



Contents lists available at ScienceDirect

European Journal of Combinatorics

journal homepage: www.elsevier.com/locate/ejc

Rainbow Turán number of clique subdivisions

Tao Jiang^{a,1}, Abhishek Methuku^{b,2}, Liana Yepremyan^{c,3}^a Department of Mathematics, Miami University, Oxford, OH 45056, US^b School of Mathematics, University of Birmingham, Birmingham, B15 2TT, UK^c Department of Mathematics, Emory University, Atlanta, GA, 30322, US

ARTICLE INFO

Article history:

Received 31 July 2021

Accepted 8 December 2022

Available online 3 January 2023

ABSTRACT

We show that for any integer $t \geq 2$, every properly edge-coloured graph on n vertices with more than $n^{1+o(1)}$ edges contains a rainbow subdivision of K_t . Note that this bound on the number of edges is sharp up to the $o(1)$ error term. This is a rainbow analogue of some classical results on clique subdivisions and extends some results on rainbow Turán numbers. Our method relies on the framework introduced by Sudakov and Tomon (2022) which we adapt to find robust expanders in the coloured setting.

© 2023 Elsevier Ltd. All rights reserved.

1. Introduction

For a graph H , the Turán number $\text{ex}(n, H)$ is the maximum number of edges in an n -vertex graph without a copy of H . Keevash, Mubayi, Sudakov and Verstraëte [6] introduced a rainbow variant of the Turán problem. In an edge-coloured graph, we say that a subgraph is *rainbow* if no two of its edges have the same colour. The rainbow Turán number $\text{ex}^*(n, H)$ is the maximum number of edges in a properly edge-coloured graph on n vertices which does not contain a rainbow copy of H . Clearly, $\text{ex}^*(n, H) \geq \text{ex}(n, H)$ for every n and H . Among other things, Keevash, Mubayi, Sudakov and Verstraëte [6] proved that if H is non-bipartite then $\text{ex}^*(n, H) = (1 + o(1))\text{ex}(n, H)$ where the asymptotic value of $\text{ex}(n, H)$ is known by celebrated Erdős–Stone–Simonovits theorem. However, if H is bipartite much less is known. Keevash, Mubayi, Sudakov and Verstraëte [6] showed

E-mail addresses: jiangt@miamioh.edu (T. Jiang), abhishek.methuku@math.ethz.ch (A. Methuku), lyeprem@emory.edu (L. Yepremyan).

¹ Research supported by National Science Foundation grant DMS-1855542.

² Research supported by the EPSRC, grant no. EP/S00100X/1.

³ Research supported by Marie Skłodowska Curie Global Fellowship, H2020-MSCA-IF-2018:846304.

<https://doi.org/10.1016/j.ejc.2022.103675>

0195-6698/© 2023 Elsevier Ltd. All rights reserved.

that $\text{ex}^*(n, C_{2k}) = \Omega(n^{1+1/k})$ for all $k \geq 2$, and conjectured that $\text{ex}^*(n, C_{2k}) = \Theta(n^{1+1/k})$. They verified this conjecture for $k \in \{2, 3\}$ and for general k , Das, Lee and Sudakov showed that for every $k \geq 2$, $\text{ex}^*(n, C_{2k}) = O(n^{\frac{(1+\varepsilon_k)\log k}{k}})$ where $\varepsilon_k \rightarrow \infty$ as $k \rightarrow \infty$. Very recently, Janzer [3] proved the conjecture by showing that $\text{ex}^*(n, C_{2k}) = O(n^{1+1/k})$.

It is well-known that a graph on n vertices without any cycle has at most $n - 1$ edges. It is then natural to ask how many edges a properly edge-coloured n -vertex graph can have if it does not contain any rainbow cycle. Keevash, Mubayi, Sudakov and Verstraëte [6] showed that there are graphs with $\Omega(n \log n)$ edges that can be properly coloured with no rainbow cycle by colouring an m -dimensional cube Q as follows. The vertices of Q are subsets of $\{1, 2, \dots, m\}$ and for any $S \subseteq \{1, 2, \dots, m\}$ and any $i \in S$, there is an edge between S and $S \setminus \{i\}$ of colour i . To see that Q has no rainbow cycle, consider any cycle C in Q and an edge e from A to $A \setminus \{i\}$ of colour i on C . Observe that there is a subpath of C connecting $A \setminus \{i\}$ back to A which does not contain e but contains at least one other edge of colour i , as A contains the element i . It follows that C is not a rainbow cycle. Moreover, since $2^m = n$, Q has $\frac{1}{2}mn = \frac{1}{2}n \log n$ edges, as desired. On the other hand, Keevash, Mubayi, Sudakov and Verstraëte showed an upper bound of $O(n^{4/3})$. Das, Lee and Sudakov [1] improved this bound as follows.

Theorem 1.1 (Das, Lee, Sudakov [1]). *If $\eta > 0$ and n is sufficiently large, then any properly edge-coloured n -vertex graph with at least $ne^{(\log n)^{\frac{1}{2}+\eta}}$ edges contains a rainbow cycle.*

Recently Janzer [3] improved the bound in Theorem 1.1 by showing that a properly edge-coloured graph on n vertices with average degree at least $c(\log n)^4$ contains a rainbow cycle, for some constant $c > 0$. The main result of our paper extends Theorem 1.1 by establishing an analogous theorem for the family of rainbow subdivisions of the clique as stated below.

Theorem 1.2. *For every integer $t \geq 2$ there exists a constant $c > 0$ such that for every integer $n \geq t$ if G is a properly edge-coloured graph on n vertices with at least $ne^{c\sqrt{\log n}}$ edges then G contains a rainbow subdivision of K_t , where each edge of K_t is replaced with a path of length at most $1300 \log^2 n$.*

Note that the problem we address here is a “rainbow” instance of a well studied class of problems that explore edge density conditions forcing the existence of a minor/subdivision of small size. A well-known result of Mader [10] states that for every integer t , there is a constant $c(t)$ such that every graph with $c(t)n$ edges contains a K_t -minor. Estimates on $c(t)$ have been improved in [8, 15, 16]. Fiorini, Joret, Theis and Wood [2] asked the following natural question: how many edges suffice to guarantee that a graph contains not only a K_t -minor, but one which has few vertices? Addressing this question, Shapira and Sudakov [13] showed that for every $\varepsilon > 0$ and $t \geq 3$, there exists n_0 such that every graph on $n \geq n_0$ vertices and $(c(t) + \varepsilon)n$ edges contains a K_t -minor of order $O_{\varepsilon,t}(\log n \log \log n)$. Montgomery [12] subsequently improved this bound on the order of the K_t -minor to the optimal bound $O_{\varepsilon,t}(\log n)$. In confirming a conjecture of Erdős, it was shown by Kostochka and Pyber [9] that any graph with $4t^2 n^{1+\varepsilon}$ edges contains a K_t -subdivision of order $O(t^2 \log t / \varepsilon)$. The bound on the order of the K_t -subdivision was improved to the optimal $O(t^2 / \varepsilon)$ by Jiang [4]. Theorem 1.2 is a rainbow version of these results. Our proof establishes an interesting connection between rainbow Turán problems and expanders, building on the method of Sudakov and Tomon used in [14].

An overview of the proof

Our method builds on the method used by Sudakov and Tomon in [14] together with some new ideas. We incorporate the minimality notion commonly used in the study of graph minors and adapt the notion of “expander” conveniently to our setting (see Definition 2.3). We show that most of the edges of a sufficiently dense graph can be covered by edge-disjoint expanders. In a properly edge-coloured expander, from any given vertex v , we can reach almost all of the other vertices by a rainbow path of poly-logarithmic length avoiding a given set of vertices and colours. Additionally, the notion of minimality ensures that the set of these “reachable” vertices induces

most of the edges in the expander (see [Lemma 2.7](#)). This additional feature of the expander is used to show the existence of a common large intersection of reachable vertices for a pair of vertices x and y in a general graph (not necessarily an expander). Eventually we are either able to join any pair of vertices by a rainbow path of poly-logarithmic length avoiding a bounded set of colours and vertices (thus allowing us to build a copy of the desired rainbow K_t -subdivision) or find a much denser subgraph (see [Lemma 3.2](#)). We then complete the proof via a density increment argument as in [\[14\]](#).

Notation and standard tools

Let G be a graph, we denote its vertex set by $V(G)$ and the edge set by $E(G)$, and let $v(G) = |V(G)|$, $e(G) = |E(G)|$. For any subset S of $V(G)$ let $e(S) = e(G[S])$. Similarly, for two disjoint subsets S and T of $V(G)$ write $e(S, T) = e(G[S, T])$. Let $d(G)$ be the average degree of G . For a properly coloured graph F , let $C(F)$ denote the set of colours used on the edges of F .

Lemma 1.3 (Chernoff bounds, [\[11\]](#)). Given a binomially distributed variable $X \in \text{Bin}(n, p)$ for all $0 < a \leq 3/2$ we have

$$\mathbb{P}[|X - \mathbb{E}[X]| \geq a \mathbb{E}[X]] \leq 2e^{-\frac{a^2}{3} \mathbb{E}[X]}.$$

2. Expanders

Definition 2.1. We say that G is d -minimal if $d(G) \geq d$, but $d(F) < d$ for every proper subgraph $F \subseteq G$.

Note that for any graph G with $d(G) \geq d$, there exists a smallest subgraph H that satisfies $d(H) \geq d$. Such a subgraph H is a d -minimal subgraph of G . Note also that, by definition, if G is d -minimal, then it is also $d(G)$ -minimal. We will utilize the minimality notion in two ways. One is based on the definition, which, roughly speaking, says that any subset of vertices spans at most the expected number of edges. The other way we utilize minimality is that whenever we delete a set of vertices we must lose at least the expected number of edges, as stated in the following lemma.

Lemma 2.2. If G is d -minimal, then for every set $S \subseteq V(G)$, the number of edges having at least one endpoint in S is at least $d|S|/2$, that is

$$e(S) + e(S, S^c) \geq \frac{d|S|}{2}.$$

In particular, $\delta(G) \geq \frac{d}{2}$.

Proof. Let $S \subseteq V(G)$. Suppose for a contradiction that $e(S) + e(S, S^c) < \frac{d|S|}{2}$. Then

$$e(S^c) > \frac{d|V|}{2} - \frac{d|S|}{2} \geq \frac{d(|V| - |S|)}{2},$$

a contradiction to the d -minimality of G . \square

Definition 2.3. Given $\varepsilon, \lambda \in (0, 1)$ and $d \geq 1$, an n -vertex graph G is called a $(d, \lambda, \varepsilon)$ -expander if G is d -minimal, and for every subset $S \subseteq V(G)$ of size at most $(1 - \varepsilon)n$, we have $d(S) \leq (1 - \lambda)d$.

Note that by the remarks after [Definition 2.1](#), if G is a $(d, \lambda, \varepsilon)$ -expander then it is also a $(d(G), \lambda, \varepsilon)$ -expander. Hence, we will often say let G be a $(d, \lambda, \varepsilon)$ -expander with $d = d(G)$.

Lemma 2.4. Let $n, d \geq 1$ and let $\lambda, \varepsilon \in (0, 1)$. Let G be a $(d, \lambda, \varepsilon)$ -expander on n vertices. Then every $S \subseteq V(G)$ with $|S| \leq (1 - \varepsilon)n$ satisfies $e(S, S^c) \geq \frac{\lambda d}{2} |S|$.

Proof. Let $S \subseteq V(G)$ and let $|S| \leq (1 - \varepsilon)n$. Since G is a $(d, \lambda, \varepsilon)$ -expander, it is d -minimal by definition. So by Lemma 2.2, $e(S) + e(S, S^c) \geq \frac{d|S|}{2}$ and $e(S) = \frac{d(S)}{2}|S| \leq \frac{(1-\lambda)d}{2}|S|$. It follows that $e(S, S^c) \geq \frac{\lambda d}{2}|S|$, as claimed. \square

Lemma 2.5. *Let n be a positive integer. Let d, ε, λ be positive reals where $d \geq 1, \varepsilon \in (0, 1)$ and $\lambda \leq \frac{\varepsilon}{2 \log n}$. Let G be a graph on n vertices with average degree d . Then G contains a $(d', \lambda, \varepsilon)$ -expander, with $d' \geq \frac{d}{2}$.*

Proof. Let $G_0 = G$. We run the following process for $i \geq 0$:

- (a) Let $d_i = d(G_i)$. Let H_i be a d_i -minimal subgraph of G_i .
- (b) If H_i is a $(d_i, \lambda, \varepsilon)$ -expander, stop. Otherwise, there exists $S_{i+1} \subseteq V(H_i)$, with $|S_{i+1}| \leq (1 - \varepsilon)|V(H_i)|$, such that $d(S_{i+1}) \geq (1 - \lambda)d_i$. Take $G_{i+1} = G[S_{i+1}]$.

Suppose that the process stopped right after H_ℓ was defined; so H_ℓ is a $(d_\ell, \lambda, \varepsilon)$ -expander. Since $|V(G_{i+1})| \leq (1 - \varepsilon)|V(G_i)|$ for every $i \geq 1$ and $|V(G_\ell)| \geq 1$, it follows that $(1 - \varepsilon)^\ell n \geq 1$, implying that $\ell \leq \frac{-\log n}{\log(1-\varepsilon)} \leq \frac{\log n}{\varepsilon}$. Since $d_{i+1} \geq (1 - \lambda)d_i$ for every $i \geq 1$ and $d_0 \geq d$, we have $d_\ell \geq (1 - \lambda)^\ell d \geq (1 - \lambda \ell)d \geq (1 - \frac{\lambda \log n}{\varepsilon})d \geq \frac{d}{2}$. \square

The next corollary says that we can iterate the previous lemma to cover almost all the edges of any graph by expanders.

Corollary 2.6. *Let n be a positive integer. Let d, ε, λ be positive reals where $d \geq 1, \varepsilon \in (0, 1)$ and $\lambda \leq \frac{\varepsilon}{2 \log n}$. Let G be a graph on n vertices with average degree d . Then G contains edge-disjoint subgraphs G_1, G_2, \dots, G_k such that for each $i \in [k]$, G_i is a $(d(G_i), \lambda, \varepsilon)$ -expander, where $d(G_i) \geq \frac{\varepsilon d}{2}$, and $e(\cup_{i=1}^k G_i) \geq (1 - \varepsilon)e(G)$.*

Proof. Let $\mathcal{F} = \{G_1, G_2, \dots, G_k\}$ be a maximal collection of edge-disjoint subgraphs such that for each $i \in [k]$, G_i is a $(d(G_i), \lambda, \varepsilon)$ -expander, where $d(G_i) \geq \frac{\varepsilon d}{2}$. Let G' be the subgraph of G consisting of edges not covered by G_1, G_2, \dots, G_k . If $e(G') \geq \varepsilon e(G)$ then $d(G') \geq \varepsilon d$. By Lemma 2.5, G' contains a $(d_{k+1}, \lambda, \varepsilon)$ -expander G_{k+1} with $d(G_{k+1}) = d_{k+1} \geq \frac{\varepsilon d}{2}$, contradicting the maximality of \mathcal{F} . It follows that $e(G') \leq \varepsilon e(G)$, as required. \square

Given an edge-coloured graph and some path P in it, we say P avoids a colour c if it does not contain any edge of colour c . More generally, we say P avoids a set of colours C if P avoids each colour $c \in C$. Analogously, we say P avoids a vertex v if it does not contain v , and it avoids a set of vertices F if it avoids each vertex $v \in F$.

Lemma 2.7. *Let $n, \ell, M \geq 1$ be integers and let d, ε, λ be reals where $d \geq 1$ and $\varepsilon, \lambda \in (0, 1)$. Suppose that $\ell = \lceil \frac{4 \log n}{\lambda} \rceil$ and $d \geq \frac{4(\ell+2M)}{\lambda}$. Suppose that G is a properly edge-coloured $(\lambda, d, \varepsilon)$ -expander on n vertices, let $x \in V(G)$ and let F be a set of at most M vertices and C be a set of at most M colours. Then there is a set of vertices U such that each vertex $u \in U$ can be reached from x by a rainbow path of length at most $\ell + 1$ avoiding C and F , and furthermore,*

- (i) $|U| \geq (1 - \varepsilon)n$, and
- (ii) $e(G[U]) \geq (1 - \lambda/2 - \varepsilon)e(G)$.

Proof. By the remarks after Definition 2.3, we may assume that $d(G) = d$. For each $i \geq 0$, let U_i be the set of vertices reachable from x by a rainbow path of length at most i avoiding C and F . For each $u \in U_i$, fix a rainbow path $P(u)$ from x to u of length at most i that avoids C and F .

Claim 1. *For any $i \leq \ell$, $e(U_{i+1}) \geq (1 - \frac{\lambda}{2}) \frac{d|U_i|}{2}$.*

Proof. Write $E_i = E(U_i, U_i^c)$, and let E_i^* be the set of edges uw , with $u \in U_i$ and $w \in U_i^c$, such that either the colour of uw appears on $P(u)$ or in C or $w \in F$. Then

$$|E_i^*| \leq (i + 2M)|U_i| \leq (\ell + 2M)|U_i| \leq \frac{\lambda d}{4}|U_i|. \quad (2.1)$$

Consider an edge $uw \in E_i \setminus E_i^*$, with $u \in U_i$ and $w \in U_i^c$. Then, by definition, $P(u) \cup uw$ is a rainbow path from x to w of length at most $i + 1$, avoiding F and C , so $w \in U_{i+1}$, implying that $uw \in E(U_i, U_{i+1} \setminus U_i)$. It follows that $E_i \setminus E_i^* \subseteq E(U_i, U_{i+1} \setminus U_i)$. Hence, for any $i \leq \ell$,

$$e(U_{i+1}) \geq e(U_i) + |E_i \setminus E_i^*| = e(U_i) + e(U_i, U_i^c) - |E_i^*| \geq \frac{d|U_i|}{2} - \frac{\lambda d|U_i|}{4} = \left(1 - \frac{\lambda}{2}\right) \frac{d|U_i|}{2},$$

where we used [Lemma 2.2](#) and (2.1). This proves the claim. \square

Claim 2. Let $i \leq \ell$. If $|U_{i+1}| \leq (1 - \varepsilon)n$, then $|U_{i+1}| \geq (1 + \frac{\lambda}{2})|U_i|$.

Proof. Suppose that $|U_{i+1}| \leq (1 - \varepsilon)n$. Since G is a $(\lambda, d, \varepsilon)$ -expander and $|U_{i+1}| \leq (1 - \varepsilon)n$, by definition, we have $e(U_{i+1}) \leq \frac{(1-\lambda)d}{2}|U_{i+1}|$. Combining this with the lower bound on $e(U_{i+1})$ obtained from [Claim 1](#), we get $|U_{i+1}| \geq \frac{1-\lambda/2}{1-\lambda}|U_i| \geq (1 + \frac{\lambda}{2})|U_i|$. \square

Suppose that $|U_\ell| \leq (1 - \varepsilon)n$. By iterating [Claim 2](#) we get $|U_\ell| \geq (1 + \frac{\lambda}{2})^\ell \geq e^{\frac{\lambda\ell}{4}} \geq e^{\log n} > n$, a contradiction. This implies that $|U_\ell| \geq (1 - \varepsilon)n$. Now, by definition, $U_{\ell+1} \supseteq U_\ell$, hence $|U_{\ell+1}| \geq (1 - \varepsilon)n$ and furthermore, by [Claim 1](#),

$$e(U_{\ell+1}) \geq \left(1 - \frac{\lambda}{2}\right) \frac{d|U_\ell|}{2} \geq (1 - \lambda/2)(1 - \varepsilon)e(G) \geq (1 - \lambda/2 - \varepsilon)e(G),$$

where in the second inequality we used that $|U_\ell| \geq (1 - \varepsilon)n$ and $e(G) = dn/2$. So it follows that $e(U_{\ell+1}) \geq (1 - \lambda/2 - \varepsilon)e(G)$. The lemma holds with $U = U_{\ell+1}$. \square

3. Proof of [Theorem 1.2](#)

Lemma 3.1. Let G be a graph with $\delta(G) \geq \frac{1}{3}d(G)$ and G' be a subgraph of G . For any $\gamma > 0$, if $e(G') \geq (1 - \gamma)e(G)$ then $v(G') \geq (1 - 3\gamma)v(G)$.

Proof. Let $V = V(G)$. Let $S = V \setminus V(G')$. Suppose for contradiction that $|S| > 3\gamma \cdot v(G)$. The number of edges in G that have at least one endpoint in S is at least

$$\frac{\delta(G)|S|}{2} \geq \frac{d(G)}{3} \cdot \frac{3\gamma \cdot v(G)}{2} > \gamma \cdot e(G).$$

This is impossible since the set of edges of G that are incident with S is disjoint from $E(G')$ and $e(G') \geq (1 - \gamma)e(G)$. \square

The following is our main lemma of the paper.

Lemma 3.2. Let $0 < \varepsilon \leq \frac{1}{40}$ be a fixed real. For all sufficiently large integers n the following holds. Let $\lambda = \frac{\varepsilon}{2 \log n}$. Let $\ell = \lceil \frac{4 \log n}{\lambda} \rceil$. Let K, L be positive integers. Let $M \geq L + \frac{12K}{\varepsilon}(\ell + 2)$. Let d be a real that satisfies $d \geq \frac{24(\ell + 2M)}{\varepsilon \lambda}$. Let G be a properly edge-coloured $(d, \lambda, \varepsilon)$ -expander, with $d = d(G)$. Suppose G contains no subgraph on at most $\frac{n}{K}$ vertices with average degree at least $\frac{ed}{6}$. Then for any two vertices $x, y \in V(G)$, any set C of at most L colours and any set F of at most L vertices in $G - \{x, y\}$, there exists a rainbow x, y -path in G of length at most $4\ell + 4$ that avoids F and C .

Proof. Since G is a $(d, \lambda, \varepsilon)$ -expander, $\delta(G) \geq d/2 = d(G)/2$ by [Lemma 2.2](#). Place each colour used in G in group 1 or 2 with probability $1/2$. For $i \in \{1, 2\}$, let G_i denote the spanning subgraph of G that consists of edges whose colour lies in group i . Using Chernoff bounds, we can show that, with positive probability, every $v \in V(G)$ satisfies $d_{G_i}(v) = (\frac{1}{2} \pm o(1))d_G(v)$. In particular, when n is sufficiently large, one can ensure that

Claim 3. For $i \in \{1, 2\}$, we have $d(G_i) \geq \frac{1}{3}d$ and $\delta(G_i) \geq \frac{1}{3}d(G_i)$.

By Corollary 2.6, for $i = 1, 2$, G_i contains edge-disjoint subgraphs $G_{i,1}, \dots, G_{i,s_i}$ such that each $j \in [s_i]$, $G_{i,j}$ is a $(d(G_{i,j}), \lambda, \varepsilon)$ -expander with $d(G_{i,j}) \geq \frac{\varepsilon d(G_i)}{2} \geq \frac{\varepsilon d}{6}$, and $e(\cup_{j=1}^{s_i} G_{i,j}) \geq (1 - \varepsilon)e(G_i)$.

Claim 4. For $i \in \{1, 2\}$, we have $s_i \leq \frac{6K}{\varepsilon}$.

Proof of Claim 4. By our discussion, for each $i \in \{1, 2\}$ and $j \in \{1, \dots, s_i\}$ we have $d(G_{i,j}) \geq \frac{\varepsilon}{6}d$. By the assumption of the lemma, we must have $v(G_{i,j}) > \frac{n}{K}$. Hence $e(G_{i,j}) = \frac{1}{2}v(G_{i,j})d(G_{i,j}) \geq \frac{1}{2} \frac{n}{K} \frac{\varepsilon d}{6} = \frac{\varepsilon e(G)}{6K}$. This implies that $s_i \leq \frac{6K}{\varepsilon}$. \square

Let S_x denote the set of vertices in G which can be reached from x by a rainbow path of length at most $\ell + 1$ avoiding F and C . Since G is a $(d, \lambda, \varepsilon)$ -expander and $d \geq \frac{4(\ell+2M)}{\lambda}$ and $|F|, |C| \leq L \leq M$, by Lemma 2.7, $|S_x| \geq (1 - \varepsilon)n$. Let $J_1 = \{j \in [s_1] : V(G_{1,j}) \cap S_x \neq \emptyset\}$. For each $j \in J_1$, fix a vertex $x_j \in S_x \cap V(G_{1,j})$ and a rainbow path P_j from x to x_j in G whose length is at most $\ell + 1$ avoiding F and C . Let $F_1 := \cup_{j \in J_1} P_j$.

Let S_y to be the set of vertices in G which can be reached from y by a rainbow path of length at most $\ell + 1$ avoiding $F \cup V(F_1)$ and $C \cup C(F_1)$.

Note that since

$$|F \cup V(F_1)|, |C \cup C(F_1)| \leq L + s_1(\ell + 2) \leq L + \frac{6K}{\varepsilon}(\ell + 2) \leq M$$

and $d \geq \frac{4(\ell+2M)}{\lambda}$, by Lemma 2.7, $|S_y| \geq (1 - \varepsilon)n$. Let $J_2 = \{j \in [s_2] : V(G_{2,j}) \cap S_y \neq \emptyset\}$. For each $j \in J_2$, fix a vertex $y_j \in S_y \cap V(G_{2,j})$ and a rainbow path Q_j from y to y_j in G whose length is at most $\ell + 1$ avoiding $F \cup V(F_1)$ and $C \cup C(F_1)$. Let $F_2 = \cup_{j \in J_2} Q_j$. Note that

$$|F \cup V(F_1) \cup V(F_2)|, |C \cup C(F_1) \cup C(F_2)| \leq L + (s_1 + s_2)(\ell + 2) \leq L + \frac{12K}{\varepsilon}(\ell + 2) \leq M. \quad (3.1)$$

For each $j \in J_1$, let $U_{1,j}$ denote the set of vertices in $G_{1,j}$ that are reachable from x_j by a rainbow path in $G_{1,j}$ of length at most $\left\lceil \frac{4 \log v(G_{1,j})}{\lambda} \right\rceil + 1 \leq \ell + 1$ that avoids $F \cup V(F_1) \cup V(F_2)$ and $C \cup C(F_1) \cup C(F_2)$. For each $j \in J_2$, let $U_{2,j}$ denote the set of vertices in $G_{2,j}$ that are reachable from y_j by a rainbow path in $G_{2,j}$ of length at most $\left\lceil \frac{4 \log v(G_{2,j})}{\lambda} \right\rceil + 1 \leq \ell + 1$ that avoids $F \cup V(F_1) \cup V(F_2)$ and $C \cup C(F_1) \cup C(F_2)$. For each $i = 1, 2, j \in J_i$, since $G_{i,j}$ is a $(d(G_{i,j}), \lambda, \varepsilon)$ -expander and $d(G_{i,j}) \geq \frac{\varepsilon d}{6} \geq \frac{4(\ell+2M)}{\lambda}$, by Lemma 2.7, we have

1. $|U_{i,j}| \geq (1 - \varepsilon)v(G_{i,j})$
2. $e(G_{i,j}[U_{i,j}]) \geq (1 - \lambda/2 - \varepsilon)e(G_{i,j}) \geq (1 - 2\varepsilon)e(G_{i,j})$.

For $i = 1, 2$, let $G'_i = \cup_{j \in J_i} G_{i,j}[U_{i,j}]$.

Claim 5. For $i = 1, 2$, we have $e(G'_i) \geq (1 - 6\varepsilon)e(G_i)$ and $v(G'_i) \geq (1 - 15\varepsilon)n$.

Proof. By our discussion above, for each $i \in \{1, 2\}, j \in J_i$, $e(G_{i,j}[U_{i,j}]) \geq (1 - 2\varepsilon)e(G_{i,j})$. Since the $G_{i,j}$'s are pairwise edge-disjoint, we have

$$e(G'_i) = \sum_{j \in J_i} e(G_{i,j}[U_{i,j}]) \geq (1 - 2\varepsilon)e(\cup_{j \in J_i} G_{i,j}). \quad (3.2)$$

By the definition of J_1 and J_2 , $V(\cup_{j \notin J_1} G_{1,j}) \subseteq V(G) \setminus S_x$ and $V(\cup_{j \notin J_2} G_{2,j}) \subseteq V(G) \setminus S_y$. Since G is d -minimal with $d = d(G)$, $e(G[V(G) \setminus S_x]) \leq d|V(G) \setminus S_x|/2 \leq d\varepsilon n/2 = \varepsilon e(G)$ and similarly $e(G[V(G) \setminus S_y]) \leq \varepsilon e(G)$. Hence, for $i = 1, 2$, $e(\cup_{j \notin J_i} G_{i,j}) \leq \varepsilon e(G) \leq 3\varepsilon e(G_i)$ and thus $e(\cup_{j \in J_i} G_{i,j}) \geq (1 - 3\varepsilon)e(G_i)$. By (3.2), we have for $i = 1, 2$

$$e(G'_i) \geq (1 - 2\varepsilon)e(\cup_{j \in J_i} G_{i,j}) \geq (1 - 2\varepsilon)(1 - 3\varepsilon)e(G_i) \geq (1 - 5\varepsilon)e(G_i).$$

By Claim 3, $\delta(G_i) \geq \frac{1}{3}d(G_i)$, so by Lemma 3.1 it follows that $v(G'_i) \geq (1 - 15\varepsilon)v(G_i)$. \square

By Claim 5 and the fact that $\varepsilon \leq \frac{1}{40}$, $V(G'_1) \cap V(G'_2) \neq \emptyset$. Let z be any vertex in $V(G'_1) \cap V(G'_2)$. Suppose $z \in U_{1,i} \cap U_{2,j}$. By definition, there exists a rainbow path P'_i in G_1 going from x_i to z and a rainbow path Q'_j in G_2 going from y_j to z , both of which have length at most $\ell + 1$ and avoid $F \cup V(F_1) \cup V(F_2)$ and $C \cup C(F_1) \cup C(F_2)$. Since G_1, G_2 are colour-disjoint, and $P'_i \subseteq G_1, Q'_j \subseteq G_2$, $C(P'_i) \cap C(Q'_j) = \emptyset$. Also, since P'_i and Q'_j avoid $F \cup F_1 \cup F_2$ while $P_i, Q_j \subseteq F_1 \cup F_2$, P'_i and Q'_j are colour-disjoint from P_i, Q_j . Finally, recall that Q_j avoids F_1 and hence $C(P_i) \cap C(Q_j) = \emptyset$. Hence, $P_i \cup P'_i \cup Q'_j \cup Q_j$ is a rainbow walk of length at most $4\ell + 4$ from x to y , that avoids F and C and hence contains a rainbow x, y -path of length at most $4\ell + 4$ that avoids F and C , as desired. \square

Corollary 3.3. For every integer $t \geq 2$, $0 < \varepsilon \leq \frac{1}{40}$ and sufficiently large n the following holds. Let $\lambda = \frac{\varepsilon}{2 \log n}$. Let $\ell = \lceil \frac{4 \log n}{\lambda} \rceil$. Let $L = \binom{t}{2}(4\ell + 4) + t$. Let K be a positive integer. Let $M \geq L + \frac{12K}{\varepsilon}(\ell + 2)$. Let d be a real that satisfies $d \geq \frac{48(\ell + 2M)}{\varepsilon \lambda}$. Let G be a properly edge-coloured n -vertex graph with $d(G) \geq d$. Suppose G contains no subgraph G' satisfying both $d(G') \geq \frac{\varepsilon d(G)}{12}$ and $v(G') \leq \frac{n}{K}$. Then G contains a rainbow subdivision of K_t , where each edge of K_t is replaced with a path of length at most $4\ell + 4$.

Proof. By Lemma 2.5, G contains a subgraph H which is a $(d(H), \lambda, \varepsilon)$ -expander with $d(H) \geq \frac{d(G)}{2}$. We may assume that H does not have a subgraph H' on at most $\frac{v(H)}{K}$ vertices with $d(H') \geq \frac{\varepsilon d(H)}{6}$. Indeed, otherwise, such a subgraph H' would be a subgraph of G on at most $\frac{n}{K}$ vertices with $d(H') \geq \frac{\varepsilon d(H)}{6} \geq \frac{\varepsilon d(G)}{12}$, contradicting our assumption about G .

Let $S = \{x_1, \dots, x_t\}$ be any set of t distinct vertices in H . Let A be a maximal collection of pairs (i, j) with $1 \leq i < j \leq t$ such that there exist paths $P_{i,j}$ for $(i, j) \in A$ satisfying the following.

1. For each $(i, j) \in A$, $P_{i,j}$ is an x_i, x_j -path of length at most $4\ell + 4$ in H such that $V(P_{i,j}) \cap \{x_1, \dots, x_t\} = \{x_i, x_j\}$.
2. The sets $V(P_{i,j}) \setminus \{x_i, x_j\}$ are pairwise disjoint over different $(i, j) \in A$ and $\cup_{(i,j) \in A} P_{i,j}$ is rainbow.

If A consists of all pairs (i, j) with $1 \leq i < j \leq t$, then the union of the paths $P_{i,j}$ forms the required subdivision of K_t and we are done. Hence, we may assume that there exist i, j such that $1 \leq i < j \leq t$ and $(i, j) \notin A$.

Let C be the set of colours used on $\cup_{(k,\ell) \in A} P_{k,\ell}$ and let $F = V(\cup_{(k,\ell) \in A} P_{k,\ell}) \cup S \setminus \{x_i, x_j\}$. Then $|C|, |F| \leq \binom{t}{2}(4\ell + 4) + t = L$. Since n is sufficiently large and $d(H) \geq \frac{d}{2}$, we may assume $v(H)$ is sufficiently large. Also, $d(H) \geq \frac{d}{2} \geq \frac{24(\ell + 2M)}{\lambda}$. Hence, by Lemma 3.2, H contains a rainbow x_i, x_j -path $P_{i,j}$ that avoids C and F and whose length is at most $4\ell + 4$. The pair (i, j) could thus be added to A , contradicting maximality of A . \square

Now we are ready to prove Theorem 1.2.

Proof of Theorem 1.2. Let $\varepsilon = \frac{1}{40}$. The choice of n_0 is not explicit; we assume it is sufficiently large with respect to t so that various inequalities below hold.

Choose $c = \max\{\sqrt{\log n_0}, 2 \log(\frac{12}{\varepsilon})\}$. By our choice of c , $e^{\frac{c}{2}} \geq \frac{12}{\varepsilon}$. We may assume $n \geq n_0$. Let $\lambda = \frac{\varepsilon}{2 \log n}$. Let $\ell = \lceil \frac{4 \log n}{\lambda} \rceil$. Let $L = \binom{t}{2}(4\ell + 4) + t$. Let $K = \lceil e^{\sqrt{\log n}} \rceil$. Let $M = L + \frac{12K}{\varepsilon}(\ell + 2)$. Note that by our choice of c and n_0 , we have $e^{\frac{c}{2} \sqrt{\log n}} \geq \frac{48(\ell + 2M)}{\varepsilon \lambda}$.

Let G be a properly edge-coloured graph on n vertices with at least $ne^{c\sqrt{\log n}}$ edges, so $d(G) \geq e^{c\sqrt{\log n}}$. Let $G_0 = G$. We define a sequence of subgraphs of G as follows. For each $i \geq 0$, suppose G_0, \dots, G_i have been defined. If G_i contains a subgraph F with $d(F) \geq \frac{\varepsilon}{12}d(G_i)$ and $v(F) \leq \frac{v(G_i)}{K}$, let $G_{i+1} = F$. If no such subgraph of G_i exists then we stop the process. Suppose the process stops when considering the graph G_m , so no such subgraph F can be found inside G_m . Then $1 \leq v(G_m) \leq \frac{n}{K^m}$. Hence $m \leq \log_K n \leq \frac{\log n}{\log e^{\sqrt{\log n}}} = \sqrt{\log n}$. Also, since $e^{\frac{c}{2}} \geq \frac{12}{\varepsilon}$ by our assumption, we have

$$d(G_m) \geq \left(\frac{\varepsilon}{12}\right)^m d(G) \geq \left(\frac{\varepsilon}{12}\right)^{\sqrt{\log n}} e^{c\sqrt{\log n}} \geq e^{\frac{c}{2}\sqrt{\log n}}.$$

By our choice of c and n_0 , we have $d(G_m) \geq e^{\frac{c}{2}\sqrt{\log n}} \geq \frac{48(\ell + 2M)}{\varepsilon \lambda'} \geq \frac{48(\ell' + 2M')}{\varepsilon \lambda'}$, where $\lambda' = \frac{\varepsilon}{2 \log v(G_m)}$, $\ell' = \lceil \frac{8 \log v(G_m)}{\lambda} \rceil$, $L' = \binom{t}{2}(4\ell' + 4) + t$, $M' = L' + \frac{12K}{\varepsilon}(\ell' + 2)$. Also, $v(G_m) > d(G_m) \geq$

$e^{\frac{c}{2}\sqrt{\log n}} > n_0$, and by the definition of G_m , G_m contains no subgraph F satisfying both $d(F) \geq \frac{\varepsilon}{12}d(G_m)$ and $v(F) \leq \frac{v(G_m)}{K}$.

So by [Corollary 3.3](#), G_m contains a rainbow subdivision of K_t , where each edge of K_t is replaced with a path of length at most $4\ell + 4 \leq 1300 \log^2 n$. This completes our proof.

4. Concluding remarks

This paper establishes that a properly edge-coloured graph on n vertices with average degree at least $(e^{c\sqrt{\log n}})$ contains a rainbow subdivision of a clique. The proof of [Theorem 1.2](#) in fact explicitly gives the following slightly stronger result. Given a positive integer $t \geq 2$, we call a properly edge-coloured graph G *rainbow t -linkable* if for any set of t vertices x_1, \dots, x_t in G there is a rainbow K_t -subdivision with x_1, \dots, x_t playing the role of the t branching vertices. For every $t \geq 2$, there is a constant c such that for every $n \geq t$, if G is a properly coloured n -vertex graph with average degree at least $e^{c\sqrt{\log n}}$, then G contains a subgraph H with average degree at least $\Omega(e^{\frac{\varepsilon}{2}\sqrt{\log n}})$ that is rainbow t -linkable.

After the initial submission of our paper, there has been further improvements on the average degree that enforces the existence of a rainbow cycle and a rainbow subdivision of a clique. For rainbow cycles, Janzer's result was improved first by Tomon [17] to $(\log n)^{2+o(1)}$ and by Kim, Lee, Liu, Tran [7] to $100(\log n)^2$. As for the family of subdivisions of a clique, in a subsequent paper with Letzter [5] we improved the average degree bound to $(\log n)^{60}$. Subsequently, Tomon [17] improved this average degree lower bound to $(\log n)^{6+o(1)}$, and finally Wang [18] improved it to $(\log n)^{2+o(1)}$. We pose the following question (also stated in [5]).

Question 4.1. *Given $t \geq 3$, what is the smallest c such that for all sufficiently large n , if G is a properly edge-coloured graph on n vertices with $\Omega(n(\log n)^c)$ many edges then G contains a rainbow subdivision of K_t ? In particular, is $c = 1$?*

In the above question, our guess of $c = 1$ is motivated by the only lower bound of order $\Omega(n \log n)$ given by the construction of a colouring of the hypercube by Keevash, Mubayi, Sudakov and Verstraëte [6] described earlier in the introduction.

Acknowledgements

We thank Shoham Letzter for carefully reading an early draft of this paper and for helpful comments. We also thank the referees for their valuable feedback which improved the presentation of the paper significantly.

References

- [1] S. Das, C. Lee, B. Sudakov, Rainbow Turán problem for even cycles, *European J. Combin.* 34 (2013) 905–915.
- [2] S. Fiorini, G. Joret, D.O. Theis, D.R. Wood, Small minors in dense graphs, *European J. Combin.* 33 (2012) 1226–1245.
- [3] O. Janzer, Rainbow Turán number of even cycles, repeated patterns and blow-ups of cycles, *Israel J. Math.* (2022) <http://dx.doi.org/10.1007/s11856-022-2380-9>.
- [4] T. Jiang, Compact topological minors in graphs, *J. Graph Theory* 67 (2011) 139–152.
- [5] T. Jiang, S. Letzter, A. Methuku, L. Yepremyan, Rainbow clique subdivisions and blow-ups, *arXiv:2108.08814*.
- [6] P. Keevash, D. Mubayi, B. Sudakov, J. Verstraëte, Rainbow Turán problems, *Combin. Probab. Comput.* 16 (2007) 109–126.
- [7] J. Kim, J. Lee, H. Liu, T. Tran, Rainbow cycles in properly edge-colored graphs, 2022, *arXiv:2211.03291*.
- [8] A. Kostochka, Lower bound of the Hadwiger number of graphs by their average degree, *Combinatorica* 4 (1984) 307–316.
- [9] A. Kostochka, L. Pyber, Small topological complete subgraphs of “dense” graphs, *Combinatorica* 8 (1988) 83–86.
- [10] W. Mader, Homomorphieeigenschaften und mittlere Kantendichte von Graphen, *Math. Ann.* 174 (1967) 265–268.
- [11] M. Molloy, B. Reed, *Graph Colouring and the Probabilistic Method*, Vol. 23, Springer Science & Business Media, 2013.
- [12] R. Montgomery, Logarithmically-small minors and topological minors, *J. Lond. Math. Soc.* 91 (2015) 71–88.
- [13] A. Shapira, B. Sudakov, Small complete minors above the extremal edge density, *Combinatorica* 35 (2015) 75–94.
- [14] B. Sudakov, I. Tomon, The extremal number of tight cycles, *Int. Math. Res. Not. IMRN* 2022 (2022) 9663–9684.
- [15] A. Thomason, An extremal function for contractions of graphs, in: *Math. Proceedings of the Cambridge Phil. Society*, Vol. 95, Cambridge University Press, 1984, pp. 261–265.
- [16] A. Thomason, The extremal function for complete minors, *J. Combin. Theory Ser. B* 81 (2001) 318–338.
- [17] I. Tomon, Robust (rainbow) subdivisions and simplicial cycles, 2022, *arXiv:2201.12309*.
- [18] Y. Wang, Rainbow clique subdivisions, 2022, *arXiv:2204.08804*.