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More on zeros and approximation of the Ising partition function

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Abstract
We consider the problem of computing the partition function $\sum_x e^{f(x)}$, where $f : \{-1, 1\}^n \longrightarrow \mathbb{R}$ is a quadratic or cubic polynomial on the Boolean cube $\{-1, 1\}^n$. In the case of a quadratic polynomial f, we show that the partition function can be approximated within relative error $0 < \epsilon < 1$ in quasi-polynomial $n^{O(\ln n - \ln \epsilon)}$ time if the Lipschitz constant of the non-linear part of f with respect to the ℓ^1 metric on the Boolean cube does not exceed $1 - \delta$, for any $\delta > 0$, fixed in advance. For a cubic polynomial f, we get the same result under a somewhat stronger condition. We apply the method of polynomial interpolation, for which we prove that $\sum_x e^{\tilde{f}(x)} \neq 0$ for complex-valued polynomials $\tilde{f}$ in a neighborhood of a real-valued f satisfying the above mentioned conditions. The bounds are asymptotically optimal. Results on the zero-free region are interpreted as the absence of a phase transition in the Lee–Yang sense in the corresponding Ising model. The novel feature of the bounds is that they control the total interaction of each vertex but not every single interaction of sets of vertices.

1. Introduction and main results

The Ising model is one of the oldest, most famous and most studied models in statistical physics; see [7] for a thorough introduction, description, results and references. In this paper, we look at the computational complexity and complex zeros of the partition function in the Ising model. This is a classical and also currently very active area of research; see [6], [9], [10], [13], [16], [14], [15], [18], [19] and [22] for some recent results.

Formally, the partition function we work with is described as follows. Let $\{-1, 1\}^n$ be the n-dimensional Boolean cube of all n-vectors $x = (\xi_1, \ldots, \xi_n)$, where $\xi_i = \pm 1$ for $i = 1, \ldots, n$, and let $f : \{-1, 1\}^n \longrightarrow \mathbb{C}$ be a function. We define the *partition function* of f as

$$S\left(e^f\right) = \sum_{x \in \{-1, 1\}^n} e^{f(x)}.$$

In this paper, we consider the cases of a quadratic

$$f(x) = \sum_{1 \leq i < j \leq n} a_{ij}\xi_i\xi_j + \sum_{i=1}^{n} b_i\xi_i \quad \text{for} \quad x = (\xi_1, \ldots, \xi_n)$$

or cubic

$$f(x) = \sum_{1 \le i < j < k \le n} c_{ijk} \xi_i \xi_j \xi_k + \sum_{1 \le i < j \le n} a_{ij} \xi_i \xi_j + \sum_{i=1}^{n} b_i \xi_i \quad \text{for} \quad x = (\xi_1, \dots, \xi_n)$$

polynomial f, and we are interested in computing (approximating) $\mathcal{S}(e^f)$ efficiently.

The case of a quadratic polynomial f is the best studied. Often, there is an underlying graph G with vertices numbered $1, \dots, n$ and the property that $a_{ij} \ne 0$ if and only if i and j span an edge of G. In this case, the variable $\xi_i = \pm 1$ is interpreted as the *spin* of a vertex i and $-f(x)$ as the *energy* of the *configuration* $x = (\xi_1, \dots, \xi_n)$. The coefficients a_{ij} describe the *interactions* of vertices i and j: the interaction is *ferromagnetic* if $a_{ij} > 0$ and *antiferromagnetic* if $a_{ij} < 0$. The coefficients b_i describe the *external field*; see [7] for a thorough discussion.

Via the method of polynomial interpolation [3], we relate the computational complexity to the absence of complex zeros of $\mathcal{S}(e^f)$ in the vicinity of real coefficients a_{ij}, b_i and c_{ijk}. For quadratic polynomials, we prove the following main result.

Theorem 1. *Let*

$$f(x) = \sum_{1 \le i < j \le n} a_{ij} \xi_i \xi_j + \sum_{i=1}^{n} b_i \xi_i \quad \text{for} \quad x = (\xi_1, \dots, \xi_n).$$

Suppose that for some $0 < \delta < 1$, we have

$$\sum_{j: \, j \ne i} \left| \Re \, a_{ij} \right| \le 1 - \delta, \quad \sum_{j: \, j \ne i} \left| \Im \, a_{ij} \right| \le \frac{\delta^2}{10} \quad \text{and} \quad \left| \Im \, b_i \right| \le \frac{\delta^2}{10}$$

for $i = 1, \dots, n$. Then

$$\sum_{x \in \{-1, 1\}^n} e^{f(x)} \ne 0.$$

A remark regarding notation: we treat indices i, j in a_{ij} as an unordered pair, so a_{ij} is the coefficient of the monomial $\xi_i \xi_j$ in the quadratic polynomial f, and the sum $\sum_{j: \, j \ne i}$ accounts for all coefficients a_{ij} of the monomials containing ξ_i. For a complex number $z = x + y\sqrt{-1}$, we denote by $\Re \, z = x$ and $\Im \, z = y$ the real and imaginary parts of z, respectively.

From Theorem 1, the by-now-standard polynomial interpolation argument (see [3], [9], [14], [15], [17]) produces an algorithm for approximating the partition function $\mathcal{S}(e^f)$ when the coefficients a_{ij} and b_i are real and satisfy the condition

$$\sum_{j: \, j \ne i} |a_{ij}| \le 1 - \delta \quad \text{for} \quad i = 1, \dots, n, \tag{1}$$

where $0 < \delta < 1$ is fixed in advance. As there is no restriction on b_i, the sum $\mathcal{S}(e^f)$ can be exponentially large in $|b_i|$. To avoid dealing with exponentially large numbers, we assume that we are provided with numbers e^{b_i} for $i = 1, \dots, n$. Then the complexity of the algorithm is quasi-polynomial: we approximate $\mathcal{S}(e^f)$ within relative error $0 < \epsilon < 1$ in $n^{O(\ln n - \ln \epsilon)}$ time. As Guus Regts explained to the authors, the approach of [17] makes our algorithm genuinely polynomial, as opposed to quasi-polynomial, provided the degree of the underlying graph G is fixed in advance.

We describe the algorithm in Section 4 and prove Theorem 1 in Section 2. Next, we relate Theorem 1 to what is known about the partition function in the Ising model.

1.1. The bounds for the zero-free region are asymptotically optimal

Let G be a graph with vertices $1, \ldots, n$. For a real number a, let us choose

$$a_{ij} = \begin{cases} a & \text{if } \{i, j\} \text{ is an edge of } G \\ 0 & \text{otherwise} \end{cases} \tag{2}$$

and let us choose $b_i = b$ for some $b \in \mathbb{C}$ for all $i = 1, \ldots, n$. Let us fix some positive integer $\Delta \geq 3$ and choose either

$$a = \frac{1}{2} \ln \frac{\Delta}{\Delta - 2}$$

(all interactions are ferromagnetic) or

$$a = \frac{1}{2} \ln \frac{\Delta - 2}{\Delta}$$

(all interactions are antiferromagnetic). We consider the partition function $\mathcal{S}\left(e^f\right)$ as a function of a complex parameter b. It is known that for a fixed Δ, as n grows and G ranges over all graphs with maximum degree Δ of a vertex, the zeros of the univariate function $b \longmapsto \mathcal{S}\left(e^f\right)$ with either choice of a can get arbitrarily close to $b = 0$; see [2], [1], [18], [6]. We have

$$\sum_{j:\ j \neq i} \left| \Re\, a_{ij} \right| \leq \frac{\Delta}{2} \ln \frac{\Delta}{\Delta - 2}. \tag{3}$$

The right-hand side approaches 1 as $\Delta \longrightarrow \infty$, which shows that '1' in the '$1 - \delta$' bound of Theorem 1 cannot be replaced by a larger number.

1.2. The bounds for approximation are asymptotically optimal in the antiferromagnetic case

As in Section 1.1, let G be a graph of maximum degree $\Delta \geq 3$, let us choose all $b_i = 0$ for all $i = 1, \ldots, n$, and let us define a_{ij} by (2). Suppose that

$$a \ < \ \frac{1}{2} \ln \frac{\Delta - 2}{\Delta}.$$

It is shown in [20] and also in [8] that the problem of approximating $\mathcal{S}\left(e^f\right)$ is NP-hard under randomized reduction. Hence unless the computational complexity hierarchy collapses, we cannot approximate $\mathcal{S}\left(e^f\right)$ in quasi-polynomial time in the class of problems where all $b_i = 0$ and

$$\sum_{j:\ j \neq i} |a_{ij}| \ \leq \ 1 + \delta \quad \text{for} \quad i = 1, \ldots, n$$

for an arbitrarily small $\delta > 0$, fixed in advance.

1.3. The ferromagnetic case is special

Suppose that $a_{ij} \geq 0$ for all i, j and that $b_i = b$ for $i = 1, \ldots, n$ and some complex parameter b. Lee and Yang [12] showed that the zeros of the univariate function $b \longmapsto \mathcal{S}\left(e^f\right)$ lie on the line $\Re\, b = 0$. If the b_i are allowed to vary, then $\mathcal{S}\left(e^f\right) \neq 0$ as long as $\Re\, b_i > 0$ for $i = 1, \ldots, n$. The ferromagnetic case is also special from the complexity point of view: Jerrum and Sinclair [11] constructed a randomized polynomial time algorithm approximating $\mathcal{S}\left(e^f\right)$ when $b_i = b$ for some real b and $i = 1, \ldots, n$. Deterministic approximation algorithms of quasi-polynomial (genuinely polynomial, if the maximum degree Δ of the underlying graph G is fixed in advance) complexity are constructed in [15], assuming

that $\Re b_i \geq \delta > 0$ for some constant $\delta > 0$, fixed in advance, and $i = 1, \ldots, n$; see also Section 7.4 of [3]. The complexity status of the approximation problem in the ferromagnetic case of $a_{ij} \geq 0$ and $b_i = 0$ by a deterministic algorithm appears to be not known.

1.4. Correlation decay

Theorem 1 can be interpreted as saying that there is no phase transition in the Lee–Yang sense [21] provided the parameters a_{ij} and b_i are real and satisfy (1) for some $0 < \delta < 1$, fixed in advance. There is a related, although not identical, concept of phase transition, based on the disappearance of correlation decay; see [7]. If we choose $b_i = 0$ for all i and define a_{ij} as in (2), then the correlation decay occurs precisely in the interval

$$\frac{1}{2} \ln \frac{\Delta - 2}{\Delta} \; < \; a \; < \; \frac{1}{2} \ln \frac{\Delta}{\Delta - 2},$$

where $\Delta \geq 3$ is the largest degree of a vertex of G; see [22], [19], and [14].

We extend the bounds of Theorem 1 to cubic polynomials on the Boolean cube $\{-1, 1\}^n$.

Theorem 2. *Let*

$$f(x) = \sum_{1 \leq i < j < k \leq n} c_{ijk} \xi_i \xi_j \xi_k + \sum_{1 \leq i < j \leq n} a_{ij} \xi_i \xi_j + \sum_{i=1}^{n} b_i \xi_i \quad \text{for} \quad x = (\xi_1, \ldots, \xi_n).$$

Suppose that for some $0 < \delta < 1/2$, we have

$$\sum_{j,k:\; j,k \neq i} \left| \Re c_{ijk} \right| + \sum_{j:\; j \neq i} \left| \Re a_{ij} \right| \; \leq \; 1 - \delta, \qquad \sum_{j,k:\; j,k \neq i} \left| \Im c_{ijk} \right| + \sum_{j:\; j \neq i} \left| \Im a_{ij} \right| \; \leq \; \frac{\delta^2}{10}$$

$$and \quad \left| \Im b_i \right| \; \leq \; \frac{\delta^2}{10}$$

for $i = 1, \ldots, n$. Then

$$\sum_{x \in \{-1,1\}^n} e^{f(x)} \neq 0.$$

Similarly to the quadratic case, we treat indices i, j, k in c_{ijk} as unordered triples of distinct numbers, so that c_{ijk} is the coefficient of the monomial $\xi_i \xi_j \xi_k$ in the cubic polynomial f, and the sum $\sum_{j,k:\; j,k \neq i}$ accounts for all coefficients c_{ijk} of the monomials containing ξ_i. As in the case of quadratic polynomials, we obtain an algorithm of quasi-polynomial $n^{O(\ln n - \ln \epsilon)}$ complexity approximating

$$\mathcal{S}\left(e^f\right) = \sum_{x \in \{-1,1\}^n} e^{f(x)}$$

within relative error $0 < \epsilon < 1$, provided the coefficients c_{ijk}, a_{ij} and b_i are real and satisfy

$$\sum_{j,k:\; j,k \neq i} \left| c_{ijk} \right| + \sum_{j:\; j \neq i} \left| a_{ij} \right| \; \leq \; 1 - \delta \quad \text{for} \quad i = 1, \ldots, n. \tag{4}$$

The bounds here are asymptotically optimal, for the trivial reason that they are asymptotically optimal in the quadratic case when $c_{ijk} = 0$ for all $1 \leq i < j < k \leq n$.

Although Theorem 1 is a particular case of Theorem 2, we first prove Theorem 1 in Section 2 and then describe in Section 3 how to adjust the proof of Theorem 1 to obtain Theorem 2. We think that this way it is easier to see the idea of the proofs. In Section 3, we also mention where our method falls short extending the bounds to polynomials of degree 4 and higher. We don't discuss the corresponding extension of the approximation algorithm of Section 4 from quadratic to cubic polynomials, as it is quite straightforward.

1.5. What's new

In the case of a quadratic polynomial, the bound (1) states that the Lipschitz constant of the non-linear part of $f : \{-1, 1\}^n \longrightarrow \mathbb{R}$ does not exceed $1 - \delta$ with respect to the ℓ^1 metric (which is twice the Hamming metric) on the Boolean cube $\{-1, 1\}^n$. This is where our approach to approximation differs from those of [9], [13], [14], [15], [19] and [22], which require a uniform bound on the strength of individual interactions a_{ij}, such as

$$\max_{i,j} |a_{ij}| \le \frac{1}{2} \ln \frac{\Delta}{\Delta - 2} \tag{5}$$

where Δ is the maximum degree of the underlying graph. As we remarked above, our results are asymptotically optimal, when $\Delta \longrightarrow \infty$. Generally, the condition (1) appears to be more robust than (5), as (1) is independent of the degree Δ and allows individual coefficients a_{ij} to be relatively large, as long as the sum for all interactions of any given vertex remains appropriately bounded. Of course, for any particular Δ, the conditions (1) and (5) are in general position, as it is easy to construct examples where one holds and the other is violated.

A novelty of our approach with respect to locating zero-free regions of $\mathcal{S}(e^f)$, compared to those of [9], [14], [6] and [18], is that we allow all parameters a_{ij} and b_i to vary: this concerns both Lee–Yang zeros [12], [6], [18] of $\mathcal{S}(e^f)$ as a function of b_i with fixed a_{ij} and Fisher zeros [14] of $\mathcal{S}(e^f)$ as a function of a_{ij} with b_i fixed. It appears that Theorem 1 is the first result establishing an asymptotically optimal zero-free region when the interactions a_{ij} are allowed to differ for different pairs $\{i, j\}$ and even to be of different signs, so that we have a mixture of ferromagnetic and antiferromagnetic interactions.

There is much less known in the case of cubic and higher degree polynomials on the Boolean cube; see though [16] and [15]. In particular, Theorem 2 might be the only result that provides asymptotically optimal bounds for cubic polynomials, although the optimality follows by a trivial reduction to the quadratic case. We note that in contrast to (1), the left-hand side of (4) is not the Lipschitz constant of the non-linear part of f but an upper bound on that Lipschitz constant.

For polynomials f of degree $d \ge 4$ on the Boolean cube, some, apparently non-optimal, estimates can be found in [4]. Roughly, estimates of [4] require the Lipschitz constant of f to be $O(1/\sqrt{d})$. It would be interesting to find out if one can get a bound uniform on d.

2. Proof of Theorem 1

2.1. Definitions and notation

We consider the Boolean cube $\{-1, 1\}^n$ of vectors $x = (\xi_1, \ldots, \xi_n)$, where $\xi_i = \pm 1$ for $i = 1, \ldots, n$. Let us choose a set $I \subset \{1, \ldots, n\}$ and numbers $\sigma_i \in \{-1, 1\}$ for $i \in I$. The set

$$F = \left\{ x \in \{-1, 1\}^n : \xi_i = \sigma_i \quad \text{for} \quad i \in I \right\}$$

is called a *face* of $\{-1, 1\}^n$. Any index $i \in \{1, \ldots, n\} \setminus I$ is called a *free index* of the face F, and the number $n - |I|$ of free indices is called the *dimension* of F and denoted $\dim F$. Indexes $i \in I$ are called *fixed*. For example, if $I = \emptyset$, then the face is the whole cube $\{-1, 1\}^n$; and if $I = \{1, \ldots, n\}$, then the face consists of a single point. Generally, a face of dimension k consists of 2^k points.

Let $g : \{-1, 1\}^n \longrightarrow \mathbb{C}$ be a function, and let $F \subset \{-1, 1\}^n$ be a face. We define the *partial sum of g on F* by

$$\mathcal{S}(g|F) = \sum_{x \in F} g(x).$$

We will use the following straightforward identity. Let F be a face of $\{-1, 1\}^n$ of dimension at least 1. Let i be a free index of F, and let $F^+ \subset F$ and $F^- \subset F$ be the faces obtained by fixing the i-th coordinate ξ_i of vectors $x \in F$ to 1 and -1, respectively. Hence $F^+ \cup F^- = F$, $F^+ \cap F^- = \emptyset$ and

$$\mathcal{S}(g|F) = \mathcal{S}(g|F^+) + \mathcal{S}(g|F^-). \tag{6}$$

Let us fix a real number $0 < \delta < 1$ and a vector $b = (b_1, \ldots, b_n)$ such that $|\Im b_i| \leq \delta^2/10$ for $i = 1, \ldots, n$. We denote by $\mathcal{U}(b; \delta)$ the set of all polynomials $f : \{-1, 1\}^n \longrightarrow \mathbb{C}$,

$$f(\xi_1, \ldots, \xi_n) = \sum_{1 \leq i < j \leq n} a_{ij}\xi_i\xi_j + \sum_{i=1}^n b_i\xi_i,$$

where

$$\sum_{j=1}^n |\Re a_{ij}| \leq 1 - \delta \quad \text{and} \quad \sum_{j=1}^n |\Im a_{ij}| \leq \frac{\delta^2}{10} \quad \text{for} \quad i = 1, \ldots, n.$$

We can view $\mathcal{U}(b; \delta)$ as a convex subset of $\mathbb{C}^{\binom{n}{2}}$ with a non-empty interior. In particular, for any $f_1, f_2 \in \mathcal{U}(b; \delta)$ and any $0 \leq \alpha \leq 1$, we have $f \in \mathcal{U}(b; \delta)$ for $f = \alpha f_1 + (1 - \alpha)f_2$. Our goal is to show that $\mathcal{S}(e^f) \neq 0$ for all $f \in \mathcal{U}(b; \delta)$.

In what follows, we view non-zero complex numbers $z \in \mathbb{C} \setminus \{0\}$ as vectors in the plane $\mathbb{R}^2 = \mathbb{C}$ and measure angles between them. We prove by induction on the dimension of a face $F \subset \{-1, 1\}^n$ that $\mathcal{S}(e^f|F) \neq 0$ and, moreover, if F^+ and F^- are faces of F defined above then the angle between complex numbers $\mathcal{S}(e^f|F^+) \neq 0$ and $\mathcal{S}(e^f|F^-) \neq 0$ is small.

We start with a simple geometric lemma.

Lemma 1. *Let $w_+, w_- \in \mathbb{C} \setminus \{0\}$ be non-zero numbers such that the angle between w_+ and w_- does not exceed some $0 \leq \theta < \pi$, and let $w = w_+ + w_-$.*

1. *We have $w \neq 0$ and, moreover,*

$$\frac{|w_+| + |w_-|}{|w|} \leq \frac{1}{\cos(\theta/2)}.$$

2. *We have*

$$\left| \Im \frac{w_+ - w_-}{w} \right| \leq \tan \frac{\theta}{2}.$$

Proof. Let v_+, v_- and v be the orthogonal projections of w_+, w_- and w, respectively, onto the bisector of the angle between w_+ and w_-. Then

$$|w| \geq |v| = |v_+ + v_-| = |v_+| + |v_-| \geq \cos(\theta/2)|w_+| + \cos(\theta/2)|w_-|$$

and Part 1 follows.

To prove Part 2, let $u_+ = w_+/w$ and $u_- = w_-/w$. Then

$$u_+ + u_- = 1$$

is real and the angle between u_+ and u_- does not exceed θ. Let

$$\alpha = \arg u_+ \quad \text{and} \quad \beta = -\arg u_-.$$

Without loss of generality, we may assume that $0 < \alpha, \beta < \pi/2$ and hence $\alpha + \beta \leq \theta$. Let

$$c = |u_+| + |u_-| \leq \frac{1}{\cos(\theta/2)},$$

where the inequality follows from Part 1. Since $\Im(u_+ + u_-) = 0$, we have

$$(\sin \alpha)|u_+| = (\sin \beta)|u_-|$$

and hence

$$|u_+| = \frac{c \sin \beta}{\sin \alpha + \sin \beta} \quad \text{and} \quad |u_-| = \frac{c \sin \alpha}{\sin \alpha + \sin \beta}.$$

We have

$$|\Im (u_+ - u_-)| = |\Im u_+| + |\Im u_-| = (\sin \alpha)|u_+| + (\sin \beta)|u_-| = \frac{2c (\sin \alpha)(\sin \beta)}{\sin \alpha + \sin \beta}.$$

Since the function $x \longmapsto 1/\sin(x)$ is convex on the interval $0 < x < \pi/2$, the minimum value of

$$\frac{\sin \alpha + \sin \beta}{(\sin \alpha)(\sin \beta)} = \frac{1}{\sin \alpha} + \frac{1}{\sin \beta}$$

subject to the constraints $0 < \alpha, \beta < \pi/2$ and $\alpha + \beta \le \theta$ is attained at $\alpha = \beta = \theta/2$. Consequently,

$$|\Im (u_+ - u_-)| \le \frac{2c \sin^2(\theta/2)}{2 \sin(\theta/2)} = c \sin(\theta/2) \le \tan \frac{\theta}{2},$$

which proves Part 2.   □

Part 1 can be extended to the sum of more than two vectors with pairwise angles not exceeding θ, for which one should require $\theta < 2\pi/3$; see Lemma 3.6.3 in [3].

Lemma 2. *Let $F \subset \{-1,1\}^n$ be a face of $\{-1,1\}^n$. Suppose that for all $f \in \mathcal{U}(b; \delta)$, we have $\mathcal{S}\left(e^f | F\right) \ne 0$ and, moreover, the following condition is satisfied: if i is a free index of F and $F^+ \subset F$ and $F^- \subset F$ are the faces obtained by setting the i-th coordinate $\xi_i = 1$ and $\xi_i = -1$, respectively, then the angle between the numbers $\mathcal{S}\left(e^f | F^+\right) \ne 0$ and $\mathcal{S}\left(e^f | F^-\right) \ne 0$ does not exceed θ for some $0 \le \theta < \pi$.*

Since $\mathcal{S}\left(e^f | F\right) \ne 0$ for all $f \in \mathcal{U}(b; \delta)$,

$$f (\xi_1, \ldots, \xi_n) = \sum_{1 \le i < j \le n} a_{ij}\xi_i\xi_j + \sum_{i=1}^{n} b_i\xi_i,$$

and the set $\mathcal{U}(b; \delta)$ is simply connected, we can choose a branch of the function $f \longmapsto \ln \mathcal{S}\left(e^f | F\right)$ for $f \in \mathcal{U}(b; \delta)$.

Let us fix two indices $1 \le p < q \le n$, at most one of which is free for F. Then

$$\left| \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}\left(e^f | F\right) \right| \le \frac{1}{\cos(\theta/2)} \quad \text{and} \quad \left| \Im \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}\left(e^f | F\right) \right| \le \tan \frac{\theta}{2}$$

for all $f \in \mathcal{U}(b; \delta)$.

Proof. Differentiating, we get

$$\frac{\partial}{\partial a_{pq}} \ln \mathcal{S}\left(e^f | F\right) = \frac{\mathcal{S}\left(\xi_p\xi_q e^f | F\right)}{\mathcal{S}\left(e^f | F\right)}.$$

Suppose first that neither of the indices p or q is free. Then the value of $\xi_p\xi_q$ is constant for all $x \in F$ and hence

$$\frac{\mathcal{S}\left(\xi_p\xi_q e^f | F\right)}{\mathcal{S}\left(e^f | F\right)} = \xi_p\xi_q,$$

from which the conclusions of the lemma follow.

Suppose now that only one of the indices p and q, say q, is free. Let $F^+ \subset F$ and $F^- \subset F$ be the faces obtained by setting the q-th coordinate $\xi_q = 1$ and $\xi_q = -1$, respectively. Applying (6), we get

$$\frac{\mathcal{S}\left(\xi_p \xi_q e^f \,|F\right)}{\mathcal{S}\left(e^f \,|F\right)} = \xi_p \frac{\mathcal{S}\left(\xi_q e^f \,|F\right)}{\mathcal{S}\left(e^f \,|F\right)} = \xi_p \left(\frac{\mathcal{S}\left(e^f \,|F^+\right)}{\mathcal{S}\left(e^f \,|F\right)} - \frac{\mathcal{S}\left(e^f \,|F^-\right)}{\mathcal{S}\left(e^f \,|F\right)} \right).$$

We apply Lemma 1 with

$$w_+ = \mathcal{S}\left(e^f \,|F^+\right) \quad \text{and} \quad w_- = \mathcal{S}\left(e^f \,|F^-\right).$$

By (6), we have

$$w = w_+ + w_- = \mathcal{S}\left(e^f \,|F\right)$$

and the angle between $w_+ \neq 0$ and $w_- \neq 0$ does not exceed θ by the assumption of the lemma. Applying Part 1 of Lemma 1, we get

$$\left| \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}\left(e^f \,|F\right) \right| = \left| \frac{w_+}{w} - \frac{w_-}{w} \right| \leq \frac{|w_+| + |w_-|}{|w|} \leq \frac{1}{\cos(\theta/2)}.$$

Applying Part 2 of Lemma 1, we conclude that

$$\left| \mathfrak{I} \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}\left(e^f \,|F\right) \right| = \left| \mathfrak{I} \frac{w_+ - w_-}{w} \right| \leq \tan\frac{\theta}{2},$$

which completes the proof.   □

Corollary 1. *Let $F \subset \{-1, 1\}^n$ be a face as in Lemma 2, and let $f \in \mathcal{U}(b; \delta)$ be a polynomial. Suppose that p is a fixed index of F, and let $g \in \mathcal{U}(b; \delta)$ be a polynomial obtained by replacing the coefficient a_{pq} in f for some q by $-a_{pq}$. Then the angle between $\mathcal{S}\left(e^f \,|F\right) \neq 0$ and $\mathcal{S}\left(e^g |F\right) \neq 0$ does not exceed*

$$\left(2 \tan\frac{\theta}{2} \right) \left| \mathfrak{R}\, a_{pq} \right| + \frac{2 \left| \mathfrak{I}\, a_{pq} \right|}{\cos(\theta/2)}.$$

Proof. For $-1 \leq t \leq 1$, let $f_t : \{-1, 1\}^n \longrightarrow \mathbb{R}$ be the polynomial obtained by replacing a_{pq} with $t a_{pq}$ in f, so $f_1 = f$ and $f_{-1} = g$ and $f_t \in \mathcal{U}(b; \delta)$ for $-1 \leq t \leq 1$. Then

$$\ln \mathcal{S}\left(e^f \,|F\right) - \ln \mathcal{S}\left(e^g |F\right) = \int_{-1}^{1} \frac{d}{dt} \ln \mathcal{S}\left(e^{f_t} |F\right)\, dt = a_{pq} \int_{-1}^{1} \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}\left(e^{f_t} |F\right)\, dt.$$

Consequently,

$$\left| \mathfrak{I} \ln \mathcal{S}\left(e^f \,|F\right) - \mathfrak{I} \ln \mathcal{S}\left(e^g |F\right) \right| \leq \left| \mathfrak{R}\, a_{pq} \right| \int_{-1}^{1} \left| \mathfrak{I} \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}\left(e^{f_t} |F\right) \right|\, dt$$

$$+ \left| \mathfrak{I}\, a_{pq} \right| \int_{-1}^{1} \left| \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}\left(e^{f_t} |F\right) \right|\, dt.$$

Applying Lemma 2, we complete the proof.   □

2.2. Proof of Theorem 1

First, we show that there is $0 < \theta < \pi/2$ such that

$$2(1 - \delta) \tan \frac{\theta}{2} + \frac{\delta^2}{5 \cos(\theta/2)} + \frac{\delta^2}{5} \leq \theta. \tag{7}$$

Indeed, we can just choose $\theta = \delta/2$. Using that

$$\cos x \geq \cos\left(\frac{1}{4}\right) \geq \frac{9}{10} \quad \text{for} \quad 0 \leq x \leq \frac{1}{4}$$

and that

$$\tan x \leq x\left(1 + \frac{x}{10}\right) \quad \text{for} \quad 0 \leq x \leq \frac{1}{4}$$

(since the function $x^{-2} (\tan x - x)$ is increasing for $x > 0$, it suffices to check the inequality for $x = 1/4$), we obtain

$$2(1 - \delta) \tan \frac{\theta}{2} + \frac{\delta^2}{5 \cos(\theta/2)} + \frac{\delta^2}{5} = 2(1 - \delta) \tan \frac{\delta}{4} + \frac{\delta^2}{5 \cos(\delta/4)} + \frac{\delta^2}{5}$$

$$\leq 2(1 - \delta)\frac{\delta}{4}\left(1 + \frac{\delta}{40}\right) + \frac{19\delta^2}{45} = \frac{\delta}{2}\left(1 - \frac{39\delta}{40} - \frac{\delta^2}{40}\right) + \frac{19\delta^2}{45} \leq \frac{\delta}{2} = \theta.$$

We prove by induction for $k = 0, \ldots, n$ the following statement.

Let $F \subset \{-1, 1\}^n$ be a face of dimension k. Then $\mathcal{S}\left(e^f | F\right) \neq 0$. Moreover, if $k > 0$ and q is a free index of F then the following holds. Let $F^+ \subset F$ and $F^- \subset F$ be the faces of F obtained by fixing the q-th variable ξ_q to $\xi_q = 1$ and $\xi_q = -1$, respectively. Then the angle between $\mathcal{S}\left(e^f | F^+\right) \neq 0$ and $\mathcal{S}\left(e^f | F^-\right) \neq 0$ does not exceed θ.

The statement clearly holds for $k = 0$. Suppose that $k > 1$; let q be a free index of F, and let $F^+, F^- \subset F$ be the corresponding faces, so $\dim F^+ = \dim F^- = k - 1$. By the induction hypothesis, we have $\mathcal{S}\left(e^f | F^+\right) \neq 0$ and $\mathcal{S}\left(e^f | F^-\right) \neq 0$. Moreover, $\mathcal{S}\left(e^f | F^-\right) = \mathcal{S}\left(e^g | F^+\right)$, where g is obtained from

$$f(\xi_1, \ldots, \xi_n) = \sum_{1 \leq i < j \leq n} a_{ij}\xi_i\xi_j + \sum_{i=1}^{n} b_i\xi_i$$

by replacing the coefficients a_{pq} for $p \neq q$ by $-a_{pq}$ and the coefficient b_q by $-b_q$. Since q is a fixed index for F^+, replacing b_q by $-b_q$ leads to multiplying $\mathcal{S}\left(e^f | F^+\right)$ by $e^{-2b_q \xi_q}$, which results in the rotation of $\mathcal{S}\left(e^f | F^+\right)$ by an angle of $2(\Im b_q) \leq \delta^2/5$ and some scaling. By Corollary 1, replacing all a_{pq} with $-a_{pq}$ leads to a rotation of $\mathcal{S}\left(e^f | F^+\right)$ by at most an angle of

$$\left(2 \tan \frac{\theta}{2}\right) \sum_{p:\, p \neq q} |\Re a_{pq}| + \frac{2}{\cos(\theta/2)} \sum_{p:\, p \neq q} |\Im a_{pq}| \leq 2(1 - \delta) \tan \frac{\theta}{2} + \frac{\delta^2}{5 \cos(\theta/2)}$$

and also some scaling. Hence the angle between $\mathcal{S}\left(e^f | F^+\right) \neq 0$ and $\mathcal{S}\left(e^f | F^-\right) = \mathcal{S}\left(e^g | F^+\right) \neq 0$ does not exceed

$$2(1 - \delta) \tan \frac{\theta}{2} + \frac{\delta^2}{5 \cos(\theta/2)} + \frac{\delta^2}{5} \leq \theta$$

by (7). Since

$$\mathcal{S}\left(e^f | F\right) = \mathcal{S}\left(e^f | F^+\right) + \mathcal{S}\left(e^f | F^-\right),$$

applying Part 1 of Lemma 1 with

$$w_+ = \mathcal{S}\left(e^f \,|F^+\right) \quad \text{and} \quad w_- = \mathcal{S}\left(e^f \,|F^-\right),$$

we conclude that

$$\mathcal{S}\left(e^f \,|F\right) \neq 0.$$

This concludes the proof of the induction step and hence of Theorem 1.   □

3. Proof of Theorem 2

First, we obtain a version of Lemma 1.

Lemma 3. *Let $v_{++}, v_{+-}, v_{-+}, v_{--} \in \mathbb{C} \setminus \{0\}$ be non-zero numbers, such that the angles*

$$\begin{aligned} \text{between} \quad & v_{++} \quad \text{and} \quad v_{+-}, \\ \text{between} \quad & v_{++} \quad \text{and} \quad v_{-+}, \\ \text{between} \quad & v_{--} \quad \text{and} \quad v_{+-}, \\ \text{between} \quad & v_{--} \quad \text{and} \quad v_{-+} \end{aligned}$$

do not exceed some $0 \leq \theta < \pi/2$, and let $v = v_{++} + v_{+-} + v_{-+} + v_{--}$.

1. *We have $v \neq 0$ and, moreover,*

$$\frac{|v_{++}| + |v_{+-}| + |v_{-+}| + |v_{--}|}{|v|} \leq \frac{1}{\cos \theta}.$$

2. *We have*

$$\left| \mathfrak{I} \frac{v_{++} - v_{+-} - v_{-+} + v_{--}}{v} \right| \leq \tan \frac{\theta}{2}.$$

Proof. Clearly, the angle between any two vectors of v_{++}, v_{+-}, v_{-+} and v_{--} does not exceed $2\theta < \pi$. Let $K \subset \mathbb{R}^2 = \mathbb{C}$ be the angle (cone) spanned by the four vectors. Hence the angular measure of K is at most 2θ.

The proof of Part 1 goes exactly as in the proof of Lemma 1.

We let $u_{++}, u_{+-}, u_{-+}, u_{--}$ and u be the orthogonal projections of the vectors $v_{++}, v_{+-}, v_{-+}, v_{--}$ and v, respectively, onto the bisector of K. Then

$$|v| \geq |u| = |u_{++}| + |u_{+-}| + |u_{-+}| + |u_{--}| \geq \cos(\theta)\left(|v_{++}| + |v_{+-}| + |v_{-+}| + |v_{--}|\right)$$

and the conclusion of Part 1 follows.

The proof of Part 2 is obtained by the application of Part 2 of Lemma 1. Namely, let

$$w_+ = v_{++} + v_{--} \quad \text{and} \quad w_- = v_{+-} + v_{-+},$$

so that

$$\frac{v_{++} - v_{+-} - v_{-+} + v_{--}}{v} = \frac{w_+ - w_-}{w_+ + w_-}.$$

We claim that the angle between w_+ and w_- does not exceed θ. Indeed, without loss of generality, we assume that one side of the angle K is spanned by v_{++}. If the other side of K is spanned by one of the vectors v_{+-} or v_{-+} or the angle is just the ray spanned by v_{++}, then the angular measure of K does not exceed θ and hence the angle between w_+ and w_- also does not exceed θ. Suppose now that the other side of K is spanned by v_{--}. Without loss of generality, we assume that listed in the clockwise order,

the vectors in K are v_{++}, v_{+-}, v_{-+} and v_{--}. We consider the following two angles: angle K_1 spanned by v_{++} and v_{-+} and angle K_2 spanned by v_{+-} and v_{--}. We note that the angular measure of K_1 does not exceed θ and the angular measure of K_2 does not exceed θ. Vector w_- lies in the angle spanned by v_{+-} and v_{-+} and hence in $K_1 \cap K_2$. On the other hand, since $K = K_1 \cup K_2$, vector w_+ lies in at least one of the angles K_1 or K_2. In either case, the angle between w_+ and w_- does not exceed θ. Applying Part 2 of Lemma 1, we conclude that

$$\left| \Im \frac{v_{++} - v_{+-} - v_{-+} + v_{--}}{v} \right| = \left| \Im \frac{w_+ - w_-}{w_+ + w_-} \right| \leq \tan \frac{\theta}{2},$$

as claimed. □

Similar to Section 2, we introduce a set $\mathcal{U}(b; \delta) \subset \mathbb{C}^{\binom{n}{3}+\binom{n}{2}}$ of cubic polynomials

$$f(x) = \sum_{1 \leq i < j < k \leq n} c_{ijk} \xi_i \xi_j \xi_k + \sum_{1 \leq i < j \leq n} a_{ij} \xi_i \xi_j + \sum_{i=1}^n b_i \xi_i \quad \text{for} \quad x = (\xi_1, \ldots, \xi_n)$$

such that

$$\sum_{j,k: \, j,k \neq i} \left| \Re c_{ijk} \right| + \sum_{j: \, j \neq i} \left| \Re a_{ij} \right| \leq 1 - \delta,$$

$$\sum_{j,k: \, j,k \neq i} \left| \Im c_{ijk} \right| + \sum_{j: \, j \neq i} \left| \Im a_{ij} \right| \leq \frac{\delta^2}{10} \quad \text{and}$$

$$\left| \Im b_i \right| \leq \frac{\delta^2}{10} \quad \text{for} \quad i = 1, \ldots, n.$$

Next, we need a refinement of Lemma 2.

Lemma 4. *Let $F \subset \{-1, 1\}^n$ be a face of $\{-1, 1\}^n$. Suppose that for all $f \in \mathcal{U}(b; \delta)$, we have $\mathcal{S}(e^f | F) \neq 0$ and, moreover, the following condition is satisfied: if $G \subseteq F$ is a face, if i is a free index of G and $G^+ \subset G$ and $G^- \subset G$ are the faces obtained by setting the i-th coordinate $\xi_i = 1$ and $\xi_i = -1$, respectively, then the angle between numbers $\mathcal{S}(e^f | G^+) \neq 0$ and $\mathcal{S}(e^f | G^-) \neq 0$ does not exceed θ for some $0 \leq \theta < \pi/2$. We pick a branch of $\ln \mathcal{S}(e^f | F)$ for $f \in \mathcal{U}(b; \delta)$.*

1. *Let us fix three indices $1 \leq p < q < r \leq n$, at most two of which are free for F. Then*

$$\left| \frac{\partial}{\partial c_{pqr}} \ln \mathcal{S}(e^f | F) \right| \leq \frac{1}{\cos \theta} \quad \text{and} \quad \left| \Im \frac{\partial}{\partial c_{pqr}} \ln \mathcal{S}(e^f | F) \right| \leq \tan \frac{\theta}{2}.$$

2. *Let us fix two indices $1 \leq p < q \leq n$, at most one of which is free for F. Then*

$$\left| \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}(e^f | F) \right| \leq \frac{1}{\cos(\theta/2)} \leq \frac{1}{\cos \theta} \quad \text{and}$$

$$\left| \Im \frac{\partial}{\partial a_{pq}} \ln \mathcal{S}(e^f | F) \right| \leq \tan \frac{\theta}{2}.$$

Proof. If there is at most one free index, the proof follows that of Lemma 2. Suppose that exactly two indices, say p and q are free, while r is fixed. Then

$$\frac{\partial}{\partial c_{pqr}} \ln \mathcal{S}(e^f | F) = \frac{\mathcal{S}(\xi_p \xi_q \xi_r e^f | F)}{\mathcal{S}(e^f | F)} = \frac{\xi_r \mathcal{S}(\xi_p \xi_q e^f | F)}{\mathcal{S}(e^f | F)}.$$

Let us define four faces $F^{++}, F^{+-}, F^{-+}, F^{--} \subset F$ as follows:

$$
\begin{aligned}
F^{++} \subset F \quad &\text{is obtained from } F \text{ by setting} \quad \xi_p = 1, \xi_q = 1; \\
F^{+-} \subset F \quad &\text{is obtained from } F \text{ by setting} \quad \xi_p = 1, \xi_q = -1; \\
F^{-+} \subset F \quad &\text{is obtained from } F \text{ by setting} \quad \xi_p = -1, \xi_q = 1; \\
F^{--} \subset F \quad &\text{is obtained from } F \text{ by setting} \quad \xi_p = -1, \xi_q = -1.
\end{aligned}
$$

We let

$$
\begin{aligned}
v_{++} &= \mathcal{S}\left(e^f \,|F^{++}\right), \\
v_{+-} &= \mathcal{S}\left(e^f \,|F^{+-}\right), \\
v_{-+} &= \mathcal{S}\left(e^f \,|F^{-+}\right), \\
v_{--} &= \mathcal{S}\left(e^f \,|F^{--}\right),
\end{aligned}
$$

so that

$$
\mathcal{S}\left(e^f \,|F\right) = v_{++} + v_{+-} + v_{-+} + v_{--}
$$

and

$$
\frac{\partial}{\partial c_{pqr}} \ln \mathcal{S}\left(e^f \,|F\right) = \xi_r \frac{v_{++} - v_{+-} - v_{-+} + v_{--}}{v_{++} + v_{+-} + v_{-+} + v_{--}}.
$$

The proof now follows by Lemma 3.   □

As in Section 2, we obtain a corollary.

Corollary 2. *Let $F \subset \{-1, 1\}^n$ be a face as in Lemma 4, and let $f \in \mathcal{U}(b; \delta)$ be a polynomial.*

1. *Suppose that r is a fixed index of F, and let $g \in \mathcal{U}(b; \delta)$ be a polynomial obtained from f by replacing the coefficient c_{pqr} in f for some p and q by $-c_{pqr}$. Then the angle between $\mathcal{S}\left(e^f \,|F\right) \neq 0$ and $\mathcal{S}\left(e^g |F\right) \neq 0$ does not exceed*

$$
\left(2 \tan \frac{\theta}{2}\right) \left|\Re\, c_{pqr}\right| + \frac{2 \left|\Im\, c_{pqr}\right|}{\cos \theta}.
$$

2. *Suppose that p is a fixed index of F, and let $g \in \mathcal{U}(b; \delta)$ be a polynomial obtained from f by replacing the coefficient a_{pq} in f for some q by $-a_{pq}$. Then the angle between $\mathcal{S}\left(e^f \,|F\right) \neq 0$ and $\mathcal{S}\left(e^g |F\right) \neq 0$ does not exceed*

$$
\left(2 \tan \frac{\theta}{2}\right) \left|\Re\, a_{pq}\right| + \frac{2 \left|\Im\, a_{pq}\right|}{\cos \theta}.
$$

Proof. The proof is similar to that of Corollary 1, only we use Lemma 4 instead of Lemma 2.   □

3.1. *Proof of Theorem 2*

First, we show that there is $0 \leq \theta < \pi/4$ such that

$$
2(1 - \delta) \tan \frac{\theta}{2} + \frac{\delta^2}{5 \cos \theta} + \frac{\delta^2}{5} \leq \theta.
$$

As in Section 2.2, we can choose $\theta = \delta/2$ (note that we assumed that $\delta < 1/2$ and hence we can use the estimates of Section 2.2 verbatim).

We prove by induction the same statement as in Section 2.2. The only difference is that if r is a free index of F and $F^+, F^- \subset F$ are the faces obtained by setting $\xi_r = 1$ and $\xi_r = -1$, respectively, then $\mathcal{S}\left(e^f | F^+\right) = \mathcal{S}\left(e^g | F^-\right)$, where g is obtained from f by changing all coefficients c_{pqr} to $-c_{pqr}$, all coefficients a_{qr} to $-a_{qr}$ and the coefficient b_r to $-b_r$. We then use Corollary 2 instead of Corollary 1. □

3.2. Polynomials of higher degree

We don't know if the straightforward extensions of Theorems 1 and 2 hold for polynomials $f :$ $\{-1, 1\}^n \longrightarrow \mathbb{C}$ of degree 4 and higher (for some non-optimal estimates in that case, see [4]). The reason our argument does not extend to the case of $\deg f \geq 4$ is that we are lacking the appropriate version of Part 2 of Lemma 3. In the case of $\deg f = 4$, that version would have looked as follows: given eight non-zero complex numbers

$$v_{+++}, \ v_{++-}, \ v_{+-+}, \ v_{+--}, \ v_{-++}, \ v_{--+}, \ v_{-+-}, \ v_{---},$$

such that the angle between any two of them whose indices differ in one position does not exceed some small number $\theta > 0$, we have

$$\left| \Im \frac{v_{+++} - v_{++-} - v_{+-+} - v_{-++} + v_{+--} + v_{-+-} + v_{--+} - v_{---}}{v_{+++} + v_{++-} + v_{+-+} + v_{-++} + v_{+--} + v_{-+-} + v_{--+} + v_{---}} \right| \overset{?}{\leq} \tan \frac{\theta}{2}.$$

However, the above inequality cannot hold. One can easily construct a counterexample by choosing $|v_{+++}| = |v_{---}| = 1$, the angle between v_{+++} and v_{---} equal 3θ, while choosing all other vectors to have a very small length. In that case, the left-hand side should be about $\tan(3\theta/2)$, and our proof cannot proceed unless we tighten the bound for the sum of the absolute values of the coefficients. This failure of the approach does not, of course, exclude a possibility that the extensions of Theorems 1 and 2 still hold for higher-degree polynomials.

4. Approximation

Let

$$f(x) = \sum_{1 \leq i < j \leq n} a_{ij} \xi_i \xi_j + \sum_{i=1}^n b_i \xi_i \quad \text{for} \quad x = (\xi_1, \ldots, \xi_n).$$

We suppose that the coefficients a_{ij} and b_i are real and satisfy

$$\sum_{j: \ j \neq i} |a_{ij}| \ \leq \ 1 - \delta \quad \text{for} \quad i = 1, \ldots, n$$

and some $0 < \delta < 1$, fixed in advance (there are no restrictions on b_i). Here we sketch an algorithm for approximating

$$\mathcal{S}\left(e^f\right) = \sum_{x \in \{-1, 1\}^n} e^{f(x)}.$$

Without loss of generality, we assume that $\delta \leq 0.1$ and that n is large enough,

$$n \ \geq \ \frac{100}{\delta^2},$$

since for smaller n the sum $\mathcal{S}\left(e^f\right)$ can be computed by brute force.

We use the method of polynomial interpolation. The gist of the method is summarized in the following lemma; see Section 2.2 of [3] and [5].

Lemma 5. *Let $U \subset \mathbb{C}$ be a connected open set containing 0 and 1. Then there is a constant $\gamma(U) > 0$ such that the following holds. Suppose that*

$$p(z) = \sum_{k=0}^{n} a_k z^k, \quad n \geq 2,$$

is a polynomial such that $p(z) \neq 0$ for all $z \in \mathbb{C}$. Then, for every $0 < \epsilon < 1$, the value of $p(1)$, up to relative error ϵ, is determined by the coefficients a_k with

$$k \leq \gamma(U)(\ln n - \ln \epsilon)$$

and, moreover, can be computed from those coefficients in $n^{O(1)}$ time.

Here we say that two complex numbers $w_1 \neq 0$ and $w_2 \neq 0$ approximate each other within relative error $0 < \epsilon < 1$ if we can write $w_1 = e^{z_1}$ and $w_2 = e^{z_2}$, where $|z_1 - z_1| \leq \epsilon$.

In our case, $U \subset \mathbb{C}$ is a sufficiently small neighborhood of the interval $[0, 1] \subset \mathbb{C}$. Our first goal is to write $\mathcal{S}\left(e^f\right)$ as $p(1)$ for a univariate polynomial of p with $\deg p = n^{O(1)}$.

4.1. Writing $\mathcal{S}\left(e^f\right)$ as $p(1)$ for a univariate polynomial $p(z)$

First, we rewrite $\mathcal{S}\left(e^f\right)$ as a polynomial in some new variables. We have

$$e^{f(x)} = \left(\prod_{1 \leq i < j \leq n} e^{a_{ij}\xi_i\xi_j}\right)\left(\prod_{i=1}^{n} e^{b_i\xi_i}\right)$$

$$= \exp\left\{-\sum_{1 \leq i < j \leq n} a_{ij}\right\}\left(\prod_{1 \leq i < j \leq n} e^{a_{ij}(\xi_i\xi_j+1)}\right)\left(\prod_{i=1}^{n} e^{b_i\xi_i}\right).$$

Let

$$c_{ij} = e^{a_{ij}/n^2} - 1 \quad \text{for} \quad 1 \leq i < j \leq n. \tag{8}$$

For given $b_1, \ldots, b_n$, we consider a polynomial

$$P_b(c) = \sum_{\xi_1,\ldots,\xi_n=\pm 1} \left(\prod_{1 \leq i < j \leq n} (c_{ij}+1)^{n^2(\xi_i\xi_j+1)}\right)\left(\prod_{i=1}^{n} e^{b_i\xi_i}\right)$$

of degree $n^3(n-1)$ in the vector c of $\binom{n}{2}$ variables c_{ij}. Thus we have

$$\mathcal{S}\left(e^f\right) = \exp\left\{-\sum_{1 \leq i < j \leq n} a_{ij}\right\} P_b(c)$$

and we want to approximate $P_b(c)$.

Next, for given c and b_i, we consider a univariate polynomial

$$p(z) = P_b(zc) = \sum_{\xi_1,\ldots,\xi_n=\pm 1} \left(\prod_{1 \leq i < j \leq n} (zc_{ij}+1)^{n^2(\xi_i\xi_j+1)}\right)\left(\prod_{i=1}^{n} e^{b_i\xi_i}\right) \tag{9}$$

of a complex variable z. Our goal is to approximate $p(1)$.

4.2. *Showing that $p(z) \neq 0$ in a neighborhood of $[0,1] \subset \mathbb{C}$*

Our next goal is to show that

$$p(z) \neq 0 \quad \text{provided} \quad |\Re z| \leq 1 + \delta^2 \quad \text{and} \quad |\Im z| \leq \frac{\delta^2}{80}. \tag{10}$$

Since $|a_{ij}| \leq 1$ and $n \geq 10$, from (8) we deduce that

$$\left| c_{ij} - \frac{a_{ij}}{n^2} \right| \leq \frac{1}{n^4} \quad \text{for all} \quad i, j \tag{11}$$

(follows from the Taylor series expansion of e^{a_{ij}/n^2}). In particular,

$$|c_{ij}| \leq \frac{2}{n^2} \quad \text{for all} \quad i, j. \tag{12}$$

From (8), we obtain

$$a_{ij} = n^2 \ln\left(1 + c_{ij}\right) \quad \text{for all} \quad i, j.$$

Let us choose an arbitrary $z \in \mathbb{C}$ such that $|z| \leq 1 + \delta^2$, and let us define

$$\widehat{a}_{ij} = \widehat{a}_{ij}(z) = n^2 \ln\left(1 + z c_{ij}\right) \quad \text{for all} \quad i, j, \tag{13}$$

where we choose the principal branch of the logarithm, so that $\widehat{a}_{ij}(0) = 0$. Let

$$\widehat{f}(x) = \sum_{1 \leq i < j \leq n} \widehat{a}_{ij} \xi_i \xi_j + \sum_{i=1}^{n} b_i \xi_i \quad \text{for} \quad x = (\xi_1, \ldots, \xi_n).$$

From Section 4.1, we have

$$p(z) = \exp\left\{ \sum_{1 \leq i < j \leq n} \widehat{a}_{ij} \right\} \mathcal{S}\left(e^{\widehat{f}} \right). \tag{14}$$

Combining (11)–(13), we obtain

$$\widehat{a}_{ij} = z a_{ij} + \eta_{ij} \quad \text{where} \quad |\eta_{ij}| \leq \frac{3}{n^2} \quad \text{for all} \quad i, j$$

(follows from the Taylor series expansion of $\ln(1 + z c_{ij})$). Consequently,

$$\sum_{j: \, j \neq i} |\Re \widehat{a}_{ij}| \leq (1 - \delta)(1 + \delta^2) + \frac{3}{n} \leq 1 - \frac{\delta}{2}.$$

In addition, if $|\Im z| \leq \delta^2/80$, then

$$\sum_{j: \, j \neq i} |\Im \widehat{a}_{ij}| \leq \frac{\delta^2}{80} + \frac{3}{n} \leq \frac{\delta^2}{40}.$$

Hence by Theorem 1 we have that $\mathcal{S}\left(e^{\widehat{f}} \right) \neq 0$, and therefore by (14), we conclude that (10) holds.

4.3. Computing $p^{(k)}(0)$

As discussed in Section 2.2 of [3] (see also [17] and [5] for some enhancements), as soon as $p(z) \neq 0$ in some neighborhood U of the interval $[0, 1] \subset \mathbb{C}$, to approximate $p(1)$ within relative error $0 < \epsilon < 1$, it suffices to compute the derivatives $p^{(k)}(0)$ for $k = \gamma(\ln \deg p - \ln \epsilon)$, where $\gamma = \gamma(U) > 0$. Since in our case $\deg p \leq n^4$, to approximate $\mathcal{S}(e^f) = p(1)$ within relative error $0 < \epsilon < 1$, it suffices to compute $p^k(0)$ for $k = O(\ln n - \ln \epsilon)$, where the implied constant in the 'O' notation depends on δ alone.

From (9), we obtain

$$p^{(k)}(0) = \sum_{\substack{(k_1,\ldots,k_s): \\ k_1+\ldots+k_s=k}} \binom{k}{k_1,\ldots,k_s} \left(2n^2\right)_{k_1} \cdots \left(2n^2\right)_{k_s} \sum_{\xi_1,\ldots,\xi_n=\pm 1} \left(\prod_{i=1}^{n} e^{b_i \xi_i}\right)$$

$$\times \sum_{\substack{\{i_1<j_1, i_2<j_2,\ldots,i_s<j_s\} \\ \xi_{i_1}\xi_{j_1}=\ldots=\xi_{i_s}\xi_{j_s}=1}} c_{i_1 j_1}^{k_1} \cdots c_{i_s j_s}^{k_s},$$

where the outermost sum is taken over all positive integer vectors $(k_1, \ldots, k_s)$ such that $k_1 + \ldots + k_s = k$, the innermost sum is taken over all s-sets of pairs $i_1 < j_1, i_2 < j_2, \ldots, i_s < j_s$, such that the products $\xi_{i_1}\xi_{j_1}, \ldots, \xi_{i_s}\xi_{j_s}$ are all equal 1's, and where

$$\left(2n^2\right)_{k_i} = 2n^2 \left(2n^2 - 1\right) \cdots \left(2n^2 - k_i + 1\right).$$

It is convenient to rewrite the above sum in the graph notation. Let K_n be the complete undirected graph with set $V = \{1, \ldots, n\}$ of vertices and set E of edges. We assign weight c_{ij} to edge $\{i, j\}$. We call a map $\sigma : V \longrightarrow \{-1, 1\}$ *consistent* on an edge $u = \{i, j\}$ if $\sigma(i) = \sigma(j)$. Then

$$p^{(k)}(0) = \sum_{\substack{(k_1,\ldots,k_s): \\ k_1+\ldots+k_s=k}} \binom{k}{k_1,\ldots,k_s} \left(2n^2\right)_{k_1} \cdots \left(2n^2\right)_{k_s} \sum_{\sigma:\, V \longrightarrow \{-1,1\}} \left(\prod_{i=1}^{n} e^{b_i \sigma(i)}\right)$$

$$\times \sum_{\substack{\{u_1,\ldots,u_s\} \subset E: \\ \sigma \text{ is consistent on each } u_1,\ldots,u_s}} c_{u_1}^{k_1} \cdots c_{u_s}^{k_s}$$

$$= \sum_{\substack{(k_1,\ldots,k_s): \\ k_1+\ldots+k_s=k}} \binom{k}{k_1,\ldots,k_s} \left(2n^2\right)_{k_1} \cdots \left(2n^2\right)_{k_s} \sum_{\{u_1,\ldots,u_s\} \subset E} c_{u_1}^{k_1} \cdots c_{u_s}^{k_s}$$

$$\sum_{\substack{\sigma:\, V \longrightarrow \{-1,1\}: \\ \sigma \text{ is consistent on each } u_1,\ldots,u_s}} \left(\prod_{i=1}^{n} e^{b_i \sigma(i)}\right).$$

For a given $s \leq k$, there are not more than $n^{2s} \leq n^{2k}$ sets of s distinct edges $u_1, \ldots, u_s$. Given such a set $u_1, \ldots, u_s$, let $W \subset V$ be the set of vertices of $u_1, \ldots, u_s$. Then $|W| \leq 2s$, and there are at most 2^s maps $\sigma_0 : W \longrightarrow \{-1, 1\}$ that are consistent on each edge $u_1, \ldots, u_s$. Finally, given such a map σ_0, we have

$$\sum_{\substack{\sigma:\, V \longrightarrow \{-1,1\}: \\ \sigma|W=\sigma_0}} \prod_{i=1}^{n} e^{b_i \sigma(i)} = \prod_{i \in W} e^{b_i \sigma_0(i)} \prod_{i \in V \setminus W} \left(e^{b_i} + e^{-b_i}\right).$$

Summarizing,

$$
p^{(k)}(0) = \sum_{\substack{(k_1,\ldots,k_s):\\ k_1+\ldots+k_s=k}} \binom{k}{k_1,\ldots,k_s} \left(2n^2\right)_{k_1} \cdots \left(2n^2\right)_{k_s} \sum_{\{u_1,\ldots,u_s\} \subset E} c_{u_1}^{k_1} \ldots c_{u_s}^{k_s}
$$
$$
\sum_{\sigma_0: W \longrightarrow \{-1,1\}} \prod_{i \in W} e^{b_i \sigma_0(i)} \prod_{i \in V \setminus W} \left(e^{b_i} + e^{-b_i}\right),
$$

where the outermost sum is taken over at most $\binom{2k-1}{k-1}$ positive integer vectors $(k_1,\ldots,k_s)$ satisfying the equation $k_1 + \ldots + k_s = k$, the next sum is taken over at most n^{2k} subsets of $s \leq k$ edges and the innermost sum is taken over at most $2^s \leq 2^k$ consistent maps σ_0 on the set W of vertices of $\{u_1,\ldots,u_s\}$. Hence the complexity of computing $p^{(k)}(0)$ is $n^{O(k)}$, and since $k = O(\ln n - \ln \epsilon)$, we obtain an algorithm of quasi-polynomial $n^{O(\ln n - \ln \epsilon)}$ complexity to approximate $\mathcal{S}\left(e^f\right)$.

Guus Regts explained to the authors that the argument of [17] allows one to obtain a genuinely polynomial, as opposed to quasi-polynomial, algorithm if the maximum degree Δ of the underlying graph G with edges $\{i,j\}$, where $a_{ij} \neq 0$ is bounded from above in advance. This is based on the observation that computing $p^{(k)}(0)$ in that case reduces to weighted enumeration of *connected* subgraphs of G with $O(k)$ vertices, which can be accomplished in polynomial time as long as Δ is fixed and $k = O(\ln n - \ln \epsilon)$.

The extension of the algorithm to cubic polynomials f is straightforward. We use then Theorem 2 instead of Theorem 1 to ascertain that the auxiliary univariate polynomial $p(z)$ does not have roots in a neighborhood of $[0,1] \subset \mathbb{C}$.

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