

NUMERICAL HOMOTOPIES FROM KHOVANSKII BASES

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ABSTRACT. We present numerical homotopy continuation algorithms for solving systems of equations on a variety in the presence of a finite Khovanskii basis. These take advantage of Anderson's flat degeneration to a toric variety. When Anderson's degeneration embeds into projective space, our algorithm is a special case of a general toric two-step homotopy algorithm. When Anderson's degeneration is embedded in a weighted projective space, we explain how to lift to a projective space and construct an appropriate modification of the toric homotopy. Our algorithms are illustrated on several examples using `Macaulay2`.

We consider the problem of computing the isolated solutions to the system

$$(1) \quad f_1(z) = f_2(z) = \cdots = f_d(z) = 0,$$

where f_1, \dots, f_d are general members of a finite-dimensional vector space V of rational functions on a complex algebraic variety X of dimension d . Kaveh-Khovanskii [18, 19] and Lazarsfeld-Mustață [21] show that the number of solutions is the normalized volume of the Newton-Okounkov body associated to V . The accompanying theory extends many uses of Newton polytopes from toric varieties to general algebraic varieties. This theory lends itself to algorithms when V has a finite Khovanskii basis [20].

The evaluation of functions in V induces the rational Kodaira map $\varphi: X \dashrightarrow \mathbb{P}(V^*)$. The solutions to System (1) are the pull backs of the points of a linear section $\varphi(X) \cap L$ along φ . When V has a finite Khovanskii basis, Anderson [2] shows that (the closure of) $\varphi(X)$ has a flat degeneration to a toric variety associated to the Newton-Okounkov body of V . We describe numerical algorithms for computing a linear section based on this toric degeneration and the polyhedral homotopy algorithm [15, 28]. Solving System (1) then requires computing the pull back of the linear section.

Our numerical algorithms for computing a linear section are based on homotopy continuation [23]. This approach uses path tracking from numerical analysis to compute the solutions to a target system F given all solutions to a start system G along with a homotopy interpolating the two systems. Anderson's flat toric degeneration gives a homotopy where the start system is a linear section of a toric variety and the target system is a linear section of $\varphi(X)$.

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Flatness guarantees that the number of solutions to the start and target systems are equal. Thus the homotopy is optimal in the sense that no extraneous paths are tracked.

Our start system is a linear section of a toric variety, which may be solved using the optimal polyhedral homotopy algorithm [15, 28]. Beyond those derived from a finite Khovanskii basis, there are many instances in which a projective variety has a flat degeneration into a toric variety. In Section 1, we describe an optimal toric two-step homotopy algorithm for solving systems given a toric degeneration in an ambient projective space. We also present examples of such flat degenerations into toric varieties.

When the Khovanskii basis is a subset of V , Anderson’s degeneration may be embedded in the projective space $\mathbb{P}(V^*)$. In Section 2, we present the Khovanskii homotopy algorithm, which uses this embedding and the toric two-step homotopy to solve System (1). For a general Khovanskii basis, Anderson’s degeneration may only be embedded in a weighted projective space and System (1) is not a pull back of a general linear section. In Section 3, we describe how to adapt the toric algorithm to this general case of a Khovanskii basis.

We end each section with a concrete example to illustrate our techniques and algorithms. These examples are computed with Macaulay2 scripts [13], which are archived on GitHub:

<https://github.com/EliseAWalker/KhovanskiiHomotopy/>

We use the `NumericalAlgebraicGeometry` package [22] to call the packages `Bertini` [3] and `PHCpack` [27] for user-defined homotopies and the polyhedral homotopy, respectively. We discuss practical issues that arise from using these software packages in Section 4.

1. HOMOTOPY CONTINUATION AND TORIC DEGENERATIONS

Numerical homotopy continuation computes the solutions to a system F of polynomial equations given the solutions to a related system G . This method uses numerical path tracking along a homotopy, which is a family of systems containing both F and G . We begin by reviewing homotopy continuation and then discuss how flat families are a source of homotopies. When a flat family is a degeneration into a toric variety and the system G is a linear section of that toric variety, we describe the toric two-step homotopy algorithm whose first step is a polyhedral homotopy.

1.1. Numerical homotopy continuation. Numerical homotopy continuation is a method for solving a system $F(x) = 0$ of polynomial equations [23]. It uses a one-parameter family $H(x; t)$ (in $t \in \mathbb{C}_t$) of polynomial systems called a *homotopy*. Numerical homotopy continuation mandates that the *start system* $G(x) := H(x; 0)$ has known solutions and the solutions to the *target system* $F(x)$ are among those to $H(x; 1)$. We further require that $H(x; t)$ defines a curve C in $\mathbb{C}_x^n \times \mathbb{C}_t$ (or $\mathbb{P}_x^n \times \mathbb{C}_t$) with $t = 0$ a regular value of the projection $C \rightarrow \mathbb{C}_t$. These assumptions imply that there are enough solutions to the start system so that along a general path in \mathbb{C}_t , the solutions to the start system $H(x; 0)$ deform to the solutions to the system $H(x; 1)$. The homotopy is *optimal* if every solution to the start system deforms to a distinct solution to the target system so that no extraneous paths are tracked.

Given a homotopy H , let γ be a general arc in \mathbb{C}_t between 0 and 1. The restriction of H to γ is a family of smooth arcs. Standard numerical path tracking algorithms starting with solutions to G can compute the set of all solutions to $H(x; 1)$, which includes all solutions

to F . When $t = 1$ is also a regular value of the projection $C \rightarrow \mathbb{C}_t$ and $H(x; 1) = F(x)$, the homotopy is optimal.

Several software packages implement numerical homotopy continuation methods. These include `Bertini` [3], `NumericalAlgebraicGeometry` [22], `HomotopyContinuation.jl` [5], `HOM4PS` [7], and `PHCpack` [27]. The first three implement user-defined homotopies, and our computational examples use the user-defined homotopy method provided in `Bertini`. The last three packages implement the polyhedral homotopy method [15, 28], and we use `PHCpack` for solving systems of sparse polynomials coming from linear sections of toric varieties.

1.2. Homotopies from flat families. Suppose that $X \subset \mathbb{P}^n$ is a subvariety of dimension d . A *linear section* of X is a transverse intersection $X \cap L$ where $L \subset \mathbb{P}^n$ is a linear subspace of codimension d so that $X \cap L$ consists of $\deg X$ points.

Let $\mathcal{X} \subset \mathbb{P}^n \times \mathbb{C}$ be a variety with a surjective map $\pi: \mathcal{X} \rightarrow \mathbb{C}$. Then π realizes \mathcal{X} as a family of projective varieties over \mathbb{C} where a point $t \in \mathbb{C}$ corresponds to the fiber $\mathcal{X}_t := \pi^{-1}(t) \subset \mathbb{P}^n$. There is an open subset $U \subset \mathbb{C}$ such that \mathcal{X} is flat over U . Flatness is an algebraic property which captures the geometric notion that the fibers \mathcal{X}_t vary continuously with $t \in U$ [11, Chapter 6]. For example, the fibers of a flat family all have the same dimension and degree.

Suppose that the fibers of a flat family \mathcal{X} over $U \subset \mathbb{C}$ have dimension d and that $0, 1 \in U$. Let $L \subset \mathbb{P}^n$ be a general linear subspace of codimension d which meets both \mathcal{X}_0 and \mathcal{X}_1 transversally so that $\mathcal{X}_0 \cap L$ and $\mathcal{X}_1 \cap L$ are linear sections. Let $H(x; t)$ be finitely many polynomials defining \mathcal{X} and d linear forms defining L . We call $H(x; t)$ a *linear section homotopy*.

Proposition 1 (Linear section homotopy). *A linear section homotopy $H(x; t)$ is an optimal homotopy with start system $\mathcal{X}_0 \cap L$ and target system $\mathcal{X}_1 \cap L$.*

Proof. Let C be the union of components of $\mathcal{X} \cap L$ that contain both $\mathcal{X}_0 \cap L$ and $\mathcal{X}_1 \cap L$. Since these intersections are zero-dimensional, C is a curve. Furthermore, both $t = 0$ and $t = 1$ are regular values of the projection $C \rightarrow \mathbb{C}_t$. Thus, H is a homotopy. Flatness implies that $\mathcal{X}_0 \cap L$ and $\mathcal{X}_1 \cap L$ have the same number of points so that the homotopy H is optimal. \square

A linear section is part of a *witness set*, which is a fundamental data structure in numerical algebraic geometry [4]. Specifically, a witness set for a d -dimensional variety $X \subset \mathbb{P}^n$ is a triple $(F, L, X \cap L)$ where F is a set of homogeneous polynomials (forms) on \mathbb{P}^n defining X , L is a set of d general linear forms defining a linear subspace (which is also written L), and $X \cap L$ is the corresponding linear section.

In the linear section homotopy in Proposition 1, L is a fixed general linear space and the variety \mathcal{X}_t moves. Our algorithms sometimes require linear spaces which are not general. For this, we use a homotopy where the variety is fixed, but the linear section moves, which is described in the following basic algorithm for moving a witness set:

Algorithm 2 (Witness Set Homotopy).

Input: A witness set $(G, L, X \cap L)$ for X and a codimension d linear subspace L' such that $X \cap L'$ is finite.

Output: The points of $X \cap L'$.

Do:

- (i) Let $\mathcal{H} := (G, tL' + (1-t)L)$, a homotopy with start system $X \cap L$ and target system $X \cap L'$.
- (ii) Use path tracking starting from the points of $X \cap L$ to compute the points of $X \cap L'$.

1.3. Toric degenerations. A *toric degeneration* $\mathcal{X} \subset \mathbb{P}^n \times \mathbb{C}_t$ is a flat family over \mathbb{C}_t whose special fiber \mathcal{X}_0 is a toric variety (see [8] for details on toric varieties). Given a toric degeneration, the linear section homotopy leads to the toric two-step homotopy which we describe in Algorithm 3.

A vector $\alpha = (a_1, \dots, a_d) \in \mathbb{Z}^d$ is the exponent of a Laurent monomial $z^\alpha := z_1^{a_1} \cdots z_d^{a_d}$, which is a function on the algebraic torus $(\mathbb{C}^\times)^d$, where $\mathbb{C}^\times := \mathbb{C} \setminus \{0\}$. Suppose that \mathcal{A} is a $d \times (n+1)$ integer matrix whose columns $\{\alpha_0, \dots, \alpha_n\}$ are a finite set of $n+1$ integer vectors. The *toric Kodaira map* $\varphi_{\mathcal{A}}: (\mathbb{C}^\times)^d \rightarrow \mathbb{P}^n$ is defined by $\varphi_{\mathcal{A}}(z) := [z^{\alpha_0}, \dots, z^{\alpha_n}]$, and the *toric variety* $X_{\mathcal{A}}$ is the closure of its image. The homogeneous ideal of $X_{\mathcal{A}}$ is spanned by the following set of binomials [26, Chapter 4]:

$$\{x^u - x^v : \sum \alpha_i u_i = \sum \alpha_i v_i \text{ and } \sum u_i = \sum v_i\}.$$

We describe a variant of this Kodaira map for translated toric varieties. The torus $(\mathbb{C}^\times)^{n+1}$ acts on \mathbb{P}^n by independently scaling each coordinate. This action factors through the quotient of $(\mathbb{C}^\times)^{n+1}$ by its diagonal, $\Delta\mathbb{C}^\times$. This quotient is the dense torus \mathbb{T} of \mathbb{P}^n . For a point $p \in \mathbb{T}$, let $p.X_{\mathcal{A}}$ be the translation of the toric variety $X_{\mathcal{A}}$ by p . We note that $p \in p.X_{\mathcal{A}}$ and that $p.X_{\mathcal{A}} = p'.X_{\mathcal{A}}$ for any $p' \in p.X_{\mathcal{A}} \cap \mathbb{T}$. The ideal of $p.X_{\mathcal{A}}$ is spanned by the following set of binomials, which depend on p :

$$\{p^v x^u - p^u x^v : \sum \alpha_i u_i = \sum \alpha_i v_i \text{ and } \sum u_i = \sum v_i\}.$$

Since $p.X_{\mathcal{A}} \simeq X_{\mathcal{A}}$, we also call $p.X_{\mathcal{A}}$ a toric variety and the ideal of $p.X_{\mathcal{A}}$ a *toric ideal*. Writing $p = [p_0, \dots, p_n]$, the corresponding toric Kodaira map for $p.X_{\mathcal{A}}$ is

$$(2) \quad \varphi_{p, \mathcal{A}}(z) = [p_0 z^{\alpha_0}, \dots, p_n z^{\alpha_n}].$$

A linear section $p.X_{\mathcal{A}} \cap L$ of the toric variety $p.X_{\mathcal{A}}$ pulls back along $\varphi_{p, \mathcal{A}}$ to the following system of sparse polynomials on $(\mathbb{C}^\times)^d$ whose monomials have exponents in \mathcal{A} :

$$(3) \quad g_1(z) = g_2(z) = \cdots = g_d(z) = 0.$$

The polyhedral homotopy algorithm is an optimal homotopy for solving this system of polynomials [15, 28].

Let $\mathcal{X} \rightarrow \mathbb{C}_t$ be a toric degeneration with d -dimensional toric special fiber $p.X_{\mathcal{A}} = \mathcal{X}_0$. A general linear subspace L of codimension d gives linear sections $p.X_{\mathcal{A}} \cap L$ and $\mathcal{X}_1 \cap L$. We combine the linear section homotopy of Proposition 1 with the polyhedral homotopy to obtain the toric two-step homotopy algorithm for computing the points of the linear section $\mathcal{X}_1 \cap L$. Let $\mathcal{G}_{\mathcal{A}}$ be System (3), which is given by the pull back of L along $\varphi_{p, \mathcal{A}}$.

Algorithm 3 (Toric two-step homotopy algorithm).

Input: A toric degeneration $\mathcal{X} \subset \mathbb{P}^n \times \mathbb{C}_t$ with $\mathcal{X}_0 = p.X_{\mathcal{A}}$ a toric variety and a general linear space $L \subset \mathbb{P}^n$ of codimension equal to the dimension of \mathcal{X}_1 .

Output: All points of the linear section $\mathcal{X}_1 \cap L$.

Do:

- (i) Compute the system $G_{\mathcal{A}}$ by pulling L back along the Kodaira map $\varphi_{p,\mathcal{A}}$.
- (ii) Use the polyhedral homotopy to solve $G_{\mathcal{A}}$.
- (iii) Use $\varphi_{p,\mathcal{A}}$ to obtain the points of the linear section $p.X_{\mathcal{A}} \cap L$.
- (iv) Use the linear section homotopy (Proposition 1) beginning with the points of $p.X_{\mathcal{A}} \cap L$ to obtain the points of the linear section $\mathcal{X}_1 \cap L$.

The discussion preceding Algorithm 3 justifies the following theorem:

Theorem 4. *Algorithm 3 is an optimal homotopy algorithm for computing $\mathcal{X}_1 \cap L$.*

Remark 5. Algorithm 3 can be applied to compute $\mathcal{X}_1 \cap L$ when the definition of a toric degeneration is relaxed so that \mathcal{X}_0 is a union of toric varieties (see Remark 8 for examples). The points in a general linear section $\mathcal{X}_0 \cap L$ in Algorithm 3 may be computed from systems of sparse polynomials for each toric component of \mathcal{X}_0 . \diamond

1.4. Examples of toric degenerations. We present three examples of toric degenerations. Example 6 is the motivating example for this paper. Example 9 illustrates an alternate source of toric degenerations. Example 10 is an explicit application of Algorithm 3.

Example 6. Weight degenerations induced by a \mathbb{C}^\times -action on \mathbb{P}^n are a source of toric degenerations. Anderson [2] constructs a toric weight degeneration given a Khovanskii basis. The SAGBI homotopy [14] is also based on a toric weight degeneration.

We review the construction in [11, Section 15.8] of flat families from \mathbb{C}^\times -actions. Let $w \in \mathbb{Z}^{n+1}$ be a weight and define an action of the torus \mathbb{C}^\times on \mathbb{P}^n by

$$(x, t) \in \mathbb{P}^n \times \mathbb{C}^\times \mapsto t.x := [x_0 t^{-w_0}, \dots, x_n t^{-w_n}] \in \mathbb{P}^n.$$

The dual action on functions is $t.f(x) := f(t^{-1}.x)$, and it induces an action on polynomials. For a polynomial $f = \sum c_\alpha x^\alpha$,

$$(4) \quad t. \left(\sum c_\alpha x^\alpha \right) = \sum c_\alpha x^\alpha t^{w \cdot \alpha},$$

where $w \cdot \alpha$ is the usual dot product. (To compare this to [11, Section 15.8], let $w = -\lambda$.) Let $w(f)$ be the minimum value of $w \cdot \alpha$ for $c_\alpha \neq 0$. We define

$$(5) \quad f_t := (t.f) t^{-w(f)} = f_w + t g,$$

where the *initial form* f_w of f is the sum of its terms $c_\alpha x^\alpha$ where $w \cdot \alpha = w(f)$, and g is a polynomial in the variables t, x_0, \dots, x_n .

Let $X \subset \mathbb{P}^n$ be a projective variety with ideal I . Define $\mathcal{X}^w \subset \mathbb{P}^n \times \mathbb{C}$ to be the Zariski closure of the family of translates of X , thus

$$\mathcal{X}^w := \overline{\{(x, t) \in \mathbb{P}^n \times \mathbb{C}^\times : x \in t.X\}} \subset \mathbb{P}^n \times \mathbb{C}_t.$$

For $t \neq 0$, we observe that $\mathcal{X}_t^w = t.X$ and has ideal $\langle f_t : f \in I \rangle$. The following result establishes the flatness of this family:

Proposition 7 ([11, Theorem 15.17]). *The family $\mathcal{X}^w \rightarrow \mathbb{C}_t$ is flat. The fiber at $t = 0$ is the scheme with ideal*

$$I_w = \langle f_w : f \in I \rangle.$$

The proof uses a Gröbner basis \mathcal{G} for I with respect to a weighted term order \leq with weight $-w$ so that f_w consists of the \leq -leading terms for f .¹

Suppose that the family \mathcal{X}^w has Gröbner basis $\mathcal{G}_t := \{g_t : g \in \mathcal{G}\}$. Then a Gröbner basis for I_w is obtained by setting $t = 0$ in \mathcal{G}_t . The scheme at $t = 0$ may be neither reduced nor irreducible. If this scheme is a toric variety, then the weight degeneration $\mathcal{X}^w \rightarrow \mathbb{C}_t$ is a toric degeneration. \diamond

Remark 8. Homotopy algorithms using weight degenerations appearing in the literature include the homotopy for solving the Kuramoto equations [6] and the Gröbner homotopy [14]. In these examples, I_w is a square-free monomial ideal so that the special fiber is a union of linear spaces. Such degenerations can be handled by Algorithm 3, see Remark 5. \diamond

Example 9. Algebraic statistics gives examples of toric degenerations [17] which do not come from a weight degeneration. Let G be a graph with vertex set $[m] := \{1, \dots, m\}$ and edge set $E \subset \binom{[m]}{2}$. For each $i \in [m]$, let a_i be a parameter. For each $\{i, j\} \in E$, let $x_{ij} = x_{ji}$ and define p_{ij} and p_{ji} via the formula

$$p_{rs} := x_{rs}(1 + a_r - ta_s).$$

These polynomials give a map $p: \mathbb{C}^{|E|} \times \mathbb{C}^m \times \mathbb{C}_t \rightarrow \mathbb{P}^{2|E|-1} \times \mathbb{C}_t$ whose image is the family \mathcal{QS} of *quasi-symmetry models*. This family contains two known quasi-symmetry models, the Pearsonian quasi-symmetry model at $t = 1$ and the toric quasi-symmetry model at $t = 0$. Polynomials associated to cycles in G generate the ideal of the family \mathcal{QS} . In the proof of this fact, one step is to show that this family is flat.

The family of quasi-symmetry models when G is a 3-cycle is the family of hypersurfaces defined by the cubic

$$(6) \quad P := (1 + t + t^2)(p_{12}p_{23}p_{31} - p_{21}p_{32}p_{13}) + t(p_{12}p_{23}p_{13} + p_{12}p_{32}p_{31} + p_{21}p_{23}p_{31} - p_{12}p_{32}p_{13} - p_{21}p_{23}p_{13} - p_{21}p_{32}p_{31}).$$

The fiber \mathcal{QS}_0 at $t = 0$ is the toric variety defined by the binomial $p_{12}p_{23}p_{31} - p_{21}p_{32}p_{13}$.

The family of quasi-symmetry models \mathcal{QS} for a graph is typically not a weight degeneration. In particular, the family defined in Equation (6) is not a weight degeneration. Indeed, in each of the eight terms of P , exactly one of p_{ij} or p_{ji} occurs, so the terms correspond to the vertices of a cube. For any weight w , P_w consists of the sum of terms identified with some face of the cube. Since the polynomial defining \mathcal{QS}_0 corresponds to a diagonal of the cube, it is not of the form P_w , for any w . \diamond

Example 10. We present an example of a weight degeneration and use it to illustrate Algorithm 3. Let $X \subset \mathbb{P}^7$ be the closure of the image of the map $\varphi: \mathbb{C}^3 \rightarrow \mathbb{P}^7$ given by

$$\varphi(x, y, z) = [1, x, y, z, xz, yz, x(xz + y), y(xz + y)].$$

This subvariety has degree six and its ideal I has nine generators:

$$\begin{aligned} & x_1x_3 - x_0x_4, \quad x_2x_3 - x_0x_5, \quad x_1x_2 - x_0x_6 + x_1x_4, \quad x_2^2 - x_0x_7 + x_3x_6 - x_4^2, \\ & x_2x_6 - x_1x_7, \quad x_2x_5 - x_3x_7 + x_4x_5, \quad x_1x_5 - x_3x_6 + x_4^2, \quad x_2x_4 - x_1x_5, \quad x_5x_6 - x_4x_7. \end{aligned}$$

¹We use $-w$ because the leading form in the weighted term order \leq_w for $w \in \mathbb{Z}^{n+1}$ is the sum of terms with highest w -weight, which is opposite our convention from valuations.

Let $w = (-2, -1, -1, -1, 0, 0, 0, 0)$. We use Equation 5 to compute the ideal of \mathcal{X}^w . The following thirteen polynomials form a Gröbner basis \mathcal{G}_t for \mathcal{X}^w with respect to the weighted term order \leq_{-w} :

$$\begin{aligned} & \underline{x_1x_3 - x_0x_4}, \quad \underline{x_2x_3 - x_0x_5}, \quad \underline{x_1x_2 - x_0x_6} + tx_1x_4, \quad \underline{x_2^2 - x_0x_7} + tx_3x_6 - t^2x_4^2, \\ & \underline{x_2x_6 - x_1x_7}, \quad \underline{x_2x_5 - x_3x_7} + tx_4x_5, \quad \underline{x_1x_5 - x_3x_6} + tx_4^2, \quad \underline{x_2x_4 - x_1x_5}, \quad \underline{x_5x_6 - x_4x_7}, \\ & \underline{x_0x_6^2 - x_1^2x_7} - tx_1x_4x_6, \quad \underline{x_0x_5^2 - x_3^2x_7} + tx_3x_4x_5, \\ & \underline{x_0x_4x_5 - x_3^2x_6} + tx_3x_4^2, \quad \underline{x_3x_6^2 - x_1x_4x_7} - tx_4^2x_6. \end{aligned}$$

The leading terms with respect to \leq_{-w} are underlined, and these binomials generate the ideal I_w . This ideal is the toric ideal of the image of the map $\varphi_{\mathcal{A}}(x, y, z) = [1, x, y, z, xz, yz, xy, y^2]$ given by the lowest order monomials in φ . For the toric ideal statement, observe that if we set $(x_0, x_1, x_2, x_3) = (1, x, y, z)$, then the first four underlined binomials in \mathcal{G}_t express x_4, \dots, x_7 as the monomials in x, y, z appearing in $\varphi_{\mathcal{A}}$. The exponent vectors of $\varphi_{\mathcal{A}}$ are the columns of the matrix \mathcal{A} in Figure 1.

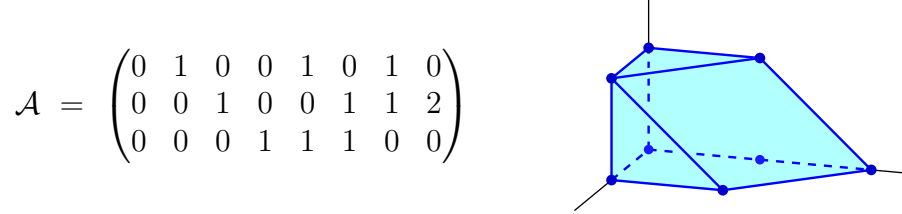


FIGURE 1. The weight vectors for the toric Kodaira map $\varphi_{\mathcal{A}}$ are the columns of matrix \mathcal{A} . The Newton polytope is the convex hull of these vectors.

Let $L \subset \mathbb{P}^7$ be the linear subspace of codimension three whose defining equations are $\ell_i = \sum c_{ij}x_j$ for $i = 1, 2, 3$, where $C = (c_{ij})$ is the 3×8 matrix

$$C = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & -2 & 3 & -4 & 5 & -6 & 7 & -8 \\ 2 & 3 & 5 & 7 & 11 & 13 & 17 & 19 \end{pmatrix}.$$

The subspace L meets both $\overline{\varphi(\mathbb{C}^3)} = \mathcal{X}_1$ and $\overline{\varphi_{\mathcal{A}}((\mathbb{C}^\times)^3)} = \mathcal{X}_0$ transversally in six points.

We follow the steps of Algorithm 3 to compute $\mathcal{X}_1 \cap L$. We first compute the sparse system $G_{\mathcal{A}}$ in Step (i) of Algorithm 3 to arrive at the system

$$\begin{aligned} 1 + x + y + z + xz + yz + xy + y^2 &= 0 \\ 1 - 2x + 3y - 4z + 5xz - 6yz + 7xy - 8y^2 &= 0 \\ 2 + 3x + 5y + 7z + 11xz + 13yz + 17xy + 19y^2 &= 0. \end{aligned}$$

In Step (ii) of Algorithm 3, we compute the six solutions of the system $G_{\mathcal{A}}$, one of which is $\zeta = (-1.33613, 1.51406, -1.22871)$. The image $\varphi_{\mathcal{A}}(\zeta)$ in \mathbb{P}^7 is

$$[1, -1.33613, 1.51406, -1.22871, 1.64171, -1.86035, -2.02298, 2.29239].$$

In Step (iii) of Algorithm 3, we compute the images of these six solutions under $\varphi_{\mathcal{A}}$, which forms the points of $\mathcal{X}_0 \cap L$. Therefore, the images of these points are the solutions to the

start system for the linear section homotopy given by $H(x, t) = (\mathcal{G}_t, L)$. In Step (iv), these solutions are followed from $t = 0$ to $t = 1$, computing the six points of the linear section $\mathcal{X}_1 \cap L$. One point of $\mathcal{X}_1 \cap L$ is

$$[1, -0.689522, 0.928435, -1.35986, 0.937652, -1.26254, -1.28671, 1.73254]. \quad \diamond$$

2. KHOVANSKII BASES AND THE KHOVANSKII HOMOTOPY

Let X be a complex variety and $V \subset \mathbb{C}(X)$ be a finite-dimensional complex vector space of rational functions on X . The closure of the image of X under the Kodaira map $\varphi_V: X \dashrightarrow \mathbb{P}(V^*)$ has homogeneous coordinate ring $R(V)$ generated by V . When this ring has a finite Khovanskii basis contained in V , Anderson's toric degeneration embeds in $\mathbb{P}(V^*)$ as a weight degeneration. We use this degeneration in the Khovanskii homotopy algorithm (Algorithm 14) to compute a linear section $\varphi_V(X) \cap L$.

We review the theory of Newton-Okounkov bodies and Khovanskii bases and then describe how to produce an embedding of Anderson's toric degeneration into $\mathbb{P}(V^*)$ when the Khovanskii basis is a subset of V . We also show how to compute a Kodaira map of the toric special fiber. With the embedding and toric Kodaira map, Algorithm 3 becomes an effective method to compute linear sections. In Section 3, we explain how to modify this method for the general case when the Khovanskii basis is not a subset of V .

2.1. Valuations, Khovanskii bases, and Newton-Okounkov bodies. We recall the key definitions and properties of Khovanskii bases from [20]. Suppose that X is a d -dimensional complex variety with function field $\mathbb{C}(X)$. Let \succ be a total order on \mathbb{Z}^d so that \mathbb{Z}^d is an ordered abelian group. A \mathbb{Z}^d -valuation on $\mathbb{C}(X)$ is a surjective group homomorphism $\nu: \mathbb{C}(X)^\times \mapsto \mathbb{Z}^d$ satisfying the property that for all $f, g \in \mathbb{C}(X)$ and $c \in \mathbb{C}^\times$,

$$\nu(f + g) \succeq \min\{\nu(f), \nu(g)\} \quad \text{and} \quad \nu(c) = 0.$$

By convention, $\nu(0) = \infty$, $\infty \succeq \alpha$, and $\alpha + \infty = \infty$ for all $\alpha \in \mathbb{Z}^d$. Since $\dim X = d$, ν is a surjection, and \mathbb{C} is algebraically closed, it follows that if $f, g \in \mathbb{C}(X)^\times$ with $\nu(f) = \nu(g)$, then there is a unique $c \in \mathbb{C}^\times$ with $\nu(f - cg) \succ \nu(f)$.

Let V be a finite-dimensional complex vector subspace of $\mathbb{C}(X)$. We assume that the image of V^\times under ν generates \mathbb{Z}^d (see Remark 11). We write $R(V)$ for the graded ring $\bigoplus_{k \geq 0} V^k s^k$, where $V^k \subset \mathbb{C}(X)$ is the subspace spanned by all k -fold products of elements in V and s is a formal variable recording the grading. A nonzero element $f \in R(V)^\times$ is the sum of its homogeneous components,

$$f = f_k s^k + \cdots + f_1 s + f_0,$$

where $f_k \neq 0$ and $f_i \in V^i$ for all i . We extend the valuation ν to $R(V)$ by defining $\nu(f) := (\nu(f_k), k) \in \mathbb{Z}^d \oplus \mathbb{N}$. We also extend \succeq to $\mathbb{Z}^d \oplus \mathbb{N}$, where $(\alpha, k) \succ (\beta, l)$ if $k < l$ or else $k = l$ and $\alpha \succ \beta$ in the order on \mathbb{Z}^d . The direction of the inequality in $k < l$ is chosen to be consistent with $\nu(f) = (\nu(f_k), k)$ defining a valuation.

We write $S(V, \nu)$ for the image $\{\nu(f) : f \in R(V)^\times\}$ of $R(V)^\times$ under ν . This is a submonoid of $\mathbb{Z}^d \oplus \mathbb{N}$. The closure of the convex hull of $S(V, \nu)$ in $\mathbb{R}^d \times \mathbb{R}$ is the cone $\text{cone}(V)$. Its base $\text{NO}_V := \text{cone}(V) \cap (\mathbb{R}^d \times \{1\})$ is the *Newton-Okounkov body* of V . The Newton-Okounkov body carries a considerable amount of information about $R(V)$, see [19, 21]. For example,

the number of solutions to System (1) where $f_1, \dots, f_d \in V$ are general (in this case, we say that System (1) is *drawn* from V) is the normalized volume of NO_V .

A *Khovanskii basis* [20] for V is a linearly independent set $\mathcal{B} \subset R(V)$ whose image under ν generates $S(V, \nu)$. We assume that the elements of \mathcal{B} are homogeneous so that for $b \in \mathcal{B}$ with $\nu(b) = (\alpha, k)$, $b \in V^k s^k$. Necessarily, \mathcal{B} generates $R(V)$ and $\mathcal{B} \cap Vs$ is a basis for Vs . We observe that $S(V, \nu)$ is finitely generated if and only if V has a finite Khovanskii basis. When V has a finite Khovanskii basis, Anderson [2] shows that NO_V is a rational polytope and that there exists a flat degeneration $\mathcal{X} \rightarrow \mathbb{C}_t$ of $\mathcal{X}_1 \simeq \text{Proj}(R(V))$ to the toric variety $\mathcal{X}_0 \simeq \text{Proj}(\mathbb{C}[S(V, \nu)])$.

The valuation ν on $R(V)$ induces a filtration on $R(V)$ by finite-dimensional subspaces indexed by elements $(\alpha, k) \in S(V, \nu)$. We let

$$\begin{aligned} R(V)_{(\alpha, k)} &:= \{f \in R(V) : \nu(f) \succeq (\alpha, k)\}, \text{ and} \\ R(V)_{(\alpha, k)}^+ &:= \{f \in R(V) : \nu(f) \succ (\alpha, k)\}. \end{aligned}$$

Since $(\alpha, k) \in S(V, \nu)$, these subspaces satisfy $R(V)_{(\alpha, k)}/R(V)_{(\alpha, k)}^+ \simeq \mathbb{C}$. Anderson's flat degeneration comes from the degeneration of the filtered algebra $R(V)$ to its associated graded algebra

$$\text{gr } R(V) := \bigoplus_{(\alpha, k) \in S(V, \nu)} R(V)_{(\alpha, k)}/R(V)_{(\alpha, k)}^+ \simeq \mathbb{C}[S(V, \nu)].$$

The toric fiber \mathcal{X}_0 of Anderson's degeneration is $\text{Proj}(\text{gr } R(V))$, and the isomorphism $\mathcal{X}_0 \simeq \text{Proj}(\text{gr } R(V))$ uses the isomorphism $\text{gr } R(V) \simeq \mathbb{C}[S(V, \nu)]$.

Kaveh and Manon give a method to compute a finite Khovanskii basis for V with respect to a valuation ν [20, Algorithm 2.18]. We take a finite Khovanskii basis as an input to our algorithms.

2.2. The Kodaira map and embedding the degeneration. To use Anderson's toric degeneration \mathcal{X} in Algorithm 3, \mathcal{X} must be embedded in a projective space. Suppose that we are given a finite Khovanskii basis \mathcal{B} for V such that $\mathcal{B} \subset Vs$. Therefore, \mathcal{B} is a basis for Vs , by definition.

Let $X^\circ \subset X$ be the open subset of points of X where no function from V has a pole, and some function in V is nonzero. Evaluation of functions from V at a point $z \in X^\circ$ gives a nonzero linear map $\text{ev}_z(f) := f(z)$ on V . Therefore, ev_z is a point in the projective space $\mathbb{P}(V^*)$, where V^* is the space of linear functions $V \rightarrow \mathbb{C}$. Thus the map $z \mapsto \text{ev}_z$ induces a map $X^\circ \rightarrow \mathbb{P}(V^*)$, which is called the *rational Kodaira map* $\varphi_V : X \dashrightarrow \mathbb{P}(V^*)$. If we write $\mathcal{B} = \{b_0s, \dots, b_ns\}$, then a Kodaira map can be explicitly written as $\varphi_{\mathcal{B}} : z \in X^\circ \mapsto [b_0(z), \dots, b_n(z)] \in \mathbb{P}^n \simeq \mathbb{P}(V^*)$.

Remark 11. Our algorithms compute the points of $\varphi_{\mathcal{B}}(X^\circ) \cap L$. Given these points, the solutions to System (1) on X° are their pull backs along $\varphi_{\mathcal{B}}$. When the Kodaira map is not injective, we follow Améndola and Rodriguez [1] and note that these pull backs may be computed from the linear section and the points in a single general fiber of the Kodaira map.

Consequently, we assume that the Kodaira map is an injection and replace X by its birational copy $\text{Proj}(R(V))$, which is the closure of $\varphi_{\mathcal{B}}(X^\circ)$ in \mathbb{P}^n . In this case, $X = X^\circ$, V

generates the function field $\mathbb{C}(X)$ of X , and the image of V^\times under ν generates \mathbb{Z}^d . Thus the assumption that $X = \text{Proj}(R(V))$ implies that the image of V^\times under ν generates \mathbb{Z}^d . \diamond

We recall the embedding of Anderson's toric degeneration \mathcal{X} into $\mathbb{P}(V^*)$ [20, Section 2.2]. We let $\mathcal{A} := \nu(\mathcal{B})$ be the $(d+1) \times (n+1)$ matrix whose i^{th} column is $\nu(b_{i-1}s)$ for the Khovanskii basis $\mathcal{B} = \{b_0s, \dots, b_ns\} \subset Vs$. We note that the last row of \mathcal{A} is $(1, \dots, 1)$. We define a partial order $>_{\mathcal{A}}$ on \mathbb{Z}^{n+1} where $\beta >_{\mathcal{A}} \alpha$ if $\mathcal{A}\alpha \succ \mathcal{A}\beta$ in \mathbb{Z}^{d+1} . The initial form $\text{in}_{\mathcal{A}}(f)$ of a polynomial f with respect to $>_{\mathcal{A}}$ is the sum of all terms $c_\alpha x^\alpha$ which minimize $\mathcal{A}\alpha$.

The ideal $I_{\mathcal{B}}$ of $X = \varphi_{\mathcal{B}}(X)$ is the kernel of the map $\mathbb{C}[x_0, \dots, x_n] \rightarrow R(V)$ which takes x_i to $b_i s$. We define $\text{in}_{\mathcal{A}}(I_{\mathcal{B}})$ to be the ideal generated by $\text{in}_{\mathcal{A}}(f)$ for $f \in I_{\mathcal{B}}$. Anderson [2, Lemma 8] shows that there exists $w \in \mathbb{Z}^{d+1}$ such that if $\leq_{-w\mathcal{A}}$ is the weighted term order on $\mathbb{C}[x_0, \dots, x_n]$ induced by $-w\mathcal{A}$, then the leading term ideal $\text{lt}_{-w\mathcal{A}}(I_{\mathcal{B}})$ of $I_{\mathcal{B}}$ equals $\text{in}_{\mathcal{A}}(I_{\mathcal{B}})$. Let w be such a weight vector and \mathcal{G} denote a Gröbner basis for $I_{\mathcal{B}}$ with respect to a total order induced from the term order $\leq_{-w\mathcal{A}}$. The leading forms of elements of \mathcal{G} with respect to $\leq_{-w\mathcal{A}}$ generate $\text{in}_{\mathcal{A}}(I_{\mathcal{B}})$.

Let $g = \sum_{\alpha} c_{\alpha} x^{\alpha}$ be a polynomial in \mathcal{G} , and define $w(g) := \min\{w\mathcal{A}\alpha : c_{\alpha} \neq 0\}$. Using Formula (5) (with $w\mathcal{A}$ in place of w), we construct

$$(7) \quad g_t = \sum_{\alpha} c_{\alpha} x^{\alpha} t^{w\mathcal{A}\alpha - w(g)}.$$

Let $\mathcal{G}_t := \{g_t : g \in \mathcal{G}\}$. At $t = 0$, \mathcal{G}_0 generates $\text{in}_{\mathcal{A}}(I_{\mathcal{B}})$ and at $t = 1$, $\mathcal{G}_1 = \mathcal{G}$ generates $I_{\mathcal{B}}$.

Finally, we define $I_{\mathcal{A}}$ to be the kernel of the map $\mathbb{C}[x_0, \dots, x_n] \rightarrow \text{gr } R(V)$ which takes x_i to $\overline{b_i s} \in R(V)_{(\nu(b_i), 1)} / R(V)_{(\nu(b_i), 1)}^+$. We note that $I_{\mathcal{A}}$ is a toric ideal, and by [20, Theorem 2.17], $I_{\mathcal{A}} = \text{in}_{\mathcal{A}}(I_{\mathcal{B}})$. Thus the toric weight degeneration can be embedded into $\mathbb{P}^n \simeq \mathbb{P}(V^*)$.

Proposition 12 ([2, Theorem 1]). *Let X be a variety and $V \subset \mathbb{C}(X)$ a finite-dimensional space of functions which has a finite Khovanskii basis $\mathcal{B} \subset Vs$. Then the family $\mathcal{X} \rightarrow \mathbb{C}_t$ defined by \mathcal{G}_t is flat and embeds into \mathbb{P}^n as the weight degeneration of $\mathcal{X}_1 = \text{Proj}(R(V)) = \varphi_{\mathcal{B}}(X)$ induced by $w\mathcal{A}$. In particular, $\mathcal{X}_0 \simeq \text{Proj}(\mathbb{C}[S(V, \nu)])$ and \mathcal{X} is a toric degeneration.*

We now discuss the relationship between $I_{\mathcal{A}}$ and $I_{\mathcal{B}}$. For $u \in \mathbb{N}^{n+1}$, we write \mathcal{B}^u for the product $\prod (b_i s)^{u_i}$ of elements in the Khovanskii basis. Since $\nu(\mathcal{B}^u) = \mathcal{A}u$, when $\mathcal{A}u = \mathcal{A}v$ for some $u, v \in \mathbb{N}^{n+1}$, $\nu(\mathcal{B}^u) = \nu(\mathcal{B}^v)$ and there is a unique $c \in \mathbb{C}^\times$ such that

$$\mathcal{A}u = \mathcal{A}v \prec \nu(\mathcal{B}^u - c\mathcal{B}^v) \quad \text{and} \quad \mathcal{B}^u - c\mathcal{B}^v \in R(V)_{\mathcal{A}u}^+.$$

Since the last row of \mathcal{A} is $(1, \dots, 1)$, both \mathcal{B}^u and $c\mathcal{B}^v \in V^k s^k$ for some k and their difference is homogeneous.

The subduction algorithm [20, Algorithm 2.11] rewrites this difference as a homogeneous polynomial of degree k in the elements of the Khovanskii basis,

$$\mathcal{B}^u - c\mathcal{B}^v = h(b_0s, b_1s, \dots, b_ns).$$

In particular, $g := x^u - cx^v - h(x_0, \dots, x_n) \in I_{\mathcal{B}}$ with initial form $x^u - cx^v \in I_{\mathcal{A}}$. Applying Formula (7), we have that

$$g_t = x^u - cx^v - t^r h_t,$$

where $r = w(h) - w(g) > 0$.

Remark 13. We recall that the torus $\mathbb{T} = (\mathbb{C}^\times)^{n+1}/\Delta\mathbb{C}^\times \simeq (\mathbb{C}^\times)^n$ is the set of points in \mathbb{P}^n with nonzero coordinates. A Kodaira map for the toric fiber \mathcal{X}_0 has the form $\varphi_{p,\mathcal{A}}$, as in Formula (2), for any $p \in \mathbb{T} \cap \mathcal{X}_0$. We provide a construction of such a point.

Let $x^u - cx^v \in I_{\mathcal{A}}$. Then $\mathcal{A}u = \mathcal{A}v$, so that $u - v \in \ker(\mathcal{A})$. Restricting this binomial to $\mathcal{X}_0 \cap \mathbb{T}$ results in the equation $c = x^{u-v}$. The constant c depends upon $u - v \in \ker(\mathcal{A})$, and we write c_{u-v} for c . Thus a point $p \in \mathcal{X}_0 \cap \mathbb{T}$ satisfies equations of the form

$$c_u = p^u$$

for $u \in \ker(\mathcal{A})$. While every $u \in \ker(\mathcal{A})$ gives such an equation, an independent set of equations is given by a basis u_1, \dots, u_{n-d} for $\ker(\mathcal{A})$. The corresponding equations, $c_{u_i} = p^{u_i}$ for $i = 1, \dots, n-d$, define $\mathcal{X}_0 \cap \mathbb{T}$ as a subvariety of \mathbb{T} .

To obtain a point of $\mathcal{X}_0 \cap \mathbb{T}$, we construct d additional equations to these $n-d$ equations, as follows: Since $\mathbb{1} := (1, \dots, 1)$ is a row of \mathcal{A} , $\ker(\mathcal{A}) \subset \ker(\mathbb{1})$, which is a rank n sublattice of \mathbb{Z}^{n+1} . Let $v_1, \dots, v_d \in \ker(\mathbb{1})$ be vectors such that $u_1, \dots, u_{n-d}, v_1, \dots, v_d$ are independent. Choose nonzero constants $c_{v_1}, \dots, c_{v_d} \in \mathbb{C}^\times$ and consider the system of binomials

$$c_{u_i} - p^{u_i} = 0 = c_{v_j} - p^{v_j} \quad \text{for } i = 1, \dots, n-d \text{ and } j = 1, \dots, d.$$

This system defines a finite set of points $p \in \mathcal{X}_0 \cap \mathbb{T}$. An algorithm for solving such a system of binomials is given in [15, Lemma 3.2], which involves computing the Smith normal form of the matrix whose columns are $u_1, \dots, u_{n-d}, v_1, \dots, v_d$. We observe that only one solution is needed to obtain a Kodaira map. \diamond

2.3. Khovanskii homotopy. The procedure described in Section 2.2, combined with the toric two-step homotopy algorithm, Algorithm 3, forms the Khovanskii homotopy algorithm for computing the points of a linear section $\varphi_V(X) \cap L$.

Algorithm 14 (Khovanskii homotopy algorithm).

Input: A finite-dimensional subspace $V \subset \mathbb{C}(X)$ for a variety $X = \text{Proj}(R(V))$ of dimension d , a finite Khovanskii basis $\mathcal{B} \subset V$ for V , and a general linear subspace $L \subset \mathbb{P}^n$ of codimension d .

Output: All points in the linear section $\varphi_V(X) \cap L \subset \mathbb{P}(V^*)$.

Do:

- (i) Compute $I_{\mathcal{B}} = \ker(\mathbb{C}[x_0, \dots, x_n] \rightarrow R(V))$ where $x_i \mapsto b_i s$.
- (ii) Compute a weight vector w using [2, Lemma 2] so that $\text{lt}_{-w\mathcal{A}}(I_{\mathcal{B}}) = \text{in}_{\mathcal{A}}(I_{\mathcal{B}})$, where \mathcal{A} is the matrix of valuations of \mathcal{B} .
- (iii) Compute a Gröbner basis \mathcal{G} for $I_{\mathcal{B}}$ using the weight $-w\mathcal{A}$.
- (iv) Construct the homotopy \mathcal{G}_t using Formula (7).
- (v) Construct the Kodaira map $\varphi_{p,\mathcal{A}}$ for \mathcal{X}_0 by following Remark 13.
- (vi) Return the output $\varphi_V(X) \cap L$ of Algorithm 3 with input \mathcal{G}_t and L .

Theorem 15. *Algorithm 14 is an optimal homotopy algorithm for computing all points of $\varphi_V(X) \cap L$.*

The correctness of Algorithm 14 follows from the discussion in Section 2.2.

Remark 16. In many cases, Algorithm 14 is applied to systems of functions where a finite Khovanskii basis is explicitly known from the theory (see Example 17). In this case, we not

only have the data for the finite Khovanskii basis \mathcal{B} , but also some or all of the data for Steps (i), (ii), and (iii) of Algorithm 14. \diamond

Example 17. We illustrate Algorithm 14 and Remark 16 on a continuation of Example 10. In [2, Section 6.4], Anderson considers a particular three-dimensional Bott-Samelson variety X for $GL(3, \mathbb{C})$ and an ample line bundle \mathcal{L} on X . In local coordinates (x, y, z) for X , the vector space V of global sections of \mathcal{L} has basis $\{1, x, y, z, xz, yz, x(xz + y), y(xz + y)\}$.

Anderson uses a valuation ν induced by the monomial valuation on $\mathbb{C}[x, y, z]$ defined by $\nu(f) = (a, b, c)$, where $x^a y^b z^c$ is the monomial of f that is minimal in the degree lexicographic order with $x > y > z$. The image $\mathcal{B} = \{1s, xs, ys, zs, xzs, yzs, x(xz + y)s, y(xz + y)s\}$ of this basis in Vs forms a Khovanskii basis for V . The corresponding matrix of valuations is

$$\mathcal{A} = \nu(\mathcal{B}) = \begin{pmatrix} 0 & 1 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 2 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix},$$

which is the matrix of Figure 1 after appending the row $\mathbb{1}$ for the exponents of s . The Newton-Okounkov body of V is also displayed in Figure 1.

Anderson provides the Khovanskii basis \mathcal{B} for Algorithm 14, and Example 10 gives the general linear section L . For Step (i), generators of $I_{\mathcal{B}}$ are the generators of I in Example 10. The weight vector $w = (1, 1, 1, -2)$ suffices for Step (ii). The vector $w\mathcal{A} = (-2, -1, -1, -1, 0, 0, 0, 0)$ appears as the weight in Example 10. The computations in Steps (iii) and (iv) are supplied by the elements in \mathcal{G}_t in Example 10. Finally, for Step (v), the toric Kodaira map $\varphi_{\mathcal{A}}$ is also given in Example 10. \diamond

3. THE KHOVANSKII HOMOTOPY FOR WEIGHTED PROJECTIVE SPACE

When a Khovanskii basis \mathcal{B} for V contains elements of degree greater than 1, Anderson's toric degeneration naturally embeds into a weighted projective space [2]. We explain how to lift the degeneration to a toric degeneration in ordinary projective space and use the toric two-step homotopy (Algorithm 3) to compute a linear section $\varphi_V(X) \cap L$ of the image of X under the Kodaira map $\varphi_V: X \dashrightarrow \mathbb{P}(V^*)$.

3.1. Weighted projective spaces. We recall the construction and some basic properties of weighted projective space, see [9]. Suppose that $a = (a_0, \dots, a_{n+m})$ is a vector of mutually relatively prime positive integers. The *weighted projective space* \mathbb{P}_a^{n+m} is $\text{Proj}(\mathbb{C}[x_0, \dots, x_{n+m}])$, where the grading on $\mathbb{C}[x_0, \dots, x_{n+m}]$ is induced by setting the degree of x_j to a_j . Equivalently, \mathbb{P}_a^{n+m} is the quotient of $\mathbb{C}^{n+m+1} \setminus \{0\}$ by the \mathbb{C}^\times -action where $t \cdot (x_0, \dots, x_{n+m}) = (t^{a_0}x_0, \dots, t^{a_{n+m}}x_{n+m})$, for $t \in \mathbb{C}^\times$. We may also construct \mathbb{P}_a^{n+m} as a quotient of \mathbb{P}^{n+m} . To see this, let $\Delta\mathbb{C}^\times \subset (\mathbb{C}^\times)^{n+m+1}$ be the diagonal embedding of \mathbb{C}^\times and let G_a be the image of the following product of groups of roots of unity in the dense torus $(\mathbb{C}^\times)^{n+m+1}/\Delta(\mathbb{C}^\times)$ of \mathbb{P}^{n+m} .

$$\text{Hom}\left(\prod_j \mathbb{Z}/a_j\mathbb{Z}, \mathbb{C}^\times\right) = \prod_j \text{Hom}(\mathbb{Z}/a_j\mathbb{Z}, \mathbb{C}^\times) \subset (\mathbb{C}^\times)^{n+m+1}.$$

Thus G_a acts faithfully on \mathbb{P}^{n+m} . As the a_j are mutually relatively prime, G_a is isomorphic to this product of groups of roots of unity. Let $\pi: \mathbb{P}^{n+m} \rightarrow \mathbb{P}_a^{n+m}$ be the quotient map by this G_a -action, which is a finite map of degree $|G_a| = \prod a_j$.

The weighted projective spaces that appear in the Khovanskii homotopy have the following special form: Let $W = \bigoplus_{k \geq 1} W_k$ be a finite-dimensional positively-graded vector space with $\dim W_1 = n+1 \geq 1$. Let $t \in \mathbb{C}^\times$ act on W_k as multiplication by t^{-k} , which gives a \mathbb{C}^\times -action on W . Identifying the dual space W^* with $\bigoplus_k W_k^*$, in the dual action, $t \in \mathbb{C}^\times$ acts on W_k^* as multiplication by t^k . Then the quotient of $W^* \setminus \{0\}$ by \mathbb{C}^\times is a weighted projective space.

We explicitly describe this weighted projective space. Suppose that $\dim W = n+m+1$, and let $a = (a_0, \dots, a_{n+m})$ be a vector in which each $k \in \mathbb{N}$ occurs $\dim W_k$ times. Then $(W^* \setminus \{0\})/\mathbb{C}^\times$ is isomorphic to \mathbb{P}_a^{n+m} , and we write $\mathbb{P}_a(W^*)$ for this quotient. The isomorphism depends upon the choice of an ordered basis for W^* which is a union of bases for each nontrivial summand W_k^* such that $a_j = k$ when the j th basis element lies in W_k^* . This choice of basis identifies W^* with \mathbb{C}^{n+m+1} , and allows us to define an action of G_a on the projective space \mathbb{P}^{n+m} with quotient map $\pi: \mathbb{P}^{n+m} \rightarrow \mathbb{P}_a(W^*)$ as in the first paragraph above. We remark that there is no natural identification of $\mathbb{P}(W^*)$ with \mathbb{P}^{n+m} that is compatible with the map π , unless $\dim W_k \leq 1$ for all $k > 1$.

Let us write V for W_1 . Under the \mathbb{C}^\times -action given by the weight a , the composition $V \hookrightarrow W \twoheadrightarrow V$ of the inclusion with the projection onto V is the identity and each map is \mathbb{C}^\times -equivariant. Taking linear duals gives the equivariant composition $V^* \hookrightarrow W^* \twoheadrightarrow V^*$, and this induces the composition $\mathbb{P}(V^*) \hookrightarrow \mathbb{P}_a(W^*) \dashrightarrow \mathbb{P}(V^*)$. We obtain ordinary projective space $\mathbb{P}(V^*)$ because $t \in \mathbb{C}^\times$ acts as multiplication by t on V^* . We write pr_a for the projection map $\mathbb{P}_a(W^*) \dashrightarrow \mathbb{P}(V^*)$, which is undefined on the image of the annihilator of V in $\mathbb{P}_a(W^*)$. In addition, we write pr for the composition $pr_a \circ \pi$. We summarize these maps in the following commutative diagram:

$$\begin{array}{ccc} G_a \curvearrowright \mathbb{P}^{n+m} & & \\ \pi \downarrow & \searrow pr & \\ \mathbb{P}_a^{n+m} \simeq \mathbb{P}_a(W^*) & \dashrightarrow pr_a & \mathbb{P}(V^*) \simeq \mathbb{P}^n. \end{array}$$

Let $X \subset \mathbb{P}(V^*)$ and $Z \subset \mathbb{P}_a(W^*)$ be varieties such that pr_a is an isomorphism between Z and X . In this case, a linear section $X \cap L$ pulls back along pr_a to $Z \cap pr_a^{-1}(L)$. We remark that the subvariety $pr_a^{-1}(L)$, which is given by d forms that are linear in x_0, \dots, x_n , is not general. For example, $pr_a^{-1}(L)$ includes $\mathcal{V}(x_0, \dots, x_n)$, which contains the singular locus of $\mathbb{P}_a(W^*)$. We let $U \subset \mathbb{P}_a(W^*)$ be the open subset over which π is a covering space. For $u \in U$, G_a acts freely on the fiber $\pi^{-1}(z)$. The following lemma relates $Z \cap pr_a^{-1}(L)$ to $X \cap L$:

Lemma 18. *Let $Z \subset \mathbb{P}_a(W^*)$ be a subvariety of dimension d such that $Z \cap U$ is dense in Z and pr_a is an isomorphism between Z and $X := pr_a(Z)$. Let $Y := \pi^{-1}(Z) \subset \mathbb{P}^{n+m}$ be its inverse image. Suppose that $L \subset \mathbb{P}(V^*)$ is a general linear subspace of codimension d . Then,*

- (i) $Z \cap pr_a^{-1}(L)$ is transverse and $pr_a: Z \cap pr_a^{-1}(L) \rightarrow X \cap L$ is a bijection.
- (ii) $Y \cap pr^{-1}(L)$ is transverse and $\pi: Y \cap pr^{-1}(L) \rightarrow Z \cap pr_a^{-1}(L)$ is a $|G_a|$ to 1 surjection.

- (iii) For any component Y' of Y , $\pi: Y' \cap pr^{-1}(L) \rightarrow Z \cap pr^{-1}(L)$ is a $|Stab_{G_a}(Y')|$ to 1 surjection.

We note that $Y = \pi^{-1}(Z)$ may not be irreducible. Each irreducible component, however, maps surjectively onto Z .

Proof. We address transversality after establishing the set-theoretic assertions. For $x \in X \cap L$, let z be the unique point of Z with $pr_a(z) = x$. Since $z \in pr_a^{-1}(L)$, this completes the proof of the first statement.

Let $z \in Z \cap pr_a^{-1}(L)$. By our assumptions, $Z \cap pr_a^{-1}(L) \subset U$, so $z \in U$. Then $\pi^{-1}(z) \subset Y \cap \pi^{-1}pr_a^{-1}(L) = Y \cap pr^{-1}(L)$. The second statement follows as $\pi: Y \rightarrow Z$ is $|G_a|$ to 1 over points of U .

For the third statement, we observe that $pr^{-1}(L)$ is invariant under the G_a -action. Therefore, for all $g \in G_a$, $g.(Y' \cap pr^{-1}(L)) = (g.Y') \cap pr^{-1}(L)$. The claim follows from the second statement and a counting argument.

For transversality, let $x \in X \cap L$. As L is general, this intersection is transverse and the forms defining L generate the maximal ideal in the local ring of X at x . Transversality in the first statement follows since the map pr_a is an isomorphism between Z and X and $pr_a^{-1}(L)$ is defined by the same forms as L . Transversality in the second statement also follows, since the maximal ideal of Y at y is generated by the pull back of the maximal ideal of Z at $\pi(y)$ and $pr_a^{-1}(x) \in U$. \square

While $pr^{-1}(L)$ is a linear subspace, it is not general. We need a result similar to Lemma 18 for a general linear subspace $\Lambda \subset \mathbb{P}^{n+m}$. We note that since Λ is general, $\pi^{-1}(\pi(\Lambda))$ consists of a union of $|G_a|$ linear subspaces.

Lemma 19. *Let $Z \subset \mathbb{P}_a(W^*)$ be a subvariety of dimension d such that $Z \cap U$ dense in Z . Let $Y := \pi^{-1}(Z) \subset \mathbb{P}^{n+m}$ be its inverse image, and suppose that $\Lambda \subset \mathbb{P}^{n+m}$ is a general linear subspace of codimension d . Then, $\pi: Y \cap \Lambda \rightarrow Z \cap \pi(\Lambda)$ is a bijection.*

Proof. Since Λ is general, $Z \cap \pi(\Lambda) \subset U$. Suppose that $q, q' \in Y \cap \Lambda$ are in the same fiber of π , and let $g \in G_a$ be defined by $q' = g.q$. Since Y is G_a -invariant, we have $q' \in Y \cap (g.\Lambda)$. Since Λ is general, $Y \cap \Lambda \cap (g.\Lambda)$ is empty unless g is the identity. Therefore, $q = q'$, and we conclude that π is injective on $Y \cap \Lambda$.

This map is also surjective. If $p \in Z \cap \pi(\Lambda)$, then there is a point $q \in \pi^{-1}(p) \cap \Lambda$. As $Y = \pi^{-1}(Z)$, it contains $\pi^{-1}(p)$ and thus $q \in Y \cap \Lambda$ and $\pi(q) = p$. \square

3.2. Khovanskii bases and the degeneration. Let X be a d -dimensional complex variety and $V \subset \mathbb{C}(X)$ a finite-dimensional complex vector subspace. Suppose that the image of V^\times under ν generates \mathbb{Z}^d and V has a finite Khovanskii basis \mathcal{B} such that $\mathcal{B} \not\subset V_s$. For each $k \in \mathbb{N}$, let $W_k s^k := \text{Span}(\mathcal{B} \cap V^k s^k) \subset V^k s^k$ be the span of the elements of \mathcal{B} of homogeneous degree k . We define $W := \bigoplus_{k \geq 1} W_k$ where $V = W_1$ and construct the corresponding weighted projective space as in Section 3.1. Anderson's toric degeneration [2] naturally embeds into $\mathbb{P}_a(W^*)$. The weighted projective space $\mathbb{P}_a(W^*)$ is needed (rather than $\mathbb{P}(V^*)$) to accommodate the generators of $\text{gr } R(V) \simeq \mathbb{C}[S(V, \nu)]$ which are not in V , as these are needed for embedding the toric fiber.

We introduce coordinates by ordering the elements of $\mathcal{B} = \{b_0 s^{a_0}, \dots, b_{n+m} s^{a_{n+m}}\}$ where $a_0 = \dots = a_n = 1$, and for $n < j \leq n+m$, $a_j > 1$. Necessarily, $\{b_0, \dots, b_n\} \subset V$, since Vs generates $R(V)$. Then, for each $n < j \leq n+m$, there is a homogeneous polynomial $h_j \in \mathbb{C}[z_0, \dots, z_n]$ of degree a_j such that $b_j = h_j(b_0, \dots, b_n)$.

Using the Khovanskii basis \mathcal{B} , the Kodaira map to $\mathbb{P}_a(W^*)$ from $X = \text{Proj}(R(V))$ is $\varphi_{\mathcal{B}}: z \mapsto [b_0(z), \dots, b_{n+m}(z)]$. Since, for $n < j \leq n+m$, $b_j = h_j(b_0, \dots, b_n)$, the image of $\varphi_{\mathcal{B}}$ is a graph over the the image of φ_V in $\mathbb{P}(V^*) \subset \mathbb{P}_a(W^*)$.

The constructions of $I_{\mathcal{B}}$, \mathcal{A} , w , $\text{in}_{\mathcal{A}}(I_{\mathcal{B}})$, and \mathcal{G}_t from Section 2.2 all carry over to this general case since all of these ideals are a -homogeneous. Collectively, they embed Anderson's toric degeneration into the weighted projective space $\mathbb{P}_a(W^*)$. The special fiber \mathcal{X}_0 is a toric variety with ideal \mathcal{G}_0 and toric Kodaira map $\varphi_{p, \mathcal{A}}$, where $p \in \mathcal{X}_0 \cap \mathbb{T}_a$ (as before, the torus $\mathbb{T}_a \subset \mathbb{P}_a(W^*)$ consists of those points with nonzero coordinates).

We pull back the embedded toric degeneration $\mathcal{X} \subset \mathbb{P}_a(W^*) \times \mathbb{C}_t$ along π to obtain a flat family $\mathcal{Y} \subset \mathbb{P}^{n+m} \times \mathbb{C}_t$ that is a toric degeneration in the sense of Remarks 5 and 8 as \mathcal{Y} or \mathcal{Y}_0 may not be irreducible. We explain how the equations defining the family \mathcal{Y} may be obtained. Let $\mathbb{C}[y_0, \dots, y_{n+m}]$ be the homogeneous coordinate ring of the projective space \mathbb{P}^{n+m} . The map $\pi: \mathbb{P}^{n+m} \rightarrow \mathbb{P}_a(W^*)$ corresponds to the map $\pi^*: \mathbb{C}[x_0, \dots, x_{n+m}] \rightarrow \mathbb{C}[y_0, \dots, y_{n+m}]$ induced by $x_i \mapsto y_i^{a_i}$. Let

$$(8) \quad \mathcal{F}_t := \{\pi^*(g_t) : g_t \in \mathcal{G}_t\}$$

be the pull back of the equations \mathcal{G}_t for the embedded degeneration $\mathcal{X} \rightarrow \mathbb{C}_t$. Then $\mathcal{Y} = \mathcal{V}(\mathcal{F}_t) \subset \mathbb{P}^{n+m} \times \mathbb{C}_t$. This lifted family $\mathcal{Y} \rightarrow \mathbb{C}_t$ is the fiberwise pull back of Anderson's toric degeneration $\mathcal{X} \rightarrow \mathbb{C}_t$ along the finite map π , where G_a acts on \mathcal{Y} fiberwise.

3.3. Weighted Khovanskii homotopy. We explain how to use the embedded degeneration \mathcal{X} in $\mathbb{P}_a(W^*)$ to compute the linear section $\varphi_V(X) \cap L$. Since $\mathcal{X}_1 = \varphi_{\mathcal{B}}(X)$, it is natural to propose to compute $\mathcal{X}_1 \cap pr_a^{-1}(L)$ using an adaptation of the linear section homotopy to weighted projective space by following points of $\mathcal{X}_0 \cap pr_a^{-1}(L)$ along Anderson's degeneration. Unfortunately, $pr_a^{-1}(L)$ is not sufficiently general for the toric special fiber in Anderson's degeneration.

To avoid this problem, we pull back the toric degeneration \mathcal{X} along π to \mathcal{Y} and use a linear section homotopy to compute the linear section $\mathcal{Y}_1 \cap pr^{-1}(L)$. Since $pr^{-1}(L)$ is not a general linear subspace, we instead choose a general linear subspace $\Lambda \subset \mathbb{P}^{n+m}$ of codimension d . Next, we use Algorithm 3 to compute $\mathcal{Y}_1 \cap \Lambda$, which is a witness set for \mathcal{Y}_1 . Then, we use the witness set homotopy (Algorithm 2) to compute $\mathcal{Y}_1 \cap pr^{-1}(L)$. Finally, $\varphi_V(X) \cap L$ is computed as $pr(\mathcal{Y}_1 \cap pr^{-1}(L))$.

Algorithm 20 (Weighted Khovanskii homotopy algorithm).

Input: A finite-dimensional subspace $V \subset \mathbb{C}(X)$ for a variety $X = \text{Proj}(R(V))$ of dimension d , finite Khovanskii basis $\mathcal{B} \not\subset Vs$ for V , and a general linear subspace $L \subset \mathbb{P}(V^*)$ of codimension d .

Output: Points in the linear section $\varphi_V(X) \cap L$ in the projective space $\mathbb{P}(V^*)$.

Do:

- (i) Follow Steps (i) through (v) of Algorithm 14, *mutatis mutandis*: The ideal $I_{\mathcal{B}}$ is the kernel of the map $\mathbb{C}[x_0, \dots, x_{n+m}] \rightarrow R(V)$ where $x_i \mapsto b_i s^{a_i}$.
- (ii) Pull back the family \mathcal{X} along π to compute the family \mathcal{Y} defined by \mathcal{F}_t , see Definition (8).
- (iii) Compute Kodaira maps for each irreducible component of \mathcal{Y}_0 .
- (iv) Let $\Lambda \subset \mathbb{P}^{n+m}$ be a general linear subspace of codimension d and use Algorithm 3 to compute $\mathcal{Y}_1 \cap \Lambda$.
- (v) Use Algorithm 2 to compute $\mathcal{Y}_1 \cap pr^{-1}(L)$.
- (vi) Return $\varphi_V(X) \cap L = pr(\mathcal{Y}_1 \cap pr^{-1}(L))$.

Remark 21. We discuss Step (iii) of Algorithm 20. As \mathcal{Y}_0 may consist of several components and $\mathcal{Y}_0 = \pi^{-1}(\mathcal{X}_0)$, the group G_a acts transitively on these components. Moreover, each component is a projective toric variety $X_{q,\mathcal{C}}$ for a point $q \in \mathcal{Y}_0 \cap \mathbb{T}$ and all have the same set of exponents, which are the columns of matrix \mathcal{C} . We explain how to compute both q and \mathcal{C} .

From Step (v) of Algorithm 14, we have a toric Kodaira map $\varphi_{p,\mathcal{A}}: (\mathbb{C}^\times)^d \rightarrow \mathbb{P}_a^{n+m}$ such that $\mathcal{X}_0 = X_{p,\mathcal{A}}$. The image includes the point $p \in \mathcal{X}_0 \cap \mathbb{T}_a$. Points $q \in \pi^{-1}(p)$ are obtained by taking all a_j -th roots of the coordinate p_j of p , for all j ,

$$\pi^{-1}(p) = \{q \in \mathbb{P}^{n+m} : q_j^{a_j} = p_j \text{ for } j = 0, \dots, n+m+1\}.$$

It remains to determine the exponents \mathcal{C} for $\pi^{-1}(X_{\mathcal{A}})$. As in Remark 13, we have a basis $u_1, \dots, u_{n+m-d} \in \mathbb{Z}^{n+m+1}$ for $\ker(\mathcal{A})$. These vectors give equations $x^{u_i} = 1$ for $X_{\mathcal{A}} \cap \mathbb{T}_a$. Applying π^* substitutes $y_j^{a_j}$ for x_j and gives equations for $\pi^{-1}(X_{\mathcal{A}}) \cap \mathbb{T}$,

$$(9) \quad y^{v_i} = 1 \quad i = 1, \dots, n+m-d,$$

where v_i is obtained from u_i by multiplying its j th coordinate by a_j .

The System (9) for $\pi^{-1}(X_{\mathcal{A}}) \cap \mathbb{T}$ leads to equations for $Y := \pi^{-1}(X_{\mathcal{A}})$, which form a lattice ideal [12, Section 2] for the lattice K spanned by $\{v_1, \dots, v_{n+m-d}\}$. That is, $Y = \mathcal{V}\langle y^\alpha - y^\beta \mid \alpha - \beta \in K \rangle$.

Let $(\gamma_1, \dots, \gamma_d, \mathbb{1})$ be a basis for the annihilator of K in \mathbb{Z}^{n+m+1} . Suppose that \mathcal{C} is the $d \times (n+m+1)$ matrix whose rows are $\gamma_1, \dots, \gamma_d$. Then $X_{\mathcal{C}}$ is the component of Y containing the identity $\mathbb{1} \in \mathbb{T}$. We remark that \mathcal{C} may be computed from v_1, \dots, v_{n+m-d} using the Hermite normal form. All Kodaira maps needed in Step (iii) of Algorithm 20 can then be computed by translations. \diamond

Remark 22. The number of components of \mathcal{Y} or of \mathcal{Y}_0 impacts the number of Kodaira maps needed in Step (iii) of Algorithm 20. Reductions in the number of Kodaira maps may significantly improve the efficiency of the algorithm.

When \mathcal{Y} is known to be reducible, this structure may be exploited, as Statement (iii) of Lemma 18 implies that it is enough to apply Algorithm 20 to a single component of \mathcal{Y} . In particular, the map π sends the curves in a linear section of one component onto $\mathcal{X} \cap pr^{-1}(L)$.

When \mathcal{Y}_0 has fewer than $|G_a|$ components, then there are redundant Kodaira maps constructed in Remark 21. More precisely, the number of redundant maps is the number of points of $\pi^{-1}(p)$ in a component of \mathcal{Y}_0 . We provide details on computing non-redundant

Kodaira maps, assuming, as in Remark 21, that $\mathbb{1} \in \mathcal{Y}_0$. The general case is obtained by translation. Let

$$\text{sat}(K) := \{w \in \mathbb{Z}^{n+m+1} : rw \in K \text{ for some } 0 \neq r \in \mathbb{Z}\}$$

be the saturation of K and $M = \ker(\mathbb{1}) \subset \mathbb{Z}^{n+m+1}$. We note that $\text{sat}(K) \subset M$. We identify \mathbb{T} with $\text{Hom}(M, \mathbb{C}^\times)$ so that $\mathcal{Y}_0 \cap \mathbb{T} = \text{Hom}(M/K, \mathbb{C}^\times)$, as these are the points satisfying System (9). The component of $\mathcal{Y}_0 \cap \mathbb{T}$ containing the identity $\mathbb{1} \in \mathbb{T}$ is $\text{Hom}(M/\text{sat}(K), \mathbb{C}^\times)$, and the group of components of $\mathcal{Y}_0 \cap \mathbb{T}$ is $\text{Hom}(\text{sat}(K)/K, \mathbb{C}^\times)$. Hence, the elements of $\text{Hom}(\text{sat}(K)/K, \mathbb{C}^\times)$ generate Kodaira maps to distinct components of \mathcal{Y}_0 . \diamond

Proof of correctness of Algorithm 20. We need only show that the tracked paths provide enough points to compute $\varphi_V(X) \cap L$. By Statement (iii) of Lemma 18, for each t , the map $\pi : \mathcal{Y}_t \cap \Lambda \rightarrow \mathcal{X}_t \cap \pi(\Lambda)$ is a bijection.

The polyhedral homotopy correctly computes the points of $\mathcal{Y}_0 \cap \Lambda$. By Theorem 4, Algorithm 3 correctly computes the points of $\mathcal{Y}_1 \cap \Lambda$. Since the solution paths of the homotopy $\mathcal{X} \cap \pi(\Lambda)$ are disjoint, the solution paths of $\mathcal{Y} \cap \Lambda$ lie above paths of $\mathcal{X} \cap \pi(\Lambda)$. In fact, by Statement (iii) of Lemma 18, π is a bijection between these sets of paths. Therefore, there is a bijection between the ends of the homotopy paths of $\mathcal{Y} \cap \Lambda$ and points in $\mathcal{X}_1 \cap \pi(\Lambda)$. The correctness of the final computation then follows from the correctness of Algorithm 2. \square

Example 23. Let V be the space of cubic polynomials in $\mathbb{C}[x, y]$ which vanish at the points $(4, 4), (-3, -1), (-1, -1)$ and $(3, 3)$. (This example is related to the example of [10, Section 5.1], which considers quartics vanishing at these points.) Then V is six-dimensional with a basis:

$$\begin{aligned} \{b_0, \dots, b_5\} = \{ & \underline{xy} - y^2 + x - y, \underline{x^2} - y^2 + 4x - 4y, \underline{y^3} - 6y^2 + 5y + 12, \\ & \underline{xy^2} - 6y^2 - x + 6y + 12, \underline{x^2y} - 6y^2 - 4x + 9y + 12, \underline{x^3} - 6y^2 - 13x + 18y + 12\}. \end{aligned}$$

We compute a general linear section of $X = \text{Proj}(R(V))$ in $\mathbb{P}(V^*) = \mathbb{P}^5$ with Algorithm 20. Let \succeq be the order on \mathbb{Z}^2 where $(a, b) \succeq (c, d)$ if $a + b < c + d$ or else $a + b = c + d$ and $a < c$. Define a valuation ν on $\mathbb{C}(X) = \mathbb{C}(x, y)$ as follows: for $f \in \mathbb{C}[x, y]$, $\nu(f) = (a, b)$ where (a, b) is the \succeq -minimal exponent of a term of f . This order and valuation ν are compatible with the **grevlex** order \leq on $\mathbb{C}[x, y]$ with $x > y$ in that $(a, b) \succeq (c, d)$ if and only if $x^a y^b \leq x^c y^d$. Using the subduction algorithm, as implemented in the unreleased Macaulay2 package **SubalgebraBases** [25] applied to $\{b_0s, \dots, b_5s\}$, we obtain a Khovanskii basis $\mathcal{B} = \{b_0s, \dots, b_5s, b_6s^2, b_7s^3\}$ with two additional generators, where

$$\begin{aligned} b_6 &:= \underline{xy^3} - y^4 + 10x^2y - 26xy^2 + 16y^3 + 10x^2 - 15xy + 5y^2 + 12x - 12y, \text{ and} \\ b_7 &:= \underline{10x^4y} - 49x^3y^2 + 89x^2y^3 - 71xy^4 + 21y^5 + 10x^4 - 18x^3y - 18x^2y^2 \\ &\quad + 50xy^3 - 24y^4 + 31x^3 - 83x^2y + 73xy^2 - 21y^3 + 24x^2 - 48xy + 24y^2. \end{aligned}$$

The corresponding matrix of valuations is

$$\mathcal{A} = \nu(\mathcal{B}) = \begin{pmatrix} 1 & 2 & 0 & 1 & 2 & 3 & 1 & 4 \\ 1 & 0 & 3 & 2 & 1 & 0 & 3 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 2 & 3 \end{pmatrix}.$$

The Newton-Okounkov body, as displayed in Figure 2, is obtained by intersecting the cone generated by the columns of \mathcal{A} with the hyperplane where the third coordinate is 1. The vertices $(1/2, 3/2)$ and $(4/3, 1/3)$ come from the initial (underlined) terms of b_6 and b_7 . While they are not integers, the Newton-Okounkov body has normalized volume 5, which is the degree of X . We may interpret this volume as follows: Two cubics drawn from V meet in $5 = 3^2 - 4$ points outside the base locus $\mathcal{V}(V) = \{(4, 4), (-3, -1), (-1, -1), (3, 3)\}$.

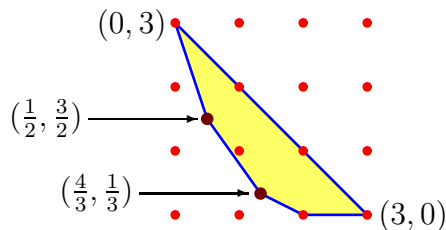


FIGURE 2. Newton Okounkov body for the space of cubic polynomials vanishing at $(4, 4)$, $(-3, -1)$, $(-1, -1)$ and $(3, 3)$.

The weight $w = (-6, -5, 0)$ is compatible with the **grevlex** order \leq on \mathcal{B} in that for $b \in \mathcal{B}$, the \leq -leading term has lowest w -weight, so that $\text{lt}_{\leq} b = b_w$. Choosing a term order on $\mathbb{C}[x_0, \dots, x_7]$ that is compatible with $w\mathcal{A}$, we use **Macaulay2** to compute a Gröbner basis \mathcal{G} for $I_{\mathcal{B}}$. This basis consists of 17 polynomials which are $a := (1, 1, 1, 1, 1, 1, 2, 3)$ -homogeneous. Let \mathbb{C}^\times act on a -homogeneous polynomials using $w\mathcal{A}$ in place of w in Formula (4). Then we compute $\mathcal{G}_t := \{g_t : g \in \mathcal{G}\}$ as in Formula (5), which defines a flat family $\mathcal{X} \subset \mathbb{P}_a^7 \times \mathbb{C}_t$ with toric special fiber \mathcal{X}_0 . This family pulls back along $\pi : \mathbb{P}^7 \rightarrow \mathbb{P}_a^7$ to a family $\mathcal{Y} \subset \mathbb{P}^7 \times \mathbb{C}_t$. The pull back \mathcal{Y}_0 of \mathcal{X}_0 under π is a toric variety as it is irreducible. From Remarks 13 and 21, a Kodaira map for \mathcal{Y}_0 is

$$\begin{aligned} \varphi_{p, \mathcal{A}} : (\mathbb{C}^*)^2 &\longrightarrow \mathbb{P}^7 \\ z &\longmapsto [z_1^6 z_2^3, z_1^4 z_2^3, z_1^6, z_1^4, z_1^2, 1, z_1^7 z_2^3, \sqrt[3]{10} z_1^6 z_2^4]. \end{aligned}$$

The polyhedral homotopy finds 30 points in $\mathcal{Y}_0 \cap \Lambda$. An application of the toric two-step algorithm (Algorithm 3) tracks these points to $\mathcal{Y}_1 \cap \Lambda$ with no paths diverging. Then, the witness set homotopy algorithm (Algorithm 2) moves Λ to $pr^{-1}(L)$ and finds the points of $\mathcal{Y}_1 \cap pr^{-1}(L)$. These 30 points project under $\pi : \mathbb{P}^7 \rightarrow \mathbb{P}_a^7$ to five points in $\mathcal{X}_1 \cap pr^{-1}(L)$. Finally, applying the map $pr : \mathbb{P}^7 \rightarrow \mathbb{P}^5 = \mathbb{P}(V^*)$ gives all five points in $\varphi_V(X) \cap L$. \diamond

4. PRACTICAL CONSIDERATIONS

We discuss how to compute a finite Khovanskii basis as well as options for tracking over-terminated homotopy systems.

4.1. Computing a Khovanskii basis. Whether or not a given vector space V of functions has a finite Khovanskii basis is generally not known and depends on the choice of valuation. Given V and a valuation ν , the subduction algorithm [20, Algorithm 2.18] terminates and returns a finite Khovanskii basis, when one exists. We only know of implementations when V is a space of polynomials and ν is induced by a term order [13, 25]. These also compute a SAGBI basis [16, 24].

4.2. Homotopy continuation for overdetermined systems. Algorithms 14 and 20 generate a homotopy (\mathcal{G}_t, L) from a Gröbner basis \mathcal{G} defining \mathcal{X} . This is not typically square in that it has more equations than variables. As most implementations of homotopy continuation, including the user-defined homotopy in *Bertini*, require square systems, we need to choose a square subsystem for tracking from $t = 0$ to $t = 1$.

Typically, a square subsystem is obtained by taking linear combinations of elements in a given system. There is an alternative for equations \mathcal{G}_t from a toric degeneration. Let \mathcal{A} be the matrix of exponents defining the Kodaira map for the toric special fiber, \mathcal{X}_0 . The intersection $\mathcal{X}_0 \cap \mathbb{T}$ with the dense torus of \mathbb{P}^n is the complete intersection defined by binomials $x^{u_i} - c_i x^{v_i}$ for $i = 1, \dots, n-d$ such that $\{u_i - v_i : i = 1, \dots, n-d\}$ form a basis for $\ker(\mathcal{A})$. The points of $\mathcal{X}_0 \cap L$ are smooth isolated solutions to the square system given by these binomials and the linear forms defining L . If we choose $\mathcal{F}_t \subset \mathcal{G}_t$ to consist of $n-d$ elements whose leading binomials are $x^{u_i} - c_i x^{v_i}$, then (\mathcal{F}_t, L) is a square subsystem of (\mathcal{G}_t, L) which defines curves containing $\mathcal{X}_0 \cap L$ and, therefore, is sufficient for homotopy continuation.

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