

INTERSECTING THE SIDES OF A POLYGON

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ABSTRACT. Consider the map S which sends a planar polygon P to a new polygon $S(P)$ whose vertices are the intersection points of second-nearest sides of P . This map is the inverse of the famous pentagram map. In this paper we investigate the dynamics of the map S . Namely, we address the question of whether a convex polygon stays convex under iterations of S . Computer experiments suggest that this almost never happens. We prove that indeed the set of polygons which remain convex under iterations of S has measure zero, and moreover it is an algebraic subvariety of codimension two. We also discuss the equations cutting out this subvariety, as well as their geometric meaning in the case of pentagons.

1. INTRODUCTION

Let P be a planar polygon, and let $S(P)$ be the polygon whose vertices are the intersection points of second-nearest sides of P ; see Figure 1. The map S is the inverse of the celebrated *pentagram map*, defined by R. Schwartz [6] and studied by many others. Indeed, as can be seen in Figure 1, one can recover P from $S(P)$ by intersecting consecutive shortest diagonals (i.e. diagonals connecting second-nearest vertices), which is precisely the definition of the pentagram map. In what follows, we denote the pentagram map by D , so that $S = D^{-1}$ and $D = S^{-1}$.

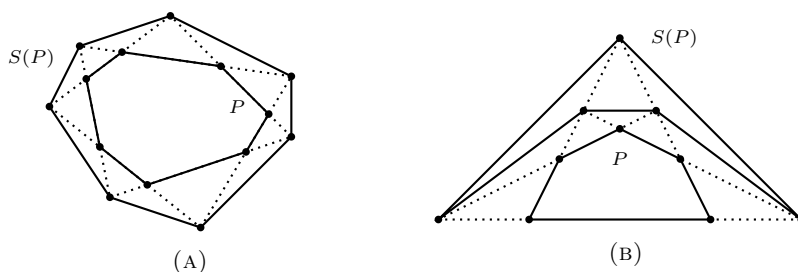


FIGURE 1. The inverse pentagram map

As can be seen from Figure 1, the image of a convex polygon under the inverse pentagram map S does not have to be convex. Moreover, computer experiments show that a randomly chosen polygon becomes non-convex after several applications of S ; see Figure 2. The goal of the present paper is to determine necessary

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FIGURE 2. The orbit of a polygon under the inverse pentagram map. Each iteration is normalized by means of an appropriate affine transformation.

and sufficient conditions on a convex polygon P which guarantee that all successive images of P under S , i.e. $S(P), S(S(P)), \dots$, are convex polygons. Our main result is that the set of polygons with this property has zero measure and moreover is a codimension two algebraic surface. Furthermore, we present explicit equations describing this surface, i.e. explicit conditions on P which ensure that all its successive images under the map S are convex.

We will state our main result in two different forms: in terms of a certain vector d_P associated with the polygon P , and in terms of a certain operator G_P also associated with P .

To define d_P , consider a convex planar n -gon P in the affine plane. Assume that the vertices of P are labeled in cyclic order by residues modulo n . Denote by d_i the vector connecting the vertices $i-1$ and $i+1$, and let $\mathcal{A}_i$ be the area of the triangle cut out by the line joining those vertices; see Figure 3. Then d_P is defined by

$$(1) \quad d_P := \sum_{i=1}^n \frac{d_i}{\mathcal{A}_i}.$$

The first way to formulate our main result is as follows: for a convex polygon P , all successive images of P under the inverse pentagram map S are convex if and only if $d_P = 0$. Since this is equivalent to two algebraic equations on the coordinates of vertices, it follows that the set of convex n -gons P such that $S^k(P)$ is convex for every $k \geq 0$ is a codimension two algebraic surface in the $2n$ -dimensional space of all convex n -gons.

Another way to state our result is in terms of a certain operator introduced by M. Glick [2]. This operator is defined as follows. Assume that we are given a polygon P in the projective plane $\mathbb{P}^2$. Let $v_i \in \mathbb{R}^3$ be the vector of homogeneous coordinates of the i 'th vertex of P . Then Glick's operator $G_P: \mathbb{R}^3 \rightarrow \mathbb{R}^3$ is defined

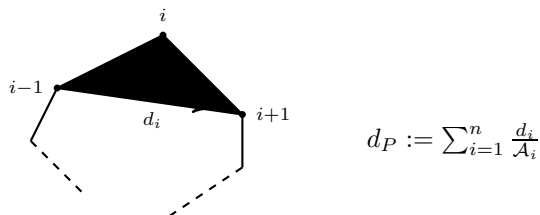


FIGURE 3. To the construction of the vector d_P

by

$$(2) \quad G_P(v) := nv - \sum_{i=1}^n \frac{v_{i-1} \wedge v \wedge v_{i+1}}{v_{i-1} \wedge v_i \wedge v_{i+1}} v_i.$$

Observe that the right-hand side does not change under rescaling of v_i 's, so this is indeed a well-defined operator. Furthermore, as any operator in the 3-space, Glick's operator can be interpreted as a projective mapping $\mathbb{P}^2 \rightarrow \mathbb{P}^2$.

Theorem 1.1. *Consider a convex planar polygon P in the affine plane with at least five vertices. Then the following conditions are equivalent:*

- (1) *All successive images of the polygon P under the inverse pentagram map S are convex.*
- (2) *The associated vector d_P defined by (1) is zero.*
- (3) *The associated Glick's operator G_P defined by (2), regarded as a projective mapping, is affine.*

We prove this theorem by establishing equivalences $2 \Leftrightarrow 3$ and $1 \Leftrightarrow 3$. The proof of $2 \Leftrightarrow 3$ is a direct computation which we outline in Section 2. As for $1 \Leftrightarrow 3$, that part is based on more subtle properties of Glick's operator and its connection with pentagram dynamics. We discuss these properties in Section 3, after which we complete the proof in Section 4.

Remark 1.2. Since the condition $d_P = 0$ is equivalent to two algebraic equations in terms of vertices of P , it follows from Theorem 1.1 that the set of polygons which remain convex under iterations of S is an algebraic subvariety of codimension two. The latter can also be proved without using the vector d_P , as outlined in Remark 4.6.

In the case of pentagons, there is one more condition equivalent to the above three:

Theorem 1.3. *For pentagons, the three conditions of Theorem 1.1 are equivalent to the following one:*

- (4) *The conic inscribed in P and the conic circumscribed about P are concentric.*

Apparently, there should be an elementary way of proving this by establishing equivalence $2 \Leftrightarrow 4$: the inscribed and circumscribed conics are concentric if and only if $d_P = 0$. However, we are not aware of such a proof. What we do instead is directly prove the equivalence $1 \Leftrightarrow 4$. The main ingredient of the proof is E. Kasner's theorem on pentagons.

Remark 1.4. S. Tabachnikov [7] proved that Kasner's theorem holds for all Poncelet polygons (i.e. polygons inscribed in conic and circumscribed about a conic). Therefore, Theorem 1.3 should be true for such polygons too. We do not consider this case here so as to not encumber the exposition.

Remark 1.5. Glick (personal communication) suggested the following geometric interpretation of the condition $d_P = 0$, valid regardless of the number of vertices. Say that two polygons P and Q are *dual with respect to the line at infinity* if sides of P are parallel to shortest diagonals of Q , while sides of Q are parallel to shortest diagonals of P ; see Figure 4. Then it turns out that a polygon P admits such a dual if and only if $d_P = 0$.

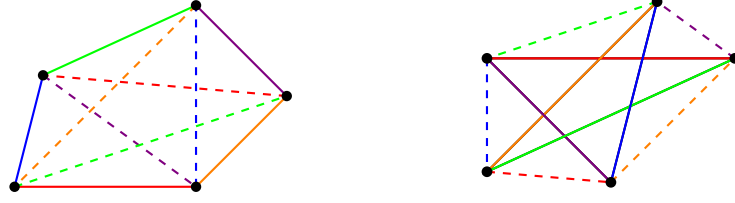


FIGURE 4. Polygons dual with respect to the line at infinity. Lines of the same color and style are parallel to each other.

2. WHEN IS GLICK'S OPERATOR AFFINE?

We start the proof of Theorem [1.1](#) by showing the equivalence $2 \Leftrightarrow 3$: Glick's operator's G_P is affine if and only $d_P = 0$. Given a polygon in the affine plane, denote by (x_i, y_i) the Cartesian coordinates of its vertices. We assume that the vertices are labeled in the counter-clockwise order. As in Section [1](#), denote by d_i the vector connecting the vertices $i - 1$ and $i + 1$, and let $\mathcal{A}_i$ be the area of the triangle cut out by the line joining those vertices, as shown in Figure [3](#). Then a straightforward calculation shows that Glick's operator G_P is given by

$$(3) \quad G_P \begin{pmatrix} x \\ y \\ z \end{pmatrix} = n \begin{pmatrix} x \\ y \\ z \end{pmatrix} - \sum_{i=1}^n \frac{1}{\mathcal{A}_i} \left(\det \begin{pmatrix} x & y & d_i \end{pmatrix} - z \cdot \det \begin{pmatrix} x_{i-1} & x_{i+1} \\ y_{i-1} & y_{i+1} \end{pmatrix} \right) \begin{pmatrix} x_i \\ y_i \\ 1 \end{pmatrix}$$

From this formula it follows that the preimage of the line at infinity under G_P is given in homogeneous coordinates x, y, z by

$$(4) \quad \det \begin{pmatrix} x & y & d_P \end{pmatrix} = z \left(n + \sum_{i=1}^n \frac{1}{\mathcal{A}_i} \det \begin{pmatrix} x_{i-1} & x_{i+1} \\ y_{i-1} & y_{i+1} \end{pmatrix} \right).$$

Now assume that G_P is affine. Then the line at infinity is mapped to itself, so the above equation implies $d_P = 0$. Conversely, assume that $d_P = 0$. Then the coefficient of z in the above equation is easily seen to be invariant under coordinate transformations, and by choosing coordinates in such a way that the origin is in the interior of the polygon, one shows that the coefficient of z is always positive and in particular does not vanish. At the same time, since $d_P = 0$, the coefficients of x and y vanish, so the preimage of the line at infinity is the line at infinity, as desired.

3. GLICK'S OPERATOR AND THE PENTAGRAM MAP

In this section we discuss properties of Glick's operator G_P and its connection with the dynamics of the pentagram map D . This is mainly an overview of [\[2\]](#), but for some of the results we present a refined version.

Proposition 3.1. *Let P be a convex polygon in the affine plane with at least five vertices. Then the following is true:*

- (1) *For the dual polygon P^* , one has $G_{P^*} = G_P^*$ (recall that the dual polygon is the polygon in the dual projective plane whose vertices are the sides of initial polygon).*
- (2) *Let $\text{conv}(P)$ be the convex hull of vertices of P , and let $\text{int}(\text{conv}(P))$ be its interior. Then $G_P(\text{conv}(P)) \subset \text{int}(\text{conv}(P))$. In particular, the projective mapping G_P is well-defined at all points of $\text{conv}(P)$ (i.e. none of those points belong to the kernel of G_P when the latter is regarded as an operator in the 3-space).*
- (3) *The intersection $\bigcap_{k \geq 0} \text{conv}(D^k(P)) = \bigcap_{k \geq 0} \text{int}(\text{conv}(D^k(P)))$ is a single point, known as the limit point of the pentagram orbit $D^k(P)$. This point is a fixed point of G_P . Moreover, it is the only fixed point of G_P in $\text{conv}(P)$.*

Proof. (1) See [2, Proposition 3.3].

- (2) See [2, Proposition 4.1]. The statement of that proposition is that $G_P(\text{conv}(P)) \subset \text{conv}(P)$, but it is actually proved that $G_P(\text{conv}(P)) \subset \text{int}(\text{conv}(P))$.

- (3) The existence of the limit point follows from [6, Theorem 3.1]. Denote that point by X . Then, by [2, Proposition 1.2], we have $G_P(X) = X$. So it remains to show that there are no other fixed points in $\text{conv}(P)$. Assume that $Y \neq X$ is another fixed point, $G_P(Y) = Y$. Note that Y must be in the interior of P , since $G_P(\text{conv}(P)) \subset \text{int}(\text{conv}(P))$. Let L be the line in $\mathbb{RP}^2$ through X and Y . Then G_P restricts to and defines a Moebius transformation of L . Furthermore, G_P^2 preserves each of the two connected components of $L \setminus \{X, Y\}$. Now, let W and Z be the intersection points of L with the boundary of the polygon, as shown in Figure 5. Then, since $G_P(\text{conv}(P)) \subset \text{int}(\text{conv}(P))$, it follows that G_P^2 maps W inside the interval WX and Z inside the interval YZ . At the same time, this Moebius transformation preserves X and Y . But Moebius transformations with these properties do not exist. Indeed, for any non-trivial Moebius transformation which has two fixed points and preserves both of the intervals between those points, one of the fixed points must be attractive and the other one repelling. So, if W is moved by G_P^2 towards X , then Z should be moved away from Y . Thus, it is indeed not possible that G_P has two fixed points in $\text{conv}(P)$. □

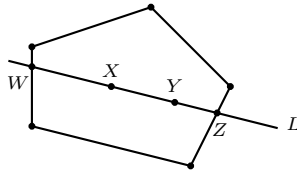


FIGURE 5. To the proof of Proposition 3.1

4. CONVEXITY PERSISTS IF AND ONLY IF GLICK'S OPERATOR IS AFFINE

We now prove the equivalence $1 \Leftrightarrow 3$: for a convex polygon P , all its successive images under the inverse pentagram map S are convex if and only if the projective mapping G_P is affine. To that end, we will reformulate the former condition in terms of the dual polygon. The following is well-known.

Proposition 4.1. *Projective duality intertwines the inverse pentagram map S with the direct pentagram map $D = S^{-1}$. In other words, $S(P)^* = D(P^*)$.*

Proof. The sides of $S(P)^*$ are vertices of $S(P)$, which, by definition of S , are intersections of second-nearest sides of P . At the same time, the sides of $D(P^*)$ are shortest diagonals of P^* , i.e. diagonals connecting second-nearest vertices. But second-nearest vertices of P^* are second-nearest sides of P , so the diagonals connecting them are precisely the intersections of second-nearest sides of P . So, $S(P)^*$ and $D(P^*)$ have the same sides and hence coincide, as claimed. $\square$

We now reformulate the convexity condition for the polygon $S^k(P)$ in terms of its dual $D^k(P^*)$. Note that it is in general not true that the dual of a convex polygon is convex. In fact, this statement does not even make sense, since the definitions of convexity in the initial and dual planes require different structures. To define convexity in the initial plane, one needs to pick an affine chart, which is determined by choosing a *line*. Likewise, convexity in the dual plane also becomes well-defined upon choice of a line, i.e. a *point* in the initial plane. So, in order to be able to talk about convexity in both the initial and dual planes, in the initial plane one needs to pick a line (which defines an affine chart) and a point (which defines an affine chart in the dual plane). The following is a folklore result.

Lemma 4.2. *Consider a projective plane $\mathbb{P}^2$ with a fixed line L and point O . Assume that P is a polygon in that plane which is convex in the affine chart $\mathbb{P}^2 \setminus L$ and contains the point O in its interior. Then the dual polygon P^* is convex in the affine chart $(\mathbb{P}^2)^* \setminus O$ and contains the point L in its interior.*

Remark 4.3. The choice of a point O turns $\mathbb{P}^2 \setminus L$ into a vector space. In that setting, the lemma can be reformulated by saying that the dual of a convex polygon containing the origin is also a convex polygon containing the origin. This is in fact true for any convex set, with an appropriate definition of duality.

Proof of Lemma 4.2. Figure 6 shows a polygon P in the affine chart $\mathbb{P}^2 \setminus L$. Let A be an arbitrary vertex of P . Consider the straight line interval connecting O to A . Then none of the sides of P non-adjacent to A intersect that interval, and neither does the line at infinity L . Therefore, all the sides of P non-adjacent to A , as well as the line L belong to the same connected component in the space of lines

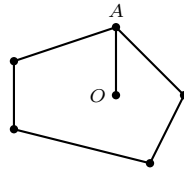


FIGURE 6. To the proof of Lemma 4.2

in $\mathbb{P}^2 \setminus \{O, A\}$. But this means that all vertices of P^* not adjacent to the side A , as well as the point L , belong to the same connected component of $(\mathbb{P}^2)^* \setminus (O \cup A)$. In other words, these points lie on the same side of the line A in the affine plane $(\mathbb{P}^2)^* \setminus O$, and since this holds for every side A of P^* , this precisely means that P^* is convex and contains the point L in its interior. $\square$

Now, consider a convex polygon P in the affine plane, and let L be the line at infinity. Fix a point O in the interior of P . Then, by Lemma 4.2, the dual polygon P^* is convex in the affine chart $(\mathbb{P}^2)^* \setminus O$ and contains the point L in its interior.

Proposition 4.4. *For any given $k > 0$, the polygon $S^k(P)$ is convex if and only if its dual $D^k(P^*)$ contains the point $L \in (\mathbb{P}^2)^*$ in its interior.*

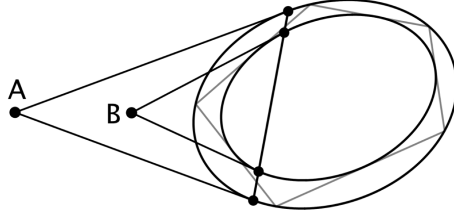
Proof. Assume that the polygon $S^k(P)$ is convex. Observe that $D^k(S^k(P)) = P$, so $S^k(P)$ contains the polygon P and hence the point O in its interior. Therefore, since $S^k(P)$ is convex and contains the point O , by Lemma 4.2 the polygon $D^k(P^*) = S^k(P)^*$ contains the point L . Conversely, assume that $D^k(P^*)$ contains the point L . Note that since the pentagram map D preserves convexity, and P^* is convex, we also have that $D^k(P^*)$ is convex. Therefore, since $D^k(P^*)$ is convex and contains the point L , by Lemma 4.2 we have that $S^k(P) = D^k(P^*)^*$ is convex too, as desired. $\square$

Corollary 4.5. *Each of the polygons $S^k(P)$, where $k > 0$, is convex if and only if the limit point of the sequence $D^k(P^*)$ is the line at infinity L .*

Proof. The limit point is the unique point in the intersection $\bigcap_{k>0} \text{conv}(D^k(P^*))$, so $L \in \text{conv}(D^k(P^*))$ for every $k > 0$ if and only if the limit point is L . $\square$

Now, to complete the proof of Theorem 1.1, it suffices to show that the limit point of the pentagram orbit $D^k(P^*)$ is L if and only if the projective mapping G_P is affine. First, assume that L is the limit point for $D^k(P^*)$. Then, L , viewed as a point in the dual plane, is fixed by the mapping G_{P^*} (Proposition 3.1, Item 3). But in view of the equality $G_{P^*} = G_P^*$ (Proposition 3.1, Item 1), this is equivalent to saying that L , viewed as the line in the initial plane, is invariant under G_P . So, G_P preserves the line at infinity and hence is affine. Conversely, assume that the mapping G_P is affine. Then the line at infinity L , viewed as a point in the dual plane, is a fixed point of G_{P^*} . And since L is inside P^* , it follows from Proposition 3.1, Item 3 that L must be the limit point of $D^k(P^*)$, as desired. Thus, Theorem 1.1 is proved.

Remark 4.6. As pointed out by the referee, one can use Corollary 4.5 to directly prove that the set of polygons which remain convex under iterations of S is of codimension two. The idea is as follows. Let $\mathcal{P}_n$ be the space of convex n -gons. Consider the map $\zeta: \mathcal{P}_n \rightarrow (\mathbb{P}^2)^*$ which takes a polygon P to the limit point of the sequence $D^k(P^*)$. Then, by Corollary 4.5, the set of polygons which remain convex under iterations of S is the preimage under ζ of the line at infinity. Assuming the map ζ is smooth (which follows, for example, from Glick's construction), it must be a submersion because it commutes with the projective group action. So the preimage of the line at infinity (which is a single point in $(\mathbb{P}^2)^*$) is a codimension two submanifold, as desired.

FIGURE 7. To the definition of the map $R_P(A) := B$

5. CASE STUDY: PENTAGONS

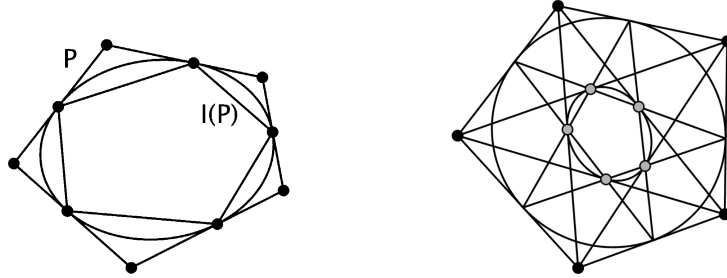
We now prove the equivalence $1 \Leftrightarrow 4$ (i.e. Theorem [1.3](#)): for a convex pentagon P , all its successive images under the inverse pentagram map S are convex if and only if the conic inscribed in P and the conic circumscribed about P are concentric. The strategy of the proof is as follows. Consider a pentagon P , and let its inscribed and circumscribed conics be given by homogeneous quadratic forms Q_I and Q_C respectively. Let $R_P := Q_I^{-1}Q_C$. Note that since the forms Q_I and Q_C are defined up to a constant factor, R_P is well-defined as a projective map $\mathbb{P}^2 \rightarrow \mathbb{P}^2$. The geometric meaning of that map is the following: it takes a point $A \in \mathbb{P}^2$ to point $B \in \mathbb{P}^2$ such that the polar line of A with respect to the circumscribed conic is the same as the polar line of B with respect to the inscribed conic; see Figure [7](#). We will show that the operator R_P has all the same properties as Glick's operator G_P . From that it follows that persistence of convexity is equivalent to R_P being affine. But R_P is affine precisely when the inscribed and circumscribed conics are concentric.

To describe the properties of R_P , we use the following classical theorem of Kasner. Let D , as before, be the pentagram map, and let I be the map which sends a pentagon P to a new pentagon whose vertices are the tangency points of the sides of P and the inscribed conic; see left picture in Figure [8](#).

Theorem 5.1 (Kasner [3](#)). *The operations D and I on pentagons commute: $DI = ID$; see right picture in Figure [8](#).*

Corollary 5.2. *For any pentagon P , one has $R_{D(P)} = R_P$.*

Proof. Consider Figure [9](#). Observe that the sides of P are polar to the vertices of $I(P)$ with respect to the inscribed conic and to the vertices of $I^{-1}(P)$ with

FIGURE 8. Definition of the operation I and Kasner's theorem.

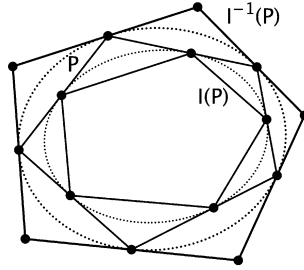


FIGURE 9. To the proof of Corollary 5.2

respect to the circumscribed conic. Therefore, $R_P(I^{-1}(P)) = I(P)$. But since the map I commutes with projective transformations, this is the same as to say that $R_P(P) = I^2(P)$. At the same time, we have

$$R_{D(P)}(P) = D^{-1}(R_{D(P)}(D(P))) = D^{-1}(I^2(D(P))) = I^2(P),$$

where the last equality follows from Kasner's theorem. Thus, the mappings R_P and $R_{D(P)}$ both map P to $I^2(P)$, while $R_{D(P)}^{-1} \circ R_P$ maps P to itself. But the only projective automorphism of a generic pentagon is the identity, so we must have $R_{D(P)} = R_P$ for generic, and hence, by continuity, for all polygons. $\square$

Remark 5.3. One can also use a more precise version of Kasner's theorem which takes into account labeling of vertices to show that R_P and $R_{D(P)}$ map every vertex of P to the same vertex of $I^2(P)$ and hence coincide.

Proposition 5.4. *Let P be a convex pentagon in the affine plane. Then the associated operator R_P has all the properties listed in Proposition 3.1. Namely, the following is true:*

- (1) *For the dual pentagon P^* , one has $R_{P^*} = R_P^*$.*
- (2) *One has $R_P(\text{conv}(P)) \subset \text{int}(\text{conv}(P))$.*
- (3) *The limit point $\bigcap_{k \geq 0} \text{conv}(D^k(P))$ of the pentagram orbit $D^k(P)$ is a fixed point of R_P . Moreover, it is the only fixed point of R_P in $\text{conv}(P)$.*

Proof. (1) The conic inscribed in P^* is the dual of the conic circumscribed about P , i.e. it is given by the quadratic form Q_C^{-1} . Likewise, the conic circumscribed about P^* is given by Q_I^{-1} . So,

$$R_{P^*} = (Q_C^{-1})^{-1} Q_I^{-1} = Q_C Q_I^{-1} = (Q_I^{-1} Q_C)^* = R_P^*.$$

- (2) Any point in the convex hull of P is inside the circumscribed conic (i.e. in the contractible component of the complement to that conic in $\mathbb{P}^2$), or belongs to that conic. Therefore, the polar line of that point with respect to the circumscribed conic is either outside or tangent to that conic, and the point polar to that line with respect to the inscribed conic is inside the inscribed conic and hence inside P . So, $R_P(\text{conv}(P)) \subset \text{int}(\text{conv}(P))$, as claimed.

- (3) By Corollary 5.2, we have $R_P = R_{D^k(P)}$, so $R_P(\text{conv}(D^k(P))) = R_{D^k(P)}(\text{conv}(D^k(P)))$, which, by above, is a subset of $\text{conv}(D^k(P))$. Therefore,

$$\begin{aligned} R_P \left(\bigcap_{k \geq 0} \text{conv}(D^k(P)) \right) \\ \subset \bigcap_{k \geq 0} R_P(\text{conv}(D^k(P))) \subset \bigcap_{k \geq 0} \text{conv}(D^k(P)), \end{aligned}$$

meaning that the limit point of the pentagram orbit $D^k(P)$ is indeed fixed by R_P . As for the uniqueness of the fixed point, it follows from other properties in the same way as for G_P . □

From this proposition it follows that for a convex pentagon P , all its successive images under the map S are convex if and only if the associated projective mapping R_P is affine. Indeed, the proof of the corresponding statement for G_P is based on Proposition 3.1, and since that proposition also holds for R_P , the result follows. Now, it remains to show that R_P is affine if and only if the conic inscribed in P and the conic circumscribed about P are concentric. To that end, recall that the center of the conic is, by definition, the point polar to the line at infinity with respect to that conic. Therefore, the image of the line at infinity under the map R_P is the line polar with respect to the inscribed conic to the center of the circumscribed conic. Thus, the centers coincide precisely when R_P preserves the line at infinity, i.e. is affine, as desired.

Remark 5.5. Implication $4 \Rightarrow 1$ (P is a convex pentagon with concentric inscribed and circumscribed conics $\Rightarrow$ all pentagons $S^k(P)$ are convex) can also be established as follows. First assume that the inscribed and circumscribed conics of P are *confocal*. Then, as shown in [4], there exists an affine transformation D_P taking P to its pentagram image $D(P)$. Therefore, we have $S^k(P) = D_P^{-k}(P)$, and $S^k(P)$ is convex for every k . Now assume that the inscribed and circumscribed conics of P are concentric, but not necessarily confocal. Then there exists a (generally speaking, defined over $\mathbb{C}$) affine transformation Q which makes those concentric conics confocal. Then, since $Q(P)$ has confocal inscribed and circumscribed conics, all the pentagons $S^k(Q(P))$ are convex, and the same holds for $S^k(P) = Q^{-1}(S^k(Q(P)))$, as claimed.

Remark 5.6. Instead of proving the equivalence $1 \Leftrightarrow 4$, we could have proved $3 \Leftrightarrow 4$, i.e. G_P is affine $\Leftrightarrow R_P$ is affine, as follows. First, recall that by Clebsch's theorem for any pentagon P there exists a projective transformation D_P such that $D_P(P) = D(P)$; see e.g. [6, Theorem 2.1]. Furthermore, any pentagon is projectively equivalent to its dual; see e.g. [1, Proposition 5]. So, since $I(P)$ is polar to P and hence projectively equivalent to P^* , there exists a projective transformation such that $I(P) = I_P(P)$. It is then a direct corollary of Kasner's theorem that transformations D_P and I_P commute, cf. [5, Corollary 7]. Therefore, D_P commutes with $R_P = I_P^2$. Further, as observed in [2], D_P coincides, as a linear operator in 3-space, with $G_P - 3 \cdot \text{Id}$, so R_P commutes with G_P . Now, assume that R_P is affine. Then the line at infinity L is a fixed point of the dual operator R_P^* . Moreover, by Proposition 5.4, Item 3, it is the only fixed point of R_P^* in the interior of P^* . Then, since G_P^* commutes with R_P^* and preserves the interior of P^* , it must

preserve L , which means that G_P is affine. Likewise, if G_P is affine, then R_P is affine too, as desired.

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