

Geometric Characterization of the H -property for Step-graphons

Mohamed-Ali Belabbas¹ and Xudong Chen^{2,*}

Abstract—In a recent paper [1], we have exhibited a set of conditions that are necessary for the H -property to hold for the class of step-graphons. In this paper, we prove that these conditions are essentially sufficient.

I. INTRODUCTION AND MAIN RESULT

In [1], we introduced the so-called H -property for a graphon W — roughly speaking, it is the property that a graph G sampled from W admits a Hamiltonian decomposition asymptotically (as the order of G goes to infinity) almost surely (*a.a.s.*). In [1], we have exhibited a set of conditions that were necessary for the H -property to hold for the class of step-graphons. We show in this paper that these conditions are also essentially (in a sense made precise below) sufficient and, moreover, that the H -property is a “zero-one” property.

A. Motivation

The line of research addressed in this sequence of papers is rooted in *structural system theory*, and investigates structural properties under random graph models described by graphons.

Structural system theory deals with the problem of understanding when a given network topology can sustain a prescribed system property. Typical such properties are controllability [2], [3] and stability [4], [5]. In more detail, consider a network of n agents x_1, \dots, x_n , whose communication topology is described by a directed graph (or simply digraph) $G = (V, E)$, with the nodes v_1, \dots, v_n representing the agents and directed edges $v_i v_j$ indicating the information flow (with the convention that a directed edge $v_i v_j$ indicates that agent x_j can access state information from x_i). Given the digraph G , a system dynamics $\dot{x}(t) = f(x(t))$ is said to be *compatible* with G if the dynamics of $x_i(t)$ depend only on its incoming neighbors in G :

$$v_j v_i \notin E \implies \frac{\partial f_i(x)}{\partial x_j} = 0,$$

where f_i describes the dynamics for agent x_i . Denote by Σ_G the set of differentiable dynamics compatible with G .

¹M.-A. Belabbas is with the Department of Electrical and Computer Engineering and the Coordinated Science Laboratory, University of Illinois, Urbana-Champaign. Email: belabbas@illinois.edu.

²X. Chen is with the Department of Electrical and Systems Engineering, Washington University in St. Louis. Email: cxudong@wustl.edu.

*Corresponding author: X. Chen. The two authors contributed equally to the manuscript in all categories.

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Given a desired system property \mathcal{S} , we say that the digraph G *sustains* \mathcal{S} if there exists a dynamics $f \in \Sigma_G$ satisfying the property \mathcal{S} .

Following our previous work [1], we focus on the existence of Hamiltonian decompositions in the digraph describing the communication topology of a network system. A Hamiltonian decomposition is a set of disjoint cycles that cover all nodes of the digraph. Their existence underlies a number of important properties pertinent to structural system theory, such as structural controllability [3] and structural stability [4], [6]. For example, it was shown in [6] that for G a *symmetric* digraph, the set Σ_G contains *exponentially* stable dynamical systems *if and only if* G contains a self-loop and a Hamiltonian decomposition. Hence, Hamiltonian decomposition are the key enabler of exponential stability for network dynamics. Hamiltonian decompositions also play an essential role in structural controllability for continuum ensembles of linear control systems [3].

Understanding the behavior of systems properties over random network topologies provides a wealth of insights [7]. For example, given a null random graph model, the probability that a network structure can sustain a desired system property \mathcal{S} tells us whether the given property is rare or frequent amongst topologies. The knowledge of whether there is an abundance or scarcity of network systems displaying the property \mathcal{S} is thus a critical component in the decision of a network manager to deploy expensive network systems operating in uncertain and/or adversarial environment (and since the H -property will be shown to be a zero-one property, abundance or scarcity can be understood as almost all or almost none under the graphon model). Furthermore, in cases such as social networks, having a random model for the topology is actually the natural option [8], as any estimate of a social network graph is bound to be affected by random (graph-valued) noise. In selecting a random graph model, the relevant aspects are: (1) universality and flexibility of the model (does the model cover a broad range of generic scenarios?) and (2) analytical and computational tractability. Graphons, which have emerged in the past decade as a powerful tool to understand large graphs [9], provide a fair amount of modeling flexibility while being tractable.

B. On the H -property

We start by recalling the definitions of a graphon and its sampling procedure. A *graphon* is a symmetric, measurable function $W : [0, 1]^2 \rightarrow [0, 1]$. Step-graphons, along with their partitions, are defined below:

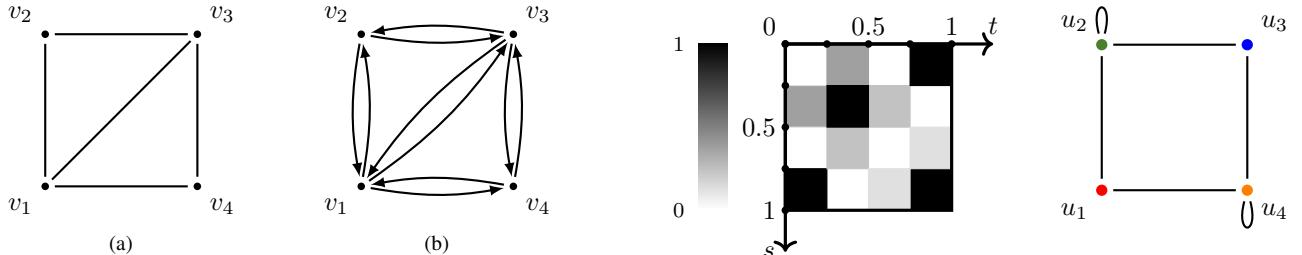


Fig. 1: (a): An undirected graph G on 4 nodes. (b): The directed graph \tilde{G} obtained from G by replacing every undirected with two oppositely oriented edges. The cycle $D_1 = v_1v_2v_3v_4v_1$ forms a Hamiltonian decomposition of \tilde{G} . The two node disjoint cycles $D_2 = v_1v_2v_1$ and $D_3 = v_3v_4v_3$ also form a Hamiltonian decomposition of \tilde{G} .

Definition 1 (Step-graphon and its partition). A graphon W is a **step-graphon** if there exists an increasing sequence $0 = \sigma_0 < \sigma_1 < \dots < \sigma_q = 1$ such that W is constant over each rectangle $[\sigma_i, \sigma_{i+1}) \times [\sigma_j, \sigma_{j+1})$ for all $0 \leq i, j \leq q-1$. We call $\sigma = (\sigma_0, \sigma_1, \dots, \sigma_q)$ a **partition** for W .

Graphons can be used to sample undirected graphs. Other uses of graphons in system theory as limits of adjacency matrices can be found in [10]–[12]. In this paper, we denote by $G_n \sim W$ graphs G_n on n nodes sampled from a graphon W . The sampling procedure was introduced in [9], [13] and is reproduced below: Let $\text{Uni}[0, 1]$ be the uniform distribution on $[0, 1]$. Given a graphon W , a graph $G_n = (V, E) \sim W$ on n nodes is obtained as follows:

- 1) Sample $y_1, \dots, y_n \sim \text{Uni}[0, 1]$ independently. We call y_i the *coordinate* of node $v_i \in V$.
- 2) For any two distinct nodes v_i and v_j , place an edge $(v_i, v_j) \in E$ with probability $W(y_i, y_j)$.

It should be clear that if $0 \leq p \leq 1$ is a *constant* and $W(s, t) = p$ for all $(s, t) \in [0, 1]^2$, then $G_n \sim W$ is an Erdős-Rényi random graph with parameter p . Consequently, graphons can be seen as a way to introduce *inhomogeneity* in the edge densities between different pairs of nodes.

Let W be a graphon and $G_n \sim W$. In the sequel, we use the notation $\tilde{G}_n = (V, \tilde{E})$ to denote the *directed* version of G_n , defined by the edge set

$$\tilde{E} := \{v_i v_j, v_j v_i \mid (v_i, v_j) \in E\}. \quad (1)$$

In words, we replace an undirected edge (v_i, v_j) with two directed edges $v_i v_j$ and $v_j v_i$. The directed graph \tilde{G}_n is said to have a *Hamiltonian decomposition* if it contains a subgraph H , with the same node set of \tilde{G}_n , such that H is a node disjoint union of directed cycles. See Fig. 1 for illustration.

We now have the following definition:

Definition 2 (H -property). Let W be a graphon and $G_n \sim W$. Then, W has the **H -property** if \tilde{G}_n has a Hamiltonian decomposition a.a.s., i.e.,

$$\lim_{n \rightarrow \infty} \mathbb{P}(\tilde{G}_n \text{ has a Hamiltonian decomposition}) = 1. \quad (2)$$

We will see below that the H -property is essentially a “zero-one” property in a sense that the probability on the left hand side of (2) converges to either 0 or 1.

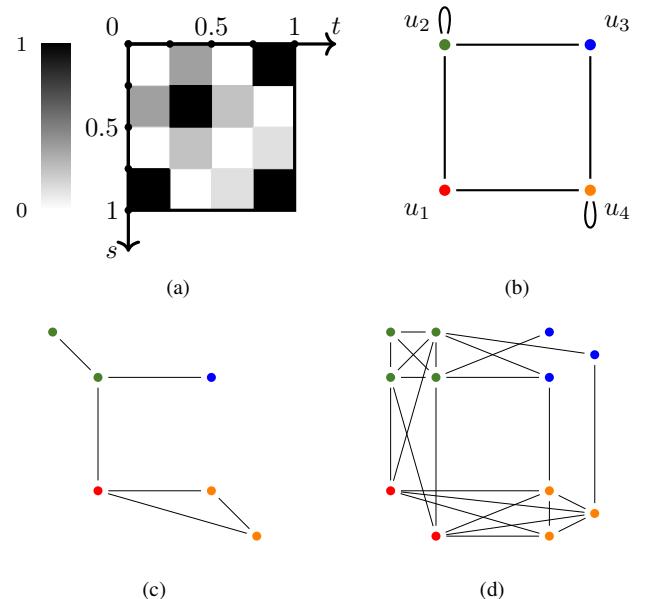


Fig. 2: (a): A step-graphon W with the partition $\sigma = (0, 0.25, 0.5, 0.75, 1)$, with the value coded by the gray scale on the left. (b): The associated skeleton graph S . (c) and (d): Graphs G_n sampled from W for $n = 6$ and $n = 12$, respectively.

C. Key objects

We present three key objects associated with a step-graphon, namely, its concentration vector, its skeleton graph, and its associated edge polytope, all of which were introduced in [1].

Definition 3 (Concentration vector). Let W be a step-graphon with partition $\sigma = (\sigma_0, \dots, \sigma_q)$. The associated **concentration vector** $x^* = (x_1^*, \dots, x_q^*)$ has entries defined as follows: $x_i^* := \sigma_i - \sigma_{i-1}$, for all $i = 1, \dots, q$.

It should be clear from the sampling procedure above that the concentration vector describes the proportion of sampled nodes in each interval $[\sigma_i, \sigma_{i+1})$ on average.

Given a step-graphon, its support can be described by a graph, which we call *skeleton graph*:

Definition 4 (Skeleton graph). To a step-graphon W with a partition $\sigma = (\sigma_0, \dots, \sigma_q)$, we assign the undirected graph $S = (U, F)$ on q nodes, with $U = \{u_1, \dots, u_q\}$ and edge set F defined as follows: there is an edge between u_i and u_j if and only if W is non-zero over $[\sigma_{i-1}, \sigma_i) \times [\sigma_{j-1}, \sigma_j)$. We call S the **skeleton graph** of W for σ .

We illustrate the relationship between a step-graphon and its skeleton graph in Figure 2.

Without loss of generality and for ease of presentation, we will consider throughout this paper step-graphons W whose skeleton graphs are *connected*. Though there is no unique skeleton graph associated to a step-graphon (since there are infinitely many different partitions for W), we show in Proposition 1 that if one such skeleton graph is connected, then so are all the others. For S not connected, it is not too hard to see that the corresponding step-graphon is block-diagonal. Our results apply naturally to every connected component of S .

We decompose the edge set F of S as

$$F = F_0 \cup F_1, \quad (3)$$

where elements of F_0 are self-loops, and elements of F_1 are edges between distinct nodes. We also introduce the subset $F_2 \subseteq F_1$ of edges that are *not* incident to two nodes with self-loops.

Let $\mathcal{I} := \{1, \dots, |F|\}$ be an index set for F (so that the edges are now ordered). We decompose \mathcal{I} similarly: let $\mathcal{I}_0, \mathcal{I}_1$, and \mathcal{I}_2 index F_0, F_1 , and F_2 respectively.

To introduce the edge-polytope of S , we recall that the incidence matrix $Z = [z_{ij}]$ of S is an $|U| \times |F|$ matrix with its entries defined as follows:

$$z_{ij} := \frac{1}{2} \begin{cases} 2, & \text{if } f_j \in F_0 \text{ is a loop on node } u_i, \\ 1, & \text{if node } u_i \text{ is incident to } f_j \in F_1, \\ 0, & \text{otherwise.} \end{cases} \quad (4)$$

Owing to the factor $\frac{1}{2}$ in (4), all columns of Z are probability vectors, i.e., all entries are nonnegative and sum to one. The edge polytope of S was introduced in [14] and the definition is reproduced below (with a slight difference in inclusion of the factor $\frac{1}{2}$ of the generators z_j):

Definition 5 (Edge polytope). *Let $S = (U, F)$ be a skeleton graph and Z be the associated incidence matrix. Let z_j , for $1 \leq j \leq |F|$, be the columns of Z . The **edge polytope** of S , denoted by $\mathcal{X}(S)$, is the convex hull of the vectors z_j :*

$$\mathcal{X}(S) := \text{conv}\{z_j \mid j = 1, \dots, |F|\}. \quad (5)$$

A point $x \in \mathcal{X}(S)$ is said to be in the *relative interior* of $\mathcal{X}(S)$, denoted by $\text{int } \mathcal{X}(S)$, if there exists an open neighborhood U of x in \mathbb{R}^q (with $q = |U|$) such that $U \cap \mathcal{X}(S) \subseteq \mathcal{X}(S)$. If x is not an interior point, then it is called a *boundary point* and we write $x \in \partial \mathcal{X}(S)$.

D. Main results

Let W be a step-graphon. For a given partition σ for W , let x^* and S be the associated concentration vector and the skeleton graph (which is assumed to be connected). We say that a cycle in S is *odd* if it contains an odd number of distinct nodes (or edges); with this definition, self-loops are odd cycles. Given these, we state the following two conditions:

Condition A: The graph S has an odd cycle.

Condition B: The vector x^* belongs to $\text{int } \mathcal{X}(S)$.

The two conditions are stated in terms of a partition σ and its induced skeleton graph and edge-polytope. As mentioned earlier, there exist infinitely many partitions for a given step-graphon. However, the following proposition states that the two above conditions are *invariant* under changes of a partition.

Proposition 1. *Let W be a step-graphon. For any two partitions σ and σ' for W , let x^*, x'^* be the corresponding concentration vectors and let S, S' be the corresponding skeleton graphs. Then, the following hold:*

- 1) S is connected if and only if S' is;
- 2) S has an odd cycle if and only if S' does;

- 3) $x^* \in \mathcal{X}(S)$ (resp. $x^* \in \text{int } \mathcal{X}(S)$) if and only if $x'^* \in \mathcal{X}(S')$ (resp. $x'^* \in \mathcal{X}(S')$).

We refer the reader to Appendix A for a proof of the proposition.

We are now in a position to state the main result:

Main Theorem. *Let W be a step-graphon. If it satisfies Conditions A and B for a given (and, hence, any) partition σ , then it has the *H*-property.*

Remark 1. In our earlier work [1], we have shown that if a step-graphon W has the *H*-property, then it is necessary that Condition A and the following hold:

Condition B': The vector x^* belongs to $\mathcal{X}(S)$.

In fact, we have established there a stronger result, which states that if either Condition A or B' does not hold, then the probability that $G_n \sim W$ has a Hamiltonian decomposition converges to zero.

Note that condition B' is weaker than Condition B: Specifically, Condition B leaves out the set of step-graphons for which $x^* \in \partial \mathcal{X}(S)$, which is a set of measure zero. For step-graphons satisfying Conditions A and B', but not B, it is possible that

$$\lim_{n \rightarrow \infty} \mathbb{P}(\vec{G}_n \sim W \text{ has a Hamiltonian decomposition}) \in (0, 1).$$

We have produced explicit examples of such step-graphon in [1], [15]. \blacksquare

Outline of proof. Given a step-graphon W with skeleton graph S , and $G_n \sim W$, the sampling procedure induces a natural graph homomorphism $\pi : G_n \rightarrow S$, whereby all nodes v_j of G_n whose coordinates y_j belong to $[\sigma_{i-1}, \sigma_i)$ are mapped to u_i . With a slight abuse of notation, we will use the same letter π to denote the homomorphism $\pi : \vec{G}_n \rightarrow \vec{S}$.

Let $n_i(G_n) := |\pi^{-1}(u_i)|$ be the number of nodes whose coordinates belong to $[\sigma_{i-1}, \sigma_i)$. We call the following vector the *empirical concentration vector* of G_n :

$$x(G_n) := \frac{1}{n} (n_1(G_n), \dots, n_q(G_n)). \quad (6)$$

The proof of the Main Theorem contains three steps, outlined below, among which step 2 contains the bulk of the proof.

Step 1. The proof starts by showing how conditions A and B imply that the empirical concentration vector eventually belongs to the edge polytope. First, it should be clear that the edge polytope $\mathcal{X}(S)$ is a subset of the standard simplex Δ^{q-1} in \mathbb{R}^q ; thus, $\dim \mathcal{X}(S) \leq (q-1)$. Condition A, owing to [14], is both necessary and sufficient for the equality $\dim \mathcal{X}(S) = (q-1)$ to hold. Next, note that $nx(G_n) = (n_1(G_n), \dots, n_q(G_n))$ is a multinomial random variable with n trials and q outcomes with probabilities x_i^* , for $1 \leq i \leq q$. Then, Condition B guarantees, via Chebyshev's inequality, that $x(G_n)$ belongs to $\text{int } \mathcal{X}(S)$ a.a.s. (See the arguments around (34) for detail).

The next two steps are then dedicated to establishing the following fact:

$$\begin{aligned} x(G_n) &\in \text{int } \mathcal{X}(S) \\ &\Rightarrow G_n \text{ admits a Hamiltonian decomposition a.a.s..} \end{aligned} \quad (7)$$

Step 2. We start by working under the assumption that W is a *binary* step-graphon, i.e., $W(s, t) \in \{0, 1\}$ for almost all $(s, t) \in [0, 1]^2$. In this case, we will see that a sampled graph $G_n \sim W$ is *completely determined* by its empirical concentration vector $x(G_n)$. Consequently, our task (establishing (7)) is reduced to establishing the following:

$$\begin{aligned} x(G_n) &\in \text{int } \mathcal{X}(S) \text{ and } n \text{ is sufficiently large} \\ \Rightarrow G_n &\text{ admits a Hamiltonian decomposition surely.} \end{aligned} \quad (8)$$

The proof of (8) is constructive.

An important object that will arise therefrom is what we call the A -matrix assigned to every Hamiltonian decomposition H for \tilde{G}_n , written as a map $\rho(H) = A$.

Specifically, the matrix A is a q -by- q matrix whose ij th entry tallies the number of edges of H that go from a node in $\pi^{-1}(u_i)$ to a node in $\pi^{-1}(u_j)$ (A precise definition is in Subsection II-A and see Figure 3 for an illustration). Any such matrix is then shown to satisfy a number of enviable properties, among which we have $\rho(H)\mathbf{1} = x(G_n)$.

In a nutshell, we have just created the following sequence of maps:

$$H \mapsto \rho(H) \mapsto \rho(H)\mathbf{1} = x(G_n),$$

with the *domain* being all Hamiltonian decompositions in \tilde{G}_n , for any G_n sampled from a given binary graphon.

Now, the effort in establishing (8) is to create appropriate right-inverses (at least locally) of the maps in the above sequence, i.e., we aim to create maps $x \mapsto A(x)$ and $\tilde{\rho} : A(x) \mapsto H$ with the property that $\rho \cdot \tilde{\rho}$ is the identity map and $A(x)\mathbf{1} = x$. The map $x \mapsto A(x)$ is created in Proposition 2, Subsection II-C, and the map $\tilde{\rho}$ is created in Proposition 3, Subsection II-D. From these two subsections, it will be clear that by introducing the A -matrix as an intermediate object, we can decouple the analytic part of the proof, contained in the creation of the map $x \mapsto A(x)$, from the graph-theoretic part, contained in the creation of $\tilde{\rho}$. This will conclude the proof of (8).

Step 3. To close gap between *binary* step-graphons and general ones, we introduce here an equivalence relation on the class of step-graphons, namely, two step-graphons W_1 and W_2 are equivalent if their supports are the same. Or, said otherwise, W_1 and W_2 are equivalent if they share the same concentration vector and skeleton graph. Note, in particular, that each equivalence class $[W]$ contains a unique representative which is a *binary* step-graphon, denoted by W^s . We then establish (7) by showing that W has the H -property if and only if W^s does. In essence, we show that the H -property is decided completely by the concentration vector and the skeleton graph of a step-graphon W . The proof of this statement builds upon several classical results from random graph theory, and is presented in Subsection II-E. \blacksquare

E. Notation

We gather here key notations and conventions.

Graph theory. Let $G = (V, E)$ be an undirected graph. Graphs in this paper do *not* have multiple edges, but may have self-

loops. We denote edges by $(v_i, v_j) \in E$; if $v_i = v_j$, then we call the edge a *self-loop*. For a given node v_i , let $N(v_i) := \{v_j \in V \mid (v_i, v_j) \in E\}$ be the *neighbor set* of v_i . The degree of v_i , denoted by $\deg(v_i)$, is the cardinality of $N(v_i)$.

We will also deal with digraphs in this paper. Whether a graph is directed or undirected will be clear from the context and/or notation. We denote by $v_i v_j$ the directed edge from v_i to v_j ; we call v_j an *out-neighbor* of v_i and v_i an *in-neighbor* of v_j . Similarly, we define $N_+(v_i)$ and $N_-(v_i)$ the sets of *in-neighbors* and *out-neighbors* of v_i , respectively.

Recall that for a given undirected graph $G = (V, E)$, possibly with self-loops, we let $\tilde{G} = (V, \tilde{E})$ be the directed graph defined as in (1). Self-loops in \tilde{G} are the same as the ones in G , i.e., they are not duplicated.

A *closed walk* in a graph (or digraph) is an ordered sequence of nodes $v_1 v_2 \cdots v_k v_1$ in G (resp. \tilde{G}) so that all consecutive nodes are ends of some edges (resp. directed edges). A *cycle* is a closed walk without repetition of nodes in the sequence except the starting- and the ending-nodes. For clarity of the presentation, we use letter C to denote cycles in undirected graphs and the letter D for cycles in directed graphs.

Miscellanea. We use $\mathbf{1}$ to denote a column vector of all 1, whose dimension will be clear within the context. We write $x \leq y$ for vectors $x, y \in \mathbb{R}^q$ if the inequality holds entrywise. For a given vector $x \in \mathbb{R}^q$, we denote its ℓ_1 normalization by \bar{x} , i.e., $\bar{x} := \frac{x}{\|x\|_1}$, with the convention that $\bar{0} = 0$. Further, given the vector x , we denote by $[x]$ the vector whose entry $[x]_i$ is a closest integer to x_i for $1 \leq i \leq q$ where for the case $x_i = k + \frac{1}{2}$, with k an integer, we set $[x]_i := k$. We denote the standard simplex in \mathbb{R}^q by $\Delta^{q-1} := \{x \in \mathbb{R}^q \mid x \geq 0 \text{ and } x^\top \mathbf{1} = 1\}$. Finally, given a $q \times q$ matrix A , we denote by $\text{supp } A$ its *support*, i.e., the set of indices corresponding to its non-zero entries.

II. ANALYSIS AND PROOF OF THE MAIN THEOREM

Throughout the proof, W is a step-graphon, σ its associated partition, and $S = (U, F)$ its skeleton graph on q nodes, which has an odd cycle. Let F_0 and F_1 as in (3). We can naturally associate to them the subgraphs

$$S_0 := (U, F_0) \quad \text{and} \quad S_1 := (U, F_1). \quad (9)$$

Note that S has an odd cycle if and only if S_0 is edgewise non-empty or S_1 has an odd cycle. The lemma below states that we can consider, without loss of generality, only the latter case of S_1 containing an odd cycle.

Lemma 1. *Let W be a step-graphon. If W admits a partition σ with skeleton graph S containing an odd cycle, then W admits a partition σ' with skeleton graph S' so that the subgraph S'_1 has an odd cycle.*

The proof of the lemma can be established by using the notion of “one-step refinement” introduced in Appendix A for the partition σ : If S_1 already has an odd cycle, then there is nothing to prove. Otherwise, consider a one-step refinement on a node with self-loop in S , which will yield a cycle of length 3 in S' .

A. On the edge polytope $\mathcal{X}(S)$

Rank of $\mathcal{X}(S)$. Recall that $\mathcal{X}(S)$ is the edge-polytope of S . Similarly, we let $\mathcal{X}(S_i)$, for $i = 0, 1$, be the edge polytope (see Definition 5) of S_i , i.e., $\mathcal{X}(S_i)$ is the convex hull of the z_j 's, with $j \in \mathcal{I}_i$, where \mathcal{I}_i indexes edges in F_i .

We call x an *extremal point* of a polytope \mathcal{X} if there is no line segment in \mathcal{X} that contains x in its interior. The maximal set of extremal points is called the set of *extremal generators* for \mathcal{X} . The following result characterizes the extremal generators of $\mathcal{X}(S_0)$, $\mathcal{X}(S_1)$, and of $\mathcal{X}(S)$:

Lemma 2. *The set of extremal generators of $\mathcal{X}(S_i)$, for $i = 0, 1$, is $\{z_j \mid j \in \mathcal{I}_i\}$. The set of extremal generators of $\mathcal{X}(S)$ is $\{z_i \mid i \in \mathcal{I}_0 \cup \mathcal{I}_2\}$.*

Proof. The statement for $\mathcal{X}(S_0)$ is obvious from the definition of the corresponding z_i . For $\mathcal{X}(S_1)$, it suffices to see that the vectors z_i , for $i \in \mathcal{I}_1$, have exactly two non-negative entries, and the supports of these vectors are pairwise distinct. Hence, if $z_i = \sum_{j \in \mathcal{I}_1} c_j z_j$ with $c_j \geq 0$, we necessarily have $c_j = 0$ for $j \neq i$ and $c_i = 1$. For $\mathcal{X}(S)$, we refer the reader to [1, Proposition 1] for a proof. \blacksquare

The rank of a polytope \mathcal{X} is the dimension of its relative interior. It is known [14] that if S has q nodes, then

$$\text{rank } \mathcal{X}(S) = \begin{cases} q-1 & \text{if } S \text{ has an odd cycle,} \\ q-2 & \text{otherwise.} \end{cases} \quad (10)$$

Equivalently, we have the following result [16] on the rank of the incidence matrix Z_S of S :

$$\text{rank } Z_S = \begin{cases} q & \text{if } S \text{ has an odd cycle,} \\ q-1 & \text{otherwise.} \end{cases} \quad (11)$$

The A -matrix. Let $G_n \sim W$ and suppose that \vec{G}_n has a Hamiltonian decomposition, denoted by H .

Recall that $\pi : \vec{G}_n \rightarrow \vec{S}$ is the graph homomorphism introduced above (6). Let $n_{ij}(H)$ be the number of (directed) edges of H from a node in $\pi^{-1}(u_i)$ to a node in $\pi^{-1}(u_j)$. It is not hard to see that (see [1, Lemma 1] for a proof) for all $u_i \in U$,

$$n_i(G_n) = \sum_{u_j \in N(u_i)} n_{ij}(H) = \sum_{u_j \in N(u_i)} n_{ji}(H). \quad (12)$$

We now assign to the skeleton graph S a convex set that will be instrumental in establishing the main result:

Definition 6 (A -matrices and their set). *Let $S = (U, F)$ be an undirected graph on q nodes. We define $\mathcal{A}(S)$ as the set of $q \times q$ nonnegative matrices $A = [a_{ij}]$ satisfying the following two conditions:*

- 1) If $(u_i, u_j) \notin F$, then $a_{ij} = 0$;
- 2) $A\mathbf{1} = A^\top \mathbf{1}$, and $\mathbf{1}^\top A\mathbf{1} = 1$.

Because every defining condition for $\mathcal{A}(S)$ is affine, the set $\mathcal{A}(S)$ is a convex set.

Now, to each Hamiltonian decomposition H of \vec{G}_n , we assign the following $q \times q$ matrix:

$$\rho(H) := \frac{1}{n} [n_{ij}(H)]_{1 \leq i, j \leq q}. \quad (13)$$

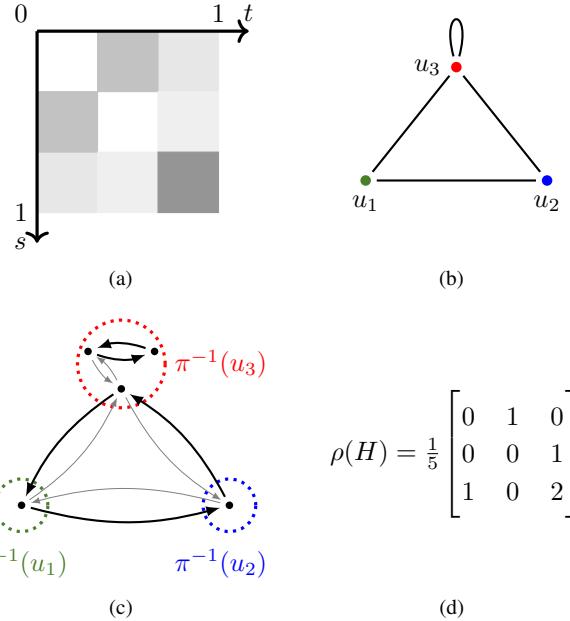


Fig. 3: (a): A step-graphon W . (b): Its skeleton graph S . (c): A digraph \vec{G}_n , with $G_n \sim W$, for $n = 5$. The edges in black form the Hamiltonian decomposition H . (d): The matrix $\rho(H)$ defined in (13).

It follows from (12) that $\rho(H) \in \mathcal{A}(S)$ and $\rho(H)\mathbf{1} = x(G_n)$. Furthermore, we have established in [1, Proposition 4] the following connection between the set $\mathcal{A}(S)$ and the edge polytope $\mathcal{X}(S)$:

$$\mathcal{X}(S) = \{x \in \mathbb{R}^q \mid x = A\mathbf{1} \text{ for some } A \in \mathcal{A}(S)\}. \quad (14)$$

We refer the reader to Figure 3 for illustration.

B. Local coordinate systems on $\mathcal{X}(S)$ and $\mathcal{X}(S_1)$

This section establishes the groundwork for the construction of the map $x \mapsto A(x)$ described in the proof outline. To this end, we first show how to express any point in a neighborhood \mathcal{U} of $x^* \in \mathcal{X}(S)$ as a *positive combination* of the columns of the incidence matrix Z_S . This amounts to solving the linear equations $Z_S \phi(x) = x$, for $x \in \mathcal{U}$, with $\phi(x)$ being continuous in x and positive. We will solve a similar problem for $y^* \in \mathcal{X}(S_1)$ and with Z_S replaced by Z_{S_1} , and we call the corresponding solution $\theta(y)$. These two maps will be put to use in the next subsection.

Construction of the map ϕ . We start with the following lemma:

Lemma 3. *Suppose that $S = (U, F)$ has an odd cycle. Then, for any $x^* \in \text{int } \mathcal{X}(S)$, there exist a closed neighborhood \mathcal{U} of x^* in the simplex and a continuous map $\phi : \mathcal{U} \rightarrow \text{int } \Delta^{|F|-1}$ such that $Z_S \phi(x) = x$ for any $x \in \mathcal{U}$.*

Proof. Because $\mathcal{X}(S)$ is finitely generated by the columns of Z_S , i.e. the z_i 's for $i \in \mathcal{I}$, and because $x^* \in \text{int } \mathcal{X}(S)$, there exists a *positive* probability vector $c := (c_1, \dots, c_{|F|})$ such that $x^* = Z_S c$. Let $\epsilon := \frac{1}{2} \min_{i \in \mathcal{I}} \{c_i\}$. Because S has at least one odd cycle, we know from (11) that Z_S is full rank. Thus, we can pick q columns, say z_1, \dots, z_q of Z_S , that form a basis of

\mathbb{R}^q . Let $B \subset \mathbb{R}^{|F|}$ be a closed ball centered at 0 with radius ϵ , and let

$$B_0 := \{y \in B \mid \mathbf{1}^\top y = 0 \text{ and } y_i = 0, \text{ for all } i > q\}.$$

The dimension of B_0 is $(q-1)$. We now define the map

$$\psi : B_0 \rightarrow \mathbb{R}^q : y \mapsto x^* + Z_S y = Z_S(y + c).$$

It should be clear that ψ is a linear bijection between B_0 and its image $\psi(B_0)$. By the construction of B and B_0 , all the entries of $(y + c)$, for $y \in B_0$, are positive and, moreover, $\mathbf{1}^\top(y + c) = 1$. It then follows that the image of ψ is a closed neighborhood of x^* inside $\mathcal{X}(S)$. We now set $\phi := \psi^{-1}$. It remains to show that $\phi(x) \in \Delta^{|F|-1}$. This holds because every column of Z_S belongs to Δ^{q-1} and so does x . Thus, from $Z_S\phi(x) = x$, we have that x is a convex combination of the columns of Z_S , which implies that $\phi(x) \in \Delta^{|F|-1}$. ■

Let $x^* \in \text{int } \mathcal{X}(S)$ and ϕ be the map given in Lemma 3. For an edge f_i of S , we let ϕ_i be the corresponding entry of ϕ . We next define two functions $\tau_i : \mathcal{U} \rightarrow \mathbb{R}^q$, for $i = 0, 1$, as follows:

$$\tau_i(x) : x \mapsto \sum_{j \in \mathcal{I}_i} \phi_j(x) z_j. \quad (15)$$

If S has no self-loops, then τ_0 is set to be the zero map. We can thus decompose $x \in \mathcal{U}$ as

$$x = \tau_0(x) + \tau_1(x).$$

We record the following simple observation for later use:

Lemma 4. *Let \mathcal{U} be as in Lemma 3. For every $x \in \mathcal{U}$, the set of indices of nonzero (positive) entries of $\tau_0(x)$ is $\{i \mid u_i \text{ has a self-loop}\}$ and, moreover, every entry of $\tau_1(x)$ is positive.*

Proof. The statement for $\tau_0(x)$ is trivial. The statement for $\tau_1(x)$ follows from the fact that $\phi(x)$ has positive entries and no row of Z_{S_1} is identically 0. ■

Construction of the map θ . For any $x \in \mathcal{U}$, let $\bar{\tau}_i(x)$, for $i = 0, 1$, be defined as follows:

$$\bar{\tau}_i(x) := \begin{cases} \tau_i(x)/\|\tau_i(x)\|_1 & \text{if } \tau_i(x) \neq 0, \\ 0 & \text{otherwise.} \end{cases}$$

Since S has odd cycle, recall that we can assume by Lemma 1 that S_1 has an odd cycle. Thus, by (10), the rank of $\mathcal{X}(S_1)$ is $(q-1)$. In particular, it implies that the relative interior of $\mathcal{X}(S_1)$ is open in Δ^{q-1} . Further, by Lemma 4, if $x \in \mathcal{U}$, then $\bar{\tau}_1(x) \in \text{int } \mathcal{X}(S_1)$.

The map θ we introduce below is akin to the map ϕ introduced in Lemma 3, but defined on a closed neighborhood of the following vector

$$\bar{x}_1^* := \bar{\tau}_1(x^*) \in \text{int } \mathcal{X}(S_1).$$

Lemma 5. *Suppose that S (and, hence, S_1) has an odd cycle; then, for the given $\bar{x}_1^* \in \text{int } \mathcal{X}(S_1)$, there exist a closed neighborhood \mathcal{V} of \bar{x}_1^* in Δ^{q-1} and a continuous map $\theta : \mathcal{V} \rightarrow \text{int } \Delta^{|F_1|-1}$ such that $Z_{S_1}\theta(y) = y$ for any $y \in \mathcal{V}$.*

The proof is entirely similar to the one of Lemma 3, and is

thus omitted.

Because ϕ and θ are both positive, continuous maps over closed, bounded domains, there exists an $\alpha \in (0, 1)$ so that

$$\begin{aligned} \phi(x) &\geq \alpha \mathbf{1} \text{ for all } x \in \mathcal{U}, \\ \theta(x) &\geq \alpha \mathbf{1} \text{ for all } x \in \mathcal{V}. \end{aligned} \quad (16)$$

On the image of $\bar{\tau}_1$. For a given $x^* \in \text{int } \mathcal{X}(S)$, the domains of ϕ and θ are closed neighborhoods \mathcal{U} and \mathcal{V} of x^* and \bar{x}_1^* , respectively. Later in the analysis, we will pick an arbitrary $x \in \mathcal{U}$ and apply θ to $\bar{\tau}_1(x)$. For this, we need that $\bar{\tau}_1(x)$ belongs to \mathcal{V} . To this end, we will shrink \mathcal{U} so that $\bar{\tau}_1(\mathcal{U}) \subseteq \mathcal{V}$ and thus the composition $\theta\bar{\tau}_1$ is well defined. In fact, we have the stronger statement:

Lemma 6. *Let $\alpha > 0$ be given as in (16). There exist a closed neighborhood $\mathcal{U}' \subseteq \mathcal{U}$ of x^* and a positive $\epsilon < \frac{1}{4}\alpha$, such that*

$$\overline{\tau_1(x) + \eta} = \frac{\tau_1(x) + \eta}{\|\tau_1(x) + \eta\|_1} \in \mathcal{V},$$

for any $x \in \mathcal{U}'$ and for any $\eta \in \mathbb{R}^q$ with $\|\eta\|_\infty \leq \epsilon$.

Proof. Let \mathcal{V}' be a closed ball centered at \bar{x}_1^* and contained in the interior of \mathcal{V} . Then, it is known [17, Theorem 4.6] that there exists an $\epsilon' > 0$ such that the ϵ' -neighborhood of \mathcal{V}' , with respect to the infinity norm, is contained in the interior of \mathcal{V} . Let $\mathcal{U}' := \bar{\tau}_1^{-1}(\mathcal{V}')$ and ϵ be sufficiently small so that

$$\frac{(8+4\alpha)\epsilon}{q\alpha^2} < \min \left\{ \epsilon', \frac{1}{4}\alpha \right\}. \quad (17)$$

We claim that the above-defined \mathcal{U}' and ϵ are as desired.

Since $\bar{\tau}_1$ is continuous and since \mathcal{V}' is a closed ball centered at \bar{x}_1^* , \mathcal{U}' is a closed neighborhood of x^* . Now, pick an arbitrary $x \in \mathcal{U}'$. For ease of notation, we set $x_1 := \tau_1(x)$ and $\bar{x}_1 = \bar{\tau}_1(x)$ for the remainder of this proof. Then,

$$\begin{aligned} \|\overline{x_1 + \eta} - \bar{x}_1\|_\infty &= \left\| \frac{x_1 + \eta}{\|x_1 + \eta\|_1} - \frac{x_1}{\|x_1\|_1} \right\|_\infty \\ &= \left\| \frac{(\|x_1\|_1 - \|x_1 + \eta\|_1)x_1 + \|x_1\|_1\eta}{\|x_1 + \eta\|_1\|x_1\|_1} \right\|_\infty \\ &\leq \frac{\|x_1\|_1 - \|x_1 + \eta\|_1}{\|x_1 + \eta\|_1\|x_1\|_1} \|x_1\|_\infty + \frac{\|\eta\|_\infty}{\|x_1 + \eta\|_1} \\ &\leq \frac{\|x_1\|_1 - \|x_1 + \eta\|_1}{\|x_1 + \eta\|_1\|x_1\|_1} + \frac{\|\eta\|_\infty}{\|x_1 + \eta\|_1} \end{aligned} \quad (18)$$

where we used the fact that $\|x_1\|_\infty \leq 1$ to obtain the last inequality. To further evaluate (18), we first note that

$$| \|x_1\|_1 - \|x_1 + \eta\|_1 | \leq \|\eta\|_1 \leq q\|\eta\|_\infty \leq q\epsilon.$$

Next, by (4), (15) and (16), every entry of x_1 is greater than $\frac{1}{2}\alpha$, so $\|x_1\|_1 \geq \frac{1}{2}q\alpha$. Moreover, since $\epsilon < \frac{1}{4}\alpha$,

$$\|x_1 + \eta\|_1 \geq \|x_1\|_1 - \|\eta\|_1 \geq \frac{1}{2}q\alpha - q\epsilon \geq \frac{1}{4}q\alpha.$$

Finally, using (17), we can proceed from (18) and obtain that

$$\|\overline{x_1 + \eta} - \bar{x}_1\|_\infty \leq \frac{8\epsilon}{q\alpha^2} + \frac{4\epsilon}{q\alpha} = \frac{(8+4\alpha)\epsilon}{q\alpha^2} < \epsilon',$$

which implies that $\overline{x_1 + \eta}$ belongs to the ϵ' -neighborhood of \mathcal{V}' and, hence, to \mathcal{V} . ■

Remark 2. From now on, to simplify the notation, we denote by \mathcal{U} the set \mathcal{U}' of Lemma 6. \blacksquare

C. Construction of the map $x \mapsto A(x)$

To construct the map, we first specify its domain, which will be a subset of \mathcal{U} . If $x \in \mathbb{R}^q$ is an empirical concentration vector of some $G_n \sim W$ for a step-graphon W , then nx necessarily has integer entries. Define a subset of \mathcal{U} as follows:

$$\mathcal{U}^* := \{x \in \mathcal{U} \mid nx \in \mathbb{Z}_+^q \text{ for some } n \in \mathbb{Z}_+\}, \quad (19)$$

where \mathbb{Z}_+ is the set of positive integers.

Since the analysis for the H -property will be carried out in the asymptotic regime $n \rightarrow \infty$, the relevant empirical concentration vectors are those of G_n for n large. To this end, let $\alpha \in (0, 1)$ be such that (16) is satisfied and ϵ be as in Lemma 6. We have the following definition:

Definition 7. Given an $x \in \mathcal{U}^*$, $n \in \mathbb{Z}_+$ is **paired with** x if $n > \frac{8}{\alpha\epsilon}$ and nx is integer valued.

With the above, we now state the main result of this subsection:

Proposition 2. Let $S = (U, F)$ be a connected undirected graph, with at least one odd cycle, and $\vec{S} = (U, \vec{F})$ be its directed version. Then, there exist a map $A : \mathcal{U}^* \mapsto \mathcal{A}(S)$ and a positive number \underline{a} such that for any $x \in \mathcal{U}^*$ and for any n paired with x , the following hold:

- 1) $A(x)\mathbf{1} = x$;
- 2) $nA(x)$ is integer-valued and $n \text{diag } A(x)$ has even entries;
- 3) $n\|\text{diag } A(x) - \tau_0(x)\|_\infty \leq 1$, where τ_0 is defined in (15);
- 4) $n|a_{ij}(x) - a_{ji}(x)| \leq 1$ for all $1 \leq i, j \leq q$;
- 5) For any $u_i u_j \in \vec{F}$, $a_{ij}(x) > \underline{a}$.

Note that item 5 and the fact that $A(x) \in \mathcal{A}(S)$ imply $\text{supp } A(x) = \vec{S}$, i.e., $a_{ij} \neq 0$ if and only if $u_i u_j \in \vec{F}$.

The proof of Proposition 2 is constructive. It will rely on a few technical facts we establish here. Let $x \in \mathcal{U}^*$ and n be paired with x . Define

$$\tau'_0(x) := \frac{2}{n} \left[\frac{n}{2} \tau_0(x) \right] \quad \text{and} \quad \tau'_1(x) := x - \tau'_0(x), \quad (20)$$

where we recall that the operation $[\cdot]$ is the integer-rounding operation, introduced in the notation of Section I. The vector $n\tau'_0(x)$ is then the vector with *even* entries closest to the entries of $n\tau_0(x)$. Next, we define

$$n'_0 := n\|\tau'_0(x)\|_1 \quad \text{and} \quad n'_1 := n\|\tau'_1(x)\|_1 \quad (21)$$

Recall that $\bar{\tau}'_0$ and $\bar{\tau}'_1$ are the ℓ_1 normalization of τ'_0 and τ'_1 , respectively. It should be clear from the construction that $n'_0 + n'_1 = n$ and

$$\bar{\tau}'_i = \frac{n}{n'_i} \tau'_i, \quad \text{for } i = 0, 1.$$

For a given $x \in \mathcal{U}^*$, there obviously exist infinitely many positive integers n that are paired with x . However, the ratios n'_0/n and n'_1/n are independent of n and determined completely by x .

We also need the following lemma:

Lemma 7. The following items hold:

- 1) For $i = 0, 1$, $\text{supp } \tau'_i(x) = \text{supp } \tau_i(x)$ and, moreover, the nonzero entries of $\tau'_i(x)$ are uniformly bounded below by $\frac{1}{2}(1 - \epsilon/4)\alpha$.
- 2) The ratio n'_0/n is bounded below by $(1 - \epsilon/8)\alpha|F_0|$. The ratio n'_1/n is bounded below by $3q\alpha/8$.
- 3) Let \mathcal{V} be defined as in Lemma 5. Then, $\bar{\tau}_1(x) \in \mathcal{V}$.

We provide a proof of the lemma in Appendix B.

With the lemma above, we now establish Proposition 2:

Proof of Proposition 2. We start by defining two matrix-valued functions $A_0(x)$ and $A_1(x)$ so that for any $x \in \mathcal{U}^*$, $A_i(x) \in \mathcal{A}(S)$ and $A_i(x)\mathbf{1} = \bar{\tau}'_i(x)$. We will then let $A(x)$ be the convex combination of these two matrices given by

$$A(x) = \frac{n'_0}{n} A_0(x) + \frac{n'_1}{n} A_1(x). \quad (22)$$

Since $\mathcal{A}(S)$ is convex, it will then follow that $A(x) \in \mathcal{A}(S)$.

Construction of A_0 . The matrix $A_0(x)$ is simply given by

$$A_0(x) := \text{diag } \bar{\tau}'_0(x). \quad (23)$$

By Lemma 4, $\text{supp } \tau_0(x)$ is constant over \mathcal{U} . By the first item of Lemma 7, $\text{supp } \tau_0(x) = \text{supp } \tau'_0(x)$ for all $x \in \mathcal{U}^*$. It then follows that $\text{supp } A_0(x)$ is also constant over \mathcal{U}^* . By the same item, the nonzero entries of $A_0(x)$ are uniformly bounded below by a positive constant.

Construction of A_1 . The construction is more involved than the one of A_0 , and requires to first define the intermediate matrix A'_1 . To this end, recall that Z_{S_1} is the edge-incidence matrix of $S_1 = (U, F_1)$, obtained by removing the self-loops of S , and that θ is the map given in Lemma 5, i.e., $Z_{S_1}\theta(y) = y$ for all $y \in \mathcal{V}$. Given an edge $f = (u_i, u_j) \in F_1$, we denote by $\theta_f(y)$ the corresponding entry of $\theta(y)$. By item 3 of Lemma 7, $\bar{\tau}'_1(x)$ belongs to \mathcal{V} , which is the domain of θ . Now, we define the symmetric matrix $A'_1(x) = [a'_{1,ij}(x)] \in \mathbb{R}^{q \times q}$ as follows:

$$a'_{1,ij}(x) := \begin{cases} \frac{1}{2}\theta_f(\bar{\tau}'_1(x)) & \text{if } f = (u_i, u_j) \in F_1, \\ 0 & \text{otherwise.} \end{cases} \quad (24)$$

In particular, the diagonal of A'_1 is 0, and so will be the diagonal of A_1 as shown below. From the definition of the incidence matrix Z_{S_1} and (24), we have that $A'_1(x)\mathbf{1} = Z_{S_1}\theta(\bar{\tau}'_1(x))$. By Lemma 5, $Z_{S_1}\theta(\bar{\tau}'_1(x)) = \bar{\tau}'_1(x)$. It then follows that

$$A'_1(x)\mathbf{1} = \bar{\tau}'_1(x), \quad (25)$$

and, hence, $\mathbf{1}^\top A'_1(x)\mathbf{1} = \mathbf{1}^\top \bar{\tau}'_1(x) = 1$. Furthermore, since $A'_1(x)$ is symmetric, $A'_1(x)\mathbf{1} = A'_1(x)^\top \mathbf{1}$. We thus have that $A'_1(x) \in \mathcal{A}(S)$. Since θ_f is positive for every $f \in F_1$, $\text{supp } A'_1(x)$ is constant over \mathcal{U} . Moreover, by (16), the nonzero entries of $A'_1(x)$ are uniformly bounded below by $\frac{1}{2}\alpha$.

Next, we use A'_1 to construct A_1 . There are two cases; one is straightforward and the other is more involved:

Case 1: $n'_1 A'_1(x)$ is integer-valued. Set $A_1(x) := A'_1(x)$.

Case 2: $n'_1 A'_1(x)$ is not integer-valued. In this case, we appeal to the result [18, Theorem 2]: There, we have shown that there

exists a matrix $A_1(x) = [a_{1,ij}(x)]$ in $\mathcal{A}(S)$, with

$$A_1(x)\mathbf{1} = A'_1(x)\mathbf{1} = \bar{\tau}'_1(x) \quad (26)$$

such that $A_1(x)$ has the same support as $A'_1(x)$ and

$$n'_1 a_{1,ij}(x) = \lfloor n'_1 a'_{1,ij}(x) \rfloor \quad \text{or} \quad n'_1 a_{1,ij}(x) = \lceil n'_1 a'_{1,ij}(x) \rceil.$$

In particular, $n'_1 A_1(x)$ is integer-valued. Because $a'_{1,ij}(x) = a'_{1,ji}(x)$, it follows that

$$n'_1 |a_{1,ij}(x) - a_{1,ji}(x)| \leq 1, \quad \forall 1 \leq i, j \leq q. \quad (27)$$

Moreover, if $a'_{1,ij}(x) > 0$, then

$$\begin{aligned} a_{1,ij}(x) > a'_{1,ij}(x) - \frac{1}{n'_1} &\geq \alpha - \frac{1}{n'_1} \geq \alpha - \frac{1}{n} \frac{n}{n'_1} \\ &\geq \alpha - \frac{8}{3nq\alpha} > \left(1 - \frac{1}{12q}\right) \alpha. \end{aligned} \quad (28)$$

where second to the last inequality follow from item 2 of Lemma 7 and the last inequality follows from the hypothesis on n (specifically $n > \frac{8}{\alpha\epsilon}$) from the statement and the condition that $\epsilon < \alpha/4$ from Lemma 6.

Proof that A satisfies the five items of the statement.

- 1) From (23), $A_0(x)\mathbf{1} = \bar{\tau}'_0(x)$. For A_1 , it was shown that $A_1(x)\mathbf{1} = \bar{\tau}'_1(x)$ in (25) and (26) for Case 1 and Case 2, respectively. Since A is the convex combination of A_0 and A_1 given in (22), it follows that

$$A(x)\mathbf{1} = \frac{n'_0}{n} \bar{\tau}'_0(x) + \frac{n'_1}{n} \bar{\tau}'_1(x) = x. \quad (29)$$

- 2) By the construction of A in (22) and the definitions of A_0 and A_1 , the diagonal of $nA(x)$ is

$$n'_0 A_0(x) = n'_0 \text{Diag } \bar{\tau}'_0(x) = n \text{Diag } \tau'_0(x).$$

By (20), all the entries of $n\tau'_0(x)$ are even.

- 3) Using (20) again, we have that

$$-\mathbf{1} \leq n(\tau_0(x) - \tau'_0(x)) \leq \mathbf{1},$$

from which it follows that

$$n \|\text{diag } A(x) - \tau_0(x)\|_\infty = n \|\tau'_0(x) - \tau_0(x)\| \leq 1.$$

- 4) The off-diagonal entries $a_{ij}(x)$ of $A(x)$ are those of $\frac{n'_1}{n} A_1(x)$, which we denoted by $\frac{n'_1}{n} a_{1,ij}(x)$. Thus,

$$n |a_{ij}(x) - a_{ji}(x)| = n'_1 |a_{1,ij}(x) - a_{1,ji}(x)| \leq 1,$$

where the last inequality is (27).

- 5) *Case 1: S does not have a self-loop.* In this case, $A(x) = A_1(x)$. By construction of A_1 , $\text{supp } A_1(x) = \vec{S}$. If $A_1(x)$ is obtained via case 1 above, then, as argued after (25), its nonzero entries are bounded below by $\alpha/2$. Otherwise, A_1 is obtained via case 2 and its nonzero entries are lower bounded as shown in (28).

Case 2: S has at least one self-loop. In this case,

$$\text{supp } A(x) = \text{supp } A_0(x) \cup \text{supp } A_1(x).$$

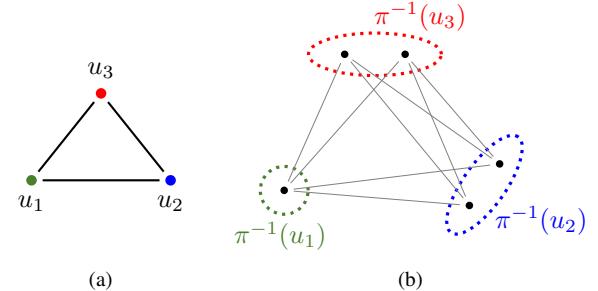


Fig. 4: (a): An undirected graph S on three nodes. (b): A complete S -multipartite graph $M(\omega, S)$ with $\omega = (1, 2, 2)$.

By construction of A_0 and item 1 of Lemma 7,

$$\begin{aligned} \text{supp } A_0(x) &= \text{supp } \text{Diag } \tau'_0(x) \\ &= \text{supp } \text{Diag } \tau_0(x) = \vec{S}_0, \end{aligned} \quad (30)$$

where the last equality follows from Lemma 4. Moreover, the nonzero entries of $A_0(x)$ are bounded below by $\frac{1}{2}(1 - \epsilon/4)\alpha$. Also, by construction of A_1 ,

$$\text{supp } A_1(x) = \vec{S}_1. \quad (31)$$

Thus, by (30) and (31), $\text{supp } A(x) = \vec{S}$. Finally, we verify that the nonzero entries of $A(x)$ are uniformly bounded below by a positive number. By item 2 of Lemma 7, n'_0/n and n'_1/n are uniformly bounded below by positive numbers (note that $|F_0| \geq 1$ in the current case). Thus, using (22), the nonzero entries of $A(x)$ are also uniformly lower bounded by a positive number.

This completes the proof. \blacksquare

D. Constructing a Hamiltonian decomposition from $A(x)$

In this subsection, we construct the map $\tilde{\rho} : A(x) \mapsto H$ announced in Section I, where $A(x)$ will be taken from the statement of Proposition 2 and H is a Hamiltonian decomposition in $G_n \sim W$, with x its empirical concentration vector. Throughout this subsection, we assume that W is a *binary step-graphon*, i.e., W is valued in $\{0, 1\}$.

Graphs sampled from a binary step-graphon have rather rigid structures as we will describe below. We refer to them as *S-multipartite graphs*, see also Figure 4:

Definition 8 (*S*-multipartite graph). *Let $S = (U, F)$ be an undirected graph, possibly with self-loops. An undirected graph G is an ***S*-multipartite graph** if there exists a graph homomorphism $\pi : G \rightarrow S$, so that*

$$(v_i, v_j) \in E \Rightarrow (\pi(v_i), \pi(v_j)) \in F.$$

Further, G is a **complete *S*-multipartite graph** if

$$(v_i, v_j) \in E \Leftrightarrow (\pi(v_i), \pi(v_j)) \in F.$$

Let G be an arbitrary complete *S*-multipartite graph with $S = (U, F)$ and set $n_i := |\pi^{-1}(u_i)|$ for $i = 1, \dots, q$. It should be clear that G is completely determined by S and the vector $w := (n_1, \dots, n_q)$. We will consequently use the notation $M(w, S)$ to refer to a complete *S*-multipartite graph. Now,

returning to the case $G_n \sim W$, where W is a binary step-graphon with skeleton graph S , the empirical concentration vector $x(G_n)$ together with S then completely determine G_n as announced above.

If G is a (complete) S -multipartite graph, then \vec{G} is (complete) \vec{S} -multipartite, and we use the same notation π to denote the homomorphism. We next introduce a special class of cycles in \vec{G} :

Definition 9 (Simple cycle). *Let G be an S -multipartite graph, and $\pi : \vec{G} \rightarrow \vec{S}$ be the associated homomorphism. A directed cycle D in \vec{G} is called **simple** if $\pi(D)$ is a cycle (rather than a closed walk) in \vec{S} .*

With the notions above, we state the main result of this subsection:

Proposition 3. *Let $S = (U, F)$ be an undirected graph, possibly with self-loops. Let $G = M(nx, S)$ be a complete S -multipartite graph on n nodes, where $x \in \mathcal{U}^*$ and n is paired with x (see Definition 7). Let $A(x)$ be as in Proposition 2 and*

$$m_{ij}(x) := n \min\{a_{ij}(x), a_{ji}(x)\}, \text{ for all } 1 \leq i, j \leq q. \quad (32)$$

Then, there exists a Hamiltonian decomposition H of \vec{G} , with $\rho(H) = A(x)$, such that the following hold:

- 1) *There exist exactly $\frac{1}{2}m_{ii}(x)$ disjoint 2-cycles in H pairing $m_{ii}(x)$ nodes in $\pi^{-1}(u_i)$ for every $i = 1, \dots, q$;*
- 2) *There are at least $m_{ij}(x)$ disjoint 2-cycles in H pairing nodes in $\pi^{-1}(u_i)$ to nodes in $\pi^{-1}(u_j)$ for each $(u_i, u_j) \in F_1$;*
- 3) *There are at most $\lceil \frac{2}{3}|F| \rceil$ cycles of length three or more in H ;*
- 4) *The length of every cycle of H does not exceed $2|F|$;*
- 5) *All cycles of length at least 3 of H are simple.*

We illustrate the Proposition on an example.

Example 1. Consider a complete S -multipartite graph G for S shown in Figure 4a. Set $n_i := |\pi^{-1}(u_i)|$, for $i = 1, 2, 3$, $n := \sum_{i=1}^3 n_i$, and $x := \frac{1}{n}(n_1, n_2, n_3)$. In this case, $x \in \text{int } \Delta^2$ if and only if the n_i 's satisfy triangle inequalities $n_i + n_j > n_k$, where i, j , and k are pairwise distinct. If these inequalities are satisfied, then \vec{G} admits a Hamiltonian decomposition H , which is comprised primarily (if not entirely) of 2-cycles. We plot in Figure 5 the corresponding undirected edges of G . Specifically, there are two cases: (1) If $n_1 - n_2 + n_3$ is even, then H is comprised solely of 2-cycles as shown in Figure 5a. (2) If $n_1 - n_2 + n_3$ is odd, then H is comprised of 2-cycles and a single triangle as shown in Figure 5b.

The proof of Proposition 3 relies on a reduction argument for both the graph G_n and the matrix $A(x)$: roughly speaking, we will first remove out of \vec{G}_n a number of 2-cycles, which leads to a graph \vec{G}' of smaller size. With regards to the matrix A , this reduction leads to another matrix $A' \in \mathcal{A}(S)$ with the property that $\text{diag}(A') = 0$. Finding a Hamiltonian decomposition H for \vec{G} with $\rho(H) = A$ is then reduced to finding a Hamiltonian decomposition H' for \vec{G}' with $\rho(H') = A'$. For the arguments outlined above, we need a supporting lemma stated below, whose proof is relegated to Appendix C:

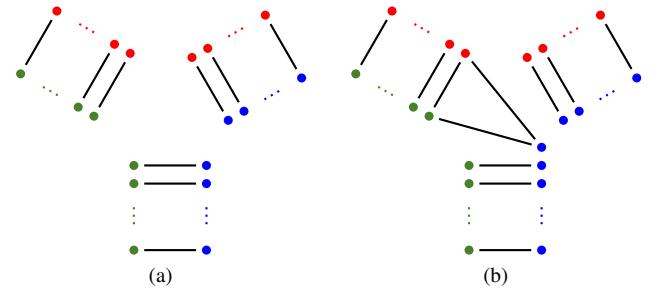


Fig. 5: Illustration of Example 1: Two complete S -multipartite graphs G with S shown in Figure 4a, where the undirected edges plotted in each subfigure give rise to a Hamiltonian decomposition H of \vec{G} . Green nodes belong to $\pi^{-1}(u_1)$, blue nodes to $\pi^{-1}(u_2)$, and red nodes to $\pi^{-1}(u_3)$.

Lemma 8. *Let n' be a nonnegative integer. Let $A' \in \mathcal{A}(S)$ be such that $n'A'$ is integer-valued, $\text{diag}(A') = 0$, and $x' := A'1$. Then, there exists a Hamiltonian decomposition H' of \vec{G}' , where $G' := M(n'x', S)$ is the complete S -multipartite graph, such that $\rho(H') = A'$ and every cycle in H' is simple.*

With the lemma above, we now establish Proposition 3:

Proof of Proposition 3. We construct a Hamiltonian decomposition H with the desired properties in two steps. We will fix x in the proof and, to simplify the notation, we omit writing the argument x for $a_{ij}(x)$, $m_{ij}(x)$, and $A(x)$.

Step 1. We claim that the following selection of 2-cycles out of \vec{G}_n is feasible:

- For every self-loop $(u_i, u_i) \in F_0$, $m_{ii} = na_{ii}$ is an even integer, and we select m_{ii} pairwise distinct nodes in $\pi^{-1}(u_i)$ that form $m_{ii}/2$ disjoint 2-cycles.
- For every $(u_i, u_j) \in F_1$, we select m_{ij} distinct nodes in $\pi^{-1}(u_i)$ and m_{ij} distinct nodes in $\pi^{-1}(u_j)$, to form m_{ij} disjoint 2-cycles (so the total number of such 2-cycles is $\sum_{(u_i, u_j) \in F_1} m_{ij}$).

The above selection is feasible because (1) G_n is a complete S -multipartite graph and, thus, there is an edge between any pair of nodes in $\pi^{-1}(u_i)$, $\pi^{-1}(u_j)$ provided that $(u_i, u_j) \in F_1$ and, (2) $A1 = x$ which implies that $\sum_{j=1}^q m_{ij} \leq n_i$ and, hence, we can always pick the required number of distinct nodes.

Let V' be the set of remaining nodes in G , i.e., V' is obtained by removing out of V the $\sum_{(u_i, u_j) \in F_1} m_{ij}$ nodes picked in step 1. If V' is the empty set, we let H be the union of the disjoint 2-cycles just exhibited. It should be clear that H is a Hamiltonian decomposition of \vec{G} . We claim that H satisfies the desired properties. To see this, let $A' := \rho(H)$ and $x' := A'1$. Then, $A' \in \mathcal{A}(S)$ and, by construction of H , $na'_{ij} = m_{ij}$. On the one hand, since V' is empty, $\|nx'\|_1 = \|nx\|_1 = n$ and, hence, the sum of the entries of A' is equal to the sum of the entries of A . On the other hand, since $m_{ij} \leq na_{ij}$, we have that $A' = \frac{1}{n}[m_{ij}]_{ij} \leq A$. It then follows that $A' = A$ and $x' = x$. Furthermore, items 1 and 2 follow from Step 1, respectively, and items 3, 4, and 5 hold trivially.

Step 2. We now assume that V' is non-empty. Let G' be the subgraph of G induced by V' . We exhibit below a Hamiltonian

decomposition H' for \vec{G}' such that the cycles in H' , together with the 2-cycles constructed in Step 1, yield a desired H . Additionally, we will show that all the cycles of length at least 3 in H' satisfy items 3, 4, and 5.

To construct the above mentioned H' , we will appeal to Lemma 8. To this end, let n'_i be the number of nodes in $\pi^{-1}(u_i) \cap V'$, i.e.,

$$n'_i = n_i - \sum_{u_j \in N(u_i)} m_{ij}.$$

Let $n' := \sum_{i=1}^q n'_i$ be the total number of nodes in G' , and

$$x' := \frac{1}{n'} [n'_1; \dots; n'_q].$$

It follows that $G' = M(n'x', S)$.

Because H has to satisfy $\rho(H) = A$ with $A\mathbf{1} = x$, H should contain na_{ij} edges from nodes in $\pi^{-1}(u_i)$ to nodes in $\pi^{-1}(u_j)$. Since m_{ij} such edges have already been accounted for by the 2-cycles created in Step 1, we need an additional n'_{ij} edges, where

$$n'_{ij} := \begin{cases} na_{ij} - m_{ij} & \text{if } (u_i, u_j) \in F, \\ 0 & \text{otherwise.} \end{cases} \quad (33)$$

Note that by (32), $n'_{ii} = 0$ for all $i = 1, \dots, q$. Correspondingly, we define a $q \times q$ matrix as follows:

$$A' := \frac{1}{n'} [n'_{ij}].$$

Because $m_{ij} = m_{ji}$ for all $(u_i, u_j) \in F$ and because $A\mathbf{1} = A^\top \mathbf{1}$, we obtain that

$$\sum_{j=1}^q n'_{ij} = \sum_{j=1}^q n'_{ji} = n'_i, \quad \forall i = 1, \dots, q.$$

Thus, $A' \in \mathcal{A}(S)$ and, by construction, $\text{diag } A' = 0$ and $A'\mathbf{1} = x'$, so A' satisfies the conditions in the statement of Lemma 8.

By Lemma 8, there exists a Hamiltonian decomposition H' of \vec{G}' such that $\rho(H') = A'$, $A'\mathbf{1} = x'$, and all cycles in H' are simple. Now, let H be the union of H' and the 2-cycles obtained in Step 1. Then,

$$\rho(H) = \frac{1}{n} [m_{ij} + n'_{ij}] = \frac{1}{n} [na_{ij}] = A,$$

where the second equality follows from (33). Moreover, since $\text{diag } A' = 0$, there is no 2-cycle in H' connecting pairs of nodes in $\pi^{-1}(u_i)$ for any $i = 1, \dots, q$. Thus, for each $i = 1, \dots, q$, H contains exactly $\frac{1}{2}m_{ii}$ disjoint two-cycles pairing m_{ii} nodes in $\pi^{-1}(u_i)$.

It now remains to show that all the cycles of length at least three in H' satisfy items 3 and 4. To do so, we first provide an upper bound on n'_i : Using items 3 and 4 of Proposition 2, we have that $na_{ij} - m_{ij} \leq 1$. Thus,

$$n'_i \leq n_i - \sum_{u_j \in N(u_i)} (na_{ij} - 1) = n_i - n_i + \deg(u_i) = \deg(u_i),$$

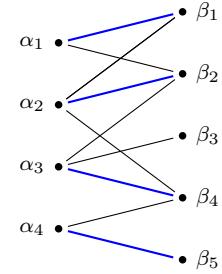


Fig. 6: A bipartite graph with $V_L = \{\alpha_1, \dots, \alpha_4\}$ and $V_R = \{\beta_1, \dots, \beta_5\}$. The blue edges form a left-perfect matching.

where $\deg(u_i)$ is the degree of u_i in S . Since

$$\sum_{i=1}^q \deg(u_i) \leq 2|F|,$$

there are at most $2|F|$ nodes in \vec{G}' . Consequently, the length of any cycle in H' is bounded above by $2|F|$ and, moreover, there exist at most $\lceil \frac{2}{3}|F| \rceil$ cycles of length three or more in H' . This completes the proof. ■

Remark 3. The fact that item 2 of the proposition provides a lower bound for the number of 2-cycles instead of an exact number can be understood as follows: The Hamiltonian decomposition H' of \vec{G}' , introduced in Step 2 of the above proof, may contain additional 2-cycles pairing nodes from $\pi^{-1}(u_i)$ to $\pi^{-1}(u_j)$ for $(u_i, u_j) \in F$. ■

E. Proof of the Main Theorem

In Subsection II-D, we dealt with the construction of a Hamiltonian decomposition in a graph \vec{G}_n sampled from a *binary* step graphon. We will now extend the result to a general step-graphon W , for which the existence of an edge between a pair of nodes is not a sure event. This will then complete the proof of the Main Theorem.

To do so, we first recall some known facts about bipartite graphs. An undirected graph $B = (V, E)$ is called *bipartite* if its node set can be written as the union of two disjoint sets $V = V_L \cup V_R$ so that there does not exist an edge between two nodes in V_L or V_R . Equivalently, a bipartite graph can be viewed as an S -multipartite graph where S is a graph with two nodes connected by a single edge. We refer to elements of V_L and V_R as *left-* and *right-nodes*, respectively. A *left-perfect matching* P in B is a set of edges so that each left-node is incident to *exactly* one edge in P , and each right-node is incident to *at most* one edge in P . See Figure 6 for an illustration. Similarly, we define a *right-perfect matching* by swapping the roles of left- and right-nodes. One can easily see that a left-perfect (resp. right-perfect) matching exists only if $|V_L| \leq |V_R|$ (resp. $|V_L| \geq |V_R|$).

Further, we denote by $B(n_1, n_2, p)$ an Erdős-Rényi random bipartite graph, with n_1 left-nodes, n_2 right-nodes, and edge probability p for all edges between left- and right-nodes.

We need the following fact:

Lemma 9. Let n_1 and n_2 be positive integers such that $\frac{1}{\kappa} \leq \frac{n_2}{n_1} \leq \kappa$, where $\kappa \geq 1$ is a constant. Let $n := n_1 + n_2$ and

$p \in (0, 1)$. Then, it holds a.a.s. that the random bipartite graph $B(n_1, n_2, p)$ is connected and contains a left-perfect (resp. right-perfect) matching if $n_2 \geq n_1$ (resp. $n_1 \geq n_2$).

The above lemma is certainly well known. For completeness of presentation, we present a proof in Appendix D.

We now return to the proof of Main Theorem. For the given step-graphon W , we fix a partition σ , and let x^* be the associated concentration vector and S be the skeleton graph. We now consider a sequence of graphs $G_n \sim W$, with $n \rightarrow \infty$. We show below that the Hamiltonian decomposition H for \vec{G}_n described in Proposition 3 exists a.a.s..

Denote by W^s the *saturation* of W : it is the binary step-graphon defined as

$$W^s(s, t) = 1 \iff W(s, t) \neq 0.$$

We similarly construct a *saturated* version of $G_n = (V_n, E_n) \sim W$, denoted by $G_n^s = (V_n, E_n^s)$, as follows: There is an edge $(v_\ell, v_k) \in E_n^s$ if and only if $(\pi(v_\ell), \pi(v_k)) \in F$. Said otherwise, the node set of G_n and G_n^s are the same, but the edges in G_n^s are obtained using the *binary* step-graphon W^s . It should be clear that $G_n \subseteq G_n^s = M(nx(G_n), S)$, where we recall that $x(G_n)$ is the empirical concentration vector of G_n defined in (6).

Let \mathcal{U} be the closed neighborhood of x^* mentioned in Remark 2. Let \mathcal{E}_0 be the event that the empirical concentration vector $x(G_n)$ of G_n belongs to \mathcal{U} . By Chebyshev's inequality, we have that

$$\mathbb{P}(\|x(G_n) - x^*\| > \epsilon) \leq \frac{c}{n^2 \epsilon^2}, \quad (34)$$

which implies that \mathcal{E}_0 is almost sure as $n \rightarrow \infty$. Thus, we can assume in the sequel that \mathcal{E}_0 is true, i.e., the analysis and computation carried out below are conditioned upon \mathcal{E}_0 .

Note that $nx(G_n)$ is always integer-valued. Since $x(G_n) \in \mathcal{U}$ by assumption, we let n be sufficiently large so that n is paired with $x(G_n)$ (see Definition 7). We can thus appeal to Proposition 2 to obtain a matrix $A(x(G_n))$, and to Proposition 3 to obtain a corresponding Hamiltonian decomposition H of \vec{G}_n^s . We now demonstrate that the same H exists a.a.s. in \vec{G}_n , up to re-labeling of the nodes of \vec{G}_n . The proof comprises two parts: In part 1, we show that the cycles in H whose lengths are greater than 2 exist a.a.s. in \vec{G}_n and, then, in part 2, we show that the 2-cycles of H do as well.

Part 1: On cycles of length greater than 2. For clarity of presentation, we denote by $\pi^s : G_n^s \rightarrow S$ the graph homomorphisms associated with G_n^s . For any path $u_1 \dots u_k$ in S , since G_n^s is complete S -multipartite, there surely exists a path $v_1 \dots v_k$ in G_n^s so that $\pi^s(v_i) = u_i$. The following result shows that such a path exist in G_n a.a.s..

Lemma 10. *Let $u_1 \dots u_k$ be a path in S . Then, it is a.a.s. that there exists a path $v_1 \dots v_k$ in G_n , with $\pi(v_i) = u_i$.*

Proof. Since the closed set \mathcal{U} is in the interior of Δ^{q-1} , there exists a $\kappa \geq 1$ such that for all $x \in \mathcal{U}$,

$$\frac{1}{\kappa} \leq \frac{x_i}{x_j} \leq \kappa, \text{ for all } 1 \leq i, j \leq q.$$

Thus, by conditioning on \mathcal{E}_0 , we have that

$$\frac{1}{\kappa} \leq \frac{x_i(G_n)}{x_j(G_n)} = \frac{|\pi^{-1}(u_i)|}{|\pi^{-1}(u_j)|} \leq \kappa, \text{ for all } 1 \leq i, j \leq q.$$

It then follows that the subgraphs of G_n induced by $\pi^{-1}(u_i) \cup \pi^{-1}(u_{i+1})$ are bipartite and satisfy the hypothesis of Lemma 9, for $1 \leq i \leq k-1$. Hence, it is a.a.s. that all of these bipartite graphs are connected. We now pick an arbitrary node $v_1 \in \pi^{-1}(u_1)$; by the above arguments, we can find $v_2 \in \pi^{-1}(u_2)$ so that $(v_1, v_2) \in G_n$ a.a.s.. Iterating this procedure, we obtain the path in G_n sought. ■

Now, let D_1, \dots, D_m be the cycles in H whose lengths are greater than 2, and C_1, \dots, C_m be the corresponding undirected cycles in G_n^s . From items 3 and 4 of Proposition 3, the number m of these cycles, as well as their lengths, are each uniformly bounded above by constants *independent* of n .

Let \mathcal{E}_1 be the event that the cycles C_1, \dots, C_m exist in G_n ; more precisely, it is the event that there exist disjoint cycles C'_i in G_n such that $\pi(C'_i) = \pi^s(C_i)$ for all $i = 1, \dots, m$. We have the following lemma:

Lemma 11. *The event \mathcal{E}_1 is true a.a.s..*

Proof. Let \mathcal{E}_{11} be the event that there exists a cycle $C'_1 \in G_n$ with $\pi(C'_1) = \pi^s(C_1)$. We show that \mathcal{E}_{11} holds a.a.s.. To start, we write explicitly $\pi^s(C_1) = u_1 \dots u_k u_1$. Since C_1 is simple, $u_1 \dots u_k u_1$ is a *cycle* in S . By Lemma 10, there exist a.a.s. nodes $v_i \in \pi^{-1}(u_i)$, for $1 \leq i \leq k-1$, such that $v_1 \dots v_{k-1}$ is a path in G_n .

In order to obtain the cycle C'_1 , it remains to exhibit a node $v_k \in \pi^{-1}(u_k)$ that is connected to both v_1 and v_{k-1} in G_n . We claim that such a node exists with probability at least

$$1 - (1 - p_{1k} p_{k-1,k})^{|\pi^{-1}(u_k)|} \quad (35)$$

where $p_{ij} > 0$ is the value of the step-graphon W over the rectangle $[\sigma_{i-1}, \sigma_i] \times [\sigma_{j-1}, \sigma_j]$. The claim holds because the probability that *no* node of $\pi^{-1}(u_k)$ connects to *both* v_1 and v_{k-1} is given by $(1 - p_{1k} p_{k-1,k})^{|\pi^{-1}(u_k)|}$. Thus, the probability of the complementary event is given by (35).

Next, recall that $\underline{a} > 0$ is the uniform lower bound for the nonzero entries of $A(x)$, for all $x \in \mathcal{U}^*$, introduced in item 5 of Proposition 2. Because $x(G_n) = A(x(G_n))\mathbf{1}$, every entry of $x(G_n)$ is bounded below by \underline{a} as well, so

$$|\pi^{-1}(u_i)| \geq \underline{a}n, \text{ for all } i = 1, \dots, q. \quad (36)$$

Thus, the expression (35) can be lower bounded by

$$1 - (1 - p_{1k} p_{k-1,k})^{|\pi^{-1}(u_k)|} \geq 1 - (1 - p^2)^{\underline{a}n},$$

where $p := \min\{p_{ij} \mid (u_i, u_j) \in F\} > 0$. Note that the right-hand-side of the above equation converges to 1 as $n \rightarrow \infty$, so \mathcal{E}_{11} is true a.a.s..

Let $n' := n - |C_1|$. Conditioning on the event \mathcal{E}_{11} , we let $G'_{n'}$ be the subgraph of G_n induced by the nodes not in C'_1 . Similarly as above, we have that there is a cycle C'_2 , with $\pi(C'_2) = \pi(C_2)$, in $G'_{n'}$ a.a.s. (note that $n \rightarrow \infty$ implies $n' \rightarrow \infty$). Iterating this argument for finitely many steps, we have that \mathcal{E}_1 is true a.a.s.. ■

In the sequel, we condition on the event \mathcal{E}_1 and let D'_1, \dots, D'_m be the directed cycles in \vec{G}_n corresponding to D_1, \dots, D_m in H of \vec{G}_n^s .

Part 2: On 2-cycles. Let $n' := n - \sum_{i=1}^m |C'_i|$, and $G'_{n'}$ be the subgraph of G_n induced by the nodes that do *not* belong to any of the above cycles C'_i , and $G'_{n'}^s$ be its saturation. By removing the cycles D_i out of H , we obtain a Hamiltonian decomposition H' for $\vec{G}_{n'}^s$, which is comprised only of 2-cycles. It now suffices to show that H' appears, up to relabeling, in $\vec{G}'_{n'}$ *a.a.s.*

Let $V_{ij} \subset \pi^{-1}(u_i)$ be the set of nodes paired to nodes in $\pi^{-1}(u_j)$ by H' in $\vec{G}_{n'}^s$. Since H' is a Hamiltonian decomposition, $\pi^{-1}(u_i)$ can be expressed as the disjoint union of the V_{ij} 's, for u_j such that $(u_i, u_j) \in F$. By items 1 and 2 of Proposition 3, the cardinality of V_{ij} , which is the same as the cardinality of V_{ji} , is at least $m_{ij} := n \min\{a_{ij}, a_{ji}\}$. Because the nonzero a_{ij} 's are bounded below by \underline{a} by item 5 of Proposition 2, we have that $m_{ij} \geq \underline{a}n$.

Suppose that u_i has a self-loop; then, we let K_i be the subgraph of $G'_{n'}$ induced by the nodes V_{ii} . The graph K_i is an Erdős-Rényi graph with parameter $p_{ii} > 0$ and, by item 1 of Proposition 3, $|V_{ii}| = m_{ii}$ is an even integer. Since $n \rightarrow \infty$ implies that $m_{ii} \rightarrow \infty$, K_i has a perfect matching *a.a.s.*. This holds because one can split the node set V_{ii} into two disjoint subsets of equal cardinality and apply Lemma 9. In other words, it is *a.a.s.* that there are $m_{ii}/2$, for $i = 1, \dots, q$, disjoint 2-cycles in $\vec{G}'_{n'}$ pairing nodes in $\pi^{-1}(u_i)$.

Suppose that (u_i, u_j) is an edge between two distinct nodes; then, we let $B_{ij} := B(|V_{ij}|, |V_{ji}|, p_{ij})$ be the bipartite graph in $G'_{n'}$ induced by $V_{ij} \cup V_{ji}$ (recall from above that $|V_{ij}| = |V_{ji}|$). Let \mathcal{E}_{ij} be the event that B_{ij} has a perfect matching. Since $m_{ij} \geq \underline{a}n$, by Lemma 9, the event \mathcal{E}_{ij} holds *a.a.s.* and, hence, it is *a.a.s.* that there are $|V_{ij}|$ disjoint 2-cycles in $\vec{G}'_{n'}$ pairing nodes from V_{ij} to V_{ji} .

Since there are finitely many edges in S , by the above arguments, we conclude that H' appears in $\vec{G}'_{n'}$ *a.a.s.*. This completes the proof. \blacksquare

III. CONCLUSIONS

Hamiltonian decompositions underlie a wide range of structural properties of control systems, such as stability and ensemble controllability. We say that a graphon W satisfies the H -property if graphs $G_n \sim W$ have a Hamiltonian decomposition almost surely. In a series of papers, of which this is the second, we exhibited necessary and sufficient conditions for the H -property to hold for the class of step-graphons. These conditions are geometric and revealed the fact that H -property depends only on concentration vector and skeleton graph of W . When these two objects are given, one can reconstruct a step-graphon modulo the exact value of W on its support, thus giving rise to an equivalence relation on the space of step-graphons. We showed that the H -property is essentially a “zero-one” property of the equivalence classes. The case of general graphons will be addressed in future work.

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APPENDIX A ANALYSIS AND PROOF OF PROPOSITION 1

We first have some preliminaries about *refinements* of partitions: given a partition σ , a *refinement* σ' of σ , denoted by $\sigma \prec \sigma'$, is any sequence that has σ as a proper subsequence. For example, $\sigma' = (0, 1/2, 3/4, 1)$ is a refinement of $\sigma = (0, 1/2, 1)$. Given a step-graphon W , if σ is a partition for W , then so is σ' .

We say that σ' is a *one-step refinement* of σ if it is a refinement with $|\sigma'| = |\sigma| + 1$. Any refinement of σ can be obtained by iterating one-step refinements. To fix ideas, and without loss of generality, we consider the refinement of $\sigma = (\sigma_0, \dots, \sigma_q, \sigma_*)$ to $\sigma' = (\sigma_0, \dots, \sigma_q, \sigma_{q+1}, \sigma_*)$ with $\sigma_q < \sigma_{q+1} < \sigma_*$. If $S = (U, F)$, then $S' = (U', F')$, the skeleton graph of W for σ' , is given by

$$\begin{cases} U' = U \cup \{u_{q+1}\}, \\ F' = F \cup \{(u_i, u_{q+1}) \mid (u_i, u_q) \in F\} \\ \quad \cup \{(u_{q+1}, u_{q+1}) \mid (u_q, u_q) \in F\}. \end{cases} \quad (37)$$

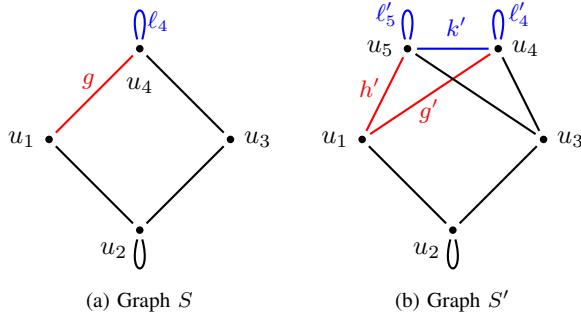


Fig. 7: The graph S' on the right is obtained from the left S via a one-step refinement. The node u_4 in S is split into u_4 and u_5 . Because $g = (u_1, u_4)$ is an edge in S , there exist two edges $g' = (u_1, u_4)$ and $h' = (u_1, u_5)$ in S' . Because u_4 has a self-loop ℓ_4 in S , both u_4 and u_5 have self-loops in S' , denoted by ℓ_4' and ℓ_5' , respectively. In addition, we have the edge $k' = (u_4, u_5)$ in S' .

In essence, the node u_{q+1} is a copy of the node u_q . If there is a loop (u_q, u_q) in F , then u_q and u_{q+1} are also connected and each has a self-loop. See Figure 7 for illustration. We say that a one-step refinement *splits a node* (here, u_q).

We now prove Proposition 1:

Proof of Proposition 1. Let σ and σ' be as given in the statement of the proposition. It should be clear that there exists another partition σ'' which is a refinement of both σ' and σ and that σ'' can be obtained via a sequence of one-step refinements starting with either σ' or σ . Thus, combining the arguments at the beginning of the section, we can assume, without loss of generality, that σ' is a *one-step* refinement of σ obtained by splitting the node $u_1 \in U$.

Let x^* and x'^* be the concentration vectors for σ and σ' , S and S' be the corresponding skeleton graphs, and Z and Z' be the corresponding incidence matrices. Note that Z' has one more row than Z does due to the addition of the new node u_{q+1} ; here, we let the last row of Z' correspond to that node. It should be clear that Z' contains Z as a submatrix. For clarity of presentation, we use f (resp. f') to denote edges of S (resp. S'). Since the graph S can be realized as a subgraph of S' in a natural way, we will write on occasion $f' \in F$ if f' is an edge of S .

We now prove the invariance of each item listed in the statement of Proposition 1 under one-step refinements. The proofs of the first two items are direct consequence of the definition of one-step refinement.

Proof for item 1). If S is connected, then from (37) we obtain that there exists a path from any node $u_i \in F$ to the new node u_{q+1} , so S' is also connected. Reciprocally, assume that S has at least two connected components. Then, the node u_{q+1} obtained by splitting u_q will only be connected to nodes in the same component as u_q by definition of F' .

Proof for item 2). If S has an odd cycle, then so does S' by (37). Reciprocally, we assume that S is lacking an odd cycle. We show that S' has no odd cycle. Suppose, to the contrary, that it does. The cycle must then contain the node u_{q+1} . Replacing u_{q+1} with u_q yields a closed walk of odd length in S . Since a closed walk can be decomposed edge-wise

into a union of cycles and since the length of the walk is the sum of the lengths of the constituent cycles, there must exist an odd cycle in S , which is a contradiction.

Proof for item 3). We prove each direction of the statement separately:

Part 1: $x \in \mathcal{X}(S) \Rightarrow x' \in \mathcal{X}(S')$ ($x \in \text{int } \mathcal{X}(S) \Rightarrow x' \in \text{int } \mathcal{X}(S')$). For ease of presentation, we let z_f (resp. $z'_{f'}$) be the edge of Z (resp. Z') corresponding to the element $f \in F$ (resp. $f' \in F'$), and $z_{f,i}$ be the i th entry of z_f . Because $\mathcal{X}(S)$ is the convex hull of the columns of Z , there exist coefficients $c_f \geq 0$, for $f \in F$, such that $x = \sum_{f \in F} c_f z_f$. If, further, $x \in \text{int } \mathcal{X}(S)$, then these coefficients can be chosen to be strictly positive. We will use c_f to construct $c'_{f'} \geq 0$, for $f' \in F'$, such that

$$x' = \sum_{f' \in F'} c'_{f'} z'_{f'} \quad (38)$$

and show that $x' \in \text{int } \mathcal{X}(S')$ if $x \in \text{int } \mathcal{X}(S)$.

To proceed, let F_{u_q} be the set of edges incident to node u_q in S . Similarly, let F'_{u_q} and $F'_{u_{q+1}}$ be the sets of edges incident to u_q and u_{q+1} in S' , respectively. The coefficients $c'_{f'}$ are defined as follows:

- a) If $f' \notin F'_{u_q} \cup F'_{u_{q+1}}$, then $f' \in F$. Let $c'_{f'} := c_f$.
- b) If $f' \in F'_{u_q}$ and $f' \neq (u_q, u_{q+1})$, then $f' \in F$. Let $c'_{f'} := \frac{\sigma_{q+1} - \sigma_q}{\sigma_* - \sigma_q} c_f$.
- c) If $f' = (u_q, u_{q+1})$ and $u_i \neq u_q$, then we pick the $f \in F$ such that

$$f = \begin{cases} (u_i, u_q) & \text{if } u_i \neq u_{q+1}, \\ (u_q, u_q) & \text{if } u_i = u_{q+1}. \end{cases}$$

$$\text{Let } c'_{f'} := \frac{\sigma_* - \sigma_{q+1}}{\sigma_* - \sigma_q} c_f.$$

- d) If $f' = (u_q, u_{q+1})$, then let $c'_{f'} := 0$.

With the coefficients as above, we prove entry-wise that (38) holds. First, note that because we obtain S' by splitting the last node u_q of S , the i th entry of x' , for $1 \leq i \leq q-1$, is equal to x_i , so $x'_i = x_i = (\sigma_i - \sigma_{i-1})$. For the i th entry of the right hand side of (38), we consider two cases:

Case 1: u_i is not incident to u_q in S . In this case, u_i is not incident to either u_q or u_{q+1} in S' . Consequently, $F'_{u_i} = F_{u_i}$ and $z'_{f,i} = z_{f,i}$ for all $f \in F_{u_i}$. Furthermore, by item (a), $c'_{f'} = c_f$ for any $f \in F_{u_i}$. Thus, the i th entry of the right hand side of (38) is given by

$$\sum_{f' \in F'_{u_i}} c'_{f'} z'_{f',i} = \sum_{f \in F_{u_i}} c_f z_{f,i} = x_i = \sigma_i - \sigma_{i-1}.$$

Case 2: u_i is incident to u_q in S . In this case, u_i is incident to both u_q and u_{q+1} in S' . Let $g' := (u_i, u_q)$ and $h' := (u_i, u_{q+1})$ be the corresponding edges in S' , see Fig. 7 for an illustration. Then, the i th entry of the right hand side of (38) is given by

$$\begin{aligned} \sum_{f' \in F'_{u_i}} c'_{f'} z'_{f',i} &= c'_{g'} z'_{g',i} + c'_{h'} z'_{h',i} \\ &+ \sum_{f' \in F'_{u_i} - \{g', h'\}} c'_{f'} z'_{f',i}. \end{aligned} \quad (39)$$

By items (b) and (c),

$$c'_{g'} = \frac{\sigma_{q+1} - \sigma_q}{\sigma_* - \sigma_q} c_{g'} \quad \text{and} \quad c'_{h'} = \frac{\sigma_* - \sigma_{q+1}}{\sigma_* - \sigma_q} c_{h'}.$$

Also, note that

$$z'_{g',i} = z'_{h',i} = z_{g',i} = \frac{1}{2}.$$

Thus, the sum of the first two terms on the right hand side of (39) is $c_{g'} z_{g',i}$. For the last term, note that $F'_{u_i} - \{g', h'\} = F_{u_i} - \{g'\}$. Also, by item (a) and the fact that $z'_{f,i} = z_{f,i}$ for any $f \in F_{u_i}$,

$$\sum_{f' \in F'_{u_i} - \{g', h'\}} c'_{f'} z'_{f',i} = \sum_{f \in F_{u_i} - \{g'\}} c_f z_{f,i}.$$

Combining the above arguments, we have that the right hand side of (39) is given by

$$\sum_{f \in F_{u_i}} c_f z_{f,i} = x_i = \sigma_i - \sigma_{i-1}.$$

Next, the q th entry of x' is $(\sigma_{q+1} - \sigma_q)$ and the q th entry of the right hand side of (38) is

$$\begin{aligned} \sum_{f' \in F'_{u_q}} c'_{f'} z'_{f',q} &= \frac{\sigma_{q+1} - \sigma_q}{\sigma_* - \sigma_q} \sum_{f \in F_{u_q}} c_f z_{f,q} \\ &= \frac{\sigma_{q+1} - \sigma_q}{\sigma_* - \sigma_q} x_q = \sigma_{q+1} - \sigma_q, \end{aligned}$$

where the first equality follows from the fact that

$$F'_{u_q} = F_{u_q} \cup \{(u_q, u_{q+1}) \text{ if } (u_q, u_q) \in F\},$$

items (b) and (d), and the last equality follows from the fact that $x_q = \sigma_* - \sigma_q$.

The last (i.e., $(q+1)$ th) entry of x' is $(\sigma_* - \sigma_{q+1})$. The last entry of the right hand side of (38) is given by

$$\begin{aligned} \sum_{f' \in F'_{u_{q+1}}} c'_{f'} z'_{f',q+1} &= \frac{\sigma_* - \sigma_{q+1}}{\sigma_* - \sigma_q} \sum_{f \in F_{u_q}} c_f z_{f,q} \\ &= \frac{\sigma_* - \sigma_{q+1}}{\sigma_* - \sigma_q} x_q = \sigma_* - \sigma_{q+1}, \end{aligned}$$

where the first equality follows from item (c) above. We have thus shown that (38) holds. In particular, since $c'_{f'}$ are nonnegative by construction, (38) implies that $x' \in \mathcal{X}(S')$.

It now remains to show that if $x \in \text{int } \mathcal{X}(S)$, then $x' \in \text{int } \mathcal{X}(S')$. Assuming $x \in \text{int } \mathcal{X}(S)$, if u_q does not have a self-loop in S , then the edge (u_q, u_{q+1}) does not exist in S' , so by items (a), (b), and (c), all coefficients $c'_{f'}$ are positive, which implies that $x' \in \text{int } \mathcal{X}(S')$.

We now assume that u_q has a self-loop in S . Then, $k' := (u_q, u_{q+1})$ is an edge in S' (see Fig. 7 for an illustration), and thus $c'_{k'} = 0$ per item (d) above. In this case, both u_q and u_{q+1} have self-loops in S' . Denote these two self-loops by $\ell'_q := (u_q, u_q)$ and $\ell'_{q+1} := (u_{q+1}, u_{q+1})$. By (4), we have that

$$z'_{k'} = \frac{1}{2}(z'_{\ell'_q} + z'_{\ell'_{q+1}}).$$

Since $c'_{\ell'_q}$ and $c'_{\ell'_{q+1}}$ are positive, there exists an $\epsilon > 0$ such

that $\epsilon < c'_{\ell'_q}$ and $\epsilon < c'_{\ell'_{q+1}}$. It then follows that

$$\begin{aligned} c'_{\ell'_q} z'_{\ell'_q} + c'_{\ell'_{q+1}} z'_{\ell'_{q+1}} &= 2\epsilon z'_{k'} + (c'_{\ell'_q} - \epsilon) z'_{\ell'_q} \\ &\quad + (c'_{\ell'_{q+1}} - \epsilon) z'_{\ell'_{q+1}}. \end{aligned} \quad (40)$$

Plugging in (38) the relation (40) shows that x' can be written as a convex combination of the $z'_{f'}$, for $f' \in F'$, with all positive coefficients, and thus $x' \in \text{int } \mathcal{X}(S')$.

Part 2: $x' \in \mathcal{X}(S') \Rightarrow x \in \mathcal{X}(S)$ ($x' \in \text{int } \mathcal{X}(S') \Rightarrow x \in \text{int } \mathcal{X}(S)$). Because $x' \in \mathcal{X}(S')$ (resp. $x' \in \text{int } \mathcal{X}(S')$), we can write $x' = \sum_{f' \in F'} c'_{f'} z'_{f'}$, with $c'_{f'} \geq 0$ (resp. $c'_{f'} > 0$), for all $f' \in F'$. We will use $c'_{f'}$ to construct c_f , for $f \in F$, so that

$$x = \sum_{f \in F} c_f z_f. \quad (41)$$

To this end, we define c_f as follows:

- e) If f is not incident to u_q in S , then let $c_f := c'_{f'}$.
- f) If $f = (u_i, u_q)$ and $u_i \neq u_q$, then $g' := (u_i, u_q)$ and $h' := (u_i, u_{q+1})$ are edges in S' , and let $c_f := c'_{g'} + c'_{h'}$.
- g) If $f = (u_q, u_q)$, then $k' := (u_q, u_{q+1})$, $\ell'_q := (u_q, u_q)$, and $\ell'_{q+1} := (u_{q+1}, u_{q+1})$ are edges in S' , and let $c_f := c'_{k'} + c'_{\ell'_q} + c'_{\ell'_{q+1}}$.

Note that all the coefficients c_f , for $f \in F$, defined above are nonnegative. Further, if all the $c'_{f'}$ are positive, i.e., $x' \in \text{int } \mathcal{X}(S')$, then the c_f are positive as well, which implies $x \in \text{int } \mathcal{X}(S)$ provided that (41) holds.

We now show that the coefficients given above are such that (41) indeed holds. We do so by checking that (41) holds for each entry.

For the i th entry, with $1 \leq i < q$, the left hand side of (41) is $x_i = (\sigma_i - \sigma_{i-1})$. For the right-hand side, if (u_i, u_q) is an edge in S , then g' and h' , as defined item (f), are two edges in S' and, consequently, $F'_{u_i} = F_{u_i} \cup \{h'\}$. Note that $z_{f,i} = z'_{f,i}$ for all $f \in F_{u_i}$ and

$$z'_{g',i} = z'_{g',i} = z'_{h',i} = \frac{1}{2}.$$

Thus, by items (e) and (f), we have that

$$\begin{aligned} \sum_{f \in F} c_f z_{f,i} &= c'_{g'} z_{g',i} + \sum_{f \in F_{u_i} - \{g'\}} c_f z_{f,i} \\ &= c'_{g'} z'_{g',i} + c'_{h'} z'_{h',i} + \sum_{f' \in F'_{u_i} - \{g', h'\}} c'_{f'} z'_{f',i} \\ &= \sum_{f' \in F'_{u_i}} c'_{f'} z'_{f',i} = x'_i = (\sigma_i - \sigma_{i-1}). \end{aligned}$$

Finally, for the last entry, i.e., the q th entry, the left hand side of (41) is $x_q = \sigma_* - \sigma_q$. For the right hand side of (41), we let $\ell'_q := (u_q, u_q)$ be the loop on u_q (if it exists in S) and thus have that

$$\sum_{f \in F_{u_q}} c_f z_{f,q} = c_{\ell'_q} z_{\ell'_q,q} + \sum_{f \in F_{u_q} - \{\ell'_q\}} c_f z_{f,q}. \quad (42)$$

Let k' , ℓ'_q , and ℓ'_{q+1} be the three edges in S' as defined in item (g). Note that

$$z_{\ell'_q,q} = z'_{\ell'_q,q} = z'_{\ell'_{q+1},q+1} = 2z'_{k',q} = 2z'_{k',q+1} = 1.$$

For the first term of (42), using item (g) and the above relations, we obtain

$$c_{\ell_q} z_{\ell_q, q} = c'_{\ell'_q} z'_{\ell'_q, q} + c'_{\ell'_{q+1}} z'_{\ell'_{q+1}, q+1} + c'_{k'} z'_{k', q} + c'_{k'} z'_{k', q+1}. \quad (43)$$

For each addend in the second term of (42), the edge $g = (u_i, u_q)$ in S , for some $u_i \neq u_q$, has two corresponding edges in S' , namely $g' = (u_i, u_q)$ and $h' = (u_i, u_{q+1})$. Note that

$$z_{g, q} = z'_{g', q} = z'_{h', q+1} = \frac{1}{2}.$$

Then, by item (f),

$$c_g z_{g, q} = c'_{g'} z'_{g', q} + c'_{h'} z'_{h', q+1}. \quad (44)$$

Combining (43) and (44), we obtain that

$$\begin{aligned} \sum_{f \in F_{u_q}} c_f z_{f, 1} &= \sum_{f' \in F'_{u_q}} c'_{f'} z'_{f', q} + \sum_{f' \in F'_{u_{q+1}}} c'_{f'} z'_{f', q+1} \\ &= x'_q + x'_{q+1} = (\sigma_{q+1} - \sigma_q) + (\sigma_* - \sigma_{q+1}) \\ &= \sigma_* - \sigma_q. \end{aligned}$$

This concludes the proof. \blacksquare

APPENDIX B PROOF OF LEMMA 7

Proof of item 1. From the definitions of $\tau'_i(x)$, we have that

$$-1 \leq n(\tau_0(x) - \tau'_0(x)) = n(\tau'_1(x) - \tau_1(x)) \leq 1. \quad (45)$$

If $\tau_{0,i}(x) = 0$, then $\tau'_{0,i}(x) = 0$, where $\tau_{0,i}(x)$ is the i th entry of the vector $\tau_0(x)$. Otherwise, by the definition of the incidence matrix (4) and by (15) and (16), we have that $\tau_{0,i}(x) \geq \alpha$. For the latter case, by (45) and the hypothesis on n in the statement of Proposition 2, we have that

$$\tau'_{0,i}(x) \geq \tau_{0,i}(x) - \frac{1}{n} \geq (1 - \epsilon/8)\alpha > 0. \quad (46)$$

It then follows that

$$\text{supp } \tau'_0(x) = \text{supp } \tau_0(x). \quad (47)$$

Similarly, for $\tau_1(x)$, using (4), (15), and (16), we have that $\tau_1(x) \geq \frac{1}{2}\alpha\mathbf{1}$. Then, again, by (45) and the hypothesis on n ,

$$\tau'_1(x) \geq \tau_1(x) - \frac{1}{n}\mathbf{1} \geq \frac{1}{2}(1 - \epsilon/4)\alpha\mathbf{1} > 0,$$

from which we conclude that

$$\text{supp } \tau'_1(x) = \text{supp } \tau_1(x) = \{1, \dots, q\}.$$

This concludes the proof of the first item.

Proof of item 2. By (46), (47), and Lemma 4, the ratio n'_0/n is uniformly lower bounded by

$$\begin{aligned} \frac{n'_0}{n} &= \|\tau'_0(x)\|_1 \geq (1 - \epsilon/8)\alpha |\text{supp } \tau'_0(x)| \\ &= (1 - \epsilon/8)\alpha |F_0|. \end{aligned} \quad (48)$$

To obtain a lower bound for n'_1/n , we let

$$\eta := \tau_1(x) - \tau'_1(x).$$

From (45) and the hypothesis on n , we have that

$$\|\eta\|_1 \leq \frac{q}{n} \leq \frac{1}{8}q\alpha\epsilon \leq \frac{1}{8}q\alpha.$$

It then follows that

$$\frac{n'_1}{n} \geq \|\tau_1(x)\|_1 - \|\eta\|_1 \geq \frac{1}{2}q\alpha - \frac{1}{8}q\alpha = \frac{3}{8}q\alpha. \quad (49)$$

This concludes the proof of the second item. \blacksquare

Proof of item 3. By (45) and the hypothesis on n , η as introduced above satisfies

$$\|\eta\|_\infty \leq 1/n < \alpha\epsilon/8 < \epsilon.$$

Because $\bar{\tau}'_1(x) = \overline{\tau_1(x) + \eta}$ and $\|\eta\|_\infty < \epsilon$, we have that $\bar{\tau}'_1(x) \in \mathcal{V}$ by Lemma 6. \blacksquare

APPENDIX C PROOF OF LEMMA 8

The proof is carried out by induction on n' . If $n' = 0$, then G' is the empty graph and there is nothing to prove. For the inductive step, we set $n' > 0$ and assume that the lemma holds for all $n'' < n'$, and prove it for n' .

To proceed, we use A' to turn \vec{S} into a weighted digraph on q nodes: we assign to edge $u_i u_j$ the weight a'_{ij} . Then, \vec{S} is a balanced graph, i.e.,

$$\sum_{u_j \in N_-(u_i)} a'_{ij} = \sum_{u_j \in N_+(u_i)} a'_{ji}, \quad \forall i = 1, \dots, q, \quad (50)$$

where we recall $N_-(u_i)$ and $N_+(u_i)$ are the sets of out-neighbors and in-neighbors of u_i , respectively, in \vec{S} .

Let \vec{S}' be the subgraph of \vec{S} induced by the nodes incident to the edges with nonzero weights. Then, \vec{S}' has at least one cycle. To see this, note that if \vec{S}' is acyclic, then by relabeling the nodes, the matrix A' is upper-triangular and, from the hypothesis, $\text{diag } A' = 0$. It follows that the only *nonnegative* solution $\{a'_{ij}\}$ to (50) is that all the a'_{ij} are zero, which contradicts the fact that A' is nonzero.

Since \vec{S}' is a subgraph of \vec{S} , any cycle of \vec{S}' is also a cycle of \vec{S} ; denote such cycle by $D_S := u_{i_1} \dots u_{i_k} u_{i_1}$. By construction, the weights of the edges in the cycle are positive. It thus follows from $A' \mathbf{1} = x'$ that the entries x'_{i_j} , for $j = 1, \dots, k$, are positive; together with the fact that $G' = M(n'x', S)$, it implies that the sets $\pi^{-1}(u_{i_j}) \in G'$, for $j = 1, \dots, k$, are non-empty. We next pick a node v_j from each $\pi^{-1}(u_{i_j})$. Since the nodes u_{i_1}, \dots, u_{i_k} are pairwise distinct, so are the nodes v_1, \dots, v_k . Also, since G' is a complete S -multipartite graph, $D_G := v_1 \dots v_k v_1$ is a cycle in \vec{G}' . Moreover, by construction, $\pi(D_G) = D_S$ and, hence, D_G is simple.

We let G'' be the graph obtained by removing from G' the k nodes v_1, \dots, v_k , and the edges incident to them. Then, G'' is a complete S -multipartite graph on $n'' := n' - k$ nodes. Define

$$x'' := \frac{1}{n''} (n'x' - \sum_{j=1}^k e_{i_j}),$$

where $\{e_1, \dots, e_q\}$ is the canonical basis of \mathbb{R}^q . Note that $x'' \geq 0$; indeed, $n'x'$ is integer valued and x'_{i_j} , for $j = 1, \dots, k$,

are positive. We can then write $G'' = M(n''x'', S)$. Correspondingly, we define a $q \times q$ matrix A'' as follows:

$$A'' := \frac{1}{n''} (n' A' - \sum_{j=1}^{k-1} e_{i_j} e_{i_{j+1}}^\top - e_{i_k} e_{i_1}^\top).$$

In words, to obtain $n''A''$, we decrease the ij th entry of $n'A'$, which is a positive integer, by one when the edge $u_i u_j$ is used in the cycle D_S and keep the other entries unchanged.

By construction, we have that $A'' \in \mathcal{A}(S)$, with $A''\mathbf{1} = x''$, and $n''A''$ is integer-valued. Because $n'' < n'$, we can appeal to the induction hypothesis and exhibit a Hamiltonian decomposition H'' of \vec{G}'' such that $\rho(H'') = A''$ and every cycle in H'' is simple. It is clear that adding the simple cycle D_G to H'' yields a Hamiltonian decomposition H' of \vec{G}' with desired properties. This completes the proof. ■



Xudong Chen is an Associate Professor in the Department of Electrical and Systems Engineering at Washington University in St. Louis. He obtained the B.S. degree in Electronic Engineering from Tsinghua University, Beijing, China, in 2009, and the Ph.D. degree in Electrical Engineering from Harvard University, Cambridge, Massachusetts, in 2014. He is an awardee of the 2020 Air Force Young Investigator Program, a recipient of the 2021 NSF Career Award, the recipient of the 2021 Donald P. Eckman Award, and the recipient of the 2023 A.V. Balakrishnan Early Career Award. His current research interests are in the area of control theory, stochastic processes, optimization, network science, and game theory.

APPENDIX D PROOF OF LEMMA 9

1. Proof that $B(n_1, n_2, p)$ has a left-perfect matching a.a.s.. The proof of this part relies on the following statement, which is a consequence of a stronger result of Erdős and Rényi [19]: For $p \in (0, 1)$ a constant, the random bipartite graph $B(m, m, p)$ contains a perfect matching a.a.s.. Now, without loss of generality, we assume that $n_1 \leq n_2$ and let $B(n_1, n_1, p)$ be a subgraph of $B(n_1, n_2, p)$. Since n_2/n_1 is bounded above by a constant κ , $n \rightarrow \infty$ implies that $n_1 \rightarrow \infty$. Since $B(n_1, n_1, p)$ has a (left-)perfect matching a.a.s., so does $B(n_1, n_2, p)$.

2. Proof that $B(n_1, n_2, p)$ is connected a.a.s.. It is well known (see, e.g., [20, Exercise 4.3.7]) that $B(m, m, p)$ is connected a.a.s.. We now extend the result to the general case where n_1 is not necessarily equal to n_2 . Again, we can assume without loss of generality that $n_1 \leq n_2$. Let $V_L = \{\alpha_1, \dots, \alpha_{n_1}\}$ and $V_R = \{\beta_1, \dots, \beta_{n_2}\}$ be the left- and right-node sets of B . Because $n_2/n_1 \leq \kappa$, we can choose κ subsets $V_{R,i} \subseteq V_R$, so that $|V_{R,i}| = n_1$ and $\cup_{i=1}^{\kappa} V_{R,i} = V_R$.

Denote by \mathcal{E}_i the event that the subgraph B_i of $B := B(n_1, n_2, p)$ induced by V_L and $V_{R,i}$ is disconnected and by \mathcal{E} the event that B is disconnected. Note that if every B_i is connected, then so is B . Conversely, we have that $\mathcal{E} \subseteq \cup_{i=1}^{\kappa} \mathcal{E}_i$.

Note that $B_i = (n_1, n_1, p)$ and, as argued above, $n_1 \rightarrow \infty$ as $n \rightarrow \infty$. Since B_i is connected a.a.s., $\lim_{n \rightarrow \infty} \mathbb{P}(\mathcal{E}_i) = 0$ and, hence, $\lim_{n \rightarrow \infty} \mathbb{P}(\mathcal{E}) \leq \lim_{n \rightarrow \infty} \sum_{i=1}^{\kappa} \mathbb{P}(\mathcal{E}_i) = 0$. ■



M.-A. Belabbas obtained his PhD degree in applied mathematics from Harvard University, Cambridge, Massachusetts, and his undergraduate degree from École Centrale Paris, France, and Université Catholique de Louvain, Belgium. He is currently an associate professor in the Electrical and Computer Engineering department at the University of Illinois, Urbana-Champaign and at the Coordinated Science Laboratory. His research interests are in Networked Control System, Stochastic Control and Geometric control theory.