

Threshold Ramsey multiplicity for paths and even cycles

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Abstract

The Ramsey number $r(H)$ of a graph H is the minimum integer n such that any two-coloring of the edges of the complete graph K_n contains a monochromatic copy of H . While this definition only asks for a single monochromatic copy of H , it is often the case that every two-edge-coloring of the complete graph on $r(H)$ vertices contains many monochromatic copies of H . The minimum number of such copies over all two-colorings of $K_{r(H)}$ will be referred to as the threshold Ramsey multiplicity of H . Addressing a problem of Harary and Prins, who were the first to systematically study this quantity, we show that there is a positive constant c such that the threshold Ramsey multiplicity of a path or an even cycle on k vertices is at least $(ck)^k$. This bound is tight up to the constant c . We prove a similar result for odd cycles in a companion paper.

1 Introduction

The *Ramsey number* $r(H)$ of a graph H is the minimum positive integer n such that any two-coloring of the edges of the complete graph K_n on n vertices contains a monochromatic copy of H . Determining Ramsey numbers is a challenging task and the exact value of $r(H)$ is known in only a few special cases. For example, determining the Ramsey number of K_5 , the complete graph on five vertices, is a well-known open problem.

The few non-trivial families for which the Ramsey number is known exactly include paths and cycles. To say more, we let P_k denote the path on k vertices and C_k the cycle on k vertices. The length of a path or cycle denotes its number of edges, so P_k has length $k - 1$ and C_k has length k . In 1967, Gerencsér and Gyárfás [15] determined the Ramsey number of paths, showing that

$$r(P_k) = k - 1 + \lfloor k/2 \rfloor.$$

It is simple to show that $r(C_3) = r(C_4) = 6$, while, for $k \geq 5$, Faudree and Schelp [13] and, independently, Rosta [30] proved that

$$r(C_k) = k + k/2 - 1 \text{ if } k \text{ is even and } r(C_k) = 2k - 1 \text{ if } k \text{ is odd.}$$

A more general problem is to determine the *Ramsey multiplicity* $M(H, n)$, defined to be the minimum number of monochromatic copies of H that appear in any two-edge-coloring of K_n . In particular, $M(H, n) = 0$ if and only if $n < r(H)$, so the problem of determining $M(H, n)$ does indeed generalize the problem of determining $r(H)$.

In 1962, Erdős [11] conjectured that if H is a clique, then $M(H, n)$ is asymptotically equal to the expected number of monochromatic copies of H in a uniformly random two-coloring of the edges of K_n

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and Burr and Rosta [1] later generalized this conjecture to all graphs H . While true for K_3 , a result of Goodman [16] that predates the conjecture and doubtless inspired it, Thomason [36] showed that it is already false for K_4 . Despite the failure of this attractive conjecture, the asymptotic behavior of $M(H, n)$ for fixed H and n tending to infinity has drawn considerable attention (see [2, 14, 17, 22, 23] for some examples or [4] for a survey). In particular, it is known that any bipartite graph which satisfies the well-known conjecture of Erdős–Simonovits [34] and Sidorenko [32, 33] also satisfies the Burr–Rosta conjecture, so the considerable recent progress [3, 7, 8, 9, 21, 25, 28, 35] on Sidorenko’s conjecture, as it is usually known, may also be interpreted as progress on our understanding of Ramsey multiplicity.

Besides the case where H is fixed and n tends to infinity, another much-studied problem asks for the value of $M(H, n)$ when it first becomes positive, that is, when $n = r(H)$. To distinguish it from the more general Ramsey multiplicity function $M(H, n)$, we call this value the *threshold Ramsey multiplicity*.

Definition 1 (Threshold Ramsey multiplicity). The *threshold Ramsey multiplicity* $m(H)$ of a graph H is the minimum number of monochromatic copies of H in any two-coloring of the edges of K_n with $n = r(H)$. In other words,

$$m(H) = M(H, r(H)).$$

The threshold Ramsey multiplicity was first studied systematically by Harary and Prins [20] almost fifty years ago. Their work and subsequent work by Schwenk [19] and by Piwakowski and Radziszowski [29] determines the threshold Ramsey multiplicity for all graphs with at most four vertices. However, in general, the problem of determining or even giving a non-trivial lower bound on the threshold Ramsey multiplicity appears extremely difficult. This is in part because it seems necessary to first determine the Ramsey number, a problem which is already hard, before one can say anything substantive about the threshold Ramsey multiplicity.

The only family for which $m(H)$ is known is for stars, where Harary and Prins [20] proved that $m(K_2) = 1$ and $m(K_{1,k}) = 1$ for k even, but $m(K_{1,k}) = 2k$ for $k \geq 3$ odd, surprisingly erratic behavior for such a simple family. In the same paper, Harary and Prins asked for a determination of the threshold Ramsey multiplicity for paths and cycles, probably the next simplest families after stars. The main result of this paper is an approximate answer to their question for paths and even cycles. The case of odd cycles will be discussed in the companion paper [5].

To the best of our knowledge, the only previous work concerning these questions is due to Rosta and her collaborators, who looked at the case of odd cycles. In her first paper on the subject, with Surányi [31], she obtained the exponential lower bound $m(C_k) \geq 2^{ck}$. This was later improved to a superexponential lower bound in an unpublished work. More recently, Károlyi and Rosta [24] improved the lower bound to $m(C_k) \geq k^{ck}$, which we will see below is sharp up to the constant in the exponent. However, their method has little to say about paths and even cycles, the main objects of interest in this paper, in large part because the Ramsey numbers of these graphs are significantly smaller than the Ramsey number of odd cycles of comparable size.

Our main result, proved in this paper and its companion [5], is the following.

Theorem 2. *There is a positive constant c such that, for every positive integer k , the threshold Ramsey multiplicity of paths and cycles on k vertices satisfy $m(P_k) \geq (ck)^k$ and $m(C_k) \geq (ck)^k$.*

We prove Theorem 2 for paths and even cycles in this paper, while the case of odd cycles is handled in the companion paper [5].

The bound in Theorem 2 is easily seen to be tight up to the constant c , since the total number of paths or cycles on k vertices in the complete graph with $r(P_k)$ or $r(C_k)$ vertices is at most $(c'k)^k$ for some constant c' . However, we may also pinpoint some edge-colorings which we believe to be optimal for $m(P_k)$ and $m(C_k)$. Consider the edge-coloring $\chi(a, b)$ of the complete graph on $n = a + b$ vertices with vertex set $A \cup B$, $|A| = a$ and $|B| = b$, where A and B form blue cliques and all edges between A and B are red. Let $a_0 = k - 1$ and $b_0 = \lfloor k/2 \rfloor - 1$. The coloring $\chi(a_0, b_0)$ does not contain a monochromatic P_k and gives

the tight lower bound for the Ramsey number of the path P_k . If k is even, the colorings $\chi(a_0 + 1, b_0)$ and $\chi(a_0, b_0 + 1)$ of the complete graph on $a_0 + b_0 + 1 = r(P_k)$ vertices each have exactly $k!/2$ monochromatic P_k . If k is odd, the coloring $\chi(a_0, b_0 + 1)$ of the complete graph on $a_0 + b_0 + 1 = r(P_k)$ vertices has exactly $\frac{(k-1)}{4}(k-1)!$ monochromatic P_k . Not only do these colorings show that Theorem 2 is tight up to the constant c for paths, but we conjecture that they realize the threshold Ramsey multiplicity for k sufficiently large.

Conjecture 3. *For sufficiently large k , if k is even, then $m(P_k) = k!/2$ and if k is odd, then $m(P_k) = \frac{(k-1)}{4}(k-1)!$.*

As P_k is a subgraph of C_k , the edge-coloring $\chi(a_0, b_0)$ with $a_0 = k - 1$ and $b_0 = \lfloor k/2 \rfloor - 1$ described above also does not contain a monochromatic C_k . For $k \geq 6$ even, this coloring realizes the tight lower bound on $r(C_k)$. The coloring formed from $\chi(a_0 + 1, b_0)$ by changing the color of one edge in the monochromatic blue clique of order $a_0 + 1 = k$ to red does not have a monochromatic red C_k and thus has $(k-1)!/2 - (k-2)! = \frac{(k-3)}{2}(k-2)!$ monochromatic C_k . We conjecture that for k sufficiently large this is the threshold Ramsey multiplicity for the even cycle C_k .

If k is odd, then the coloring $\chi(k-1, k-1)$ has no monochromatic C_k and realizes the tight lower bound on the Ramsey number $r(C_k)$. In this case, the edge-coloring $\chi(k, k-1)$ has all monochromatic C_k in the blue clique of order k and thus has $(k-1)!/2$ monochromatic C_k . We conjecture that for k sufficiently large this is the threshold Ramsey multiplicity for the odd cycle C_k .

Conjecture 4. *For sufficiently large k , if k is even, then $m(C_k) = \frac{(k-3)}{2}(k-2)!$ and if k is odd, then $m(C_k) = (k-1)!/2$.*

The rest of the paper is dedicated to the proof of Theorem 2 in the case of paths and even cycles. Because we focus entirely on this case, we will often use the phrase Theorem 2 as a shorthand to mean Theorem 2 for paths and even cycles. We note that we have made no attempt to optimize the value of the constant c in Theorem 2. Throughout the proof, we have also chosen to omit floor and ceiling signs whenever they are not essential.

2 Proof of Theorem 2 for paths and even cycles

Szemerédi’s regularity lemma (see Lemma 7 below) will be an important tool in our proof. Given any graph, the regularity lemma shows that there is a vertex partition of the graph into a small number of parts of almost equal size, where the bipartite graph between almost every pair of parts is random-like. This property is useful for many purposes, particularly for embedding and counting sparse subgraphs. For an excellent (though now somewhat outdated) survey, we refer the interested reader to [27].

To state the regularity lemma, we need some definitions making precise what is meant by saying that the graph between two vertex sets is “random-like”. For a pair of vertex subsets (X, Y) of a graph, let $e(X, Y)$ denote the number of pairs in $X \times Y$ that are edges and $d(X, Y) = e(X, Y)/|X||Y|$ denote the density of edges between X and Y .

Definition 5 (ϵ -regular pair). A pair of vertex subsets (X, Y) of a graph is ϵ -regular if, for all subsets $U \subset X, V \subset Y$ such that $|U| \geq \epsilon|X|$ and $|V| \geq \epsilon|Y|$, $|d(U, V) - d(X, Y)| \leq \epsilon$.

The following lemma collects some basic facts which follow easily from this definition.

Lemma 6. *If (X, Y) is an ϵ -regular pair in a graph G and $d = d(X, Y)$, then the following hold:*

- (i) *If $Y' \subset Y$ satisfies $|Y'| \geq \epsilon|Y|$, then the number of vertices in X with degree in Y' greater than $(d + \epsilon)|Y'|$ is less than $\epsilon|X|$ and the number of vertices in X with degree in Y' less than $(d - \epsilon)|Y'|$ is less than $\epsilon|X|$.*

- (ii) If $X' \subset X$ and $Y' \subset Y$ are such that $|X'| \geq \alpha|X|$ and $|Y'| \geq \alpha|Y|$, then (X', Y') is $\max(\epsilon/\alpha, 2\epsilon)$ -regular.
- (iii) Provided X and Y are disjoint, the pair (X, Y) is also ϵ -regular in the complement of G .

A partition of a set is said to be *equitable* if each pair of parts differ in size by at most one. With this definition, we can now state the regularity lemma in a standard colored form, whose equivalence to the usual form follows easily from Lemma 6(iii).

Lemma 7 (Szemerédi's regularity lemma). *For every $\epsilon > 0$ and positive integer m_0 , there exist positive integers M_0 and n_0 such that every two-edge-coloring of the complete graph K_n with $n \geq n_0$ in colors red and blue admits an equitable vertex partition $V_1 \cup \dots \cup V_M$ into M parts with $m_0 \leq M \leq M_0$ where all but at most $\epsilon \binom{M}{2}$ pairs (V_i, V_j) of parts with $1 \leq i < j \leq M$ are ϵ -regular in both the red and blue subgraphs.*

We remark that there is a strengthening of the regularity lemma, proved in [6], where each part is ϵ -regular with all but an ϵ -fraction of the other parts and each part is also ϵ -regular with itself. Working with this variant rather than Lemma 7 would allow us to simplify our proof very slightly. However, since this variant is, as yet, non-standard, we have opted to work with the usual version instead.

Once we have the partition guaranteed by the regularity lemma, it is often convenient to consider a simplified rendering of the graph called the *reduced graph* of the partition. By saying that a graph is *red/blue-multicolored*, we will mean that each edge is colored either blue, red, or both blue and red.

Definition 8. Given a red/blue-edge-colored graph G , a partition $V_1 \cup \dots \cup V_M$ of its vertex set and parameters $0 < \epsilon, d < 1$, the *reduced graph* $H = H(\epsilon, d)$ of the partition with parameters ϵ and d is the red/blue-multicolored graph with vertex set $[M]$ and a red (respectively, blue) edge between i and j if and only if (V_i, V_j) is ϵ -regular with density at least d in the red (respectively, blue) graph.

With this preliminary, we may now give a broad outline of the proof of Theorem 2.

2.1 Proof Outline

We first prove that in any red/blue edge-coloring of the complete graph K_n , there is a color and an almost spanning subset W of the vertices such that, for any two vertices of W , there are many short paths between them in the specified color. We then apply Szemerédi's regularity lemma to the subgraph of K_n induced by W , obtaining a reduced graph. If, in this reduced graph, we can find a large monochromatic matching, then we can build as many of the required paths and even cycles as we need. This case will be discussed in detail in Section 2.3.1.

If, instead, there is no sufficiently large monochromatic matching in the reduced graph, then a key stability result (Lemma 11 below) shows that the original two-colored graph G induced by the vertex set W is close to a certain shape (described in Definition 10). In this case, we can directly bound the number of paths and even cycles to complete the proof. The details of this case may be found in Section 2.3.2.

A variant of our stability lemma already appeared in the work of Gyárfás, Sárközy, and Szemerédi [18]. However, the version we need is somewhat stronger, so we include a complete proof in Section 3. One point worth noting is that we make an appeal to the regularity lemma in our statement and proof, whereas the stability lemma in [18] is proved without it. We now describe our version in more detail.

2.2 The stability lemma

The next two definitions already appear in the work of Gyárfás, Sárközy, and Szemerédi [18], though the first is stated in slightly more generality than in [18].

Definition 9 (Well-connected). A vertex subset W of a graph G is (t, l) -*well-connected* if any two vertices $u, v \in W$ are connected by at least t internally vertex-disjoint paths of length at most l . Note that any vertex in $V(G) \setminus \{u, v\}$ is allowed as an internal vertex for these paths.

We will often refer to a vertex set as being *well-connected in a particular color*, meaning that the vertex set is well-connected with respect to the graph consisting of edges in that color. For the second definition, given a graph G and disjoint vertex subsets A and B , we let $G[A]$ denote the induced subgraph of G with vertex set A and $G[A, B]$ the bipartite graph with parts A and B whose edges are the edges of G between A and B . Note that the density within a set X is given by $d(X, X) = e(X, X)/|X|^2 = 2e(X)/|X|^2$.

Definition 10 (Extremal coloring with parameter α). A two-coloring of the edges of a graph G is an *extremal coloring with parameter α* if there exists a partition $V(G) = A \cup B$ such that

- $|A| \geq (2/3 - \alpha)|V(G)|$ and $|B| \geq (1/3 - \alpha)|V(G)|$ and
- the graph $G[A]$ has density at least $(1 - \alpha)$ in some color and the bipartite graph $G[A, B]$ has density at least $(1 - \alpha)$ in the other color.

Our key stability lemma is now as follows. Roughly speaking, it says that every two-coloring of the edges of K_n is either close to an extremal coloring or the reduced graph contains a monochromatic matching covering more than $2/3$ of the vertices such that the underlying vertex set is well-connected in the same color.

Lemma 11. *For any $0 < \epsilon \leq 10^{-10}$ and $d, \lambda \geq 1000\epsilon$, there is a positive integer $M_0 = M_0(\epsilon)$ such that if n is sufficiently large in terms of ϵ , then any two-coloring of the edges of the complete graph K_n falls into at least one of the following two cases:*

- **Case 1:** *There is a positive integer $M \leq M_0$ and disjoint vertex subsets $U_1, \dots, U_m, V_1, \dots, V_m$ with $m \geq (2/3 + \lambda)M/2$ such that each $|U_i|, |V_i| \geq cn$ with $c \geq (1 - \epsilon)/M$, all pairs (U_i, V_i) are simultaneously ϵ -regular in some color with the edge density in that color at least $d - \epsilon$, and $\bigcup_{i=1}^m U_i \cup \bigcup_{i=1}^m V_i$ is $(200M, 6)$ -well-connected in the same color.*
- **Case 2:** *The coloring is an extremal coloring with parameter $1000(d + \lambda + \sqrt{\epsilon})$.*

Observe that if $\alpha \geq 2/3$, any two-coloring of the edges of a complete graph is trivially an extremal coloring with parameter α , since we may take A to be the empty set. It follows that we may assume $d, \lambda \leq 1/1000$ in Lemma 11.

2.3 Proof of Theorem 2 assuming Lemma 11

We now prove Theorem 2 by applying Lemma 11 with $d = 20\sqrt{\epsilon}$ and $\lambda = 13\sqrt{\epsilon}$.

2.3.1 Proof of Theorem 2 in the situation of Case 1 of Lemma 11

We first prove Theorem 2 for paths for edge-colorings satisfying Case 1 of Lemma 11 with the following approach. Roughly speaking, in the graph of the color given in this case, between any regular pair (U_i, V_i) with density $d(U_i, V_i) = d$, there should be many paths of length close to $2cn$. Since the bipartite graph between U_i and V_i is random-like, the count of paths of length l is roughly at least $d^l \prod_{i=0}^l (cn - \lfloor i/2 \rfloor)$. Since the union of the U_i and V_i is well-connected, any two vertices in this union are connected by many internally vertex-disjoint short paths. We can then find many long paths P_k by using the short paths guaranteed by the well-connectedness property to connect the end vertices of the paths from different pairs (U_i, V_i) . In this section, we will make this idea rigorous.

The following two lemmas show that for a regular pair (U, V) in a graph G the number of long paths starting from any vertex of large degree or between any pair of vertices of large degree in the bipartite graph $G[U, V]$ is roughly at least the expected count if $G[U, V]$ were a random graph of the same density.

Lemma 12. *Suppose (U, V) is an ϵ -regular pair of disjoint vertex subsets of a graph G such that $|U|, |V| \geq n$ and $d(U, V) = d$. If $n \geq \epsilon^{-2}$ and $d > \epsilon + \sqrt{\epsilon}$, then, for any vertex $v \in V$ with at least $(d - \epsilon)|U|$ neighbors in U and any positive integer $l \leq 2(1 - \sqrt{\epsilon})n - 1$, there are at least $(d - \epsilon - \sqrt{\epsilon})^l \prod_{i=1}^l (n - \lfloor i/2 \rfloor)$ paths of length l in $G[U, V]$ starting from v .*

Proof. Let N_j be the number of paths P of length j in $G(U, V)$ of the form $v_0 = v, v_1, \dots, v_j$ starting from v for which there are at least $(d - \epsilon)(|U| - \lfloor (j+1)/2 \rfloor)$ ways to extend the path if j is even and at least $(d - \epsilon)(|V| - \lfloor (j+1)/2 \rfloor)$ ways to extend the path if j is odd. By extending the path, we mean finding a vertex v_{j+1} that is adjacent to v_j but distinct from the vertices in P . We will prove by induction on j that for $j \leq 2(1 - \sqrt{\epsilon})n - 2$, we have $N_j \geq (d - \epsilon - \sqrt{\epsilon})^j \prod_{i=1}^j (n - \lfloor i/2 \rfloor)$, which easily implies the lemma.

Clearly $N_0 = 1$, since a path with zero edges starting from v is just v itself and it is extendable in sufficiently many ways by the degree condition on v . This is the base case of the induction.

Suppose now that we have the claimed lower bound on N_j for some $j \leq 2(1 - \sqrt{\epsilon})n - 3$ and we wish to prove the lower bound on N_{j+1} . Suppose j is even (the case where j is odd can be handled in exactly the same way). Let $P : v_0 = v, \dots, v_j$ be a path in $G(U, V)$ of length j which can be extended in at least $(d - \epsilon)(|U| - \lfloor (j+1)/2 \rfloor)$ ways. Then $v_j \in V$ and there are at least $(d - \epsilon)(|U| - \lfloor (j+1)/2 \rfloor)$ neighbors of v_j in U which are not in P . We let U' be this set of neighbors. As the pair (U, V) is ϵ -regular and $|V \setminus \{v_0, v_1, \dots, v_j\}| = |V| - \lfloor (j+2)/2 \rfloor \geq \epsilon|V|$, Lemma 6(i) implies that there are fewer than $\epsilon|U|$ vertices in U whose degree in $V \setminus \{v_0, v_1, \dots, v_j\}$ is less than $(d - \epsilon)|V \setminus \{v_0, v_1, \dots, v_j\}| = (d - \epsilon)(|V| - \lfloor (j+2)/2 \rfloor)$. Therefore, the number of vertices in U' which can be used as v_{j+1} and added to P so that this longer path is extendable in sufficiently many ways is at least

$$|U'| - \epsilon|U| \geq (d - \epsilon)(|U| - \lfloor (j+1)/2 \rfloor) - \epsilon|U| \geq (d - \epsilon)(n - \lfloor (j+1)/2 \rfloor) - \epsilon n \geq (d - \epsilon - \sqrt{\epsilon})(n - \lfloor (j+1)/2 \rfloor),$$

where the final inequality follows from the upper bound on j assumed in the lemma. Hence, $N_{j+1} \geq N_j(d - \epsilon - \sqrt{\epsilon})(n - \lfloor (j+1)/2 \rfloor)$. By the lower bound on N_j , we obtain the desired lower bound on N_{j+1} , completing the induction. $\square$

Lemma 13. *Suppose (U, V) is an ϵ -regular pair of disjoint vertex subsets of a graph G such that $|U|, |V| \geq n$ with $n \geq 5\epsilon^{-2}$ and $d(U, V) = d$ with $d > 5\sqrt{\epsilon}$. Let $u, v \in U \cup V$ be distinct vertices which are each adjacent to at least a $(d - \epsilon)$ fraction of the vertices in the other part. Suppose l is an integer with $3 \leq l \leq 2(1 - 2\sqrt{\epsilon})n$, where l is even if u and v are in the same part and l is odd if u and v are in different parts. Then the number of paths of length l in $G[U, V]$ with end vertices u and v is at least $(d - 7\sqrt{\epsilon})^{l-1}(\epsilon n) \prod_{i=1}^{l-2} (n - \lfloor i/2 \rfloor)$.*

Proof. We will focus on the case where $u \in U$ and $v \in V$. The case where u and v are in the same part can be handled similarly.

As $|N(u)| \geq (d - \epsilon)|V| \geq \epsilon|V| + 1$, we can set aside $\epsilon|V|$ neighbors of u (not including v) and remove them from V , calling this set of $\epsilon|V|$ neighbors V_0 . We will only use these vertices in the last step to connect with u . As $1 \geq d > 5\sqrt{\epsilon}$, we have $\epsilon < 1/25$. By Lemma 6(ii) with $\alpha = 1 - \epsilon$, and noting that $\max(\epsilon/(1 - \epsilon), 2\epsilon) = 2\epsilon$, the pair $(V \setminus V_0, U \setminus \{u\})$ is 2ϵ -regular.

Let l be an odd positive integer. Our aim is to give a lower bound on the number of paths of length l with end vertices u and v . Suppose that we fix a path of length $l - 3$ starting from v , say $P : w_0 = v, w_1, \dots, w_{l-3}$, such that the vertices are in $(V \setminus V_0) \cup (U \setminus \{u\})$ and there are at least $(d - 2\epsilon)(|U| - 1 - \lfloor (l - 2)/2 \rfloor)$ ways to extend the path to a vertex $w_{l-2} \in U$. Let W_P be this set of candidate vertices for w_{l-2} . Then

$$\begin{aligned} |W_P| &\geq (d - 2\epsilon)(|U| - 1 - \lfloor (l - 2)/2 \rfloor) \geq (d - 2\epsilon)(|U| - 1 - (2(1 - 2\sqrt{\epsilon})n - 2)/2) \\ &\geq (d - 2\epsilon)(|U| - (1 - 2\sqrt{\epsilon})|U|) = (d - 2\epsilon)(2\sqrt{\epsilon}|U|) \geq \epsilon|U|. \end{aligned}$$

As (U, V) is ϵ -regular, $|W_P| \geq \epsilon|U|$, and $|V_0| \geq \epsilon|V|$, the number of edges $(w_{l-2}, w_{l-1}) \in W_P \times V_0$ satisfies

$$e(W_P, V_0) \geq (d - \epsilon)|W_P||V_0| > (d - 2\epsilon)^2(|U| - \lfloor (l - 2)/2 \rfloor) \cdot (\epsilon|V|).$$

We can obtain a path of length l from v to u by beginning with the path P of length $l - 3$, followed by any pair (w_{l-2}, w_{l-1}) of adjacent vertices as above, and finally ending with u .

For any non-negative integer i , let N_i be the total number of paths $P : v_0 = v, v_1, \dots, v_i$ of even length i in the bipartite graph $G[V \setminus V_0, U \setminus \{u\}]$ starting from v for which the number of ways to extend the path is at least $(d - 2\epsilon)(|U| - 1 - \lfloor (i+1)/2 \rfloor)$. Applying Lemma 12 with n replaced by $(1 - \epsilon)n$, d replaced by $d(V \setminus V_0, U \setminus \{u\}) \geq d - \epsilon$, and ϵ replaced by 2ϵ , we deduce that the number N_{l-3} of such paths P of length $l - 3$ is at least

$$((d - \epsilon) - 2\epsilon - \sqrt{2\epsilon})^{l-3} \prod_{i=1}^{l-3} ((1 - \epsilon)n - \lfloor i/2 \rfloor),$$

where we can apply Lemma 12 since the conditions on path length, density, and the number of vertices are all satisfied. Therefore, the number of paths of length l with end vertices u and v is at least

$$\begin{aligned} & N_{l-3} \cdot (d - 2\epsilon)^2 (n - \lfloor (l-2)/2 \rfloor) \cdot (\epsilon n) \\ & \geq ((d - \epsilon) - 2\epsilon - \sqrt{2\epsilon})^{l-3} \prod_{i=1}^{l-3} ((1 - \epsilon)n - \lfloor i/2 \rfloor) \cdot (d - 2\epsilon)^2 (n - \lfloor (l-2)/2 \rfloor) \cdot (\epsilon n) \\ & \geq (d - 5\sqrt{\epsilon})^{l-1} (1 - \epsilon - \sqrt{\epsilon})^{l-3} (\epsilon n) \prod_{i=1}^{l-2} (n - \lfloor i/2 \rfloor) \\ & \geq (d - 7\sqrt{\epsilon})^{l-1} (\epsilon n) \prod_{i=1}^{l-2} (n - \lfloor i/2 \rfloor). \end{aligned}$$

The second inequality holds since $(1 - \epsilon)n - \lfloor i/2 \rfloor \geq (1 - \epsilon - \sqrt{\epsilon})(n - \lfloor i/2 \rfloor)$ for $i \leq 2(1 - 2\sqrt{\epsilon})n$. $\square$

We now prove the path case of Theorem 2 when the coloring satisfies Case 1 of Lemma 11.

Theorem 2 for paths for colorings satisfying Case 1 of Lemma 11. Fix $0 < \epsilon \leq 10^{-20}$ and let $d = 20\sqrt{\epsilon}$ and $\lambda = 13\sqrt{\epsilon}$. Suppose there are vertex subsets $U_1, \dots, U_m, V_1, \dots, V_m$ with $m = (2/3 + \lambda)M/2$ and $|V_i|, |U_i| \geq cn$ satisfying the properties of Case 1 of Lemma 11, say in color red. We may assume that n is sufficiently large in terms of c, ϵ , and M . Let $d' = d - \epsilon$, so the edge density between each pair (U_i, V_i) is at least d' . We will show that there is a constant $c' > 0$ such that the number of monochromatic paths with $k = \lceil 2(n+1)/3 \rceil$ vertices is at least $(c'k)^k$.

We give a lower bound on the number of paths with k vertices in the red graph G by first choosing a pair of anchor vertices $(v_i, u_i) \in V_i \times U_i$ for each $1 \leq i \leq m$ and then picking short paths P_i to join u_i and v_{i+1} and long paths T_i to join u_i and v_i , where we will use Lemma 13 to show that there are many paths T_i connecting u_i and v_i within $G(U_i, V_i)$ that avoid the vertices of all the short paths P_j .

From each V_i and U_i , pick vertices $v_i \in V_i$ and $u_i \in U_i$ as anchor vertices such that each is adjacent to at least a $(d' - \epsilon)$ -fraction of the vertices in the other part. Since (V_i, U_i) is ϵ -regular, there are at least $(1 - \epsilon)cn$ choices for each of v_i and u_i .

After fixing the choice of pairs of anchor vertices $(v_i, u_i)_{1 \leq i \leq m}$, we now pick a set of short disjoint paths P_i to connect u_i to v_{i+1} for each $1 \leq i \leq m - 1$. By assumption, the vertex set $\bigcup_{i=1}^m V_i \cup \bigcup_{i=1}^m U_i$ is $(200M, 6)$ well-connected. For $1 \leq i \leq m - 1$, we will greedily pick a red path P_i of length at most six to connect u_i and v_{i+1} . In total, we will pick $m - 1$ paths; together with v_1 and u_m , there will be at most $7(m - 1) + 2 \leq 7m$ vertices in all the P_i 's and anchor vertices. Since there are at least $200M$ internally vertex-disjoint paths of length at most six connecting u_i and v_{i+1} by the definition of a $(200M, 6)$ -well-connected set and $200M > 100(2/3 + \lambda)M/2 = 100m > 7m$, we can greedily choose these $m - 1$ paths such that they are vertex disjoint and internally do not use any anchor vertices.

After fixing P_i for $1 \leq i \leq m - 1$, we will use long paths T_i to connect each pair (v_i, u_i) . In each regular pair (V_i, U_i) , we remove the internal vertices of the $m - 1$ paths P_i (so that at most $7m$ vertices

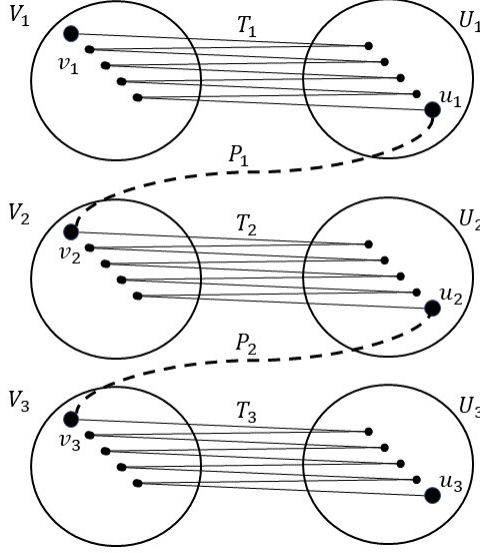


Figure 1: An illustration showing the anchor vertices v_i and u_i , the short paths P_i , and the long paths T_i used to build paths with k vertices.

are removed). Removing only a few further vertices if necessary (but not removing v_i or u_i), we may suppose that the resulting subsets $V'_i \subset V_i, U'_i \subset U_i$ satisfy

$$|V'_i| = |U'_i| = cn - 7m.$$

Since $|V'_i| \geq \epsilon|V_i|$ and $|U'_i| \geq \epsilon|U_i|$, the fact that (V_i, U_i) is ϵ -regular implies that $d(V'_i, U'_i) \geq d' - \epsilon$. Moreover, by Lemma 6(ii), as $cn \geq 14m$, the pair (V'_i, U'_i) is 2ϵ -regular. Furthermore, since $\epsilon cn > 7m$, v_i has at least

$$(d' - \epsilon)cn - 7m > (d' - 2\epsilon)(cn - 7m)$$

neighbors in U'_i and similarly for u_i . Let ℓ_0 be the largest odd integer not larger than $\lfloor 2(1 - 2\sqrt{2\epsilon})(cn - 7m) \rfloor$. By Lemma 13 with d replaced by $d' - \epsilon$, ϵ by 2ϵ , and n by $cn - 7m$, for odd $l \leq \ell_0$, the number of paths of length l connecting v_i and u_i is at least

$$\begin{aligned} (d' - \epsilon - 7\sqrt{2\epsilon})^{l-1} (2\epsilon(cn - 7m)) \prod_{i=1}^{l-2} (cn - 7m - \lfloor i/2 \rfloor) &\geq 2\epsilon(d' - 8\sqrt{2\epsilon})^{l-1} \prod_{i=0}^{l-2} (cn - 7m - \lfloor i/2 \rfloor) \\ &\geq 2\epsilon(d' - 8\sqrt{2\epsilon})^{l-1} \frac{2\pi}{e^2} ((cn - 7m)/e)^{l-1} = \frac{4\epsilon\pi}{e^2} \left((d' - 8\sqrt{2\epsilon})(cn - 7m)/e \right)^{l-1}. \end{aligned} \quad (1)$$

In the last inequality, we used the fact that $a!/b! \geq \frac{\sqrt{2\pi}}{e} \left(\frac{a}{e}\right)^{a-b}$ for positive integers $a > b$, which easily follows from the upper and lower bounds in Stirling's approximation for factorials. Thus, $\frac{a!}{(a-l_1)!} \frac{a!}{(a-l_2)!} \geq \frac{2\pi}{e^2} \left(\frac{a}{e}\right)^{l_1+l_2}$, which we applied with $a = cn - 7m$, $l_1 = \lfloor (l-2)/2 \rfloor + 1$, and $l_2 = \lfloor (l-3)/2 \rfloor + 1$.

Therefore, within each bipartite graph $G(U'_i, V'_i)$, there are many choices for the path T_i of any fixed odd length between 3 and ℓ_0 . Recall that the way we intend to build paths of length $k-1$ is by alternately concatenating T_i and P_i . If all m pairs give rise to a path T_i of length ℓ_0 , the total length of these T_i 's, which is also a lower bound on the length of the path we build, is

$$\begin{aligned} m\ell_0 &= m(\lfloor 2(1 - 2\sqrt{2\epsilon})(cn - 7m) \rfloor - 1) \\ &\geq m(2(1 - 2\sqrt{2\epsilon})(cn - 7m) - 2) \\ &\geq (2/3 + \lambda)M \cdot (1 - 2\sqrt{2\epsilon})((1 - \epsilon)n/M - 7m - 2), \end{aligned} \quad (2)$$

where the last inequality holds because $m = (2/3 + \lambda)M/2$, $2(1 - 2\sqrt{2\epsilon}) > 1$, and $c \geq (1 - \epsilon)/M$. Since $en/M > 7m + 2$ and $\lambda = 13\sqrt{\epsilon}$, (2) is bounded below by

$$(2/3 + \lambda) \cdot (1 - 2\sqrt{2\epsilon})(1 - 2\epsilon)n > 2n/3 \geq k - 1.$$

Hence, if all the paths T_i are of length exactly ℓ_0 , the length of the full path we build would be larger than $k - 1$. Since the lengths of the P_i 's are fixed, while the length of T_i can be any positive odd integer at least three and at most ℓ_0 , we will shorten some T_i to make the path be of length exactly $k - 1$. We greedily include $T_1, T_2, \dots$ such that each T_i is of length ℓ_0 , until the length of the concatenated path $T_1, P_1, T_2, P_2, \dots$ is at least $k - 1$ for the first time. If, when we stop, the total length is exactly $k - 1$, we will take all those T_i to have length ℓ_0 . Otherwise, when we stop, the total length is greater than $k - 1$. If, when we stop, the last path is T_j for some j , we will shorten the length of T_j by deleting the last few vertices from T_j ; if the last path is P_j for some j , we will shorten the length of P_j by deleting the last few vertices of P_j . In summary, there exists a properly chosen integer $m - 1 \geq m' \geq 0$ such that $T_1, \dots, T_{m'}$ are of length ℓ_0 and, after concatenating $T_1, P_1, T_2, P_2, \dots, T_{m'}, P_{m'}, T_{m'+1}$ with the length of $T_{m'+1}$ less than ℓ_0 or $T_1, P_1, T_2, P_2, \dots, T_{m'}, P_{m'}$ with a possibly shortened $P_{m'}$, we obtain a path of length $k - 1$. Using (1) to bound the number of T_i for $i \leq m'$, the total number of choices for $T_1, \dots, T_{m'}$ when fixing the anchor vertices $\{v_i, u_i\}$ for $1 \leq i \leq m'$ and P_i for $1 \leq i \leq m'$ is at least

$$\left(\frac{4\epsilon\pi}{e^2}\right)^{m'} \cdot \left((d' - 8\sqrt{2\epsilon})(cn - 7m)/(2e)\right)^{\ell_0 m' - m'}. \quad (3)$$

If the concatenated path of length $k - 1$ needs to end with a path $T_{m'+1}$ of length $1 \leq \ell' < \ell_0$, then $T_{m'+1}$ can be any path of length ℓ' alternating between $U'_{m'+1}$ and $V'_{m'+1}$ that starts with $v_{m'+1}$. By Lemma 12 with d replaced by $d' - \epsilon$, ϵ by 2ϵ , and n by $cn - 7m$, since $1 \leq \ell' < \ell_0$, the number of choices for $T_{m'+1}$ is at least

$$\begin{aligned} (d' - \epsilon - 2\epsilon - \sqrt{2\epsilon})^{\ell'} \prod_{i=1}^{\ell'} ((cn - 7m) - \lceil i/2 \rceil) &> \frac{2\pi}{e^2} \left((d' - 4\sqrt{2\epsilon})(cn - 7m - 1)/e\right)^{\ell'} \\ &> \frac{2\pi}{e^2} \left((d' - 4\sqrt{2\epsilon})(cn - 7m)/(2e)\right)^{\ell'}, \end{aligned}$$

where the first inequality is by the same estimate as in (1). Together with (3), the total number of k -vertex paths when fixing the anchor vertices $\{v_i, u_i\}$ for $1 \leq i \leq m'$ and P_i for $1 \leq i \leq m'$ is at least

$$\left(\frac{4\epsilon\pi}{e^2}\right)^{m'+1} \cdot \left((d' - 8\sqrt{2\epsilon})(cn - 7m)/(2e)\right)^{\ell_0 m' - m' + \ell'}. \quad (4)$$

Here we can assume $0 \leq \ell' < \ell_0$ to combine the two cases of whether the path of length $k - 1$ ends with $T_{m'+1}$ or not.

Since the total length of $P_1, \dots, P_{m'}$ is at most $6m'$, the total length of the T_i 's for $1 \leq i \leq m' + 1$, which is $\ell_0 m' + \ell'$, is at least $k - 1 - 6m'$. Thus, (4) is at least

$$\left(\frac{4\epsilon\pi}{e^2}\right)^{m'+1} \cdot \left((d' - 8\sqrt{2\epsilon})cn/(4e)\right)^{k-7(m'+1)} \geq \epsilon^m \left((d' - 8\sqrt{2\epsilon})cn/(4e)\right)^{k-7m}.$$

Since $m = (2/3 + \lambda)M/2$, which is a constant, there exists c' such that the expression above is at least $(c'k)^k$, completing the proof. $\square$

Theorem 2 for even cycles for colorings satisfying Case 1 of Lemma 11. Fix $0 < \epsilon \leq 10^{-20}$ and let $d = 20\sqrt{\epsilon}$ and $\lambda = 13\sqrt{\epsilon}$. The proof for even cycles is very similar to the previous proof for paths. Suppose

there are vertex subsets $U_1, \dots, U_m, V_1, \dots, V_m$ with $m = (2/3 + \lambda)M/2$ and $|V_i|, |U_i| \geq cn$ satisfying the properties of Case 1 of Lemma 11, say in color red. Let the edge density between U_i and V_i be at least $d' = d - \epsilon$. We may assume that n is sufficiently large in terms of c, ϵ , and M . We will show that there is a constant c' such that the number of monochromatic cycles with $k = \lceil 2(n+1)/3 \rceil$ vertices, with k even, is at least $(c'k)^k$.

To do this, we will find distinct vertices $v_i^1, v_i^2, w_i^1, w_i^2 \in V_i$ for $2 \leq i \leq m-1$ and $v_1^1, w_1^1 \in V_1$ and $v_m^2, w_m^2 \in V_m$ such that (all the indices are mod m and the edges considered are all in color red):

1. $v_i^1, v_i^2, w_i^1, w_i^2$ each have degree at least $(d' - \epsilon)|U_i|$ to U_i for all $1 \leq i \leq m$;
2. there is a path P_i connecting v_i^1 and v_{i+1}^2 and a path Q_i connecting w_i^1 and w_{i+1}^2 such that both P_i and Q_i have length at most six;
3. for each i , the lengths of P_i and Q_i have the same parity;
4. there is a path L_i connecting w_i^1 and w_i^2 of length four;
5. all of the paths P_i, Q_i , and L_i with $1 \leq i \leq m-1$ are vertex disjoint except where they share an end vertex.

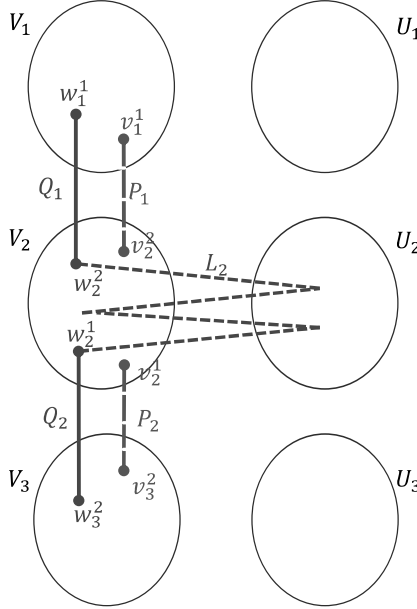


Figure 2: An illustration showing the paths P_i, Q_i , and L_i .

Suppose that we can find such vertices $v_i^1, v_i^2, w_i^1, w_i^2$ together with appropriate paths P_i, Q_i , and L_i . We now show that we are done in this case. First we remove all the internal vertices in these paths from U_i and V_i ; for $2 \leq i \leq m-1$, we also remove w_i^1, w_i^2 from V_i . This results in subsets U'_i of U_i and V'_i of V_i . By the length constraints on P_i, Q_i , and L_i in conditions 2 and 4 above, all of these paths have in total at most

$$5(m-1) + 5(m-1) + 3m < 13m \quad (5)$$

internal vertices. Thus, we have

$$|V'_i| \geq |V_i| - 13m - 2 \geq |V_i| - 15m \quad \text{and} \quad |U'_i| \geq |U_i| - 13m \geq |U_i| - 15m.$$

Let ℓ_0 be the largest even integer not larger than $\lfloor 2(1 - 2\sqrt{2\epsilon})(cn - 15m) \rfloor$. Then, in the bipartite graph $G(U'_1, V'_1)$, we will show that we can obtain many paths T_1 from w_1^1 to v_1^1 where T_1 has length $l_1 \leq \ell_0$. Clearly, T_1 is of even length since it alternates between U_1 and V_1 , eventually coming back to the side where it started. For $2 \leq i \leq m-1$, we show that we can find many paths T_i from v_i^2 to v_i^1 in $G(V'_i, U'_i)$ where T_i has even length $l_i \leq \ell_0$. In $G(U_m, V_m)$, we will find many paths T_m from v_m^2 to w_m^2 where T_m has even length $l_m \leq \ell_0$. Since $v_i^1, v_i^2, w_i^1, w_i^2$ each have large degree to U_i , they also have large degree to U'_i . As in the proof of the path case, we can use Lemma 13 applied to $G(V'_i, U'_i)$ to count the number of choices for the path T_i given l_i .

As in the previous proof, we can create a cycle by concatenating

$$T_1, P_1, T_2, P_2, \dots, T_{m-1}, P_{m-1}, T_m, Q_{m-1}, L_{m-1}, \dots, Q_2, L_2, Q_1.$$

Since the T_i and L_i are all of even length and the lengths of P_i and Q_i have the same parity by Condition 3, we obtain an even cycle.

The total length of the even cycle we build is the total length of P_i , Q_i , and L_i plus $\sum_{i=1}^m l_i$. If $l_i = 4$ for all $1 \leq i \leq m$, the total length is at most

$$6(m-1) + 6(m-1) + 4m + 4m < 20m < k.$$

On the other hand, if $l_i = \ell_0$ for all $1 \leq i \leq m$, when n is sufficiently large, the total length is at least

$$\begin{aligned} (m-1) + (m-1) + 4m + \ell_0 m &> (m-1) + (m-1) + 4m + (\lfloor 2(1 - 2\sqrt{2\epsilon})(cn - 15m) \rfloor - 1)m \\ &> \lfloor 2(1 - 2\sqrt{2\epsilon})(cn - 15m) \rfloor (2/3 + \lambda)M/2 \\ &\geq \lfloor 2(1 - 2\sqrt{2\epsilon})((1 - \epsilon)n/M - 15m) \rfloor (2/3 + \lambda)M/2 \\ &> (2/3 + \sqrt{\epsilon})n > k, \end{aligned}$$

where we used that $c \geq (1 - \epsilon)n/M$ and $\lambda = 13\sqrt{\epsilon}$. Therefore, we can reduce the lengths of some l_i , maintaining the condition that $4 \leq l_i \leq \ell_0$ are even integers for each i , to obtain a cycle of length exactly k .

Thus, the total number of even cycles of length k , having fixed the P_i , Q_i , and L_i , is at least the product of the number of choices for T_i for $1 \leq i \leq m$. Since the total length of the P_i , Q_i , and L_i is at most $6(m-1) + 6(m-1) + 4m < 16m$, the total length of the T_i , which is $\sum_{i=1}^m l_i$, is at least $k - 16m$. By a similar computation to (1) in the previous proof, the total number of even cycles of length k is therefore at least

$$\begin{aligned} \prod_{i=1}^m \left(\frac{4\epsilon\pi}{e^2} \right) (d' - 8\sqrt{2\epsilon})^{l_i-1} ((cn - 15m)/e)^{l_i-1} &= \left(\frac{4\epsilon\pi}{e^2} \right)^m (d' - 8\sqrt{2\epsilon})^{\sum_{i=1}^m l_i - m} ((cn - 15m)/e)^{\sum_{i=1}^m l_i - m} \\ &\geq \left(\frac{4\epsilon\pi}{e^2} \right)^m (d' - 8\sqrt{2\epsilon})^{k-17m} ((cn - 15m)/e)^{k-17m}, \end{aligned}$$

which is at least $(c'k)^k$ for some positive constant c' . It thus suffices to show that we can find vertices $v_i^1, v_i^2, w_i^1, w_i^2$ and paths P_i , Q_i , and L_i satisfying Conditions 1 to 5.

We will pick $v_i^1, v_i^2, w_i^1, w_i^2$ and L_i, P_i , and Q_i with the desired properties greedily. In step one, we pick four vertices $w_1^1, v_1^1, w_2^2, v_2^2$ and two paths P_1 and Q_1 . In each step $i \geq 2$ except the last one, we pick four vertices $w_i^1, v_i^1, w_{i+1}^2, v_{i+1}^2$ and three paths P_i, Q_i , and L_i . Suppose we have completed all steps $j < i$. We now need to pick $w_i^1, v_i^1 \in V_i$ and $w_{i+1}^2, v_{i+1}^2 \in V_{i+1}$.

Let distinct arbitrary vertices $v_1, v'_1, v''_1 \in V_i, v_2, v'_2, v''_2 \in V_{i+1}$ be such that v_1, v'_1, v''_1 each have degree at least $(d' - \epsilon)|U_i|$ to U_i and v_2, v'_2, v''_2 each have degree at least $(d' - \epsilon)|U_{i+1}|$ to U_{i+1} . Since (U_i, V_i) is ϵ -regular, there are at least $(1 - \epsilon)cn$ vertices in each of V_i and V_{i+1} that satisfy this degree condition from which v_1, v'_1, v''_1 and v_2, v'_2, v''_2 can be chosen. Since $\bigcup_{i=1}^m V_i \cup \bigcup_{i=1}^m U_i$ is $(200M, 6)$ well-connected in

red, the pigeonhole principle implies that there are at least $100M$ red internally-disjoint paths connecting v_1 and v_2 whose lengths are at most six and of the same parity. Label (v_1, v_2) as odd or even depending on the parity of the paths between them. We can similarly label (v'_1, v'_2) and (v''_1, v''_2) . By the pigeonhole principle again, at least two of the pairs (v_1, v_2) , (v'_1, v'_2) , and (v''_1, v''_2) have the same parity. Suppose (v_1, v_2) and (v'_1, v'_2) have the same parity, say odd. Then we let v_1 be v_i^1 , v_2 be v_{i+1}^2 , v'_1 be w_i^1 , and v'_2 be w_{i+1}^2 , noting that there are at least $100M$ internally vertex-disjoint paths connecting v_1 and v_2 of odd length at most 6 and the same for v'_1 and v'_2 . Therefore, we have at least $100M$ candidates for P_i and at least $100M$ candidates for Q_i . Since the previously chosen paths P_j, Q_j , and L_j use in total at most $7(m-1) + 7(m-1) + 5m < 19m$ vertices and $100M > 19m$, there are choices for P_i and Q_i with the desired properties.

It remains to choose L_i . We remove all the internal vertices in the previously chosen P_j, Q_j , and L_j from U_i and V_i and we also remove v_i^1 and v_i^2 from V_i . This results in U_i'' and V_i'' . By (5),

$$|V_i''| \geq |V_i| - 13m - 2 > |V_i| - 15m$$

and, similarly, $|U_i''| > |U_i| - 15m$. Furthermore, the pair (U_i'', V_i'') is 2ϵ -regular with density $d(U_i'', V_i'') \geq d' - \epsilon$.

As w_i^1 and w_i^2 are in V_i'' and each has degree at least $(d' - 2\epsilon)|U_i''|$ to U_i'' , Lemma 13 applied with n replaced by $cn - 15m$, ϵ by 2ϵ , d by $d' - \epsilon$, and l by 4 implies that there are at least

$$(d' - \epsilon - 7\sqrt{2\epsilon})^3 (2\epsilon)(cn - 15m)^2 (cn - 15m - 1) \geq \epsilon(d - 11\sqrt{\epsilon})^3 (cn)^3 / 4$$

paths of length 4 connecting w_i^1 and w_i^2 . Each vertex is in at most $3n^2$ paths with the prescribed end vertices, since n^2 is an upper bound on the number of choices for the other two internal vertices in this path and the multiplicative factor 3 indicates which of the three internal vertices our vertex is. Therefore, we have at least $\epsilon(d - 11\sqrt{\epsilon})^3 c^3 n / 12$ vertex-disjoint paths of length 4 connecting the two end vertices, each of which is a candidate for L_i , completing the proof. $\square$

2.3.2 Proof of Theorem 2 in the situation of Case 2 of Lemma 11

We begin by showing that Theorem 2 is true for paths for edge-colorings satisfying Case 2 of Lemma 11. The argument for even cycles will be almost the same. Throughout the proof, n will be assumed to be sufficiently large in terms of ϵ .

Theorem 2 for paths for colorings satisfying Case 2 of Lemma 11. Fix $0 < \epsilon \leq 10^{-20}$ and let $d = 20\sqrt{\epsilon}$ and $\lambda = 13\sqrt{\epsilon}$. We are given an extremal coloring with parameter $\alpha = 1000(d + \lambda + \sqrt{\epsilon})$, so we have a red/blue edge-coloring of a complete graph whose vertex set has a partition into subsets V and U with

$$|V| \geq (2/3 - \alpha)n, \quad |U| \geq (1/3 - \alpha)n, \quad |U| + |V| = n. \quad (6)$$

Furthermore, without loss of generality, we can assume that the red density within V is at least $1 - \alpha$ and the blue density between U and V is at least $1 - \alpha$. We want to prove that the number of monochromatic paths with k vertices for $k = \lceil 2(n+1)/3 \rceil$ is at least $(k/10)^k$.

We first perform a standard cleaning-up process, moving a few vertices between U and V , so that within V and between U and V certain degree conditions hold.

Claim 14 (Updated Extremal Coloring). *There is a partition $V' \cup U'$ satisfying the following conditions:*

- $|V'| \geq (2/3 - 3\alpha)n$, $|U'| \geq (1/3 - 2\alpha)n$.
- The red graph on V' has minimum degree at least $(2/3 - 4\alpha)|V'|$.
- The blue density between U' and V' is at least $1 - 8\alpha$.

- Each vertex in U' has blue degree to V' at least $(1/3 - 4\alpha)|V'|$.

Proof. We define V' to be the set of vertices which have red degree at least $2|V|/3$ in V and let U' be the complement of V' , noting that each vertex in U' has blue degree larger than $|V|/3 - 1$ to V . We claim that this partition has the desired properties.

We first show that most vertices of V are in V' . Suppose $|V \setminus V'| = x|V|$. Since the red density in V is at least $1 - \alpha$, we have $x \cdot 2|V|/3 + (1 - x)|V| \geq (1 - \alpha)|V|$ and, therefore, $x \leq 3\alpha$. Combining this inequality with (6), we conclude that

$$|V'| \geq (1 - 3\alpha)|V| \geq (1 - 3\alpha)(2/3 - \alpha)n > (2/3 - 3\alpha)n.$$

We next show that not many vertices in U were moved to V' . Suppose $|U \cap V'| = y|U|$. Since the red density between U and V is at most α , we have $y \cdot 2|V|/3 \leq \alpha|V|$ and so $y \leq 3\alpha/2$. This implies that $|V| + 3\alpha|U|/2 \geq |V'|$. Using that $\alpha \leq 10^{-5}$ and (6) gives

$$|V'| \leq |V| + 3\alpha/2 \cdot |V|(1/3 + \alpha)/(2/3 - \alpha) < (1 + \alpha)|V|. \quad (7)$$

We also have

$$|V'| \leq |V| + 3\alpha|U|/2 \leq (2/3 + \alpha)n + 3\alpha/2 \cdot (1/3 + \alpha)n \leq (2/3 + 2\alpha)n$$

and so $|U'| = n - |V'| \geq (1/3 - 2\alpha)n$.

Furthermore, each vertex in V' has red degree in V' at least

$$2|V|/3 - 3\alpha|V| = (2/3 - 3\alpha)|V| \geq (2/3 - 3\alpha)|V'|/(1 + \alpha) > (2/3 - 4\alpha)|V'|,$$

where the second to last inequality is by (7).

Similarly, together with (7), each vertex in U' has blue degree to V' at least

$$|V|/3 - 1 - 3\alpha|V| > \frac{1/3 - 3\alpha}{1 + \alpha}|V'| - 1 > (1/3 - 4\alpha)|V'|.$$

The blue density between U' and V' is $e(U', V')/|U'||V'|$, which is at least

$$\begin{aligned} \frac{(1 - \alpha)|V||U| - |V \setminus V'||U| - \frac{1}{3}|U \setminus U'||V|}{|U'||V'|} &= \frac{(1 - \alpha)|V||U| - |V \setminus V'||U| - \frac{1}{3}|U \setminus U'||V|}{(|U| + |V \setminus V'| - |U \setminus U'|)(|V| + |U \setminus U'| - |V \setminus V'|)} \\ &\geq \frac{(1 - \alpha)|V||U| - |V \setminus V'||U| - \frac{1}{3}|U \setminus U'||V|}{|V||U| + (|V| - |U|)|V \setminus V'| - (|V| - |U|)|U \setminus U'|}. \end{aligned}$$

We have already established that $|V \setminus V'| \leq 3\alpha|V|$ and $|U \setminus U'| = |U \cap V'| \leq 3\alpha|U|/2$. Hence, substituting for $|U \setminus U'|$ in the numerator of the last expression its maximum $3\alpha|U|/2$ and in the denominator zero decreases the fraction. Moreover, since α is sufficiently small, the last expression above is decreasing in $|V \setminus V'|$ and, therefore, minimized when $|V \setminus V'| = 3\alpha|V|$. Hence, the blue density is at least

$$\frac{(1 - 4.5\alpha)|V||U|}{|V||U| + (|V| - |U|)3\alpha|V|} > \frac{(1 - 4.5\alpha)|V||U|}{|V||U| + (\frac{2/3+\alpha}{1/3-\alpha}|U| - |U|)3\alpha|V|} > 1 - 8\alpha,$$

where we used that $\alpha \leq 10^{-5}$ is sufficiently small and $|V|/|U| \leq (2/3 + \alpha)/(1/3 - \alpha)$. $\square$

Abusing notation, we let V' be the new V and U' the new U and assume that they satisfy the properties described in Claim 14. We now wish to count the number of monochromatic paths with k vertices in this configuration. We have two cases.

Case A: Suppose $|V| \geq k = \lceil 2(n+1)/3 \rceil$. Let $V'' \subset V$ be an arbitrary subset with $|V''| = k$. The minimum red degree in V'' is at least

$$(2/3 - 4\alpha)|V| - (|V| - \lceil 2(n+1)/3 \rceil) > (2/3 - 4\alpha)|V| - 4\alpha n \geq \frac{2}{3}k - 8\alpha n \geq k/2.$$

Therefore, the red graph on V'' is a Dirac graph. By the main result of [10], the number of Hamiltonian cycles (and, hence, paths with k vertices) in the red graph on V'' is at least $k!/2^{k+o(k)}$.

Case B: Suppose $|V| < k$. In this case, $|U| \geq n - (k-1) \geq \lfloor k/2 \rfloor$. To complete the proof, we apply Lemma 15 below to the blue bipartite graph with parts U and V with $\beta = 8\alpha$ and $\delta = (1/3 - 4\alpha)|V|$. Since $|V| \geq (2/3 - 3\alpha)n \geq 3k/4$ and every vertex in U has blue degree to V at least $(1/3 - 4\alpha)|V| \geq 4\sqrt{8\alpha} \max(|V|, 2|U|)$, the conditions of Lemma 15 are satisfied. Thus, the number of monochromatic blue paths with k vertices is at least

$$0.9^{|U|} 2^{-k/2} 0.94^{k/2} \lfloor k/2 \rfloor! (3k/4)! / (k/4)! > (k/10)^k.$$

In either case, we get at least $(k/10)^k$ monochromatic paths with k vertices, as required. $\square$

For a complete bipartite graph with parts U and V , where $|V| \geq |U| \geq \lfloor k/2 \rfloor$, the number of paths with k vertices starting in V is precisely $(|V|)_{\lfloor k/2 \rfloor} (|U|)_{\lfloor k/2 \rfloor}$, where we use the standard falling factorial notation $(n)_k = n(n-1)\cdots(n-k+1)$. If the bipartite graph is not complete but just nearly complete, then, provided $|V|$ is much larger than $|U|$ and U satisfies an appropriate minimum degree condition, we can prove that there are still almost this many paths with k vertices between U and V . This can be thought of as a counting version of a special case of the blow-up lemma [26].

Lemma 15. *Let k be a sufficiently large positive integer and G a bipartite graph with parts U and V such that $|V| \geq 3k/4$ and $|U| \geq \lfloor k/2 \rfloor$, the edge density between U and V is at least $1 - \beta$ with $\beta < 10^{-4}$, and every vertex in U has degree at least $\delta \geq 4\sqrt{\beta} \max(|V|, 2|U|)$. Then the number of paths with k vertices in G starting from a vertex in V is at least*

$$\left(\frac{\delta}{4|V|} \right)^{2\sqrt{\beta}|U|} \left(1 - \frac{4\sqrt{\beta}|U|}{\delta} \right)^{k/2} (1 - 6\sqrt{\beta})^{k/2} (|U|)_{\lfloor k/2 \rfloor} (|V|)_{\lfloor k/2 \rfloor}.$$

Proof. Let U_0 be the set of vertices in U that have degree at most $(1 - \sqrt{\beta})|V|$ and $U_1 = U \setminus U_0$. The number of edges in G satisfies $(1 - \beta)|V||U| \leq e(G) \leq (1 - \sqrt{\beta})|V||U_0| + |V|(|U| - |U_0|)$, from which we obtain $|U_0| \leq \sqrt{\beta}|U|$.

We will show that there are many paths with k vertices in G alternating between V and U that start from a vertex in V . We do this by first showing that there are many sequences $L = u_1, \dots, u_{\lfloor k/2 \rfloor}$ of $\lfloor k/2 \rfloor$ distinct vertices in U that extend to many paths with k vertices in $G(U, V)$, where extending here means that we can find vertices $v_1, \dots, v_{\lfloor k/2 \rfloor}$ such that $v_1, u_1, v_2, u_2, \dots$ is a path with k vertices, where the last vertex of the path is $u_{k/2}$ if k is even and $v_{\lfloor k/2 \rfloor}$ if k is odd.

We will require that the sequences L satisfy the following property:

(P) If $u_i \in L$ is in U_0 , then $i \leq \delta/2$ and i is odd.

We first bound the number of choices for L . Note that if $i \leq \delta/2$ and even or $\lfloor k/2 \rfloor \geq i > \delta/2$, then u_i must be in U_1 . There are in total $\ell := \lfloor k/2 \rfloor - \lceil \delta/4 \rceil$ such terms. Thus, we have $(|U_1|)_{\ell}$ choices for these terms in the sequence L . For the remaining $\lceil \delta/4 \rceil$ terms, we can choose any of the remaining vertices from U , so we get $(|U| - \ell)_{\lceil \delta/4 \rceil}$ possible choices to complete the sequence, giving a total of

$$\begin{aligned}
(|U_1|)_\ell(|U| - \ell)_{\lceil \delta/4 \rceil} &\geq \left(\frac{|U_1| - \ell + 1}{|U| - \ell + 1} \right)^\ell (|U|)_\ell(|U| - \ell)_{\lceil \delta/4 \rceil} = \left(\frac{|U| - |U_0| - \ell + 1}{|U| - \ell + 1} \right)^\ell (|U|)_{\lfloor k/2 \rfloor} \\
&\geq \left(\frac{\lfloor k/2 \rfloor - |U_0| - \ell + 1}{\lfloor k/2 \rfloor - \ell + 1} \right)^\ell (|U|)_{\lfloor k/2 \rfloor} \geq \left(\frac{\delta/4 - |U_0|}{\delta/4} \right)^\ell (|U|)_{\lfloor k/2 \rfloor} \\
&\geq \left(1 - 4\sqrt{\beta}|U|/\delta \right)^{k/2} (|U|)_{\lfloor k/2 \rfloor}
\end{aligned}$$

possible sequences L .

Having picked L , we greedily choose $v_1, \dots, v_{\lfloor k/2 \rfloor}$ to complete the path. We can pick v_1 to be any neighbor of u_1 , so there are at least δ choices if $u_1 \in U_0$ and at least $(1 - \sqrt{\beta})|V|$ choices if $u_1 \in U_1$. Having already picked out $v_1, \dots, v_{j-1}$, we next show how to pick v_j . Note that (aside from the case¹ where k is odd and $j = \lceil k/2 \rceil$), this amounts to picking a common neighbor of u_{j-1} and u_j different from $v_1, \dots, v_{j-1}$. Notice that, by property (P), no two consecutive terms of L are in U_0 . We thus have two cases to consider.

In the first case, one of u_{j-1} and u_j is in U_0 . In this case, we have $j - 1 \leq \delta/2$ by property (P). The degree of the vertex from U_0 is at least δ and the degree of the other vertex, which is in U_1 , is at least $(1 - \sqrt{\beta})|V|$, so u_{j-1} and u_j have at least $\delta + (1 - \sqrt{\beta})|V| - |V| = \delta - \sqrt{\beta}|V|$ common neighbors. Hence, there are at least $\delta - \sqrt{\beta}|V| - (j - 1) \geq \delta/2 - \sqrt{\beta}|V| \geq \delta/4$ common neighbors of u_{j-1} and u_j not among $v_1, \dots, v_{j-1}$. Any of these at least $\delta/4$ vertices can be chosen for v_j .

In the second case, both u_{j-1} and u_j are in U_1 . Then u_{j-1} and u_j have at least $2(1 - \sqrt{\beta})|V| - |V| = (1 - 2\sqrt{\beta})|V|$ common neighbors, so there are at least $(1 - 2\sqrt{\beta})|V| - (j - 1) > (1 - 6\sqrt{\beta})(|V| - (j - 1))$ choices for v_j , where we used $|V| \geq \frac{3}{4}k$ and $j - 1 < k/2$.

As there are at most $|U_0|$ terms in the sequence L that belong to U_0 , there are at most $2|U_0|$ consecutive pairs in L that include a term from U_0 . Therefore, the first case happens at most $2|U_0|$ times. Observe that $\delta/4 \leq |V|/4 \leq (1 - 6\sqrt{\beta})|V|/3 \leq (1 - 6\sqrt{\beta})(|V| - (j - 1))$. From the estimates above, we therefore see that the number of ways of greedily choosing $v_1, \dots, v_{\lfloor k/2 \rfloor}$ is at least

$$(\delta/4)^{2|U_0|} (1 - 6\sqrt{\beta})^{\lfloor k/2 \rfloor - 2|U_0|} (|V|)_{\lfloor k/2 \rfloor} / |V|^{2|U_0|} \geq (\delta/(4|V|))^{2\sqrt{\beta}|U|} (1 - 6\sqrt{\beta})^{k/2} (|V|)_{\lfloor k/2 \rfloor}.$$

Hence, by counting the number of choices for L and then the number of ways of completing any given choice of L to a path with k vertices, we find that the number of paths with k vertices in G starting from a vertex in V is at least the desired bound. $\square$

We now briefly discuss how to modify the argument above to prove Theorem 2 for even cycles for colorings satisfying Case 2 of Lemma 11. Case A is identical, since we actually counted red cycles of length k in the proof. Case B is almost identical, in that we only need to modify the proof of Lemma 15 so that the conclusion guarantees many even cycles of length k (instead of just paths with k vertices) in the bipartite graph G . This amounts to also guaranteeing that v_1 is a neighbor of $u_{k/2}$ and only changes the bound slightly, so that, as in the case of paths, the number of blue cycles of length k is at least $(k/10)^k$ for k sufficiently large.

3 Proof of Lemma 11

Throughout this subsection $0 < \epsilon \leq 10^{-10}$ and $d, \lambda \geq 1000\epsilon$, as in the assumptions of Lemma 11.

¹When k is odd and $j = \lceil k/2 \rceil$, we are instead picking a neighbor of $u_{(k-1)/2}$ not among $v_1, \dots, v_{j-1}$, for which there are at least $(1 - \sqrt{\beta})|V| - (j - 1) > (1 - 3\sqrt{\beta})(|V| - (j - 1))$ choices, where we used $|V| \geq \frac{3}{4}k$ and $j = (k - 1)/2$.

3.1 Preparation

In order to prove Lemma 11, we first collect some auxiliary results, beginning with a lemma of Gyárfás, Sárközy, and Szemerédi [18] about finding a well-connected subset in any red/blue-multicolored K_n .

Lemma 16 (Lemma 4.1 in [18]). *For every positive integer t and red/blue-multicolored K_n , there exist $W \subset V(K_n)$ and a color, say red, such that $|W| \geq n - 28t$ and W is $(t, 3)$ -well-connected in the red subgraph of K_n .*

We will also use two lemmas concerning extremal numbers of matchings. The first is a classical result of Erdős and Gallai [12], which is easily seen to be tight by considering either the graph which consists of a clique on $2k + 1$ vertices and a collection of isolated vertices or the graph in which the only edges are those incident to at least one of k vertices.

Lemma 17 (Erdős and Gallai [12]). *For integers k and n with $0 \leq k \leq n/2$, if the maximum matching in an n -vertex graph G has size k , then G has at most $\max\left(\binom{2k+1}{2}, \binom{k}{2} + (n-k)k\right)$ edges.*

The second lemma about matchings that we will need is the following simple consequence of König's theorem, which says that the covering and matching numbers of a bipartite graph are equal. It is easily seen to be tight by considering the bipartite graph with k vertices in one part complete to the other part, which has n vertices, and no other edges.

Lemma 18. *If a bipartite graph has at most n vertices in each part and does not contain a matching of size larger than k , then it has at most kn edges.*

The final lemma in this section says that if a reduced graph satisfies certain properties, then the original graph it describes is well-connected.

Lemma 19. *Let H be a graph with vertex set $[h]$ in which each vertex has distance at most three from vertex 1. Let G be an h -partite graph with parts $V_1, \dots, V_h$, each of order at least N . Suppose that $0 < \alpha < 1/10$, $3\alpha < d < 1$, and, for every edge $(i, j) \in E(H)$, the pair (V_i, V_j) is α -regular in G with $d(V_i, V_j) \geq d$. Let $T \leq (d - 3\alpha)N/5$ be a positive integer. Then, for each $i \in [h]$, there is $V'_i \subset V_i$ such that the following hold:*

1. $|V'_i| \geq (1 - \alpha)|V_i|$.
2. For every edge $(i, j) \in E(H)$, the pair (V'_i, V'_j) is 2α -regular with density $d(V'_i, V'_j) \geq d - \alpha$.
3. $\bigcup_{i=1}^h V'_i$ is $(T, 6)$ -well-connected in G .

Proof. For $i \in [h]$, let $s(i)$ denote the distance of i from 1. We are given $s(i) \leq 3$ for all $i \in [h]$. For $i \in [2, h]$, let $n(i)$ be an arbitrary neighbor of i with $s(n(i)) = s(i) - 1$ and let $n(1)$ be an arbitrary neighbor of vertex 1 in H . We call $n(i)$ the *successor* of vertex i . Let D be the directed graph on $[h]$ in which each vertex i has outdegree one with $n(i)$ as its outneighbor.

For each $i \in [h]$, let V'_i be the set of vertices in V_i whose degree to $V_{n(i)}$ is at least $(d - \alpha)|V_{n(i)}|$. By Lemma 6(i), $|V'_i| \geq (1 - \alpha)|V_i|$. For each $(i, j) \in E(H)$, we have $d(V'_i, V'_j) \geq d - \alpha$ as (V_i, V_j) is α -regular and $1 - \alpha > \alpha$. Moreover, by Lemma 6(ii), since $\max(2\alpha, \alpha/(1 - \alpha)) = 2\alpha$, (V'_i, V'_j) is 2α -regular. It only remains to check Item 3 of the lemma, that is, to show that, for any two vertices $u, v \in \bigcup_{i=1}^h V'_i$, we can find T internally-disjoint paths of length at most six connecting them.

For each $i \in [h]$, there is a unique directed path P_i in D from i to 1. This path has length $s(i) \leq 3$ and the next vertex of the path is the successor of the current vertex. For each pair (a, b) of not necessarily distinct vertices of H , let W_{ab} be a walk in H from a to b formed by concatenating a walk from a to 1 in D of length two or three and a walk from b to 1 in D of length two or three. Such a walk of length two

or three from a to 1 in D is either P_a itself or formed by adding to P_a a walk of length two from 1 to its successor and back. We can similarly construct a walk of length two or three from b to 1 in D .

Let $u \in V'_a$ and $v \in V'_b$, noting that a and b may not be distinct. Let $a = a_0, \dots, a_s = b$ with $s \leq 6$ denote the vertices of the walk W_{ab} from a to b in order and let r be the length of the walk from a to 1 that makes up the first part of W_{ab} , so that $r = 2$ or 3 and $a_r = 1$. In particular, a_j is the successor of a_{j-1} for $1 \leq j \leq r$ and a_{j-1} is the successor of a_j for $r < j \leq s$.

We greedily construct T internally vertex-disjoint paths from u to v of length at most s . Each such path has at most five internal vertices. In particular, after pulling out the internal vertices of $t < T$ such paths, all but at most $5t$ of the vertices in each part remain. The remaining subset U_i of V'_{a_i} for $1 \leq i \leq s-1$ has size $|U_i| \geq |V'_{a_i}| - 5t \geq (1 - \alpha)|V_{a_i}| - 5t$. We next build a walk $u = u_0, u_1, \dots, u_s = v$ from u to v of length s with $u_i \in U_i$ for each i . If this walk is a path, it is the desired next path from u to v . Otherwise, we get the desired path by deleting some internal vertices from the walk.

The vertex u has at least $(d - \alpha)|V_{a_1}|$ neighbors in V_{a_1} , so u has at least $(d - \alpha)|V_{a_1}| - (|V_{a_1}| - |U_{a_1}|) \geq (d - 2\alpha)|V_{a_1}| - 5t$ neighbors in U_{a_1} . These neighbors are all potential choices for u_1 . As $(d - 2\alpha)|V_{a_1}| - 5t \geq \alpha|V_{a_1}|$, $d > \alpha$, the pair (V_{a_1}, V_{a_2}) is α -regular, and $U_{a_2} \subset V_{a_2}$, Lemma 6(i) implies that all but at most $\alpha|V_{a_2}|$ vertices in U_{a_2} have a common neighbor with u in U_{a_1} and thus can be chosen for u_2 . If $r = 3$, we similarly get that all but $\alpha|V_{a_3}|$ vertices in U_{a_3} can be chosen for u_3 . In either case, we get that all but at most $\alpha|V_{a_r}|$ vertices in U_r can be chosen for u_r when starting the walk from u . Similarly, working backwards from v , we get that all but at most $\alpha|V_{a_r}|$ vertices in U_r can be chosen for u_r when starting the walk from v . As $|U_r| \geq (1 - \alpha)|V_{a_r}| - 5t > 2\alpha|V_{a_r}|$, there is a vertex in U_r that can be chosen for u_r to complete the walk from u to v . Hence, we can continue the process of pulling out T internally vertex-disjoint paths from u to v , completing the proof. $\square$

3.2 Proof of Lemma 11

Consider a red/blue edge-coloring of K_n . Let M_0 be as in Szemerédi's regularity lemma, Lemma 7, with $\epsilon/2$ in place of ϵ and $m_0 = 1000/\epsilon$. Apply Lemma 16 to this edge-coloring of K_n with $t = 200M_0$ to obtain a vertex subset W with $|W| \geq n - 5600M_0$ and a color such that W is $(200M_0, 3)$ -well-connected in that color. By applying the regularity lemma, Lemma 7, with $\epsilon/2$ in place of ϵ and $m_0 = 1000/\epsilon$ to the induced edge-coloring on W , we obtain the following lemma.

Lemma 20. *For every $0 < \epsilon, d \leq 1/2$, there are positive integers M_0 and n_0 such that the following holds. For every red/blue edge-coloring of K_n with $n \geq n_0$, there is a positive integer $1000/\epsilon \leq M \leq M_0$, a vertex subset $W \subset V(K_n)$ with $|W| \geq n - 5600M_0$ which is $(200M_0, 3)$ -well-connected in either the red or the blue subgraph, and an equitable partition $W = V_1 \cup \dots \cup V_M$ such that the red/blue-multicolored reduced graph H with vertex set $[M]$ and parameters $\epsilon/2$ and d has at most $\epsilon \binom{M}{2}/2$ non-adjacent pairs.*

For the rest of this section, we fix an edge-coloring of K_n with colors red and blue and the set W , the sets $V_1, \dots, V_M$ in the equitable partition of W , and the reduced graph H guaranteed by Lemma 20. We will also assume without loss of generality that W is $(200M_0, 3)$ -well-connected in red. Let H_b be the spanning blue subgraph of H and H_r the spanning red subgraph, noting that the same edge can be in both H_b and H_r . As n is sufficiently large, for each $i \in M$, we have $|V_i| \geq \lfloor |W|/M \rfloor \geq \lfloor (n - 5600M_0)/M \rfloor \geq (1 - \epsilon/4)n/M$.

Lemma 21. *If the red subgraph H_r contains a matching with at least $(2/3 + \lambda)M$ vertices, then the conditions of Case 1 in Lemma 11 are satisfied.*

Proof. For each edge (a_i, b_i) of such a matching in H_r , consider the corresponding sets of vertices V_{a_i} and V_{b_i} . By construction, the union of these sets of vertices is a subset of a $(200M, 3)$ -well-connected set in red. Furthermore, each of the parts has size at least $(1 - \epsilon/4)n/M$ and each pair of parts corresponding to an edge of the matching in H_r is ϵ -regular of density at least d . Hence, all the conditions of Case 1 in Lemma 11 are indeed satisfied. $\square$

Lemma 22. *If the blue subgraph H_b contains a matching with at least $(2/3 + \lambda)M$ vertices and there is a vertex v such that each vertex in the matching has distance at most three from v in H_b , then the conditions of Case 1 in Lemma 11 are satisfied.*

Proof. Let S be the set of vertices with distance at most 3 from v . Then, by assumption, S contains the vertices of the blue matching. Furthermore, by Lemma 20, for every $(i, j) \in E(H_b[S])$, (V_i, V_j) is $\epsilon/2$ -regular with blue density $d(V_i, V_j) \geq d$. Thus, we can apply Lemma 19 with H being $H_b[S]$, G being the blue $|S|$ -partite graph induced on $\bigcup_{i \in S} V_i$, vertex 1 being v , $\alpha = \epsilon/2$, $T = 200M_0$, and $N = (1 - \epsilon/4)n/M$. As n is sufficiently large, the conditions of Lemma 19 are satisfied. Hence, for each $i \in S$, there is $V'_i \subset V_i$ such that $|V'_i| \geq (1 - \epsilon/2)|V_i| \geq (1 - 3\epsilon/4)n/M$, for every edge (i, j) of H_b the pair (V'_i, V'_j) is ϵ -regular in H_b with density $d(V'_i, V'_j) \geq d - \epsilon/2$, and $\bigcup_{i \in S} V'_i$ is $(200M, 6)$ -well-connected. Since the vertices of the matching in H_b are all in S , the conditions of Case 1 in Lemma 11 are satisfied. $\square$

For the rest of the section, we may therefore suppose that the largest matching in H_r has $2m < (2/3 + \lambda)M$ vertices and no subgraph of H_b with radius at most three contains a matching with $(2/3 + \lambda)M$ vertices. We will conclude that the given coloring of K_n must be an extremal coloring with parameter $\beta := 1000(d + \lambda + \sqrt{\epsilon})$, which will complete the proof of Lemma 11.

Consider a maximum matching in H_r with m edges (a_i, b_i) for $1 \leq i \leq m$, so $m < (1/3 + \lambda/2)M$. Let $A = \{a_i : i \in [m]\}$, $B = \{b_i : i \in [m]\}$, and $C = [M] \setminus (A \cup B)$, so A , B , and C form a partition of $[M]$ with $|A| = |B| = m$ and $|C| = M - 2m > (1/3 - \lambda)M$. We may assume without loss of generality that the red degree of b_i to C is at least the red degree of a_i to C . Observe that C contains no red edge as otherwise we could add it to the already constructed red matching, contradicting the fact that the chosen matching is maximum in H_r . Moreover, each a_i has red degree to C at most one, as otherwise there are red edges (a_i, c_1) and (b_i, c_2) with $c_1, c_2 \in C$ distinct and we could replace the edge (a_i, b_i) in the matching by the two edges (a_i, c_1) and (b_i, c_2) , making a larger matching in H_r and again contradicting that the red matching is of maximum size. For the rest of the proof, we fix vertex subsets A , B , and C with the properties described above.

We prove several claims along the way to establishing that the coloring is an extremal coloring with parameter β . In outline, we will first show that the parts A , B , and C each have roughly equal size by showing that m is close to $M/3$ (this will follow from the upper bound on m already given above and the lower bound on m given in Claim 1 below). We will then deduce that we have an extremal coloring by showing that either the edges in $\bigcup_{i \in A \cup C} V_i$ are almost all blue and the edges from this set to $\bigcup_{i \in B} V_i$ are almost all red or the edges in $\bigcup_{i \in A \cup B} V_i$ are almost all red and the edges from this set to $\bigcup_{i \in C} V_i$ are almost all blue. Note that this will be sufficient as W contains almost all vertices of K_n and the regularity partition of W is equitable.

Recall that all edges of H with both vertices in C are blue. Let $v \in C$ be a vertex of largest blue degree in C and $C' \subset C$ be the neighbors of v in C in the graph H_b . As there are at most $\epsilon \binom{M}{2}/2$ non-adjacent pairs in H and, hence, in the induced subgraph on C , by averaging, the vertex v is in at most $\epsilon \binom{M}{2}/|C| < \epsilon \binom{M}{2}/((1/3 - \lambda)M) < 2\epsilon M - 1$ non-adjacent pairs. Hence, $|C'| \geq |C| - 1 - (2\epsilon M - 1) = |C| - 2\epsilon M$.

Let m_1 be the size of a maximum blue matching between A and C' . Pick a blue matching between A and C' of size m_1 whose vertices consist of subsets $A_1 \subset A$ and $C_1 \subset C'$ subject to the condition that $C' \setminus C_1$ contains a blue matching as large as possible. Let $C_2 \subset C' \setminus C_1$ be the vertices of this blue matching, so the vertices in $A_1 \cup C_1 \cup C_2$ are all in C' or adjacent in blue to a vertex in C' and therefore have distance in H_b at most two from v and are all in a blue matching with $2|A_1| + |C_2|$ vertices. By construction, $|A_1| = |C_1| = m_1$. Let $C_3 = C' \setminus (C_1 \cup C_2)$, so $C' = C_1 \cup C_2 \cup C_3$ forms a partition of C' into three parts. Figure 3 illustrates the different sets.

By the choice of the maximum blue matching, there are no blue edges from $A \setminus A_1$ to $C' \setminus C_1$. Moreover, for each edge (a, c) of the maximum blue matching between A and C' (so $a \in A_1$ and $c \in C_1$), a has no blue edges to $C' \setminus C_1$ or c has no blue edges to $A \setminus A_1$. On the other hand, each vertex in A has at most one red edge to C , C contains no red edges, and there are at most $\frac{\epsilon}{2} \binom{M}{2}$ non-adjacent pairs. Comparing

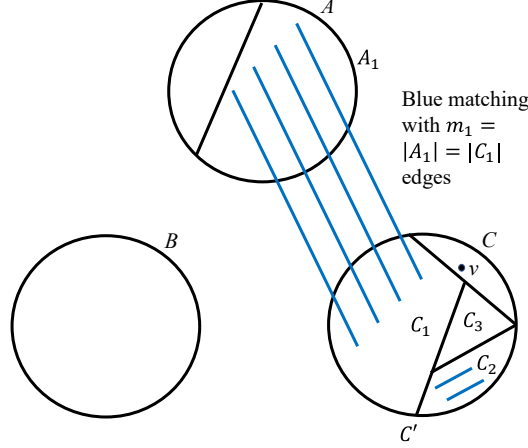


Figure 3: An illustration showing A , B , C , A_1 , C_1 , C_2 , and C_3 . C' is the set of blue neighbors of the vertex v and C_1 , C_2 , and C_3 form a partition of C' .

these upper and lower bounds on the number of non-blue pairs in $A \cup C'$ with not both vertices in A , we obtain

$$|A \setminus A_1||C' \setminus C_1| + |A_1| \min(|A \setminus A_1|, |C' \setminus C_1|) \leq \frac{\epsilon}{2} \binom{M}{2} + |A| \leq \frac{2}{7} \epsilon M^2, \quad (8)$$

where the last inequality is from $|A| \leq M$ and $M \geq m_0 = 1000/\epsilon$.

For each edge $(a, c) \in A_1 \times C_1$ in the maximum blue matching between A and C' , either a or c has blue degree to C_3 at most one, since otherwise we can replace (a, c) by two blue edges (a, c_1) and (c, c_2) with $c_1, c_2 \in C_3$, which would also give a maximum blue matching between A and C' , but would increase the size of the maximum blue matching in the remaining vertices in C' , contradicting our choice of the blue matching between A and C' . Hence, there are at least $|A_1|(|C_3| - 1)$ pairs between $A_1 \cup C_1$ and C_3 which are not blue. Moreover, for each matching edge (c_3, c_4) in the maximum blue matching in $C' \setminus C_1$ (so $c_3, c_4 \in C_2$), either c_3 or c_4 has blue degree at most one to C_3 , so there are at least $(|C_2|/2)(|C_3| - 1)$ pairs between C_2 and C_3 which are not blue. Finally, there are no blue edges in C_3 . Hence, comparing the upper and lower bounds on the number of non-blue pairs between $A_1 \cup C'$ and C_3 , we similarly obtain

$$(|A_1| + |C_2|/2 + |C_3|/2)(|C_3| - 1) \leq \frac{\epsilon}{2} \binom{M}{2} + |A_1| \leq \frac{2}{7} \epsilon M^2. \quad (9)$$

Claim 1. $m \geq (\frac{1}{3} - \lambda - 4\epsilon) M$.

Proof. Suppose, for the sake of contradiction, that $m < (\frac{1}{3} - \lambda - 4\epsilon) M$. Then $|C'| \geq |C| - 2\epsilon M = M - 2m - 2\epsilon M \geq m = |A|$. In particular, $\min(|A \setminus A_1|, |C' \setminus C_1|) = |A \setminus A_1|$. Since also $|A_1| = |C_1| = m_1$, the left-hand side of (8) simplifies to $|C'|(m - m_1)$. Hence, (8) implies that $|C'|(m - m_1) \leq \epsilon M^2/3$. As $|C'| \geq M - 2m - 2\epsilon M \geq M/3$, we obtain $m_1 \geq m - \epsilon M$.

As $|A_1| = |C_1| = m_1$ and $|C'| = |C_1| + |C_2| + |C_3|$, we have

$$|A_1| + |C_2|/2 + |C_3|/2 = (|C'| + m_1)/2 \geq (M - m - 3\epsilon M)/2 \geq M/3.$$

Hence, from (9), we similarly obtain $|C_3| \leq \epsilon M$.

Thus, the number of vertices of the blue matching of distance at most two from v is

$$\begin{aligned} 2|A_1| + |C_2| &= 2m_1 + |C'| - m_1 - |C_3| \geq m_1 + M - 2m - 2\epsilon M - \epsilon M \\ &\geq M - m - 4\epsilon M > (2/3 + \lambda)M, \end{aligned}$$

contradicting the assumption that no such large blue matching exists. $\square$

If $|C' \setminus C_1| \geq |A \setminus A_1|$, the left-hand side of (8) is equal to $|C'|(m - m_1)$. Otherwise, $|A \setminus A_1| > |C' \setminus C_1|$ and the left-hand side of (8) is equal to $m(|C'| - m_1)$. In either case, as $m, |C'| \geq \frac{2}{7}M$, we obtain from (8) that

$$|A_1| = m_1 \geq \min(m, |C'|) - \epsilon M. \quad (10)$$

Consider a maximum blue matching between A_1 and B . Let m_2 be the number of edges of this blue matching and let $A_2 \subset A_1$ and $B_2 \subset B$ be the vertices in this blue matching. Consider a blue matching between $A_1 \setminus A_2$ and C_1 that matches every vertex in $A_1 \setminus A_2$ subject to the condition that the vertices in C' not contained in this blue matching contain a blue matching of maximum possible size. Note that such a blue matching between $A_1 \setminus A_2$ and C_1 exists as there is a perfect matching between A_1 and C_1 by construction. Let C'_1 be the set of $|A_1 \setminus A_2|$ vertices in C_1 that match with a vertex in $A_1 \setminus A_2$ and $C_4 \subset C' \setminus C'_1$ consist of the vertices in the maximum blue matching in $C' \setminus C'_1$. Let $C_5 = C' \setminus (C'_1 \cup C_4)$. Figure 4 is an illustration of these sets.

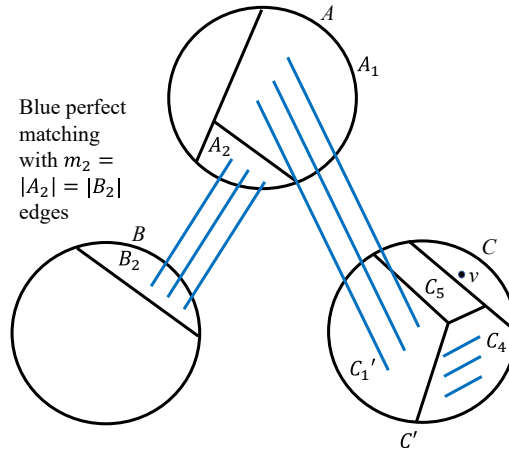


Figure 4: An illustration showing $A_2 \subset A_1$, B_2 , $C'_1 \subset C_1$, C_4 , and C_5 . The sets of blue edges represent three distinct blue matchings.

There are clearly no blue edges in C_5 . As before, for each edge (a, c) in the blue matching between $A_1 \setminus A_2$ and C'_1 , either a or c has blue degree at most one to C_5 . Similarly, for each edge (c_1, c_2) of the blue perfect matching in C_4 , either c_1 or c_2 has blue degree at most one to C_5 . Hence, the number of pairs between $(A_1 \setminus A_2) \cup C'$ and C_5 which are not blue is at least $(|C'_1| + |C_4|/2 + |C_5|/2)(|C_5| - 1)$ and at most $\frac{\epsilon}{2} \binom{M}{2} + |A_1 \setminus A_2| \leq \epsilon M^2/4 + M$ (recall that C has no red edges and each vertex in A has at most one red neighbor in C). As

$$|C'_1| + |C_4|/2 + |C_5|/2 = |C'|/2 + |C'_1|/2 \geq |C'|/2 \geq M/7,$$

we obtain that $|C_5| \leq \frac{7}{4}\epsilon M + 8 \leq 2\epsilon M$.

We have obtained a blue matching with vertex set $B_2 \cup A_1 \cup C'_1 \cup C_4$. Each vertex in this blue matching has distance in H_b at most three from v and, therefore, the number of vertices in this blue matching is less than $(2/3 + \lambda)M$. On the other hand, the number of vertices in this blue matching is at least

$$\begin{aligned} |B_2| + m_1 + |C'| - |C_5| &\geq |B_2| + \min(m, |C'|) + |C'| - 3\epsilon M \geq |B_2| + \min(m, |C|) + |C| - 7\epsilon M \\ &> |B_2| + \left(\frac{2}{3} - 2\lambda - 7\epsilon\right)M, \end{aligned}$$

where the first inequality uses (10) and the last inequality uses $2m < (2/3 + \lambda)M$ and $|C| + m = M - m$ when $m \leq |C|$ and uses $|C| > (1/3 - \lambda)M$ when $m > |C|$. We thus have the following claim.

Claim 2. $|B_2| \leq (3\lambda + 7\epsilon)M$.

In particular, as there are no blue edges between $B \setminus B_2$ and $A_1 \setminus A_2$, the graph between A and B is almost entirely red. Let $\mu := (6\lambda + 16\epsilon)M$.

Claim 3. The largest red matching with vertices in A has size less than μ or the largest red matching between B and C' has size less than 2μ .

Proof. Suppose, for the sake of contradiction, that the claim does not hold. Consider a red matching in A of size μ and let A_3 be the set of vertices of this red matching, so $|A_3| = 2\mu$. Consider a red matching between B and C' of size 2μ and let B_3 be the vertices of this matching in B . Figure 5 illustrates these sets.

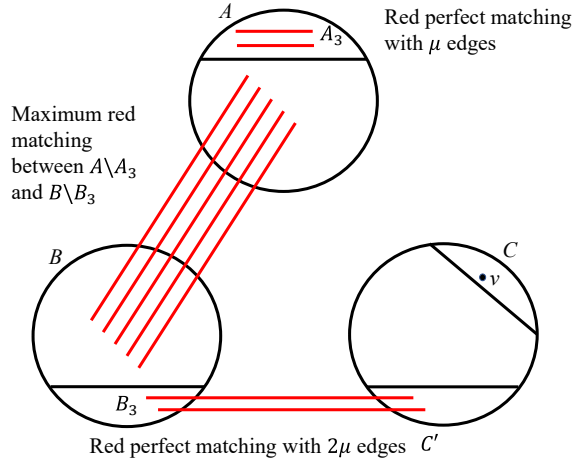


Figure 5: An illustration showing A_3 and B_3 .

Observe that $|A \setminus A_3| = |B \setminus B_3| = m - 2\mu$. Edge partition the balanced complete bipartite graph between $A \setminus A_3$ and $B \setminus B_3$ into $m - 2\mu$ perfect matchings, so each of the perfect matchings has exactly $m - 2\mu$ edges. The number of missing edges in H is at most $\frac{\epsilon}{2} \binom{M}{2} < \epsilon M^2/4$ and $m - 2\mu \geq M/4$, so the density of non-adjacent pairs between $A \setminus A_3$ and $B \setminus B_3$ in H is at most 4ϵ and there is a matching $\mathcal{M}$ in H between $A \setminus A_3$ and $B \setminus B_3$ with at least $(1 - 4\epsilon)(m - 2\mu)$ edges by Lemma 18. By Claim 2, the maximum size of a blue matching between A_1 and B is at most $(3\lambda + 7\epsilon)M$, so the maximum size of a blue matching between A and B is at most

$$(3\lambda + 7\epsilon)M + |A| - |A_1| \leq (4.5\lambda + 10\epsilon)M,$$

where the last inequality uses (10) (so that $|A| - |A_1| = m - m_1 \leq m - \min(m, |C'|) + \epsilon M$), $|C'| \geq |C| - 2\epsilon M = M - 2m - 2\epsilon M$, and $m < (1/3 + \lambda/2)M$. Hence, the matching $\mathcal{M}$ has at least $(1 - 4\epsilon)(m - 2\mu) - (4.5\lambda + 10\epsilon)M$ red edges. This red matching, together with the red matching of size μ with vertex set A_3 and the red matching between B_3 and C_3 of size 2μ , forms a red matching of size at least

$$\begin{aligned} (1 - 4\epsilon)(m - 2\mu) - (4.5\lambda + 10\epsilon)M + 3\mu &\geq m - (4.5\lambda + 12\epsilon)M + \mu \\ &\geq (1/3 - 5.5\lambda - 16\epsilon)M + \mu \\ &\geq (2/3 + \lambda)M/2, \end{aligned}$$

where the first inequality uses $m \leq (2/3 + \lambda)M/2 \leq M/2$, the second inequality follows from Claim 1, and the final inequality from the definition of μ . This contradicts the assumption that H_r has no red matching of size $(2/3 + \lambda)M/2$. $\square$

From Claim 3, the rest of the proof naturally splits into two cases.

Case 1: The largest red matching with vertices in A has size less than μ .

Let $\tau := 2\mu + 2\lambda M + 3\epsilon M + 1$. Suppose, for the sake of contradiction, that there is a blue matching between C' and B of size τ . Let $C_6 \subset C'$ be the τ vertices of C' in this blue matching. Figure 6 is an illustration.

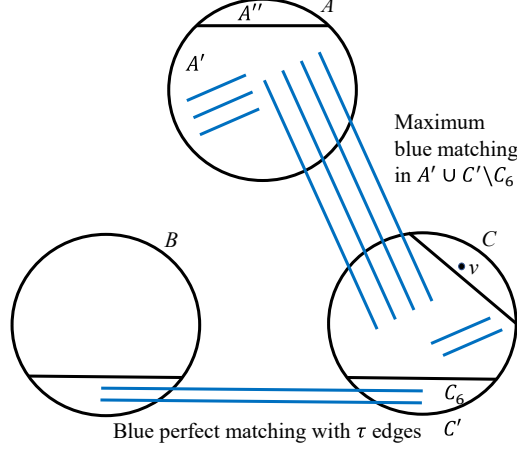


Figure 6: An illustration showing A' , A'' , and C_6 .

Let $A' \subset A$ be those vertices with at least one blue neighbor in C' and $A'' = A \setminus A'$. Each vertex in A has red degree to C' at most one and so each vertex in A'' is in at most one edge to C' . As there are at most $\epsilon M^2/4$ non-adjacent pairs in total and $|C'| \geq M/4 + 1$, we thus have $|A''|(|C'| - 1) \leq \epsilon M^2/4$ and hence $|A''| \leq \epsilon M$.

We next bound the number of pairs of vertices in $A' \cup C' \setminus C_6$ which are not blue edges. There are at most $\frac{\epsilon}{2} \binom{M}{2}$ non-adjacent pairs. There are also no red edges in C' . By Lemma 17, there are fewer than $\binom{\mu}{2} + \mu(|A'| - \mu)$ red edges in A' . Finally, as each vertex in A has red degree at most one to C , there are at most $|A'|$ red edges between A' and C' . In total, the number of pairs of vertices in $A' \cup C' \setminus C_6$ which are not blue is at most

$$\begin{aligned}
 \frac{\epsilon}{2} \binom{M}{2} + \binom{\mu}{2} + \mu(|A'| - \mu) + |A'| &\leq \mu|A'| - \mu^2/2 + |A'| + \epsilon M^2/4 \\
 &\leq 2\mu|A'| \\
 &\leq 2\mu(|A'| + |C'| - \tau - \mu - 1/2) \\
 &= 2\mu(|A'| + |C' \setminus C_6| - \mu - 1/2) \\
 &= \binom{|A'| + |C' \setminus C_6|}{2} - \binom{|A'| + |C' \setminus C_6| - 2\mu}{2}.
 \end{aligned}$$

Here we used $\binom{x}{2} - \binom{x-y}{2} = y(x - y/2 - 1/2)$ with $x = |A'| + |C' \setminus C_6|$ and $y = 2\mu$. Therefore, there are at least $\binom{|A'| + |C' \setminus C_6| - 2\mu}{2}$ blue edges with both vertices in $A' \cup C' \setminus C_6$. By Lemma 17, there is a blue matching in $A' \cup C' \setminus C_6$ spanning at least $|A'| + |C' \setminus C_6| - 2\mu - 1$ vertices. The blue matching consisting of the blue matching between C_6 and B of size τ together with the maximum blue matching in $A' \cup C' \setminus C_6$ contains only vertices of distance at most two from v and has at least

$$\begin{aligned}
 2\tau + |A'| + |C' \setminus C_6| - 2\mu - 1 &= \tau + |A'| + |C'| - 2\mu - 1 \geq \tau + |A| + |C| - 3\epsilon M - 2\mu - 1 \\
 &= \tau + M - m - 3\epsilon M - 2\mu - 1 = 2\lambda M + M - m \geq (2/3 + \lambda)M
 \end{aligned}$$

vertices, contradicting our assumption that there is no such blue matching. Hence, there is no blue matching between C' and B of size τ .

We now show that the coloring is an extremal coloring with parameter $\beta = 1000(d + \lambda + \sqrt{\epsilon})$ with the set $\bigcup_{i \in A \cup C} V_i$ almost entirely blue and the bipartite graph between $\bigcup_{i \in A \cup C} V_i$ and $\bigcup_{i \in B} V_i$ almost entirely red. We first check that the two parts of this partition are of the claimed size. It suffices to check that $|\bigcup_{i \in B} V_i|/n$ is $\frac{1}{3} \pm \beta$. Since $W \subset V(K_n)$ satisfies $|W|/n \geq 1 - \epsilon/4$ and the partition $W = \bigcup_{i \in [M]} V_i$ is equitable with $n \gg M$, it is enough to show that $|B|/M$ is $\frac{1}{3} \pm \frac{\beta}{2}$. But this follows easily, since $|B| = m$ and $\frac{1}{3} - \lambda - 4\epsilon < m/M \leq \frac{1}{3} + \lambda/2$, where the first inequality is by Claim 1.

We next show that the induced subgraph on $\bigcup_{i \in A \cup C} V_i$ is almost entirely blue. The number of non-edges in $A \cup C$ is at most $\frac{\epsilon}{2} \binom{M}{2}$. There are also no red edges in C . The number of red edges from A to C is at most $|A|$. By Lemma 17, the number of red edges in A is at most $\binom{\mu}{2} + \mu(|A| - \mu)$. Each part V_i in the equitable partition has size at most $\lceil n/M \rceil$ and the density between any $\epsilon/2$ -regular pair of parts that does not correspond to a red edge in H is at most d . Hence, the total number of edges with both vertices in $\bigcup_{i \in A \cup C} V_i$ which are red is at most

$$\left(\frac{\epsilon}{2} \binom{M}{2} + |A| + \binom{\mu}{2} + \mu(|A| - \mu) \right) \lceil n/M \rceil^2 + d \binom{n}{2} \leq \left(d + \frac{\epsilon}{2} + \frac{\mu}{M} \right) \binom{n}{2} \leq \beta \binom{|\bigcup_{i \in A \cup C} V_i|}{2},$$

where we used that $|\bigcup_{i \in A \cup C} V_i| \geq 3n/5$ and $\beta \geq 4d + 2\epsilon + 4\mu/M$.

Finally, we show that the bipartite graph between $\bigcup_{i \in A \cup C} V_i$ and $\bigcup_{i \in B} V_i$ is almost entirely red. We first bound the number of edges between $A \cup C$ and B which are not red. The number of missing edges between $A \cup C$ and B is at most $\frac{\epsilon}{2} \binom{M}{2}$. Recall that the maximum blue matching from A_1 to B is of size $|B_2|$. Moreover, the maximum blue matching between C' and B is of size less than τ and hence, by Lemma 18, the number of blue edges between $A \cup C$ and B is less than

$$\begin{aligned} & (|B_2| + |A \setminus A_1| + |C \setminus C'|) |B| + \max(|C'|, |B|) \tau \\ & \leq ((3\lambda M + 7\epsilon M) + (m - \min(m, |C'|) + \epsilon M) + 2\epsilon M) m + \tau M \\ & = (3\lambda + 10\epsilon) M m + (m - \min(m, |C'|)) m + \tau M \\ & \leq (3\lambda + 12\epsilon) M m + (m - \min(m, |C|)) m + \tau M \\ & \leq (4.5\lambda + 12\epsilon) M m + \tau M \\ & < (20\lambda + 50\epsilon) M^2, \end{aligned}$$

where we used Claim 2 and (10) in the first inequality, $|C'| \geq |C| - 2\epsilon M$ in the second inequality, $m - \min(m, |C|) \leq 1.5\lambda M$ in the third inequality, and, in the last inequality, we substituted in the values of τ and μ and used the lower bound on M . Hence, the number of blue edges between $\bigcup_{i \in A \cup C} V_i$ and $\bigcup_{i \in B} V_i$ is at most

$$\left(\frac{\epsilon}{2} \binom{M}{2} + (20\lambda + 50\epsilon) M^2 \right) \lceil n/M \rceil^2 + d \binom{n}{2} \leq (d + 20\lambda + 51\epsilon) n^2 \leq \beta \left| \bigcup_{i \in A \cup C} V_i \right| \cdot \left| \bigcup_{i \in B} V_i \right|,$$

where we used that $3n/5 \leq |\bigcup_{i \in A \cup C} V_i| \leq 7n/10$ and also that $\beta \geq 6(d + 20\lambda + 51\epsilon)$. This completes the proof in this case.

Case 2: The largest red matching between B and C' has size less than 2μ .

Our goal is to show that the coloring is an extremal coloring with parameter $\beta = 1000(d + \lambda + \sqrt{\epsilon})$ with at most a β -fraction of the edges in $\bigcup_{i \in A \cup B} V_i$ blue and at most a β -fraction of the edges between $\bigcup_{i \in A \cup B} V_i$ and $\bigcup_{i \in C} V_i$ red. To show these two parts have the desired size, it suffices to show that $\bigcup_{i \in C} V_i$ has size $(1/3 \pm \beta)n$ and this follows from a very similar computation to that in Case 1.

Let B_4 be the set of vertices in B that have at least one blue neighbor in C' , so every vertex in B_4 has distance at most two from v in blue. Since there is no red matching between $B \setminus B_4$ and C' of size 2μ ,

Lemma 18 implies that the number of red edges between $B \setminus B_4$ and C' is at most $2\mu \max(|C'|, |B \setminus B_4|)$. However, since the edges between $B \setminus B_4$ and C' are all red and there are at most $\frac{\epsilon}{2} \binom{M}{2}$ non-edges, the number of red edges between $B \setminus B_4$ and C' is at least $|B \setminus B_4||C'| - \frac{\epsilon}{2} \binom{M}{2}$. Since $|C'| \geq |C| - 2\epsilon M \geq (1/3 - \lambda - 2\epsilon)M$, it follows that $|B \setminus B_4| \leq 3\mu$. We get a blue matching with vertices of distance at most two from v by taking a maximum blue matching between A and C' (which is of size m_1) together with a maximum blue matching in B_4 , which is of size m_4 , say. Together this matching has size $m_1 + m_4$ and so, by assumption, we have $m_1 + m_4 < (2/3 + \lambda)M/2$. Hence, together with (10),

$$m_4 < (2/3 + \lambda)M/2 - m_1 \leq (2/3 + \lambda)M/2 - \min(m, |C'|) + \epsilon M \leq \frac{3}{2}\lambda M + 5\epsilon M,$$

where we used that $m \geq (\frac{1}{3} - \lambda - 4\epsilon)M$ by Claim 1 and $|C'| \geq |C| - 2\epsilon M > (\frac{1}{3} - \lambda - 2\epsilon)M$. By Lemma 17, with $k = \frac{3}{2}\lambda M + 5\epsilon M$, the number of blue edges in B_4 is thus at most $|B_4|k \leq km$. The number of pairs of vertices in B that are not in B_4 is at most $|B \setminus B_4||B| \leq 3\mu m$. Hence, there are at most $(\frac{3}{2}\lambda M + 5\epsilon M + 3\mu)m \leq 4\mu m$ blue edges in B .

We next bound the number of blue edges in A . We first claim that there is a matching (which does not have to be monochromatic) in $B_4 \cup C'$ with each edge containing at most one vertex in B_4 and with at least $|C'| + \min(|B_4|, |C'|) - 5\sqrt{\epsilon}M$ vertices. Indeed, if $|B_4| \geq |C'|$, then, since the number of edges between B_4 and C' is at least $|B_4||C'| - \frac{\epsilon}{2} \binom{M}{2}$, Lemma 18 gives a matching of size $|C'| - \epsilon M$, better than desired. On the other hand, if $|B_4| < |C'|$, Lemma 18 instead implies that there is a matching of size $|B_4| - \epsilon M$ between B_4 and C' . To complete the matching, we consider the set of at least $|C'| - |B_4|$ remaining vertices of C' and show that if q is the size of the maximum matching on this set and $2q < |C'| - |B_4| - 3\sqrt{\epsilon}M$, then $\max(\binom{2q+1}{2}, \binom{q}{2}) + (|C'| - |B_4| - q)q < \binom{|C'| - |B_4|}{2} - \epsilon \binom{M}{2}$, which would contradict Lemma 17. We may clearly assume that $|C'| - |B_4| \geq 3\sqrt{\epsilon}M$. But then

$$\begin{aligned} \binom{2q+1}{2} &= (2q+1)q < (|C'| - |B_4| - 3\sqrt{\epsilon}M + 1)(|C'| - |B_4| - 3\sqrt{\epsilon}M)/2 \\ &\leq \binom{|C'| - |B_4|}{2} - (6\sqrt{\epsilon}M - 2)(|C'| - |B_4|)/2 + 9\epsilon M^2/2 \\ &\leq \binom{|C'| - |B_4|}{2} - (6\sqrt{\epsilon}M - 2)(3\sqrt{\epsilon}M)/2 + 9\epsilon M^2/2 < \binom{|C'| - |B_4|}{2} - \epsilon \binom{M}{2} \end{aligned}$$

and

$$\begin{aligned} \binom{q}{2} + (|C'| - |B_4| - q)q &\leq q(q/2 + (|C'| - |B_4| - q)) = q(|C'| - |B_4| - q/2) \\ &< (|C'| - |B_4| - 3\sqrt{\epsilon}M)(|C'| - |B_4|)/2 \\ &= \binom{|C'| - |B_4|}{2} - (3\sqrt{\epsilon}M - 1)(|C'| - |B_4|)/2 \\ &\leq \binom{|C'| - |B_4|}{2} - (3\sqrt{\epsilon}M - 1)(3\sqrt{\epsilon}M)/2 < \binom{|C'| - |B_4|}{2} - \epsilon \binom{M}{2}, \end{aligned}$$

as required.

As there are no red edges in C' and the largest red matching between B_4 and C' has size less than 2μ , there is a blue matching in $B_4 \cup C'$ with at least $|C'| + \min(|B_4|, |C'|) - 5\sqrt{\epsilon}M - 4\mu$ vertices. If now A_1 contains a blue matching of size at least $4\mu + 3\sqrt{\epsilon}M$, then, together with the blue matching in $B_4 \cup C'$, we get a blue matching, each vertex of distance at most two from v , with the total number of vertices at least

$$8\mu + 6\sqrt{\epsilon}M + |C'| + \min(|B_4|, |C'|) - 5\sqrt{\epsilon}M - 4\mu \geq (2/3 + \lambda)M,$$

a contradiction. Note that in the inequality we used that $|C'| \geq |C| - 2\epsilon M \geq (1/3 - \lambda - 2\epsilon)M$ and, by Claim 1, that $|B_4| \geq |B| - 3\mu \geq (1/3 - \lambda - 4\epsilon)M - 3\mu$. Hence, A_1 does not contain a blue matching of

size $4\mu + 3\sqrt{\epsilon}M$. By Lemma 17, it follows that A_1 has at most $4\mu m + 3\sqrt{\epsilon}Mm$ blue edges. The number of pairs in A not in A_1 is also at most $|A \setminus A_1||A| \leq \mu m$, where we used (10) to obtain that

$$\begin{aligned} |A \setminus A_1| &\leq m - \min(m, |C'|) + \epsilon M \leq \max(0, m - (1/3 - \lambda - 2\epsilon)M) + \epsilon M \\ &< (1/3 + \lambda/2)M - (1/3 - \lambda - 2\epsilon)M + \epsilon M < \mu. \end{aligned}$$

Hence, there are at most $5\mu m + 3\sqrt{\epsilon}Mm$ blue edges in A .

We next bound the number of blue edges between A and B . First suppose, for the sake of contradiction, that there is a blue matching between A_1 and B_4 of size $2\lambda M$. Let the set of remaining vertices in A_1 be A'_1 and the set of remaining vertices in B_4 be B_5 . Pick a maximum blue matching in the union of C' and $A'_1 \cup B_5$. We claim that this second matching has size at least $|C'| - 2\epsilon M$. See Figure 7 for an illustration.

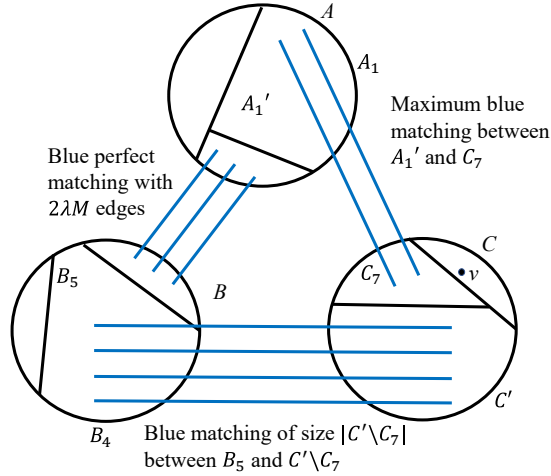


Figure 7: An illustration showing A'_1 , B_5 , and C_7 and the corresponding blue matchings.

To see this, we first build a blue matching between C' and B_5 . Since there are at most $\frac{\epsilon}{2} \binom{M}{2}$ non-edges in total and there is no red matching between B and C' of size 2μ , Lemma 18 implies that the number of blue edges between B_5 and C' is at least $|B_5||C'| - 2\mu \max(|B_5|, |C'|) - \frac{\epsilon}{2} \binom{M}{2}$. By Lemma 18 again, there is a blue matching between B_5 and C' of size at least

$$\frac{|B_5||C'| - 2\mu \max(|B_5|, |C'|) - \frac{\epsilon}{2} \binom{M}{2}}{\max(|B_5|, |C'|)} - 1 \geq \min(|B_5|, |C'|) - 2\mu - \epsilon M.$$

Let C_7 be the remaining vertices of C' that are not in this matching, noting that $|C_7|$ is significantly smaller than $|A_1|$ and so also significantly smaller than $|A'_1| = |A_1| - 2\lambda M$. Since each vertex in A has red degree at most one to C , Lemma 18 implies that there is a blue matching between A'_1 and C_7 of size at least

$$\frac{|A'_1||C_7| - |A'_1| - \frac{\epsilon}{2} \binom{M}{2}}{\max(|A'_1|, |C_7|)} - 1 \geq |C_7| - 1 - \frac{\epsilon}{2|A'_1|} \binom{M}{2} - 1 \geq |C_7| - 2 - \epsilon M.$$

Thus, we have a matching of size at least $|C' \setminus C_7| + |C_7| - 2 - \epsilon M > |C'| - 2\epsilon M$, as required. Together with the matching of size $2\lambda M$ between $A_1 \setminus A'_1$ and B_4 , we see that we have a blue matching with at least

$$4\lambda M + 2|C'| - 4\epsilon M \geq (2/3 + \lambda)M$$

vertices. But these vertices are all of distance at most two from v , a contradiction. Hence, there is no blue matching of size $2\lambda M$ between A_1 and B_4 and Lemma 18 implies that there are in total at most

$2\lambda M \max(|A_1|, |B_4|) \leq \lambda M^2$ blue edges between these two sets. As $|A \setminus A_1| \leq m - \min(m, |C'|) + \epsilon M$ by (10) and $|B \setminus B_4| \leq 3\mu$, we see that the number of blue edges between A and B is at most

$$\begin{aligned} & \lambda M^2 + |A \setminus A_1||B| + |B \setminus B_4||A| \\ & \leq \lambda M^2 + (\max(0, m - |C'|) + \epsilon M)M + 3\mu M/2 \\ & \leq \lambda M^2 + (\max(0, (1/3 + \lambda/2)M - (|C| - 2\epsilon M)) + \epsilon M)M + 1.5\mu M \\ & \leq \lambda M^2 + ((1/3 + \lambda/2)M - (1/3 - \lambda - 2\epsilon)M + \epsilon M)M + 1.5\mu M \\ & \leq (1.5\mu + 2.5\lambda M + 3\epsilon M)M < 2\mu M. \end{aligned}$$

In total, the number of blue edges in $A \cup B$ is at most $11\mu m + 3\sqrt{\epsilon}Mm$. Hence, the number of blue edges in $\bigcup_{i \in A \cup B} V_i$ is at most

$$\left(\frac{\epsilon}{2} \binom{M}{2} + 11\mu m + 3\sqrt{\epsilon}Mm \right) \lceil n/M \rceil^2 + d \binom{n}{2} \leq \left(d + \frac{11\mu}{M} + 4\sqrt{\epsilon} \right) \binom{n}{2} \leq \beta \binom{|\bigcup_{i \in A \cup B} V_i|}{2},$$

where we used that $m \leq M/2$, $|\bigcup_{i \in A \cup B} V_i| \geq 3n/5$ and $\beta \geq 3(d + 11\mu/M + 4\sqrt{\epsilon})$.

Since the largest red matching between B and C' has size less than 2μ and $|C \setminus C'| \leq 2\epsilon M$, it follows that the number of red edges between B and C is at most μM . Indeed, by Lemma 18, the number of red edges between B and C' is at most $2\mu \max(|B|, |C'|)$. Thus, the number of red edges between B and C is at most

$$\begin{aligned} 2\mu \max(|B|, |C'|) + |C \setminus C'||B| & \leq 2\mu \max(|B|, |C'|) + 2\epsilon M|B| \leq 2\mu \max(m, M - 2m) + 2\epsilon M|B| \\ & \leq 2\mu \max((1/3 + \lambda/2)M, (1/3 + 2\lambda + 8\epsilon)M) + 2\epsilon M^2 < \mu M. \end{aligned}$$

Moreover, every vertex in A has red degree at most one to C , so there are at most $|A|$ red edges between A and C . In total, we get that the number of red edges between $\bigcup_{i \in A \cup B} V_i$ and $\bigcup_{i \in C} V_i$ is at most

$$\left(\frac{\epsilon}{2} \binom{M}{2} + \mu M + |A| \right) \lceil n/M \rceil^2 + d \binom{n}{2} \leq \left(d + \frac{2\mu}{M} \right) n^2 \leq \beta \left| \bigcup_{i \in A \cup C} V_i \right| \cdot \left| \bigcup_{i \in B} V_i \right|.$$

This completes the proof in this case. $\square$

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