

A generalization of the Bollobás set pairs inequality

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Abstract

The Bollobás set pairs inequality is a fundamental result in extremal set theory with many applications. In this paper, for $n \geq k \geq t \geq 2$, we consider a collection of k families $\mathcal{A}_i : 1 \leq i \leq k$ where $\mathcal{A}_i = \{A_{i,j} \subset [n] : j \in [n]\}$ so that $A_{1,i_1} \cap \cdots \cap A_{k,i_k} \neq \emptyset$ if and only if there are at least t distinct indices i_1, i_2, \dots, i_k . Via a natural connection to a hypergraph covering problem, we give bounds on the maximum size $\beta_{k,t}(n)$ of the families with ground set $[n]$.

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1 Introduction

A central topic of study in extremal set theory is the maximum size of a family of subsets of an n -element set subject to restrictions on their intersections. Classical theorems in the area are discussed in Bollobás [2]. In this paper, we generalize one such theorem, known as the Bollobás set pairs inequality or two families theorem [3]:

Theorem 1. (Bollobás) *Let $\mathcal{A} = \{A_1, A_2, \dots, A_m\}$ and $\mathcal{B} = \{B_1, B_2, \dots, B_m\}$ be families of finite sets, such that $A_i \cap B_j \neq \emptyset$ if and only if $i, j \in [m]$ are distinct. Then*

$$\sum_{i=1}^m \binom{|A_i \cup B_i|}{|A_i|}^{-1} \leq 1. \quad (1)$$

For convenience, we refer to a pair of families \mathcal{A} and \mathcal{B} satisfying the conditions of Theorem 1 as a *Bollobás set pair*. The inequality above is tight, as we may take the pairs (A_i, B_i) to be distinct partitions of a set of size $a + b$ with $|A_i| = a$ and $|B_i| = b$ for $1 \leq i \leq \binom{a+b}{a}$.

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The latter inequality was proved for $a = 2$ by Erdős, Hajnal and Moon [5], and in general has a number of different proofs [11, 12, 14, 17, 18]. A geometric version was proved by Lovász [17, 18], who showed that if A_1, A_2, \dots, A_m and B_1, B_2, \dots, B_m are respectively a -dimensional and b -dimensional subspaces of a linear space and $\dim(A_i \cap B_j) = 0$ if and only if $i, j \in [m]$ are distinct, then $m \leq \binom{a+b}{a}$.

1.1 Main Theorem

Theorem 1 has been generalized in a number of different directions in the literature [6, 9, 13, 16, 21, 24]. In this paper, we give a generalization of Theorem 1 from the case of two families to $k \geq 3$ families of sets with conditions on the k -wise intersections. For $2 \leq t \leq k$, a *Bollobás (k, t) -tuple* is a sequence $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$ of set families $\mathcal{A}_j = \{A_{j,i} : 1 \leq i \leq m\}$ where $\bigcap_{j=1}^k A_{j,i_j} \neq \emptyset$ if and only if at least t of the indices i_1, i_2, \dots, i_k are distinct. We refer to m as the *size* of the Bollobás (k, t) -tuple. Let $[m]_{(t)}$ denote the set of sequences of t distinct elements of $[m]$ and fix a surjection $\phi : [k] \rightarrow [t]$. For $\sigma \in [m]_{(t-1)}$, set $\sigma(t) = \sigma(1)$ and define $A_{1,\sigma}(\phi) = \bigcap_{j:\phi(j)=1} A_{j,\sigma(1)}$ and, for $2 \leq j \leq t$, we define

$$A_{j,\sigma}(\phi) = \bigcap_{h:\phi(h)=j} A_{h,\sigma(j)} \setminus \bigcup_{h=1}^{j-1} A_{h,\sigma}(\phi).$$

Using this notation, we generalize (1) as follows:

Theorem 2. *Let $k \geq t \geq 2$ and $m \geq t$, let $\phi : [k] \rightarrow [t]$ be a surjection, and let $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$ be a Bollobás (k, t) -tuple of size m . Then*

$$\sum_{\sigma \in [m]_{(t-1)}} \left(\frac{|A_{1,\sigma}(\phi) \cup A_{2,\sigma}(\phi) \cup \dots \cup A_{t,\sigma}(\phi)|}{|A_{1,\sigma}(\phi)| |A_{2,\sigma}(\phi)| \dots |A_{t,\sigma}(\phi)|} \right)^{-1} \leq 1. \quad (2)$$

We show in Section 2.1 that this inequality is tight for all $k \geq t = 2$, but do not have an example to show that this inequality is tight for any $t > 2$.

For $n \geq k \geq t \geq 2$, let $\beta_{k,t}(n)$ denote the maximum m such that there exists a Bollobás (k, t) -tuple of size m consisting of subsets of $[n]$. Then (1) gives $\beta_{2,2}(n) \leq \binom{n}{\lfloor n/2 \rfloor}$ which is tight for all $n \geq 2$. Letting $H(q) = -q \log_2 q - (1-q) \log_2 (1-q)$ denote the standard binary entropy function, we prove the following theorem:

Theorem 3. *For $k \geq 3$ and large enough n ,*

$$\frac{1}{k} \leq \frac{\log_2 \beta_{k,2}(n)}{n} \leq H\left(\frac{1}{k}\right) \leq \frac{\log_2(ke)}{k}. \quad (3)$$

For $k \geq t \geq 3$ and large enough n ,

$$\frac{\log_2 e}{\binom{k}{t-1}(t+1)t^{t-1}} \leq \frac{\log_2 \beta_{k,t}(n)}{n} \leq \frac{2}{\binom{k}{t-1}(t-1)^{t-3}}. \quad (4)$$

This determines $\log_2 \beta_{k,2}(n)$ up to a factor of order $\log_2 k$ and $\log_2 \beta_{k,t}(n)$ up to a factor of order t^3 . We leave it as an open problem to determine the asymptotic value of $(\log_2 \beta_{k,t}(n))/n$ as $n \rightarrow \infty$ for any $k \geq 3$ and $t \geq 2$. A natural source for lower bounds on $\beta_{k,t}(n)$ comes from the probabilistic method – see the random constructions in Section 3.1 which establish the lower bounds in Theorem 3. To prove Theorem 3, we use a natural connection to hypergraph covering problems.

1.2 Covering hypergraphs

Theorem 1 has a wide variety of applications, from saturation problems [3, 19] to covering problems for graphs [11, 20], complexity of 0-1 matrices [23], geometric problems [1], counting cross-intersecting families [7], crosscuts and transversals of hypergraphs [24, 25, 26], hypergraph entropy [15, 22], and perfect hashing [8, 10]. In this section, we give an application of our main results to hypergraph covering problems. For a k -uniform hypergraph H , let $f(H)$ denote the minimum number of complete k -partite k -uniform hypergraphs whose union is H . In the case of graph covering, a simple connection to the Bollobás set pairs inequality (1) may be described as follows. Let $K_{n,n} \setminus M$ denote the complement of a perfect matching $M = \{x_i y_i : 1 \leq i \leq n\}$ in the complete bipartite graph $K_{n,n}$ with parts $X = \{x_1, x_2, \dots, x_n\}$ and $Y = \{y_1, y_2, \dots, y_n\}$. If H_1, H_2, \dots, H_m are complete bipartite graphs in a minimum covering of $K_{n,n} \setminus M$, then let $A_i = \{j : x_i \in V(H_j)\}$ and $B_i = \{j : y_i \in V(H_j)\}$. Setting $\mathcal{A} = \{A_i\}_{i \in [m]}$ and $\mathcal{B} = \{B_i\}_{i \in [m]}$, it is straightforward to check that $(\mathcal{A}, \mathcal{B})$ is a Bollobás set pair, and Theorem 1 applies to give

$$f(K_{n,n} \setminus M) = \min\left\{m : \binom{m}{\lceil m/2 \rceil} \geq n\right\}. \quad (5)$$

In a similar way, Theorem 2 applies to covering complete k -partite k -uniform hypergraphs. Let $K_{n,n,\dots,n}$ denote the complete k -partite k -uniform hypergraph with parts $X_i = \{x_{ij} : j \in [n]\}$ for $i \in [k]$. Let $H_{k,t}(n)$ denote the subhypergraph consisting of hyperedges $\{x_{1,i_1}, x_{2,i_2}, \dots, x_{k,i_k}\}$ such that at least t of the indices i_1, i_2, \dots, i_k are distinct, and set $f_{k,t}(n) = f(H_{k,t}(n))$. Then there is a one-to-one correspondence between Bollobás (k, t) -tuples of subsets of $[m]$ and coverings of $H_{k,t}(n)$ with m complete k -partite k -graphs. We let $\beta_{k,t}(m)$ be the maximum size of a Bollobás (k, t) -tuple of subsets of $[m]$, so that

$$f_{k,t}(n) = \min\{m : \beta_{k,t}(m) \geq n\}. \quad (6)$$

This correspondence together with Theorem 2 will be exploited to prove

$$f_{k,2}(n) \geq \min\left\{m : \binom{m}{\lceil m/k \rceil} \geq n\right\} \quad (7)$$

which is partly an analog of (5). More generally, we prove the following theorem:

Theorem 4. *For $k \geq 3$ and large enough n ,*

$$\frac{k}{\log_2(ke)} \leq \frac{1}{H(\frac{1}{k})} \leq \frac{f_{k,2}(n)}{\log_2 n} \leq k. \quad (8)$$

For $k \geq t \geq 3$ and large enough n ,

$$\binom{k}{t-1} \frac{(t-1)^{t-3}}{2} \leq \frac{f_{k,t}(n)}{\log_2 n} \leq \frac{(t+1)t^{t-1}}{\log_2 e} \binom{k}{t-1}. \quad (9)$$

The bounds on $\beta_{k,t}(n)$ in Theorem (3) follow immediately from this theorem and (6). Equation (9) gives the order of magnitude for each $t \geq 3$ as $k \rightarrow \infty$, but for $t = 2$, Equation (8) has a gap of order $\log_2 k$. From (7), we obtain $\beta_{k,2}(n) \leq \binom{n}{\lfloor n/k \rfloor}$. It is perhaps unsurprising that the asymptotic value of $f_{k,t}(n)/\log_2 n$ as $n \rightarrow \infty$ is not known for any $k > 2$, since a limiting value of $f(K_n^k)/\log_2 n$ is not known for any $k > 2$ – see Körner and Marston [15] and Guruswami and Riazanov [10].

1.3 Organization and notation

Given a subset $A \subset [n]$, let $A^c := [n] \setminus A$ be the complement of A in $[n]$. For positive integers $k \leq n$, let $(n)_{(k)} = (n)(n-1) \cdots (n-k+1)$ denote the falling factorial. This paper is organized as follows. In Section 2, we prove Theorem 2. In Section 2.1, we construct a Bollobás $(k, 2)$ -tuple which achieves equality in Theorem 2 and in Section 2.2, we construct a Bollobás $(k, 2)$ -tuple which gives the lower bound in Equation (3). The upper bound on $f_{k,t}(n)$ in Theorem 4 comes from a probabilistic construction in Section 3.1, and the proof of the lower bound on $f_{k,t}(n)$ is given in Section 3.3; we prove (7) in Section 3.2.

2 Proof of Theorem 2

Given a Bollobás set (k, t) -tuple $(\mathcal{A}_1, \dots, \mathcal{A}_k)$ with $\mathcal{A}_j = \{A_{j,i} : 1 \leq i \leq m\}$ and a surjection $\phi : [k] \rightarrow [t]$, consider $\mathcal{A}_\ell(\phi) : 1 \leq \ell \leq t$ where $\mathcal{A}_\ell(\phi) = \{A_{\ell,i}(\phi) : 1 \leq i \leq m\}$ and

$$A_{\ell,i}(\phi) = \bigcap_{h:\phi(h)=\ell} A_{h,i}.$$

It follows that $(\mathcal{A}_1(\phi), \dots, \mathcal{A}_t(\phi))$ is a Bollobás set (t, t) -tuple and hence it suffices to prove Theorem 2 in the case where $t = k$. In this setting, surjections $\phi : [k] \rightarrow [k]$ simply permute the k families and as such we suppress the notation of ϕ for the remainder of this section. One of the proofs of Theorem 1, given a Bollobás set pair, defines a collection of chains \mathcal{C}_i for $i \in [m]$ and shows that these chains are necessarily disjoint. Similarly, given a Bollobás set (k, k) -tuple, we will define a collection of chains \mathcal{C}_σ for every ordered collection σ of $(k-1)$ distinct indices of $[m]$ and show these chains are pairwise disjoint.

Let $(\mathcal{A}_1, \dots, \mathcal{A}_k)$ with $\mathcal{A}_j = \{A_{j,i} : 1 \leq i \leq m\}$ be a Bollobás set (k, k) -tuple, and set

$$X = \bigcup_{i=1}^m (A_{1,i} \cup A_{2,i} \cup \dots \cup A_{k,i})$$

with $|X| = n$. For $\sigma \in [m]_{(k-1)}$, define a subset \mathcal{C}_σ of permutations $\pi : X \rightarrow [n]$ by

$$\mathcal{C}_\sigma := \left\{ \pi : X \rightarrow [n] : \max_{x \in A_{1,\sigma}} \pi(x) < \min_{y \in A_{2,\sigma}} \pi(y) \leq \max_{y \in A_{2,\sigma}} \pi(y) < \cdots < \min_{z \in A_{k,\sigma}} \pi(z) \right\}.$$

Letting $U_\sigma := A_{1,\sigma} \cup \cdots \cup A_{k,\sigma}$, elementary counting methods give

$$|\mathcal{C}_\sigma| = \binom{n}{|U_\sigma|} |A_{1,\sigma}|! \cdots |A_{k,\sigma}|! (n - |U_\sigma|)! = n! \cdot \binom{|U_\sigma|}{|A_{1,\sigma}| \cdots |A_{k,\sigma}|}^{-1}. \quad (10)$$

We will now prove a lemma which states that $\{\mathcal{C}_\sigma\}_{\sigma \in [m]_{(k-1)}}$ forms a disjoint collection of a permutations. The general proof only works for $k \geq 4$, so we first consider $k = 3$.

Lemma 5. *If $\sigma_1, \sigma_2 \in [m]_{(2)}$ are distinct, then $\mathcal{C}_{\sigma_1} \cap \mathcal{C}_{\sigma_2} = \emptyset$.*

Proof. Seeking a contradiction, suppose there exists $\pi \in \mathcal{C}_{\sigma_1} \cap \mathcal{C}_{\sigma_2}$. After relabeling, it suffices to consider the following five cases.

- (1) $\sigma_1 = \{1, 3\}$ and $\sigma_2 = \{2, 4\}$ (2) $\sigma_1 = \{1, 3\}$ and $\sigma_2 = \{2, 3\}$
- (3) $\sigma_1 = \{1, 2\}$ and $\sigma_2 = \{1, 3\}$ (4) $\sigma_1 = \{1, 2\}$ and $\sigma_2 = \{2, 3\}$
- (5) $\sigma_1 = \{1, 2\}$ and $\sigma_2 = \{3, 1\}$.

In case (1), without loss of generality, $\max\{\pi(x) : x \in A_{1,1}\} \leq \max\{\pi(x) : x \in A_{1,2}\}$ and thus $\pi \in \mathcal{C}_{\sigma_2}$ yields

$$\max_{x \in A_{1,1}} \pi(x) \leq \max_{x \in A_{1,2}} \pi(x) < \min_{y \in A_{2,4} \setminus A_{1,2}} \pi(y).$$

Then as $A_{1,1} \cap A_{2,4} \cap A_{3,2} \neq \emptyset$, there exists $w \in A_{1,1} \cap A_{2,4} \cap A_{3,2}$. It follows that $w \notin A_{1,2}$ since if $w \in A_{1,2}$, then $w \in A_{1,2} \cap A_{2,4} \cap A_{3,2} \neq \emptyset$; a contradiction. But this yields a contradiction as

$$\pi(w) \leq \max_{x \in A_{1,1}} \pi(x) \leq \max_{x \in A_{1,2}} \pi(x) < \min_{y \in A_{2,4} \setminus A_{1,2}} \pi(y) \leq \pi(w).$$

In case (2), without loss of generality, $\max\{\pi(x) : x \in A_{1,1}\} \leq \max\{\pi(x) : x \in A_{1,2}\}$ and we recover a similar contradiction as case (1) by noting that there exists $w \in A_{1,1} \cap A_{2,3} \cap A_{3,2}$ with $w \notin A_{1,2}$.

In case (3) we may assume $\max\{\pi(x) : x \in A_{2,2} \setminus A_{1,1}\} \leq \max\{\pi(x) : x \in A_{2,3} \setminus A_{1,1}\}$ and $\pi \in \mathcal{C}_{1,3}$ yields $\max\{\pi(x) : x \in A_{2,3} \setminus A_{1,1}\} < \min\{\pi(x) : x \in A_{3,1} \setminus (A_{1,1} \cup A_{2,3})\}$. Thus

$$\max\{\pi(x) : x \in A_{2,2} \setminus A_{1,1}\} < \min\{\pi(x) : x \in A_{3,1} \setminus (A_{1,1} \cup A_{2,3})\}$$

and there exists $w \in A_{1,3} \cap A_{2,2} \cap A_{3,1}$ with $w \notin A_{1,1}$ and $w \notin A_{2,3}$. It follows that $\pi(w) < \pi(w)$, a contradiction.

In case (4), if $\max\{\pi(x) : x \in A_{1,1}\} \leq \max\{\pi(x) : x \in A_{1,2}\}$, then using $w \in A_{1,1} \cap A_{2,3} \cap A_{3,2}$ and noting $w \notin A_{1,2}$, we get a contradiction. Thus, we may assume otherwise and $\pi \in \mathcal{C}_{1,2}$ gives

$$\max_{x \in A_{1,2}} \pi(x) < \max_{x \in A_{1,1}} \pi(x) < \min_{z \in A_{3,1} \setminus (A_{1,1} \cup A_{2,2})} \pi(z).$$

This is a contradiction as there exists $w \in A_{1,2} \cap A_{2,3} \cap A_{3,1}$ with $w \notin A_{1,1}$ and $w \notin A_{2,2}$. In case (5), if $\max\{\pi(x) : x \in A_{1,1}\} \leq \max\{\pi(x) : x \in A_{1,3}\}$, then we may proceed as in the latter part of case (4) using $w \in A_{1,1} \cap A_{2,2} \cap A_{3,3}$ and $w \notin A_{2,1}$ and $w \notin A_{1,3}$ to get a contradiction. Otherwise, proceeding as in case (1) and noting there exists $w \in A_{1,3} \cap A_{2,2} \cap A_{3,1}$, but $w \notin A_{1,1}$ yields a contradiction. \square

A similar argument yields the analog of Lemma 5 to the case where $k \geq 4$.

Lemma 6. *Let $k \geq 4$. If $\sigma_1, \sigma_2 \in [m]_{(k-1)}$ are distinct, then $\mathcal{C}_{\sigma_1} \cap \mathcal{C}_{\sigma_2} = \emptyset$.*

Proof. Since $\sigma_1, \sigma_2 \in [m]_{(k-1)}$ are distinct, there exists minimal $h \in [k-1]$ so that $\sigma_1(h) \neq \sigma_2(h)$. Seeking a contradiction, suppose there exists a $\pi \in \mathcal{C}_{\sigma_1} \cap \mathcal{C}_{\sigma_2}$. Without loss of generality,

$$\max\{\pi(x) : x \in A_{h,\sigma_1}\} \leq \max\{\pi(x) : x \in A_{h,\sigma_2}\} < \min\{\pi(z) : z \in A_{k,\sigma_2}\}.$$

Now, consider a bijection $\tau : [k-1] \setminus \{h\} \rightarrow [k-1] \setminus \{1\}$ which has no fixed points. As in Lemma 5, we want to show that there exists a $w \in A_{h,\sigma_1} \cap A_{k,\sigma_2}$ and consider two separate cases.

First, suppose that $\sigma_1(h) \notin \sigma_2([k-1])$. As $|\{\sigma_1(h), \sigma_2(1), \dots, \sigma_2(k-1)\}| = k$, there exists

$$w \in A_{h,\sigma_1(h)} \cap A_{k,\sigma_2(1)} \cap \bigcap_{l \in [k-1] \setminus \{h\}} A_{l,\sigma_2(\tau(l))}. \quad (11)$$

Next, suppose that $\sigma_1(h) = \sigma_2(x)$ for some x . We now claim that $x \neq 1$. If $h = 1$, then this is trivial. If $h > 1$, then $\sigma_1(1) = \sigma_2(1)$, so $\sigma_1(h) \neq \sigma_2(1)$ since $\sigma_1(h) \neq \sigma_1(1)$. For τ as above, there exists $y \in [k-1] \setminus \{h\}$ so that $\tau(y) = x$. Taking γ distinct from $\{\sigma_2(1), \dots, \sigma_2(k-1)\} \setminus \{\sigma_2(x)\}$, $|\{\sigma_1(h), \gamma, \sigma_2(1), \dots, \sigma_2(k-1)\} \setminus \{\sigma_2(x)\}| = k$ and hence there exists

$$w \in A_{h,\sigma_1(h)} \cap A_{k,\sigma_2(1)} \cap A_{y,\gamma} \cap \bigcap_{l \in [k-1] \setminus \{y,h\}} A_{l,\sigma_2(\tau(l))}. \quad (12)$$

By construction, $w \in A_{h,\sigma_1(h)} \cap A_{k,\sigma_2(1)}$. Suppose there exists a $t \in [k-1] \setminus \{h\}$ so that $w \in A_{t,\sigma_2(t)}$. As τ has no fixed points, replacing the set in the k -wise intersection corresponding to A_t with $A_{t,\sigma_2(t)}$ in either (11) or (12), w is an element of this new k -wise intersection with $(k-1)$ distinct indices; a contradiction. If $w \in A_{h,\sigma_2(h)}$, then we may similarly replace $A_{h,\sigma_1(h)}$ with $A_{h,\sigma_2(h)}$ in the k -wise intersection in either (11) or (12) to get a contradiction. Thus, $w \notin A_{1,\sigma_2(1)} \cup \dots \cup A_{k-1,\sigma_2(k-1)}$ and hence $w \in A_{h,\sigma_1} \cap A_{k,\sigma_2}$ so that $\pi(w) < \pi(w)$; a contradiction. \square

Using Equation (10), Lemma 5, and Lemma 6, we are now able to prove Theorem 2 in the case where $t = k$. There are $n!$ total permutations, and Lemma 5 and Lemma 6 yield that each of which appears in at most one of the sets \mathcal{C}_σ for $\sigma \in [m]_{(k-1)}$. Hence, using $|\mathcal{C}_\sigma|$ in Equation (10),

$$\sum_{\sigma \in [m]_{(k-1)}} |\mathcal{C}_\sigma| = \sum_{\sigma \in [m]_{(k-1)}} n! \cdot \left(\frac{|A_{1,\sigma} \cup \dots \cup A_{k,\sigma}|}{|A_{1,\sigma}| \cdots |A_{k,\sigma}|} \right)^{-1} \leq n!$$

and thus the result follows by dividing through by $n!$.

2.1 Sharpness of Theorem 2

We give a simple construction establishing the sharpness of Theorem 2 for $k \geq t = 2$. Let $n \geq 4k$ and using addition modulo n , define $A_{1,i} = \{i\}^c$, $A_{j,i} = \{i - (j - 1), i + (j - 1)\}^c$ for $j \in [2, k - 1]$, and $A_{k,i} = \{i - k + 2, i - k + 3, \dots, i + k - 2\}$. Letting $\mathcal{A}_j = \{A_{j,i}\}_{i \in [n]}$ for all $j \in [k]$, we will show $(\mathcal{A}_1, \dots, \mathcal{A}_k)$ is a Bollobás $(k, 2)$ -tuple. Since $|A_{1,i}| = n - 1$ and $|A_{2,i} \cap \dots \cap A_{k,i}| = 1$, Theorem 2 with $t = 2$ and surjection $\phi : [k] \rightarrow [2]$ with $\phi(1) = 1$ and $\phi(i) = 2$ for $i \neq 1$ gives

$$1 \geq \sum_{i=1}^n \left(\frac{|A_{1,i}| + |A_{2,i} \cap \dots \cap A_{k,i}|}{|A_{1,i}|} \right)^{-1} = \sum_{i=1}^n \frac{1}{n} = 1.$$

By construction, for all $i \in [n]$, $A_{1,i} \cap A_{2,i} \cap \dots \cap A_{k,i} = \emptyset$. It thus suffices to show these are the only empty k -wise intersections. To this end, for $\mathbf{i} = (i_1, \dots, i_{k-1})$, define

$$A(\mathbf{i}) := A_{1,i_1} \cap \dots \cap A_{k-1,i_{k-1}}.$$

Lemma 7. *Let $\mathbf{i} = (i_1, \dots, i_{k-1})$. If $A(\mathbf{i})^c = A_{k,i_k}$, then $i_1 = \dots = i_k$.*

Proof. We proceed by induction on k where the result is trivial when $k = 2$. In the case where $k > 2$, $i_{k-1} - k + 2 = i_k + x$ for some x such that $-(k - 2) \leq x \leq (k - 2)$ and thus $i_{k-1} + (k - 2) = i_{k-1} - (k - 2) + (2k - 4) = i_k + x + (2k - 4)$.

Next, there is a y such that $-(k - 2) \leq y \leq (k - 2)$ with $i_{k-1} + (k - 2) = i_k + y$, and since $n \geq 4k$, $x + 2k - 4 = y$ with equality over \mathbb{Z} and moreover $i_{k-1} + (k - 2) = i_k + (k - 2)$ over \mathbb{Z} and hence $i_k = i_{k-1}$. Removing these elements from each set, the result then follows by induction. \square

If $A_{1,i_1} \cap \dots \cap A_{k,i_k} = \emptyset$, then as $A(\mathbf{i}) = A_{1,i_1} \cap A_{2,i_2} \cap \dots \cap A_{k-1,i_{k-1}}$,

$$\emptyset = A_{1,i_1} \cap A_{2,i_2} \cap \dots \cap A_{k-1,i_{k-1}} \cap A_{k,i_k} = A(\mathbf{i}) \cap A_{k,i_k}.$$

The result follows by noting $|A(\mathbf{i})| \geq n - (2k - 3)$, $|A_{k,i_k}| = 2k - 3$, and using Lemma 7.

2.2 An Explicit Construction

Let $k \geq 3$. An explicit construction of a Bollobás $(k, 2)$ -tuple $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$ where $|\mathcal{A}_i| = 2^n$ and each \mathcal{A}_i consists of subsets of X for $|X| = kn$ may be described as follows. Let $I_j := \{x_{j,1}, x_{j,2}, \dots, x_{j,k}\}$ and consider $X = I_1 \sqcup \dots \sqcup I_n$. Now, for each $f : [n] \rightarrow [2]$ and $j \in [k]$, define

$$A_{j,f} := \{x_{1,f(1)+j-1}, \dots, x_{n,f(n)+j-1}\}^c$$

where we work modulo k within the subscripts of I_j . It is straightforward to check that $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$ is a Bollobás $(k, 2)$ -tuple. This establishes the lower bound on $\beta_{k,2}(n)$ in Equation (3) and hence the upper bound on $f_{k,2}(n)$ in Equation (8).

3 Proof of Theorem 4

3.1 Upper bound on $f_{k,t}(n)$

We wish to find a covering of $H_{k,t}(n)$ with complete k -partite k -graphs and assume the parts of $H_{k,t}(n)$ are X_1, X_2, \dots, X_k . For each subset T of $[k]$ of size t , consider the uniformly random coloring $\chi_T : [n] \rightarrow T$. Given such a χ_T , let $Y_i \subset X_i$ be the vertices of color i for $i \in T$; that is $Y_i := \{x_{ij} : \chi(j) = i\}$ and $Y_i = X_i$ for $i \notin T$. Denote by $H(T, \chi)$ the (random) complete k -partite hypergraph with parts Y_1, Y_2, \dots, Y_k , and note that $H(T, \chi) \subset H_{k,t}(n)$. We place each $H(T, \chi)$ a total of N times independently and randomly where

$$N = \left\lfloor \frac{(t+1)t^t \log_2 n}{(k-t+1) \log_2 e} \right\rfloor$$

and produce $\binom{k}{t}N$ random subgraphs $H(T, \chi)$. For a set partition π of $[k]$, let $|\pi|$ denote the number of parts in the partition and index the parts by $[\pi]$. Given a set partition $\pi = (P_1, P_2, \dots, P_s)$, let

$$f(\pi, t) = \sum_{T \in [s]^{(t)}} \prod_{i \in T} |P_i|.$$

If U is the number of edges of $H_{k,t}(n)$ not in any of these subgraphs, then

$$\mathbb{E}(U) \leq \sum_{|\pi| \geq t} n^{|\pi|} (1 - t^{-t})^{Nf(\pi, t)} = \sum_{t \leq s \leq k} n^s \sum_{|\pi|=s} (1 - t^{-t})^{Nf(\pi, t)}. \quad (13)$$

For sufficiently large n , we claim that $\mathbb{E}(U) < 1$, which implies there exists a covering of $H_{k,t}(n)$ with at most $\binom{k}{t}N$ complete k -partite k -graphs, as required. The following technical lemma states that f is a decreasing function in the set partition lattice, and that $f(\pi, t)$ increases when we merge all but one element of a smaller part of π with a larger part of π :

Lemma 8. *Let $k \geq s \geq t \geq 2$, and let $\pi = (P_1, P_2, \dots, P_s)$ be a partition of $[k]$.*

- (i) *If π' is a refinement of π with $|\pi'| = s+1$, then $f(\pi, t) \leq f(\pi', t)$.*
- (ii) *If $|P_1| \geq |P_2| \geq 2$ and $a \in P_2$, and π' is the partition $(P'_1, P'_2, \dots, P'_s)$ of $[k]$ with $P'_1 = P_1 \cup P_2 \setminus \{a\}$ and $P'_2 = \{a\}$ and with $P'_i = P_i$ for $3 \leq i \leq s$, then $f(\pi', t) \leq f(\pi, t)$.*

The proof of Lemma 8 part (i) is in Appendix A and the proof of (ii) is similar to the proof of (i). By Lemma 8, a set partition of $[k]$ into s parts which minimizes $f(\pi, t)$ consists of one part of size $k - s + 1$ and $s - 1$ singleton parts and hence

$$\min\{f(\pi, t) : |\pi| = s\} = (k - s + 1) \binom{s-1}{t-1} + \binom{s-1}{t}. \quad (14)$$

In what follows, we denote a set partition of $[k]$ into s parts which minimizes $f(\pi, t)$ by π_s .

For n large enough, and all s where $t \leq s \leq k$, we will show

$$\frac{\sum_{|\pi|=t}(1-t^{-t})^{Nf(\pi,t)}}{\sum_{|\pi|=s}(1-t^{-t})^{Nf(\pi,t)}} \geq n^{s-t}.$$

Replacing the numerator with its largest term and each term in denominator with its largest term,

$$\frac{\sum_{|\pi|=t}(1-t^{-t})^{Nf(\pi,t)}}{\sum_{|\pi|=s}(1-t^{-t})^{Nf(\pi,t)}} \geq \frac{(1-t^{-t})^{Nf(\pi_t,t)}}{S(k,s)(1-t^{-t})^{Nf(\pi_s,t)}} = \frac{1}{S(k,s)}(1-t^{-t})^{N(f(\pi_t,t)-f(\pi_s,t))}$$

where $S(k,s)$ is the Stirling number of the second kind. Taking $n \geq S(k,s)$, we will show in Appendix B that

$$\frac{1}{S(k,s)}(1-t^{-t})^{N(f(\pi_t,t)-f(\pi_s,t))} \geq n^{s-t}. \quad (15)$$

Therefore, the index $s = t$ maximizes the right hand side of Equation (13), and hence

$$\mathbb{E}[U] \leq (k-t+1)(n^t) \sum_{|\pi|=t} (1-t^{-t})^{Nf(\pi,t)} < (k-t+1)n^t S(k,t)(1-t^{-t})^{N(k-t+1)} < 1$$

for our choice of N provided $n \geq kS(k,t)$. Thus,

$$f_{k,t}(n) \leq \binom{k}{t} \frac{(t+1)t^t \log_2 n}{(k-t+1) \log_2 e} = \frac{(t+1)t^{t-1}}{\log_2 e} \binom{k}{t-1} \log_2 n.$$

3.2 Lower bound on $f_{k,2}(n)$

In this section, we show

$$f_{k,2}(n) \geq \min\{m : \binom{m}{\lceil m/k \rceil} \geq n\}. \quad (16)$$

Let $\{H_1, H_2, \dots, H_m\}$ be a covering of $H_{k,2}(n)$ with $m = f_{k,2}(n)$ complete k -partite k -graphs. We recall $H_{k,2}(n) = K_{n,n,\dots,n} \setminus M$, where M is a perfect matching of $K_{n,n,\dots,n}$. For $i \in [k]$ and $j \in [n]$, define $A_{i,j} = \{H_r : x_{ij} \in V(H_r)\}$ and $\mathcal{A}_i = \{A_{i,j} : 1 \leq j \leq n\}$. As in (6), $(\mathcal{A}_1, \mathcal{A}_2, \dots, \mathcal{A}_k)$ is a Bollobás $(k, 2)$ -tuple of size n . For convenience, for each $i \in [k]$, let $\phi_i : [k] \rightarrow [2]$ be so that $\phi_i^{-1}(1) = \{i\}$. Taking the sum of inequality from Theorem 2 with $t = 2$ over all $i \in [k]$,

$$\sum_{i=1}^k \sum_{j=1}^n \left(\frac{|A_{1,j}(\phi_i) \cup A_{2,j}(\phi_i)|}{|A_{1,j}(\phi_i)|} \right)^{-1} \leq k. \quad (17)$$

We use this inequality to give a lower bound on $f_{k,2}(n) = m$. First we observe

$$\sum_{r=1}^m |V(H_r)| = \sum_{j=1}^n \sum_{i=1}^k |A_{i,j}| = \sum_{j=1}^n \sum_{i=1}^k |A_{1,j}(\phi_i)|. \quad (18)$$

Let ∂H denote the set of $(k-1)$ -tuples of vertices contained in some edge of a hypergraph H . Then

$$\sum_{r=1}^m |\partial H_r \cap \partial M| = \sum_{j=1}^n \sum_{i=1}^k |A_{2,j}(\phi_i)|. \quad (19)$$

Putting the above identities together,

$$\sum_{r=1}^m |V(H_r)| + \sum_{r=1}^m |\partial H_r \cap \partial M| = \sum_{j=1}^n \sum_{i=1}^k (|A_{1,j}(\phi_i)| + |A_{2,j}(\phi_i)|). \quad (20)$$

We note $|\partial H_r \cap \partial M| \leq |V(H_r)|/(k-1)$, and therefore

$$\sum_{r=1}^m |\partial H_r \cap \partial M| \leq \frac{1}{k-1} \sum_{r=1}^m |V(H_r)|. \quad (21)$$

It follows that

$$\sum_{j=1}^n \sum_{i=1}^k (|A_{1,j}(\phi_i)| + |A_{2,j}(\phi_i)|) \leq \frac{k}{k-1} \sum_{r=1}^m |V(H_r)|. \quad (22)$$

Subject to the linear inequalities (18) and (22), the left side of (17) is minimized when $kn|A_{1,j}(\phi_i)| = \sum_{r=1}^m |V(H_r)|$ and $kn(|A_{1,j}(\phi_i)| + |A_{2,j}(\phi_i)|) = (k-1)|A_{1,j}(\phi_i)|$. Since $|V(H_r)| \leq (k-1)n$ for all $r \in [m]$, (17) implies $\binom{m}{\lceil m/k \rceil} \geq n$, which gives (16). \square

3.3 Lower bound on $f_{k,k}(n)$

Let $\mathcal{H} = \{H_1, H_2, \dots, H_m\}$ be a minimal covering of $H_{k,k}(n)$ with complete k -partite k -graphs, so $m = f(H_{k,k}(n))$. Given a k -partite k -graph H , consider its 2-shadow $\partial_2(H) = \{R \subset V(H) : |R| = k-2, R \subset e \text{ for some } e \in H\}$. Let $\partial_2(\mathcal{H}) = \bigcup_{i=1}^m \partial_2(H_i)$.

Given $R \in \partial_2(\mathcal{H})$ and $H_i \in \mathcal{H}$, let $H_i(R) := \{e \in \binom{V(H_i)}{2} : e \cup R \in H_i\}$ be the possibly empty link graph of the edge R in the hypergraph H_i and let $V(H_i(R))$ be the set of vertices in the link graph. Observe that double counting yields

$$\sum_{R \in \partial_2(\mathcal{H})} \left(\sum_{i=1}^m |V(H_i(R))| \right) = \sum_{i=1}^m \left(\sum_{R \in \partial_2(H_i)} |V(H_i(R))| \right). \quad (23)$$

An optimization argument yields $|\partial_2(H_i)|$ is maximized when the parts of H_i are of equal or nearly equal maximal size. Since $|V(H_i(R))| \leq 2(n-k+2)$, the right hand side of Equation (23) is bounded above by

$$\sum_{i=1}^m \left(\sum_{R \in \partial_2(H_i)} |V(H_i(R))| \right) \leq m \cdot \binom{k}{2} \cdot \left(\frac{n}{k} \right)^{k-2} \cdot 2(n-k+2). \quad (24)$$

For a lower bound on the left hand side of Equation (23), fix $R \in \partial_2(\mathcal{H})$ and without loss of generality suppose that $R = \{x_{1,1}, \dots, x_{k-2,k-2}\}$. Let $Y = [k-1, n]$. Let $K_{Y,Y}$ be the complete bipartite graph with two distinct copies of Y and $\mathcal{M} = \{(x_{k-1,i}, x_{k,i} : i \in Y\}$ be a perfect matching in $K_{Y,Y}$. Then, $\{H_1(R), \dots, H_m(R)\}$ forms a biclique cover of $K_{Y,Y} \setminus \mathcal{M}$. Applying the convexity result of Tarjan [23, Lemma 5],

$$\sum_{i=1}^m |V(H_i(R))| \geq (n-k+2) \log_2(n-k+2).$$

Noting that $|\partial_2(\mathcal{H})| = \binom{k}{2} \binom{n}{k-2}$, the left hand side of Equation (23) is bounded below by

$$\sum_{R \in \partial_2(\mathcal{H})} \left(\sum_{i=1}^m |V(H_i(R))| \right) \geq \binom{k}{2} \binom{n}{k-2} (n-k+2) \log_2(n-k+2). \quad (25)$$

Comparing the bounds from Equation (24) and Equation (25),

$$m \geq \frac{\binom{n}{k-2} \log_2(n-k+2)}{2 \left(\frac{n}{k}\right)^{k-2}} \geq \frac{k^{k-2}}{2} \log_2 n$$

provided that n is large enough.

For $t \geq 3$ and $t < k$, the lower bound on $f_{k,t}(n)$ in Theorem 4 is obtained from the lower bounds on $f_{t-1,t-1}(n-1)$ as follows: Let $\mathcal{H} = \{H_1, H_2, \dots, H_m\}$ be a minimal covering of $H_{k,t}(n)$ with complete k -partite k -graphs, so $m = f(H_{k,t}(n))$. Given $T \in \binom{[k]}{k-t+1}$, define $H_T \subset H_{k,t}(n)$ by

$$H_T := \{\{x_{1,i_1}, \dots, x_{k,i_k}\} \in H_{k,t}(n) : i_j = 1 \ \forall j \in T\}.$$

It follows that at least $f_{t-1,t-1}(n-1)$ of the complete k -partite k -graphs in \mathcal{H} are needed to cover H_T . Moreover, for distinct $T, T' \in \binom{[k]}{k-t+1}$, the corresponding complete k -partite k -graphs from \mathcal{H} are necessarily pairwise disjoint and hence

$$f_{k,t}(n) \geq \binom{k}{k-t+1} f_{t-1,t-1}(n-1) \geq \binom{k}{t-1} \frac{(t-1)^{t-3}}{2} \log_2 n$$

provided that n is large enough.

4 Concluding remarks

- Our main theorem, Theorem 2 is tight for $t = 2$ and $k \geq 2$, as shown in Section 2.1. It would be interesting to generalize this example to $2 < t \leq k$ to determine whether Theorem 2 is tight in general. The first open case is $t = k = 3$.
- A particular case of the Bollobás set pairs inequality occurs when every set in \mathcal{A} has size a and every set in \mathcal{B} has size b , and one obtains the tight bound $|\mathcal{A}| \leq \binom{a+b}{b}$. The

generalization to Bollobás (k, t) -tuples for $k \geq 3$ is equally interesting but wide open, as are potential generalizations to vector spaces – see Lovász [17, 18].

- Orlin [20] proved that the clique cover number $cc(K_n \setminus M)$ of a complete graph K_n minus a perfect matching M is precisely $\min\{m : 2^{\binom{m-1}{\lfloor m/2 \rfloor}} \geq n\}$. Theorem 4 yields lower bounds on the clique cover number of the complement of a perfect matching M in the complete k -uniform hypergraph K_n^k :

Corollary 9. *Let $K_n^k \setminus M$ be the complement of a perfect matching in K_n^k . Then*

$$cc(K_n^k \setminus M) \geq \frac{\log_2 \frac{n}{k}}{H(\frac{1}{k})} \geq \frac{k \log_2 \frac{n}{k}}{\log_2(ke)}.$$

- It would be interesting to prove an analog of Equation (16) for $t \geq 3$. That is,

$$f_{k,t}(n) \geq \min\{m : \binom{m}{\alpha_1, \dots, \alpha_t} \geq n_{(t-1)}\}$$

for some optimal $\alpha_1, \dots, \alpha_t$. The difficulty here lies in determining effective bounds on $|A_{i,\sigma}(\phi)|$.

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A Proof of Lemma 8(i)

Let $k \geq s \geq t \geq 2$, and let $\pi = (P_1, P_2, \dots, P_s)$ be a partition of $[k]$. In this section, we will show that if π' is a refinement of π with $|\pi'| = s + 1$, then $f(\pi, t) \leq f(\pi', t)$.

Proof. Let $\pi = P_1|P_2|\dots|P_s$ and without loss of generality, $\pi' = P_x|P_y|P_2|\dots|P_s$. Setting $\mathcal{T}(\bar{1}) = \{T \in [s]^{(t)} : 1 \notin T\}$ and $\mathcal{T}'(\bar{x}, \bar{y}) = \{T \in \{x, y, 2, \dots, s\}^{(t)} : x, y \notin T\}$, it follows that

$$\sum_{T \in \mathcal{T}(\bar{1})} \prod_{i \in T} |P_i| = \sum_{T \in \mathcal{T}'(\bar{x}, \bar{y})} \prod_{i \in T} |P_i|.$$

Now, letting $\mathcal{T}(1) = \{T \in [s]^{(t)} : 1 \in T\}$ and $\mathcal{T}'(x, \bar{y}) = \{T \in \{x, y, 2, \dots, s\}^{(t)} : x \in T, y \notin T\}$ and $\mathcal{T}'(\bar{x}, y) = \{T \in \{x, y, 2, \dots, s\}^{(t)} : x \notin T, y \in T\}$, we see that

$$\sum_{T \in \mathcal{T}(1)} \prod_{i \in T} |P_i| = \sum_{T \in \mathcal{T}'(\bar{x}, y)} \prod_{i \in T} |P_i| + \sum_{T \in \mathcal{T}'(x, \bar{y})} \prod_{i \in T} |P_i|$$

since $|P_1| = |P_x| + |P_y|$. Thus letting $\mathcal{T}'(x, y) = \{T \in \{x, y, 2, \dots, s\}^{(t)} : x \in T, y \in T\}$,

$$f(\pi', t) - f(\pi, t) = \sum_{T \in \mathcal{T}'(x, y)} \prod_{i \in T} |P'_i|$$

and in particular $f(\pi, t) \leq f(\pi', t)$. □

B Proof of Equation (15)

Let $S(k, s)$ be the Stirling number of the second kind and $f(\pi)$ be as in Section 3. In this section we will show

$$\frac{1}{S(k, s)} (1 - t^{-t})^{N(f(\pi_t, t) - f(\pi_s, t))} \geq n^{s-t}.$$

Proof. First, we recall that

$$N = \left\lfloor \frac{(t+1)t^t \log_2 n}{(k-t+1) \log_2 e} \right\rfloor \quad \text{and} \quad f(\pi_s, t) = (k-s+1) \binom{s-1}{t-1} + \binom{s-1}{t}.$$

As a result, when $t \leq s < k$, a calculation yields that

$$f(\pi_{s+1}, t) - f(\pi_s, t) = (k-s) \binom{s-1}{t-2}. \quad (26)$$

Letting $n \geq S(k, t)$, after taking $\log_2(\cdot)$ on both sides of (15), it suffices to prove that

$$N \cdot \frac{f(\pi_s, t) - f(\pi_t, t)}{t^t} \left(-t^t \log_2(1 - t^{-t}) \right) \geq (s-t+1) \log_2(n). \quad (27)$$

Using the fact that $(1 - t^{-t})^{t^t} \leq e^{-1}$ and our choice of N , it suffices to show that

$$f(\pi_s, t) - f(\pi_t, t) \geq \frac{(s-t+1)(k-t+1)}{t+1}. \quad (28)$$

The inequality in (28) holds for all $k \geq s > t \geq 3$ by using (26). □