

Independent sets in subgraphs of a Shift graph

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Abstract

Erdős, Hajnal and Szemerédi proved that any subset G of vertices of a shift graph Sh_n^k has the property that the independence number of the subgraph induced by G satisfies $\alpha(\text{Sh}_n^k[G]) \geq (\frac{1}{2} - \varepsilon) |G|$, where $\varepsilon \rightarrow 0$ as $k \rightarrow \infty$. In this note we prove that for $k = 2$ and $n \rightarrow \infty$ there are graphs $G \subseteq \binom{[n]}{2}$ with $\alpha(\text{Sh}_n^2[G]) \leq (\frac{1}{4} + o(1)) |G|$, and $\frac{1}{4}$ is best possible. We also consider a related problem for infinite shift graphs.

Mathematics Subject Classifications: 05C69, 05C63

1 Introduction

For $n > k \in \mathbb{N}$ the shift graph Sh_n^k with

$$V(\text{Sh}_n^k) = \{(x_1, \dots, x_k) : 1 \leq x_1 < \dots < x_k \leq n\}$$

is a graph in which two vertices $\mathbf{x} = (x_1, \dots, x_k)$ and $\mathbf{y} = (y_1, \dots, y_k)$ are adjacent if $x_i = y_{i+1}$ for all $i \in \{1, \dots, k-1\}$ (or $y_i = x_{i+1}$ for all $i \in \{1, \dots, k-1\}$). Shift graphs were introduced by Erdős and Hajnal [3],[4] and are standard examples of graphs with large chromatic number and large odd girth. More precisely, while the odd girth of Sh_n^k is

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$2k + 1$, they proved* that Sh_n^k has chromatic number $(1 + o(1)) \log^{(k-1)} n$, where $\log^{(k-1)}$ stands for $k - 1$ times iterated $\log_2$.

Shift graphs have another interesting property: For each finite set $G \subseteq V(\text{Sh}_n^k)$ the induced subgraph $\text{Sh}_n^k[G]$ has a relatively large independent set with respect to $|G|$. In other words, the property “having a large independent subset” is hereditary for Sh_n^k . Namely, for

$$\alpha_n^k = \min \left\{ \frac{\alpha(\text{Sh}_n^k[G])}{|G|} : \emptyset \neq G \subseteq V(\text{Sh}_n^k) \right\}, \quad (1)$$

Erdős, Hajnal and Szemerédi [5, Theorem 1] proved the following.

Theorem 1 (Erdős, Hajnal, Szemerédi). *For positive integers $k < n$*

$$\alpha_n^k \geq \frac{1}{2} - \frac{1}{k}.$$

As for the upper bound, for $n \geq 2k + 1$ the shift graph Sh_n^k contains an odd cycle and so $\alpha_n^k < 1/2$. Therefore, Theorem 1 yields a lower bound which for large values of k is essentially optimal.

Nevertheless, determining the values of α_n^k for fixed k and large n seems to represent an interesting and non-trivial problem. We will concentrate our attention on the case $k = 2$. In this case the bound from Theorem 1 is not optimal, as we observe that $\alpha_n^2 \geq 1/4$ for all n , and prove a matching upper bound.

Theorem 2. $\lim_{n \rightarrow \infty} \alpha_n^2 = \frac{1}{4}$.

In [2], Czipser, Erdős and Hajnal proved that the densest independent set of the infinite graph $\text{Sh}_{\mathbb{N}}^2$ has density $1/4$ (see Section 3 for precise formulation). We complement their result by showing that the infinite shift graph $\text{Sh}_{\mathbb{N}}^2$ does not have a similar hereditary property, i.e., there exists $G \subseteq V(\text{Sh}_{\mathbb{N}}^2)$ such that any independent set in $\text{Sh}_{\mathbb{N}}^2[G]$ has density zero in G (see Theorem 7).

2 Proof of Theorem 2

Note that $\alpha_n^2 = \min \left\{ \frac{\alpha(\text{Sh}_n^2[G])}{|G|} : \emptyset \neq G \subseteq V(\text{Sh}_n^2) \right\}$ is a nonincreasing positive sequence, so the sequence $\{\alpha_n^2\}$ has a limit. Additionally, we will often view $G \subseteq V(\text{Sh}_n^2)$ as a graph with $V(G) = [n]$ and set of edges equal to G . Subsequently $|G|$ will denote both a size of G as a subset of $V(\text{Sh}_n^2)$, and the number of edges in G when it is viewed as a graph.

*In [4] authors considered infinite graphs, however their proof can be adapted for finite case (see [1] and [6] for more detailed description).

2.1 Lower bound

We first show that the value of the limit in Theorem 2 is at least $1/4$.

Claim 3. *For every set $G \subseteq V(\text{Sh}_n^2)$ we have $\alpha(\text{Sh}_n^2[G]) \geq \frac{1}{4}|G|$.*

Proof. Let $G \subseteq V(\text{Sh}_n^2)$ be given. Consider a random colouring $c : [n] \rightarrow \{r, b\}$ such that every $i \in [n]$ is coloured red/blue with probability $1/2$ independently of other elements of $[n]$.

Let G_c be a random subset of G defined by

$$G_c = \{(i, j) \in G : i < j, c(i) = b, c(j) = r\}.$$

Then such G_c is always an independent set in Sh_n^2 . Moreover, $\mathbb{P}(e \in G_c) = \frac{1}{4}$ for every $e \in G$, and so $\mathbb{E}(|G_c|) = \frac{1}{4}|G|$. Therefore $\alpha(\text{Sh}_n^2[G]) \geq \frac{1}{4}|G|$. $\square$

2.2 Upper bound

We now proceed and prove the upper bound

$$\lim_{n \rightarrow \infty} \alpha_n^2 \leq \frac{1}{4}. \quad (2)$$

In what follows for every $\varepsilon > 0$, integer d satisfying $\frac{3+\ln d}{4d} \leq \frac{\varepsilon}{2}$, and for every integer $n \geq n_0(\varepsilon, d)$ that is a multiple of 2^d , we will construct a graph $G_\varepsilon(n, d) \subseteq V(\text{Sh}_n^2)$ with

$$\alpha(\text{Sh}_n^2[G_\varepsilon(n, d)]) \leq \left(\frac{1}{4} + \varepsilon\right) |G_\varepsilon(n, d)|.$$

To be more precise, for such ε and d we inductively build $G_\varepsilon(n, d)$ satisfying

$$\frac{\alpha(\text{Sh}_n^2[G_\varepsilon(n, d)])}{|G_\varepsilon(n, d)|} \leq \left(\frac{1}{4} + \frac{3 + \ln d}{4d} + \frac{\varepsilon}{2}\right). \quad (3)$$

Since $\{\alpha_n^2\}$ is nonincreasing, (3) implies that $\lim_{n \rightarrow \infty} \alpha_n^2 \leq 1/4 + \varepsilon$, which subsequently implies (2) by letting $\varepsilon \rightarrow 0$.

While constructing $G_\varepsilon(n, d)$ we will use random bipartite graphs. Recall that if G is a graph and $X, Y \subseteq V(G)$ then $G[X, Y]$ is a graph consisting of edges of G with one vertex in X and another in Y . Finally let $e_G(X, Y) = |E(G[X, Y])|$ and we will omit subscript when G is obvious from the context. The following claim can be easily verified by considering a random graph and so the proof of Claim 4 is postponed to Appendix.

Claim 4. *For $\varepsilon > 0$ and $d \in \mathbb{N}$ there is $n_0 = n_0(\varepsilon, d)$ such that for all $n \geq n_0$ that are divisible by 2^d the following holds. Let $[n] = S \cup L$, where $S = \{1, \dots, \frac{n}{2}\}$ and $L = [n] \setminus S$. There exists a bipartite graph $B_\varepsilon(n, d)$ with bipartition $V(B_\varepsilon(n, d)) = S \sqcup L$ such that*

$$(i) \quad |B_\varepsilon(n, d)| = \frac{n^2}{2^{d+1}}.$$

(ii) for all $X \subseteq S$ and $Y \subseteq L$

$$e(X, Y) = \frac{1}{2^{d-1}} |X| |Y| \pm \frac{\varepsilon n^2}{2^{d+2}}.$$

Construction of $G_\varepsilon(n, d)$.

Definition 5. For every even n let $G_\varepsilon(n, 1)$ be such that

$$G_\varepsilon(n, 1) = \{(i, j) : 1 \leq i \leq \frac{n}{2} < j \leq n\},$$

i.e., $G_\varepsilon(n, 1)$ is a complete balanced bipartite graph.

For $d \in \mathbb{N}$ define graph $G_\varepsilon(n, d)$ recursively for all sufficiently large[†] n such that $2^d | n$. Let $[n] = S \cup L$, where $S = \{1, \dots, \frac{n}{2}\}$ and $L = [n] \setminus S$. Then define

$$G_\varepsilon(n, d) = G_\varepsilon(S, d-1) \cup G_\varepsilon(L, d-1) \cup B_\varepsilon(n, d),$$

where $G_\varepsilon(S, d-1) = G_\varepsilon(\frac{n}{2}, d-1)$, $V(G_\varepsilon(L, d-1)) = L$ and $G_\varepsilon(L, d-1) \cong G_\varepsilon(\frac{n}{2}, d-1)$, and $B_\varepsilon(n, d)$ is a graph guaranteed by Claim 4.

To summarize, every $G_\varepsilon(n, d) = G$ satisfies the following properties (with $S_n = \{1, \dots, \frac{n}{2}\}$ and $L_n = [n] \setminus S_n$):

(i) $e_G(S_n, L_n) = \frac{n^2}{2^{d+1}}.$

(ii) for all $X \subseteq S_n$ and $Y \subseteq L_n$

$$e(X, Y) = \frac{1}{2^{d-1}} |X| |Y| \pm \frac{\varepsilon n^2}{2^{d+2}}.$$

(iii) $G[S_n] \cong G[L_n] = G_\varepsilon(\frac{n}{2}, d-1)$

Using properties (i) and (iii) and induction on d it is easy to verify that for all $d \in \mathbb{N}$ and n divisible by 2^d

$$|G_\varepsilon(n, d)| = d \frac{n^2}{2^{d+1}}. \quad (4)$$

We will now proceed with proving (3). First let $G \subseteq V(\text{Sh}_n^2)$ and let $I \subseteq G$ be an independent set in Sh_n^2 . In other words there is no $1 \leq i < j < k \leq n$ with both (i, j) and (j, k) in I . One can observe that for each such $I \subseteq G$ there exists a 2-colouring $c : [n] \rightarrow \{r, b\}$ with $c(i) = r$ and $c(j) = b$ whenever $(i, j) \in I$, and then

$$I \subseteq G_c = \{(x, y) \in G : x < y, c(x) = b, c(y) = r\}. \quad (5)$$

[†] $n \geq 2^i n_0(\varepsilon, d-i)$ for all $i \in \{0, 1, \dots, d-2\}$, where $n_0(\varepsilon, d-i)$ is the number provided by Claim 4.

Therefore, in order to prove (3) we will show that for $G = G_\varepsilon(n, d)$ and any $c : [n] \rightarrow \{r, b\}$

$$\frac{|G_c|}{|G|} \leq \frac{1}{4} + \frac{3 + \ln d}{4d} + \frac{\varepsilon}{2}. \quad (6)$$

For the rest of our calculation let ε be fixed. We will now prove (6) by induction on d . In order to make use of recursive structure of $G_\varepsilon(n, d)$ we will prove a version of (6) with an additional assumption that $|\{i : c(i) = b\}| = \alpha n$.

To that end for $d \in \mathbb{N}$, $\alpha \in [0, 1]$ and $n \geq n_0(\varepsilon, d)$ let

$$f_d^\alpha(n) = d \cdot \max_c \left\{ \frac{|G_c|}{|G|} : G = G_\varepsilon(n, d), |\{i : c(i) = b\}| = \alpha n \right\}. \quad (7)$$

We will prove the following estimate on $f_d^\alpha(n)$.

Claim 6. *For every $d \in \mathbb{N}$, $\alpha \in [0, 1]$ and $n \geq n_0(\varepsilon, d)$*

$$f_d^\alpha(n) \leq (d + 3)(\alpha - \alpha^2) + \frac{1}{4} \ln d + \frac{d\varepsilon}{2}.$$

From (7) it follows that for $G = G_\varepsilon(n, d)$ and any colouring c we have

$$\frac{|G_c|}{|G|} \leq \max_{\alpha \in [0, 1]} \frac{f_d^\alpha(n)}{d}.$$

Then by Claim 6 we get

$$\frac{|G_c|}{|G|} \leq \frac{1}{4} \frac{d + 3}{d} + \frac{\ln d}{4d} + \frac{\varepsilon}{2},$$

establishing (6) and (3). Hence it remains to prove Claim 6 in order to finish the proof of the upper bound.

Proof of Claim 6. We prove a slightly stronger inequality for all $n \geq n_0(\varepsilon, d)$

$$f_d^\alpha(n) \leq (d + 3)(\alpha - \alpha^2) + \frac{1}{4} \sum_{i=3}^{d+1} \frac{1}{i} + \frac{d\varepsilon}{2}. \quad (8)$$

The proof is by induction on d . For $d = 1$ recall that $G = G_\varepsilon(n, 1)$ is a complete bipartite graph between S_n and L_n . Let $c : [n] \rightarrow \{r, b\}$ be such that for $B = \{i : c(i) = b\}$ we have $|B| = \alpha n$. Then in view of (5) the maximum value of $|G_c|$ is achieved when $B = [\alpha n]$ and so

$$f_1^\alpha(n) = \begin{cases} 2\alpha, & \alpha \in [0, \frac{1}{2}] \\ 2 - 2\alpha, & \alpha \in [\frac{1}{2}, 1]. \end{cases}$$

Now it is easy to verify that $f_1^\alpha \leq 4(\alpha - \alpha^2)$ for all $\alpha \in [0, 1]$, establishing (8) in the case $d = 1$.

To prove inductive step let $G = G_\varepsilon(n, d)$ and let $c : [n] \rightarrow \{r, b\}$ be such that for $B = \{i : c(i) = b\}$ we have $|B| = \alpha n$. As before, let $S = \{1, \dots, \frac{n}{2}\}$ and $L = [n] \setminus S$. Let

B_S , B_L , R_S and R_L denote the set of blue and red vertices in S and L respectively. We will further refine our analysis by assuming that $|B_S| = x\frac{n}{2}$ with some $x \in [0, 2\alpha]$. Since $|B_S| + |R_S| = \frac{n}{2}$, $|B_S| + |B_L| = |B| = \alpha n$, and $|B_L| + |R_L| = \frac{n}{2}$, we have $|R_S| = (1-x)\frac{n}{2}$, $|B_L| = (2\alpha-x)\frac{n}{2}$, and consequently $|R_L| = (1-2\alpha+x)\frac{n}{2}$ (see Figure 1). Then

$$|G_c| = |G_c[S]| + |G_c[L]| + e_G(B_S, R_L). \quad (9)$$

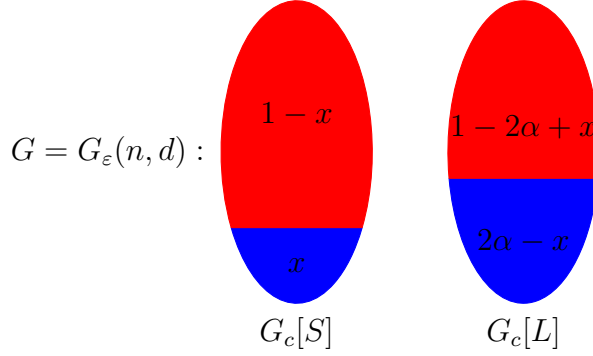


Figure 1: Proportions of red and blue vertices in $G_c[S]$ and $G_c[L]$.

Now, by (iii) $G[S] = G_\epsilon(\frac{n}{2}, d-1)$ and we assumed $|B_S| = x\frac{n}{2}$, so

$$|G_c[S]| \stackrel{(7)}{\leq} \frac{f_{d-1}^x(\frac{n}{2})}{d-1} |G[S]| \stackrel{(4)}{=} \frac{n^2}{2^{d+2}} f_{d-1}^x \left(\frac{n}{2} \right). \quad (10)$$

Similarly, since $|B_L| = (2\alpha-x)\frac{n}{2}$ we have

$$|G_c[L]| \leq \frac{n^2}{2^{d+2}} f_{d-1}^{2\alpha-x} \left(\frac{n}{2} \right). \quad (11)$$

And finally, since $G = G_\epsilon(n, d)$,

$$e_G(B_S, R_L) \stackrel{(ii)}{\leq} \frac{1}{2^{d-1}} |B_S| |R_L| + \frac{\epsilon n^2}{2^{d+2}} = \frac{n^2}{2^{d+1}} \left(x(1-2\alpha+x) + \frac{\epsilon}{2} \right). \quad (12)$$

Combining (9) with (10), (11), and (12) we obtain

$$|G_c| \leq \frac{n^2}{2^{d+1}} \left(\frac{1}{2} \left(f_{d-1}^x \left(\frac{n}{2} \right) + f_{d-1}^{2\alpha-x} \left(\frac{n}{2} \right) \right) + x(1-2\alpha+x) + \frac{\epsilon}{2} \right).$$

Finally, $|G| = |G_\epsilon(n, d)| \stackrel{(4)}{=} d \frac{n^2}{2^{d+1}}$ and so by (7) we deduce that

$$f_d^\alpha(n) \leq \max_{x \in \mathbb{R}} \left\{ \frac{1}{2} \left(f_{d-1}^x \left(\frac{n}{2} \right) + f_{d-1}^{2\alpha-x} \left(\frac{n}{2} \right) \right) + x(1-2\alpha+x) + \frac{\epsilon}{2} \right\}. \quad (13)$$

The last inequality allows us to incorporate induction hypothesis. In particular, by induction hypothesis we have

$$f_{d-1}^x\left(\frac{n}{2}\right) \leq (d+2)(x-x^2) + \frac{1}{4} \sum_{i=3}^d \frac{1}{i} + \frac{(d-1)\varepsilon}{2},$$

$$f_{d-1}^{2\alpha-x}\left(\frac{n}{2}\right) \leq (d+2)(2\alpha-x)(1-2\alpha+x) + \frac{1}{4} \sum_{i=3}^d \frac{1}{i} + \frac{(d-1)\varepsilon}{2},$$

and these two inequalities together with (13), after some simple but tedious algebraic manipulations yield

$$f_d^\alpha(n) \leq \max_{x \in \mathbb{R}} \left\{ -(d+1)x^2 + (1+2\alpha(d+1))x + (d+2)(\alpha-2\alpha^2) + \frac{1}{4} \sum_{i=3}^d \frac{1}{i} + \frac{d\varepsilon}{2} \right\}.$$

In other words $f_d^\alpha(n) \leq \max_{x \in \mathbb{R}} \{g(x)\}$, where $g(x) = ax^2 + bx + c$ with $a = -(d+1)$. Since $a < 0$ we have $\max_{x \in \mathbb{R}} g(x) = g(\frac{-b}{2a}) = c - \frac{b^2}{4a}$. Therefore after another set of algebraic manipulations we obtain

$$f_d^\alpha(n) \leq \max_{x \in \mathbb{R}} \{g(x)\} \leq (d+3)(\alpha-\alpha^2) + \frac{1}{4} \sum_{i=3}^{d+1} \frac{1}{i} + \frac{d\varepsilon}{2},$$

finishing the proof of the inductive step and Claim 6. □

3 Infinite graphs

Recall that Theorem 2 states

$$\lim_{n \rightarrow \infty} \min \left\{ \frac{\alpha(\text{Sh}_n^2[G])}{|G|} : \emptyset \neq G \subseteq V(\text{Sh}_n^2) \right\} = \frac{1}{4}. \quad (14)$$

On the other hand, considering $I = \{(i, j) : 1 \leq i \leq \frac{n}{2} < j \leq n\}$ we clearly have $\alpha(\text{Sh}_n^2) \geq \lfloor \frac{n^2}{4} \rfloor$. Moreover $\lfloor \frac{n^2}{4} \rfloor$ is optimal, since any graph $G \subseteq V(\text{Sh}_n^2)$ with $|G| \geq \lfloor \frac{n^2}{4} \rfloor + 1$ contains a triangle and hence such G is not an independent set in Sh_n^2 . Therefore,

$$\lim_{n \rightarrow \infty} \frac{\alpha(\text{Sh}_n^2)}{|\text{Sh}_n^2|} = \frac{1}{2}. \quad (15)$$

It may be interesting to note that infinite version of (15) was considered by Czipser, Erdős and Hajnal [2] who proved that if I is independent set in countable shift graph $\text{Sh}_{\mathbb{N}}^2$, then the density of I does not exceed $1/4$, i.e.

$$\liminf_{n \rightarrow \infty} \frac{|I \cap \binom{[n]}{2}|}{\binom{n}{2}} \leq \frac{1}{4}. \quad (16)$$

(Here $\frac{1}{4}$ is clearly optimal, since $I = \{(i, j) : i < j, i \text{ odd}, j \text{ even}\}$ is independent in $\text{Sh}_{\mathbb{N}}^2$.)

To complete this discussion we provide an infinite variant of (14).

Theorem 7. *There is $G \subseteq V(\text{Sh}_{\mathbb{N}}^2)$ such that if I is an independent set in $\text{Sh}_{\mathbb{N}}^2[G]$, then*

$$\liminf_{n \rightarrow \infty} \frac{|I \cap \binom{[n]}{2}|}{|G \cap \binom{[n]}{2}|} = 0.$$

Proof. Consider an infinite ordered tree G with $V(G) = \mathbb{N}$, and with vertices labeled v_i^j , where j denotes the “level” L_j that vertex v_i^j belongs to and i denotes the order in which vertices are listed on the level.

Consider a labeling of vertices of G by integers satisfying $v_i^j < v_{i'}^{j'}$ if $j < j'$ and $v_i^j < v_{i'}^j$ if $i < i'$ such that for all v_i^j the finite set $N^+(v_i^j)$ of all children of v_i^j forms an interval (and these intervals on the level L_{j+1} follow the order of their parents on L_j , see Figure 2). Finally we will assume that for all v_i^j

$$|N^+(v_i^j)| \geq 2^j \sum_{v < v_i^j} |N^+(v)|. \quad (17)$$

Now, let $I \subseteq G$ be an infinite independent set in $\text{Sh}_{\mathbb{N}}^2$ and let $(v_k^{j-1}, v_i^j) \in I$, where v_k^{j-1} and v_i^j are parent and child respectively. Let $w = \max\{N^+(v_i^j)\}$ be the largest son of v_i^j and let $W = \{1, \dots, w\}$ (see Figure 2). Then

$$G[W] = \bigcup_{v \leq v_i^j} \{(v, u) : u \in N^+(v)\}. \quad (18)$$

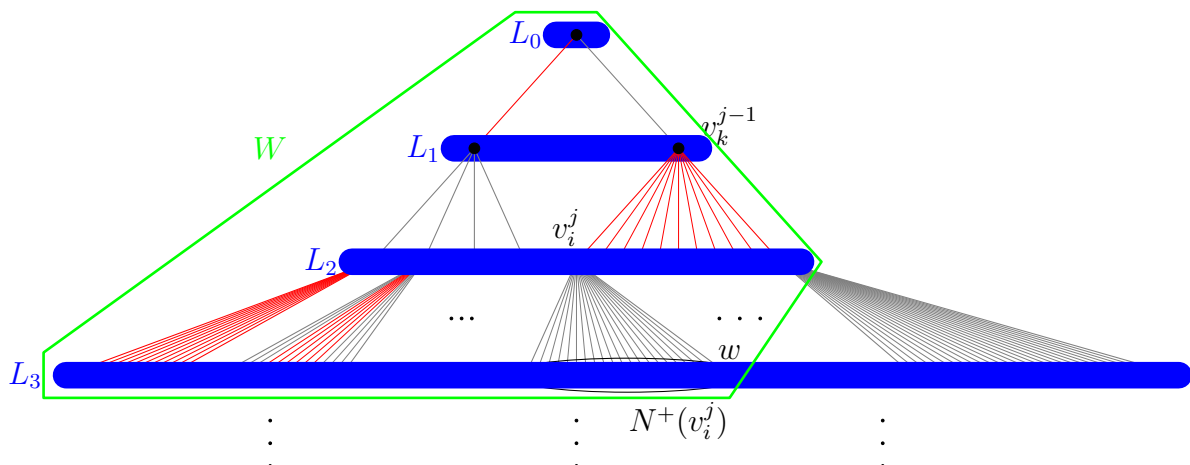


Figure 2: Infinite tree G , vertices are ordered top to bottom, left to right. Edges of I are labeled with red.

In particular in view of (17)

$$|G[W]| \geq |\{(v_i^j, u) : u \in N^+(v)\}| = |N^+(v_i^j)|. \quad (19)$$

On other hand, since $(v_k^{j-1}, v_i^j) \in I$ and I is independent set in $\text{Sh}_{\mathbb{N}}^2$, the set I does not contain any other edge incident to v_i^j . Consequently,

$$|I[W]| \leq \left| \bigcup_{v < v_i^j} \{(v, u) : u \in N^+(v)\} \right| \stackrel{(17)}{\leq} 2^{-j} |N^+(v_i^j)|. \quad (20)$$

In view of (19) and (20) we have $|I[W]|/|G[W]| \leq \frac{1}{2^j}$. Now, since I is infinite there are edges $(v_k^{j-1}, v_i^j) \in I$ with sufficiently large j , hence the ratio $|I[W]|/|G[W]|$ can be made arbitrary small, finishing the proof. $\square$

4 Concluding remarks

In [5] it was proved[‡] that for any n, k

$$\alpha_n^k \geq \begin{cases} \frac{1}{2} - \frac{1}{k}, & \text{if } k \text{ is even,} \\ \frac{1}{2} - \frac{1}{2k}, & \text{if } k \text{ is odd.} \end{cases} \quad (21)$$

It remains an open problem to determine for any $k \geq 3$ the exact value of $\lim_{n \rightarrow \infty} \alpha_n^k$. For $k = 4$ we were able to improve the constant in the lower bound (21) from $\frac{1}{4}$ to $\frac{3}{8}$ [§] and for $k = 3$ we believe that estimate in (21) is sharp.

Problem 8. Show that $\lim_{n \rightarrow \infty} \alpha_n^3 = \frac{1}{3}$.

Finally, all of the results in this paper can be reformulated in terms of subgraphs with no increasing paths of length two. For instance, Theorem 2 implies that for any $\varepsilon > 0$ there exists an vertex-ordered graph G such that if $G' \subseteq G$ with $|G'| \geq (\frac{1}{4} + \varepsilon) |G|$, then G' contains an increasing path of length two, i.e. there are $i < j < k$ with $(i, j), (j, k) \in G'$. One can ask similar questions for longer increasing paths.

Problem 9. For any $\varepsilon > 0$ does there exist an ordered graph G such that if $G' \subseteq G$ with $|G'| \geq (\frac{1}{3} + \varepsilon) |G|$, then G' contains an increasing path of length three?

Note that in regards to Problem 9, one can consider a random coloring c of $V(G)$ with colors $\{0, 1, 2\}$ and define G' to be the collection of all $(i, j) \in E(G)$ with $i < j$ and $c(i) < c(j)$. Then such G' on average contains $\frac{1}{3}|G|$ edges and has no increasing paths of length three, motivating the constant $\frac{1}{3}$ in the problem.

[‡]the result follows from the proof of Theorem 1 in [5]

[§] $\alpha(\text{Sh}_n^4[G]) \geq \frac{3}{8}|G|$ can be proved by considering a random colouring $c : [n] \rightarrow \{0, 1\}$ and forming an independent set in Sh_n^4 by taking hyperedges of G of form 1000, 1110, or $x01y$ for some $x, y \in \{0, 1\}$.

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Appendix

Proof of Claim 4. Let $B_\varepsilon(n, d) = G$, where G is a random graph between S and L obtained by selecting a random subset of size $\frac{n^2}{2^{d+1}}$ without replacement from $K_{S,L}$ (complete bipartite graph between S and L). Then G satisfies (i) and we will show that G satisfies (ii) almost surely.

For every $X \subseteq S$ and $Y \subseteq L$, $e(X, Y) = e_G(X, Y)$ is distributed as a hypergeometric random variable $H\left(\frac{n^2}{4}, \frac{n^2}{2^{d+1}}, |X||Y|\right)$ with expectation $\frac{1}{2^{d-1}}|X||Y|$. Let $B_{X,Y}$ be the event that

$$\left|e_G(X, Y) - \frac{1}{2^{d-1}}|X||Y|\right| > \frac{\varepsilon n^2}{2^{d+2}},$$

i.e., $B_{X,Y}$ is the event that (ii) fails for given X and Y .

We will use a concentration inequality for hypergeometric random variables (this version is a corollary of Theorem 2.10 and inequalities (2.5),(2.6) of Janson, Luczak, Rucinski [7]).

Theorem 10. *Let $Z \sim H(N, m, k)$ be a hypergeometric random variable with the expectation $\mu = \frac{mk}{N}$, then for $t \geq 0$*

$$\mathbb{P}(|Z - \mu| > t) \leq 2 \exp\left(\frac{-t^2}{2(\mu + t/3)}\right).$$

For a given $X \subseteq L$ and $Y \subseteq R$, as a consequence of Theorem 10 with $Z = e_G(X, Y)$, $t = \frac{\varepsilon n^2}{2^{d+2}}$ and $\mu = \frac{1}{2^{d-1}}|X||Y| \leq \frac{n^2}{2^{d+1}}$ we get

$$\mathbb{P}(B_{X,Y}) = e^{-\Omega(n^2)},$$

where constant in $\Omega()$ term depends on ε and d only. Therefore,

$$\mathbb{P}\left(\bigcup_{X,Y} B_{X,Y}\right) \leq \sum_{X,Y} \mathbb{P}(B_{X,Y}) \leq 2^n e^{-\Omega(n^2)} = o(1).$$

In particular, $\mathbb{P}(G \text{ satisfies (ii)}) = \mathbb{P}\left(\bigcap_{X,Y} \overline{B_{X,Y}}\right) = 1 - o(1)$. Hence, G almost surely satisfies (ii). $\square$