



# Packing $(1, 1, 2, 4)$ -coloring of subcubic outerplanar graphs

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## ABSTRACT

For  $1 \leq s_1 \leq s_2 \leq \dots \leq s_k$  and a graph  $G$ , a  $\text{packing } (s_1, s_2, \dots, s_k)$ -coloring of  $G$  is a partition of  $V(G)$  into sets  $V_1, V_2, \dots, V_k$  such that, for each  $1 \leq i \leq k$  the distance between any two distinct  $x, y \in V_i$  is at least  $s_i + 1$ . The  $\text{packing chromatic number}$ ,  $\chi_p(G)$ , of a graph  $G$  is the smallest  $k$  such that  $G$  has a packing  $(1, 2, \dots, k)$ -coloring. It is known that there are trees of maximum degree 4 and subcubic graphs  $G$  with arbitrarily large  $\chi_p(G)$ . Recently, there was a series of papers on packing  $(s_1, s_2, \dots, s_k)$ -colorings of subcubic graphs in various classes. We show that every 2-connected subcubic outerplanar graph has a packing  $(1, 1, 2)$ -coloring and every subcubic outerplanar graph is packing  $(1, 1, 2, 4)$ -colorable. Our results are sharp in the sense that there are 2-connected subcubic outerplanar graphs that are not packing  $(1, 1, 3)$ -colorable and there are subcubic outerplanar graphs that are not packing  $(1, 1, 2, 5)$ -colorable. We also show subcubic outerplanar graphs that are not packing  $(1, 2, 2, 4)$ -colorable and not packing  $(1, 1, 3, 4)$ -colorable.

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## 1. Introduction

For a non-decreasing sequence  $S = (s_1, s_2, \dots, s_k)$  of positive integers, a  $\text{packing } S$ -coloring of a graph  $G$  is a partition of  $V(G)$  into sets  $V_1, \dots, V_k$  such that, for each  $1 \leq i \leq k$ , the distance between any two distinct  $x, y \in V_i$  is at least  $s_i + 1$ . A  $\text{packing } k$ -coloring is a packing  $(1, 2, \dots, k)$ -coloring. The  $\text{packing chromatic number}$ ,  $\chi_p(G)$  (we will use the abbreviation PCN for short), of a graph  $G$  is the minimum  $k$  such that  $G$  has a packing  $k$ -coloring.

Packing  $k$ -coloring was introduced in 2008 by Goddard, Hedetniemi, Hedetniemi, Harris and Rall [20] motivated by frequency assignment problems in broadcast networks. There are more than 50 papers on the topic (e.g. [1, 5, 6, 8–17, 19, 23, 24, 27, 28] and references in them). In particular, Fiala and Golovach [14] proved that finding the PCN of a graph is NP-complete even in the class of trees. Sloper [27] showed that the infinite complete ternary tree (every vertex has 3 child vertices) has unbounded PCN.

The question whether PCN is bounded in the class of subcubic graphs was discussed in several papers (e.g., in [10, 11, 19, 25, 27]) and answered in the negative in [2]. Brešar and Ferme [5] later provided an explicit family of subcubic graphs with unbounded PCN. This stimulated studying subclasses of subcubic graphs with bounded PCN.

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One of the studied classes was the class of subdivisions of subcubic graphs. Recall that the *subdivision*,  $D(G)$ , of a graph  $G$  is the graph obtained from  $G$  by replacing each edge by a path of length two. In particular, Gastineau and Togni [19] asked whether  $\chi_p(D(G)) \leq 5$  for every subcubic graph  $G$  and Brešar, Klavžar, Rall, and Wash [10] later conjectured this.

**Conjecture 1** (Brešar, Klavžar, Rall, and Wash [10]). *Let  $G$  be a subcubic graph. Then  $\chi_p(D(G)) \leq 5$ .*

In 2018 [3] it was shown that the PCN of the subdivision of every subcubic graph is at most 8. Gastineau and Togni [19] pointed out at the following connection between the bounds on packing  $S$ -colorings of a graph  $G$  and the packing  $S$ -colorings of  $D(G)$ :

**Proposition 2** (Gastineau and Togni [19], Proposition 1). *Let  $G$  be a graph and  $S = (s_1, \dots, s_k)$  be a non-decreasing sequence of integers. If  $G$  is packing  $S$ -colorable, then  $D(G)$  is packing  $(1, 2s_1 + 1, \dots, 2s_k + 1)$ -colorable.*

They [19] also proved that the Petersen graph is not packing  $(1, 1, k, k')$ -colorable for any  $k, k' \geq 2$  and suggested the study of packing  $(1, 1, 2, 2)$ -coloring to approach **Conjecture 1**. Brešar, Klavžar, Rall, and Wash [10] later verified that the Petersen graph admits a packing  $(1, 2, 3, 4, 5)$ -coloring. **Proposition 2** implies that if one can prove every subcubic graph except the Petersen graph is packing  $(1, 1, 2, 2)$ -colorable then  $\chi_p(D(G)) \leq 5$  for every subcubic graph. Gastineau and Togni [19] also asked the question whether the stronger claim holds: every subcubic graph except the Petersen graph is packing  $(1, 1, 2, 3)$ -colorable.

The problem whether every subcubic graph except the Petersen graph is packing  $(1, 1, 2, 2)$ -colorable is intriguing by itself. Some subclasses of subcubic graphs were shown to have such a coloring. In particular, Brešar, Klavžar, Rall, and Wash [10] showed that if  $G$  is a generalized prism of a cycle, then  $G$  is packing  $(1, 1, 2, 2)$ -colorable if and only if  $G$  is not the Petersen graph. Very recently, Liu, Liu, Rolek, and Yu [26] proved that every subcubic graph with maximum average degree less than  $\frac{30}{11}$  is  $(1, 1, 2, 2)$ -colorable and thus confirmed **Conjecture 1** for subcubic graph  $G$  with  $mad(G) < \frac{30}{11}$ .

Many similar colorings have also been considered (e.g. [3, 7, 18, 19, 21, 22]). In particular, Gastineau and Togni [19] showed that subcubic graphs are packing  $(1, 2, 2, 2, 2, 2, 2)$ -colorable and packing  $(1, 1, 2, 2, 2)$ -colorable. They also showed that every 3-irregular (has no adjacent vertices of degree 3) subcubic graph is packing  $(1, 2, 2, 2)$ -colorable and packing  $(1, 1, 2)$ -colorable. Gastineau [16] showed that determining whether a subcubic bipartite graph is packing  $(1, 1, 2)$ -colorable and whether a subcubic graph is  $(1, 1, 2)$ -colorable are both NP-complete. In [3] it was proved that subcubic graphs are packing  $(1, 1, 2, 2, 3, 3, k)$ -colorable with color  $k$  used at most once for each integer  $k \geq 4$ , and that 2-degenerate subcubic graphs are packing  $(1, 1, 2, 2, 3, 3)$ -colorable. Borodin and Ivanova [4] proved that every subcubic planar graph with girth at least 23 has a packing  $(2, 2, 2, 2)$ -coloring.

Packing colorings of subclasses of subcubic outerplanar graphs were first studied by Gastineau, Holub, and Togni [17], who showed upper bounds for PCN of 2-connected subcubic outerplanar graphs with conditions on the number of internal faces. Recently, Brešar, Gastineau and Togni [7] proved that the PCN of any 2-connected bipartite subcubic outerplanar graph is bounded by 7, which gives a partial answer to the question posed in several papers concerning the boundedness of the PCN in the class of planar subcubic graphs. Moreover, they proved that every triangle-free subcubic outerplanar graph has a packing  $(1, 2, 2, 2)$ -coloring (and thus a packing  $(1, 1, 2, 2)$ -coloring) and their result is sharp in the sense that there exists a subcubic outerplanar graph with no triangles that is not packing  $(1, 2, 2, 3)$ -colorable. They also showed that every bipartite outerplanar graph admits a packing  $S$ -coloring for  $S = (1, 3, \dots, 3)$ , where 3 appears  $\Delta$  (maximum degree) times. Their result is sharp in the sense that if one of the integers 3 is replaced by 4 in the sequence  $S$ , then there exist outerplanar bipartite graphs that do not admit a packing  $S$ -coloring. The following two interesting questions were also suggested by Brešar, Gastineau and Togni [7] for future research.

**Question 3** (Brešar, Gastineau and Togni [7]). *Is the PCN bounded in the class of 2-connected outerplanar subcubic graphs and is the PCN bounded in the class of 2-connected bipartite planar subcubic graphs?*

In this paper, we prove that every 2-connected subcubic outerplanar graph is packing  $(1, 1, 2)$ -colorable and every subcubic outerplanar graph is packing  $(1, 1, 2, 4)$ -colorable. Our results are sharp in the sense that there is a 2-connected subcubic outerplanar graph  $G$  that is not  $(1, 1, 3)$ -colorable (see **Example 5**) and not  $(1, 2, 2)$ -colorable (see **Example 6**); and there are subcubic outerplanar graphs that are neither  $(1, 2, 2, 4)$ -colorable, nor  $(1, 1, 3, 4)$ -colorable; there are also subcubic outerplanar graphs that are not  $(1, 1, 2, 5)$ -colorable.

**Theorem 4.** *Every 2-connected subcubic outerplanar graph  $G$  is packing  $(1, 1, 2)$ -colorable.*

**Example 5.** Let  $G$  be the graph obtained by starting with a four cycle  $C$  such that  $V(C) = \{u_1, u_2, u_3, u_4\}$  and  $u_iu_{i+1} \in E(G)$  for  $1 \leq i \leq 4$  ( $5 \equiv 1$ ). Then we add a path  $u_1v_1u_2$  of length two from  $u_1$  to  $u_2$  and a path  $u_3v_2u_4$  from  $u_3$  to  $u_4$ .

Assume  $G$  has a packing  $(1, 1, 3)$ -coloring. Since both  $u_1, u_2, v_1$  and  $u_3, u_4, v_2$  form triangles, at least one vertex of  $u_1, u_2, v_1$  and one vertex of  $u_3, u_4, v_2$  are colored with 3 respectively. But the diameter of  $G$  is 3, a contradiction.

**Example 6.** Let  $G$  be the same graph as used in **Example 5**. We show that  $G$  is not  $(1, 2, 2)$ -colorable.

Assume  $G$  has a packing  $(1, 2, 2)$ -coloring. Say the colors are  $1, 2_a, 2_b$ . Since both  $u_1, u_2, v_1$  and  $u_3, u_4, v_2$  form triangles, all three colors are used exactly once on  $u_1, u_2, v_1$  and  $u_3, u_4, v_2$  respectively. By symmetry, we assume that  $u_2$  is colored with a color in  $\{2_a, 2_b\}$ , say  $2_a$ . This contradicts the fact that  $2_a$  is used on the triangle  $u_3v_2u_4$ .

**Theorem 7.** Every subcubic outerplanar graph has a packing  $(1, 1, 2, 4)$ -coloring  $f$  such that

(A) Color 4 is used at most once within each block.

(B) if  $v$  is a vertex of degree at most 2 and is colored with 2, then there is no vertex of color 4 within distance two from  $v$ .

By [Proposition 2](#), we also have the following immediate corollary, which confirms [Conjecture 1](#) for subcubic outerplanar graphs.

**Corollary 8.** If  $G$  is a subcubic outerplanar graph, then  $\chi_p(D(G)) \leq 5$ . Moreover, if  $H$  is a 2-connected subcubic outerplanar graph, then  $\chi_p(D(G)) \leq 4$ .

**Proof.** [Proposition 2](#) implies that if  $G$  is packing  $(1, 1, 2, 2)$ -colorable then  $D(G)$  is packing  $(1, 3, 3, 5, 5)$ -colorable, which implies a packing  $(1, 2, 3, 4, 5)$ -coloring of  $D(G)$ . Similarly, since  $H$  is packing  $(1, 1, 2)$ -colorable,  $D(G)$  is packing  $(1, 3, 3, 5)$ -colorable and thus  $(1, 2, 3, 4)$ -colorable.  $\square$

The result of Liu, Liu, Rolek, and Yu [26] implies that every subcubic planar graph with girth at least 8 is packing  $(1, 1, 2, 2)$ -colorable. We would also like to ask the following questions.

**Question 9.** Is it true that every subcubic planar graph is packing  $(1, 1, 2, 2)$ -colorable?

**Question 10.** Is it true that every subcubic 2-connected outerplanar graph is packing  $(1, 2, 2, 2)$ -colorable?

## 2. Notation and preliminaries

We use  $N_G^d(u)$  to denote the set of all vertices that are at distance  $d$  from  $u$ . We will work with *outerplane* graphs, that is, outerplanar graphs with a fixed drawing where all vertices are on the outer face.

A *block* of a graph  $G$  is an inclusion maximal subgraph with no cut vertices. By definition, each block is either 2-connected or a  $K_2$ . In the former case, we call the block *nontrivial*. A block in a graph  $G$  is *pendant* if it contains at most one cut vertex of  $G$ .

Given an outerplane graph  $G$ , the *weak dual graph*,  $\mathcal{T}(G)$ , is the graph that has a vertex for every bounded face of the embedding, and an edge for every pair of bounded faces sharing at least one edge. Below when we say “face” we will mean an internal face, unless we explicitly say “outer face”. For a face  $F$  in an outerplane graph  $G$ , we denote by  $C(F)$  the chordless cycle in  $G$  that bounds  $F$ . It is well known that

$$\text{a plane graph is outerplane if and only if its weak dual is a forest.} \quad (1)$$

By an  $i$ -face we will mean a face of length  $i$ . In view of (1), we say that an internal face  $F$  in a outerplane graph  $G$  is *pendant*, if either  $F$  corresponds to a leaf in  $\mathcal{T}(G)$  and  $C(F)$  contains no cut vertices of  $G$  or  $C(F)$  induces a pendant block in  $G$ .

**Claim 11.** Each 3-face in a 2-connected subcubic outerplane graph is pendant.

**Proof.** Let  $F$  be a 3-face with  $C(F) = xyz$  in a 2-connected subcubic outerplane graph  $G$ . If, say  $d(x) = 2$ , then the edges  $xy$  and  $xz$  are on the boundary of the outer face, and so  $F$  is pendant.

So, suppose  $d(x) = d(y) = d(z) = 3$ . Let  $x'$ ,  $y'$  and  $z'$  be the neighbors of  $x$ ,  $y$  and  $z$  respectively outside of  $\{x, y, z\}$  (some of them can coincide). Since  $G$  is 2-connected, all  $x'$ ,  $y'$  and  $z'$  are in the same component of  $G - \{x, y, z\}$ . But then  $G$  contains a  $K_4$ -minor, which implies that  $G$  is not outerplane, a contradiction.  $\square$

## 3. Proof of Theorem 4

All our  $(1, 1, 2)$ -colorings will use the colors  $1_a$ ,  $1_b$  and 2.

**Proof.** Suppose the theorem is false and  $G$  is a smallest 2-connected subcubic outerplane graph that has no  $(1, 1, 2)$ -coloring. Let  $n = |V(G)|$ . Then  $n \geq 4$ , since otherwise we can color all vertices of  $G$  with different colors.

**Claim 12.** Each pendant face in  $G$  is a 3-face.

**Proof.** Suppose  $F$  is a pendant face in  $G$  with  $C(F) = u_1u_2 \dots u_ku_1$  where  $k \geq 4$ . If  $V(G) = \{u_1, u_2, \dots, u_k\}$  then we color  $u_1$  with 2 and the remaining vertices alternately with  $1_a$  and  $1_b$ . So suppose  $d(u_1) = d(u_k) = 3$  and  $d(u_2) = d(u_3) = \dots = d(u_{k-1}) = 2$ . Let  $G' = G - \{u_2, \dots, u_{k-1}\}$ . Then  $G'$  is a 2-connected outerplane graph. By the minimality of  $G$ ,  $G'$  has a packing  $(1, 1, 2)$ -coloring  $f$ .

We can extend  $f$  to the vertices  $u_2, \dots, u_{k-1}$  by coloring them alternately with  $1_a$  and  $1_b$ , unless  $k$  is odd and  $\{f(u_1), f(u_k)\} = \{1_a, 1_b\}$ . In this exceptional case, assuming  $f(u_1) = 1_a$ ,  $f(u_k) = 1_b$ , we let  $f(u_2) = 1_b$ ,  $f(u_3) = 2$ , and color  $u_4, u_5, \dots, u_{k-1}$  alternately with  $1_a$  and  $1_b$ . Thus when  $k \geq 4$ , in all cases, we get a packing  $(1, 1, 2)$ -coloring of  $G$ , a contradiction.  $\square$

**Claim 13.** *G has no face of even length.*

**Proof.** Suppose  $G$  has a face  $F_0$  with  $C(F_0) = u_1u_2\dots u_ku_1$ , where  $k$  is even. Let  $G_1, \dots, G_\ell$  be the components of  $G - \{u_1, u_2 \dots u_k\}$ . Since  $G$  is 2-connected and outerplane, each  $G_i$  has exactly two neighbors on  $C(F_0)$ , and these neighbors are consecutive on  $C(F_0)$ . Let these neighbors be  $u_{j(i)}$  and  $u_{j(i)+1}$ , and let the neighbors of  $u_{j(i)}$  and  $u_{j(i)+1}$  in  $V(G_i)$  be  $v_i$  and  $v'_i$  (possibly,  $v'_i = v_i$ ). If  $v'_i \neq v_i$  and  $v_i v'_i \notin E(G)$  then we define  $G'_i = G_i + v_i v'_i$ , otherwise we let  $G'_i = G_i$ .

With these definitions, if  $v'_i = v_i$  then, by [Claim 11](#),  $G'_i$  is a single vertex; otherwise,  $G'_i$  is a 2-connected outerplane graph. So by the minimality of  $G$ , each  $G'_i$  has a packing  $(1, 1, 2)$ -coloring  $f'_i$  such that

$$\text{if } v'_i \neq v_i, \text{ then } f'_i(v'_i) \neq f'_i(v_i). \quad (2)$$

We now define a  $(1, 1, 2)$ -coloring  $f$  of  $G$  as follows:

- (i) For  $1 \leq j \leq \frac{k}{2}$ , we let  $f(u_{2j-1}) = 1_a$  and  $f(u_{2j}) = 1_b$ .
- (ii) If  $1 \leq i \leq \ell$  and  $v'_i = v_i$ , then  $G'_i$  is a single vertex  $v_i$  and we let  $f(v_i) = 2$ .
- (iii) If  $1 \leq i \leq \ell$  and  $v'_i \neq v_i$ , then by (2) and the fact that  $\{f(u_{j(i)}), f(u_{j(i)+1})\} = \{1_a, 1_b\}$ , we can switch the names of the colors  $1_a$  and  $1_b$  in  $f'_i$  so that  $f'_i(v_i) \neq f(u_{j(i)})$  and  $f'_i(v'_i) \neq f(u_{j(i)+1})$ . In this case, we let  $f(v) = f'_i(v)$  for each  $v \in V(G_i)$ .

By construction,  $f$  is a packing  $(1, 1, 2)$ -coloring on  $G$ , since the vertices of color 2 in different  $G_i$  are at distance at least 3.  $\square$

If  $G$  has only one face apart from the outer face, then  $G$  is an odd cycle, say  $u_1u_2\dots u_{2k+1}u_1$ , and we can color its vertices apart from  $u_{2k+1}$  alternately with  $1_a$  and  $1_b$  and let  $f(u_{2k+1}) = 2$ . Thus, suppose  $G$  has at least two faces. Let  $F_0$  be a pendant face corresponding to an end vertex in a longest path in  $\mathcal{T}(G)$ . By [Claim 12](#),  $C(F_0)$  is a 3-cycle.

Let  $F'_0$  be the unique face adjacent to  $F_0$ . By our choice of  $F_0$ ,

$$F'_0 \text{ is adjacent to at most one non-pendant face.} \quad (3)$$

If  $|C(F'_0)| = 3$ , then, by [Claim 11](#),  $G = K_4 - e$ . In this case, we color the two vertices of degree 2 in  $G$  with  $1_a$  and the remaining two vertices with  $1_b$  and 2.

Thus, we may assume  $F'_0$  is a face with an odd length  $k \geq 5$ . Since  $F'_0$  is an odd face and each face adjacent to  $F'_0$  (apart from the outer face) shares exactly two vertices with  $F'_0$ , at least one vertex in  $C(F'_0)$  has degree two in  $G$ . Fix one such vertex, say  $w_1$ . Let  $C(F'_0) = w_1w_2\dots w_kw_1$ .

If one of  $w_2, w_k$  has degree two, say  $w_2$ , then we delete  $w_1, w_2$  and add the edge  $w_3w_k$  (it is not in  $E(G)$  since  $F'_0$  has length at least five). This results in a 2-connected subcubic outerplane graph  $G'$  with fewer vertices. By the minimality of  $G$ ,  $G'$  has a packing  $(1, 1, 2)$ -coloring  $f'$ . If  $2 \notin \{f'(w_3), f'(w_k)\}$ , say  $f'(w_3) = 1_a$  and  $f'(w_k) = 1_b$ , then we color  $w_1, w_2$  with  $1_a, 1_b$ . If  $2 \in \{f'(w_3), f'(w_k)\}$ , say  $f'(w_3) = 2$  and  $f'(w_k) = 1_b$ , then we color  $w_1, w_2$  with  $1_a, 1_b$ .

Thus, we may assume that both neighbors of  $w_1$ , i.e.,  $w_2$  and  $w_k$ , have degree three. Let  $G_1, \dots, G_\ell$  be the components of  $G - \{w_1, w_2 \dots w_k\}$ . Since  $k \geq 5$  and  $d(w_2) = d(w_k) = 3$ ,  $\ell \geq 2$ . As in the proof of [Claim 13](#), each  $G_i$  has exactly two neighbors on  $C(F_0)$ , and these neighbors are consecutive on  $C(F_0)$ . Let these neighbors be  $w_{j(i)}$  and  $w_{j(i)+1}$ , and let the neighbors of  $w_{j(i)}$  and  $w_{j(i)+1}$  in  $V(G_i)$  be  $v_i$  and  $v'_i$  (possibly,  $v'_i = v_i$ ). We can rename  $G_i$ s so that  $j(1) < j(2) < \dots < j(\ell)$ . By (3) and [Claim 12](#), at most one of  $G_1, \dots, G_\ell$  is not a single vertex. By the symmetry between  $w_2$  and  $w_k$ , we may assume that  $G_1$  is a single vertex.

We start coloring by letting  $f(w_2) = 2$  and coloring the remaining vertices of  $C(F'_0)$  alternately with  $1_a$  and  $1_b$ . Then let  $f(v_1) = 1_b$  and color the unique vertex in each other single-vertex  $G_i$  with 2. If  $G - \{w_1, w_2 \dots w_k\}$  has no larger components, then we are done. Otherwise, suppose  $G_{i_0}$  is the unique “large” component of  $G - \{w_1, w_2 \dots w_k\}$ . If  $v_{i_0}v'_{i_0} \notin E(G)$ , then we define  $G'_{i_0} = G_{i_0} + v_{i_0}v'_{i_0}$ , otherwise we let  $G'_{i_0} = G_{i_0}$ . By the minimality of  $G$ ,  $G'_{i_0}$  has a packing  $(1, 1, 2)$ -coloring  $f'$  such that  $f'(v'_{i_0}) \neq f'(v_{i_0})$ . As in the proof of [Claim 13](#), the facts that  $\{f(w_{j(i_0)}), f(w_{j(i_0)+1})\} = \{1_a, 1_b\}$  and  $f'(v'_{i_0}) \neq f'(v_{i_0})$ , we can switch the names of the colors  $1_a$  and  $1_b$  in  $f'$  so that  $f'(v_{i_0}) \neq f(w_{j(i_0)})$  and  $f'(v'_{i_0}) \neq f(w_{j(i_0)+1})$ . After that, we let  $f(v) = f'(v)$  for each  $v \in V(G_{i_0})$ .

So, we obtain a packing  $(1, 1, 2)$ -coloring of  $G$ . This contradicts the choice of  $G$  and proves the theorem.

#### 4. Proof of Theorem 7

By a *feasible* coloring of  $G$  we call a coloring of  $G$  with colors  $1_a, 1_b, 2, 4$  such that the distance between two vertices of color  $i_x$  is at least  $i + 1$  for all  $i \in \{1, 2, 4\}$  and  $x \in \{a, b\}$ , and  $f$  satisfies conditions (A) and (B) of [Theorem 7](#).

Suppose, [Theorem 7](#) fails and  $G$  is a smallest outerplane graph not admitting a feasible coloring. Clearly,

$$G \text{ is connected and } \delta(G) \geq 2. \quad (4)$$

It follows that every pendant block is nontrivial. So if  $G$  has only one non-trivial block, then it has no other blocks. In this case,  $G$  has a packing  $(1, 1, 2)$ -coloring by [Theorem 4](#). Hence we may assume that  $G$  has at least two blocks, and thus at least two pendant blocks (which are nontrivial).

**Claim 14.** *Each pendant face in  $G$  is a 3-face.*

**Proof.** Suppose  $F$  is a pendant face with  $C(F) = u_1u_2 \dots u_ku_1$  where  $k \geq 4$ . If  $F$  does not contain a cut vertex of  $G$ , then we can repeat the proof of [Claim 12](#) word by word. Note that when we color  $u_3$  with 2 in the second paragraph of [Claim 12](#), it is at distance at least two from  $u_1$  and  $u_k$  respectively, and thus condition (B) in [Theorem 7](#) is satisfied. So, suppose  $u_1$  is a cut vertex of  $G$ , and its neighbor outside of  $C(F)$  is  $v$ . Recall that by the definition of pendant faces, in this case  $C(F)$  induces a pendant block in  $G$ .

By the minimality of  $G$ ,  $G - \{u_1, u_2, \dots, u_k\}$  has a feasible coloring  $f$ . We extend  $f$  to  $G$  as follows. First choose  $f(u_1) \in \{1_a, 1_b\} - f(v)$ . By symmetry, assume  $f(u_1) = 1_b$ . If  $k$  is even, then we can color  $u_2, \dots, u_k$  alternately with  $1_a$  and  $1_b$ . Otherwise, let  $f(u_2) = 1_a, f(u_3) = 2$  and color  $u_4, \dots, u_k$  alternately with  $1_b$  and  $1_a$ . In both cases, we obtain a feasible coloring of  $G$ , a contradiction.  $\square$

Let  $G_0$  be one of the pendant blocks. Let the cut edge connecting  $G_0$  and  $G - G_0$  be  $u_1v_1$  with  $v_1 \in V(G_0)$ . Let  $F_0$  be the face in  $G_0$  containing  $v_1$  with  $C(F_0) = v_1v_2 \dots v_kv_1$ . Let  $N(u_1) = \{u_2, u_3, v_1\}$ . By the minimality of  $G$ , graph  $G' = G - G_0$  has a feasible coloring  $f$ .

**Case 1:**  $G_0$  is a cycle. By [Claim 14](#),  $k = 3$ .

**Case 1.1:**  $f(u_1) \in \{1_a, 1_b\}$ , say  $f(u_1) = 1_a$ . We let  $f(v_1) = 1_b, f(v_2) = 1_a$  and  $f(v_3) = 2$ .

**Case 1.2:**  $f(u_1) = 2$ . Since  $f$  is a feasible coloring of  $G'$  and  $d_{G'}(u_1) \leq 2$ , by (B) in the claim of [Theorem 7](#), there is no vertex in  $G - G_0$  colored with 4 within distance two from  $u_1$ . Then we let  $f(v_1) = 1_a, f(v_2) = 1_b$  and  $f(v_3) = 4$ .

**Case 1.3:**  $f(u_1) = 4$ . If  $2 \in \{f(u_2), f(u_3)\}$ , say  $f(u_2) = 2$  and  $f(u_3) = 1_a$ , then we recolor  $u_1$  with  $1_b$  and obtain Case 1.1. Thus we may assume that  $\{f(u_2), f(u_3)\} = \{1_a, 1_b\}$ . In this case, Then we let  $f(v_1) = 2, f(v_2) = 1_b$  and  $f(v_3) = 1_a$ .

**Case 2:**  $F_0$  is adjacent in  $\mathcal{T}(G_0)$  only to pendant faces. Let these faces be  $F_1, \dots, F_\ell$  ordered so that the indices of the vertices in  $V(F_i) \cap V(F_0)$  are larger than the indices of the vertices in  $V(F_j) \cap V(F_0)$  if and only if  $i > j$ . Suppose the common vertices of  $C(F_0)$  and  $C(F_1)$  are  $v_p$  and  $v_{p+1}$ . By [Claim 14](#), each  $F_i$  is a 3-face. If  $k$  is even, we can choose  $f(v_1) \in \{1_a, 1_b\} - f(u_1)$ , then color alternately with  $1_a$  and  $1_b$  all vertices  $v_2, \dots, v_k$ , and for each  $i = 1, \dots, \ell$ , color the unique vertex  $w_i \in C(F_i) - C(F_0)$  with 2. So we may assume  $k$  is odd.

**Case 2.1:**  $f(u_1) \neq 2$ . Let  $f(v_p) = 2$  and color alternately with  $1_a$  and  $1_b$  all vertices in  $V(F_0) - v_p$  so that  $f(v_1) \neq f(u_1)$ . For all  $2 \leq i \leq \ell$ , color  $w_i$  with 2 and choose  $f(w_i) \in \{1_a, 1_b\} - f(v_{p+1})$ .

**Case 2.2:**  $f(u_1) = 2$ . As in Case 1.2, within distance two of  $u_1$  there is no vertex in  $G - G_0$  colored with 4. We color vertices in  $G_0$  almost as in Case 2.1, except we color  $v_p$  with 4. Observe that conditions (A) and (B) in the claim of [Theorem 7](#) hold for the new coloring, and that  $v_p$  is at distance at least 2 from  $u_1$  which in turn is at distance at least 5 from other vertices of color 4.

**Case 3:**  $F_0$  is adjacent to some non-pendant face. Let  $R$  be a pendant face of  $G_0$  that has the largest distance from  $F_0$  in the weak dual of  $G_0$ . By the description of Case 3, this distance is at least two. Let  $R_0$  be the face that  $R$  is adjacent to. By the choice of  $R$ ,  $R_0$  is adjacent to only one non-pendant face, say  $R'_0$ .

Let  $C(R_0) = x_1x_2 \dots x_rx_1$  and  $V(R_0) \cap V(R'_0) = \{x_1, x_2\}$ . Let  $R_1, \dots, R_m$  be the pendant faces that are adjacent to  $R_0$  arranged in the order of  $C(R_0)$ . Recall that by [Claim 14](#), each pendant  $R_i$  is a 3-face. Assume that for  $i = 1, \dots, m$ ,  $V(R_i) \cap V(R_0) = \{x_{q_i}, x_{q_i+1}\}$  and  $V(R_i) \setminus V(R_0) = \{y_i\}$ .

**Case 3.1:**  $R'_0$  has only three vertices; say the common neighbor of  $x_1$  and  $x_2$  in  $R'_0$  is  $x_0$ . Then by our construction,  $\bigcup_{i=0}^m V(R_i) \cup \{x_0\}$  comprises  $V(G_0)$ , and  $x_0$  is the vertex in  $G_0$  that is adjacent to  $u_1$ . By the minimality of  $G$ ,  $G'$  has a feasible coloring  $f$ . Recall that the neighbor of  $x_0$  in  $V(G - G_0)$  is  $u_1$ .

**Case 3.1.1:**  $f(u_1) \in \{1_a, 1_b\}$ . Then we color  $x_1$  with 2 and  $x_0, x_2, x_3, \dots, x_r$  alternately with  $1_a$  and  $1_b$  so that  $f(x_0) \neq f(u_1)$ . After that, we let  $f(y_m) = 4$ , and  $f(y_i) = 2$  for all  $1 \leq i \leq m-1$ . Then the coloring will be a packing  $(1, 1, 2, 4)$ -coloring and the conditions (A) and (B) will hold.

**Case 3.1.2:**  $f(u_1) = 2$ . Since  $d_{G'}(u_1) \leq 2$ , the distance from  $u_1$  to a vertex of color 4 in  $G'$  is at least 3. Then we color  $x_1$  with 4 and color  $x_{q_m+1}$  with 2,  $x_0, x_2, x_3, \dots, x_{q_m}, y_m$  alternately with  $1_a$  and  $1_b$ ,  $x_{q_m+2}, \dots, x_r$  (if  $q_m+1 < r$ ) alternatively with  $1_a$  and  $1_b$ . After that, if  $m \geq 2$  then we let  $f(y_i) = 2$  for all  $1 \leq i \leq m-1$ .

**Case 3.1.3:**  $f(u_1) = 4$ . If  $\{f(u_2), f(u_3)\} \neq \{1_a, 1_b\}$ , then we can recolor  $u_1$  with a color in  $\{1_a, 1_b\} \setminus \{f(u_2), f(u_3)\}$  and get Case 3.1.1. Otherwise,  $2 \notin \{f(u_2), f(u_3)\}$ . By [Theorem 4](#),  $G_0$  has a packing  $(1, 1, 2)$ -coloring  $f_0$ . Since the vertices  $x_0, x_1, x_2$  have degree 3 in  $G$ , the coloring  $f \cup f_0$  will be a packing  $(1, 1, 2, 4)$ -coloring of  $G$  satisfying (A) and (B).

**Case 3.2:**  $R'_0$  has at least four vertices. For  $j = 1, 2$ , let  $x'_j$  be the neighbor of  $x_j$  on  $C(R'_0)$  distinct from  $x_{3-j}$ . By the case,  $x'_2 \neq x'_1$ . Let  $G''$  be obtained from  $G - \bigcup_{i=0}^m V(R_i)$  by adding edge  $x'_1x'_2$  if this edge is not in  $G$ . By the minimality of  $G$ , the subcubic outerplane graph  $G''$  has a feasible coloring  $f$ . Since  $x'_1x'_2 \in E(G'')$ , by symmetry, we may assume  $f(x'_2) \neq 2$ . We color  $x_{q_1}$  with 2 and the remaining vertices of  $R_0$  alternately with  $1_a$  and  $1_b$  so that  $f(x_1) \neq f(x'_1)$  and hence  $f(x_2) \neq f(x'_2)$ . We can provide the two last inequalities because if  $f(x'_1), f(x'_2) \in \{1_a, 1_b\}$ , then  $f(x'_1) \neq f(x'_2)$ .

After that, we choose  $f(y_1) \in \{1_a, 1_b\} - f(x_{q_1+1})$  and let  $f(y_j) = 2$  for  $j = 2, 3, \dots, m$ . We obtain a feasible coloring of  $G$ , a contradiction. This proves [Theorem 7](#).

## 5. Sharpness of the bound

To show the sharpness of our result that every subcubic outerplanar graph is  $(1, 1, 2, 4)$ -colorable, we show that there are subcubic outerplanar graphs that are not  $(1, 1, 2, 5)$ -colorable, not  $(1, 2, 2, 4)$ -colorable, and not  $(1, 1, 3, 4)$ -colorable.

We first show that there is a subcubic outerplanar graph that is not  $(1, 1, 2, 5)$ -colorable.

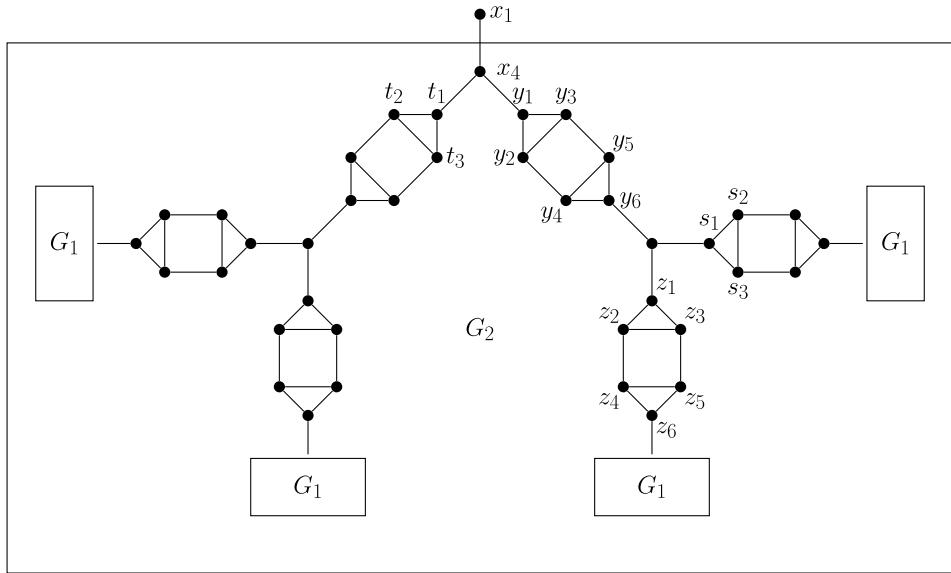
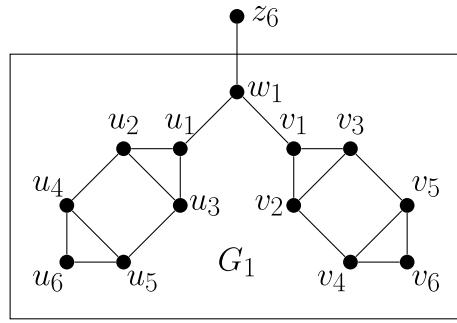
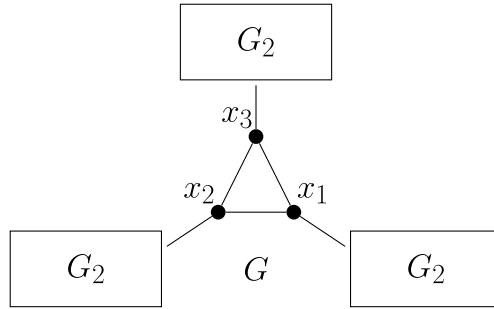
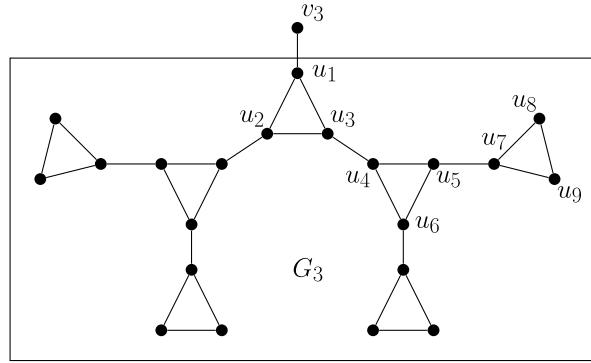
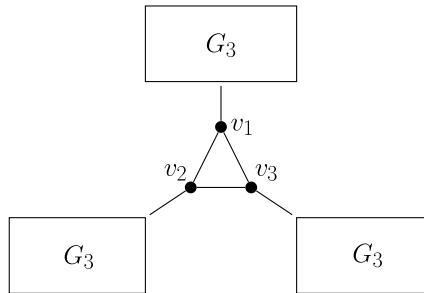
Fig. 1. Gadget  $G_2$ .Fig. 2. Gadget  $G_1$ .

Fig. 3. The construction.

**Example 15.** Our construction is the graph  $G$  in Fig. 3, where each of the gadgets  $G_2$  is the graph in Fig. 1 without the vertex  $x_1$  (the graph surrounded by the rectangle), where each of the gadgets  $G_1$  used in  $G_2$  is the graph in Fig. 2 without the vertex  $z_6$  (the graph surrounded by the rectangle). We show that  $G$  is not packing  $(1, 1, 2, 5)$ -colorable.

**Claim 16.** In any packing  $(1, 1, 2, 5)$ -coloring of  $G_1$  vertex  $z_6$  cannot be colored with 5.

**Proof.** Suppose  $z_6$  is colored with 5. Then no vertex in  $G_1$  can be colored with 5 since the farthest vertices from  $z_6$  are  $u_6, v_6$ , each of them is at distance 5 from  $z_6$ . Then exactly one vertex of each of the four triangles,  $u_1u_2u_3, u_4u_5u_6, v_1v_2v_3$ ,

Fig. 4. Gadget  $G_3$ .Fig. 5. construction  $H$ .

$v_4v_5v_6$  is colored with 2. But the only way to use color 2 in  $u_1u_2u_3$  and  $u_4u_5u_6$  is to color vertices  $u_1$  and  $u_6$  with 2, and the only way to use color 2 in  $v_1v_2v_3$  and  $v_4v_5v_6$  is to color vertices  $v_1$  and  $v_6$  with 2, which is impossible since  $u_1$  and  $v_1$  are at distance two.  $\square$

Suppose  $G$  has a packing  $(1, 1, 2, 5)$ -coloring  $f$ . Then

each of triangles in  $G$  has a vertex of color 2 or 5. (5)

In particular, a vertex in  $\{x_1, x_2, x_3\}$  is colored with 2 or 5. By symmetry, we may assume  $f(x_1) \in \{2, 5\}$ . By [Claim 16](#) applied to the top of [Fig. 1](#), we then have  $f(x_1) = 2$ . Since triangles  $y_1y_2y_3$  and  $t_1t_2t_3$  are too close to each other to both have a vertex of color 5, in view of (5) one of them has a vertex of color 2. By symmetry, we may assume it is  $y_1y_2y_3$ . Since  $y_1$  is at distance two from  $x_1$ , one of  $y_2$  and  $y_3$ , say  $y_2$ , is colored with 2. Then  $\{y_4, y_5, y_6\}$  does not have vertices of color 2, and hence it has a vertex of color 5. By [Claim 16](#) applied to right part of [Fig. 1](#), this vertex is not  $y_6$  and thus belongs to  $\{y_4, y_5\}$ . Then both triangles  $s_1s_2s_3$  and  $z_1z_2z_3$  have to use color 2, and since we cannot use 2 at  $z_1$  and  $s_1$  at the same time, we may assume by symmetry that  $z_2$  is colored with 2. This implies we need to use 5 at a vertex of the triangle  $z_4z_5z_6$  and this vertex must be  $z_6$  since  $z_4$  and  $z_5$  are at distance 5 from  $y_4$ , which contradicts [Claim 16](#).

Now we show there is a graph that is not  $(1, 2, 2, 4)$ -colorable and not  $(1, 1, 3, 4)$ -colorable.

**Example 17.** Our construction is the graph  $H$  in [Fig. 5](#), where each of the gadgets  $G_3$  is the graph in [Fig. 4](#) without the vertex  $v_3$  (the graph surrounded by the rectangle). We now show that  $H$  is not  $(1, 2, 2, 4)$ -colorable and not  $(1, 1, 3, 4)$ -colorable.

**Claim 18.**  $H$  is not  $(1, 2, 2, 4)$ -colorable.

**Proof.** Let the colors be  $1, 2_a, 2_b, 4$ .

**Case 1:** The three colors used on  $v_1, v_2, v_3$  are  $1, 2_a, 2_b$ . Say  $v_3$  is colored with  $2_b$  and  $v_2$  is colored with  $2_a$ . Then the color 4 must be used on the triangle  $u_1u_2u_3$  and this vertex cannot be  $u_1$  since otherwise we cannot color the triangle in the gadget hang on  $v_2$  which corresponds to  $u_1u_2u_3$  (we need to use color 4 on this triangle as well). We assume by symmetry that  $u_3$  is colored with 4. Then one of  $u_5$  and  $u_6$  is colored with a color in  $\{2_a, 2_b\}$ , say  $u_5$  is colored with  $2_a$ . But then we cannot use color  $2_a$  and color 4 on the triangle  $u_7u_8u_9$ , a contradiction.

**Case 2:** The three colors used on  $v_1, v_2, v_3$  are  $1, 2_a, 4$ . Say  $v_2$  is colored with 4 and  $v_3$  is colored with  $2_a$ . Then the vertices  $u_1, u_2, u_3$  have to choose colors from  $\{1, 2_b\}$ , a contradiction.

**Case 3:** The three colors used on  $v_1, v_2, v_3$  are  $2_a, 2_b, 4$ . Say  $v_2$  is colored with 4 and  $v_3$  is colored with  $2_a$ . Similarly to Case 2, we reach a contradiction.  $\square$

**Claim 19.**  $H$  is not  $(1, 1, 3, 4)$ -colorable.

**Proof.** Let the colors be  $1_a, 1_b, 3, 4$ . Since  $v_1v_2v_3$  is a triangle, at least one of the vertices  $v_1, v_2, v_3$ , say  $v_3$ , is colored with a color in  $\{3, 4\}$ .

**Case 1:**  $v_3$  is colored with 3. Then  $u_1$  cannot be colored with 4 since otherwise the triangle in the gadget hang on  $v_2$  which corresponds to  $u_1u_2u_3$  can only choose colors from  $\{1_a, 1_b\}$ , a contradiction. Thus one of the vertices in  $\{u_2, u_3\}$  is colored with 4, say  $u_3$ . Since  $v_3$  and  $u_4$  have distance 3,  $u_4$  cannot be colored with 3 and one of  $u_5, u_6$  is colored with 3, say  $u_5$ . But then the vertices  $u_7, u_8, u_9$  have to choose colors from  $\{1_a, 1_b\}$ , a contradiction.

**Case 2:**  $v_3$  is colored with 4. Then one of the vertices in  $\{u_4, u_5, u_6\}$  is colored with 3. But then the vertices  $u_7, u_8, u_9$  have to choose colors from  $\{1_a, 1_b\}$ , a contradiction.  $\square$

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