

Three early problems on size Ramsey numbers

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Abstract

The size Ramsey number of a graph H is defined as the minimum number of edges in a graph G such that there is a monochromatic copy of H in every two-coloring of $E(G)$. The size Ramsey number was introduced by Erdős, Faudree, Rousseau, and Schelp in 1978 and they ended their foundational paper by asking whether one can determine up to a constant factor the size Ramsey numbers of three families of graphs: complete bipartite graphs, book graphs (obtained by adding many common neighbors to the vertices of a clique), and starburst graphs (obtained by adding many pendant edges to each vertex of a clique). In this paper, we completely resolve the latter two questions and make substantial progress on the first by determining the size Ramsey number of $K_{s,t}$ up to a constant factor for all $t = \Omega(s \log s)$.

1 Introduction

Given two graphs G and H , we say that G is *Ramsey for H* if every two-coloring of the edges of G contains a monochromatic copy of H . Graph Ramsey theory is mainly concerned with determining which graphs G are Ramsey for a given H . In particular, of central concern is the *Ramsey number* $r(H)$ of H , defined as the minimum number of vertices in a graph G which is Ramsey for H .

In this paper, we study the *size Ramsey number* $\hat{r}(H)$, defined as the minimum number of edges in a graph G which is Ramsey for H . The size Ramsey number was introduced by Erdős, Faudree, Rousseau, and Schelp [16] in 1978. They proved several bounds on size Ramsey numbers, noting, for example, the basic inequality

$$\hat{r}(H) \leq \binom{r(H)}{2}$$

and presenting a proof, due to Chvátal, that this bound is tight when H is a complete graph. They ended their paper with four questions, asking for the asymptotic order of $\hat{r}(H)$ as H ranges over four specific families of graphs. In this paper, we fully resolve two of these questions and make substantial progress on a third. The fourth question, about the size Ramsey number of paths, was resolved by Beck [2], who proved the surprising result that $\hat{r}(P_n) = \Theta(n)$ for the path P_n with n vertices. This breakthrough inspired many of the subsequent developments in the field, such as the classic papers [3, 18, 21, 24, 29] and the more recent results in [4, 5, 6, 12, 13, 14, 15, 19, 20, 23, 26].

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The first question asked by Erdős, Faudree, Rousseau, and Schelp [16] was about $\hat{r}(K_{s,t})$ for $s \leq t$. They proved the bounds

$$\Omega(st2^s) \leq \hat{r}(K_{s,t}) \leq O(s^2t2^s),$$

with the lower bound only holding for $t = \Omega(s^2)$. However, in a later paper [17], Erdős and Rousseau proved the lower bound $\hat{r}(K_{s,t}) = \Omega(st2^s)$ for all $s \leq t$.¹ More recently, Pikhurko [28] found an asymptotic formula for $\hat{r}(K_{s,t})$ for all fixed s and $t \rightarrow \infty$, which, in particular, implies that $\hat{r}(K_{s,t}) = \Theta(s^2t2^s)$ for t sufficiently large in terms of s . However, Pikhurko's technique provides no quantitative estimate on how large t must be for such a bound to hold.

Our first main result is an improved lower bound on $\hat{r}(K_{s,t})$ which holds for all $s \leq t$.

Theorem 1.1. *For all $s \leq t$,*

$$\hat{r}(K_{s,t}) = \Omega\left(s^{2-\frac{s}{t}}t2^s\right).$$

In particular, if $t \geq (1 + \delta)s$ for any fixed $\delta > 0$, then we get a power saving over the earlier lower bound of $\Omega(st2^s)$. Moreover, once $t = \Omega(s \log s)$, the bound is tight up to a constant factor.

Corollary 1.2. *If $t = \Omega(s \log s)$, then*

$$\hat{r}(K_{s,t}) = \Theta(s^2t2^s).$$

The second question raised by Erdős, Faudree, Rousseau, and Schelp [16] concerned book graphs. Given positive integers k and n , the *book graph* $B_n^{(k)}$ consists of n copies of K_{k+1} , glued along a common K_k . Equivalently, it can be described as the join of a K_k and an independent set of order n . This K_k is called the *spine* and the vertices of the independent set are called *pages*. Ramsey numbers of book graphs play a central role in Ramsey theory, because all known techniques for proving upper bounds on the diagonal Ramsey numbers $r(K_t)$ rely on induction schemes that repeatedly use bounds on $r(B_n^{(k)})$ for appropriately chosen $k < t$ and n . These Ramsey numbers have received considerable attention of late, beginning with work of the first author [7], who asymptotically determined $r(B_n^{(k)})$ for k fixed and $n \rightarrow \infty$, and continuing with work of the authors [9, 10] giving alternative proofs and exploring variations of the basic question. Regarding size Ramsey numbers, Erdős, Faudree, Rousseau, and Schelp [16] proved that

$$\Omega(k^2n^2) \leq \hat{r}(B_n^{(k)}) \leq O(16^kn^2)$$

for n sufficiently large in terms of k . Thus, while they were able to prove that the dependence on n is quadratic, there was a massive gap between the lower and upper bounds for the dependence on k . Our second main result closes this gap, determining $\hat{r}(B_n^{(k)})$ up to a constant factor for n sufficiently large in terms of k .

Theorem 1.3. *For every fixed $k \geq 2$ and all sufficiently large n ,*

$$\hat{r}(B_n^{(k)}) = \Theta(k2^kn^2).$$

¹They only state their result for $s = t$, but the proof carries through for all $s \leq t$. We present their proof, in this greater generality, in Section 2.

The third question raised in [16] concerns graphs which we call starburst graphs (they appear to have not been previously named in the literature). For positive integers k and n , the *starburst graph* $S_n^{(k)}$ is obtained from K_k by adding n pendant edges to every vertex of K_k ; thus, it has $kn + k$ vertices. Erdős, Faudree, Rousseau, and Schelp [16] proved² that if k is fixed and n is sufficiently large, then

$$\Omega(k^3 n^2) \leq \hat{r}(S_n^{(k)}) \leq O(k^4 n^2).$$

Thus, in this case, there is only a $\Theta(k)$ gap between the upper and lower bounds. Our final main result shows that the lower bound is tight up to the constant factor.

Theorem 1.4. *For every fixed $k \geq 2$ and all sufficiently large n ,*

$$\hat{r}(S_n^{(k)}) = \Theta(k^3 n^2).$$

The proofs of our main theorems are all relatively short, but they employ a surprising array of different techniques. In Theorem 1.1, the main new idea is to use a random coloring with a hypergeometric distribution between certain vertices and their higher degree neighbors, rather than the uniform distribution that is usually used in Ramsey-theoretic lower bound constructions. The lower bound in Theorem 1.3 uses a degree-based random coloring, where the probability an edge is red depends on the degrees of its endpoints, while the upper bound uses some of the regularity techniques that were recently developed for studying the ordinary Ramsey numbers of books [7, 9, 10]. Finally, Theorem 1.4 is proved by examining the properties of an appropriate random graph. We discuss each of these techniques, both at a high level and in detail, in the relevant sections.

The rest of the paper is organized as follows. We prove Theorem 1.1 in Section 2, Theorem 1.3 in Section 3, and Theorem 1.4 in Section 4, concluding with some further remarks and open problems. We use $\log$ throughout to denote the base 2 logarithm and $\ln$ for the natural logarithm. For the sake of clarity of presentation, we systematically omit floor and ceiling signs whenever they are not crucial. For the same reason, we make no serious attempt to optimize any of the constants appearing in our results.

2 Complete bipartite graphs

In order to obtain some intuition for our proof of Theorem 1.1, it is helpful to briefly review the proofs of the existing bounds on $\hat{r}(K_{s,t})$, namely,

$$\Omega(st2^s) \leq \hat{r}(K_{s,t}) \leq O(s^2 t 2^s)$$

for $s \leq t$. We first consider the upper bound, due to Erdős, Faudree, Rousseau, and Schelp [16].

Proposition 2.1 ([16]). *For all $s \leq t$, $\hat{r}(K_{s,t}) \leq 4es^2 t 2^s$.*

Proof. Let G be a complete bipartite graph with one part A of order $2s^2$ and the other part B of order $2et2^s$. We claim that G is Ramsey for $K_{s,t}$, which implies the desired result since G has $4es^2 t 2^s$ edges. Consider a red/blue coloring of $E(G)$. Call a vertex in B *red* if at least half its

²They only included the proof for the weaker lower bound $n^2/2$, but claimed the bound shown. For completeness, we include a proof of the lower bound in Section 4.

incident edges are red and *blue* otherwise. Without loss of generality, we may assume that there are at least $et2^s$ red vertices in B . Each such vertex contributes at least $\binom{s^2}{s}$ red stars $K_{1,s}$ whose central vertex is in B . Since there are exactly $\binom{2s^2}{s}$ s -tuples in A , one of them must appear as the set of leaves of such a star at least

$$et2^s \frac{\binom{s^2}{s}}{\binom{2s^2}{s}} \geq et2^s \cdot 2^{-s} \left(1 - \frac{1}{s^2}\right) \cdots \left(1 - \frac{s-1}{s^2}\right) \geq et \exp\left(-2 \sum_{i=1}^{s-1} \frac{i}{s^2}\right) \geq t$$

times, yielding a red $K_{s,t}$, where we used the inequality $1 - x \geq e^{-2x}$, valid for $x \in [0, \frac{1}{2}]$. $\square$

By being a little more careful in the proof of Proposition 2.1, one can improve the bound by a factor of $8 + o(1)$ to $\hat{r}(K_{s,t}) \leq (e/2 + o(1))s^2t2^s$, where the $o(1)$ term tends to 0 as $s \rightarrow \infty$. This can be done by picking A of order $s^2/2$, B of order $(e + o(1))t2^s$, using both red and blue stars out of vertices of B and applying convexity, though we also need to be a little more careful with the inequalities. This stronger upper bound is also present in [16] and it follows from [28, Theorem 4.6] that it is asymptotically tight, i.e., that

$$\hat{r}(K_{s,t}) = \left(\frac{e}{2} + o(1)\right)s^2t2^s$$

as long as t is sufficiently large in terms of s , where the $o(1)$ term tends to 0 as $s \rightarrow \infty$.

The lower bound argument, essentially due to Erdős and Rousseau [17], is somewhat more involved. Since $\hat{r}(K_{s,t}) \geq \hat{r}(K_{s-2,t})$, it suffices to prove the following statement.

Proposition 2.2 ([17]). *For all $t \geq s + 2$, $\hat{r}(K_{s,t}) \geq st2^s/100$.*

Proof. Since the desired result is obvious for $s = 1$, we will assume throughout that $s \geq 2$. Let G be a graph with q edges. We first give an upper bound for the number of copies of $K_{s,t}$ in G . Suppose G has N vertices $v_1, \dots, v_N$ and let $d_i = \deg(v_i)$. Assume, without loss of generality, that $d_1 \geq d_2 \geq \dots \geq d_N$. We recall the simple fact that $d_i \leq 2q/i$ for all i , which follows by double-counting the number of edges incident to $\{v_1, \dots, v_i\}$: it is at least $id_i/2$, but at most q , from which the bound follows. Let N_i denote the number of copies of $K_{s,t}$ in G where v_i is the last (i.e., maximum index) vertex in the side of order s . Then

$$N_i \leq \binom{i-1}{s-1} \binom{d_i}{t} \leq \left(\frac{ei}{s}\right)^s \left(\frac{2eq}{it}\right)^t \leq \frac{(2e^2q)^t}{s^s t^t i^{t-s}}. \quad (1)$$

Therefore, the total number of $K_{s,t}$ in G is at most

$$\sum_{i=s}^N N_i \leq \frac{(2e^2q)^t}{s^s t^t} \sum_{i=s}^{\infty} \frac{1}{i^{t-s}} \leq \frac{(2e^2q)^t}{s^s t^t} \int_{s-1}^{\infty} \frac{1}{x^{t-s}} dx = \left(\frac{2e^2q}{st}\right)^t \frac{s-1}{t-s-1} \left(\frac{s}{s-1}\right)^{t-s} \leq \left(\frac{8e^2q}{st}\right)^t,$$

where we may start the sum at $i = s$ as i is the maximum index on an s -set of vertices and we use our assumptions that $t \geq s + 2$ and $s \geq 2$ to evaluate the integral. In the final inequality, we use the bounds $\frac{s-1}{t-s-1} \leq s \leq 2^s \leq 2^t$ and $\frac{s}{s-1} \leq 2$, which hold since $t \geq s + 2$ and $s \geq 2$.

Now, suppose we color the edges of G uniformly at random. Each copy of $K_{s,t}$ is monochromatic with probability 2^{1-st} , so the expected number of monochromatic $K_{s,t}$ is at most $2(8e^2q/(st2^s))^t$. If we plug in $q = st2^s/100$, we find that the expected number of monochromatic $K_{s,t}$ is less than 1, which yields the desired result. $\square$

Note that in the upper bound argument, the worst case happens when all vertices in B are incident with the same number of red and blue edges. However, in the lower bound argument, we colored all the edges uniformly at random, meaning that, for any given $b \in B$, the number of edges incident to b that are colored red is binomially distributed. Usually, the difference between a binomial distribution and a hypergeometric distribution (where we condition on having equally many red and blue edges) is fairly minor, but it can have a significant effect when coloring, for instance, a complete bipartite graph between vertex sets A and B where A is small. This suggests using a slightly different procedure for the lower bound, where we randomly color the edges from low-degree to high-degree vertices hypergeometrically rather than uniformly. This intuition is indeed at the heart of our proof of Theorem 1.1, although the actual coloring we use is somewhat more complicated: roughly speaking, we first dyadically partition according to degree and then use independent hypergeometric variables for each dyadic interval.

In the course of the proof, we will need the following simple analytic lemma.³

Lemma 2.3. *Let m and x be positive integers and let $x_1, \dots, x_m$ be non-negative integers summing to x . If $x \geq 2m$, then*

$$\sum_{k=1}^m 2^{-k}(x_k^2 - x_k) \geq \frac{x^2}{2^{m+3}}.$$

Proof. Let $P \subseteq [m]$ denote the set of $k \in [m]$ with $x_k \geq 1$. Then we have that

$$\begin{aligned} 2^{m+1} \sum_{k=1}^m 2^{-k}(x_k^2 - x_k) &= 2^{m+1} \sum_{k \in P} 2^{-k}(x_k^2 - x_k) \\ &\geq 2^{m+1} \sum_{k \in P} 2^{-k}(x_k - 1)^2 && [y^2 - y \geq (y - 1)^2 \text{ for } y \geq 1] \\ &\geq \left(\sum_{k \in P} 2^{-k}(x_k - 1)^2 \right) \left(\sum_{k \in P} 2^k \right) && [\sum_{k \in P} 2^k \leq \sum_{k=1}^m 2^k \leq 2^{m+1}] \\ &\geq \left(\sum_{k \in P} (x_k - 1) \right)^2 && [\text{Cauchy-Schwarz}] \\ &\geq (x - m)^2 \geq \frac{x^2}{4}. \end{aligned}$$

Dividing by 2^{m+1} gives the desired result. $\square$

We are now ready to proceed with the proof of Theorem 1.1. Once again, since $\hat{r}(K_{s,t}) \geq \hat{r}(K_{s-2,t})$, it suffices to prove the following statement.

Theorem 2.4. *For all $t \geq s + 2$, $\hat{r}(K_{s,t}) \geq s^{2-\frac{s}{t}} t 2^s / 10000$.*

Proof. The desired result already follows from Proposition 2.2 for $s < 100$, so we henceforth assume that $s \geq 100$.

Let $L = \log \frac{s^2 t}{50(t-s)}$. We note, for future convenience, that $2^L \leq s^3 \leq t^3$, since $\frac{t}{t-s} \leq s$ for $t \geq s + 2$. Additionally, since $s \geq 100$, we have that $s \geq 6 \log s \geq 2L$.

³We are grateful to Mehtaab Sawhney for a suggestion that greatly simplified the proof of this lemma.

Let $q = s^{2-\frac{s}{i}} t 2^s / 10000$ and let G be a q -edge graph with vertices $v_1, \dots, v_N$ and degree sequence $d_1 \geq \dots \geq d_N$, where $d_i = \deg(v_i)$. By adding isolated vertices to G if necessary, we may assume that $L < \log(N+2)$.

For every $1 \leq \ell \leq L$, we define the interval of vertices $I_\ell = \{v_i : 2^\ell - 1 \leq i \leq 2^{\ell+1} - 2\}$. Let $A = \bigcup_{\ell=1}^L I_\ell$ and $B = V(G) \setminus A$. We color all edges inside B uniformly at random. All remaining edges have at least one endpoint in A and we color these hypergeometrically, as follows. For every vertex $v_j \in V(G)$ and every $\ell \leq \min\{L, \lceil \log(j+2) - 1 \rceil\}$, we pick a set $R_{j,\ell} \subset I_\ell$ uniformly at random among all sets with exactly $|I_\ell|/2$ elements, making these choices independently for all j and ℓ . Then, for every $v_i \in I_\ell$ with $i < j$ that is adjacent to v_j , we color the edge $v_i v_j$ red if $i \in R_{j,\ell}$ and blue otherwise.⁴

One key observation we use several times is that, for every set F of edges in G , the probability they are all red is at most $2^{-|F|}$. To see this, we first note that it is true if F consists of edges whose left endpoint lies in some fixed interval I_ℓ and whose right endpoint is some fixed vertex v_j . Indeed, in this case, the result follows immediately from the negative correlation property of the hypergeometric distribution: the probability that such an F is monochromatic red equals the probability that the set of left endpoints is a subset of $R_{j,\ell}$, which is at most $2^{-|F|}$. We can partition F into subsets of this form, plus some additional edges lying in B . The events that each of these subsets is monochromatic red are all independent, so the probability that F is monochromatic red equals the product of the probabilities that each is monochromatic red, which is at most $2^{-|F|}$.

Now suppose that $S \subseteq A$ is a set of vertices with $|S| = s$. Let $s_\ell = |S \cap I_\ell|$, so that $s = s_1 + \dots + s_L$. Let v_i be the rightmost vertex in S (i.e., i is the maximal index appearing in S) and let $L^* = \lfloor \log(i+1) \rfloor$. By the choice of L^* , as $v_i \in S \subseteq A$, we have $v_i \in I_{L^*}$. We then have that $s_\ell = 0$ if $\ell > L^*$, since there can be no element in $S \cap I_\ell$ for $\ell > L^*$ by the maximality of i .

For $1 \leq \ell \leq L^*$ and for a vertex v_j with $j > i$, the probability that v_j is monochromatic red to $S \cap I_\ell$ is exactly

$$\begin{aligned} \frac{|I_\ell|/2}{|I_\ell|} \frac{|I_\ell|/2 - 1}{|I_\ell| - 1} \dots \frac{|I_\ell|/2 - (s_\ell - 1)}{|I_\ell| - (s_\ell - 1)} &\leq 2^{-s_\ell} \left(1 - \frac{1}{|I_\ell|}\right) \left(1 - \frac{2}{|I_\ell|}\right) \dots \left(1 - \frac{s_\ell - 1}{|I_\ell|}\right) \\ &\leq 2^{-s_\ell} \exp\left(-2^{-\ell} \binom{s_\ell}{2}\right) \\ &= 2^{-s_\ell} e^{-2^{-\ell-1}(s_\ell^2 - s_\ell)}. \end{aligned}$$

Therefore, the probability that v_j is monochromatic red to S is at most

$$2^{-s} \exp\left(-\frac{1}{2} \sum_{\ell=1}^{L^*} 2^{-\ell} (s_\ell^2 - s_\ell)\right).$$

We now apply Lemma 2.3 with $x = s$ and $m = L^*$, which we may do since $s \geq 2L \geq 2L^*$, to conclude that the probability v_j is monochromatic red to S is at most

$$2^{-s} \exp\left(-\frac{1}{2} \frac{s^2}{2^{L^*+3}}\right) \leq 2^{-s} \exp\left(-\frac{s^2}{32i}\right).$$

⁴Note that, as described, the neighborhood of v_j in I_ℓ isn't exactly hypergeometrically distributed, since we first pick a set of "possible" red neighbors hypergeometrically, but then only color red those "possible" neighbors which are truly neighbors in G . Though it doesn't really matter, this choice makes the analysis slightly simpler.

There is also a slight subtlety when $v_j \in A$, in that, if $v_j \in I_\ell$ for some interval I_ℓ , we only wish to describe the color of the edges from v_j to the previous vertices $v_i \in I_\ell$; this is why we add the condition $i < j$.

As in the proof of Proposition 2.2, we let N_i denote the number of copies of $K_{s,t}$ whose last vertex from the side of order s is v_i . Similarly, for $v_i \in A$ and $0 \leq x \leq t$, let $N_{i,x}$ be the number of copies of $K_{s,t}$ whose last vertex from the side of order s is i and where exactly x vertices from the side of order t have index greater than i . Let $y = t - x$. Then

$$N_{i,x} \leq \binom{i-1}{s-1} \binom{i-s}{y} \binom{d_i}{x} \leq \left(\frac{ei}{s}\right)^s \left(\frac{ei}{y}\right)^y \left(\frac{2eq}{ix}\right)^x \leq \frac{(2e^2)^t q^x}{s^s y^y x^x i^{x-y-s}},$$

where we take $0^0 = 1$ so that the bound also holds when $x = 0$ or $x = t$. For each of the x vertices to the right of v_i , they have probability at most $2^{-s} e^{-s^2/32i}$ of being monochromatic red to all s chosen vertices. Additionally, by the negative correlation property mentioned earlier, the remaining y vertices each have probability at most 2^{-s} of being monochromatic red to all s chosen vertices.

Therefore, the probability that some $K_{s,t}$ counted by $N_{i,x}$ is monochromatic red is at most

$$\frac{(2e^2)^t q^x \cdot 2^{-st} e^{-s^2 x/32i}}{s^s y^y x^x i^{x-y-s}} = \frac{(2e^2)^t i^s}{s^s 2^{st}} \exp\left(x \ln q - \left(y \ln y + x \ln x + (x-y) \ln i + \frac{s^2 x}{32i}\right)\right), \quad (2)$$

where we let $0 \ln 0 = 0$, agreeing with our earlier convention that $0^0 = 1$. We claim that this probability is maximized when $x = t$. To see this, let

$$f(x) = x \ln q - \left((t-x) \ln(t-x) + x \ln x + (2x-t) \ln i + \frac{s^2 x}{32i}\right).$$

For $x \leq t-1$, we can compute that

$$f'(x) = \ln q - \left(-\ln(t-x) + \ln x + 2 \ln i + \frac{s^2}{32i}\right) \geq \ln q - \left(\ln \frac{xi^2}{t-x} + \frac{s}{32}\right) \geq \ln q - \ln(4ts^6 e^{s/32}),$$

since we have $s \leq i \leq 2^{L+1} \leq 2s^3$ and $t-x \geq 1$. Note that for $s \geq 100$, we have that $2^s e^{-s/32} \geq (10s)^5$. Therefore, by our choice of q , we see that

$$\frac{q}{4ts^6 e^{s/32}} = \frac{s^{2-\frac{s}{i}} t 2^s / 10000}{4ts^6 e^{s/32}} > \frac{2^s e^{-s/32}}{(10s)^5} \geq 1$$

and thus $f'(x) > 0$ for all $x \in [0, t-1]$.

This shows that f is maximized on this interval at $x = t-1$. Additionally, we have that

$$f(t) - f(t-1) = \ln q - \left(t \ln t - (t-1) \ln(t-1) + 2 \ln i + \frac{s^2}{32i}\right) \geq \ln q - \ln(4ets^6 e^{s/32}) \geq 0,$$

again by our choice of q , where we use the fact that

$$t \ln t - (t-1) \ln(t-1) = \int_{t-1}^t (1 + \ln x) dx \leq 1 + \ln t = \ln(et).$$

This implies that $f(x) \leq f(t)$ for all $0 \leq x \leq t$, which shows that, as claimed, the probability in (2) is maximized for $x = t$. Therefore, for $v_i \in A$, the probability that some $K_{s,t}$ counted by N_i is monochromatic red is at most

$$\sum_{x=0}^t \frac{(2e^2)^t q^x \cdot 2^{-st} e^{-s^2 x/32i}}{s^s y^y x^x i^{x-y-s}} \leq (t+1) \frac{(2e^2 q)^t 2^{-st} e^{-s^2 t/32i}}{s^s t^t i^{t-s}} \leq \frac{(4e^2 q)^t}{2^{st} s^s t^t} \exp\left(-\left(\frac{s^2 t}{32i} + (t-s) \ln i\right)\right).$$

Let $g(i) = \frac{s^2 t}{32i} + (t-s) \ln i$ and note that

$$g'(i) = \frac{t-s}{i} - \frac{s^2 t}{32i^2}.$$

This shows that $g(i)$ is minimized when $i = \frac{s^2 t}{32(t-s)}$ and, therefore, the minimum value of g is $(t-s) \ln \frac{es^2 t}{32(t-s)}$. Plugging this in, the probability that some $K_{s,t}$ counted by N_i is monochromatic red is at most

$$\frac{(4e^2 q)^t}{2^{st} s^s t^t} \left(\frac{32(t-s)}{es^2 t} \right)^{t-s} \leq \left(\frac{4e^2 q}{st2^s} \right)^t \left(\frac{50(t-s)}{st} \right)^{t-s}.$$

Summing this up, the probability that some $K_{s,t}$ whose side of order s lies in A is monochromatic red is at most

$$\sum_{i=s}^{2^{L+1}} \left(\frac{4e^2 q}{st2^s} \right)^t \left(\frac{50(t-s)}{st} \right)^{t-s} \leq 2t^3 \left(\frac{4e^2 q}{st2^s} \right)^t \left(\frac{50(t-s)}{st} \right)^{t-s} \leq \left(\frac{100q}{st2^s} \right)^t \left(\frac{50(t-s)}{st} \right)^{t-s},$$

since $2^{L+1} \leq 2t^3 \leq 3^t$ for $t \geq 6$. For all the remaining $K_{s,t}$, we use the negative correlation property to see that each one is monochromatic red with probability at most 2^{-st} . Adding this up over all $i > 2^L$ and using our bound $N_i \leq \frac{(2e^2 q)^t}{s^s t^t i^{t-s}}$ from (1), we find that the probability any of these copies is monochromatic red is at most

$$\sum_{i=2^L+1}^N \frac{(2e^2 q)^t}{s^s t^t i^{t-s}} 2^{-st} \leq \frac{(2e^2 q)^t}{s^s t^t 2^{st}} \int_{2^L}^{\infty} \frac{1}{x^{t-s}} dx = \left(\frac{2e^2 q}{st2^s} \right)^t \frac{2^L}{t-s-1} \left(\frac{s}{2^L} \right)^{t-s} \leq \left(\frac{100q}{st2^s} \right)^t \left(\frac{s}{2^L} \right)^{t-s},$$

using our assumption $t \geq s+2$ to evaluate the integral and again using the bound $2^L \leq t^3 \leq 3^t$. In total, the probability that any $K_{s,t}$ is monochromatic red is at most

$$\left(\frac{100q}{st2^s} \right)^t \left(\left(\frac{50(t-s)}{st} \right)^{t-s} + \left(\frac{s}{2^L} \right)^{t-s} \right).$$

We now recall our choice of $2^L = \frac{s^2 t}{50(t-s)}$, so that both terms in parentheses are equal. Therefore, in total, we find that the probability of a monochromatic red $K_{s,t}$ is at most

$$\left(\frac{100q}{st2^s} \right)^t 2 \left(\frac{50(t-s)}{st} \right)^{t-s} < \left(\frac{100q}{st2^s} \right)^t \left(\frac{100}{s} \right)^{t-s} = 100^{-s} < \frac{1}{2},$$

by plugging in $q = s^{2-\frac{s}{t}} t^s / 10000$. By symmetry, the same estimate holds for the probability of a blue $K_{s,t}$, which yields the desired conclusion. $\square$

3 Book graphs

We prove the lower and upper bounds in Theorem 1.3 separately. We begin with the lower bound, which is somewhat simpler.

3.1 The lower bound

In this section, we use the Chernoff bound in the following form.

Lemma 3.1. *Let $X_1, \dots, X_t$ be independent random variables taking values in $\{0, 1\}$ and let $X = \sum X_i$. If $\mathbb{E}[X] \leq n/2$, then*

$$\Pr(X \geq n) \leq e^{-n/6}.$$

Proof. It follows from Corollary 2.4 and Theorem 2.8 in [22] that if $x \geq c\mathbb{E}[X]$ for some $c > 1$, then $\Pr(X \geq x) \leq e^{-c'x}$, where $c' = \ln c - 1 + \frac{1}{c}$. Plugging in $c = 2$ (so that $x = n$), we find that $c' = \ln 2 - \frac{1}{2} > \frac{1}{6}$, which yields the claimed bound. $\square$

We also need the following simple analytic result, estimating a certain integral that arises in the proof.

Lemma 3.2. *For every integer $k \geq 1$,*

$$\int_0^1 \ln \left(1 - \frac{y^{1/k}}{2} \right) dy \leq \frac{1}{k} - \ln 2.$$

Proof. Denote the integral by I . We begin by changing variables to $u = y^{1/k}$, so that $du = \frac{1}{k} y^{(1-k)/k} dy = \frac{1}{k} u^{1-k} du$ or, equivalently, $dy = ku^{k-1} du$. We get that

$$I = \int_0^1 ku^{k-1} \ln \left(1 - \frac{u}{2} \right) du.$$

We now integrate by parts, letting $dg = ku^{k-1} du$ and $f = \ln(1 - \frac{u}{2})$. Then $g = u^k$ and $df = du/(u - 2)$, so that

$$I = \left[u^k \ln \left(1 - \frac{u}{2} \right) \right]_0^1 - \int_0^1 \frac{u^k}{u - 2} du = -\ln 2 + \int_0^1 \frac{u^k}{2 - u} du.$$

This integral can be explicitly evaluated in terms of the beta function, but for our purposes it suffices to bound it: we note that $1/(2 - u) \leq 1$ since $u \leq 1$ and, therefore,

$$\int_0^1 \frac{u^k}{2 - u} du \leq \int_0^1 u^k du = \frac{1}{k+1} \leq \frac{1}{k}.$$

Thus, $I \leq \frac{1}{k} - \ln 2$, as claimed. $\square$

We now prove our lower bound on $\hat{r}(B_n^{(k)})$. We only state the result for $k \geq 25$, which implies the lower bound in Theorem 1.3 by picking the implicit constant appropriately to deal with $2 \leq k \leq 24$.

Theorem 3.3. *For every $k \geq 25$ and every $n \geq 300k^2$, $\hat{r}(B_n^{(k)}) \geq k2^k n^2/1200$.*

Proof. Let G be a graph with at most $q = k2^k n^2/1200$ edges, with vertices $v_1, \dots, v_N$. We order the vertices so that $\deg(v_1) \geq \dots \geq \deg(v_N)$. We may assume that G has no isolated vertices, so that $N \leq 2q$.

We fix $s = k/3$. For $0 \leq i \leq s$, let $V_i = \{v_{in/10+1}, \dots, v_{(i+1)n/10}\}$. Thus, V_0 consists of the $n/10$ vertices of G of highest degree, V_1 consists of the next $n/10$ vertices of highest degree, and so on. For every $1 \leq i \leq s$, every vertex in V_i has degree at most

$$\deg(v_{in/10}) \leq \frac{2q}{in/10} = \frac{20q}{in} = \frac{k2^k n}{60i} = \frac{s2^k n}{20i} =: D_i.$$

Let $U = V(G) \setminus (V_0 \cup \dots \cup V_s)$. On $V_0 \cup \dots \cup V_s$, we use a Turán coloring: we color all edges inside each V_i red and all edges between V_i and V_j blue for all $0 \leq i < j \leq s$. All remaining edges are colored randomly, as follows. First, every edge inside U is colored red or blue with probability $\frac{1}{2}$. Second, every edge between V_i and U is colored red with probability $p_i = \frac{1}{2}(i/s)^{1/k}$ and blue with probability $1 - p_i$. All these random choices are made independently. Note that since $p_0 = 0$, the edges between V_0 and U are actually colored blue deterministically.

If Q is a copy of K_k in G , let $\text{ext}(Q)$ denote the number of extensions of Q to a monochromatic K_{k+1} ; equivalently, $\text{ext}(Q)$ is the number of pages in the largest monochromatic book whose spine is Q . Note that $\text{ext}(Q) = 0$ if Q is not itself monochromatic. We next argue that $\mathbb{E}[\text{ext}(Q)] \leq n/2$ for every choice of Q . We then use the Chernoff bound and the union bound to conclude that with positive probability there is no monochromatic $B_n^{(k)}$ (and, thus, that there is a coloring with no monochromatic $B_n^{(k)}$).

First, we deal with red books. Let Q be a copy of K_k in G , which we think of as a potential spine of a red book. Suppose first that $Q \cap V_0 \neq \emptyset$. Since $B_{n/2}^{(k)}$ is connected and all edges between V_0 and $V(G) \setminus V_0$ are blue, any red copy of $B_{n/2}^{(k)}$ with a vertex in V_0 must actually be entirely contained within V_0 . But $|V_0| = n/10 < n/2$, so any red $B_{n/2}^{(k)}$ cannot intersect V_0 .

Next, suppose that $Q \cap V_i \neq \emptyset$ for some $i \in [s]$. Since all edges between V_i and V_j are blue for $i \neq j$, any red book whose spine is Q must lie in $V_i \cup U$. Let $r = |V_i \cap Q|$ be the number of vertices in the spine which lie in V_i . For a vertex $v \in V_i \setminus Q$ which is adjacent in G to every vertex of Q , the probability that v forms a red extension of Q equals p_i^{k-r} . Similarly, for $u \in U \setminus Q$ adjacent to all vertices of Q , the probability that u forms a red extension of Q is $p_i^r \cdot (\frac{1}{2})^{k-r} = 2^{-k}(2p_i)^r$. If $r < k$, then the number of choices for such a vertex $u \in U \setminus Q$ is at most D_s , since any such vertex must be a common neighbor of the non-empty set $U \cap Q$ and so, in particular, must lie in the neighborhood of some fixed vertex in U . Therefore, by the definitions of p_i and D_i ,

$$\mathbb{E}[\text{ext}(Q)] \leq p_i^{k-r}|V_i| + 2^{-k}(2p_i)^r D_s \leq \frac{n}{10} + \left(\frac{i}{s}\right)^{r/k} \cdot \frac{n}{20} \leq \left(\frac{1}{10} + \frac{1}{20}\right)n \leq \frac{n}{2}.$$

If, instead, $r = k$, then there are most D_i choices for a vertex which is joined to every vertex of Q , so that

$$\mathbb{E}[\text{ext}(Q)] \leq |V_i| + p_i^k D_i \leq \frac{n}{10} + \frac{i}{s} \cdot \frac{sn}{20i} = \left(\frac{1}{10} + \frac{1}{20}\right)n \leq \frac{n}{2}.$$

Finally, we consider the case where $Q \subseteq U$. Since every vertex of U has degree at most D_s , there are at most D_s common neighbors of Q in U , each of which forms a red extension of Q with probability 2^{-k} . Moreover, for all $i \in [s]$, each common neighbor of Q in V_i forms a red extension of Q with probability p_i^k . Therefore,

$$\mathbb{E}[\text{ext}(Q)] \leq 2^{-k} D_s + \sum_{i=1}^s p_i^k |V_i| = \frac{n}{20} + \frac{n}{10} \sum_{i=1}^s 2^{-k} \frac{i}{s} \leq \left(\frac{1}{20} + \frac{s}{10 \cdot 2^k}\right)n \leq \frac{n}{2},$$

since $s \leq k \leq 2^k$. This shows that for every choice of Q the expected number of red extensions of Q is at most $n/2$.

Next, we deal with blue books, so let Q be a potential spine of a blue book. Since each V_i is colored monochromatically in red, Q contains at most one vertex from each V_i . Let $R = \{0 \leq i \leq s : V_i \cap Q \neq \emptyset\}$ be the set of i such that Q contains a vertex from V_i and let $r = |R|$. Then all the pages of a blue book whose spine is Q must lie in $U \cup \bigcup_{j \notin R} V_j$. Since $r \leq s + 1 < k$, at least one vertex of Q must be in U , so the number of common neighbors of Q in U is at most D_s . Each of these common neighbors forms a blue extension of Q with probability $2^{r-k} \prod_{i \in R} (1 - p_i)$. Additionally, for $j \notin R$, the probability that a common neighbor of Q in V_j yields a blue extension of Q equals $(1 - p_j)^{k-r}$. Therefore,

$$\mathbb{E}[\text{ext}(Q)] \leq D_s 2^{r-k} \prod_{i \in R} (1 - p_i) + \sum_{j \notin R} (1 - p_j)^{k-r} |V_j| = \frac{n}{20} \prod_{i \in R} (2 - 2p_i) + \frac{n}{10} \sum_{j \notin R} (1 - p_j)^{k-r}. \quad (3)$$

By definition, we see that $0 \leq p_i \leq \frac{1}{2}$ for all i . Moreover, for $j \geq 1$, we have that

$$p_j \geq \frac{1}{2} \left(\frac{1}{s} \right)^{1/k} \geq \frac{1}{2} \left(\frac{1}{2^k} \right)^{1/k} = \frac{1}{4},$$

since $s \leq k \leq 2^k$. Additionally, $r \leq s + 1 \leq k/2$, so $k - r \geq k/2$. Therefore,

$$\sum_{j \notin R} (1 - p_j)^{k-r} \leq \sum_{j \notin R} (1 - p_j)^{k/2} \leq 1 + \sum_{j=1}^s (1 - p_j)^{k/2} \leq 1 + \sum_{j=1}^s \left(\frac{3}{4} \right)^{k/2} \leq 1 + k \cdot \left(\frac{3}{4} \right)^{k/2} \leq 2,$$

since $k(3/4)^{k/2} \leq 1$ for $k \geq 25$. This shows that the second term in (3) is at most $n/5$. For the first term, we recall that since $p_i \leq \frac{1}{2}$, every term in the product is at least 1. Therefore, this product is maximized when $R = \{0, 1, \dots, s\}$. In that case, we have

$$\prod_{i=0}^s (2 - 2p_i) = 2^{s+1} \exp \left(\sum_{i=0}^s \ln(1 - p_i) \right) = 2^{s+1} \exp \left(\sum_{i=0}^s \ln \left(1 - \frac{1}{2} \left(\frac{i}{s} \right)^{1/k} \right) \right).$$

To estimate this sum, we use the fact that the function $f(x) = \ln(1 - \frac{1}{2}(\frac{x}{s})^{1/k})$ is decreasing to write

$$\sum_{i=0}^s f(i) = \sum_{i=1}^s f(i) = \sum_{i=1}^s \int_{i-1}^i f(i) dx \leq \sum_{i=1}^s \int_{i-1}^i f(x) dx = \int_0^s f(x) dx.$$

For this integral, we change variables to $y = x/s$ and $dy = dx/s$ to find that

$$\int_0^s f(x) dx = \int_0^s \ln \left(1 - \frac{1}{2} \left(\frac{x}{s} \right)^{1/k} \right) dx = s \int_0^1 \ln \left(1 - \frac{y^{1/k}}{2} \right) dy \leq s \left(\frac{1}{k} - \ln 2 \right),$$

where the inequality is from Lemma 3.2. Putting all this together, we see that

$$\prod_{i \in R} (2 - 2p_i) \leq \prod_{i=0}^s (2 - 2p_i) \leq 2^{s+1} \exp \left(s \left(\frac{1}{k} - \ln 2 \right) \right) = 2e^{1/3}.$$

Thus, the first term in (3) is at most $\frac{e^{1/3}}{10}n < \frac{n}{5}$. Plugging all this into (3), we conclude that for every choice of a potential blue spine Q ,

$$\mathbb{E}[\text{ext}(Q)] \leq \frac{n}{5} + \frac{n}{5} \leq \frac{n}{2}.$$

Thus, for every k -tuple Q of vertices, the expected number of pages in a monochromatic book whose spine is Q is at most $n/2$. The random variable $\text{ext}(Q)$ can be expressed as the sum of independent $\{0, 1\}$ -valued random variables. Therefore, Lemma 3.1 implies that, for any fixed Q , we have $\Pr(\text{ext}(Q) \geq n) \leq e^{-n/6}$. Recalling that the number of vertices in G is $N \leq 2q$, we see that the total number of choices for Q is at most

$$\binom{N}{k} \leq N^k \leq (k2^k n^2)^k \leq \exp(k^2 + 2k \ln n),$$

since $k2^k < e^k$ for $k \geq 25$. If $n \geq 300k^2$, then $\frac{n}{6} - k^2 \geq \frac{n}{8}$ and $\ln n \leq \sqrt{n}$, so

$$k^2 + 2k \ln n - \frac{n}{6} \leq 2k\sqrt{n} - \frac{n}{8} < 0,$$

since $\sqrt{n} > 16k$. Therefore, taking a union bound over all choices of Q , we see that, with positive probability, none of them will have at least n monochromatic extensions, showing that there is a coloring of G with no monochromatic $B_n^{(k)}$. $\square$

By being slightly more careful in the proof of Theorem 3.3, we can show that the bound already holds when $n \geq Ck \log k$ for an appropriate constant C . To accomplish this, we count the expected number of k -cliques Q in G that extend to a monochromatic $B_n^{(k)}$ as

$$\sum_Q \Pr(Q \text{ is monochromatic}) \Pr(\text{ext}(Q) \geq n \mid Q \text{ is monochromatic}).$$

We then split this sum according to how many vertices of Q are in U . Indeed, if we let t be the number of vertices in $Q \cap U$, then Lovász's version of the Kruskal–Katona theorem [27, Exercise 13.31] tells us that, provided $t \geq 2$, the number of ways of choosing the t vertices of Q in U is at most $(2q)^{t/2}/t!$. By using this bound and carefully estimating the probability that Q is monochromatic, one can obtain the claimed result.

However, such effort may be for naught, because this bound may not be optimal in this regime. Indeed, $\hat{r}(B_n^{(k)}) \geq \hat{r}(K_k) = \binom{r(K_k)}{2}$. If the Ramsey number $r(k)$ is at least $(\sqrt{2} + \varepsilon)^k$ for some fixed $\varepsilon > 0$, as some have conjectured, then the lower bound in Theorem 3.3 is not sharp for $n = 2^{o(k)}$.

This lower bound argument points towards which host graph to use for the upper bound. Namely, it should have roughly kn vertices of very high degree, each around $2^k n$, while all other vertices should have substantially lower degree. This suggests taking the host graph to be a large book, with a spine of order around kn and about $2^k n$ pages, which is indeed what we do. However, the proof that such a large book is indeed Ramsey for $B_n^{(k)}$ is quite involved, requiring substantial input from recent work on Ramsey numbers of books [7, 9, 10]. This proof is the content of the next subsection.

3.2 The upper bound

For the upper bound in Theorem 1.3, we need to recall some facts related to Szemerédi's regularity lemma. Given vertex subsets X, Y in a graph G , we denote by $e(X, Y)$ the number of pairs in $X \times Y$ that are edges of G and define the *edge density* by $d(X, Y) = e(X, Y)/(|X||Y|)$. We write $d(X)$ if $X = Y$ and $d(x, Y)$ if $X = \{x\}$ is a singleton. A pair (X, Y) is called ε -regular if, for all $X' \subseteq X, Y' \subseteq Y$ with $|X'| \geq \varepsilon|X|, |Y'| \geq \varepsilon|Y|$, we have $|d(X', Y') - d(X, Y)| \leq \varepsilon$. Similarly, we say that a set X is ε -regular if the pair (X, X) is ε -regular. A partition of the vertex set of a graph is called *equitable* if all the parts have orders that differ by at most 1. We use the following strengthened version of Szemerédi's regularity lemma, proved in [9, Lemma 2.1].

Lemma 3.4. *For every $\varepsilon > 0$ and $M_0 \in \mathbb{N}$, there is some $M = M(\varepsilon, M_0) \geq M_0$ such that, for every graph G , there is an equitable partition $V(G) = V_1 \sqcup \dots \sqcup V_m$ into $M_0 \leq m \leq M$ parts such that the following hold:*

1. *Each part V_i is ε -regular and*
2. *For every $1 \leq i \leq m$, there are at most εm values $1 \leq j \leq m$ such that the pair (V_i, V_j) is not ε -regular.*

To complement the regularity lemma, we have the following consequence of the counting lemma, proved in [9, Corollary 2.6]. It is designed to count monochromatic extensions of cliques and thus estimate the size of monochromatic books. Given a copy Q of K_k in some graph G , we say that $u \in V(G)$ *extends* Q if u is a common neighbor of all k vertices of Q .

Lemma 3.5. *Fix $k \geq 2$ and let $\varepsilon, \delta \in (0, \frac{1}{2})$ be parameters with $\varepsilon \leq \delta^{3k^2}$. Let X be a set of vertices in a graph G and suppose that X is ε -regular with edge density at least δ . Then X contains at least one K_k . Moreover, if Q is a randomly chosen copy of K_k in X , then, for any $u \in V(G)$,*

$$\Pr(u \text{ extends } Q) \geq d(u, X)^k - \delta.$$

The final tool we need concerns good configurations, which we now define.

Definition 3.6. Let $k \geq 2$ and $\varepsilon, \delta \in (0, 1)$ be parameters and let $C_1, \dots, C_k$ be disjoint vertex subsets in a graph G whose edges have been two-colored. We say that $C_1, \dots, C_k$ form a (k, ε, δ) -*good configuration* if the following conditions hold:

- $C_1 \cup \dots \cup C_k$ induces a complete subgraph of G , meaning that there is an edge of G between any two vertices of $C_1 \cup \dots \cup C_k$.
- Each C_i is ε -regular in red and has internal red density at least δ .
- Each pair (C_i, C_j) for $i \neq j$ is ε -regular in blue and has blue density at least δ .

The following lemma, which shows that colorings which contain good configurations also contain large monochromatic books, was arguably the main idea in [7], but was first stated explicitly in [9]. The statement differs slightly from that in [9, Lemma 3.3], but it is easy to check that the proof carries through.

Lemma 3.7. *Fix $k \geq 2$, $0 < \delta \leq 2^{-k-1}$, and $0 < \varepsilon \leq \delta^{3k^2}$. Suppose that the edges of a graph G have been two-colored and let $C_1, \dots, C_k$ be a (k, ε, δ) -good configuration in G . If the vertices in $C_1 \cup \dots \cup C_k$ have t common neighbors in G , then the coloring contains a monochromatic $B_{2^{-k-1}t}^{(k)}$.*

With these preliminaries, we are ready to prove the upper bound in Theorem 1.3.

Theorem 3.8. *For every $k \geq 2$ and every sufficiently large n , $\hat{r}(B_n^{(k)}) \leq k2^{k+3}n^2$.*

Proof. Let $G = B_N^{(K)}$, where $K = 2kn$ and $N = 2^{k+1}n$, so that G has $\binom{K}{2} + NK \leq 2k^2n^2 + k2^{k+2}n^2 \leq k2^{k+3}n^2$ edges. We claim that G is Ramsey for $B_n^{(k)}$ if n is sufficiently large with respect to k . To prove this, fix a red/blue coloring of $E(G)$. Let $\delta = \min\{2^{-k-1}, 1/(10k^3)\}$ and $\varepsilon = \delta^{3k^2}$.

We apply Lemma 3.4 to the red graph on the subgraph of G induced by the spine, with ε as above and $M_0 = 10k^2$. We obtain an equitable partition $V_1 \sqcup \dots \sqcup V_m$ of the spine with $M_0 \leq m \leq M = M(\varepsilon, M_0)$ parts, where each V_i is ε -regular and, for every i , at most εm of the pairs (V_i, V_j) are not ε -regular. Note that, since the colors are complementary, if a pair is ε -regular in red, then it is also ε -regular in blue.⁵ Call a part V_i *red* if at least half its internal edges are red and *blue* otherwise. Without loss of generality, suppose that $m' \geq m/2$ of the parts are red and reindex so that these red parts are $V_1, \dots, V_{m'}$. We form a reduced graph F with vertices $v_1, \dots, v_{m'}$, connecting v_i to v_j in F if (V_i, V_j) is ε -regular and has blue density at least δ .

First suppose that some vertex in F , say v_1 , has degree at most $(1 - \frac{1}{k-1})m'$. This means that there are at least $\frac{m'}{k-1} - 1$ parts V_j with $2 \leq j \leq m'$ such that (V_1, V_j) is either irregular or has blue density less than δ . By the second condition in Lemma 3.4, the number of irregular pairs (V_1, V_j) is at most $\varepsilon m \leq 2\varepsilon m'$. Therefore, if we let J be the set of those $2 \leq j \leq m'$ such that (V_1, V_j) is ε -regular and has blue density less than δ , we find that

$$|J| \geq \frac{m'}{k-1} - 1 - 2\varepsilon m' = \left(\frac{1}{k-1} - \frac{1}{m'} - 2\varepsilon \right) m' \geq \left(\frac{1}{k-1} - \frac{2}{5k^2} \right) m' \geq \left(\frac{1}{k} + 3k\delta \right) m',$$

using the facts that $m' \geq m/2 \geq M_0/2 = 5k^2$ and $\varepsilon \leq \delta \leq 1/(10k^3)$, as well as $\frac{1}{k-1} - \frac{1}{k^2} \geq \frac{1}{k}$.

Let $U = \bigcup_{j \in J} V_j$. Since the partition is equitable,

$$|U| \geq \frac{|J|}{m} K - m \geq \left(\frac{1}{k} + 3k\delta \right) \frac{K}{2} - M = \left(\frac{1}{k} + 3k\delta - 2\frac{M}{K} \right) \frac{K}{2} \geq \left(\frac{1}{k} + 2k\delta \right) \frac{K}{2},$$

where we used that $m \leq M$ and $2M/K \leq k\delta$ for n (and thus K) sufficiently large.

Recall that V_1 is ε -regular and has red density at least $\frac{1}{2} \geq \delta$. Thus, by Lemma 3.5, if Q is a random red K_k in V_1 , then, for any $u \in U$, the probability that u extends Q is at least $d_R(u, V_1)^k - \delta$, where d_R denotes edge density in the red graph. Adding this up over all $u \in U$, we find that the expected number of red extensions of Q is at least

$$\sum_{u \in U} \left(d_R(u, V_1)^k - \delta \right) \geq \left(\frac{1}{|U|} \sum_{u \in U} d_R(u, V_1) \right)^k |U| - \delta |U| \geq \left((1 - \delta)^k - \delta \right) |U| \geq (1 - 2k\delta) |U|,$$

where the first inequality follows from convexity of the function $x \mapsto x^k$ and the second uses the fact that $d_R(V_1, V_j) \geq 1 - \delta$ for every $j \in J$. Plugging in our lower bound for $|U|$, we find that the expected number of red extensions of Q is at least

$$(1 - 2k\delta) \left(\frac{1}{k} + 2k\delta \right) \frac{K}{2} \geq \left(\frac{1}{k} + 2k\delta - 2k\delta \right) \frac{2kn}{2} = n.$$

⁵This is true if the parts in the pair are disjoint. If not, one can choose a smaller ε to guarantee the same end.

Thus, there exists a red K_k in V_1 with at least n red extensions, that is, a red $B_n^{(k)}$.

Thus, we may assume that every vertex in F has degree greater than $(1 - \frac{1}{k-1})m'$. By Turán's theorem, this implies that F contains a K_k . Let the vertices of this K_k be $v_{i_1}, \dots, v_{i_k}$ and let $C_1 = V_{i_1}, \dots, C_k = V_{i_k}$. Then $C_1, \dots, C_k$ form a (k, ε, δ) -good configuration, since each C_i is ε -regular with red density at least $\frac{1}{2} \geq \delta$ and all pairs (C_i, C_j) with $i \neq j$ correspond to edges of F and so are ε -regular with blue density at least δ . Moreover, the vertices of $C_1 \cup \dots \cup C_k$ all lie in the spine of G and, hence, they have $N = 2^{k+1}n$ common neighbors. Thus, by Lemma 3.7, our coloring must contain a monochromatic $B_n^{(k)}$. $\square$

4 Starburst graphs

4.1 The lower bound

Recall that the starburst graph $S_n^{(k)}$ is obtained from K_k by adding n pendant edges to every vertex of K_k . As mentioned in the introduction, although Erdős, Faudree, Rousseau, and Schelp [16] asserted the lower bound $\hat{r}(S_n^{(k)}) = \Omega(k^3 n^2)$ for sufficiently large n , they did not include a proof, so we begin with this lower bound. In fact, we record the following more general lower bound, where $\hat{r}(H_1, H_2)$ denotes the minimum number of edges in a graph G such that every red/blue coloring of $E(G)$ contains a red copy of H_1 or a blue copy of H_2 .

Proposition 4.1. *Let H_1 be a connected graph with $n+1$ vertices and maximum degree Δ and let H_2 be a graph with chromatic number $k+1$. Then $\hat{r}(H_1, H_2) \geq \binom{k}{2}\Delta n$.*

Since the graph $S_n^{(k)}$ has $kn+k$ vertices, maximum degree $n+k-1$, and chromatic number k , we immediately get the following corollary, which clearly implies the lower bound in Theorem 1.4 for $k \geq 3$ (see [16, Fact A] for the $k=2$ case).

Corollary 4.2. *For every $k \geq 2$ and $n \geq 1$,*

$$\hat{r}(S_n^{(k)}) \geq \binom{k-1}{2}(n+k-1)(kn+k-1).$$

In particular, $\hat{r}(S_n^{(k)}) \geq (\frac{1}{2} - o(1))k^3 n^2$, where the $o(1)$ term tends to 0 as $k \rightarrow \infty$.

Proof of Proposition 4.1. Let G be an N -vertex graph with fewer than $\binom{k}{2}\Delta n$ edges. By iteratively deleting vertices of maximum degree, we may order the vertices of G as $v_1, \dots, v_N$, where v_i is a maximum-degree vertex in the induced subgraph $G_i = G[v_i, \dots, v_N]$ for every i . For $1 \leq i \leq N$, let $D(i)$ be the degree of v_i in G_i and declare $D(i) = 0$ if $i > N$. By our choice of ordering, we have that $D(1) \geq D(2) \geq \dots$.

Suppose for the moment that $D(jn) \geq (k-j)\Delta$ for all $1 \leq j \leq k-1$. Then the number of edges in G is

$$\sum_{i=1}^N D(i) \geq \sum_{j=1}^{k-1} \sum_{i=(j-1)n+1}^{jn} D(i) \geq \sum_{j=1}^{k-1} n D(jn) \geq \Delta n \sum_{j=1}^{k-1} (k-j) = \binom{k}{2} \Delta n,$$

where the second inequality uses that $D(i) \geq D(jn)$ for all $i \leq jn$. This contradicts our assumption that the number of edges in G is less than $\binom{k}{2}\Delta n$. So we may set $j^* \leq k-1$ to be the smallest

positive integer such that $D(j^*n) < (k - j^*)\Delta$. We now create a partition of $V(G)$ into $V_1 \sqcup \dots \sqcup V_k$, as follows.

For $1 \leq j \leq j^*$, let $V_i = \{v_{(j-1)n+1}, \dots, v_{jn}\}$ and let $G' = G \setminus (V_1 \cup \dots \cup V_{j^*})$. We consider a $\max(k - j^*)$ -cut of G' , that is, a partition of $V(G')$ into $k - j^*$ parts which maximizes the number of edges between the parts. Let this partition of G' be $V_{j^*+1}, \dots, V_k$. Note that $G' = G_{j^*n+1}$, which implies that every vertex in G' has degree at most $D(j^*n) < (k - j^*)\Delta$. Moreover, in the $\max(k - j^*)$ -cut, every vertex has at least as many neighbors in every other part as in its own part, for otherwise we could increase the number of edges going between parts by moving some vertex to a different part. Thus, for every $j^* + 1 \leq j \leq k$, every vertex in V_j has fewer than $(k - j^*)\Delta / (k - j^*) = \Delta$ neighbors in its own part V_j .

In summary, each part V_j for $1 \leq j \leq j^*$ has exactly n vertices and each part V_j for $j^* + 1 \leq j \leq k$ induces a subgraph of G with maximum degree less than Δ . We color the edges inside each part red and all edges between parts blue. Then the blue graph is k -partite and so cannot contain a copy of H_2 . On the other hand, the red graph is the disjoint union of $G[V_1], \dots, G[V_{j^*}]$, so, since H_1 is connected, any red H_1 must appear in $G[V_j]$ for some j . But it cannot appear in $G[V_j]$ for $j \leq j^*$ since H_1 has more than n vertices and it cannot appear in $G[V_j]$ for $j > j^*$ since H_1 has maximum degree Δ and $G[V_j]$ has maximum degree less than Δ . Therefore, G is not Ramsey for H . $\square$

4.2 The upper bound

To prove the upper bound on $\hat{r}(S_n^{(k)})$, we need to recollect some facts about random graphs. We use $G(N, p)$ to denote the Erdős–Rényi random graph with N vertices where each edge is chosen independently with probability p , saying that an event happens *with high probability (w.h.p.)* if the probability it holds in $G(N, p)$ tends to 1 as N tends to ∞ .

The first result we need is a version of Turán’s theorem relative to random graphs. This is a well-studied topic, culminating in works by Conlon–Gowers [11] and Schacht [31], who determined the threshold for Turán’s theorem to hold in this setting (see also [1, 30] for subsequent proofs using the method of hypergraph containers). However, we only need a rather weak version, so the following earlier result of Szabó and Vu [32], with explicit bounds on the error probability, more than suffices for our purposes. For a graph G , we use $\text{ex}(G, K_k)$ to denote the maximum number of edges in a K_k -free subgraph of G .

Theorem 4.3 ([32, Theorem 1.2]). *For every $k \geq 4$ and $\varepsilon > 0$, there exist constants $C, c > 0$ such that the following holds for all sufficiently large N . If $p \geq CN^{-2/(2k-3)}$ and $G \sim G(N, p)$, then*

$$\Pr \left(\text{ex}(G, K_k) \leq \left(1 - \frac{1}{k-1} + \varepsilon \right) p \binom{N}{2} \right) \geq 1 - 2^{-cpN^2}.$$

The following lemma records the facts about $G(N, p)$ used in our proof of the upper bound on $\hat{r}(S_n^{(k)})$. We use the notation $x = y \pm z$ to mean that x lies in the interval $[y - z, y + z]$. Recall too that $d(X, Y)$ denotes the edge density between vertex sets X, Y .

Lemma 4.4. *Fix $k \in \mathbb{N}$ and $p, \delta \in (0, 1)$. The following facts all hold w.h.p. in $G \sim G(N, p)$:*

- (a) *Every vertex $v \in V(G)$ has degree $(p \pm \delta)N$.*
- (b) *For any two (not necessarily distinct or disjoint) sets $X, Y \subseteq V(G)$ with $|X|, |Y| \geq p^{8k}N$, $d(X, Y) = p \pm \delta$.*

- (c) For every $S \subseteq V(G)$ with $|S| \geq p^{8k}N$, the number of $v \in V(G)$ with fewer than $(p - \delta)|S|$ neighbors in S is at most $p^{8k}N$. Similarly, the number of $v \in V(G)$ with more than $(p + \delta)|S|$ neighbors in S is at most $p^{8k}N$.
- (d) For every $S \subseteq V(G)$ with $|S| \geq N/3$, every K_k -free subgraph of $G[S]$ has at most $(1 - \frac{1}{2k})p \binom{|S|}{2}$ edges.

Proof. It suffices to prove that each property holds w.h.p., since a simple application of the union bound then implies that they all hold w.h.p. simultaneously.

- (a) The expected degree of each vertex is $p(N - 1)$, so, by the Chernoff bound (e.g., [22, equation (2.12)]), for any fixed $v \in V(G)$,

$$\Pr(\deg(v) \neq (p \pm \delta)N) \leq \Pr\left(|\deg(v) - \mathbb{E}[\deg(v)]| \geq \frac{\delta}{2}N\right) \leq 2 \exp\left(-2 \left(\frac{\delta}{2}\right)^2 N\right),$$

where we use that N is sufficiently large in terms of p and δ to write $p(N - 1) = pN \pm \frac{\delta}{2}N$. We now take a union bound over all N choices of v .

- (b) This also follows easily from an application of the Chernoff bound and the union bound. For details, see, e.g., [25, Corollary 2.3].
- (c) This follows immediately from part (b). Indeed, suppose we let X consist of all vertices with fewer than $(p - \delta)|S|$ neighbors in S . Then $d(X, S) < p - \delta$, which contradicts part (b) if $|X| \geq p^{8k}N$. The second statement follows identically.
- (d) For any fixed $S \subseteq V(G)$ with $|S| \geq N/3$, we can apply Theorem 4.3 with $\varepsilon = \frac{1}{2k}$ to conclude that the probability there is a K_k -free subgraph of $G[S]$ with more than $(1 - \frac{1}{2k})p \binom{|S|}{2}$ edges is at most $2^{-cp(N/3)^2} = 2^{-\Omega_{p,k}(N^2)}$. Taking a union bound over the fewer than 2^N choices for S , we see that the required conclusion holds provided N is sufficiently large with respect to p and k . $\square$

The following result completes the proof of Theorem 1.4. Note that we can assume $k \geq 4$, since the required result for $k \in \{2, 3\}$ follows from the $k = 4$ case.

Theorem 4.5. *For every $k \geq 4$ and all sufficiently large n , $\hat{r}(S_n^{(k)}) \leq 10^5 k^3 n^2$.*

Proof. Fix $p = 1/(10k)$, $\delta = p^2$, and $N = 1000k^2n$. Then, for n sufficiently large, $G \sim G(N, p)$ satisfies all the properties in Lemma 4.4 w.h.p., so we may fix an N -vertex graph G with these properties. In particular, by Lemma 4.4(a), G has at most $pN^2 = 10^5 k^3 n^2$ edges, so it suffices to show that G is Ramsey for $S_n^{(k)}$.

Suppose, then, that the edges of G have been red/blue colored. Let V_R denote the set of vertices of G whose blue degree is less than $3n$, let V_B denote the set of vertices of G whose red degree is less than $3n$, and let V_0 consist of all vertices with red and blue degree at least $3n$. At least one of V_R, V_B , and V_0 has at least $N/3$ vertices.

Suppose first that $|V_R| \geq N/3$. By Lemma 4.4(b), the number of edges in V_R is at least $(p - \delta) \frac{|V_R|^2}{2}$. Moreover, each vertex in V_R has blue degree at most $3n$, which means that the number of blue edges in V_R is at most

$$\frac{3}{2}n|V_R| = \frac{3}{2} \cdot \frac{N}{1000k^2}|V_R| \leq \frac{9}{2000k^2}|V_R|^2 \leq \delta \frac{|V_R|^2}{2}.$$

Therefore, the red graph on V_R has at least $(p-2\delta)\frac{|V_R|^2}{2} \geq (1-2p)p\binom{|V_R|}{2}$ edges. Since $2p = \frac{1}{5k} < \frac{1}{2k}$, Lemma 4.4(d) implies that V_R contains a red K_k . Let the vertices of this red K_k be $v_1, \dots, v_k$. Each v_i has at least $(p-\delta)N$ incident edges, of which at most $3n$ are colored blue. Let X_i be the set of red neighbors of v_i , apart from v_j for $j \neq i$, so that

$$|X_i| \geq (p-\delta)N - 3n - k \geq \frac{p}{2}N = 50kn.$$

Now, we let Y_1 be an arbitrary subset of X_1 with $|Y_1| = n$. Inductively, we let Y_i be an arbitrary subset of $X_i \setminus (Y_1 \cup \dots \cup Y_{i-1})$ of order n , which we can do since $|X_i \setminus (Y_1 \cup \dots \cup Y_{i-1})| \geq 50kn - (i-1)n \geq n$. By doing this for every i , we find disjoint sets $Y_1, \dots, Y_k$ of red neighbors of $v_1, \dots, v_k$, respectively, which yields a red $S_n^{(k)}$. By interchanging the roles of the colors, the same argument yields a blue $S_n^{(k)}$ if $|V_B| \geq N/3$.

So we may now suppose that $|V_0| \geq N/3$. Our goal is to apply the Erdős–Szekeres neighborhood-chasing argument to build a monochromatic K_k inside V_0 , while also maintaining large sets of red and blue neighbors for each vertex of this K_k , so that we can complete it to a monochromatic copy of $S_n^{(k)}$. Formally, we inductively construct three sequences: a sequence $v_1, \dots, v_{2k}$ of vertices in V_0 , a sequence $C_1, \dots, C_{2k} \in \{R, B\}$ of colors, and a sequence $Z_1, \dots, Z_{2k}$ of disjoint vertex sets, with the property that each Z_i consists of $\frac{3}{2}n$ red and $\frac{3}{2}n$ blue neighbors of v_i . Additionally, we consider the sequence $N_1 \supseteq N_2 \supseteq \dots$, where $N_i = V_0 \cap N_{C_1}(v_1) \cap N_{C_2}(v_2) \cap \dots \cap N_{C_i}(v_i)$ consists of the common neighborhoods in V_0 of the i vertices $v_1, \dots, v_i$ in the colors $C_1, \dots, C_i$. We inductively maintain the properties that $|N_i| \geq p^{2i}|V_0|$ and $v_{i+1} \in N_i$ for all $1 \leq i < 2k$.

To begin, by Lemma 4.4(c), at most $p^{8k}N < |V_0|$ vertices in V_0 have fewer than $(p-\delta)|V_0|$ neighbors in V_0 , so let $v_1 \in V_0$ be an arbitrary vertex with at least $(p-\delta)|V_0|$ neighbors in V_0 . Among the edges between v_1 and V_0 , let C_1 be the majority color. Finally, since $v_1 \in V_0$, it has at least $3n$ red neighbors and at least $3n$ blue neighbors, so we let Z_1 be an arbitrary set of $\frac{3}{2}n$ red and $\frac{3}{2}n$ blue neighbors of v_1 . As indicated above, let $N_1 = V_0 \cap N_{C_1}(v_1)$ be the set of neighbors of v_1 in V_0 in the color C_1 . By construction,

$$|N_1| \geq \frac{p-\delta}{2}|V_0| \geq \frac{p}{4}|V_0| \geq p^2|V_0|,$$

since v_1 has at least $(p-\delta)|V_0|$ neighbors in V_0 , of which at least half have color C_1 . This sets up the base case of our induction.

Inductively, suppose that for some $i < 2k$ we have defined $v_1, \dots, v_i, C_1, \dots, C_i$, and $Z_1, \dots, Z_i$ and we wish to construct v_{i+1}, C_{i+1} , and Z_{i+1} . As above, let $N_i = V_0 \cap N_{C_1}(v_1) \cap \dots \cap N_{C_i}(v_i)$, where the induction hypothesis and $i < 2k$ implies that $|N_i| \geq p^{2i}|V_0| \geq p^{4k}N$. By Lemma 4.4(c), the number of vertices with fewer than $(p-\delta)|N_i|$ neighbors in N_i is at most $p^{8k}N$. Similarly, if we let $Z = Z_1 \cup \dots \cup Z_i$, then $|Z| = 3in \geq p^{8k}N$, so the number of vertices with more than $(p+\delta)|Z|$ neighbors in Z is at most $p^{8k}N$. Since $2p^{8k}N < p^{4k}N \leq |N_i|$, we may pick some $v_{i+1} \in N_i$ with at least $(p-\delta)|N_i|$ neighbors in N_i and at most $(p+\delta)|Z|$ neighbors in Z . Let C_{i+1} be the majority color among all edges between v_{i+1} and N_i , so that

$$|N_{i+1}| = |N_i \cap N_{C_{i+1}}(v_{i+1})| \geq \frac{p-\delta}{2}|N_i| \geq p^2|N_i| \geq p^{2(i+1)}|V_0|,$$

maintaining the inductive hypothesis. Finally, to define Z_{i+1} , we recall that the number of neighbors of v_{i+1} in Z is at most

$$(p+\delta)|Z| \leq 2p|Z| = \frac{6in}{10k} < \frac{12kn}{10k} < \frac{3}{2}n.$$

Since $v_{i+1} \in V_0$, it has at least $3n$ blue neighbors and at least $3n$ red neighbors. Hence, avoiding the neighbors of v_{i+1} that are in Z , we can still find a set Z_{i+1} of $\frac{3}{2}n$ red and $\frac{3}{2}n$ blue neighbors of v_{i+1} which is disjoint from $Z_1, \dots, Z_i$.

To conclude, we assume, without loss of generality, that at least half of $C_1, \dots, C_{2k}$ are red. That is, there are distinct $j_1, \dots, j_k \in [2k]$ such that $C_{j_1}, \dots, C_{j_k}$ are all red. Then $v_{j_1}, \dots, v_{j_k}$ form a red K_k , since, for any $j > i$, $v_j \in N_{j-1} \subseteq N_{C_i}(v_i)$. Moreover, the sets $Z_{j_1}, \dots, Z_{j_k}$ are disjoint and each contains $\frac{3}{2}n$ red neighbors of $v_{j_1}, \dots, v_{j_k}$, respectively. By deleting at most $k < \frac{n}{2}$ vertices from each Z_{j_i} , we can ensure that each Z_{j_i} is also disjoint from $\{v_{j_1}, \dots, v_{j_k}\}$, while still containing at least n red neighbors of v_{j_i} for each i . This yields a red copy of $S_n^{(k)}$, as desired. $\square$

Observe that Theorem 1.4, which states that $\hat{r}(S_n^{(k)}) = \Theta(k^3 n^2)$ for n sufficiently large, actually requires that n is at least exponential in k . This is because

$$\hat{r}(S_n^{(k)}) \geq \hat{r}(K_k) = \binom{r(K_k)}{2} \geq 2^k.$$

A careful analysis of the proof of Theorem 4.5 shows that the upper bound $\hat{r}(S_n^{(k)}) = O(k^3 n^2)$ holds when $n \geq k^{Ck}$ for an appropriate constant C . This particular proof, using an Erdős–Rényi random graph, cannot give a better bound on n . However, a different random graph model can be used to show that $\hat{r}(S_n^{(k)}) = O(k^3 n^2)$ holds already when $n \geq 2^{Ck}$. Indeed, consider the random graph model $G(N, p, \omega)$ obtained by iteratively picking a uniformly random subset of $[N]$ of order ω , adding a complete subgraph on those vertices, and repeating this process until the edge density is at least p . The global structure of $G(N, p, \omega)$ is rather close to that of $G(N, p)$ for an appropriate range of parameters, but the two models differ because $G(N, p, \omega)$ has larger cliques than $G(N, p)$. In order to prove the desired result, we again take $N = \Theta(k^2 n)$ and $p = \Theta(1/k)$ and let $\omega = N^{1/8}$. One can then follow the same proof technique as in Theorem 4.5, partitioning the vertex set of an edge-colored $G(N, p, \omega)$ into V_R, V_B, V_0 as before. The main difference from the proof of Theorem 4.5 is that one must apply Turán’s theorem and the Erdős–Szekeres argument inside the intersection of one of these sets with an appropriately chosen clique of order ω .

5 Concluding remarks

Although we have made substantial progress on the three questions asked by Erdős, Faudree, Rousseau, and Schelp, several interesting open problems remain. First, for complete bipartite graphs $K_{s,t}$, though we have shown that $\hat{r}(K_{s,t}) = \Theta(s^2 t^2 s)$ for $t = \Omega(s \log s)$, we suspect that a similar bound may hold whenever $s \leq t$. That is, we have the following conjecture.

Conjecture 5.1. *For all $s \leq t$, $\hat{r}(K_{s,t}) = \Theta(s^2 t^2 s)$. In particular, $\hat{r}(K_{t,t}) = \Theta(t^3 2^t)$.*

Secondly, for book graphs $B_n^{(k)}$, where we determined $\hat{r}(B_n^{(k)})$ up to a constant factor once n is sufficiently large in terms of k , it would be interesting to understand how large is sufficient. As discussed near the end of Section 3.1, our lower bound can be made to work once $n \geq Ck \log k$, but the upper bound relies on an application of Szemerédi’s regularity lemma, resulting in tower-type bounds. In our earlier paper [9], we were able to avoid applying the regularity lemma when looking at the ordinary Ramsey number of books and it is possible that a similar technique could yield the upper bound $\hat{r}(B_n^{(k)}) = O(k 2^k n^2)$ for n bounded by a constant-height tower function of k .

However, it would be interesting to determine whether the same bound holds for n bounded by a single-exponential function of k , say.

Thirdly, our determination of the asymptotic order of $\hat{r}(S_n^{(k)})$ for starburst graphs $S_n^{(k)}$ raises the question of when the general lower bound $\hat{r}(H) = \Omega(\chi(H)^2 \Delta(H) |V(H)|)$ in Proposition 4.1, which holds for any connected graph H , is asymptotically tight.

Question 5.2. *For which families of graphs H does one have $\hat{r}(H) = \Theta(\chi(H)^2 \Delta(H) |V(H)|)$?*

The reason this question may be interesting is that the proof of Proposition 4.1 uses a version of the Turán coloring. Indeed, we find an appropriate partition of the host graph G into $\chi(H) - 1$ parts and then color all internal edges red and all edges between parts blue. The number of parts guarantees that there is no blue copy of H , while the choice of partition ensures that no copy of H is contained in any of the red parts. One can therefore view Question 5.2 as a size Ramsey variant of the so-called Ramsey goodness problem (see, e.g., [8, Section 2.5]), which asks when the Turán coloring is extremal for ordinary Ramsey numbers. While it is not completely clear that Question 5.2 is the correct size Ramsey analogue of this problem, it may be a good first step.

It is also natural to consider multicolor size Ramsey numbers. Formally, given a graph H and an integer $q \geq 2$, let $\hat{r}(H; q)$ denote the minimum number of edges in a graph G with the property that no matter how the edges of G are colored in q colors, there is always a monochromatic copy of H . Most of the results and proofs in this paper can be adapted to deal with the multicolor case. For example, one can show that if $q \geq 2$ is fixed and $t = \Omega(s \log s)$, then

$$\hat{r}(K_{s,t}; q) = \Theta(s^2 t q^s),$$

with the upper bound again following from a simple averaging argument in an unbalanced complete bipartite graph and the lower bound from a hypergeometric random coloring. Note that here and below all implicit constants may depend on the number of colors q .

For starburst graphs, one can show that if $q \geq 3$ is fixed and n is sufficiently large in terms of k , then

$$\hat{r}(S_n^{(k)}; q) = \Theta(k \cdot n^2 \cdot r(K_k; q - 1)^2),$$

where $r(K_k; q - 1)$ is the $(q - 1)$ -color (ordinary) Ramsey number of K_k . Although it may not look like it at first, this is a natural generalization of the two-color case, since the one-color Ramsey number of K_k is simply k . Here, the lower bound is proved analogously to Proposition 4.1, except that rather than the Turán coloring, one uses a blowup of a coloring on $r(K_k; q - 1) - 1$ vertices with no monochromatic K_k . For the upper bound, one again uses $G \sim G(N, p)$, but with $p = \Theta(1/k)$ and $N = \Theta(k \cdot n \cdot r(K_k; q - 1))$. We partition the vertices of G according to which of the q color classes they have degree at least $3n$ in and then pass to the largest of these 2^q sets. If every vertex in this set has at least $3n$ neighbors in all colors, we can again use the Erdős–Szekeres argument to build a monochromatic $S_n^{(k)}$. If not, then there is at least one sparse color in this large set. One applies Szemerédi’s regularity lemma inside this set, followed by Turán’s theorem and Ramsey’s theorem on the reduced graph to find a non-sparse color and k parts which are pairwise dense and regular in this color. Finally, one can greedily pick out one vertex from each part, inductively maintaining the ability to extend this to a monochromatic copy of $S_n^{(k)}$.

However, it is not clear how to extend our results on book graphs to more than two colors. A natural conjecture is the following.

Conjecture 5.3. Fix $q \geq 3$. For every $k \geq 2$ and all sufficiently large n ,

$$\hat{r}(B_n^{(k)}; q) = \Theta\left(q^k \cdot n^2 \cdot r(K_k; q - 1)\right).$$

We feel that this is a natural conjecture, because, in the two-color case, both the lower and upper bounds are fundamentally determined by a “Turán part” of order $\Theta(kn)$ and a “random part” of order $\Theta(2^k n)$ and the natural multicolor analogues of these would have orders $\Theta(r(K_k; q - 1)n)$ and $\Theta(q^k n)$, respectively. However, we are not confident of Conjecture 5.3 even in the case $q = 3$ and can prove neither the upper nor the lower bound.

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