

COUNTEREXAMPLES TO A CONJECTURE OF HARRIS ON HALL RATIO*

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Abstract. The Hall ratio of a graph G is the maximum value of $v(H)/\alpha(H)$ taken over all non-null subgraphs $H \subseteq G$. For any graph, the Hall ratio is a lower-bound on its fractional chromatic number. In this note, we present various constructions of graphs whose fractional chromatic number grows much faster than their Hall ratio. This refutes a conjecture of Harris.

Key words. fractional colorings, independent sets, Hall ratio, random graphs

AMS subject classifications. 05C15, 05C69, 05C72

DOI. 10.1137/18M1229420

1. Introduction. A graph G is k -colorable if its vertices can be colored with k colors so that adjacent vertices receive different colors. The minimum integer k such that G is k -colorable is called *the chromatic number* of G , and it is denoted by $\chi(G)$.

Various refinements and relaxations of the chromatic number have been considered in the literature. One of the classical and most studied ones is the *fractional chromatic number*, which we denote by $\chi_f(G)$; see section 2.1 for its definition.

A basic averaging argument reveals that $\chi_f(G) \geq v(G)/\alpha(G)$, where $v(G)$ and $\alpha(G)$ are the number of vertices and the size of a largest independent set in G , respectively. Moreover, since $\chi_f(G) \geq \chi_f(H)$ for a subgraph $H \subseteq G$, it holds that

$$\chi_f(G) \geq \frac{v(H)}{\alpha(H)} \quad \text{for every non-null } H \subseteq G.$$

We define $\rho(G)$ —the *Hall ratio* of a graph G —to be the best lower-bound obtained

*Received by the editors November 28, 2018; accepted for publication (in revised form) April 24, 2020; published electronically July 18, 2022.

<https://doi.org/10.1137/18M1229420>

Funding: This project has received funding from the European Union’s Horizon 2020 Research and Innovation Programme under the Marie Skłodowska-Curie grant agreement 800607. This work was partially supported by NSF-DMS grant 1604458, Collaborative Research: Rocky Mountain Great Plains Graduate Research Workshops in Combinatorics, and by NSA grant H98230-18-1-0017, the 2018 and 2019 Rocky Mountain Great Plains Graduate Research Workshop in Combinatorics. Research of the second author is partially supported by NSF grant DMS-1600390. Research of the third author is partially supported by a grant from the Simons Foundation, 353292. Research of the fourth author is supported by NSERC grant 418520. Research of the fifth author is partially supported by NSF grant DMS-1600483. Research of the sixth author was partially supported by a CRM-ISM fellowship.

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in this way, i.e.,

$$\rho(G) := \max_{\emptyset \neq H \subseteq G} \frac{v(H)}{\alpha(H)}.$$

How tight is $\rho(G)$ as a lower-bound for $\chi_f(G)$? In 2009, Johnson [11] suggested that there are graphs G where the value of $\chi_f(G)/\rho(G)$ is unbounded. In earlier versions of [7] (see [8, Conjecture 6.2]), Harris explicitly conjectured the opposite.

CONJECTURE 1. *There exists C such that $\chi_f(G) \leq C \cdot \rho(G)$ for every graph G .*

In 2016, Barnett [2] constructed graphs showing that if such a constant C exists, then $C \geq 343/282 \sim 1.216$, improving an earlier bound 1.2 [3]. Our first result refutes Conjecture 1.

THEOREM 2. *There exists P_0 such that for every $P \geq P_0$, there is a graph G with $\rho(G) \leq P$ and $\chi_f(G) > P^2/33$.*

The proof of Theorem 2 is very short and simple, modulo some standard results about random graphs. The following two theorems strengthen Theorem 2 at the expense of somewhat more technical proofs.

THEOREM 3. *There exists P_0 such that for every $P \geq P_0$ there is a K_5 -free graph G with $\rho(G) \leq P$ and $\chi_f(G) > P^2/82$.*

THEOREM 4. *There exists P_0 such that for all $P \geq P_0$ there is a graph G with $\rho(G) \leq P$ and $\chi_f(G) \geq e^{\ln^2(P)/5}$.*

This note is organized as follows. In section 2, we recall definitions and properties of the fractional chromatic number, and Erdős–Rényi random graphs. Proofs of our results are in section 3. We conclude the note in section 4 with related open problems.

2. Definitions and preliminaries. The join of two graphs G_1 and G_2 , which we denote by $G_1 \wedge G_2$, is obtained by taking vertex-disjoint copies of G_1 and G_2 and adding all the edges between $V(G_1)$ and $V(G_2)$. More generally, for graphs $G_1, G_2, \dots, G_\ell$, we write $\bigwedge_{i=1}^\ell G_i$ to denote $(\bigwedge_{i=1}^{\ell-1} G_i) \wedge G_\ell$.

For a graph H on the vertex-set $\{1, \dots, \ell\}$ and a collection of ℓ vertex-disjoint graphs $G_1, \dots, G_\ell$, we define $H\{G_1, \dots, G_\ell\}$ to be the graph obtained by taking a union $G_1, \dots, G_\ell$, and, for every edge $ij \in E(H)$, adding all the edges between $V(G_i)$ and $V(G_j)$. Note that if $G_1 \cong \dots \cong G_\ell$, then $H\{G_1, \dots, G_\ell\}$ corresponds to the composition (also known as the lexicographic product) of G and H . Also, observe that

$$K_\ell\{G_1, \dots, G_\ell\} = \bigwedge_{i=1}^\ell G_i.$$

2.1. Fractional chromatic number. We present a definition of the fractional chromatic number based on a linear programming relaxation of an integer program computing the ordinary chromatic number. For a graph G , let $\mathcal{I}(G)$ be the set of all independent sets. Let FRACC be the following linear program:

$$\text{FRACC} \left\{ \begin{array}{ll} \text{Minimize} & \sum_{I \in \mathcal{I}(G)} x_I \\ \text{subject to} & \sum_{\substack{I \in \mathcal{I}(G) \\ v \in I}} x_I \geq 1 \quad \text{for } v \in V(G); \\ & x_I \geq 0 \quad \text{for } I \in \mathcal{I}(G). \end{array} \right.$$

Furthermore, let FRACD be the following program, which is the dual of FRACC:

$$\text{FRACD} \begin{cases} \text{Maximize} & \sum_{v \in V(G)} y(v) \\ \text{subject to} & \sum_{v \in I} y(v) \leq 1 \quad \text{for } I \in \mathcal{I}(G); \\ & y(v) \geq 0 \quad \text{for } v \in V(G). \end{cases}$$

Since these two linear programs are dual of each other, the LP-duality theorem ensures that they have the same value, which we denote by $\chi_f(G)$.

Let us now mention a different way to introduce the fractional chromatic number. As we have already mentioned, $\alpha(G) \geq v(G)/\chi_f(G)$. Moreover, the lower-bound stays valid even in the setting where the vertices have weights, and we measure the size of an independent set by the proportion of the weight it occupies rather than its cardinality.

More precisely, let $G = (V, E)$ be a graph and $w : V \rightarrow \mathbb{R}_+$ a weight function. Let $\alpha(G, w)$ be the maximum sum of the weights of the vertices that form an independent set, i.e.,

$$\alpha(G, w) := \max_{I \in \mathcal{I}} \sum_{v \in I} w(v).$$

If we rescale an optimal solution of FRACC by a factor $1/\chi_f(G)$ and interpret it as a probability distribution on $\mathcal{I}$, the linearity of expectation yields that

$$\alpha(G, w) \geq \mathbb{E}_I \left[\sum_{v \in I} w(v) \right] = \sum_{v \in V} w(v) \cdot \sum_{\substack{I \in \mathcal{I}(G) \\ v \in I}} \frac{x_I}{\chi_f(G)} \geq \frac{\sum_{v \in V} w(v)}{\chi_f(G)}.$$

On the other hand, any optimal solution of FRACD yields a weight function w_0 for which the bound is tight, i.e., $\alpha(G, w_0) = \sum_{v \in V} w_0(v)/\chi_f(G)$. Therefore,

$$\chi_f(G) = \sup_{w: V \rightarrow [0,1]} \frac{\sum_{v \in V} w(v)}{\alpha(G, w)}.$$

Note that the Hall ratio can be viewed as an integral version of the above, since

$$\rho(G) = \max_{w: V \rightarrow \{0,1\}} \frac{\sum_{v \in V} w(v)}{\alpha(G, w)}.$$

For other possible definitions of the fractional chromatic number, see [15].

We finish this section with a straightforward generalization of the fact that the fractional chromatic number of the composition of two graphs is equal to the product of their fractional chromatic numbers.

PROPOSITION 5. *Let H be a graph with the vertex-set $\{1, \dots, \ell\}$, and let $G_1, \dots, G_\ell$ be graphs. It holds that $\chi_f(H) \cdot \min_{i \in [\ell]} \chi_f(G_i) \leq \chi_f(H\{G_1, \dots, G_\ell\})$. In particular, $\chi_f(\bigwedge_{i=1}^\ell G_i) \geq \ell \cdot \min_{i \in [\ell]} \chi_f(G_i)$.*

Proof. Without loss of generality, we may assume that $V(G_i) = \{1, \dots, v(G_i)\}$. Let $w_1^H, \dots, w_\ell^H$ be any optimal solution of the dual program FRACD for H , and, for every $i \in [\ell]$, let $w_1^i, \dots, w_{v(G_i)}^i$ be any optimal solution of FRACD for G_i .

Let $G := H\{G_1, \dots, G_\ell\}$. For a vertex $(i, j) \in V(G)$, where $i \in [\ell]$ and $j \in [v(G_i)]$, we set $y_{i,j} := w_i^H \cdot w_j^i$. It holds that

$$\sum_{(i,j) \in V(G)} y_{i,j} = \sum_{i \in [\ell]} w_i^H \cdot \sum_{j \in V(G_i)} w_j^i = \sum_{i \in [\ell]} w_i^H \cdot \chi_f(G_i) \geq \chi_f(H) \cdot \min_{i \in [\ell]} \chi_f(G_i).$$

We claim that $(y_{i,j})$, where $(i,j) \in V(G)$, is a feasible solution of FRACD for G .

Indeed, fix any $I \in \mathcal{I}(G)$. For $i \in [\ell]$, let $I_i := \{j \in [v(G_i)] : (i,j) \in I\}$. Since $I_i \in \mathcal{I}(G_i)$, it holds that

$$\sum_{j \in I_i} w_i^H \cdot w_j^i = w_i^H \cdot \sum_{j \in I_i} w_j^i \leq w_i^H.$$

On the other hand, the set $I_H := \{i \in [\ell] : \exists (i,j) \in I\}$ is independent in H . Therefore,

$$\sum_{(i,j) \in I} y_{i,j} = \sum_{i \in I_H} \sum_{j \in I_i} y_{i,j} = \sum_{i \in I_H} w_i^H \cdot \sum_{j \in I_i} w_j^i \leq \sum_{i \in I_H} w_i^H \leq 1. \quad \square$$

We note that a similar composition of optimal solutions of FRACC yields $\chi_f(H\{G_1, \dots, G_\ell\}) \leq \chi_f(H) \cdot \max_{i \in [\ell]} \chi_f(G_i)$, but we will never need this bound. However, we will use the following analogue of this bound for proper colorings.

PROPOSITION 6. *Let H be a graph with the vertex-set $\{1, \dots, \ell\}$, and let $G_1, \dots, G_\ell$ be graphs. It holds that $\chi(H) \cdot \max_{i \in [\ell]} \chi(G_i) \geq \chi(H\{G_1, \dots, G_\ell\})$.*

Proof. Let $k := \max_{i \in [\ell]} \chi(G_i)$, and let d be a proper $\chi(H)$ -coloring of H . Next, for every $i \in [\ell]$, let $c_i : V(G_i) \rightarrow [k]$ be a proper k -coloring of G_i . It is straightforward to verify that assigning each vertex $v \in V_i$ a color $(d(i), c_i(v))$ yields a proper coloring of $\chi(H\{G_1, \dots, G_\ell\})$ using $\chi(H) \cdot k$ colors. $\square$

Finally, the following observation is going to be useful in the next section.

OBSERVATION 7. *If every $H \subseteq G$ has at most $|V(H)|$ edges, then $\chi(G) \leq 3$.*

Proof. Without loss of generality, we may assume G is connected. Since $|E(G)| \leq |V(G)|$, the graph G contains at most one cycle, and hence it is 3-colorable. $\square$

2.2. Sparse Erdős–Rényi random graphs. Let $G_{n,p}$ be a random graph on $\{1, 2, \dots, n\}$ where each pair of vertices forms an edge independently with probability p . We now recall some well-known properties of $G_{n, \frac{D}{n}}$ we are going to use.

PROPOSITION 8. *There exists C_0 such that for every $C \geq C_0$ the following is true: There exists $n_0 = n_0(C) \in \mathbb{N}$ such that for every $n \geq n_0$ there is an n -vertex triangle-free graph $G = G^1(n, C)$ with the following properties:*

- (A) $1.001 \cdot C > \chi(G) \geq \chi_f(G) \geq \frac{n}{\alpha(G)} > C$, and
- (B) for all $k \leq \sqrt{\ln n}$, every k -vertex subgraph of G is 3-colorable.

Proof. Suppose that C and n are sufficiently large, and let $D > 1$ be such that $C = \frac{D}{2 \cdot \ln D}$. By [6] and [13], a random graph $G_{n, \frac{D}{n}}$ satisfies with high probability $\alpha(G_{n, \frac{D}{n}}) < n/C$ and $\chi(G_{n, \frac{D}{n}}) < 1.001 \cdot C$, respectively.

Next, the expected number of subgraphs H in $G_{n, \frac{D}{n}}$ with $v(H) \leq \sqrt{\ln n}$ and more than $v(H)$ edges is at most

$$\sum_{k=3}^{\sqrt{\ln n}} 2^{k^2} \cdot n^k \cdot \left(\frac{D}{n}\right)^{k+1} \leq \sqrt{\ln n} \cdot \frac{D^{\sqrt{\ln n}+1}}{n^{1-\ln 2}} = O(n^{-0.3}).$$

By Markov's inequality, $G_{n, \frac{D}{n}}$ has no such H with high probability; hence the property (B) follows from Observation 7.

Finally, Schürger [16] showed that the number of triangles in $G_{n, \frac{D}{n}}$ converges to the Poisson distribution with mean $\Theta(D^3)$. Therefore, $G_{n, \frac{D}{n}}$ is triangle-free with probability $e^{-\Theta(D^3)} > 0$. Note that a similar estimate can also be deduced using the FKG inequality [9, Theorem 2.12]. $\square$

3. Counterexamples to Conjecture 1. We start with a simple construction of a sequence of graphs for which $\chi_f(G) \gg \rho(G)$. Each graph G is the join of the graphs $G^1(n_i, C)$ of very different orders.

Proof of Theorem 2. Let C_0 be the constant from Proposition 8, and let $P_0 := 8C_0$.

Given $P \geq P_0$, let $\ell := \lfloor P/4 \rfloor$, $C := P/8$, and $n_1 := n_0(C)$ from Proposition 8. For all $j \in [\ell - 1]$, let $n_{j+1} := \lceil e^{2 \cdot n_j^2} \rceil$, and, for all $i \in [\ell]$, let $G_i := G^1(n_i, C)$. We set $G := \bigwedge_{i=1}^{\ell} G_i$.

By Proposition 5, $\chi_f(G) > \ell \cdot C > P^2/33$. It only remains to prove that $\rho(G) \leq P$, i.e., that $\alpha(G[X]) \geq v(G[X])/P$ for every $X \subseteq V(G)$.

Fix $X \subseteq V(G)$, and let $X_i := V(G_i) \cap X$ for $i \in [\ell]$. We split the indices into two categories, *small* and *big*, based on $|X_i|$ with respect to $v(G_i) = n_i$. Specifically, let

$$\mathcal{S} := \left\{ i \in [\ell] : |X_i| < \sqrt{\ln n_i} \right\} \quad \text{and} \quad \mathcal{B} := [\ell] \setminus \mathcal{S}.$$

Next, let H_S and H_B be the subgraphs of G induced by $\bigcup_{i \in \mathcal{S}} X_i$ and $\bigcup_{i \in \mathcal{B}} X_i$, respectively, and let v_s and v_b be their respective orders. In both of these subgraphs, we can find quite large independent sets.

CLAIM 9. H_S has an independent set of size at least $4v_s/3P$.

Proof. Fix $i \in \mathcal{S}$ such that $|X_i|$ is maximized. Note that $|X_i| \geq v_s/|\mathcal{S}|$. The property (B) of G_i established in Proposition 8 yields that $G[X_i]$ is 3-colorable, and hence its largest color class has size at least

$$\frac{v_s}{3|\mathcal{S}|} \geq \frac{v_s}{3\ell} \geq \frac{4v_s}{3P},$$

which finishes the proof. $\square$

CLAIM 10. H_B has an independent set of size at least $4v_b/P$.

Proof. Let m be the largest element of $\mathcal{B}$. Since G_m is (0.51ℓ) -colorable, $G[X_m]$ contains an independent set of size at least $1.9 \cdot |X_m|/\ell$. If $m = 1$, then $|X_m| = v_b$. On the other hand, if $m \geq 2$, then

$$1.9 \cdot |X_m| \geq |X_m| + 0.9 \cdot \sqrt{\ln n_m} > |X_m| + 1.2 \cdot n_{m-1} > |X_m| + \sum_{i=1}^{m-1} n_i \geq v_b.$$

We conclude that H_B has an independent set of size at least $v_b/\ell \geq 4v_b/P$. $\square$

If $v_s \geq 3|X|/4$, then we find an independent set of size at least $|X|/P$ in H_S by Claim 9. Otherwise, $v_b \geq |X|/4$, and Claim 10 guarantees an independent set in H_B of size at least $|X|/P$. $\square$

3.1. K_5 -free and iterated constructions. As we have already noted in section 2, the graph $G = \bigwedge_{i=1}^{\ell} G_i$ constructed in Theorem 2 can be equivalently viewed as $K_{\ell}\{G_1, G_2, \dots, G_{\ell}\}$. An adaptation of the proof of Theorem 2 will show that replacing K_{ℓ} by a graph from Proposition 8 yields another graph G^2 with $\chi_f(G^2) \sim (\rho(G^2))^2$. However, as all the graphs involved in the composition are now triangle-free, G^2 will be K_5 -free.

But we do not need to stop here. Since we have now much better control of the chromatic numbers of small subgraphs in G^2 than in the original graph G ,

replacing the graphs $G_i = G^1(n_i, C)$ in the composition by n_i -vertex variants of G^2 yields a graph G^3 with $\chi_f(G^3) \sim (\rho(G^3))^3$. Repeating this procedure k times leads to a construction of a graph G^{k+1} with $\chi_f(G^{k+1}) \sim (\rho(G^{k+1}))^{k+1}$.

In order to present our proofs of Theorems 3 and 4, we need to introduce some additional notation. Let us start with recalling Knuth's up-arrow notation

$$a \uparrow^{(k)} b = \begin{cases} a^b & \text{if } k = 1, \\ 1 & \text{if } k \geq 1 \text{ and } b = 0, \\ a \uparrow^{(k-1)} (a \uparrow^{(k)} (b-1)) & \text{otherwise,} \end{cases}$$

where $a, b, k \in \mathbb{N}$, and its inverse $a \downarrow^{(k)} n$, which is the largest integer b such that $n \geq a \uparrow^{(k)} b$. Using this, we define the following Ackermann-type function $F_k(b)$ and its inverse $f_k(b)$:

$$F_k(b) := 2 \uparrow^{(k)} b \quad \text{and} \quad f_k(b) := 2 \downarrow^{(k)} b.$$

Note that $F_1(b) = 2^b$ and $f_1(b) = \lfloor \log_2(b) \rfloor$, and for every $k \in \mathbb{N}$ it holds that $F_k(1) = 2$ and $F_k(2) = 4$. The functions also satisfy the following properties.

FACT 11. *For every $k \in \mathbb{N}$, the following hold:*

1. $f_k(f_k(F_{k+1}(n+2))) = F_{k+1}(n)$ for every $n \in \mathbb{N}$,
2. $f_{k+1}(4M) < f_k(f_k(M))$ for every $M \geq F_k(F_k(7))$, and
3. $\sum_{b=0}^n F_k(b) < F_k(n+1)$ for every $n \in \mathbb{N}$.

For a proof, see Appendix A. We are now ready to present the main lemma.

LEMMA 12. *Let C_0 be the constant from Proposition 8. For every $k \in \mathbb{N}$ and $C \geq C_0$ there is $n_0 := n_0(k, C)$ such that for all $n \geq n_0$ there is an n -vertex $K_{2^{k+1}}$ -free graph $G := G^k(n, C)$ with the following properties:*

- $\chi_f(G) \geq C^k$,
- $\rho(G) \leq 1.001 \cdot 3^k \cdot C$, and
- $G[W]$ is 3^k -colorable for every $W \subseteq V(G)$ such that $|W| \leq f_k(f_k(n))$.

Proof. For any fixed $C \geq C_0$, we proceed by induction on k . As the case $k = 1$ follows by letting $G := G^1(n, C)$ from Proposition 8, we may assume $k \geq 2$.

Let M be the smallest positive integer such that $f_k(4M) \leq f_{k-1}(f_{k-1}(M))$. Note that $M \leq F_{k-1}(F_{k-1}(7))$ by the second property of Fact 11. We set $n_0(k, C) := \max\{M, F_k(4 \cdot n_0(k-1, C))\}$. Given $n \geq n_0(k, C)$, we define m to be the largest integer such that

$$m + \sum_{i=2}^m F_k(m+3i-6) \leq n.$$

Note that $F_k(4m-1) > n$, as otherwise the third property of Fact 11 yields

$$(m+1) + \sum_{i=2}^{m+1} F_k(m+3i-6) \leq F_k(4m-1) \leq n,$$

contradicting the maximality of m . Therefore,

$$F_k(4m) > F_k(4m-1) > n_0(k, C) \geq F_k(4n_0(k-1, C)),$$

and hence $m > n_0(k-1, C)$. We set $b_1 := m$, and $b_i := F_k(m+3i-6)$ for every $i = 2, 3, \dots, m-1$. Finally, we set $b_m := n - \sum_{i < m} b_i$.

Let $H := G^1(m, C)$, and let $G_i := G^{k-1}(b_i, C)$ for all $i \in [m]$. We define $G := H\{G_1, G_2, \dots, G_m\}$. Clearly, the graph G contains no K_{2^k+1} . In the following three claims, we show that G has the desired three properties.

CLAIM 13. $\chi_f(G) \geq C^k$.

Proof. By the induction hypothesis, $\chi_f(H) \geq C$ and $\chi_f(G_i) \geq C^{k-1}$ for all $i \in [m]$. Therefore, Proposition 5 yields the desired lower-bound on $\chi_f(G)$. $\square$

CLAIM 14. $\rho(G) \leq 1.001 \cdot 3^k \cdot C$.

Proof. Fix an $X \subseteq V(G)$. Our aim is to show that $\alpha(G[X]) \geq |X| / (1.001 \cdot 3^k \cdot C)$. For $i \in [m]$, let X_i be $X \cap V(G_i)$. As in the proof of Theorem 2, let

$$\mathcal{S} := \{i \in [m] : |X_i| \leq f_{k-1}(f_{k-1}(b_i))\} \quad \text{and} \quad \mathcal{B} := [m] \setminus \mathcal{S}.$$

First, suppose the case $|\bigcup_{i \in \mathcal{S}} X_i| \geq |X|/3$. By the definition of $\mathcal{S}$ and the properties of G_i , every subgraph $G[X_i]$, where $i \in \mathcal{S}$, has an independent set of size at least $|X_i|/3^{k-1}$. On the other hand, $\chi(H) < 1.001 \cdot C$, so the projection of at least one of the color classes of the optimal coloring of H on $\bigcup_{i \in \mathcal{S}} X_i$ contains an independent set of size at least

$$\sum_{i \in \mathcal{S}} \frac{|X_i|}{3^{k-1}} \cdot \frac{1}{1.001 \cdot C} \geq \frac{|X|}{1.001 \cdot 3^k \cdot C}.$$

Now suppose $|\bigcup_{i \in \mathcal{B}} X_i| \geq 2|X|/3$, and let z be the maximum index in $\mathcal{B}$. If $z = 1$, then $|X_1| \geq 2|X|/3$. On the other hand, if $z \geq 2$, then

$$f_{k-1}(f_{k-1}(b_z)) \geq f_{k-1}(f_{k-1}(F_k(m + 3z - 6))) = F_k(m + 3z - 8) \geq \sum_{i < z} b_i,$$

where the equality and the last inequality follow from the first and the third properties of Fact 11, respectively. Therefore, $|X_z| \geq |X|/3$ and $G_z[X_z]$ contains an independent set of the sought size by $\rho(G_z) \leq 1.001 \cdot 3^{k-1} \cdot C$. $\square$

CLAIM 15. $G[W]$ is 3^k -colorable for every $W \subseteq V$ with $|W| \leq f_k(f_k(n))$.

Fix a set $W \subseteq V$ of size at most $f_k(f_k(n))$. First, let $Z := \{i : W \cap V(G_i) \neq \emptyset\}$. Clearly, $|Z| \leq |W| \leq f_k(f_k(n))$. Since $f_k(n) \leq 4m$ and $f_k(x) \ll \log_2 \log_2(x/4)$, we conclude that $|Z| \leq \log_2 \log_2(v(H))$. Therefore, there exists a proper 3-coloring of the induced subgraph $H[Z]$.

By the second property of Fact 11, for every $i \in [m]$ it holds that

$$|V(G_i) \cap W| \leq |W| \leq f_k(4m) \leq f_{k-1}(f_{k-1}(m)) \leq f_{k-1}(f_{k-1}(b_i)).$$

Therefore, the induction hypothesis yields that each $V(G_i) \cap W$ induces a 3^{k-1} -colorable subgraph of G , and hence $\chi(G[W]) \leq 3^k$ by Proposition 6. $\square$

Theorem 3 is a direct consequence of Lemma 12 applied with $k = 2$. It remains to establish Theorem 4.

Proof of Theorem 4. Let $P_0 := (2C_0)^2$. Given $P \geq P_0$, let $C := \sqrt{P/1.001}$ and $k := \lfloor \log_3 C \rfloor$. Applying Lemma 12 with k and C yields an $n_0(k, C)$ -vertex graph G with $\rho(G) \leq P$ and $\chi_f(G) \geq C^{\lfloor \log_3 C \rfloor} > e^{0.9 \cdot \ln^2(C)} > e^{\ln^2(P)/5}$. $\square$

4. Concluding remarks. We presented various constructions of graphs where the fractional chromatic number grows much faster than the Hall ratio, which refuted Conjecture 1. It is natural to ask whether the conclusion in Conjecture 1 can be relaxed and the fractional chromatic number of a graph is always upper-bounded by some function of its Hall ratio.

QUESTION 16. *Is there a function $g : \mathbb{R} \rightarrow \mathbb{R}$ such that $\chi_f(G) \leq g(\rho(G))$ for every graph G ?*

Theorem 4 shows that if such a function g exists, then $g(x) \geq e^{\ln^2(x)/5}$. While preparing our manuscript, we learned that Dvořák, Ossona de Mendez, and Wu [4] constructed graphs with Hall ratio at most 18 and arbitrary large fractional chromatic number. Therefore, the answer to Question 16 is no.

Conjecture 1 was partially motivated by another conjecture of Harris concerned with fractional colorings of triangle-free graphs, which was inspired by a famous result of Johansson [10] (for a recent short proof, see [14]) stating that $\chi(G) = O(\Delta/\ln \Delta)$ for every triangle-free graph G with maximum degree Δ .

CONJECTURE 17 ([7, Conjecture 6.2]). *There is C such that $\chi_f(G) \leq C \cdot d/\ln d$ for every triangle-free d -degenerate graph G .*

A classical result of Ajtai, Komlós, and Szemerédi [1] and an averaging argument together yield that $\rho(G) = O(d/\ln d)$ for G and d as above. Therefore, if Conjecture 1 could be recovered in the triangle-free setting, it would immediately yield the sought bound on χ_f in Conjecture 17.

QUESTION 18. *Is there C such that $\chi_f(G) \leq C \cdot \rho(G)$ for every triangle-free graph G ?*

In [11], it was been mentioned that the sequence of Mycelski graphs might provide a negative answer to Question 18, but we still do not know. For K_5 -free graphs, Theorem 3 shows that the answer is definitely negative. As a possibly simpler question, does the answer stay negative in the case of K_4 -free graphs?

QUESTION 19. *Is there C such that $\chi_f(G) \leq C \cdot \rho(G)$ for every K_4 -free graph G ?*

Let us conclude with an additional motivation for studying Conjecture 17. Esperet, Kang, and Thomassé [5] conjectured that dense triangle-free graphs must contain dense induced bipartite subgraphs.

CONJECTURE 20 ([5, Conjecture 1.5]). *There exists $C > 0$ such that any triangle-free graph with minimum degree at least d contains an induced bipartite subgraph of minimum degree at least $C \cdot \ln d$.*

Erdős–Rényi random graphs of the appropriate density show that the bound would be, up to the constant C , best possible. A relation between the fractional chromatic number and induced bipartite subgraphs proven in [5, Theorem 3.1] shows that if Conjecture 17 holds, then Conjecture 20 holds as well. Very recently, Kwan, Letzter, Sudakov, and Tran [12] proved a slightly weaker version of Conjecture 20 where the bound $C \cdot \ln n$ is replaced by $C \cdot \ln n / \ln \ln n$.

Appendix A. Proof of Fact 11. The definitions of f_k and F_{k+1} readily yield that $f_k(F_{k+1}(n+1)) = F_{k+1}(n)$. Therefore, $f_k(f_k(F_{k+1}(n+2))) = F_{k+1}(n)$, proving the first property.

For every $k, n \in \mathbb{N}$, a straightforward induction yields that $F_k(n) \geq n+1$. This in turn implies that $F_{k+1}(n) = F_k(F_{k+1}(n-1)) \geq F_k(n) \geq 2^n$. Similarly, for all $k \in \mathbb{N}$, the functions $F_k(\cdot)$ and $f_k(\cdot)$ are monotone nondecreasing. Therefore, for all $k \in \mathbb{N}$ and $n \geq 7$, it holds that

$$F_{k+1}(n) = F_k(F_k(F_k(F_k(F_{k+1}(n-4))))) \geq 2^{F_k(F_k(2^{n-4})+1)} \geq 4 \cdot F_k(F_k(n+1)+1).$$

Since $F_k(f_k(M)+1) > M \geq F_k(f_k(M))$, we assert that $f_{k+1}(4M) < f_k(f_k(M))$ for

all $M \geq F_k(F_k(7))$. Indeed, otherwise

$$4M \geq F_{k+1}(f_{k+1}(4M)) \geq F_{k+1}(f_k(f_k(M))) \geq 4 \cdot F_k(F_k(f_k(f_k(M)) + 1) + 1) > 4M,$$

a contradiction. This concludes the proof of the second property.

The last property is proven by induction on k . Indeed, the case $k = 1$ is the sum of a geometric progression. If $k \geq 2$, then by the induction hypothesis

$$\sum_{b=0}^n F_{k+1}(b) = \sum_{b=0}^n F_k(F_{k+1}(b-1)) \leq \sum_{i=0}^{F_{k+1}(n-1)} F_k(i) < F_k(F_{k+1}(n-1) + 1).$$

However, the right-hand side is at most $F_k(F_k(F_{k+1}(n-1))) = F_{k+1}(n+1)$.

Acknowledgments. The authors thank Zdeněk Dvořák for sharing the manuscript of [4] with them, and Józsi Balogh, Sasha Kostochka, and Yani Pehova for fruitful discussions at the beginning of the project. The authors also thank the anonymous referees for carefully reading the manuscript and for their valuable comments, which greatly improved the presentation of the results.

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