polyhedral when  $|E| \geq 4$ , and its inequality and extreme ray descriptions are incompletely known [6], [7].

Significant progresses have been achieved on difference aspects of the topic [8]–[10] [11]–[13]. However, much less has been written considering the instance when we only need the list of all fundamental inequalities linking those entropies of only certain subsets. Furthermore, it is only this restriction to only a certain collection of subsets that is often required in the multiterminal information theory applications. Hence, in this article we are aiming to study the restriction of the entropy region to only those dimensions associated with a certain given collection of subsets.

To describe this concept properly, consider a collection of subsets  $\mathcal{F} \subseteq 2^E$ , and define the projection,

$$\text{proj}_{\mathcal{F}} \overline{\Gamma_E^*} = \left\{ h \in \mathbb{R}^{\mathcal{F}} \mid \exists \mathbf{h} \in \overline{\Gamma_E^*}, h(A) = \mathbf{h}(A) \forall A \in \mathcal{F} \right\} \quad (12)$$

Again, most applied problems, for instance network coding capacity regions [14], depend on  $\Gamma_E^*$  only through its projection  $\text{proj}_{\mathcal{F}} \overline{\Gamma_E^*}$  for some smaller collection of subsets  $\mathcal{F} \subsetneq 2^E$ . Even worse, even the outer bounds such as the Shannon outer bound have both a number of inequalities and dimensions that grow exponentially in  $|E|$ . Significantly, while  $\overline{\Gamma_E^*}$  is non-polyhedral for  $|E| \geq 4$ , and these outer bounds have complexities that are exponential in  $|E|$ , there is the possibility that its projection  $\text{proj}_{\mathcal{F}} \overline{\Gamma_E^*}$  can in some instances be polyhedral and, further, have a number of inequalities and dimensions that grow linearly in  $|E|$ . This paper proves that indeed, many non-trivial such instances exist.

## II. PRELIMINARIES

To understand the general idea behind the method the paper utilizes to prove these existence of these cases, let  $E = E_1 \cup E_2$ , and suppose that  $\mathcal{F} = 2^{E_1} \cup 2^{E_2}$ . In the trivial case that  $E_1 \cap E_2 = \emptyset$ , it is direct to show that,

$$\text{proj}_{2^{E_1} \cup 2^{E_2}} \overline{\Gamma_E^*} = \overline{\Gamma_{E_1}^*} \bigcap \overline{\Gamma_{E_2}^*} \quad (13)$$

The more general case, when there's a collection of subsets  $E_1, \dots, E_m$  such that  $E_i \cap E_j \neq \emptyset$  for some  $i, j \in [1, m]$ , is more complicated. The main contribution of the paper is summarized as the following theorem.

### A. Main Results

**Theorem 1.** Let  $n \in \mathbb{N}$ , consider a collection of  $m = 2n$  subsets  $E_{2i+1} = \{3i+1, 3i+2, 3i+3\}$ ,  $E_{2i+2} = \{3i+2, 3i+3, 3i+4\}$ ,  $\forall i \in [0, n-1]$ . Let  $\mathcal{F} = \bigcup_{i=1}^m 2^{E_i}$ , then  $\text{proj}_{\mathcal{F}} \overline{\Gamma_E^*}$  is a **polyhedral cone** with:

- 1) **explicitly known inequality descriptions.**
- 2) a number of inequalities that grows **linearly** in  $|E|$ .

**Theorem 1** is proved based on the following two theorems.

**Theorem 2.** Let  $E_1, E_2$  two ground sets with such that  $|E_1 \cap E_2| = 1$ , then

$$\text{proj}_{2^{E_1} \cup 2^{E_2}} \overline{\Gamma_E^*} = \overline{\Gamma_{E_1}^*} \bigcap \overline{\Gamma_{E_2}^*} \quad (14)$$

**Theorem 3.** If  $E_1 = \{1, 2, 3\}$ ,  $E_2 = \{2, 3, 4\}$  and  $E = E_1 \cup E_2 = \{1, 2, 3, 4\}$ , then,

$$\text{proj}_{2^{E_1} \cup 2^{E_2}} \overline{\Gamma_E^*} = \overline{\Gamma_{E_1}^*} \bigcap \overline{\Gamma_{E_2}^*} \quad (15)$$

### B. Outer bounds on $\text{proj}_{\bigcup_{i=1}^m 2^{E_i}} \overline{\Gamma_E^*}$

A natural outer bound of  $\text{proj}_{\bigcup_{i=1}^m 2^{E_i}} \overline{\Gamma_E^*}$  would concatenate the inequality descriptions for  $\overline{\Gamma_{E_i}^*}$ ,  $i \in [1, m]$  to get,

$$\text{proj}_{\bigcup_{i=1}^m 2^{E_i}} \overline{\Gamma_E^*} \subseteq \bigcap_{i=1}^m \overline{\Gamma_{E_i}^*} \quad (16)$$

A tighter outer bound forces there to be a way to select the distributions to match on their marginals,

$$\text{proj}_{\bigcup_{i=1}^m 2^{E_i}} \overline{\Gamma_E^*} \subseteq \left( \bigcap_{i=1}^m \overline{\Gamma_{E_i}^*} \right) \bigcap \overline{\mathcal{P}_{E_1, \dots, E_m}^{\overline{\Gamma_E^*}}} \quad (17)$$

where,

$$\overline{\mathcal{P}_{E_1, \dots, E_m}^{\overline{\Gamma_E^*}}} = \left\{ \overline{\mathbf{f}} \in \bigcap_{i=1}^m \overline{\Gamma_{E_i}^*} \mid \begin{array}{l} \exists \mathbf{h}_i^{(n)} \in \Gamma_{E_i}^*, i \in [1, m] : \\ 1) \lim_{n \rightarrow \infty} \mathbf{h}_i^{(n)} = \overline{\mathbf{h}_i}, i \in [1, m] \\ 2) \overline{\mathbf{f}}(A) = \overline{\mathbf{h}_i}(A), A \subseteq E_i, i \in [1, m] \\ 3) \forall i \in [1, m], \exists p_{E_i}^{(n)} \in \hat{\mathbf{h}}_{E_i}^{-1}(\mathbf{h}_i^{(n)}), \text{ s.t.} \\ \quad \forall i, j \in [1, m] \text{ if } E_i \cap E_j \neq \emptyset, \\ \quad \text{then } p_{E_i}^{(n)} \stackrel{M}{=} p_{E_j}^{(n)}, \forall n \in \mathbb{N} \end{array} \right\} \quad (18)$$

### C. Marginal problem and Running Intersection Property

**Definition 1.** (The marginal problem on PMF) Consider a collection of marginal PMFs  $p_{E_1}, \dots, p_{E_m}$ , the marginal problem wonders if there exists a joint PMF  $p_E$  with  $E = \bigcup_{i=1}^m E_i$ , such that,

$$M(p_E, E_i) = p_{E_i}, \forall i \in [1, m] \quad (19)$$

Defined above is the marginal problem on PMF, which has received research interests in different literature [15]–[18]. The problem is trivial when the collection of PMFs are defined on disjoint ground sets, in which case the stochastic product of the marginals give us a valid joint PMF. However, the problem is not trivial anymore when the ground sets are not disjoint. One obvious necessary condition for such joint PMF to exist for a marginal problem is pairwise compatibility, meaning for any pair of  $E_i$  and  $E_j$ ,  $M(p_{E_i}, E_i \cap E_j) = M(p_{E_j}, E_i \cap E_j)$  if  $E_i \cap E_j \neq \emptyset$ .

So the question becomes when does the pairwise compatibility implies a global compatibility. In fact, using the theory of hypergraphs, Beeri et al. [19](see also [20]) established a theorem on this problem stating that,

**Theorem 4.** (Beeri et al. [19]) The following two conditions are equivalent:

1) For any collection of pairwise compatible marginal PMFs  $p_{E_1}, \dots, p_{E_m}$  over  $(E_1, \dots, E_m)$ , exists a joint PMF  $p_E$ , over  $E = \bigcup_{i=1}^m E_i$  such that,

$$M(p_E, E_i) = p_{E_i}, \quad \forall i \in [1, m] \quad (20)$$

2) The collection of ground sets  $E_1, \dots, E_m$  satisfy Running Intersection Property (**RIP**).

with the running intersection property defined as,

**Definition 2.** (Running Intersection Property (**RIP**)) A sequence of subsets  $E_1, E_2, \dots, E_m$  of a finite ground set  $E$  satisfies the **Running Intersection Property** if for every  $k \in [2, m]$  the intersection of  $E_k$  with  $\bigcup_{j=1}^{k-1} E_j$  is contained in one of these previous subsets, that is,

$$\forall k \in [2, m], \exists i \in [1, k-1] \text{ s.t. } E_k \cap \left( \bigcup_{j=1}^{k-1} E_j \right) \subseteq E_i \quad (21)$$

### III. THEOREMS AND PROOFS

*Proof.* (sketch of the proof of **Theorem 2**) The containment

$$\text{proj}_{2^{E_1} \cup 2^{E_2}}(\overline{\Gamma_E^*}) \subseteq \overline{\Gamma_{E_1}^*} \bigtimes \overline{\Gamma_{E_2}^*} \quad (22)$$

is straightforward, so to prove the theorem is suffices to prove that

$$\overline{\Gamma_{E_1}^*} \bigtimes \overline{\Gamma_{E_2}^*} \subseteq \text{proj}_{2^{E_1} \cup 2^{E_2}}(\overline{\Gamma_E^*}) \quad (23)$$

The group characterizable entropy region is helpful to prove this containment - by **Corollary 4.1** in [21], we know that,

$$\overline{\text{con}}(\Upsilon_E) = \overline{\Gamma_E^*} \quad (24)$$

In fact, using a rational approximation of the conic combination coefficients, one can show that  $\overline{\text{con}}(\Upsilon_E) = \overline{\text{ray}}(\Upsilon_E)$ . Thus we must show that for any pair of sequences of scaled group characterizable entropy vectors  $\alpha_k \mathbf{h}_1^k \in \text{ray}(\Upsilon_E)$  and  $\alpha_k^k \mathbf{h}_2^k \in \text{ray}(\Upsilon_E)$  with  $\lim_{k \rightarrow \infty} \alpha_k^k \mathbf{h}_1^k(A) = \lim_{k \rightarrow \infty} \alpha_k^k \mathbf{h}_2^k(A), \forall A \subseteq E_1 \cap E_2$ , we can construct a sequence of global entropy vectors  $\alpha^k \mathbf{h}^k$  such that  $\lim_{k \rightarrow \infty} \alpha^k \mathbf{h}^k(A) = \lim_{k \rightarrow \infty} \mathbf{h}_1^k(A), \forall A \subseteq E_1$  and  $\lim_{k \rightarrow \infty} \alpha^k \mathbf{h}^k(A) = \lim_{k \rightarrow \infty} \mathbf{h}_2^k(A), \forall A \subseteq E_2$ . Dirichlet's approximation theorem enables one to construct a rational approximation with a common denominator for  $\alpha_1^k, \alpha_2^k$  and  $\alpha_1^k \mathbf{h}_1(A) - \alpha_2^k \mathbf{h}_2(A)$ , as  $\frac{n_1}{q}, \frac{n_2}{q}, \frac{e}{q}$ , respectively that enables  $|n_1 \mathbf{h}_1^k(A) - n_2 \mathbf{h}_2^k(A) - e|$  to be made arbitrarily small. Repeat the group structure for  $\mathbf{h}_i^k$   $n_i$  times for  $i \in \{1, 2\}$ , and build a map from the cosets for group associated with  $A$  for  $i = 1$  to those for  $i = 2$  that is one to one until running out of cosets for  $i = 1$ , mapping the remaining cosets to a deterministic extra symbol, drawing  $2^e$  times uniformly from the selected coset for  $i = 2$ . The subset entropies of this construction, using the coset index from  $i = 1$  as the random variable for the overlapping element, provide the required entropies in the limit.  $\blacksquare$

**Lemma 1.** Given a collection of subsets  $E_1, \dots, E_m$  with  $E = \bigcup_{i=1}^m E_i$  such that for any  $i, j \in [1, m]$

$$|E_i \cap E_j| = \begin{cases} 1, & |i - j| = 1 \\ 0, & |i - j| > 1 \end{cases} \quad (25)$$

then,

$$\text{proj}_{\bigcup_{i=1}^m 2^{E_i}}(\overline{\Gamma_E^*}) = \bigcap_{i=1}^m \overline{\Gamma_{E_i}^*} \quad (26)$$

*Proof.* (proof of **Lemma 1**) Prove by induction, **Theorem 2** gives the base case that,

$$\text{proj}_{2^{E_1} \cup 2^{E_2}}(\overline{\Gamma_E^*}) = \overline{\Gamma_{E_1}^*} \bigtimes \overline{\Gamma_{E_2}^*} \quad (27)$$

Next for the inductive step, let  $E'_k = \bigcup_{i=1}^k E_i$  and assume that

$$\text{proj}_{\bigcup_{i=1}^k 2^{E_i}}(\overline{\Gamma_{E'_k}^*}) = \bigcap_{i=1}^k \overline{\Gamma_{E_i}^*} \quad (28)$$

To prove it holds when adding  $E_{k+1}$ , let  $E'_{k+1} = E'_k \cup E_{k+1}$ , then applying **Theorem 2** on the pair of sets  $E_{k+1}$  and  $E'_k$  we have,

$$\text{proj}_{2^{E'_k} \cup 2^{E_{k+1}}}(\overline{\Gamma_{E'_{k+1}}^*}) = \overline{\Gamma_{E'_k}^*} \bigtimes \overline{\Gamma_{E_{k+1}}^*} \quad (29)$$

Now project both sides of (29) down to  $\bigcup_{i=1}^{k+1} 2^{E_i}$ , we have

$$\begin{aligned} & \text{proj}_{\bigcup_{i=1}^{k+1} 2^{E_i}}(\text{proj}_{2^{E'_k} \cup 2^{E_{k+1}}}(\overline{\Gamma_{E'_{k+1}}^*})) \\ &= \text{proj}_{\bigcup_{i=1}^{k+1} 2^{E_i}}(\overline{\Gamma_{E'_{k+1}}^*}) \\ &= \text{proj}_{\bigcup_{i=1}^{k+1} 2^{E_i}}(\overline{\Gamma_{E'_k}^*} \bigtimes \overline{\Gamma_{E_{k+1}}^*}) \\ &= \text{proj}_{\bigcup_{i=1}^{k+1} 2^{E_i}}(\overline{\Gamma_{E'_k}^*}) \bigtimes \overline{\Gamma_{E_{k+1}}^*} \\ &= \left( \bigcap_{i=1}^k \overline{\Gamma_{E_i}^*} \right) \bigtimes \overline{\Gamma_{E_{k+1}}^*} \\ &= \bigcap_{i=1}^{k+1} \overline{\Gamma_{E_i}^*} \end{aligned} \quad (30)$$

$\blacksquare$

**Lemma 2.** The  $\bigtimes$  is both associative and commutative.

*Proof.* Straightforward.  $\blacksquare$

*Proof.* (proof of **Theorem 3**) To show that,

$$\text{proj}_{2^{E_1} \cup 2^{E_2}}(\overline{\Gamma_E^*}) \subseteq \overline{\Gamma_{E_1}^*} \bigtimes \overline{\Gamma_{E_2}^*} \quad (31)$$

Pick an arbitrary vector  $\bar{v} \in \text{proj}_{2^{E_1} \cup 2^{E_2}}(\overline{\Gamma_E^*})$ , then we can find a vector  $\bar{h} \in \overline{\Gamma_E^*}$  such that  $\bar{v} = \text{proj}_{2^{E_1} \cup 2^{E_2}}(\bar{h})$ , which means,

$$\bar{v}(A) = \bar{h}(A), \forall A \subseteq E_1 \text{ or } A \subseteq E_2 \quad (32)$$

By definition of  $\overline{\Gamma_E^*}$ ,  $\bar{h}$  is a limit point of  $\Gamma_E^*$ , so we can find a sequence of entropic vectors  $\mathbf{h}^{(n)} \in \Gamma_E^*$  each of which is associated with a PMF  $p_E^{(n)}$ . Create from  $\mathbf{h}^{(n)}$  and  $\bar{h}$  the following terms,



So from **Theorem 3** we know that  $\text{proj}_{2^{E_{2i+1}} \cup 2^{E_{2i+2}}}(\overline{\Gamma_{B_i}^*}) = \overline{\Gamma_{E_{2i+1}}^*} \bowtie \overline{\Gamma_{E_{2i+2}}^*}$ .  $\blacksquare$

Now we have everything we need to prove **Theorem 1**.

*Proof.* (proof of **Theorem 1**) From **Theorem 5** we know that

$$\text{proj}_{\mathcal{F}}(\overline{\Gamma_E^*}) = \bigcap_{i=0}^{n-1} \left( \overline{\Gamma_{E_{2i+1}}^*} \bowtie \overline{\Gamma_{E_{2i+2}}^*} \right) \quad (51)$$

given that for each  $i \in [0, n-1]$ , both  $E_{2i+1}$  and  $E_{2i+2}$  have cardinality 3, we know that

$$\overline{\Gamma_{E_{2i+1}}} = \overline{\Gamma_{E_{2i+1}}^*} \quad (52)$$

$$\overline{\Gamma_{E_{2i+2}}} = \overline{\Gamma_{E_{2i+2}}^*} \quad (53)$$

Plugging the above two equations into (51) we have

$$\text{proj}_{\mathcal{F}}(\overline{\Gamma_E^*}) = \bigcap_{i=0}^{n-1} \left( \overline{\Gamma_{E_{2i+1}}} \bowtie \overline{\Gamma_{E_{2i+2}}} \right) \quad (54)$$

So either  $\overline{\Gamma_{E_{2i+1}}}$  or  $\overline{\Gamma_{E_{2i+2}}}$  involved in (54) is polyhedral, the **explicit** inequality representation of  $\text{proj}_{\mathcal{F}}(\overline{\Gamma_E^*})$  is nothing but the stacking of all inequalities defining  $\overline{\Gamma_{E_{2i+1}}}$  and  $\overline{\Gamma_{E_{2i+2}}}$  together, which gives,

$$\mathbf{h}(E_i \setminus \{j\}) \leq \mathbf{h}(E_i), \forall j \in E_i, \forall i \in [1, m] \quad (55)$$

$$\begin{aligned} \mathbf{h}(\{j, k\} \cup A) + \mathbf{h}(A) &\leq \mathbf{h}(\{j\} \cup A) + \mathbf{h}(\{k\} \cup A), \\ \forall A \subseteq E_i \setminus \{j, k\}, \forall \{j, k\} \subseteq E_i, \forall i \in [1, m] \end{aligned} \quad (56)$$

So the number of these inequalities (including replicas) is,

$$m(3 + \binom{3}{2}2^{3-2}) = 9m = 6(|E| - 4) + 18 \quad (57)$$

The number of these inequalities after excluding replicas is,

$$\frac{m}{2} \left( 2(3 + \binom{3}{2}2^{3-2}) - 3 \right) = \frac{15m}{2} = 5(|E| - 4) + 15 \quad (58)$$

which means that the number of inequalities defining  $\text{proj}_{\mathcal{F}}(\overline{\Gamma_E^*})$  grows **linearly** in  $|E|$ .  $\blacksquare$

#### IV. DISCUSSION AND FUTURE WORK

This manuscript proved several instances of systems of subsets  $\mathcal{F}$ , for which  $\text{proj}_{\mathcal{F}}(\overline{\Gamma_E^*})$  is equal to the pasting of a collection of small entropy regions together, including several in which  $E$  grows arbitrarily large, and, further, several for which, utilizing the known inequalities for the entropy region on two and three random variables, a polyhedral description of this projection of the entropy region can be provided. Additionally, it is shown that not only may this set be polyhedral, but also that it may have a number of inequalities that grows only linearly in  $|E|$ , which provides a stark contrast to the behavior of, for example, the commonly used Shannon outer bound  $\Gamma_E$  which requires a number of inequalities that grows exponentially. A companion manuscript [1] reviews how key problems often expressed in terms of  $\overline{\Gamma_E^*}$  or its bounds, only depend on it through  $\text{proj}_{\mathcal{F}}(\overline{\Gamma_E^*})$ , so this difference in description complexity can be highly useful in applications. Current and future work is cataloging further collections of systems of subsets  $\mathcal{F}$  enabling this type of decomposition.

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