

# Sampling for Remote Estimation of the Wiener Process over an Unreliable Channel

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In this paper, we study a sampling problem where a source takes samples from a Wiener process and transmits them through a wireless channel to a remote estimator. Due to channel fading, interference, and potential collisions, the packet transmissions are unreliable and could take random time durations. Our objective is to devise an optimal causal sampling policy that minimizes the long-term average mean square estimation error. This optimal sampling problem is a recursive optimal stopping problem, which is generally quite difficult to solve. However, we prove that the optimal sampling strategy is, in fact, a simple threshold policy where a new sample is taken whenever the instantaneous estimation error exceeds a threshold. This threshold remains a constant value that does not vary over time. By exploring the structure properties of the recursive optimal stopping problem, a low-complexity iterative algorithm is developed to compute the optimal threshold. This work generalizes previous research by incorporating both transmission errors and random transmission times into remote estimation. Numerical simulations are provided to compare our optimal policy with the zero-wait and age-optimal policies.

CCS Concepts: • Networks  $\rightarrow$  Network performance evaluation; *Network performance modeling*; • Mathematics of computing  $\rightarrow$  Information theory.

Additional Key Words and Phrases: Remote estimation, unreliable channel, optimal multiple stopping times ACM Reference Format:

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## 1 INTRODUCTION

Several applications in real-time control of systems involving sensor networks, such as autonomous driving, military networks, intelligent manufacturing, etc., involve sampling and remote estimation of information. For example, in military systems, status information about the instantaneous speed and position of the vehicles, channel conditions, and targets changes over time. In order to ensure that the system performs efficiently, reliably, and safely, the controller(s) has to obtain accurate estimates of the current status of the system from nearby sensors. This involves judicious sampling of the information in order to minimize the estimation error. Designing an optimal sampling strategy is a hard problem, because some easy strategies, such as continuous sampling, are infeasible due to the limited energy resources and can be far from optimality due to the transmission delay, limited channel capacity, etc.

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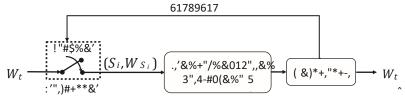


Fig. 1. System model.

To that end, in this paper, we aim to embark on a sampling problem of a wireless network, as is illustrated in Fig. 1. The sampler takes the sample of a continuous-time source process and transmits the sample to a remote estimator. The continuous-time source process is modeled as the Wiener process  $\mathbb{Z}_2$ , which helps describe the dynamics of sensors measuring quantities like movement, providing insights into how these quantities change over time. The Wiener process, also commonly referred to as Brownian motion [18], is one of the best known Lévy process, that features stationary and independent increments. It finds widespread applications in various fields such as pure and applied mathematics, economics, quantitative finance, evolutionary biology, and physics. The Wiener process  $\mathbb{Z}_2$  has the following key properties: (i)  $\mathbb{Z}_0 = 0$ ; (ii)  $\mathbb{Z}_1$  is continuous; (iii)  $\mathbb{Z}_2$  has independent increments; (iv)  $\mathbb{Z}_2 - \mathbb{Z}_1 \sim \mathbb{N}$  (0,2 - 2) for  $0 \leq \mathbb{Z} \leq \mathbb{Z}$ , where  $\mathbb{N}$  denotes the normal distribution. The remote estimator, in turn, provides a minimum mean square estimation error (MMSE) estimate  $\mathbb{Z}_2$  based on the received samples. The core objective is to control the sequence of sampling times to minimize the estimation error  $\mathbb{Z}_2 - \mathbb{Z}_2$ , specifically, aiming at optimizing the long term average of MMSE.

Organized according to the sampling strategies and optimization metrics, our review of related works encompasses three distinct perspectives.

# 1.1 Related Works

Signal-aware sampling with reliable transmissions. There have been several studies on sampling for remote estimation, e.g., in [19, 21, 30–32], where the sampling times depend on the source process (signal-aware). A nice survey paper is included in [13]. In [30], the authors consider the

Wiener process as the source process and provide an exact solution to minimize the estimation error. According to the optimal solution, the sampler should wait until the instantaneous estimation error exceeds a threshold, and the threshold is given explicitly. A similar result was developed in [21] by extending [30] from the Wiener process to the Ornstein Uhlenbeck (OU) process. The optimal threshold retains its simplicity, remaining a root of a closed-form equation.

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Further exploration, as in [32], delves into an asymmetric sensor-controller remote system. In this scenario, there are random transmission times in both directions. At the sensor, the sampling time is a stopping time based on the evolution of the Wiener process, and at the controller, the sampling time depends on the information sent from the sensor. The authors yield precise optimal solutions, noting the potential existence of multiple thresholds for the sensor's optimal stopping time. Joint optimality designs on the sampling and the estimation, concerning the Wiener process or the autoregressive process, is investigated in [6, 11].

To summarize, except [6], these previous studies on sampling assume reliable transmissions. However, in a variety of wireless systems, channel errors may occur due to fading, and the transmission times of a packet could be random. While packet drops are considered in [6], a time-slotted system is considered, which assumes that the total transmission time is the same as the transmission instance (one time slot). In contrast, in this paper, our model allows for both packet errors and random transmission times. Moreover, we enable the selection of real-valued transmission

instances.

Signal-agnostic sampling. When the sampling times are independent of the Wiener process (signal-agnostic), the MMSE is equal to the age of information [30]. More generally, the MMSE is a function of the age of information under a linear time invariant system [7, 15]. Thus, our study is closely related to numerous studies on age-based sampling, e.g., in [1, 2, 12, 17, 22–24, 28]. Age of information, or simply age, is a metric to evaluate data freshness. Age at current time  $\mathbb P$  is defined as  $\Delta(\mathbb P) = \mathbb P - \mathbb P$  ( $\mathbb P)$ , where  $\mathbb P$  ( $\mathbb P)$  is the generation time of the latest delivered sample. Age has gained much popularity in the recent decade and has contributed to various remote control systems such as sensor networks, UAV navigation, and semantic communication. A recent literature review on the age is provided in [35].

In [2], the paper studies sampling energy harvesting sources with a unit battery buffer under an erasure channel. In the case of a single source, it provides an optimal sampling policy without feedback. With perfect feedback, an optimal policy is offered among the policies that may wait only when the previous transmission is successful. In [1], the paper solves explicit optimal solutions for an energy harvesting source with finite buffer sizes, where the arrived energy can fill up the whole buffer or fill up incrementally. In [28], the authors relate autocorrelation, remote estimation, and mutual information to the nonlinear age penalty functions, and provide an optimal sampling policy under sampling rate constraint. In [24], when the source process is a multidimensional Gaussian diffusion process, and the estimator is the Kalman Filter, the expected square estimation error is an increasing function of the age. For a general non-decreasing age penalty function, the optimal sampling policy has a threshold structure under unreliable channel conditions and random transmission delay. An extended sampling scenario where the sampler can transmit the sample before receiving the feedback is studied in [17].

However, compared to the signal-aware sampling policies, signal-agnostic counterparts exhibit suboptimal performance in terms of minimizing the estimation error. Numerical results in [30] validate that the optimal signal-aware sampling policy can achieve less than half of the long term average MMSE than that of the age-optimal sampling policy. This is intuitive, due to the criticality

of the content of information within remote monitoring systems, such as the pedestrian intentions in vehicular networks or target locations in UAV navigations.

Aoll-optimal scheduling. Recently, researchers have studied signal-aware policies to optimize a new metric: the age of incorrect information (Aoll) [8, 14, 16]. Aoll incorporates both the content of information (estimation error) and the freshness of information (data freshness). Aoll was first advanced in [16], serving as a cornerstone for subsequent research. In the context of a finite symmetric Markov source, [16] provides the transmission strategy with a focus on minimizing the Aoll, displaying low computational complexity. In [14], the authors employ dynamic programming to minimize the Aoll under a binary Markovian source and exponential channel delay distribution. Meanwhile, the paper in [8] extends [14] to a general transmission time distribution, showing that it is optimal to always transmit whenever the channel is idle and the Aoll is not zero.

Although these studies focus on content-aware transmission strategies, they all focus on a finite state Markov source under a discrete-time system. These scenarios restrict transmission choices between transmit and idle at the beginning of each time slot. Instead, we consider an unbounded and continuous-time Markov process, enabling the selection of real-valued transmission instances.

### 1.2 Our Contributions

In comparison to these three prevailing perspectives, in this paper, we consider a scenario of minimizing the estimation error of the Wiener process. Specifically, we (i) embrace a signal-aware sampling policy and (ii) accommodate an unreliable channel with a random transmission time. Our contributions expand on [30] by considering an unreliable channel, and [24] by allowing sampling time dependence on the content of the Wiener process. Our problem belongs to a semi-Markov decision problem and is difficult to solve. There have been solutions for some special cases. In the first case where the channel is reliable (e.g., [21, 28, 30, 32]), the original problems are reduced to a single sample problem, which can be further solved by convex optimizations or optimal stopping rules. However, these methods do not hold in our case because our new problem is decoupled to a recursive optimal stopping problems with multiple samples . Similarly, our work is different from

[24], because this problem is decoupled to a discounted MDP, and each action of the MDP is not a stopping time. Nonetheless, we are able to circumvent these challenges and solve the optimal sampling problem. The main contributions of this paper are stated as follows:

• We provide an exact solution to our optimal sampling problem. The optimal sampling strategy has a simple structure: each sampling time is a stopping time that takes the sample when the instantaneous estimation error exceeds a threshold. The optimal threshold remains the same, independent of the Wiener process value and whether the last transmission failed or not. Moreover, the optimal threshold can be solved efficiently, e.g., by using a two layer bisection search algorithm. Our results hold for general distributions of the transmission delay and arbitrary probability of the i.i.d. transmission failure. To solve our recursive optimal stopping problems, we developed new approaches. We provide an exact value function to the value iteration problem. Specifically, we solve a sequence of optimal stopping problems, where the action value function implies taking an action at the first sample and taking the optimal stopping times at the remaining samples. The technical tools

used to establish the results include (a) the strong Markov property and Martingale properties of the Wiener process, (b) Shiryaev's free boundary method for solving optimal stopping problems.

- When the sampling time does not depend on the Wiener process, the expected square
  estimation error is equal to the age [30], and our original problem is equivalent to an age
  minimization problem. We provide the exact solution as well. The sampler takes the sample
  when the age first exceeds a threshold. This result also improves [24, Theorem 1] by
  removing the assumption of the regenerative process.
- Numerical simulations are provided to validate our results. An interesting observation is that
  when the channel is highly unreliable, our optimal policy still performs much better than the
  age-optimal and zero-wait policies.

### 2 MODEL AND FORMULATION

# 2.1 System Model and MMSE Estimator

We consider a continuous-time status update system as is depicted in Fig. 1, where a sampler takes the sample from the Wiener process  $\mathbb{Z}_{\mathbb{Z}}$  and transmits to a destination through an unreliable channel. The destination provides an estimate  $\mathbb{Z}$  based on the samples that have been successfully delivered. The extended setting from a reliable channel to an unreliable channel is one of the key features of our study.

We use  $\mathbb{Z} \in \{^1, 2, \ldots\}$  to indicate the number of samples generated by the sampler. The  $\mathbb{Z}$ th sample is generated at time  $\mathbb{Z}_2$  and is transmitted through the unreliable channel. The sample contains the sampling time  $\mathbb{Z}_2$  and the sample value  $\mathbb{Z}_2$ . The unreliable channel has an i.i.d. transmission failure, and we denote  $\mathbb{Z} \in [0, 1)$  as the probability of failure (i.e., the channel condition is *OFF*).

2

The channel also has an i.i.d. transmission time  $\mathbb{Z}_{\mathbb{Z}}$ , and we have  $\mathbb{E}[\mathbb{Z}_{\mathbb{Z}}] < \infty$ . The transmission time and the channel condition are mutually independent. In this paper, we also assume that the

<sup>&</sup>lt;sup>1</sup> Also, our problem is significantly different from that with instantaneous transmission time, e.g., [11], because even if there is no sampling rate constraint, the zero-wait sampling policy is not optimal.

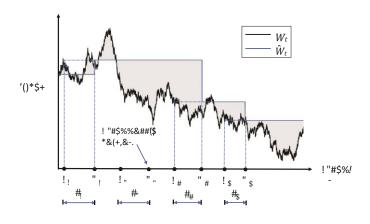


Fig. 2. A sample path of the Wiener process \$\mathbb{Z}\_2\$ and the MMSE \$\mathbb{Z}\_2\$ over time \$\mathbb{Z}\$. At \$\mathbb{Z}\_1\$, \$\mathbb{Z}\_3\$, \$\mathbb{Z}\_4\$, the sample is successfully delivered, so \$\mathbb{Z}\_2\$ is updated to be \$\mathbb{Z}\_{21}\$, \$\mathbb{Z}\_{23}\$, \$\mathbb{Z}\_{24}\$, respectively. At \$\mathbb{Z}\_2\$, the sample is not successfully \$\hat{\text{delivered}}\$, so \$\mathbb{Z}\_2\$ remains unchanged.

transmission time is lower bounded, i.e., there exists  $\mathbb{Z} > 0$  (which can be sufficiently small) such that  $\mathbb{Z}_{\mathbb{Z}} \geq \mathbb{Z}$ . The  $\mathbb{Z}$ th sample is delivered to the destination at time  $\mathbb{Z}_{\mathbb{Z}}$ , where  $\mathbb{Z}_{\mathbb{Z}} = \mathbb{Z}_{\mathbb{Z}} + \mathbb{Z}_{\mathbb{Z}}$ . At the delivery time  $\mathbb{Z}_{\mathbb{Z}}$ , the destination knows the outcome of the transmission of the  $\mathbb{Z}$ th sample. Only if the transmission was successful, the destination receives the sample message ( $\mathbb{Z}_{\mathbb{Z}}, \mathbb{Z}_{\mathbb{Z}}$ ). In addition, at  $\mathbb{Z}_{\mathbb{Z}}$ , the destination then sends an acknowledgment back to the sampler, informing whether the transmission of the  $\mathbb{Z}$ th sample was successful or not. We assume that the transmission process of the acknowledgment is instantaneous and error free. Note that the sampler always generates a sample after it receives feedback, i.e.,  $\mathbb{Z}_{\mathbb{Z}+1} \geq \mathbb{Z}_{\mathbb{Z}}$ . Otherwise, the generated sample will be queued for waiting to be transmitted, and the queued sample is staled compared to the fresh sample.

The estimator (destination) also provides a minimum mean square error (MMSE) estimator 2 abased on the successfully received samples until time 2.

We denote the random variable  $\mathbb{Z}$  as the index of the latest sample that is *successfully* delivered to the destination by the time  $\mathbb{Z}$ . In the special case of a reliable channel, each sample is successfully delivered, so we have  $\mathbb{Z} = \mathbb{Z}$ ; otherwise,  $\mathbb{Z} \leq \mathbb{Z}$ . The latest (and thus freshest) sample the destination

has received during  $\mathbb{Z} \in [\mathbb{Z}_p, \mathbb{Z}_{p+1})$  is  $(\mathbb{Z}_p, \mathbb{Z}_{p^2})$ . Using the strong Markov property of the Wiener Proc. ACM Meas. Anal. Comput. Syst., Vol. 7, No. 3, Article 60. Publication date: December 2023.

process [25, Eq. (4.3.27)], the MMSE estimator 2 ais expressed as

$$2^{2} = E[2_{2} \mid 2_{2}, 2_{2}] = E[2_{2} \mid 2_{2}] = 2_{2}, 2 \in [2_{2}, 2_{2}+1).$$
 (1)

- - -

A sample path of  $\mathbb{Z}_2$ ,  $\mathbb{Z}_2$ , and the estimation error  $\mathbb{Z}_2 - \mathbb{Z}_2$  are depicted in Fig. 2. In this figure, the 2nd sample is not successfully delivered. Thus, when  $\mathbb{Z} \in [\mathbb{Z}_2, \mathbb{Z}_3]$ , the estimator  $\mathbb{Z}_2$  is still  $\mathbb{Z}_2$ , not  $\mathbb{Z}_2$ . In other words,  $\mathbb{Z}_2 = 2$ , but  $\mathbb{Z}_2 = 1$ . This is one of the key differences from the previous studies with the reliable channel case, e.g., [21, 30–32].

# 2.2 Sampling Times and Problem Formulation

We will control the sequence of causal sampling times  $\mathbb{Z}_{\mathbb{Z}}$ 's to minimize the estimation error. We will consider two types of sampling time: (i) the sampling time depends on the Wiener process (signalaware sampling) and (ii) the sampling time is independent of the Wiener process (signalagnostic sampling).

2.2.1 Signal-aware Sampling. When the sampling time  $\mathbb{Z}_{\mathbb{Z}}$  depends on the Wiener process,  $\mathbb{Z}_{\mathbb{Z}}$  is a stopping time, i.e.,  $\mathbb{Z}_{\mathbb{Z}}$  satisfies:

$$+ + \{\mathbb{Z}_{\mathbb{Z}} < \mathbb{Z}\} \in F(\mathbb{Z}), F(\mathbb{Z}) \triangleq \cap_{\mathbb{Z} > \mathbb{Z}} \mathbb{Z}(\mathbb{Z}_{\mathbb{Z}}, \mathbb{Z} \in [0, \mathbb{Z}]). \tag{2}$$

$$( ) - F() +$$

Here,  $\mathbb{Z}[2]_1, \ldots, \mathbb{Z}_2$  is the  $\mathbb{Z}$  field generated by the random variables  $\mathbb{Z}_1, \ldots, \mathbb{Z}_2$ , and  $\mathbb{Z}$  is a filtration, i.e., a non-decreasing and right-continuous family of  $\mathbb{Z}$ —field available to the sampler at time  $\mathbb{Z}$ . Intuitively, the sampling time  $\mathbb{Z}_2$  not only depends on the history information prior to  $\mathbb{Z}_2$ —1, but also depends on the evolution of the Wiener process starting from  $\mathbb{Z}_2$ —1.

Then, we define the sampling policies. The policy space  $\Pi_{\text{signal-aware}}$  is defined as the collection of causal policies  $\mathbb{Z} = \mathbb{Z}_1, \mathbb{Z}_2, \ldots$  such that: (i)  $\mathbb{Z}_2$  satisfies the condition (2), and  $\mathbb{Z}_2 \geq \mathbb{Z}_{2^{n-1}}$ ; (ii) For each  $\mathbb{Z}_2$ , the waiting time  $\mathbb{Z}_2 - \mathbb{Z}_{2^{n-1}}$  is bounded by a stopping time that is independent of the history information before  $\mathbb{Z}_{2^{n-1}}$ . In addition, this bounded stopping time  $\mathbb{Z}^{\infty}$  satisfies  $\mathbb{E}_2 \mathbb{Z}_2 = \mathbb{Z}$ 

2.2.2Signal-agnostic Sampling. When the sampling time is independent of the Wiener process, we then define the collection of policies  $\square_1, \square_2, \ldots$  such that: (i)  $\square_2$  satisfies  $\square_2 \ge \square_{2-1}$ ; (ii) For each  $\square$ ,  $\square_2 - \square_{2-1}$  is bounded by a finite 2nd moment random variable that is independent of the history information before  $\square_{2-1}$ .

Note that for any finite 2nd moment random variable  $\mathbb{Z}$ , we have  $\mathbb{E} \mathbb{Z}^4 = 3\mathbb{E} \mathbb{Z}^2 < \infty$ .

Therefore, ∏signal-agnostic ⊂∏signal-aware

2.2.3 Problem Formulation. Our objective in this paper is to optimize the long-term average mean square estimation error (MSE) for both signal-aware and signal-agnostic cases:

1 
$$\int$$
 ②

mse<sub>opt</sub> = inf lim sup \_E (② -② -② )²②② . (3) ② ←□ ② →∞ ② 0

We aim to find a sampling policy  $\mathbb{Z}$  from the set  $\Pi$  of all causal policies, in order to minimize the MSE. The value  $mse_{opt}$  is also called the optimal objective value. Problem (3) is typically hard to solve due to the following reasons. (i) Problem (3) is an infinite horizon undiscounted semi-Markov decision problem with an uncountable state space. (ii) For the case of signal-aware sampling, each action (sampling time) is a stopping time.

## 3 MAIN RESULTS

# 3.1 Optimal Signal-aware Sampling Policy

We first break down the time-horizon problem (3) into a series of optimal sampling subproblems. Each of these subproblems determines the optimal sampling times between  $2^-$  and  $2^-$ , where  $2^-$  represents the time of the  $2^-$ th successful delivery.

Lemma 1. Solving the problem (3) is the same as solving a series of equivalent optimal sampling subproblems, where the 2th subproblem is given by

where 
$$= mse_{ont}$$
.

 $<sup>^2</sup>$ It is the upper bound stopping time that is independent of the history information before  $^2$  $_{\square}$ -1, not all of the stopping times. For example, the upper bound stopping time is to stop where the estimation error  $^2$  $_{\square}$  exceeds a sufficiently large value. Setting this stopping time as an upper bound is reasonable, because we want to minimize the estimation error. Then, this stopping time is independent of the history information before  $^2$  $_{\square}$ -1.

<sup>&</sup>lt;sup>3</sup>If the condition (ii) does not hold, then the term  $\lim_{\mathbb{Z}\to\infty} \mathbb{E}[\mathbb{Z}_{\mathbb{Z}(\mathbb{Z})}]/\mathbb{Z}$  may not be 0, where  $\mathbb{Z}(\mathbb{Z})$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}<\mathbb{Z}$ , and  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}<\mathbb{Z}$ , and  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}<\mathbb{Z}$ , and  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}<\mathbb{Z}$ , and  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}<\mathbb{Z}$ , and  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}<\mathbb{Z}$ , and  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}<\mathbb{Z}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  is the largest number  $\mathbb{Z}$  such that  $\mathbb{Z}_{\mathbb{Z}}$  is the largest number  $\mathbb{Z}$  is the largest number  $\mathbb$ 

Lemma 1 is a restatement of Lemma 4 in Section 5. We note that, by choosing  $\mathbb{Z} = \mathsf{mse}_{\mathsf{opt}}$ , the sequence of linearized optimal stopping subproblems (4) have the same solution as the original problem (3). Note that these subproblems are independent and thus *equivalent*. In other words, we only need to solve one subproblem (4) regardless of  $\mathbb{Z}$ , and  $\mathbb{Z}$  ( $\mathbb{Z}$ ,  $\mathbb{Z}$ ) remains the same for any given

However, problem (4) is still hard to solve. Let □ be the total number of transmission attempts

is a repeated optimal stopping problem with continuous-time control and a continuous state space. This

is the key technical challenge of our study. To the extend of our knowledge, this type of problems has not been addressed before. One limiting case of problem (4) was studied in [30, Eq. 47], where there exists no transmission errors and hence  $\mathbb{Z}_{\mathbb{Z}} = 1$ .

We develop a value iteration algorithm that can find the optimal stopping times for solving problem (4). To that end, we define a sequence of optimal stopping problems:

$$\mathbb{Z}_{\mathbb{Z}}\left(\mathbb{Z},\mathbb{Z}\right) \triangleq \inf \mathbb{E}\left(\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z}^{2}\right)^{2}\mathbb{Z} - \mathbb{Z}\left(\mathbb{Z}_{\mathbb{Z}} + 1 - \mathbb{Z}_{\mathbb{Z}} + 1 - \mathbb{Z}_{\mathbb{Z}} + 1 - \mathbb{Z}_{\mathbb{Z}}\right)$$

$$\mathbb{Z} \in \Pi \qquad \mathbb{Z}^{1 - \min(\mathbb{Z}_{\mathbb{Z}}\mathbb{Z})}$$

$$\mathbb{Z}_{\mathbb{Z}} \qquad -\mathbb{Z}^{2} \mathbb{Z} \qquad = \mathbb{Z},\mathbb{Z} = 1,2,\ldots.$$

$$\mathbb{Z}_{1 - \min(\mathbb{Z}_{\mathbb{Z}}\mathbb{Z})} \qquad \mathbb{Z}_{1 - \min(\mathbb{Z}_{\mathbb{Z}}\mathbb{Z})}$$

$$\mathbb{Z}_{1 - \min(\mathbb{Z}_{\mathbb{Z}}\mathbb{Z})} \qquad \mathbb{Z}_{1 - \min(\mathbb{Z}_{\mathbb{Z}}\mathbb{Z})}$$

Hence,  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ,  $\mathbb{Z}$ ) determines the optimal solution for at most the last  $\mathbb{Z}$  transmission attempts in problem 4. The principle of backward induction implies that  $\mathbb{Z}_{\mathbb{Z}}$  satisfies the *value iteration algorithm*:

$$\mathbb{P}_0(\mathbb{P}, \mathbb{P}) \triangleq 0,$$

$$\mathbb{P}_{\mathbb{P}^+1}(\mathbb{P}, \mathbb{P}) \triangleq \inf \mathbb{P}(\mathbb{P}; \mathbb{P}) + \mathbb{P} \mathbb{E} \left[ \mathbb{P}_{\mathbb{P}} \left( \mathbb{P} + \mathbb{P}_{\mathbb{P}^+1}, \mathbb{P} \right) \right], \mathbb{P} = 0, 1, 2, \dots,$$
(6)

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where  $\mathbb{Z} + \mathbb{Z}_{\mathbb{Z} + \mathbb{Z}}$  is the estimation error after a stopping time  $\mathbb{Z}$  and the transmission time  $\mathbb{Z}$ . And the per-stage cost function  $\mathbb{Z}(\mathbb{Z};\mathbb{Z})$  is defined as the square estimation error minus  $\mathbb{Z}$  from the last delivery time to the next delivery time with a stopping time  $\mathbb{Z}$ :

The following theorem provides an exact solution to (6), which is the key contribution in this paper:

Theorem 1. The sequence of optimal stopping times  $\mathbb{Z}_{\delta}$ 's to problem (6) is given as follows:

 $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) is the unique positive root of the free boundary differential equation:

2<sub>2</sub> (2, 2) is updated as:

$$2_0(2, 2) = 0,$$

1

$$\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}, \mathbb{Z}) = \mathbb{Z}(\mathbb{Z}, \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}), \mathbb{Z}) + \mathbb{Z}\mathbb{E}_{\mathbb{Z}}[\mathbb{Z}_{\mathbb{Z}}-1(\max\{|\mathbb{Z}|, \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})\} + \mathbb{Z}_{\mathbb{Z}}, \mathbb{Z})], \mathbb{Z} = 1, 2, \dots, (10)$$

Fig. 3. The evolution of the root function:  $\mathbb{Z}^{\mathbb{Z}}(\mathbb{Z},\mathbb{Z}) + \mathbb{I}_{3\mathbb{Z}^3} - \mathbb{Z}$  over  $\mathbb{Z}$ , with  $\mathbb{Z} = \mathbb{I}_{1,2,4}^{2}$ . In this example, we set  $\mathbb{Z} = 11.0, \mathbb{Z} = 0.3$ , and a constant transmission delay  $\mathbb{Z} = 6$ . It is easy to see that  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})$ , which is the positive root of  $\mathbb{Z}^{\mathbb{Z}}(\mathbb{Z},\mathbb{Z}) + \mathbb{I}_{3,3}^{2} - \mathbb{Z}^{2}$ , is decreasing in  $\mathbb{Z}$ . the function  $\mathbb{Z}(\mathbb{Z},\mathbb{Z},\mathbb{Z})$  is equal to

1

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$$\mathbb{Z}(\mathbb{Z}, \mathbb{Z}, \mathbb{Z}) = \mathbb{Z} \mathbb{Z}^2 + \mathbb{E} \mathbb{Z} \mathbb{Z}^2 - \mathbb{E} \mathbb{Z} \mathbb{Z}^2 - \mathbb{E} \mathbb{Z} \mathbb{Z}^2 + \mathbb{E} \mathbb{Z}^2 \mathbb{Z}^2 = \mathbb{Z}^2 \mathbb{Z}^2 = \mathbb{Z}^2 \mathbb{Z}^2 = \mathbb{Z}^$$

(11) Moreover, the sequence  $\{\mathbb{P}_{\mathbb{P}}(\mathbb{P})\}_{\mathbb{P}}$  is decreasing and thus convergent.

The proof of Theorem 1 is provided in Section 5.3. Theorem 1 implies that each optimal stopping time  $\mathbb{Z}_2$  is a hitting time that will stop when the estimation error exceeds a threshold  $\mathbb{Z}_2$  ( $\mathbb{Z}$ ). The threshold  $\mathbb{Z}_2$  ( $\mathbb{Z}$ ) is chosen by the free boundary method [25], where the optimal value function  $\mathbb{Z}_2$  ( $\mathbb{Z}$ ) should be continuously differentiable on  $\mathbb{Z} \in \mathbb{R}$ . Since  $\mathbb{Z}_2$  ( $\mathbb{Z}$ ) is decreasing and convergent,  $\mathbb{Z}_2$  is also convergent.

In addition, the optimal threshold  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) can be solved efficiently. In Theorem 5 of Section 5.3, we showed that the root of the free boundary method in (9) is equivalent to:

2 1 2 13

Interestingly,  $\mathbb{Z}_{2}$  ( $\mathbb{Z}_{2}$ ,  $\mathbb{Z}_{2}$ ) =  $\mathbb{Z}_{2}$  ( $\mathbb{Z}_{2}$ ,  $\mathbb{Z}_{2}$ ) |  $\mathbb{Z}_{2}$  =  $\mathbb{Z}_{2}$  ( $\mathbb{Z}_{2}$ ) |  $\mathbb{Z}_{2}$  =  $\mathbb{Z}_{2}$  =  $\mathbb{Z}_{2}$  ( $\mathbb{Z}_{2}$ ) |  $\mathbb{Z}_{2}$  =  $\mathbb{Z}_{2}$  =  $\mathbb{Z}_{2}$  ( $\mathbb{Z}_{2}$ ) |  $\mathbb{Z}_{2}$  =  $\mathbb{Z$ 

 $\mathbb{Z}_0$  ( $\mathbb{Z}$ ,  $\mathbb{Z}$ ) = 0, and the function  $\mathbb{Z}_2^{\mathbb{Z}}$  ( $\mathbb{Z}$ ,  $\mathbb{Z}$ ) is updated as

?

1

$$+2\mathbb{E}_{\mathbb{Z}} \qquad \mathbb{Z}^{\mathbb{Z}}_{-1}(\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}},\mathbb{Z})^{1}|\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}}|\geq \mathbb{Z}_{-1}(\mathbb{Z})+ \qquad \mathbb{Z}(\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}})-\mathbb{Z}(\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}}) = \mathbb{Z}(\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}}) = \mathbb{Z}(\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}})$$

Because (13) contains only an expectation over  $\mathbb{Z}_{\mathbb{Z}}$  without derivatives, computing  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ,  $\mathbb{Z}$ ) is easy.  $\mathbb{Z}$ 1 3

We also showed that  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) + 3.2 -22 is strongly convex for  $\mathbb{Z}$  > 0. Thus, we only need logarithm time complexity to solve  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) for each  $\mathbb{Z}$  in (12), such as bisection search or Newton's method.

Fig. 3 illustrates some intuitive properties of  $\mathbb{Z}_2$  ( $\mathbb{Z}$ ) and its root function,  $\mathbb{Z}_2$  ( $\mathbb{Z}$ ,  $\mathbb{Z}$ ) + 3 $\mathbb{Z}$  - 2 $\mathbb{Z}$ 

Further, 🛮 converges linearly to 🗈. To illustrate, we first define a norm. Let us pick any value 🗈

( $\mathbb{P}$ ) = max( $\mathbb{P}$ ,  $\mathbb{P}^{-2}$ ), where  $\mathbb{P}^{-1}$  can take any positive

with  $\mathbb{Z} < \mathbb{Z} < 1$ , and denote a weight function  $\mathbb{Z}$ 

$$+$$
  $^{2}\sqrt{2}$   $+$   $_{-}$   $+$   $_$ 

P (2)

 $\|\cdot\|$  of a function  $\mathbb{Z}$  ( $\mathbb{Z}$ ) is defined as  $\|\mathbb{Z}\| = \sup_{\mathbb{Z} \to \mathbb{Z}} (\mathbb{Z})$ . We have the following result: Lemma 2.  $\|\mathbb{Z}_{\mathbb{Z}}(\cdot, \mathbb{Z}) - \mathbb{Z}(\cdot, \mathbb{Z})\| \le \mathbb{Z}\|\mathbb{Z}_{\mathbb{Z}^{-1}}(\cdot, \mathbb{Z}) - \mathbb{Z}(\cdot, \mathbb{Z})\|$ .

Lemma 2 is restated in Lemma 9 at Section 5.4. Since  $\mathbb{Z}_{\mathbb{Z}}$  is also convergent, each of the optimal stopping (waiting) times in (4) should also be a hitting time with the threshold  $\mathbb{Z}(\mathbb{Z}) = \lim_{\mathbb{Z} \to \infty} \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})$ . We finally conclude the following result:

Theorem 2. An optimal sampling solution 2 's to the series of problem (4) is:

$$\mathbb{Z}_{\mathbb{Z}+1} = \inf \mathbb{Z} \ge \mathbb{Z}_{\mathbb{Z}} : |\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z} \cap \mathbb{Z}| \ge \mathbb{Z}(\mathbb{Z}) , \mathbb{Z} = 0, 1, 2, \dots,$$

$$(14)$$

where  $\mathbb{Z}(\mathbb{Z})$  is the limit of the sequence  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})$ 's, and  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})$  can be computed by solving (9), or more efficiently, by solving (12) and (13).

The proof of Theorem 2 is provided in Section 5.4.

Theorem 2 illustrates an important property of an optimal sampling policy for a given parameter  $\mathbb{Z}$ . Note that  $\|\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z}\|_{\mathbb{Z}}$  is the estimation error at the current time  $\mathbb{Z}$ . Theorem 2 implies that the optimal sampling policy given in (14) has a simple structure. The optimal policy is a *threshold type*: the sampler may wait until the instantaneous estimation error  $\|\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z}\|_{\mathbb{Z}}$  exceeds the threshold  $\mathbb{Z}(\mathbb{Z})$ . Specifically, if the estimation error at the initial time  $\mathbb{Z}_{\mathbb{Z}}$  exceeds  $\mathbb{Z}(\mathbb{Z})$ , then it is optimal to immediately transmit the sample. The optimal threshold  $\mathbb{Z}(\mathbb{Z})$  is independent of the evolution of the Wiener process.

After solving (4) with a given  $\mathbb{Z}$ , we will finally determine the optimal objective value  $\mathbb{Z} = \mathsf{mse}_{\mathsf{opt}}$ .

Note that in (4),  $\mathbb{Z}_{\mathbb{Z}} = -\mathbb{Z}_{\mathbb{Z}} = -\mathbb{Z}_{\mathbb{Z}}$  has the same distribution as  $\mathbb{Z}_{\mathbb{Z}}$ , where  $\mathbb{Z}$  has the same distribution as

the i.i.d. transmission delay 2 's. Then, we have the following result:

Theorem 3.  $\square$  = mseopt is the root of

$$\mathsf{E}\left[\mathbb{P}\left(\mathbb{P}_{\mathbb{P}},\,\mathbb{P}\right)\right]=0,\tag{15}$$

where mseopt is the optimal objective value of (3).

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Theorem 3 is shown in Lemma 3 at Section 5. Combining Theorem 2 and Theorem 3, we finally provide the optimal solution to (3).

Moreover, we showed that we can also use a low complexity algorithm, such as bisection search, to compute the root of  $\mathbb{Z}$ . So in conclusion, we can efficiently solve  $\mathbb{Z}(\mathsf{mse}_{\mathsf{opt}})$  and  $\mathsf{mse}_{\mathsf{opt}}$  with low complexity, which is provided in Algorithm 1:

- Line 4—10 in Algorithm 1 is an *inner layer* update to efficiently compute the optimal threshold  $\mathbb{Z}(\mathbb{Z})$  and the function  $\mathbb{Z}(\mathbb{Z},\mathbb{Z})$  for a given  $\mathbb{Z}$  (corresponding to Theorem 1 and Theorem 2). In Line 5, due to Lemma 2, we only need a logarithm number of iterations. In Line 8, since the root function in (12) is strongly convex, we only need a simple Newton's method to obtain  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})$ .
- Line 2,3,11 serves as an *outer* layer that uses a simple bisection method to determine the root of ② (corresponding to Theorem 3).

In the special case where  $\mathbb{Z} = 0$ , it is easy to observe that  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}) = \mathbb{Z}_1(\mathbb{Z})$ , and  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}, \mathbb{Z}) = \mathbb{Z}_1(\mathbb{Z}, \mathbb{Z}) = \mathbb{Z}_1(\mathbb{Z}, \mathbb{Z})$  for all  $\mathbb{Z} = 1, 2, \ldots$  As a result, the optimal threshold  $\mathbb{Z}(\mathbb{Z}) = \mathbb{Z}_1(\mathbb{Z})$ , and the optimal value function  $\mathbb{Z}(\mathbb{Z}, \mathbb{Z}) = \mathbb{Z}(\mathbb{Z}, \mathbb{Z}_1, \mathbb{Z})$ . By (12),  $\mathbb{Z}_1(\mathbb{Z}) = p3(\mathbb{Z} - \mathbb{E}[\mathbb{Z}])$ . Therefore, Theorem 2 and 3 reduces to the following corollary:

Corollary 1. Suppose that  $\mathbb{Z} = 0$ , then an optimal solution  $\mathbb{Z}_3$ 's to problem (3) satisfies:

Algorithm 1: Bisection method for solving the optimal threshold 2(mse<sub>opt</sub>) and mse<sub>opt</sub>

```
Given \mathbb{Z}_1 small, \mathbb{Z}_2 large, \mathbb{Z}_1 < \mathbb{Z}_2, and tolerance \mathbb{Z}_1, \mathbb{Z}_2 small.
                     repeat
2
3
                     Set \mathbb{Z}_0(\mathbb{Z},\mathbb{Z}) (\mathbb{Z},\mathbb{Z}) = 0
                     Set iteration number 2 = [-\log_2]
5
                     for 2 = 1 : 2
6
                     Update 2 (2, 2) in (13)
                     Solve 2₂ (2) in (12)
8
                     Update 2 (2, 2) in (10)
9
                     end for
10
                     if E^{\mathbb{P}}[\mathbb{P}_{\mathbb{P}}(\mathbb{P}_{\mathbb{P}},\mathbb{P})] < 0: \mathbb{P}_2 = \mathbb{P}. else \mathbb{P}_1 = \mathbb{P}
11
```

until 
$$22 - 21 < 22$$

return  $2 (2)$ 

where 2 is the root of

E h p E i 
$$2(2, 3(2 - [2]), 2) = 0.$$
 (17)

Moreover,  $\mathbb{Z} = mse_{opt}$  is the optimal objective value of (3).

The optimal policy provided in Corollary 1 is the same as that of [30, Theorem 1]. In addition, we have also improved [30, Theorem 1] by removing the assumption of the regenerative process. The optimal sampling policy provided in Corollary 1 is a threshold type on the instantaneous estimation error, and the optimal threshold is given in closed-form.

There are several variations of Corollary 1 with a reliable channel case  $\mathbb{Z}=0$ . In [21], the paper changes the source process to be the Ornstein-Uhlenbeck process and shows that the optimal threshold is a root of the closed-form equation. The model where the source can reset the Wiener process is described in [32]. Theorem 2 and Theorem 3 are different from these studies by generalizing to an i.i.d. unreliable channel scenario ( $\mathbb{Z} \geq 0$ ). Note that the last transmission may be successful or failed for each sample. However, in Theorem 2 and Theorem 3, each sampling time follows the same threshold type with the same threshold  $\mathbb{Z}(\mathbb{Z})$ , regardless of whether the last transmission failed or not.

The expression (14) in Theorem 2 implies that our optimal policy relies on the value of the Wiener process at the sampling time of the successfully delivered sample,  $\mathbb{Z}_{\mathbb{Z}}$ , but may not on  $\mathbb{Z}_{\mathbb{Z}}$ . This is

also a key difference from the case of a reliable channel (2 = 0), e.g., [21, 30, 32] and Corollary 1.

# 3.2 Optimal Signal-agnostic Sampling Policy with Sampling Rate Constraint

Finally, we turn to the signal-agnostic case and provide the exact solution to Problem (3). Using [30], for any signal-agnostic policy, we have

$$\mathbb{P}\left(\mathbb{P}_{\mathbb{P}} - \mathbb{P}_{\mathbb{P}}\right) = \mathbb{P}\left(\mathbb{P}_{\mathbb{P}}, \mathbb{P}_{\mathbb{P}} + 1\right). \tag{18}$$

In other words, when the sampling time does not depend on the Wiener process, the expected square estimation error MMSE is equal to the *age of information*. So our MSE-optimal sampling problem (Problem (3)) is equivalent to the age-optimal sampling problem. Problem (3) is equivalent

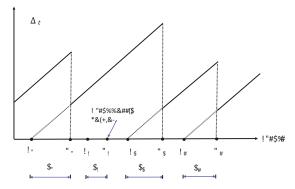


Fig. 4. Evolution of the age  $\Delta \mathbb{Z}$  over time  $\mathbb{Z}$ .

to

Age of information  $\Delta_{\mathbb{Z}}$ , or simply the age, is a metric for evaluating the data freshness. As is mentioned in (18), the age  $\Delta_{\mathbb{Z}}$  is defined as the time elapsed since the freshest delivered sample is generated [29]. If a fresh sample is successfully delivered to the estimator, the age decreases to the system time of the sample. Otherwise, the age increases linearly in time. A sample path of the age  $\Delta_{\mathbb{Z}}$  is depicted in Fig 4.

We then have the following result:

Theorem 4. An optimal solution  $\mathbb{Z}_{\mathbb{Z}}$ 's to the problem (19) is provided as:

☑ is the root of

E 
$$222 - 2E \max(2 - 2, \underline{\hspace{1cm}}) = 0,$$
 (21)

1 - 2

"  $E[2]$  #  $E[2]$ 

where  $2^{'}$   $\int_{\mathbb{R}^2=1}^{6} \mathbb{R}^2=1$   $\mathbb{R}^2=1$   $\mathbb{R}^2=1$ 

delay 2g's.

Theorem 4 provides the same sampling policy as that of [24, Theorem 1]. But we slightly improve [24, Theorem 1] by removing its assumption of the regenerative process. The proof of this improvement is provided in Appendix H.

Different from Theorem 2, the optimal sampling policy is a threshold policy on the age, or equivalently, the MMSE, instead of the instantaneous estimation error. Note that the age keeps increasing over time if there is no successful delivery. As a result, if the previous transmission

failed, the age is always larger than the optimal threshold  $\ensuremath{\mathbb{Z}}$  – \_\_\_\_- $\ensuremath{\mathbb{Z}}$ . Therefore, Theorem 4 tells that

if the previous transmission is successful, the sampler may wait for some time until the current age

exceeds the threshold  $\ensuremath{\mathbb{Z}}$  –  $\ensuremath{-}$  –  $\ensuremath{\mathbb{Z}}$  . If the previous transmission failed, the sampler chooses zerowait.

This is another key difference from the optimal signal-aware sampling policy in Theorem 2. In Theorem 2, due to the randomness of the Wiener process, each sampler may need to wait, regardless of the outcome of the previous transmission. In addition, since there is only one waiting time

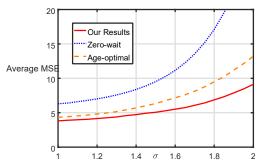


Fig. 5. Average MSE versus  $\mathbb{Z}$ , where the channel delay is lognormal distributed with the parameter  $\mathbb{Z}$ . As  $\mathbb{Z}$  increases, the channel delay distribution is more heavy-tailed. The probability of i.i.d. transmission failure  $\mathbb{Z}$  = 0.65.

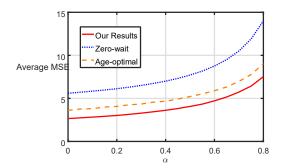


Fig. 6. Average MSE versus the probability of i.i.d. transmission failure  $\mathbb{Z}$ , where the channel delay is lognormal distributed with the parameter  $\mathbb{Z} = 1.5$ .

between two successful deliveries, the optimal objective value  $\mathbb{Z}$  is the root of the closed form expression (21). But the root function of  $\mathbb{Z}$  for the signal-aware case in (15) is not closed-form. Instead, as is illustrated in Theorem 2 and Algorithm 1, we need to construct a sequence of functions  $\mathbb{Z}_{\mathbb{Z}}$ 's to approach the root function.

#### 4 SIMULATION

In this section, we will compute the long term average MMSE (average MSE) of the following three sampling policies:

- 1. Our Results: our optimal sampling policy, which is the solution to problem (3), provided in Theorem 1—3. The average MSE is then computed in Algorithm 1. It waits until the estimation error exceeds a threshold.
- 2. Zero-wait: The source transmits a sample once it receives the feedback, i.e.,  $\mathbb{Z}_{\mathbb{R}^+1} = \mathbb{Z}_{\mathbb{R}}$ . This simple policy can achieve the maximum throughput and the minimum delay. However, even in the case of a reliable channel, it may not optimize the age of information [34] or optimize the estimation error [30]. In our study with an unreliable channel, Theorem 4 implies that the zero-wait policy does not optimize the age. Moreover, Theorem 1—3 imply that the zero-wait policy does not optimize the estimation error.

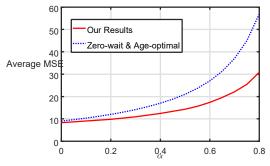


Fig. 7. Average MSE versus  $\mathbb{Z}$ , where the channel delay is a constant with the delay  $\mathbb{Z} = 6$ .

3. Age-optimal: This policy is provided in Theorem 4, restated in [24, Theorem 1], and the averageMSE is computed by [24, Algorithm 1]. Age-optimal policy achieves the optimal average age. It waits until the age (i.e., MMSE  $E[(\mathbb{Z}_2 - \mathbb{Z}_2)^2]$ ) exceeds a threshold.

We will follow the same network system as is illustrated in Section 2 and Fig. 1. We consider two scenarios about the delay distribution of the unreliable channel: heavy-tailed distribution (e.g., lognormal distribution) and short-tailed distribution (e.g., constant).

In the first scenario, we assume that the channel delay follows a lognormal distribution. The lognormal random variable with scale parameter 2 is expressed as 2 /E[2 ], where 2 is the

standard normal random variable. Fig. 5 illustrates the relationship between the average MSE of the four sampling policies with parameter ② of lognormal channel delay, given a discount factor ② (probability of failure of the channel). The numerical results validate that our proposed policy always achieves the lowest average MSE. Note that as ② increases, the lognormal distribution of the channel becomes more heavy-tailed. We observe that the zero-wait policy is far from optimality, and the age-optimal policy also grows much quicker than our optimal policy. Therefore, our optimal policy substantially outperforms the age-optimal and zero-wait policies when the channel delay becomes heavy-tailed. Fig. 6 plots the evolutions of the average MSE with the change of ② given that the parameter ② = 1.5. From our observation, the zero-wait policy is always far from our optimal policy.

In the second scenario, we assume that the channel delay is a constant. Fig. 7 depicts the evolution of the average MSE of different policies with the change of  $\mathbb{Z}$ . Note that the age-optimal policy is equivalent to the zero-wait policy when the delay is a constant, as is shown in [24, Corollary 3]. We observe that when the channel connectivity is more reliable ( $\mathbb{Z}$  very small), then the zero-wait policy is only slightly inferior to the optimal policy. However, as  $\mathbb{Z}$  increases, the zero-wait policy becomes far from optimality. The intuitive reason is as follows: since the Wiener process oscillates, with a nontrivial probability, our optimal policy waits at each sample, no matter whether the last transmission failed or not. Compared to the zero-wait policy, such a quite different sampling strategy leads to much improvement for the average MSE. This is the newly observed phenomenon that has not been found in the previous studies, e.g., [21, 24, 28, 30].

In summary, our optimal policy can perform much better than the zero-wait and the age-optimal policy when either (i) the transmission time is heavy-tailed, or (ii) the transmission time is lighttailed, and the channel is highly unreliable.

### 5 PROOF OF MAIN RESULTS

In this section, we provide the proof for efficiently solving the optimal signal-aware policy for (3). In Section 5.1, we first show that there exists an optimal policy such that the inter-sampling time of the successfully delivered packet is i.i.d. Thus, the long term average MMSE in (3) is equal to the average MMSE only between the two successful delivery times. In Section 5.2, after linearizing, the reduced problem is equivalent to optimizing a discrete time discounted problem with multiple stopping times (27). This new problem a strict generalization to a discrete time discounted MDP, where each action is extended to be a stopping time. To solve (27), in Section 5.3, we first speculate that the optimal policy and its optimal value function satisfy the Bellman equation. Then, we use a value iteration algorithm to approach the optimal value function, where each iteration is an optimal stopping problem. Interestingly, we analytically solve the optimal stopping time for each iteration, which is a key technical contribution in this paper. Finally, in Section 5.4, we use the contraction mapping property to show that the optimal value function of the value iteration

П

algorithm convergences linearly to that of the Bellman equation. Thus, we exactly solve (27). This ends the proof.

- 5.1 Reducing to a Single-epoch Problem
- 5.1.1 Replacing the subscript  $\mathbb{Z}$  by  $(\mathbb{Z},\mathbb{Z})$ . The proof relies on the number of successfully delivered samples and the number of samples attempted for a successful delivery. These messages cannot be easily described in  $\{\mathbb{Z}_2, \mathbb{Z}_2, \mathbb{Z}_2\}$ 's by using only one subscript  $\mathbb{Z}$ . Therefore, for notational simplicity, throughout Section 5, we will replace  $\mathbb{Z}_2, \mathbb{Z}_2, \mathbb{Z}_2$  by  $\mathbb{Z}_2, \mathbb{Z}_2, \mathbb{Z}_2,$

By (1), the MMSE estimator 2 sis expressed as

$$2^{2} = 2_{22-1,22-1}, 2 \in [2_{2}-1,$$

8.1.2 Reducing to a Single-epoch Problem. We aim to show that solving the original problem (3) can be reduced to solving the optimal sampling times  $\mathbb{Z}_{\mathbb{Z},1},\mathbb{Z}_{\mathbb{Z},2},\ldots$  within an epoch  $\mathbb{Z}$  over a subset of the policy space  $\Pi_{\text{Signal-aware}}$ . We denote such the subset  $\Pi_{\mathbb{Z}}$  as a collection of sampling times  $\mathbb{Z}_{\mathbb{Z},1},\mathbb{Z}_{\mathbb{Z},2},\ldots$  within epoch  $\mathbb{Z}$  such that each inter-sampling time  $\{\mathbb{Z}_{\mathbb{Z},2},\mathbb{Z}_{\mathbb{Z},2},\ldots\}$  is independent of the history information before  $\mathbb{Z}_{\mathbb{Z}-1},\mathbb{Z}_{\mathbb{Z}-1}$ . The following result shows that our average cost problem (3) reduces to a single epoch problem (with arbitrary index  $\mathbb{Z}$ ) that contains possibly multiple samples from one successful delivery time until the next successful delivery time.

Proposition 1. There exists an optimal policy for the problem (3) such that  $\{20,20-20-1,20-1\}$  are i.i.d. Moreover, problem (3) is equivalent to

$$E^{h} \int_{\mathbb{R},\mathbb{R}} \left( - \mathbb{R}_{1,\mathbb{R}} \right)^{2} i$$

$$\mathbb{R} \quad \mathbb{R} \quad \mathbb{R} \quad \mathbb{R}$$

$$\mathbb{R} \quad \mathbb{R} \quad \mathbb{R} \quad \mathbb{R}$$

$$\mathbb{R} \quad \mathbb{R} \quad$$

Proof. See Appendix B.

Proposition 1 implies that to solve the long term average MMSE problem (3), we can solve a problem with only a single epoch. Each sampling decision in this epoch is independent of the history information prior to the final sampling time of the previous epoch. Proposition 1 is motivated by

[28, 30] under a reliable channel. In these studies, the original problem is reduced to an average MMSE problem between two delivery times (a single sample problem). One of the key reasons is that at each delivery time, the estimation error is updated and is independent of the history information before the last sampling time. But in our unreliable case, at a failed delivery time, the estimation error is not updated and is still correlated to that history information. Thus, our single epoch problem cannot be further reduced to a single sample problem. In addition, we also improve [28, 30] by removing the assumption of the regenerative process. A similar result to Proposition 1 is presented in [2] with an unreliable channel and signal-agnostic sampling, without the assumption of the regenerative process. We also generalize [2] since our sampling time depends on the Wiener process.

Although we have reformulated the long term average MMSE problem (3) into an average MMSE problem within a single epoch (23), problem (23) is still hard to solve. This is because it contains a fraction and thus is a repeated semi-MDP.

5.2 Reformulating as a Multiple Stopping Times Problem: an Extension to a Discounted MDP In this section, we will linearize problem (23) and reformulate it as a discounted cost and repeated Markov decision process (MDP), where each action is a stopping time.

Let us define a minimization problem with a parameter  $\mathbb{Z} \in \mathbb{R}$ :

Here,  $\mathbb{Z} = (\mathbb{Z}_{2,1}, \mathbb{Z}_{2,2}, \ldots)$ . By Dinkelbach's method [9], we have

Lemma 3. (i)  $h(2) \leq 0$  if and only if  $mse_{opt} \leq 2$ . (ii) When  $2 = mse_{opt}$ , the solution to (23) and (24) are equivalent.

Therefore, to solve (23), we will solve  $h(mse_{opt}) = 0$ .

We denote 22,2 as the waiting time for the 2th sample in epoch 2. Then,

Then, combined with (25) and the strong Markov property of the Wiener process, given that

As a result, (25) and (26) give:

Lemma 4. An optimal solution to (23) given that  $\mathbb{Z}_{\mathbb{Z}^2-1,\mathbb{Z}^2-1}-\mathbb{Z}_{\mathbb{Z}^2-1,\mathbb{Z}^2-1}=\mathbb{Z}$ ,  $\mathbb{Z}\in\mathbb{R}$  satisfies

Here,  $\mathbb{Z}$  ( $\mathbb{Z}$ ) is the total cost of the optimal policy, which is also called the *optimal value function*. And  $\mathbb{Z}_2$  ( $\mathbb{Z}$ ) is the total cost of a policy, which is also called the *action value function* with a policy  $\mathbb{Z}$ .

For any policy  $\mathbb{Z}$ , the action value function  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) in (28) is further written as

where the state values  $2 = 1, 2 \dots$  satisfy

$$\mathbb{C}_{\mathbb{C}^{+1}} = \mathbb{C}_{\mathbb{C}} + \mathbb{C}_{\mathbb{C}^{n}} + \mathbb{C}_{\mathbb{C}^{n}}, \mathbb{C} = 1, 2, \dots,$$
 (30)

②(②:②), also called a per stage cost function, is the expected integration of square estimation error 4 minus mseopt from the last

delivery time to the next delivery time, where the initial estimation error is 2, and the sampler's waiting time is  $\mathbb{Z}$ .  $\mathbb{Z}(\mathbb{Z};\mathbb{Z})$  is defined as:

$$\mathbb{P}(\mathbb{P};\mathbb{P}) = \mathbb{E} \qquad (\mathbb{P} + \mathbb{P}_{\mathbb{P}}) \,\mathbb{P}\mathbb{P} - \mathsf{mse}_{\mathsf{opt}}(\mathbb{P} + \mathbb{P}) \,, \tag{31}$$

Note that  $\mathbb{Z}_2(\mathbb{Z})$  represents the expected cost of square estimation error minus a constant mse<sub>opt</sub> within an epoch. In an epoch, if the transmission is successful with probability  $1 - \mathbb{Z}$ , then the system will stop. Thus, the system state will enter a "stopping" set with 0 cost; If the transmission fails with probability  $\mathbb{Z}$ , the system state will enter the next transmission with a per-stage cost  $\mathbb{Z}$ . Therefore,

$$\mathbb{P}^{2}(\mathbb{P}) = \tilde{\mathbb{Q}}\mathbb{P}^{-1}\mathbb{E}[\mathbb{P}(\mathbb{P}^{2},\mathbb{P},\mathbb{P})|\mathbb{P}^{-1} = \mathbb{P},$$

$$\mathbb{P}^{2} = 1$$
(32)

which is proven in [24, Appendix F]. The 2th stage state 2 2 implies that all the previous 2 – 1

transmissions failed, and the coefficient 2 is the probability of 2 1 consecutive failures.

Equations (27)—(32) imply that problem (27) belongs to a discounted cost problem with multiple stopping times, or in other words, a *repeated MDP*, because there are multiple waiting times

 $\mathbb{Z}_{\mathbb{Z},1},\mathbb{Z}_{\mathbb{Z},2},\ldots$ , and each waiting time is a stopping time. Suppose that each waiting time is not a stopping time, i.e., the waiting time policy chooses a real value that is independent of the Wiener process. Then, problem (27) is reduced to a discrete time discounted cost MDP [3]. This is because:

(i) the state at each stage  $\mathbb Z$  is the estimation error at the  $\mathbb Z-1$ th delivery time,  $\mathbb Z_{\mathbb Z}$  (when  $\mathbb Z=1$ ,  $^{\sim}=(32)$ ). (ii) The action at each stage  $\mathbb Z$  is the waiting time for the  $\mathbb Z$ th sample,  $\mathbb Z_{\mathbb Z}$ . (iii) The  $\mathbb Z_1$   $\mathbb Z_2$  state transition is provided in (30). (iv) The cost function is defined in (31).

Note that the waiting times  $\mathbb{Z}_{\mathbb{Z},1},\mathbb{Z}_{\mathbb{Z},2},\ldots$  are correlated. Thus, despite that we have linearized the problem (23) into a multiple stopping time problem (27), problem (27) still faces the curse of dimensionality.

5.3 Analytical Solution to the Value Iteration (35) for the Multiple Stopping Times Problem (27) In the special case where each waiting time  $\mathbb{Z}_{2,1}, \mathbb{Z}_{2,2}, \ldots$  is not a stopping time, the optimal policy and the optimal value function to the discounted MDP satisfies the Bellman equation [5, Chapter

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9]. The advantage of the Bellman equation is that it turns the MDP with correlated waiting times into an optimization problem over a single waiting time and thus helps reduce the complexity of

<sup>4</sup>For comparison,  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) is the expected integration of square estimation error minus mse<sub>opt</sub> from the last delivery time to the next *succsssful* delivery time.

the MDP. Suppose that we can propose a waiting time decision  $\mathbb{Z}(\mathbb{Z}), \mathbb{Z} \in \mathbb{R}$  and the action value function of the stationary policy  $\mathbb{Z}, \mathbb{Z}, \ldots$  that is the unique solution to the Bellman equation. Then, the policy  $\mathbb{Z}, \mathbb{Z}, \ldots$  is an optimal policy.

Similar to the previous MDP case, we believe that the optimal policy and the optimal value 5 function of our repeated

MDP (27) still satisfies the Bellman equation . Because except that each waiting time is extended to be a stopping time, our repeated MDP (27) has the same components as that of a discounted MDP. The Bellman equation for our repeated MDP (27) is defined as follows:

where 2 is the set of stopping times on the Wiener process 2 such that

where  $F(\mathbb{Z}) = \bigcap_{\mathbb{Z} \to \mathbb{Z}} (\mathbb{Z}_p, \mathbb{Z} \in [0,\mathbb{Z}])$ . In (33),  $\mathbb{Z} + \mathbb{Z}_{\mathbb{Z} + \mathbb{Z}}$  is the next state of estimation error, after a stopping time  $\mathbb{Z}$  and a channel delay  $\mathbb{Z}$ .

$$\mathbb{Z}_0(\mathbb{Z}) \triangleq 0$$
,

$$\mathbb{Z}_{\mathbb{Z}+1}(\mathbb{Z}) \triangleq \mathbb{Z} \ \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}) = \inf \mathbb{Z}(\mathbb{Z};\mathbb{Z}) + \mathbb{Z} \mathbb{E} \left[ \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}+\mathbb{Z}}) \right], \ \mathbb{Z} = 0, 1, 2, \dots$$
(35)

We also denote  $\mathbb{Z}_1, \mathbb{Z}_2, \ldots$  as the optimal stopping time of the problem (35) when  $\mathbb{Z} = 1, 2, \ldots$ , respectively. Then,  $\mathbb{Z}_2$  ( $\mathbb{Z}$ ) =  $\mathbb{Z}$  O( $\mathbb{Z}$ ) is the discounted integrated cost from the first delivery

time (the last transmission was successful) until at most the  $\mathbb{Z}$ th delivery time, where the  $\mathbb{Z}$ th transmission implies that previous  $\mathbb{Z}-1$  transmissions have failed. In addition, the waiting times for the  $\mathbb{Z}$  transmissions are  $\mathbb{Z}_1,\mathbb{Z}_2,\ldots,\mathbb{Z}_{\mathbb{Z}}$ , respectively. Note that  $\mathbb{Z}$  ( $\mathbb{Z}$ ) is the discounted cost about

infinite number of transmissions. Thus, our objective is to exactly solve (35) by figuring out  $\mathbb{Z}_1,\mathbb{Z}_2$ , . . . and show that  $\mathbb{Z}$   $\mathbb{Q}$   $\mathbb{Z}$   $\mathbb{Z}$   $\mathbb{Z}$   $\mathbb{Z}$   $\mathbb{Z}$   $\mathbb{Z}$   $\mathbb{Z}$   $\mathbb{Z}$  .

?

5.3.1 Candidate Solutions to (35). We speculate that each optimal stopping time  $\mathbb{Z}_1, \mathbb{Z}_2, \ldots$  for (35) is a hitting time, or in other words, threshold type, defined as follows:

$$\mathbb{Z}_{\mathbb{Z}} = \inf\{\mathbb{Z} : |\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}}| \ge \mathbb{Z}_{\mathbb{Z}}\}, \, \mathbb{Z} = 1, 2, \dots,$$

$$\mathbb{Z} \ge 0$$
(36)

where  $\mathbb{Z}$ , called the initial state, is the estimation error at the  $\mathbb{Z}$  – 1th delivery time  $\mathbb{Z}_{\mathbb{Z}}$ –1 ( $\mathbb{Z}$  = 1 implies that the last transmission was successful, and the delivery time is  $\mathbb{Z}_{\mathbb{Z}-1,\mathbb{Z}-1}$ ). Next, we aim to find out the sequence of the optimal thresholds  $\mathbb{Z}_1,\mathbb{Z}_2,\ldots$ 

Let us define a function [2] (2) as follows:

$$\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}) = \mathbb{Z}(\mathbb{Z}; 0) + \mathbb{Z}E\left[\mathbb{Z}_{\mathbb{Z}-1}(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}})\right].$$
 (37)

Intuitively,  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})$  is the action value function that chooses 0 waiting time at the first stage, incurs the cost  $\mathbb{Z}(\mathbb{Z}; 0)$ , and chooses the optimal waiting times at the remaining  $\mathbb{Z}-1$  stages. Since the

speculated optimal waiting time (36) is a hitting time,  $\mathbb{Z}$  ( $\mathbb{Z}$ ) =  $\mathbb{Z}$  ( $\mathbb{Z}$ ) if  $|\mathbb{Z}| \geq \mathbb{Z}$ . In addition, we provide an alternative expression of  $\mathbb{Z}(\mathbb{Z};\mathbb{Z})$ :

Lemma 5.

Moreover, if 2 is a hitting time with a threshold 2 given the initial value 2. i.e.,  $2 = \inf_{2 \ge 0} \{2 : |2 + 2 = 1 \}$ , then we have

$$\mathbb{Z}(\mathbb{Z};\mathbb{Z}) = \mathbb{Z}(\mathbb{Z}, mse_{opt}), \tag{39}$$

This statement is technically true if we can show that our action space is a Borel space (We call ② as a Borel space if there exists a complete separable metric space ② and a Borel subset ② ∈ B<sub>ℤ</sub> such that ② is homeomorphic to ② ) [5, Chapter 9]. Examples of a Borel space are R and any real-valued intervals. For showing that our action space is a Borel space, we leave to our future studies.

(45)

where  $\mathbb{P}(\mathbb{P},\mathbb{P}, mse_{opt})$  is defined in (11).

Then, our problem (35) is augmented as the sequence of standard optimal stopping problem [25, Chapter 1]:

where

 $\mathbb{P}_{2}\left(\mathbb{P}+\mathbb{P}_{2},\mathbb{P}+\mathbb{P}_{2}\right)\triangleq\mathbb{P}^{2}\left(\mathbb{P}+\mathbb{P}_{2},\mathbb{P}+\mathbb{P}_{2}\right)+\mathbb{P}\mathbb{E}\left[\mathbb{P}_{2}-1\left(\mathbb{P}+\mathbb{P}_{2}+\mathbb{P}_{2}\right)\right],\tag{41}$ 

$$2^{\sim}(2 + 2_{\mathbb{Z}}, 2 + 2_{\mathbb{Z}}) \triangleq 2 + 2_{\mathbb{Z}} + E[2](2 + 2_{\mathbb{Z}})$$
 mse<sub>opt</sub>, (42)

$$\mathbb{Z}_{\mathbb{Z}} \longrightarrow \mathbb{Z}_{\mathbb{Z}})^2 - \mathsf{mse}_{\mathsf{opt}} \mathbb{Z} \mathbb{Z}. \tag{43}$$

By Lemma 5, for any 2, we have  $2(2;2) = E[2^{(2+2)},2] = E[2^{(2+2)},2] = E[2^{(2+2)},2] + 2] = 2.$ 

According to [20, Chapter 10] and [25, Section 8], the free boundary method implies that the optimal objective function  $\mathbb{Z}$   $(\mathbb{Z},\mathbb{Z})$  should satisfy

1 
$$\mathbb{Z}_2$$

2

-  $\mathbb{Z}_2(\mathbb{Z},\mathbb{Z}) + \mathbb{Z}$ 

-  $\mathsf{mse}_{\mathsf{opt}} = 0, \mathbb{Z} \in (-\mathbb{Z}_2,\mathbb{Z}_2),$ 

2  $\mathbb{Z}^2$ 

The first equation (44) tells that in the continuation set  $(-\mathbb{Z}_2,\mathbb{Z}_2)$ , the infinitesimal operator of  $\mathbb{Z}_2^{\infty}(\mathbb{Z},\mathbb{Z})$  is zero. In the second equation (45), at the stopping set  $(-\infty, -\mathbb{Z}_2] \cup [\mathbb{Z}_2, \infty)$ , the stopping time  $\mathbb{Z}_2$  is zero. The third equation (46) implies that  $\mathbb{Z}_2^{\infty}(\mathbb{Z},\mathbb{Z})$  should be continuously differentiable at the boundary points  $\mathbb{Z} = \pm \mathbb{Z}_2$ . These three equations are then simplified to:

1 
$$-2 \cdot (2) + 2 - mse_{opt} = 0, 2 \in (-2 \cdot , 2 \cdot ),$$
 (47)

2
$$\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}) = \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}), \mathbb{Z} \in (-\infty, -\mathbb{Z}_{\mathbb{Z}}] \cup [\mathbb{Z}_{\mathbb{Z}}, \infty), \tag{48}$$

$$\mathbb{Z}(\mathbb{Z}) = \mathbb{Z}'(\mathbb{Z}) \qquad . \tag{49}$$

? ?=±?<sub>P</sub> ? ?=±?<sub>P</sub>

following results for deriving the sequence  $\mathbb{Z}_1, \mathbb{Z}_2, \ldots$ :

 ${\Bbb Z}$   ${\Bbb Z}''$  ( ${\Bbb Z}_{\Bbb Z}$ ). Combined with Lemma 5, we provide Lemma 6. For all  ${\Bbb Z}=1,2,\ldots$  we have that:

$$\begin{array}{ccc}
2 & 2 & 3 \\
\mathbb{Z}\mathbb{P}(\mathbb{P}) &= \mathbb{L}(\mathbb{P}, \mathbb{P}_{\mathbb{P}}, mse_{opt}) &= -\mathbb{L} & +2mse_{opt}\mathbb{P}.
\end{array}$$
(50)

If |?| > ?, then

The optimal threshold 2 is the positive solution to

Moreover,  $2^{n}(2)$  and  $2^{n}(2)$  are continuous.

- (b) 22 (0) = 0, and 22 (2) + 22 −  $2mseopt \ge 0$  for all  $2 \in [22, \infty)$ .
- (c)  $22'''(2) \ge 0$ , and  $22'''(2) + 42 \ge 0$  for all  $2 \ge 0$ .
- (d) The sequence of thresholds  $\mathbb{Z}_1, \mathbb{Z}_2, \ldots$  is bounded with  $\mathbb{Z}_n \leq p$ 3mse<sub>opt</sub> and is decreasing, thus converges.

Proof. See Appendix D.

5.3.2 Optimality of the Candidate Solution to (35). We finally validate that the hitting time (36) is the optimal solution. Combined with Lemma 6, we have the following result:

Theorem 5. (a) An optimal sequence of waiting times  $\begin{bmatrix} 2 \\ 1 \end{bmatrix} \begin{bmatrix} 2 \\ 2 \end{bmatrix} \cdots$  for  $\begin{bmatrix} 35 \end{bmatrix}$  satisfies  $\begin{bmatrix} 36 \end{bmatrix}$ , and each threshold  $\mathbb{Z}_2$  is the positive root of  $\begin{bmatrix} 52 \\ 2 \end{bmatrix}$ , where  $\mathbb{Z}_0$  ( $\mathbb{Z}$ ) =  $\begin{bmatrix} 0 \\ 2 \end{bmatrix}$  ( $\mathbb{Z}$ ) is updated by (51),  $\mathbb{Z}_0$  ( $\mathbb{Z}$ ) =  $\begin{bmatrix} 0 \\ 2 \end{bmatrix}$ , and  $\mathbb{Z}$  ( $\mathbb{Z}$ ) is updated by (50)(51).

(b) The function  $\mathbb{Z}_{\mathbb{Z}} (\mathbb{Z}) + \mathbb{Z}_3 \mathbb{Z}^3 - \mathbb{Z}^3 = 0$  mse<sub>opt</sub> in  $\mathbb{Z}$  is convex for  $\mathbb{Z} \ge 0$  and strongly convex for  $\mathbb{Z} > 0$ . Therefore, the positive root of  $\mathbb{Z}_{\mathbb{Z}}$  is unique. In addition,  $\mathbb{Z}_{\mathbb{Z}}$  decreases and thus converges.

Theorem 5 (b) is directly shown by Lemma 6. It remains to show that the exact solution provided in Theorem 5 (a) is optimal to the value iteration problem (35).

Proof of Theorem 5 (a). we obtain the two following results:

Lemma 7. We have  $\mathbb{Z}^{\sim}$  ( $\mathbb{Z},\mathbb{Z}$ )  $\leq$   $\mathbb{Z}^{\sim}$  ( $\mathbb{Z},\mathbb{Z}$ ) for any ( $\mathbb{Z},\mathbb{Z}$ )  $\in$  R2 and the iteration number  $\mathbb{Z}$  = 1, 2, . . . .

Proof. See Appendix E.

Definition 1. A function f(w,q) is excessive if  $E[2^{\infty}(2+22,2+22)] \leq 2^{\infty}(2,2)$  for all  $2 \geq 0$  and  $(2,2) \in \mathbb{R}^2$ .

Lemma 8. The negative value function  $-\mathbb{Z} \ \ \mathbb{Z}(\mathbb{Z},\mathbb{Z})$  is excessive for any  $(\mathbb{Z},\mathbb{Z}) \in \mathbb{R}^2$  and the iteration number  $\mathbb{Z} = 1, 2, \ldots$ 

Proof. See Appendix F.

By Lemma 7 and Lemma 8, using Corollary to Theorem 1 in [27, Section 3.3.1], we have that the stopping times  $\mathbb{Z}_1, \mathbb{Z}_2, \ldots$  in (36) are optimal to (40), thus are optimal to (35). This completes the proof of Theorem 5 (a).

5.4 Linear Convergence of Value Iteration to the Repeated MDP (27)

In this section, we will show that the optimal value functions  $\mathbb{Z}_2$  of the value iteration algorithm (35) converge linearly to the optimal value function  $\mathbb{Z}$  of our problem in (27). We have the following result:

Lemma 9. (a) Suppose that the continuation set of  $\mathbb{Z}$  is is bounded by  $\mathbb{Z}$ , i.e., if  $\mathbb{Z}^2 \geq \mathbb{Z}$ , then  $\mathbb{Z} = 0$ .

Then,  $\mathbb{Z}(\mathbb{Z}) \leq \mathbb{Z}$ .

- (b) The function  $\mathbb{Z}_2(\mathbb{Z}) = \mathbb{Z}^2_0(\mathbb{Z})$  satisfies the contraction mapping property, i.e.,  $\|\mathbb{Z}^2_0\| < \infty$ , and  $\|\mathbb{Z}^{2+1}^0 \mathbb{Z}^{2^0}\| \le \mathbb{Z}^2 \|\mathbb{Z}^0\|$ . The Bellman operator  $\mathbb{Z}$  is defined in (33).
- (c)  $2 (2) = 2 \circ 0(2)$  is the unique solution to the Bellman equation  $2 = 2 \circ 2 \circ 0(2)$  (33) (with  $||2|| < \infty$ ). Further,  $||2| \circ 0 2|| \le 2 ||2| \circ 10 2||$ .

By Lemma 9(c),  $\mathbb{Z}$  is the unique solution to the Bellman equation  $\mathbb{Z} = \mathbb{Z}$   $\mathbb{Z}$ . Therefore,  $\mathbb{Z}$  is the optimal value function for the problem (27). Due to the linear convergence of  $\mathbb{Z}_2$  to  $\mathbb{Z}$ , Lemma 6(d) implies that the optimal stopping time for (27) is also a hitting time, where the optimal threshold is  $\mathbb{Z} = \lim_{\mathbb{Z} \to \infty} \mathbb{Z}_{\mathbb{Z}}$ . This completes the proof of Theorem 2. In addition, Lemma 3 implies that  $\operatorname{mse}_{\operatorname{opt}}$  is the solution to  $\operatorname{E}[\lim_{\mathbb{Z} \to \infty} \mathbb{Z} \times \mathbb{Z} \times \mathbb{Z}] = 0$ . These statements combined with Theorem 5 completes the solution to the problem (3).

### 5.5 Discussion

In this section, we compare our proof and technical contributions with some related works and discuss some interesting future directions.

5.5.1 Special Case 1: Reliable Channel [30]. In the special case of a reliable channel ( $\mathbb{Z} = 0$ ),  $\mathbb{Z} = 1$ . The problem (27) is then reduced to:

$$\mathbb{Z}(\mathbb{Z}) \triangleq \inf \mathbb{Z}(\mathbb{Z};\mathbb{Z}). \tag{53}$$

The problem (27) for general  $\mathbb{Z} \ge 0$  is a repeated MDP, because we need to determine multiple correlated waiting times in an epoch, and each waiting time is a stopping time. However, when  $\mathbb{Z} = 0$ , the problem (53) reduces to an MDP, or in other words, an optimal stopping problem with a single waiting time. Note that solving (53) is still nontrivial. We speculate that the optimal waiting time  $\mathbb{Z}$  is a hitting time. Using Lemma 6 (a), the optimal threshold  $\mathbb{Z}$  is the positive root of

which is  $\mathbb{Z} = p3(mse_{opt} - E[\mathbb{Z}])$ . By Theorem 5, the speculated waiting time is optimal. This implies the final result Corollary 1 ([30, Theorem 1]).

Similar studies with a reliable channel are also indicated, e.g., in [21, 28, 32]. The key insight is to solve an optimal stopping time like (53). Our study with an unreliable channel is different from these studies, because we need to solve a problem with multiple correlated stopping times (27). To solve this, we need to analytically solve a value iteration algorithm (35) that includes a sequence of optimal stopping problems. Compared to (53), for each iteration  $\mathbb{Z}$ , our optimal

stopping problem is more challenging to solve, because the optimal value function is a more complicated expression that contains a summation of ② correlated samples.

5.5.2 Special Case 2: Signal-agnostic Sampling [24]. When the sampling time is independent of the Wiener process, each waiting time takes a nonnegative value based on the timing history information, but not the evolution of the Wiener process. The previous problem (27) is reduced from a discounted and repeated MDP to a discounted MDP. The study in [24] has shown that the optimal policy is a threshold policy on the age (i.e., MMSE).

Since the optimal signal-aware sampling policy is different from the optimal signal-aware sampling policy, the proof of solving our problem (27) is different from that of the discounted MDP in [24]. The authors in [24] solve their problems as follows: (i) they first propose a threshold based waiting decision  $\mathbb{Z}(\mathbb{Z}) = \max(\text{ageopt} - \mathbb{Z} - \mathbb{E}[\mathbb{Z}]/(1-\mathbb{Z}), 0)$ , where  $\mathbb{Z}$  is the age state, and age is the optimal average age; (ii) then they show that  $\mathbb{Z}$  and its value function are the unique opt

solution to the Bellman equation:  $\mathbb{P}_{agnostic}(\mathbb{P}) = \inf_{\mathbb{P}} \ge 0 \ \mathbb{P}_{agnostic}(\mathbb{P},\mathbb{P}) + \mathbb{E}[\mathbb{P}_{agnostic}(\mathbb{P} + \mathbb{P} + \mathbb{P})]$ , where  $\mathbb{P}_{agnostic}(\mathbb{P},\mathbb{P}) \triangleq [\mathbb{P}_{agnostic}(\mathbb{P},\mathbb{P}) + \mathbb{P}_{agnostic}(\mathbb{P},\mathbb{P})]$ .

However, such the proof ideas cannot be applied to our case, due to the following challenges that do not appear in [24]: (i) Since each waiting time is a stopping time, solving (27) faces the curse of dimensionality. For example, when  $\mathbb{Z} = 0$ , (27) reduces to (53), but (53) is still an optimal stopping problem. In the signal-agnostic case, (53) is reduced to a convex optimization problem [28, Lemma 7], thus is much easier to solve; (ii)  $\ln [24]$ , the Bellman equation is solvable. Since  $\mathbb{Z}$  is threshold type on the age, it is optimal to wait ( $\mathbb{Z} > 0$ ) only when the last transmission was successful. Thus, the optimal value function is a closed-form expression:  $\mathbb{Z}_{\operatorname{agnostic}}(\mathbb{Z}) = \operatorname{Eh} \int \mathbb{Z} + \mathbb{Z}(\mathbb{Z}) + \mathbb{Z}$ 

?

age ( $\mathbb{Z}(\mathbb{Z}) + \mathbb{Z}'$ )i, where  $\mathbb{Z}'$  is given in Theorem 4. Since the Bellman equation is a minimization  $_{opt}$  over nonnegative values, solving the Bellman equation is the same as comparing a few closed-form expressions. In our case, however, it is hard to compare, because the optimal value function  $\mathbb{Z}(\mathbb{Z})$  is not closed-form. This is due to the randomness of the Wiener process, and we may wait for each sample.

5.5.3 Future Direction 1: Non i.i.d. Channel Failure. When the channel failure is extended from i.i.d. to Markovian, we still believe that the statements in Section 5.1 and Section (5.2) are correct. However, there is a key difference in Section 5.3: the problem (27) (32) is changed to be

$$\mathbb{P}(\mathbb{P};\mathbb{P},1) + (1-\mathbb{P}')\mathbb{E}(\mathbb{P}^2;\mathbb{P},2)|\mathbb{P}^1 = \mathbb{P} + \tilde{\mathbb{O}}(\mathbb{P}^2;\mathbb{P},\mathbb{P})|\mathbb{P}^2 = \mathbb{P},$$

2 (2) = inf 2=22,1,22,2,...

2=3

(55)

'where 2 is the self

transition probability from 22 state to 22 state, and 2 is the self transition probability from 22 state to 22 state. In the non i.i.d. case where  $1 - 2 \neq 2$ , Problem (55) has a changing discount factor. Thus, the Bellman equation and the value iteration algorithm are not well-defined, making this new problem challenging to solve.

5.5.4 Future Direction 2: Non i.i.d. Transmission Delay. Suppose that we consider a Markovian transmission delay. Then, the waiting time not only should depend on the evolution of the Wiener process, but also should depend on the last transmission delay. This is because the last transmission delay effects the next transmission delay. Therefore, the value iteration (35) should be extended as:

$$\mathbb{Z}_{\mathbb{Z}+1,\text{markov}}(\mathbb{Z},\mathbb{Z}) = \inf \mathbb{Z}(\mathbb{Z};\mathbb{Z}) + \mathbb{Z} \mathbb{E} \mathbb{Z}_{\mathbb{Z},\text{markov}}(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}} + \mathbb{Z},\mathbb{Z}), \mathbb{Z} = 0, 1, 2, \dots,$$
 (56)

# 6 CONCLUSION

In this paper, we provide a sampling policy to minimize the mean square estimation error, where the sampler generates the sample at the source and transmits it to the remote estimator over a time-varying channel. We show that the optimal sampling policy is a threshold policy on the instantaneous estimation error, and the threshold is computed efficiently. The curse of dimensionality that originates from the randomness of the Wiener process, channel conditions, and the channel delay is circumvented. We believe that the proof of our main results provides an insight about how to solve a problem with discounted and multiple stopping times.

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## A PROOF OF LEMMA 5

We denote  $\mathbb{Z}_2 \triangleq \mathbb{Z} + \mathbb{Z}_2$  as the Wiener process starting from the initial state  $\mathbb{Z}_0 = \mathbb{Z}$ . Using the definition of  $\mathbb{Z}(\mathbb{Z};\mathbb{Z})$  in (31), we have

∫ 2+2 ( 2

22;2) 
$$\triangleq$$
E 22 27 - mse<sub>opt</sub>(2 + 2)

0

$$\int 2 \qquad \int 2 + 2$$
=
E 22 22 + 22 27 - mse<sub>opt</sub>(2 + 2) . (57)

Using the strong Markov property of the wiener process,  $\{\mathbb{Z}_{\mathbb{Z}}+\mathbb{Z},\mathbb{Z}\geq 0\}$  has the same distribution as  $\{\mathbb{Z}_{\mathbb{Z}}+\mathbb{Z}_{\mathbb{Z}},\mathbb{Z}\geq 0\}$ . The second term of (57) turns to:

the last equality holds because the delay ② is independent of ②₂. By [18, Theorem 2.5.1], 1/3②₂3 –

Combined with (58) and (59), we finally get (38).

Before showing that  $\mathbb{Z}(\mathbb{Z};\mathbb{Z}) = \mathbb{Z}(\mathbb{Z},\mathbb{Z}, \text{mse}_{\text{opt}})$ , we need Lemma 10:

Lemma 10. If a finite stopping time  $\mathbb{Z}$  satisfies that  $\{\mathbb{Z}_{\mathbb{Z}}, 0 \leq \mathbb{Z} \leq \mathbb{Z}\}$  is bounded, then

Proof. By [18, Theorem 2.5.1] and [10, Theorem 8.5.1],  $1/3\mathbb{Z}_2 - 0 \mathbb{Z}_2$  is a martingale for any given positive value  $\mathbb{Z}$ . Note that for any  $\mathbb{Z} = 1, 2, \ldots, \mathbb{Z} \wedge \mathbb{Z}$  is obviously bounded. Then, we have

Since  $\mathbb{Z}$  is finite,  $\mathbb{Z} \wedge \mathbb{Z} \to \mathbb{Z}$ ,  $\mathbb{Z}_{\mathbb{Z}^{N}} \to \mathbb{Z}_{\mathbb{Z}}$  almost surely. Since  $\{\mathbb{Z}_{\mathbb{Z}}, 0 \leq \mathbb{Z} \leq \mathbb{Z}\}$  is bounded, using Dominated convergence theorem [26, Theorem 5.3.3],

Using Monotone Convergence Theorem [26, Theorem 5.3.1],  $E [2 \land 2] \rightarrow E [2]$ . This leads to

$$\lim_{\mathbb{Z} \to \infty} \mathbb{Z} = 0, \qquad (63)$$

$$\mathbb{Z} \to \infty$$

$$0 \qquad 0$$

where 2 is an upper bound of  $\{2, 0 \le 2 \le 2\}$ . So we have

Now we start to prove  $\mathbb{Z}(\mathbb{Z};\mathbb{Z}) = \mathbb{Z}(\mathbb{Z},\mathbb{Z}, \, \text{mse}_{\text{opt}})$ . If  $|\mathbb{Z}| > \mathbb{Z}$ , then  $\mathbb{Z} = 0$ , and  $\mathbb{E}[\mathbb{Z}] \mathbb{Z}^2 = \mathbb{E}[\mathbb{Z}] \mathbb{Z}^2$ . Therefore,

1 
$$\mathbb{Z}(\mathbb{Z};\mathbb{Z}) = \mathbb{E}[\mathbb{Z}] \mathbb{Z}^2 + \mathbb{E}[\mathbb{Z}] \mathbb{Z}^2 - \mathbb{E}[\mathbb{Z}] \text{ mse}_{\text{opt}} = \mathbb{Z}(\mathbb{Z},\mathbb{Z}, \text{ mse}_{\text{opt}}).$$
 (65)

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If 
$$|\mathbb{Z}| \leq \mathbb{Z}$$
, then E  $\mathbb{Z}^2$ . By [10, Theorem 8.5.5], E  $[\mathbb{Z}] = \mathbb{Z}^2 - \mathbb{Z}^2$ . So we have 
$$\mathbb{Z}^2 = \mathbb{Z}^2 - \mathbb{Z}^2 = \mathbb{Z}^2 - \mathbb{Z}^2 = \mathbb{Z}^2$$

Since  $\mathbb{Z}_0 = \mathbb{Z}$  and the Wiener process has strong Markov property, the first term becomes

The last equality holds due to Lemma 10 and [30, Lemma 3]. From [18, Theorem 2.49], we have

$$2 - 2$$
  $2 \pm 2$  with probability , (68)  $=$   $-2 - 2$  with probability .

Then, we have

$$\frac{\mathbb{Z} + \mathbb{Z}}{\mathbb{Z}} = (\mathbb{Z} - \mathbb{Z})^{3} - (\mathbb{Z} + \mathbb{Z})^{3} = -2\mathbb{Z} (\mathbb{Z}_{2} - \mathbb{Z}_{2}), \tag{69}$$

$$\frac{\mathbb{Z} + \mathbb{Z}}{\mathbb{Z}_{3}} = (\mathbb{Z} - \mathbb{Z})^{3} + \mathbb{Z}_{3} = -2\mathbb{Z} (\mathbb{Z}_{2} - \mathbb{Z}_{2}), \tag{69}$$

$$\mathbb{E} \mathbb{Z}_{3} = \mathbb{Z}_{3} + \mathbb{Z}_{3} = \mathbb{Z}_{3} \times \mathbb{Z}_{3} + \mathbb{Z}_{3} \times \mathbb{Z}_{3} + \mathbb{Z}_{3} = \mathbb{Z}_{3} \times \mathbb{Z}_{3} + \mathbb{Z}_{3} \times \mathbb{Z}_{3} \times \mathbb{Z}_{3} \times \mathbb{Z}_{3} \times \mathbb{Z}_{3} \times \mathbb{Z}_{3} + \mathbb{Z}_{3} \times \mathbb{Z}_{3} \times$$

22

This gives

22

22

6

Therefore, if |2| < 2,

This ends our proof.

# B PROOF OF PROPOSITION 1

The proof is modified from [2], but we strictly extends [2] in two-folds: (i) we consider the square estimation error  $(\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z}_{\mathbb{Z}})^2$ , which is a more complicated metric than the age  $\Delta\mathbb{Z}$  considered in [2]. Note that in the special case where the sampling time is independent of the Wiener process,  $\mathbb{E}[(\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z}_{\mathbb{Z}})^2] = \Delta\mathbb{Z}$ ; (ii) The process  $(\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z}_{\mathbb{Z}})^2$  of two consecutive epochs are correlated, while in [2], the process  $\Delta\mathbb{Z}$  of that are independent.

We denote  $2\mathbb{Z} \triangleq 2\mathbb{Z}_{\mathbb{Z}} - 2\mathbb{Z}_{-1,\mathbb{Z}_{-1}}$  as the inter sampling time of the  $2\mathbb{Z}$ th epoch. We also denote  $H_2$  as the history information of sampling times, transmission times and the Wiener process until  $2\mathbb{Z}_{\mathbb{Z}}$ 

□

Then, by the definition of signal-aware in Section 2.2,  $\square$  is bounded by a stopping time, denoted by

( ) =  $\int \mathbb{Z}$  (  $-\hat{}$  )  $^2$  . We denote  $\mathbb{Z}_2 = \int \mathbb{Z}^2 \mathbb{Z}^2$  ( $\mathbb{Z}_2 - \mathbb{Z}^2$  )  $^2\mathbb{Z}_2 = \mathbb{Z}^2$  For simplicity, let us denote  $\mathbb{Z}$  2  $\mathbb{Z}_2$  2  $\mathbb{$ 

 20.00
 ( - 10.1,2 )2
 ( )

 20
 20
 20, and 2 2
 as the largest epoch number 2 such that 2 < 2, i.e., the</td>

number of successful samples attempted until 2. Then, we have

$$\widetilde{O} 1 \leq (^{2}) \widetilde{O} 1 \in [_{2} \times _{1}]$$

$$\widetilde{C} = (^{2}) \widetilde{O} 1 \in [_{2} \times _{2}]$$

$$\widetilde{C} = (^{2}) \widetilde{C} = (^{2})$$

This tells that

$$\frac{\int_{\infty\mathbb{Z}=1}^{\mathbb{Z}} \mathbb{Z}_{2} \leq \mathbb{Z}(\mathbb{Z}) - 1}{\leq \frac{\mathbb{Z}}{\leq 1}} \leq \frac{\int_{\infty\mathbb{Z}=1}^{\mathbb{Z}} \mathbb{Z}_{2} \leq \mathbb{Z}(\mathbb{Z}) + 1}{\mathbb{Z}}.$$
(75)

Then, we have the following lemma:

Lemma 11.

$$E \stackrel{?}{=} ( \qquad ) + \stackrel{?}{=} (?) + 1$$

$$\lim \underline{\qquad \qquad } 0. \qquad (76)$$

$$? \rightarrow \infty \qquad ?$$

Lemma 11 tells that the "residual terms"  $\mathbb{Z}_{\mathbb{Z}(\mathbb{Z})}/\mathbb{Z}$  and  $\mathbb{Z}_{\mathbb{Z}(\mathbb{Z})+1}/\mathbb{Z}$  vanishes as time  $\mathbb{Z}$  goes to infinity. Therefore, instead of  $\mathbb{Z}(\mathbb{Z})$ , we can analyze  $\mathbb{Z}_{\mathbb{Z}}$ . We have

" 
$$\infty$$
 #
$$E \ [? (?)] \ 1 \ \tilde{O} \ 1 \ lim \ sup \ \_= \ lim \ sup \ \_E \ ?? \ ? \rightarrow \infty ? ?$$

Here we denote  $\mathbb{Z}_{\mathbb{Z}} = \mathbb{Z}$  ( $\mathbb{Z}$ ) + 1 for simplicity. We denote  $\mathbb{Z}_{\mathbb{Z},\mathbb{Z}}$ ,  $\mathbb{Z}_{\mathbb{Z},\mathbb{Z}}$  as the integral of ( $\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z}_{\mathbb{Z}}$ )<sup>2</sup> and  $\mathbb{Z}$  -  $\mathbb{Z}_{\mathbb{Z}-1}$  between  $\mathbb{Z}_{\mathbb{Z}-1}$  and  $\mathbb{Z}$ th delivery time of  $\mathbb{Z}$ th epoch, respectively, given that there are  $\mathbb{Z}$  transmissions at  $\mathbb{Z}$ th epoch. Then,

\*(H )

2min.

**⊵**−1

=EH2-1 hE 22,212 ≤22 H2-1i

Condition (2) is because  $2 \le 2 \pmod{(i.e., 2-1)}$  (i.e.,  $2-1 \le 2 \pmod{2}$ ) is fixed given  $1 \le 2 \pmod{2}$ . Similarly,

E 22,212 ≤22

$$=E_{H_{2}-1}E_{2} + H_{2}-1 + I_{2} \leq I_{2}H_{2}-1.$$
 (79)

Í2∞=1 22-1 (1-2)Eh22,2H2-1i

We then find out the lower bound of  $\int_{-\infty}^{\infty} \mathbb{Z} = 1 \mathbb{Z}_2^{-1} \mathbb{Z}_2^{-1} \mathbb{Z}_2^{-1} \mathbb{Z}_2^{-1}$ , equals to  $\mathbb{Z}_{min}$ , in the following equations (80). Here, condition (i) is due to monotone convergence theorem, and condition (ii) is due to (78). The

value 
$$\mathbb Z$$
  $\mathbb Z-1$  is the minimum of the fraction  $(\mathbb Z-\mathbb Z) \in \mathbb Z-1$  is the minimum of the fraction  $(\mathbb Z-\mathbb Z) \in \mathbb Z-1$  is the minimum  $\mathbb Z$   $\mathbb Z$ 

Divide  $\mathbb{Z}$  on both sides and take the limit of  $\mathbb{Z}$ , then we can get  $\lim_{\mathbb{Z} \to \infty} \mathbb{Z}^{-1} \to \mathbb{Z}$  by  $\lim_{\mathbb{Z} \to \infty} \mathbb{Z} = \mathbb{Z}$  by  $\lim_{\mathbb{Z} \to \infty} \mathbb{$ 

policy that solves (23). This ends our proof of Proposition 1.

## C PROOF OF LEMMA 11

Denote 2 ==

Using [30, Lemma 3], for any finite stopping time 2, we have

 $0\mathbb{Z}_2^2\mathbb{Z}_2$ , where  $\mathbb{Z}_2^2$  is a stopping time upper bound denoted in Appendix B, independent of  $\mathbb{Z}_2$ . Then,

According to strong Markov property of the Wiener process, 22 m +22, 22 - 22 is independent of 22 m

Using [18, Theorem 2.44 and Theorem 2.48], for any finite stopping time  $\mathbb{Z}$ ,  $\mathbb{E}$   $\mathbb{Z}$  <sup>2</sup> =  $\mathbb{E}$  [ $\mathbb{Z}$ ] and

?

The first equality holds because  $\mathbb{Z}^{\sim}\mathbb{P}$ ,  $\mathbb{Z}^{\sim}\mathbb{P}+1$ ,  $\mathbb{Z}^{\sim}\mathbb{P}+1$  are independent of H<sub>B</sub>-1. By (83),  $\mathbb{Z}$  (0) <  $\infty$ ,  $\mathbb{Z}$  ( $\mathbb{Z}$ ) is monotone decreasing, and  $\mathbb{Z}$  ( $\mathbb{Z}$ )  $\rightarrow$  0 as  $\mathbb{Z}$   $\rightarrow$   $\infty$ . We trivially set  $\mathbb{Z}-1$ ,  $\mathbb{Z}-1$  =  $\mathbb{Z}_{0,\mathbb{Z}^0}$  = 0, so we have  $\mathbb{Z}$  ( $\mathbb{Z}$ )  $\geq$  0 and  $\mathbb{Z}_0$  = 0. We have

So we have E
$$E\left[\mathbb{Z}\left(\mathbb{Z}\right)+1\right] \xrightarrow{\mathbb{Z}} \mathbb{Z}\left(\mathbb{Z}\right) = \mathbb{Z}\left(\mathbb{Z}\left(\mathbb{Z}\right)\right) = \mathbb{Z}\left(\mathbb{Z}\left(\mathbb{Z}\right)\right) = \mathbb{Z}\left(\mathbb{Z}\left(\mathbb{Z}\right)\right) + 1 = \mathbb{Z}_{\mathbb{Z}_{2}-1,\mathbb{Z}_{2}-1}}\left(\mathbb{Z}\right) + 1. \tag{88}$$

$$\mathbb{Z}=1 \qquad \mathbb{Z}=2 \qquad \mathbb$$

Note that  $\mathbb{Z}$  ( $\mathbb{Z}$ ) vanishes to 0 as  $\mathbb{Z} \to \infty$ . Following the same steps as [33, Appendix C1], we have  $\int_{\mathbb{Z}} \mathbb{Z} \left( \mathbb{Z} - \mathbb{Z} \right) \mathbb{Z} \mathbb{E} \left[ \mathbb{Z} \left( \mathbb{Z} \right) \right] / \mathbb{Z} \to 0 \text{ as } \mathbb{Z} \to \infty. \text{ This ends our proof.}$ 

#### D PROOF OF LEMMA 6

In this appendix, for simplicity, we will replace the per stage cost 2(2,2, mse<sub>opt</sub>) by 2(2,2).

## D.1 Preliminary

Definition 2. Let  $\mathbb{Z}_1, \dots, \mathbb{Z}_{\mathbb{Z}}$  as an i.i.d. sequence with the same distribution as the channel delay  $\mathbb{Z}$ , and  $\mathbb{Z}_1 \dots \mathbb{Z}_{\mathbb{Z}}$  as any nonnegative sequence. For any real value  $\mathbb{Z}$ , we denote the event  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) =  $\{|\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}_1}| \geq \mathbb{Z}_1, |\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}_1} + \mathbb{Z}_{\mathbb{Z}_2}| \geq \mathbb{Z}_2, \dots, |\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}_1} + \dots + \mathbb{Z}_{\mathbb{Z}_n}| \geq \mathbb{Z}_n\}$ . If  $\mathbb{Z} = 0$ , we denote  $\mathbb{Z}_0(\mathbb{Z})$  as simply the whole set. Denote  $\mathbb{Z}_{\mathbb{Z}_1 + \mathbb{Z}_{\mathbb{Z}_1} + \dots + \mathbb{Z}_{\mathbb{Z}_{n-1}}}(\mathbb{Z})$  as the conditional probability density function (pdf) of  $\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}_1} + \dots + \mathbb{Z}_{\mathbb{Z}_{n+1}}$  with the condition  $\mathbb{Z}_{\mathbb{Z}_n}^1(\mathbb{Z})$ , multiplied by a probability  $\mathbb{P}(\mathbb{Z}_n \setminus \mathbb{Z})$ . In other words,

Note that  $\mathbb{Z}_{2+2^{n}}(\mathbb{Z})$  is equal to the pdf of  $\mathbb{Z}$  + $\mathbb{Z}_{2}$  at  $\mathbb{Z}$ , since  $\mathbb{Z}_{0}(\mathbb{Z})$  is the whole set.

Lemma 12. Suppose that  $E [2] < \infty$ , and there exists 2 > 0 (which can be arbitrary small), such that  $2 \ge 2$ . Then, the following conditions hold.

(a) For any  $\mathbb{Z} \in \mathbb{R}$ ,  $\mathbb{Z}_{2n}$  ( $\mathbb{Z}$ ) is continuously differentiable in  $\mathbb{Z}$ . In addition,  $\mathbb{Z}_{2n}$  ( $\mathbb{Z}$ ),  $\mathbb{Z}_{2n}$  ( $\mathbb{Z}$ )

are bounded, thus  $\mathbb{Z}_{\mathbb{Z}_2}(\mathbb{Z})$ ,  $\mathbb{Z}_2^{'}$  (2) are both uniformly continuous.

- (b) Almost surely,  $12^{\circ}(2)$ ,  $22^{\circ}(2+\Delta 2) \rightarrow 12^{\circ}(2)$ .
- (c) For all  $\mathbb{Z}, \mathbb{Z} \geq 0$ , we have  $\mathbb{Z}_2 + \mathbb{Z}_{21} + \cdots + \mathbb{Z}_{222+1}$  ( $\mathbb{Z}$ )  $\geq \mathbb{Z}_2 + \mathbb{Z}_{21} + \cdots + \mathbb{Z}_{222+1}$  ( $\mathbb{Z}$ ) for all  $\mathbb{Z} = 0, 1, \ldots$  In addition,  $\mathbb{Z}$

Proof. We first show Lemma 12(a). Note that 2 ≥ 2, and 2 is normally distributed with variance

2. Thus, 
$$\mathbb{Z}_{\mathbb{Z}^n}(\mathbb{Z}) = 1/\sqrt{\mathbb{Z}^{2-0.5\mathbb{Z}_2}/\mathbb{Z}}$$
, and  $\mathbb{Z}_{\mathbb{Z}^n}(\mathbb{Z})$  is bounded. Also,  $\mathbb{Z}^{n'}$  (2) =  $-\mathbb{Z}/\mathbb{Z}^{1.5}\mathbb{Z}^{-0.5\mathbb{Z}_2}/\mathbb{Z}$ ,  $\mathbb{Z}^n$  (2) =  $-\mathbb{Z}/\mathbb{Z}^{1.5}\mathbb{Z}^{-0.5\mathbb{Z}_2}/\mathbb{Z}$ ,  $\mathbb{Z}^n$  (2) =  $-\mathbb{Z}/\mathbb{Z}^{1.5}\mathbb{Z}^{-0.5\mathbb{Z}_2}/\mathbb{Z}$  2 /2 2 are still bounded and continuous. So

 $\emptyset \mid ? + ?_{?1} + \cdots + ?_{?2} \mid \geq ? \cap \mid ? + \Delta ? + ?_{?1} + \cdots + ?_{?2} \mid < ?$ 

We finally show Lemma 12(c) by induction. Note that the initial condition holds because  $2 \times (2 - 2) \ge 2 \times (-2 - 2)$ . Suppose that the hypothesis holds. Then, by (90), for any  $2 \times 2 \ge 0$ ,

$$\begin{bmatrix}
22+221+\cdots+222+1 & (2) & = & 22+221+\cdots+222 & (2) \times 2222+1 & (2-2)22, \\
|2|+221+\cdots+222+222+1 & (2) & = & (2+221+\cdots+222+1 & (-2)
\end{bmatrix}$$

$$\begin{bmatrix}
22+221+\cdots+222+1 & (2) & = & (2+221+\cdots+222+1 & (-2)
\end{bmatrix}$$

$$\begin{bmatrix}
22+221+\cdots+222 & (2) \times (2222+1 & (2-2) - 2222+1 & (-2-2))22
\end{bmatrix}$$

$$\begin{bmatrix}
22+221+\cdots+222 & (2) \times (2222+1 & (2-2) - 2222+1 & (-2-2))22
\end{bmatrix}$$

$$\begin{bmatrix}
(22+221+\cdots+222 & (2) - 22+221+\cdots+222 & (-2)) \times (2222 & (2-2) - 2222 & (-2-2))22 \ge 0, \\
1 & 2 & 1 & 2 & 2+1
\end{bmatrix}$$

$$\begin{bmatrix}
(22+221+\cdots+222 & (2) - 22+221+\cdots+222 & (-2)) \times (2222 & (2-2) - 2222 & (-2-2))22 \ge 0, \\
1 & 2 & 1 & 2 & 2+1
\end{bmatrix}$$

which ends the proof of the first claim. The last equation holds because

To show that [2] [2]+[2]1+...[2]2 ([2]) is continuous and bounded in [2], note that

$$\int \int \mathbb{R} e^{-\frac{\pi}{2}} e^{-\frac{\pi}{2}} e^{-\frac{\pi}{2}} = \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{1}) \qquad \mathbb{R}_{\mathbb{R}_{2}} (\mathbb{R}_{2}) \\
= \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{2}) = \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{2}) \\
= \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{2} - \mathbb{R}_{1}) = \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{2}) = \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{2}) \\
= \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{2} - \mathbb{R}_{1}) = \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{2} - \mathbb{R}_{2}) = \mathbb{R} e^{-\frac{\pi}{2}} (\mathbb{R}_{2} - \mathbb{R}_{2})$$

In the above expression, only the final term  $2_{22}(2-2-2_1-\cdots-2_{2-1})$  is related to 2 and this ' ( ) term is continuously differentiable in 2. Also, 2 2 is bounded, and the above expression is

bounded. Therefore,

2=1

Since 2" (2) is bounded, (2) is bounded and continuous in 2. This ends the proof of Lemma 12.

'(2), we need the following lemma:

For the property of  $\ensuremath{\mathbb{Z}}$ 

(a)  $\mathbb{F}(^{\mathbb{Z}})$  is continuous for all  $^{\mathbb{Z}}$ .

Lemma 13.

(b) The functions  $\mathbb{Z}_2(\mathbb{Z})$ ,  $\mathbb{Z}_2(\mathbb{Z})$  are upper bounded by some functions, respectively, such that for any given  $\mathbb{Z}$ , we have  $\mathbb{Z}_2(\mathbb{Z}) < \infty$  and  $\mathbb{Z}_2(\mathbb{Z} + \mathbb{Z}_2) < \infty$ .

(c) If  $|2| < 22+1 \cdot 22'+1(2) = 222(2,22+1)$ . If |2| > 22+1, we have

$$= 2 \cdot 2 \cdot (2 \cdot 2 \cdot 2 + 1) + 2 \cdot 2 \cdot 2 \cdot (2 + 2 \cdot 2 \cdot 1 \cdot 2 \cdot 2) + 2 \cdot 2 \cdot E \cdot h2 \cdot (2 + 2 \cdot 2 \cdot 1 + 2 \cdot 2 \cdot 2)^{1} |2 + 2 \cdot 2 \cdot 2 \cdot 1$$

$$...$$

$$= 2 \cdot 2 \cdot 2 \cdot (2 \cdot 2 \cdot 2 \cdot 1) + 2 \cdot E \cdot 2 \cdot 2 \cdot (2 + 2 \cdot 2 \cdot 1 + \dots + 2 \cdot 2 \cdot 2 \cdot 1 + 1 - 2) \qquad \text{22.2} (2) \qquad (98)$$

the event  $\mathbb{Z}_{2,1}(\mathbb{Z})$  is the whole set, and the events  $\mathbb{Z}_{2,\mathbb{Z}}(\mathbb{Z})$  for  $\mathbb{Z}=2,\cdots,\mathbb{Z}$  are defined as

$$\mathbb{Z}_{\mathbb{Z},\mathbb{Z}}(\mathbb{Z}) = \{ | \mathbb{Z} + \mathbb{Z}_{\mathbb{Z}_1} | \geq \mathbb{Z}_{\mathbb{Z}}, \cdots, | \mathbb{Z} + \mathbb{Z}_{\mathbb{Z}_1} + \cdots + \mathbb{Z}_{\mathbb{Z}_{\mathbb{Z}}-1} | \geq \mathbb{Z}_{\mathbb{Z}} + 2 - \mathbb{Z} \}, \mathbb{Z} = 1, 2, 3, \cdots.$$

$$(99)$$

Proof. Note that

2(2,22+1)

When  $\mathbb{Z} = 1$ , the free boundary method (47)—(49) implies that  $\mathbb{Z}_1$  is the positive root of  $-2/3\mathbb{Z}$  +

 $2 \text{mse}_{\text{opt}} \mathbb{Z} = 2 \text{E} [\mathbb{Z}] \mathbb{Z}$ , which is  ${}^p 3 \text{(mse}_{\text{opt}} - \text{E} [\mathbb{Z}])$ . Then,  $\mathbb{Z}_1(\mathbb{Z}) = \mathbb{Z}(\mathbb{Z}, \mathbb{Z}_1)$ . By (49),  $\mathbb{Z}_1'(\mathbb{Z})$  is continuous at  $\mathbb{Z} = \pm \mathbb{Z}_1$ , thus continuous at  $\mathbb{Z} \in \mathbb{R}$ . For any given  $\mathbb{Z}$ ,  $\mathbb{Z}_1(\mathbb{Z})$  is bounded by  $\mathbb{E} [\mathbb{Z}] \mathbb{Z}_1'(\mathbb{Z})$  is bound

 $\infty$ , condition (b) holds. Condition (c) trivially holds because we have already set  $\mathbb{Z}_0(\mathbb{Z}) = 0$ .

Now we suppose that the hypothesis holds at  $\mathbb{Z}$ . We will show condition (a)—(c) for the case  $\mathbb{Z} + 1$ . Since function  $\mathbb{Z}_2$  is even, and  $\mathbb{Z}_2$  has a symmetric pdf, we have

that  $\mathbb{Z}_{2}+1(\mathbb{Z})$  is continuous. When  $|\mathbb{Z}| < \mathbb{Z}_{2}+1$ , it is easy to find that

Note that the event  $|\mathbb{Z} + \mathbb{Z}_{21} + \dots \mathbb{Z}_{26}| = \mathbb{Z}$  has zero probability for all index  $\mathbb{Z}$  and real value  $\mathbb{Z}$ . Utilizing the hypothesis that  $\mathbb{Z}_2$  ( $\mathbb{Z}_2$ ) is continuous, condition (b),  $\mathbb{Z}(\mathbb{Z}_2 \mathbb{Z}_2 + 1)$  is continuous, we have Further, when  $|\mathbb{Z}| > \mathbb{Z}_2 + 1$ , by the definition in (37),  $\mathbb{Z}_2 + 1(\mathbb{Z}) = \mathbb{Z}_2 + 1(\mathbb{Z})$  and

$$\frac{2}{2}$$

$$\frac{2}$$

Here, $\mathbb{Z}$  is a number that is between 0 and  $\Delta \mathbb{Z}$ . The third equation holds because  $\mathbb{Z}_{\mathbb{Z}'}$  ( $\mathbb{Z}$ ) is well-defined. The last equation holds due to dominated convergence theorem and the hypothesis conditions (a),(b).

When  $|\mathbb{Z}| < \mathbb{Z}_{\mathbb{Z}+1}$ , we have  $\mathbb{Z}_{\mathbb{Z}'+1}(\mathbb{Z}) = \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z},\mathbb{Z}_{\mathbb{Z}+1})$ . Thus, we directly get condition (c). By the free '( $\mathbb{Z}$ ) is continuous at  $|\mathbb{Z}| = \mathbb{Z}_{\mathbb{Z}+1}$ , thus condition (a) holds. In addition, boundary method (49),  $\mathbb{Z}$ 

?

note that

where  $\mathbb{Z}_1,\mathbb{Z}_2$  are bounded values irrelevant to  $\mathbb{Z}$ . Thus, combined with  $E[\mathbb{Z}^2] < \infty$ , condition (b)

holds. This ends the proof of lemma 13.

Lemma 13 implies that  $\mathbb{Z}_{2}'(\mathbb{Z})$ ,  $\mathbb{Z}_{2}'+1(\mathbb{Z})$  are well-defined. Also, (103) implies that we can inter+1 change the derivative and expectation of  $\mathbb{Z}_{2}$  ( $\mathbb{Z}+\mathbb{Z}_{2}$ ), i.e.,

?

\_\_\_E [
$$\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}})$$
] = E  $\mathbb{Z}^{2}(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}})$ . (104)

### D.2 Proof of Lemma 6(a)

Equations (50) and (51) in Lemma 6(a) are easily shown by (102) and (103) in the proof of Lemma 13. According to (103), to show that  $\mathbb{Z}_{\mathbb{Z}'}(\mathbb{Z}),\mathbb{Z}_{\mathbb{Z}''+1}(\mathbb{Z})$  are continuous, it is equivalent to show  $_{1}$  that  $_{\mathbb{Z}_{2}}$  E  $_{\mathbb{Z}_{2}}$  E  $_{\mathbb{Z}_{2}}$  E  $_{\mathbb{Z}_{2}}$  E  $_{\mathbb{Z}_{2}}$  B  $_{\mathbb{Z}_{2}}$  and  $_{\mathbb{Z}_{3}}$  E  $_{\mathbb{Z}_{2}}$  E  $_{\mathbb{Z}_{2}}$  ( $\mathbb{Z}_{3}$  + $\mathbb{Z}_{2}$ ) are continuous. To show this, we should analyze

the derivative of each term  $\mathbb{E}[2\mathbb{Z}(\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}^1}+\cdots+\mathbb{Z}_{\mathbb{Z}},\mathbb{Z}_{\mathbb{Z}^2}+\mathbb{Z}_{\mathbb{Z}})^{1}]$  in (98). We look at any odd polynomial function  $\mathbb{E}(\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}^1}+\cdots+\mathbb{Z}_{\mathbb{Z}^2})$  with  $\mathbb{E}(\mathbb{Z})=\mathbb{E}(\mathbb{Z}^3)$ .

We are interested in analyzing the derivative

$$1 \lim \underline{\mathbb{L}} \mathbb{E} \mathbb{E} \left( \mathbb{E} + \Delta \mathbb{E} + \mathbb{E}_{\mathbb{P}_1} + \cdots \mathbb{E}_{\mathbb{P}_2} \right)^{\mathbf{1}}_{\mathbb{P}_{\mathbb{P}^{-1}}} \left( \mathbb{E} + \Delta \mathbb{E} \right) - \mathbb{E} \left( \mathbb{E} + \mathbb{E}_{\mathbb{P}_1} + \cdots \mathbb{E}_{\mathbb{P}_2} \right)^{\mathbf{1}}_{\mathbb{P}_2 - 1} \left( \mathbb{E} \right) . \tag{105} \Delta \mathbb{E} \rightarrow 0 \Delta \mathbb{E}$$

For simplicity, we utilize the event 22-1(2) from Definition 2. We partition the whole set to 2 sets:

$$21(2) \triangleq |2 + 221| < 21$$

First, Lemma 12(b), Lemma 13 and dominated convergence theorem give

$$= 1 \text{ lim } \mathsf{E} \Delta \mathsf{D} \left( \mathsf{D} + \mathsf{D}_{\mathsf{D}_1} + \cdots \cdot \mathsf{D}_{\mathsf{D}_{\mathsf{D}}} \right)^{\mathbf{1}}_{\mathsf{D}_{\mathsf{D}-1}} \left( \mathsf{D} + \Delta \mathsf{D} \right) = \mathsf{E} \left( \mathsf{D} + \mathsf{D}_{\mathsf{D}_1} + \cdots \cdot \mathsf{D}_{\mathsf{D}_{\mathsf{D}}} \right)^{\mathbf{1}}_{\mathsf{D}_{\mathsf{D}-1}} \left( \mathsf{D} \right)$$

Δ2→0 Δ2

1 lim \_\_E Δ2 (
$$2 + 2e_1 + \cdots + 2e_2$$
 ),  $21,22 ∈ \{1, 2, \cdots, 2 - 1\}$ . (107) Δ2  $\rightarrow$  0 Δ2

Similarly, for any  $\mathbb{Z} \in \{1, 2, \ldots, \mathbb{Z} - 1\}$ ,

1 lim 
$$\Delta E \Delta P$$
 ( $P + P_{P_1} + \cdots P_{P_n}$ )1  $P_{P_{n-1}}(P)$ ,  $P_n(P + \Delta P)$ 

Δ2→0 2

 $\Delta \mathbb{P} \rightarrow 0 \Delta \mathbb{P}$ 

for simplicity.  $\[ \mathbb{Z}_1, \ldots, \mathbb{Z}_{\mathbb{R}} \]$  are arbitrary finite numbers.

Therefore,

$$\lim_{\Delta \to 0} \Delta \mathbb{P} \left( \mathbb{P} + \mathbb{P}_1 + \cdots \mathbb{P}_{\mathbb{P}_0} \right) \mathbb{P}_{\mathbb{P}^{-1}} \left( \mathbb{P} + \Delta \mathbb{P} \right), \mathbb{P}_0 \left( \mathbb{P} \right) + \mathbb{1}_{\mathbb{P}^{-1}} \left( \mathbb{P} \right), \mathbb{P}_0 \left( \mathbb{P} + \Delta \mathbb{P} \right)$$

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$$= \frac{2}{2+2} + \frac{2}{2} +$$

$$-\frac{\mathbb{P}_{\mathbb{P}^{1}}}{\mathbb{P}^{1}}+\mathbb{P}_{\mathbb{P}^{1}}+\cdots\mathbb{P}_{\mathbb{P}^{2}}\left(-\frac{\mathbb{P}_{\mathbb{P}^{1}}}{\mathbb{P}^{2}}\right)E^{\frac{\mathbb{P}_{\mathbb{P}^{1}}}{\mathbb{P}^{2}}}+\frac{\mathbb{P}_{\mathbb{P}^{2}}}{\mathbb{P}^{2}}+\mathbb{P}_{\mathbb{P}^{2}}$$

h i

which is a constant term multiplied by  $2\mathbb{Z}+2\mathbb{Z}+\cdots\mathbb{Z}$  ( $2\mathbb{Z}$ )  $+2\mathbb{Z}+2\mathbb{Z}+\cdots\mathbb{Z}$  ( $-2\mathbb{Z}$ ), where  $2\mathbb{Z}+2\mathbb{Z}+\cdots\mathbb{Z}$  ( $2\mathbb{Z}$ ) 1

is defined in Definition 2. The last equation in (112) holds because  $\mathbb{Z}_{\mathbb{Z}}$  is symmetric and  $\mathbb{Z}$  (·) is odd . Recall that  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z},\mathbb{Z})$  contains two odd polynomial terms 2E  $[\mathbb{Z}]$   $\mathbb{Z}$  and  $\mathbb{Z}$  and  $\mathbb{Z}$  and  $\mathbb{Z}$   $\mathbb{Z}$ 

8

related to  $\mathbb{Z}$ . Therefore, the derivative of  $\mathbb{E}[\mathbb{Z}\mathbb{Z}(\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}^1}+\cdots+\mathbb{Z}_{\mathbb{Z}},\mathbb{Z}_{\mathbb{Z}^1}+\mathbb{Z}_{\mathbb{Z}})^1$  that appears in (98) is expressed as the sum of forms (106), (107), and (112). The value of (107) is 0. According to Lemma 12(c), the value of (112) is a constant multiplied by  $\mathbb{Z}_{\mathbb{Z}^1+\mathbb{Z}^2}+\cdots+\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})$ , a continuously differentiable function in  $\mathbb{Z}$  for some parameter  $\mathbb{Z}$ . For (106), note that the term of (106) is continuous in  $\mathbb{Z}$ . We can take the derivative and apply the previous calculations (105)—(112) again . Then, the term

# D.3 Proof of Lemma 6(b),(c)

We then use induction to prove Lemma 6(b),(c). Let us denote ② = mse<sub>opt</sub> − E [②], Note that the free

boundary method implies that  $\mathbb{Z}_1$  3 $\mathbb{Z}$ , and  $\mathbb{Z}$   $\mathbb{Z},\mathbb{Z}_1$  is continuously differentiable. In addition,  $2\mathsf{E}$  [ $\mathbb{Z}$ ]  $\mathbb{Z} > \mathbb{Z}_1$ ,

$$2 (2,2) = 1$$
 (113)

$$-2\mathbb{Z}_2 + 2 \text{mseopt}$$
  $0 \le \mathbb{Z} < \mathbb{Z}_1$ . (114)

$$[2,2](2,2) = -42 0 \le 2 < 2_1,$$

Then, we have 2 (2) = 2E [2] and for all  $2 \ge 2$ 1,

For all  $\mathbb{Z} \geq 0$ ,

$$\mathbb{Z}_{1}(\mathbb{Z}) - (-4\mathbb{Z}) = 0 + 4\mathbb{Z} \ge 0. \tag{116}$$

This satisfies the initial condition of Lemma 6(b),(c). By Lemma 13,

Note that if ② (·) is even, then the first term of the last equation in (112) becomes ☑ + □ 2 + □ 2 ■ (22)-22+22 +---22

9Despite that 2 (2) becomes an even polynomial function with 2 (22), except the minor sign change of the last equality of (112) as described in the previous footnote, the calculations (105)—(112) remain the same.

≤∆ Δ≤ Since 2 is continuous, there exists 0 2 or 2 2 0, such that ? 1

We have shown that 22 = [22 (2 +22)] is continuous, so

"( + +  $\Delta$ )

the proof of Lemma 6(a) described in Appendix D.2, both 22 2 22 and 22 22 22 are bounded by a finite random variable. Therefore, using dominated convergence theorem,

By Lemma 12, the two remaining events vanishes as  $\Delta \mathbb{Z} \to 0$ . Thus, for small  $\Delta \mathbb{Z}$ , we have

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|₽₂−₽ |≤|Δ₽ |,₽≠₽≥

By (119),(120), we have an interesting result:

221(2) = 2E[2] + E[2]'(2 + 2). +

Then, we consider the third derivative:

For this derivation, there exists  $0 \le \mathbb{Z} \le \Delta \mathbb{Z}$  or  $\Delta \mathbb{Z} \le \mathbb{Z} \le 0$ , such that

???

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2+1<sub>2</sub>

$$E''(+\Delta +)-''(+)$$

"'(2 + 22 +  $\Delta$ 2) and 22"'(2 + 22 +  $\Delta$ 2) are bounded by a finite random variable. Therefore, 2

using dominated convergence theorem (similar to (119)),

=E lim 
$$\mathbb{Z}^{'''}$$
(? + 2 +?2)  $\times$  1|2+ $\Delta$ 2+22|> $_{\mathbb{Z}^2}$ ,|2+22|> $_{\mathbb{Z}^2}$ 

$$=E ??"(? +??) \times 1|?+??| >??$$

We then discuss the two remaining events. If  $\Delta \mathbb{Z} > 0$ ,

$$= 2e''(2e) + 2e^2 - 2mse_{opt} + 2(\Delta e) \qquad \underline{\qquad} 1 \ 1e^2 - 2\Delta e^2 < 2e - 2e, \qquad (126) \ \Delta e$$

If  $\Delta \mathbb{Z} < 0$ ,

1

...

2 "(2 + 
$$\Delta$$
2 +22) - 22"(2 +22) 1|2+ $\Delta$ 2+22|>22, |2+22|<22

 $\Delta$ 2

=  $2$ 2"(-22) +  $2$ 22-2 - 2mse<sub>opt</sub> +  $2$ ( $\Delta$ 2)

1

...

2 "(2 +  $\Delta$ 2 +22) -  $2$ 2"(2 +22) 1|2+ $\Delta$ 2+22|<22, |2+22|>22

$$= - \mathbb{Z}_{\mathbb{Z}}^{n}(\mathbb{Z}_{\mathbb{Z}}) + 2\mathbb{Z}_{\mathbb{Z}}^{2} - 2\mathsf{mse}_{\mathsf{opt}} + \mathbb{Z}(\Delta\mathbb{Z}) \qquad \underline{\qquad} 1 \, 1\mathbb{Z}_{\mathbb{Z}} - \mathbb{Z} < \mathbb{Z}_{\mathbb{Z}} - \mathbb{Z} - \Delta\mathbb{Z}. \tag{129} \Delta\mathbb{Z}$$

Therefore,

The last inequality of (130) holds due to the induction hypothesis of  $2^{n}$  (·) and Lemma 12. Combining

(125) and (130), for all  $2 \ge 0$ ,

Δ?

$$\begin{aligned} & = 22^{m} + 1(2) = 0 + 2E \ 22^{m} (2 + 22) 1_{2+22} | > 22 + 22^{m} (2 + 22) 1_{2+22} | < 22 \\ & = 22^{m} (22) + 2222 - 2mseopt (22 (-22 - 2)) - 222 (22 - 2)) \end{aligned}$$

$$& \geq 2E \ 22^{m} (2 + 22) 1_{2+22} | > 22 + 22^{m} (2 + 22) 1_{2+22} | < 22 \\ & = 2E \ 22^{m} (2 + 22) 1_{2+22} | > 22 - 4(2 + 22) 1_{2+22} | < 22 \\ & = 2E \ 22^{m} (2 + 22) 1_{2+22} | > 22 - 4(2 + 22) 1_{2+22} | < 22 \\ & = 2E \ 22^{m} (2 + 22) + 4(2 + 22) | 2 + 22 | > 22 - 4(2 + 22) 1_{2+22} | < 22 \\ & = 2E \ 22^{m} (2 + 22) + 4(2 + 22) | 2 + 22 | > 22 - 4(2 + 22) 1_{2+22} | < 22 \\ & = 2E \ 22^{m} (2 + 22) + 4(2 + 22) | 2 + 22 | > 22 - 4(2 + 22) 1_{2+22} | < 22 \\ & = 2E \ 22^{m} (2 + 22) + 4(2 + 22) | 2 + 22 | > 22 - 4(2 + 22) 1_{2+22} | < 22 \\ & = 2E \ 22^{m} (2 + 22) + 4(2 + 22) | 2 + 22 | > 22 - 4(2 + 22) | 2 + 22 | > 22 - 4(2 + 22) | < 22 | > 22 - 22 | > 22 \\ & = 2E \ 22^{m} (2 + 22) + 4(2 + 22) | 2 + 22 | > 22 - 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22 | > 22$$

Note that 2'''(2) + 42 is an odd function, and by hypothesis,  $22'''(2) + 42 \ge 0$  for all  $2 \ge 0$ . By

Finally, (131) and (132) give

$$22'''(2) \ge 2E - 4(2 + 22) 1 | 2 + 22 | \ne 22$$

$$= 2E [-4(2 + 22)]$$

$$= -422$$

$$\ge -42.$$
(133)

The last inequality is strict if  $\mathbb{Z} > 0$ . This ends the proof of Lemma 6(c).

Let us define  $^{\mathbb{Z}}_{\mathbb{Z}+1}(^{\mathbb{Z}}) = ^{\mathbb{Z}}_{\mathbb{Z}+1}(^{\mathbb{Z}}) + 2/3^{\mathbb{Z}}_{3} - mse_{opt}\mathbb{Z}$  for simplicity. Then,  $^{\mathbb{Z}}_{\mathbb{Z}}$  and  $^{\mathbb{Z}}_{\mathbb{Z}}$   $^{\mathbb{Z}}_{\mathbb{Z}}$  1 for  $\mathbb{Z} \in [0, \infty)$ . This implies that  $\mathbb{Z}_{\mathbb{Z}}+1(\mathbb{Z})$  is convex in  $\mathbb{Z} \in [0, \infty)$  and strictly convex in  $\mathbb{Z} > 0$ . By Lemma 6(a),  $\mathbb{Z}$ 

 $\mathbb{Z}'+1(\mathbb{Z})$  is continuous and odd. Thus,  $\mathbb{Z}_{\mathbb{Z}}+1(0)=0$ . By the definition of free boundary

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 $2 \in (0, 2+1)$ . This ends the proof of Lemma 6(b). D.4

Proof of Lemma 6(d)

Now we show Lemma 6(d). We now use induction to show that  $\mathbb{Z}_2 \leq p^2$  3mse<sub>opt</sub> for all  $\mathbb{Z} = 1, 2, \cdots$ .

Note that  $\mathbb{Z}_1 = p3\mathbb{Z} = p3 \text{ (mse}_{opt} - \mathbb{E} \left[ \mathbb{Z} \right] ) \leq p3 \text{ mse}_{opt}$ . The second threshold  $\mathbb{Z}_2$  is the root of  $\frac{2}{0} = \mathbb{Z} - 2\mathbb{Z} + \mathbb{Z}$   $3 \qquad \qquad \mathbb{E} \mathbb{Z} \mathbb{Z} \mathbb{Z} \left( \mathbb{Z} + \mathbb{Z} \mathbb{Z}_1, \mathbb{Z}_1 \right) .$  3

Note that  $2mse_{opt}$ ?  $2s is positive at <math>0 \le 2 \le p3mse_{opt}$ . Therefore, if  $2 \le p3mse_{opt}$ ,  $2s = p3mse_{opt}$ , 2s =

is always positive at  $\mathbb{Z} \geq 0$ . Since  $\mathbb{Z}_1 \leq p$ 3mse<sub>opt</sub>,  $\mathbb{Z}_2\mathbb{Z}(\mathbb{Z},\mathbb{Z}_1) \geq 0$  for all  $\mathbb{Z} \geq 0$ . Recall that  $\mathbb{Z}_2\mathbb{Z}(\mathbb{Z},\mathbb{Z})$  is an odd function on  $\mathbb{Z}$  for any  $\mathbb{Z} \geq 0$ . Therefore, utilizing the same analysis as (132), for all  $\mathbb{Z} \geq 0$ ,

The first term  $\sqrt[3]{2} - 2\mathbb{Z} > 0$  for all  $\mathbb{Z} \ge p3mse_{opt}$ . To keep the equation (134) holds, we have  $\mathbb{Z}_2 \le p3mse_{opt}$ .

Suppose that  $\mathbb{Z}_2, \dots, \mathbb{Z}_2 \leq p$ 3mse<sub>opt</sub>. Now, we will show that  $\mathbb{Z}_{2^{+1}} \leq p$ 3mse<sub>opt</sub>. Note that  $\mathbb{Z}_{2^{+1}}$  is the root of

?=2

Since the hypothesis tells that  $22 \le p_{3mse_{opt}}$ , E  $22(2 + 22_{1}, 22)$ . To show that  $22 + 1 \le p_{3mse_{opt}}$ .

 $p_{\text{3mse}_{\text{opt}}}$ , it is sufficient to show that for  $2 = 2, \dots, 2$ ,

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$$\mathbb{E} \, \mathbb{E}_{\mathbb{Z}^{2}}(\mathbb{Z} + \mathbb{Z}_{21} + \cdots \times \mathbb{Z}_{22}, \mathbb{Z}_{2} + 1 - \mathbb{Z}) \, \mathbf{1}_{\{|\mathbb{Z} + \mathbb{Z}_{1}| \geq \mathbb{Z}_{2}, \cdots, |\mathbb{Z} + \mathbb{Z}_{21} + \cdots + \mathbb{Z}_{22-1}| \geq \mathbb{Z}_{2+2-2}\}} \geq 0. \tag{137}$$

Since  $\mathbb{Z}_{\mathbb{R}^{+1}-\mathbb{Z}} \leq p3mse_{opt}$ ,  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z},\mathbb{Z}_{\mathbb{R}^{+1}-\mathbb{Z}}) \geq 0$  for all  $\mathbb{Z} \geq 0$ . Therefore, the inequality (137) is shown by Lemma 12(c) and that  $\mathbb{Z}_{\mathbb{Z}}(\mathbb{Z},\mathbb{Z})$  is an odd function for any  $\mathbb{Z} \geq 0$ .

Now, we will jointly show that  $\mathbb{Z}^{1}$  +  $(\mathbb{Z}) \geq \mathbb{Z}^{2}$  ( $\mathbb{Z}$ ), and  $\mathbb{Z}^{1}$  +  $\mathbb{Z}^{2}$ . First,  $\mathbb{Z}_{1} = p3\mathbb{Z}, \mathbb{Z}_{1}$  ( $\mathbb{Z}$ ) =

2E [2]  $\mathbb{Z},\mathbb{Z}_{2}$  ( $\mathbb{Z}_{2}$ ) +  $\mathbb{Z}_{2}$  [ $\mathbb{Z}_{2}$ ( $\mathbb{Z}_{2}$ ) +  $\mathbb{Z}_{2}$  [ $\mathbb{Z}_{2}$ ( $\mathbb{Z}_{2}$ )]. Since  $\mathbb{Z}_{1} \leq \mathsf{p3mse}_{\mathsf{opt}}$ , E [ $\mathbb{Z}_{2}$ ( $\mathbb{Z}_{2}$ + $\mathbb{Z}_{2}$ , $\mathbb{Z}_{1}$ )]  $\geq 0$ ,

directly have 2. For simplicity, let us define 22

$$\mathbb{P}_{\mathbb{P}}(\mathbb{P}) = \mathbb{P}^* \qquad \mathbb{P}^3 + 2\mathsf{mse}_{\mathsf{opt}}\mathbb{P}). \tag{138}$$

If  $\mathbb{Z}_2 > \mathbb{Z}_1$ , then we have  $\mathbb{Z}_1(\mathbb{Z}_2) \leq \mathbb{Z}_2(\mathbb{Z}_2) = 0$ , which contradicts to  $\mathbb{Z}_1(\mathbb{Z}) > 0$  for  $\mathbb{Z} > \mathbb{Z}_1$ . Therefore,  $\mathbb{Z}_2 \leq \mathbb{Z}_1$ .

Then, suppose that 2

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 $_{\mathbb{Z}^{'}}(\mathbb{Z})$  is odd. If  $|\mathbb{Z}+\mathbb{Z}_{\mathbb{Z}}|\geq\mathbb{Z}_{\mathbb{Z}^{-1}}$ , utilizing the same analysis as (132), we have Note that  $\mathbb{Z}$ 

$$\mathsf{E}\left( \mathsf{PP'}\left( \mathsf{P} + \mathsf{PP} \right) - \mathsf{PP'} - \mathsf{1} \left( \mathsf{P} + \mathsf{PP} \right) \right) \mathsf{1} \left| \mathsf{P} + \mathsf{PP} \right| \geq \mathsf{PP} - \mathsf{1}$$

$$= \mathbb{E} \left( \mathbb{Z}_{2}^{'} \left( \mathbb{Z} + \mathbb{Z}_{2} \right) - \mathbb{Z}_{2}^{'} - 1 \left( \mathbb{Z} + \mathbb{Z}_{2} \right) \right)^{1} | \mathbb{Z} + \mathbb{Z}_{2} | \ge \mathbb{Z}_{2-1} \ge 0. \tag{140}$$

If  $|\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}}| < \mathbb{Z}_{\mathbb{Z}}$ , we have  $\mathbb{Z}_{\mathbb{Z}'}(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}}) = 2 \text{mse}_{\text{opt}}(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}}) - 2/3(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}})^3$ , irrelevant to  $\mathbb{Z}$ . Thus

$$E(2^{\circ}(2+2^{\circ})-2^{\circ}-1(2+2^{\circ}))1|2+2^{\circ}|<2^{\circ}$$

$$= \mathbb{E} \left[ \mathbb{Z}^{2} \left( \mathbb{Z} + \mathbb{Z}_{2} \right)^{1} \right] \mathbb{Z} + \mathbb{Z}_{2} \left[ \mathbb{Z} - \mathbb{Z}^{2} \right]^{-1} \left( \mathbb{Z} + \mathbb{Z}_{2} \right)^{1} \mathbb{Z} + \mathbb{Z}_{2} = 0. \tag{141}$$

If  $2 \le |2 + 2 \le | < 2 \le 1$ ,

$$E(2^{\circ}(2+2)-2^{\circ}-1(2+2))12^{\circ} \le |2+2^{\circ}| < 2^{\circ}-1$$

$$=E(2^{\circ}(2+2)-2^{\circ}-1(2+2))1_{22\leq|2+22|<22-1}$$

$$= \mathbb{E}\left(\mathbb{Z}_2^{'}(\mathbb{Z} + \mathbb{Z}_2) - 2\mathsf{mse}_{\mathsf{opt}}(\mathbb{Z} + \mathbb{Z}_2) + \mathbb{Z}_2^{\mathsf{p}}\right) + \mathbb{Z}_2^{\mathsf{p}} + \mathbb{Z}_2^{\mathsf{p}} + \mathbb{Z}_2^{\mathsf{p}} = \mathbb{Z}_2^{\mathsf{p}} + \mathbb{Z}_2^{\mathsf{p}} = \mathbb{Z}_2^{\mathsf{p}} + \mathbb{Z}$$

The last inequality holds because  $\mathbb{Z}_{\mathbb{Z}}$  ( $\mathbb{Z}$ ) is an odd function and non-negative for  $\mathbb{Z} \geq \mathbb{Z}_{\mathbb{Z}}$ . Inserting  $\mathbb{Z}'+1(\mathbb{Z}) \geq \mathbb{Z}_{\mathbb{Z}'}$  ( $\mathbb{Z}$ ) for  $\mathbb{Z} \geq 0$ .

(140),(141),(142) into (139), we finally have 2

Recall that the free boundary method (49) implies that  $\mathbb{Z}_{\mathbb{Z}^{1}}$  is the root of  $\mathbb{Z}_{\mathbb{Z}^{1}}$   $\mathbb{Z}_{\mathbb{Z}^{2}}$   $\mathbb{Z}_{\mathbb{Z}^{2}}$ 

 $\mathbb{Z}'(\mathbb{Z}_{\mathbb{Z}^+1}) \leq \mathbb{Z}_{\mathbb{Z}'+1}(\mathbb{Z}_{\mathbb{Z}^+1}) = 0, \text{ which contradicts to } \mathbb{Z}_{\mathbb{Z}'}(\mathbb{Z}) > 0 \text{ for } \mathbb{Z} > \mathbb{Z}_{\mathbb{Z}}. \text{ Therefore, then we have } \mathbb{Z}$ 

 $\mathbb{Z}_{2+1} \leq \mathbb{Z}_{2}$ , and we have that  $\{\mathbb{Z}_{2}\}_{2}$  is decreasing.

## E PROOF OF LEMMA 7

By (45),  $\mathbb{Z} \otimes (\mathbb{Z}, \mathbb{Z}) = \mathbb{Z} \otimes (\mathbb{Z}, \mathbb{Z})$  if  $|\mathbb{Z}| > \mathbb{Z}_{\mathbb{Z}}$ . It remains to show that  $\mathbb{Z} \otimes (\mathbb{Z}, \mathbb{Z}) \leq \mathbb{Z} \otimes (\mathbb{Z}, \mathbb{Z})$  for  $|\mathbb{Z}| \leq \mathbb{Z} \otimes (\mathbb{Z})$  (by symmetry, we will assume  $\mathbb{Z} \geq 0$ ).

Define  $\mathbb{Z}(\mathbb{Z}) \triangleq \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z},\mathbb{Z}) - \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z},\mathbb{Z}) = \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z}) - \mathbb{Z}_{\mathbb{Z}}(\mathbb{Z})$ . It is easy to see that  $\mathbb{Z}(\mathbb{Z}) = 0$ , and  $\mathbb{Z}(\mathbb{Z})$  is not a function of  $\mathbb{Z}$ .

By Lemma 6, we have shown that  $2\mathbb{Z}'(\mathbb{Z}) = -2/3\mathbb{Z}^3 + 2 \text{mse}_{\text{opt}}\mathbb{Z}$ ,  $2\mathbb{Z}''(\mathbb{Z}) = -2\mathbb{Z}^2 + 2 \text{mse}_{\text{opt}}$ ,  $2\mathbb{Z}'''(\mathbb{Z}) = -4\mathbb{Z}$  if  $|\mathbb{Z}| \leq \mathbb{Z}$ . Therefore, by Lemma 6(c),  $2\mathbb{Z}'''(\mathbb{Z}) \geq 0$  for  $2\mathbb{Z} \in [0,\mathbb{Z}]$ . This implies that  $2\mathbb{Z}'(\mathbb{Z})$  is convex in  $2\mathbb{Z} \in [0,\mathbb{Z}]$ . Since  $2\mathbb{Z}'(0) = 2\mathbb{Z}'(2\mathbb{Z}) = 0$ , we have  $2\mathbb{Z}'(2) \leq 0$  for  $2\mathbb{Z} \in [0,\mathbb{Z}]$ . Note that  $2\mathbb{Z}'(2\mathbb{Z}) = 0$ . So  $2\mathbb{Z}'(2\mathbb{Z}) = 0$ .

increasing in  $\mathbb{Z} \in [0, \mathbb{Z}_2]$ , and thus  $\mathbb{Z}$  ( $\mathbb{Z}$ )  $\geq 0$  for  $\mathbb{Z} \in [0, \mathbb{Z}_2]$ . This implies that  $\mathbb{Z}^{\sim}$  ( $\mathbb{Z}$ ,  $\mathbb{Z}$ )  $-\mathbb{Z}^{\sim}$  ( $\mathbb{Z}$ ,  $\mathbb{Z}$ )  $\geq 0$ , which completes our proof.

## F PROOF OF LEMMA 8

 $\mathbb{Z}^{\sim}$  ( $\mathbb{Z}$ , $\mathbb{Z}$ ) is continuously differentiable, and twice condinuously differentiable except at ( $\pm \mathbb{Z}_{\mathbb{Z}}$ , $\mathbb{Z}$ ).

However, since the Lebesgue measure of reaching  $(\pm 2 \, \mathbb{Z}, \mathbb{Z})$  is zero, the values  $(\pm 2 \, \mathbb{Z} \times 2 \, \mathbb$ 

In Lemma 6, it is easy to see that  $\mathbb{Z}_2 \mathbb{Z}_2 (\mathbb{Z}, \mathbb{Z}) = \mathbb{Z}_2 (\mathbb{Z})$ , not a function of  $\mathbb{Z}$ , and  $\mathbb{Z}_2 \mathbb{Z}_2 (\mathbb{Z}, \mathbb{Z}) = \mathbb{Z}(\mathbb{Z})$ . Therefore, for any given time  $\mathbb{Z}$ ,

$$\int \mathbb{P} \qquad \qquad 2$$

$$E \qquad \qquad \tilde{\mathbb{P}} \mathbb{P} \mathbb{P} \left( \mathbb{P} + \mathbb{P}_{\mathbb{P}}, \mathbb{P} + \mathbb{P}_{\mathbb{P}} \right) \mathbb{P} < \infty. \tag{143}$$

∫ ☑ 2

The integral ☑ ② ② is increasing in ②. Using Itô's formula [18, Theorem 7.14], almost surely,

$$\mathbb{E}_{\mathbb{Z}} \mathbb{Z}^{\sim}_{\mathbb{Z}} (\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}}, \mathbb{Z} + \mathbb{Z}_{\mathbb{Z}}) \mathbb{Z} \mathbb{Z} = 0. \tag{145}$$

Therefore,

(146)

If 
$$|\mathbb{Z} + \mathbb{Z}_2| < \mathbb{Z}_2$$
, according to Lemma 6(a) we have  $\mathbb{Z}_2 \mathbb{Z}_2$  ( $\mathbb{Z} + \mathbb{Z}_2$ ) =  $-2/3(\mathbb{Z} + \mathbb{Z}_2)^3 + \mathbb{Z}_2$ 

 $2mse_{opt}(\mathbb{Z} + \mathbb{Z}_{\mathbb{Z}})$ , and  $\mathbb{Z}_{\mathbb{Z}} \mathbb{Z} \mathbb{Z} \mathbb{Z} \mathbb{Z} \mathbb{Z} \mathbb{Z} + \mathbb{Z}_{\mathbb{Z}} \mathbb{Z} + \mathbb{Z}_{\mathbb{Z}} \mathbb{Z} + \mathbb{Z}_{\mathbb{Z}} \mathbb{Z} + 2mse_{opt}$  (correspond to the first equation of free boundary method (44)). Therefore,

$$(2 + 22)_2 - mseopt + 222 2^2 (2 + 22, 2 + 22) = 0.$$
 (147)

If  $| \mathbb{Z} + \mathbb{Z}_2 | \geq \mathbb{Z}_2$ , according to Lemma 6(a),(b), we get

$$2((2 + 2 2) 2 - mse_{opt}) + 2 2 2 2 2 (2 + 2 2, 2 + 2 2)$$

$$= 2((2 + 2 2)^{2} - mse_{opt}) + 2 2 (2 + 2 2) \geq 0.$$

$$= E(^{2}, 2)$$

$$= (2 + 2 2)^{2} - mse_{opt}) + 2 2 (2 + 2 2) \geq 0.$$

$$= (148)$$

$$= (2 + 2 2)^{2} - mse_{opt}) + 2 2 (2 - 2) = 0.$$

$$= (2 + 2 2)^{2} - mse_{opt}) + 2 2 2 (2 - 2) = 0.$$

Applying to (146), we get

G PROOF OF LEMMA 9

Note that for simplicity, we have set 2 = 2 + 2 = 3 as a Wiener process that starts from 2 = 2.

2
When 2 < 2,

2 -

When 2 < 2, 2 = 2, and we have

This ends the proof of Lemma 9(a).

Note that  $\mathbb{Z}(2;\mathbb{Z})$  is bounded in  $\mathbb{Z}\in[-\mathbb{Z}, \mathbb{Z}]$ , and  $\mathbb{Z}(2;\mathbb{Z})=2\mathbb{E}[\mathbb{Z}]\mathbb{Z}^2+2\mathbb{Z}\mathbb{Z}\mathbb{Z}$  for  $|\mathbb{Z}|^2\geq\mathbb{Z}$ . Therefore, there exists  $\mathbb{Z}>0$  such that  $||\mathbb{Z}(2;\mathbb{Z})||\leq\mathbb{Z}$ . Recall that we denote  $\mathbb{Z}_{\mathbb{Z}+1}$  as the state value at  $\mathbb{Z}+1$ th stage with  $\mathbb{Z}_{1}=\mathbb{Z}$ , and  $\mathbb{Z}_{2+1}=\mathbb{Z}_{2}+\mathbb{Z}_{2}+\mathbb{Z}_{2}$  for a stopping time  $\mathbb{Z}'$ . Then,

We have shown that each optimal stopping times for solving  $\mathbb{Z}(\mathbb{Z}^20(\mathbb{Z}))$  are some hitting times with bounded and convergent thresholds, so each stopping time belongs to the assumption of Lemma 9(a). We have

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Thus, both  $\|\mathbb{Z}^0\|$  and are finite. For any stopping time within the assumption of Lemma 9(a),

This gives that

Take the minimum for left and right side of (154) over all the stopping times  $\ensuremath{\mathbb{Z}}$  with bounded

threshold 2, then

By symmetry,

$$2 O(2)$$
 $2 O(2)$ 
 $2 O(2)$ 

Therefore,

$$\|20+10-200\| \le 2\|200-20-10\| \cdot \cdot \cdot \le 20\|20\|.$$
 (157)

This completes the proof of Lemma 9(b). Due to (157), the sequence  $\{\mathbb{Z}^2 \circ (\mathbb{Z})\}$  is a Cauchy sequence, and thus  $\mathbb{Z} \circ (\mathbb{Z})$  converges pointwise to  $\mathbb{Z} \circ (\mathbb{Z})$ , which is also measurable, and we have shown

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that  $\|\mathbb{Z}^{\infty}0(\mathbb{Z})\| < \infty$ . Therefore, using [4, pp. 47],  $\|\mathbb{Z}^{\mathbb{Z}}0-\mathbb{Z}^{\infty}0\| \to 0$ . We replace  $\mathbb{Z}^{\mathbb{Z}}0(\mathbb{Z})$  by  $\mathbb{Z}^{\infty}0(\mathbb{Z})$  in (154) and use symmetry, we then find that

$$\|22 \infty 0 - 220\| \le 2\|2 \infty 0 - 22-10\| \to 0.$$
 (158)

Therefore,  $2*=2\infty0$  is the solution to the Bellman equation 2\*=2\*, and the 2-convergence rate is immediately given. The solution is unique: If there exists any other measurable function

 $\|\mathbb{P}(\mathbb{P})\| < \infty$ ), we replace  $\mathbb{P}^{\mathbb{P}}(\mathbb{P})$  by  $\mathbb{P}(\mathbb{P})$  and replace

that satisfies the Bellman equation (with

 $\mathbb{Z}^{2}$ \_10 by  $\mathbb{Z}$  ( $\mathbb{Z}$ ) in (154), and we have

$$\|2 \infty 0 - 2^{\sim}\| = \|22 \infty 0 - 22^{\sim}\| \le 2\|2 \infty 0 - 2^{\sim}\|, \tag{159}$$

 $_{\infty}$  ~. These completes the proof of Lemma 9(c). which implies that  $\mathbb{Z}$  0 =  $\mathbb{Z}$ 

### H PROOF OF THEOREM 4

We denote  $\Pi \mathbb{Z}_{signal-agnostic} \subset \Pi_{signal-agnostic}$  as a collection of sampling times  $\mathbb{Z}_{2,1}, \mathbb{Z}_{2,2}, \cdots$  at  $\mathbb{Z}$ th epoch such that the inter-sampling times  $\mathbb{Z}_{2,1} - \mathbb{Z}_{2-1,\mathbb{Z}_{2,2}} - \mathbb{Z}_{2-1,\mathbb{Z}_{2,2}} - \mathbb{Z}_{2-1,\mathbb{Z}_{2,2}}, \ldots$  are independent of the history information before  $\mathbb{Z}_{2-1,\mathbb{Z}_{2-1}}$ . Note that the subscripts  $(\mathbb{Z}, 1)$ ,  $(\mathbb{Z}, 2)$ ,  $\ldots$  are illustrated in Section 5.1. Similar to Proposition 1, we have the following result:

Proposition 2. There exists an optimal policy in  $\Pi_{\mathbb{Z}, signal-agnostic}$  such that  $\{ \mathbb{Z}_{\mathbb{Z}^m} - \mathbb{Z}_{\mathbb{Z}^{-1}, \mathbb{Z}^{m-1}} \}_{\mathbb{Z}}$  are i.i.d. Moreover, problem (3) when  $\Pi = \Pi_{signal-agnostic}$  is equivalent to the following problem:

$$E \ h \int \mathbb{P}_{\mathbb{R},\mathbb{P}_{\mathbb{R}}} \qquad i$$

$$\Delta \mathbb{P} \mathbb{P}$$

$$\mathbb{P}_{\mathbb{P}^{-1},\mathbb{P}_{\mathbb{P}^{-1}}}$$

$$mse^{opt} = \mathbb{P}_{\text{CNB}, signol-agnostic} \text{inf} \quad E \ \mathbb{P} \mathbb{P}_{\mathbb{P}^{-1}} = \mathbb{P}^{-1},\mathbb{P}_{\mathbb{P}^{-1}} \ . \tag{160}$$

The proof of Proposition 2 is a special case of (thus included in) the proof of Proposision (1) and is omitted. Problem (160) has a much simpler form to Problem (23) because (i) the sampling

times are independent of the Wiener process, and (ii) we replace the square estimation error ( $\mathbb{Z}_2$   $-\mathbb{Z}_2$ ) by the linear age  $\Delta_{\mathbb{Z}}$ , the time period between  $\mathbb{Z}$  and the sampling time  $\mathbb{Z}_{\mathbb{Z}-1,\mathbb{Z}^{2}-1}$ . By (160) and [24, Section V.B], we complete the proof of Theorem 4.

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 $_{10}$   $_{\infty}$  (2) to be measurable. Here, we do not require 22  $_{\odot}$  0